

# Lecture 7: Geometric Modeling and Visualization

## Boundary & Finite Element Meshed Models III: Topologically Accurate Non-Linear Elements

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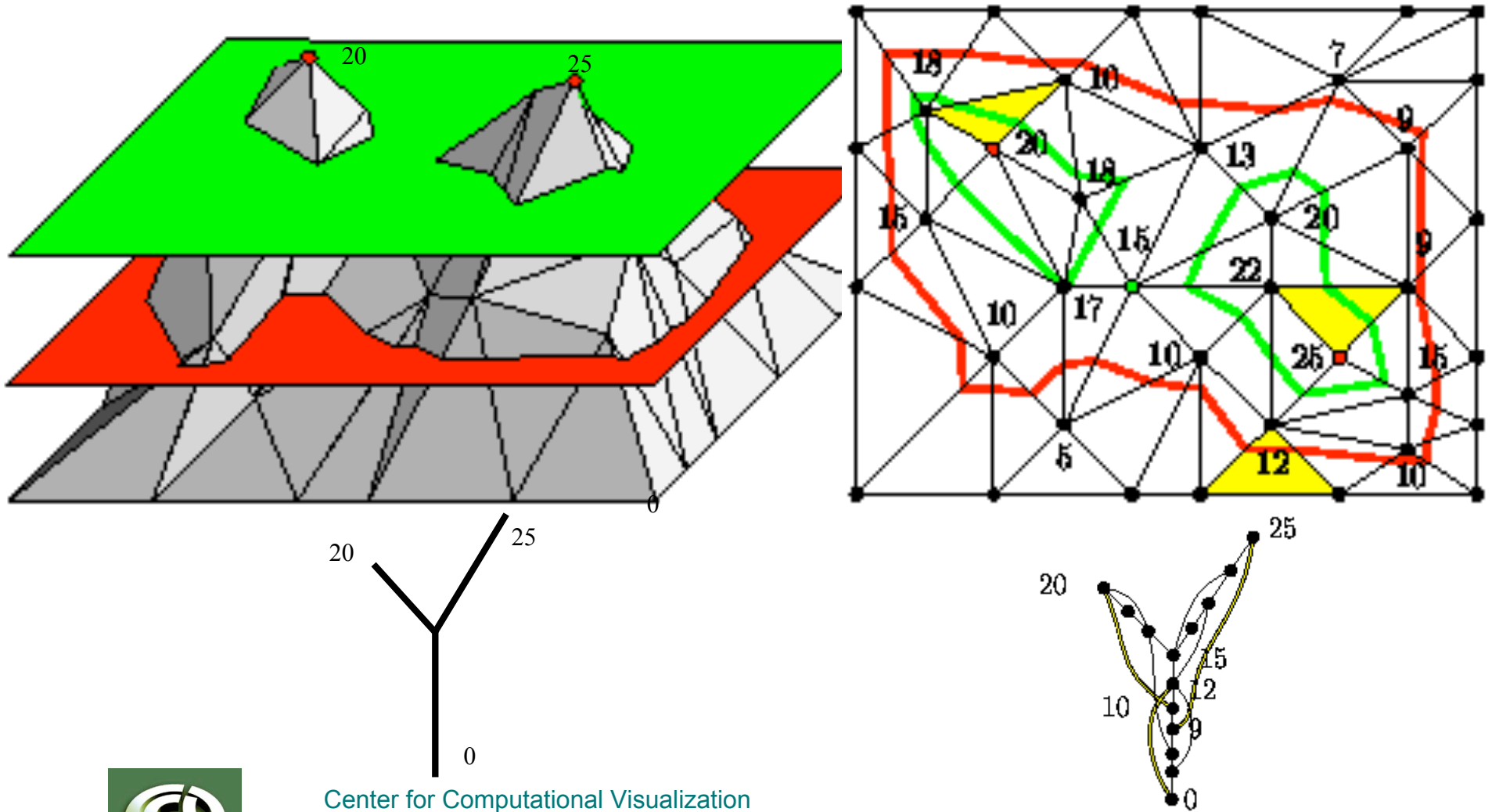


Center for Computational Visualization  
Institute of Computational and Engineering Sciences  
Department of Computer Sciences

University of Texas at Austin

October 2007

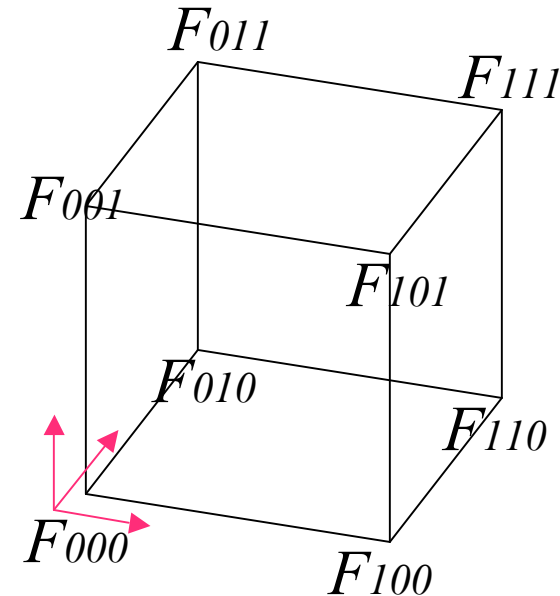
# Contouring: Capturing the Topology and Geometry of Zero Sets



# Isosurface of Trilinear Function

- Trilinear Function

$$\begin{aligned} F(x,y,z) &= F_{000}(1-x)(1-y)(1-z) \\ &+ F_{001}(1-x)(1-y)z \\ &+ F_{010}(1-x)y(1-z) \\ &+ F_{011}(1-x)yz \\ &+ F_{100}x(1-y)(1-z) \\ &+ F_{101}x(1-y)z \\ &+ F_{110}xy(1-z) \\ &+ F_{111}xyz \end{aligned}$$



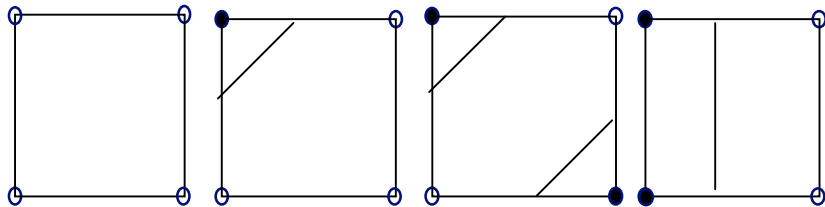
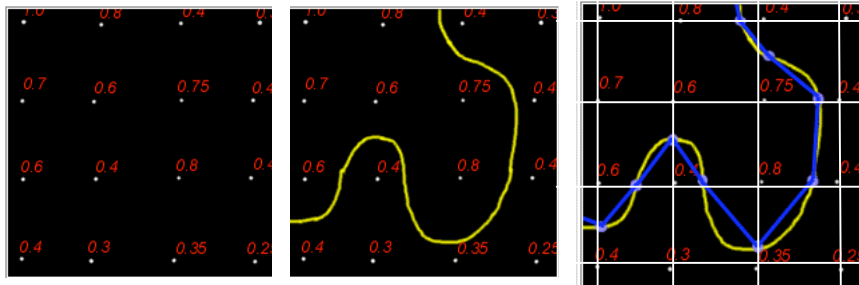
- Bilinear Function

$$\begin{aligned} F^f(x,y) &= F_{00}(1-x)(1-y) + F_{01}(1-x)y \\ &+ F_{10}x(1-y) + F_{11}xy \end{aligned}$$

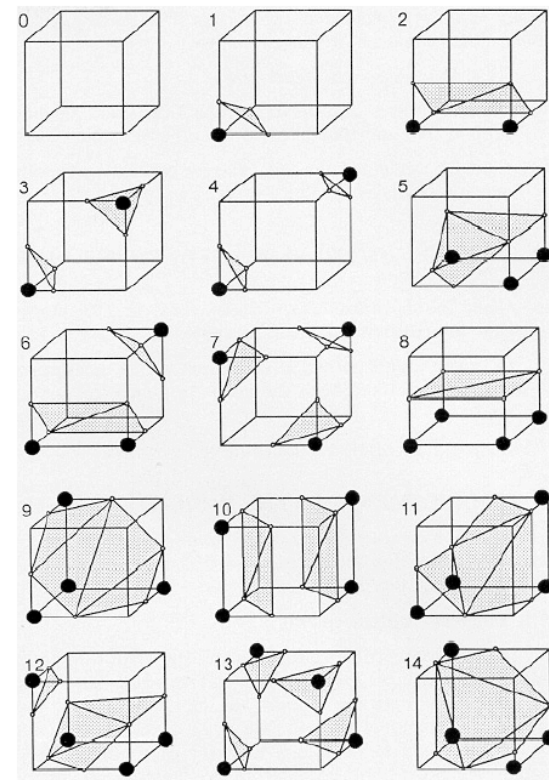


# Marching Cubes (MC) : Triangular Approximation

- 2D rectangle



- 3D cube :  
15 Cases for Triangulation





# Saddle Points Computation

- Face Saddle Point

$$F(x,y) = ax + by + cxy + d \quad (\text{bilinear interpolant})$$

$$\text{First derivatives : } F_x = a + cy = 0, F_y = b + cx = 0$$

$$\text{Saddle point } S = (-b/c, -a/c)$$

- Body Saddle Point

$$F(x,y,z) = a + ex + cy + bz + gxy + fxz + dyz + hxyz$$

First derivatives = 0 :

$$F_x = e + gy + fz + hyz = 0$$

$$F_y = c + gx + dz + hxz = 0$$

$$F_z = b + fx + dy + hxy = 0$$



# Face and Body Saddle Points

- We obtain saddle points :

$$x = -\frac{c+dz}{g+hz}$$

$$y = \frac{k_0+k_1z}{k_2}$$

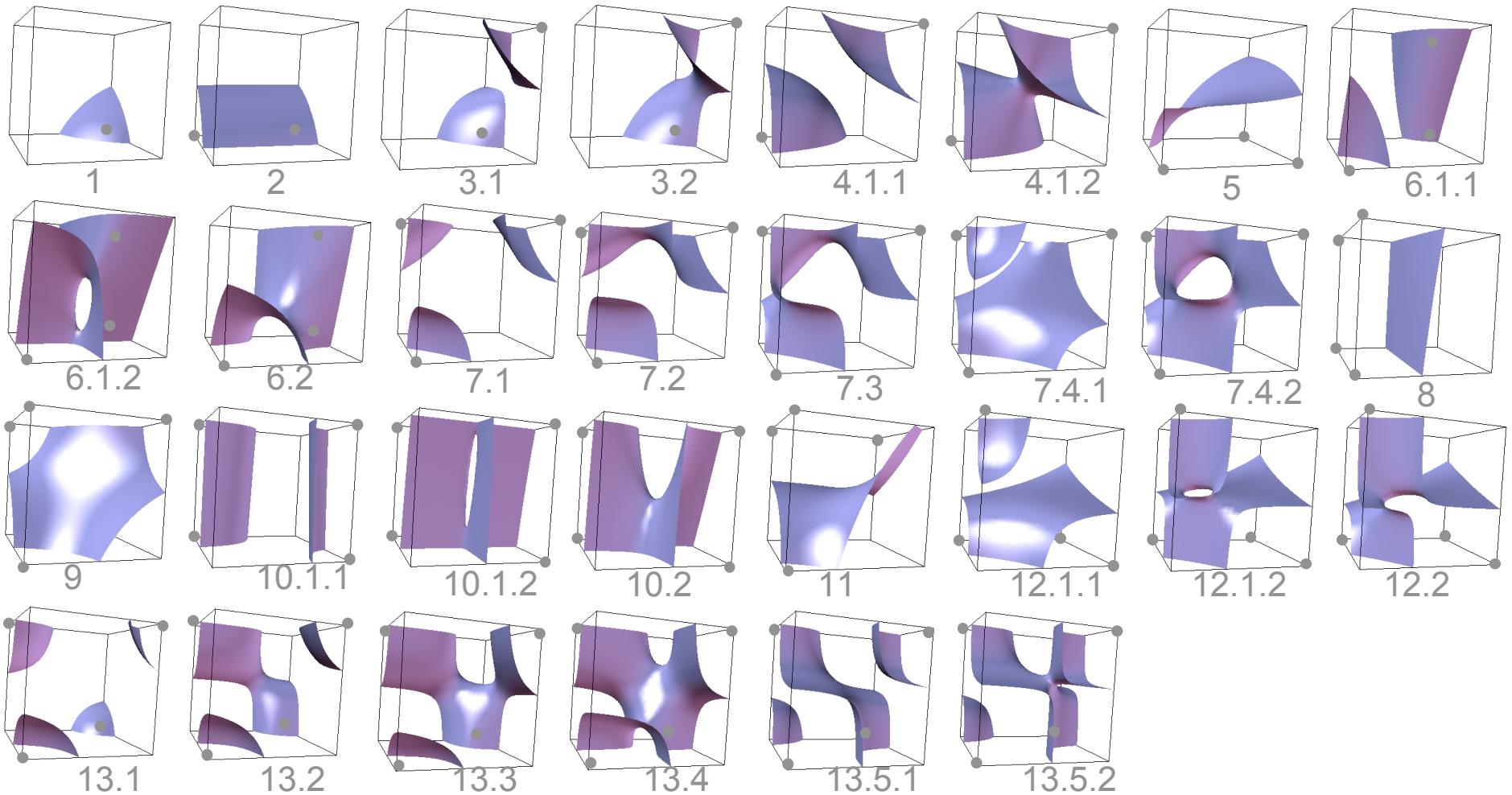
$$z = -\frac{g}{h} \pm \frac{\sqrt{g^2k_1^2 - hk_1^{1/2}(ek_2+gk_0)}}{h}$$

$$k_0 = cf - bg, k_1 = df - bh, k_2 = dg - ch$$

- saddle point outside the cube → discard  
( only case 13.5 has more than one valid body saddle point. )



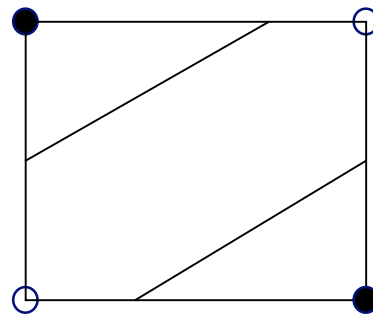
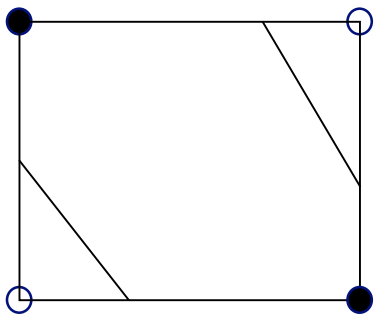
# Trilinear Isosurface Topology 31 cases



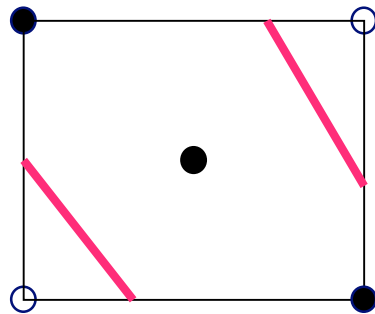
# Decision on Contour Topology ( Nielson 92 : Asymptotic Decider )

- Resolving Face Ambiguity

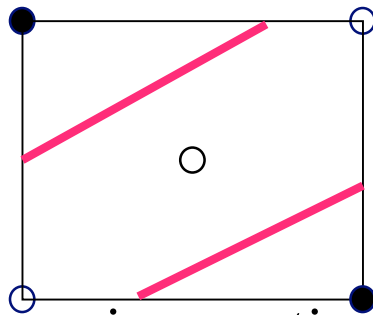
- Ambiguity ( face saddle )



- Decision based on the value  $s$  of saddle point



*$s$  is positive*

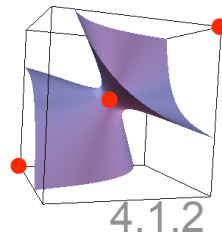
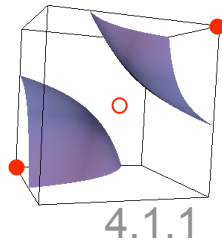


*$s$  is negative*



# Decision on Contour Topology ( Natarajan 94 )

- Resolving Internal Ambiguity
  - Ambiguity ( body saddle )



–Decision based on the value  $s$  of saddle point

- (i)  $s$  is positive  $\rightarrow$  tunnel
- (ii)  $s$  is negative  $\rightarrow$  two pieces



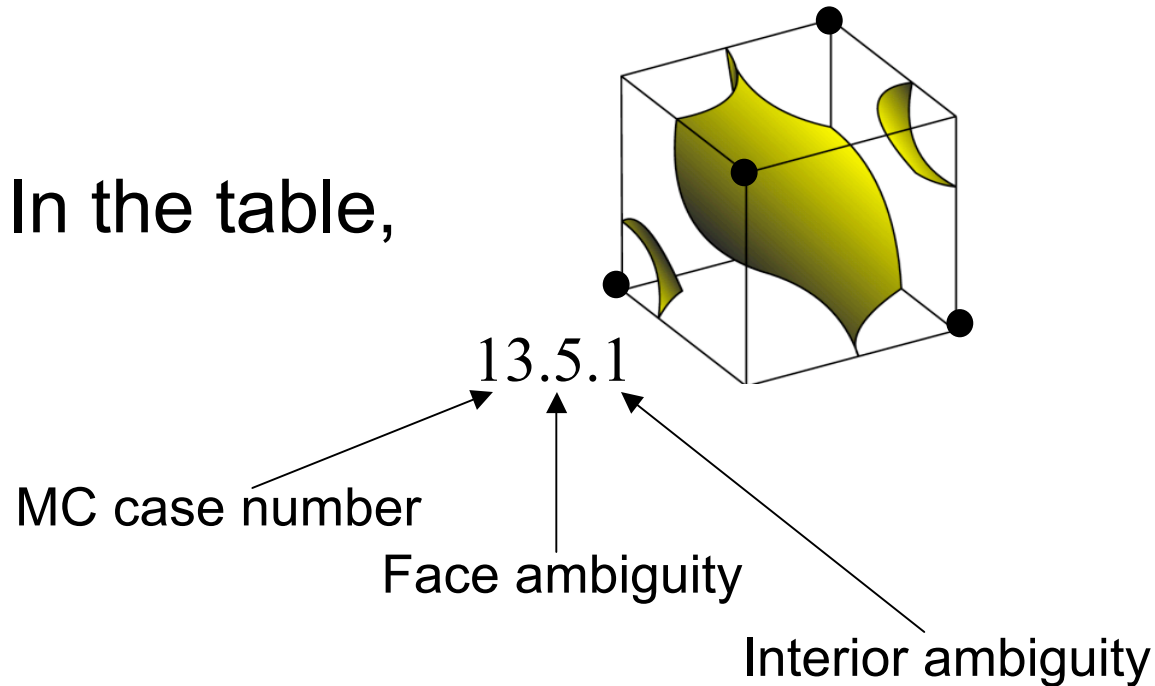
# Contour Topology Decision

- Trilinear isosurface connectivity is determined by sign configuration of saddle points and 8 corner vertices of a cube
- Marching Cubes : Consider only 8 corner vertices. Additional Ambiguity problems exist



# 31 Cases

- In the table,

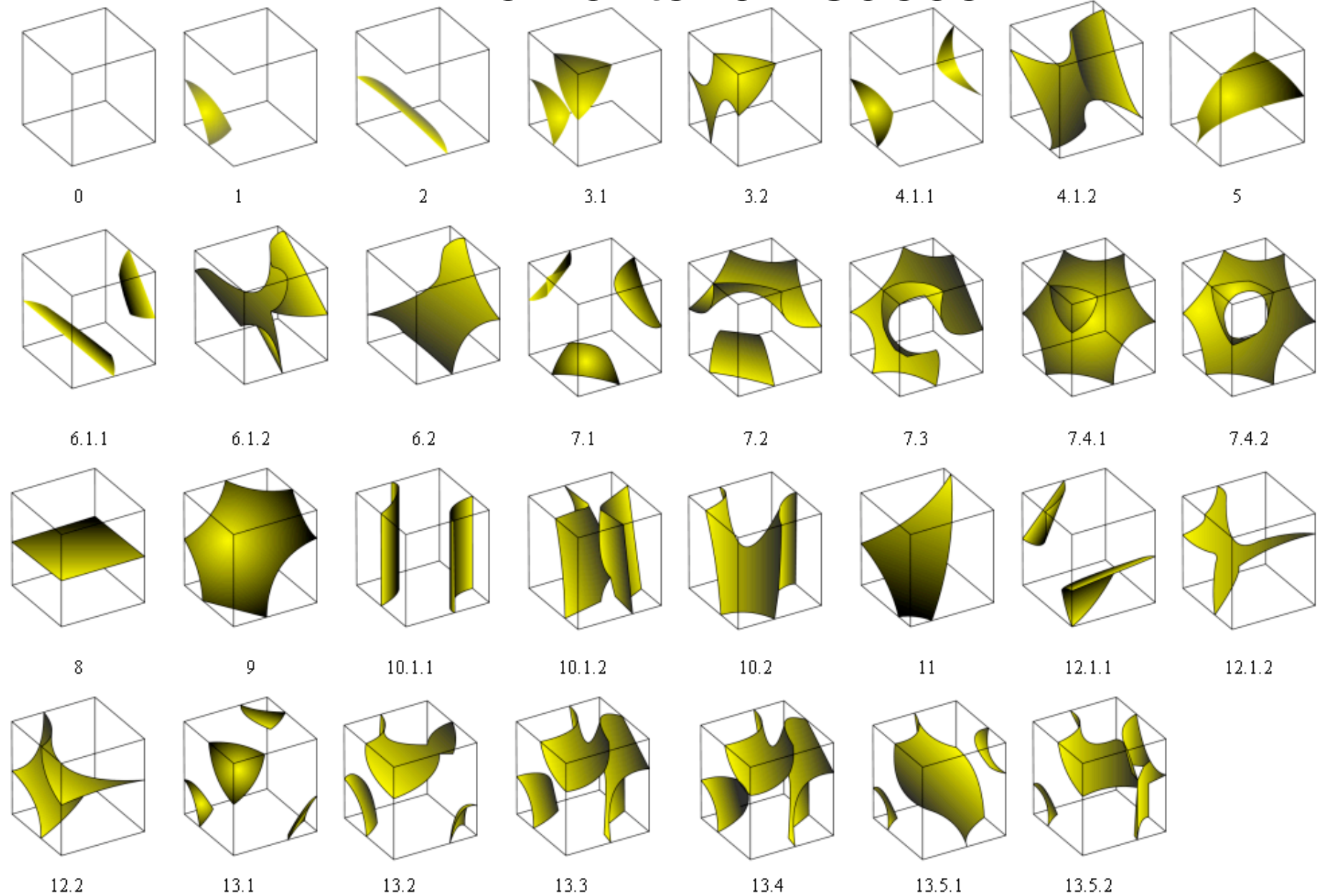


- MC : 15 (further reduced to 14) cases based on vertex coloring (symm).
- 31 cases ← (vertex coloring , face ambiguity , internal ambiguity)

Symmetry of different configurations are used to reduce the cases.

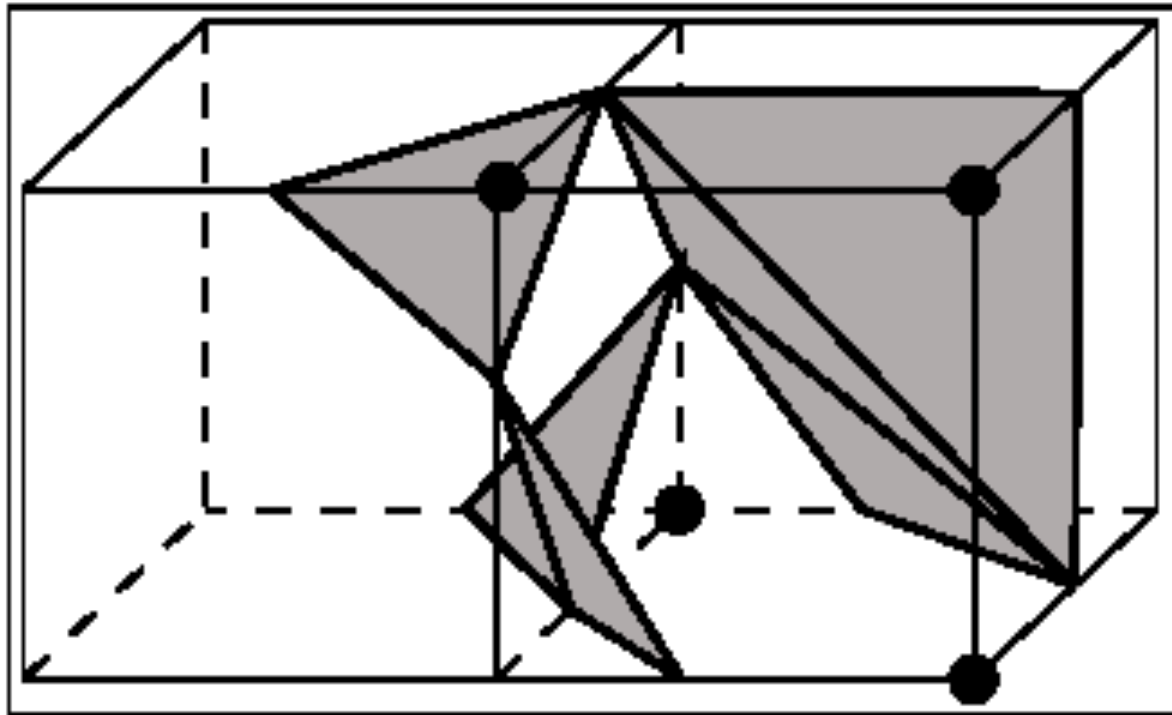


# Trilinear Implicit Surface Boundary Elements: 31 Cases





# Triangulation Ambiguity



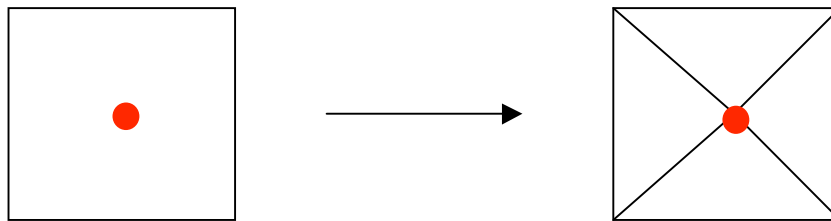
<wrong surface>

- Saddle points play important roles in determining contour connectivity

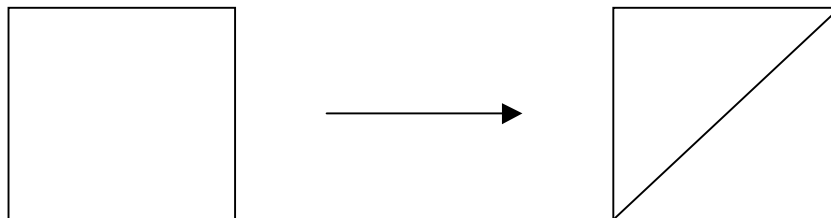


# Topology Preserving Tetrahedral Decomposition

- 2D case
  - If there is a saddle point

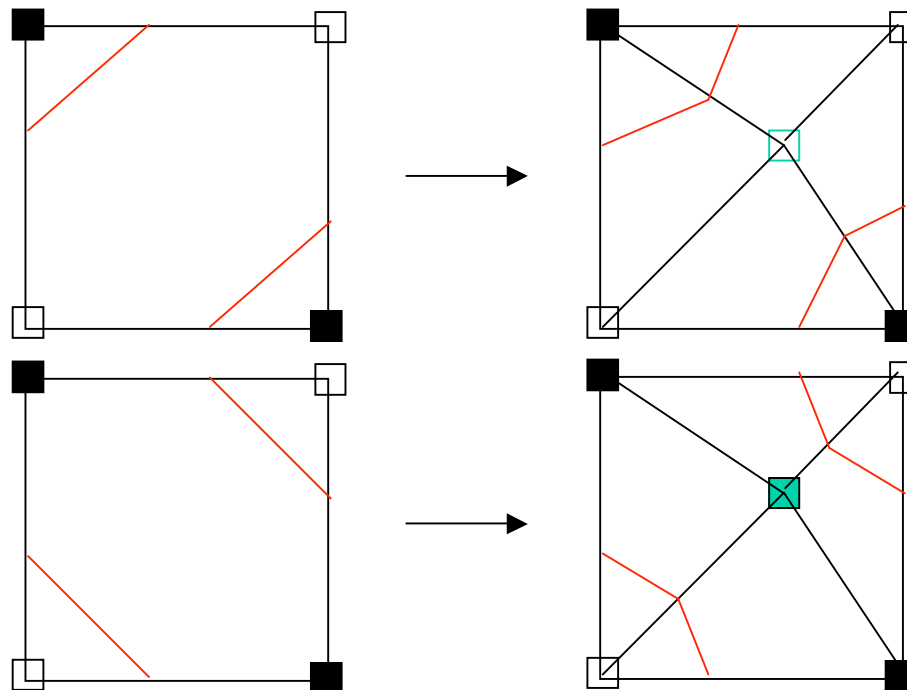


- If there is no saddle point



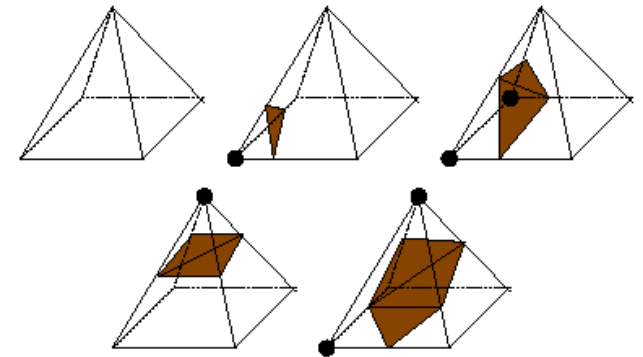
# Cell Decomposition Method

- Decompose a cell when a saddle point affects the contour connectivity



# Main Decomposition Rule for Trilinear Cell with Topological Ambiguity

- If isosurface has a tunnel
  - With a body saddle point generate six pyramids with the cube faces
  - Further decompose pyramids that have face ambiguity into four tetrahedra
- If isosurface has no tunnel
  - Choose a face saddle and generate five pyramids with remaining faces
  - Further decompose pyramids that have face ambiguity into four tetrahedra
- Case 13 is an exception

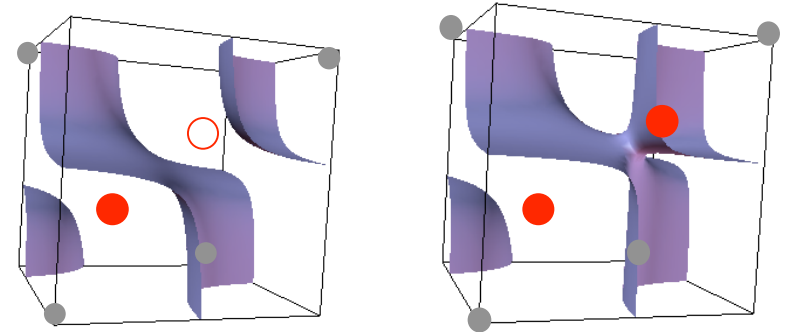


< pyramid triangulation >

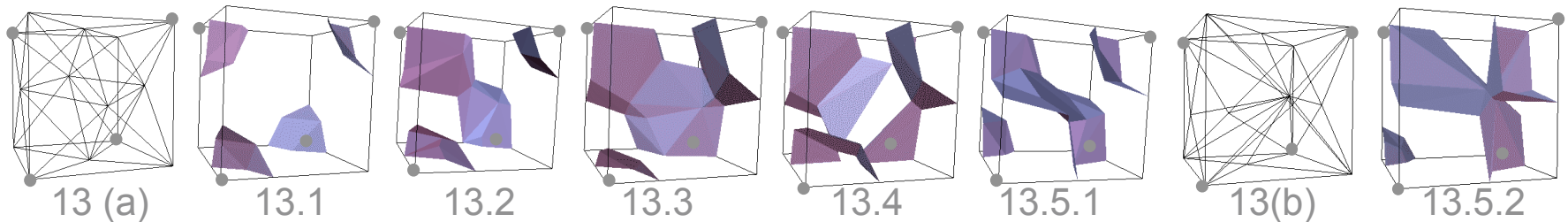


# Complicated Topology

- Case 15 of MC and # 13
  - the most complicated case in geometry and topology
  - involve

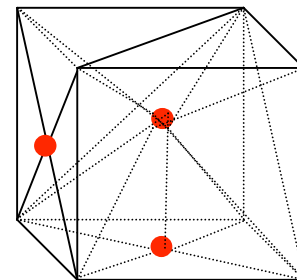
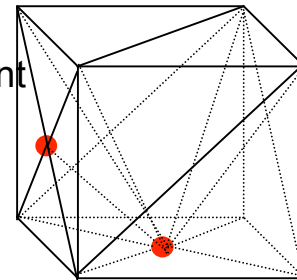
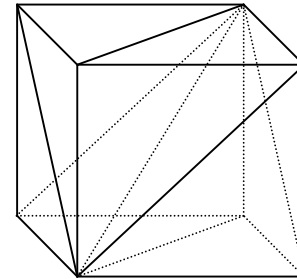


Face saddles for each face and two body saddles



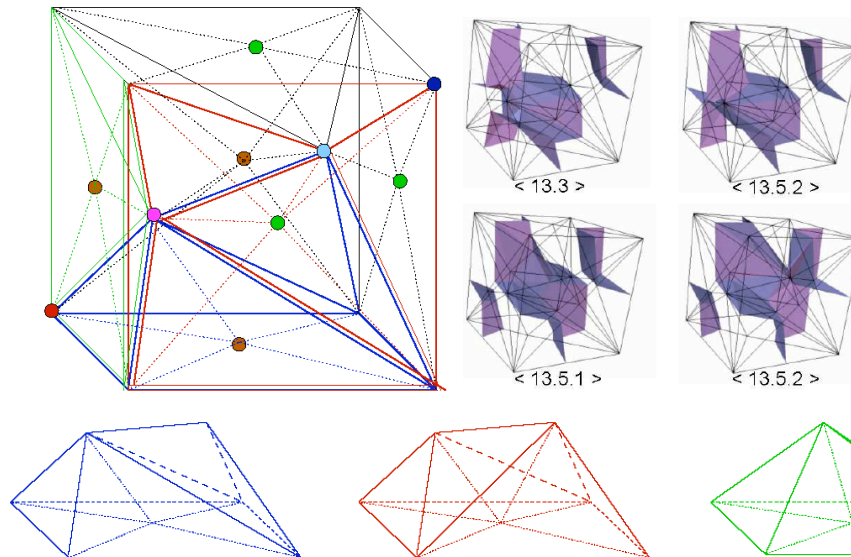
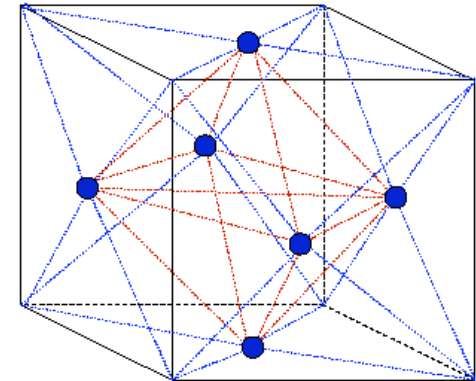
# Topology Preserving Tetrahedral Decomposition

- 3D case (  $S_b$  and  $S_f = \#$  of body and face saddles )
  - (i)  $S_b = 0 ; S_f = 0$ 
    - Standard decomposition ( 6 tet )
  - (ii)  $S_b = 0 \ \& \ 1 \leq S_f \leq 4$ 
    - Decompose a face with a face saddle into 4 tris
    - Decompose a face without a face saddle into 2 tris
    - Choose one face saddle and connect it to each face to form 5 pyramids. Each pyramid decomposed into four or two tets (Choice of 2nd largest face saddle point when 3 or 4 face saddles present)
  - (iii)  $S_b = 1 \ \& \ 1 \leq S_f \leq 4$ 
    - Connect a body saddle to each face to form 6 pyramids
    - Each pyramid decomposed into four or two tets, depending on presence or absence of face saddles



# Topology Preserving Tetrahedral Decomposition (#13)

- (iv)  $S_b = 0$  &  $S_f = 6$   
24 tetrahedral split
- (v)  $1 \leq S_b \leq 2$  &  $S_f = 6$



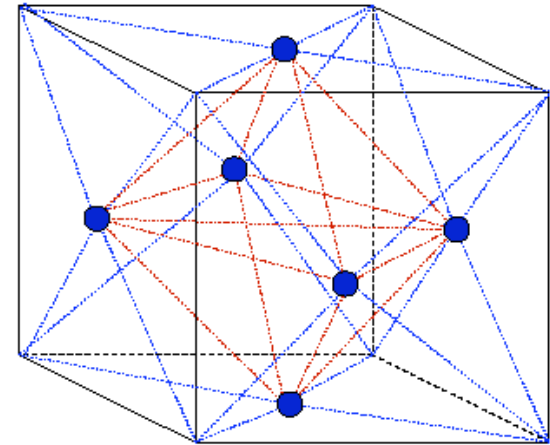
Decompose into  
pyramids & prisms  
and then further split  
into tetrahedra



## #13 Case 4: $S_b = 0$ & $S_f = 6$

- Connect the 6 face saddles forming an 8 triangular-facet diamond, which is split into 4 tetrahedra
- 12 tetrahedra are created by joining 2 vertices of an edge with 2 face saddles of the faces incident at the edge
- 8 additional tetrahedra are created by connecting each facet of the diamond to the 8 vertices of the cube

Overall 24 tetrahedral split





# Case 4:1 $\leq S_b \leq 2$ & $S_f = 6$

Order the saddle points in increasing order of saddle values and into 3 small face saddles, small body saddle, big body saddle, and 3 big face saddles

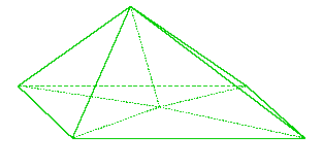
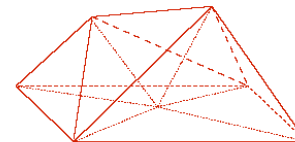
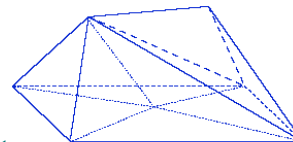
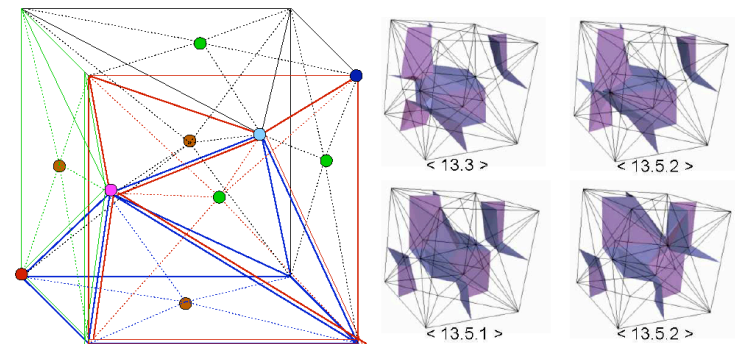
Small/Big corner vertex is a vertex adjacent to the three faces containing small/big face saddles

–  $S_b = 2$

Decompose into 2 Pyramids and 4 Truncprisms.

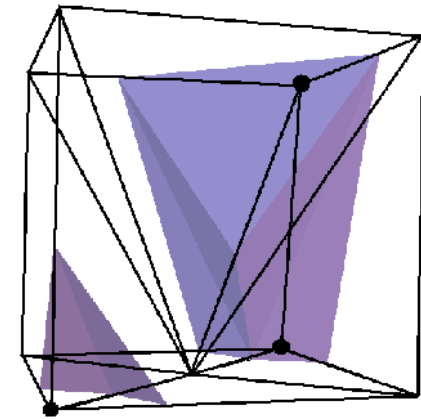
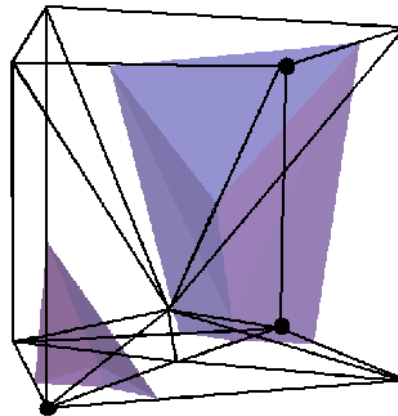
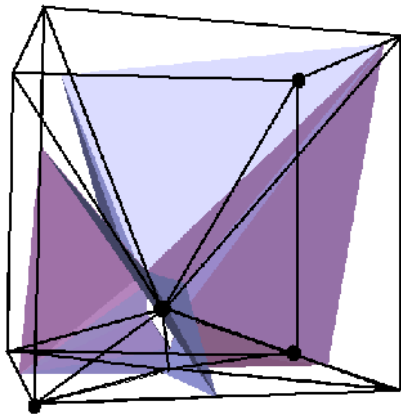
–  $S_b = 1$

Decompose into 1 Pyramid and 4-Truncprisms



# Cell Decomposition Method

- Disambiguate internal topology

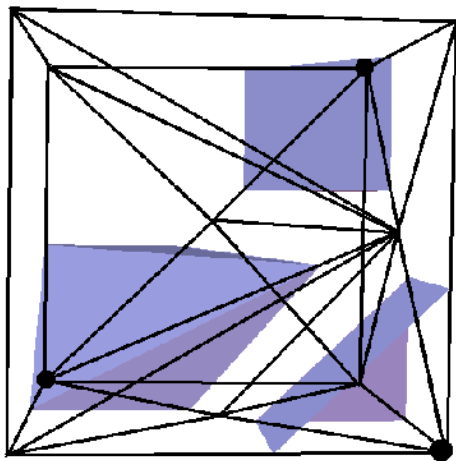


Body saddle can be ignored  
when no tunnel exists

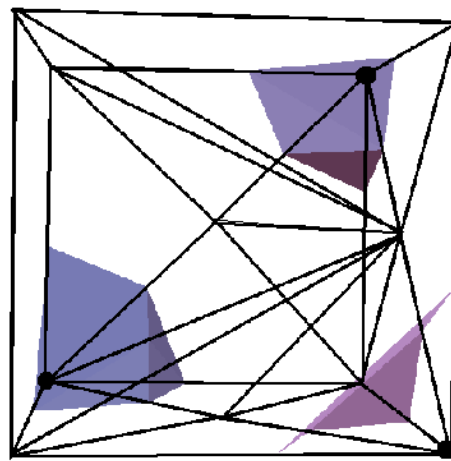


# Geometric Improvement

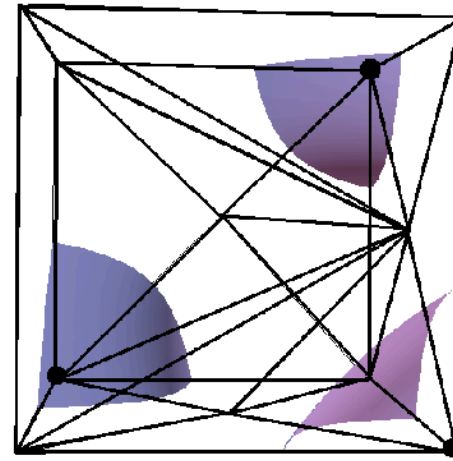
- Compute true intersection between an edge and isosurface



<linear interpolation  
Along edge >



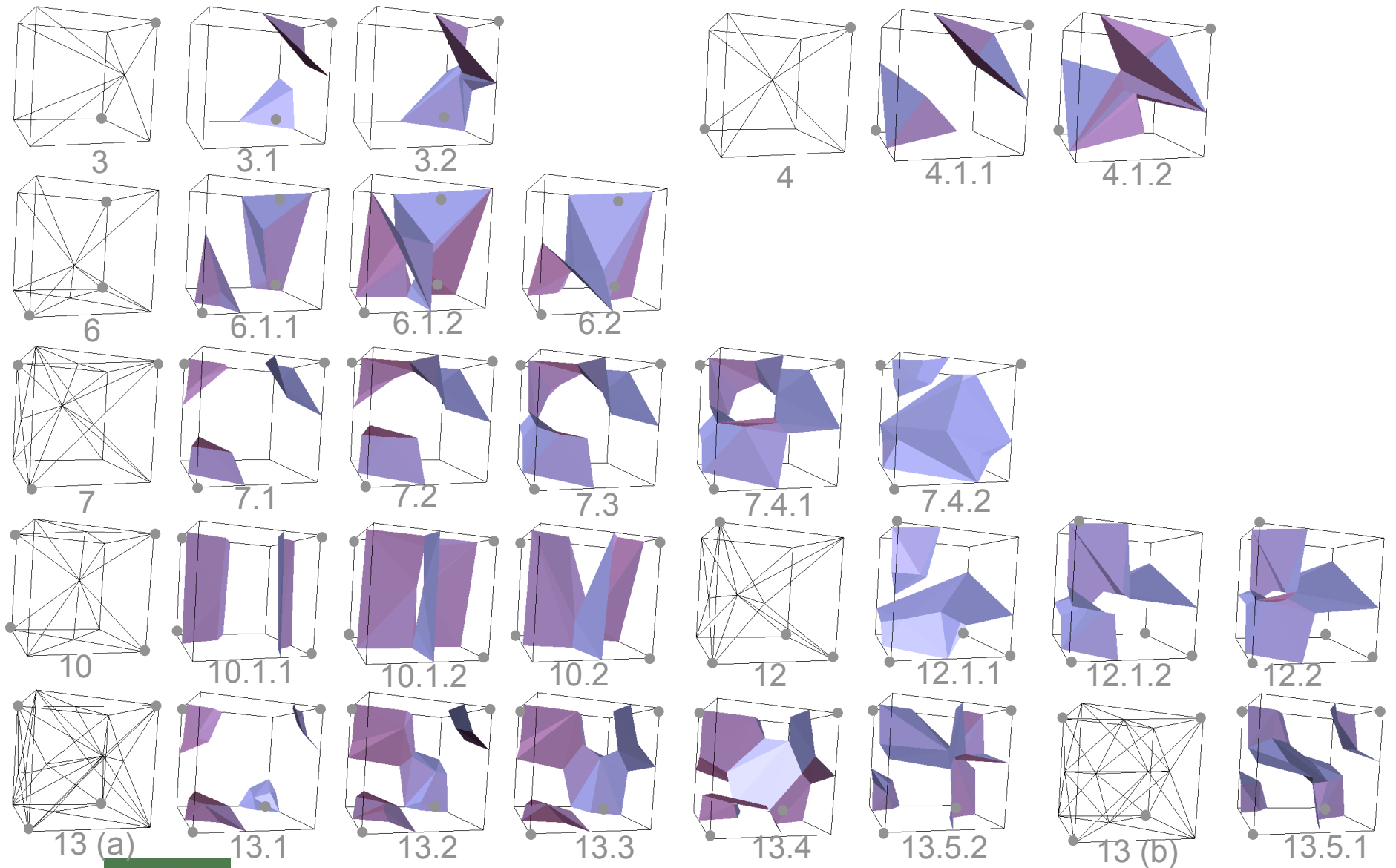
<intersection between  
edge and true isosurface >



<trilinear isosurface >

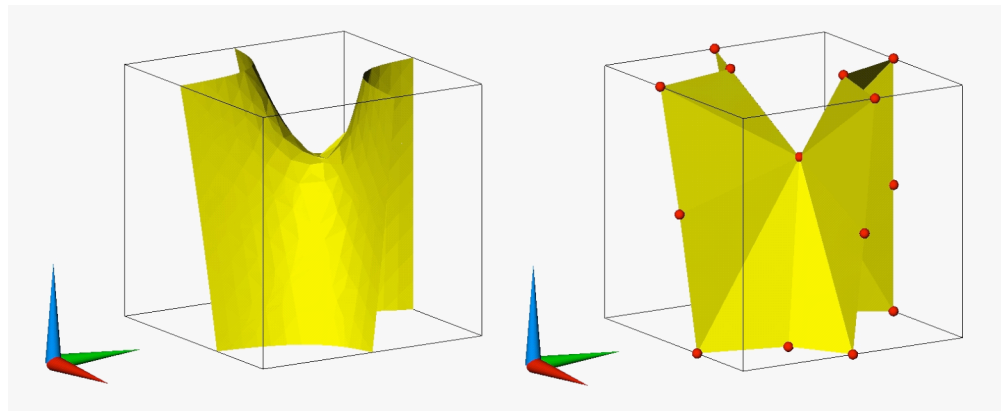
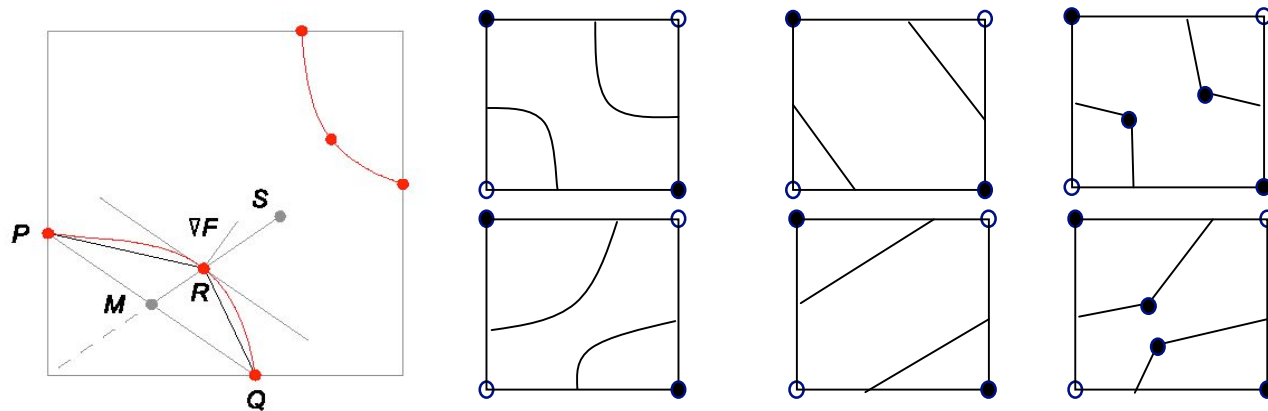


# Results



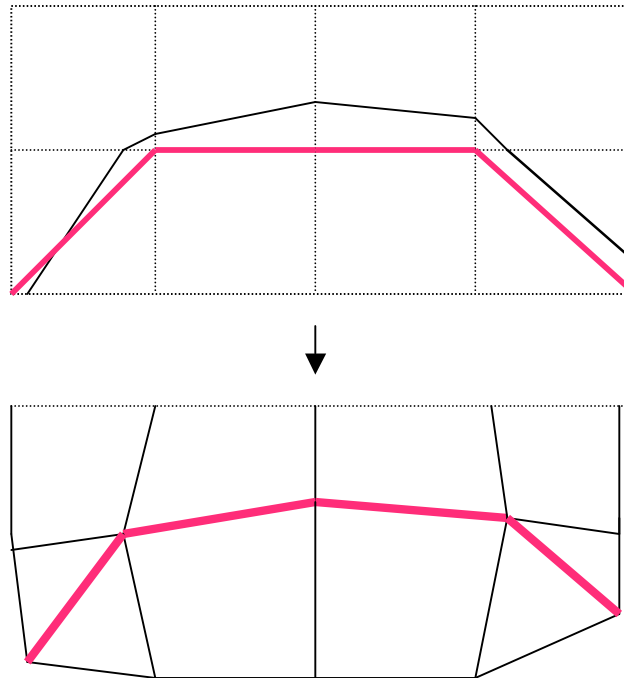
# Geometric Approximations

- Better approximation of trilinear interpolant
  - Adding a shoulder and inflection points

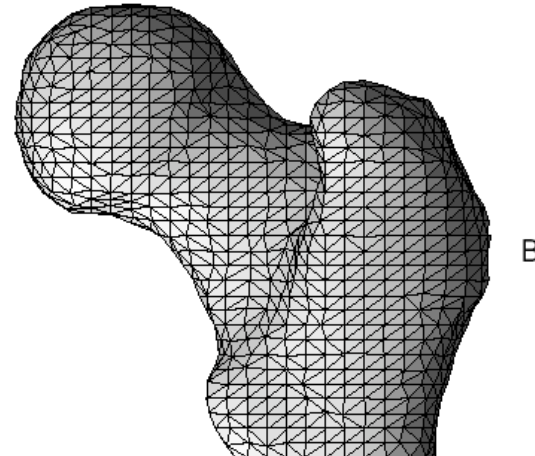
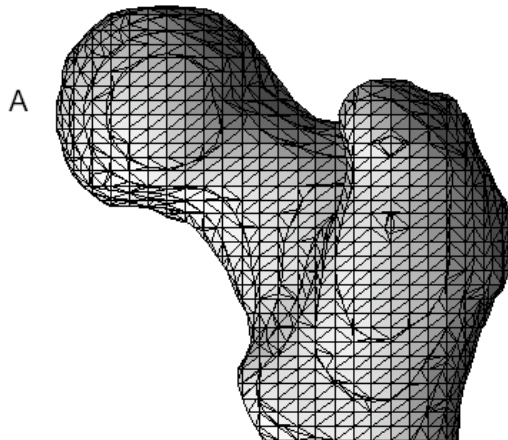
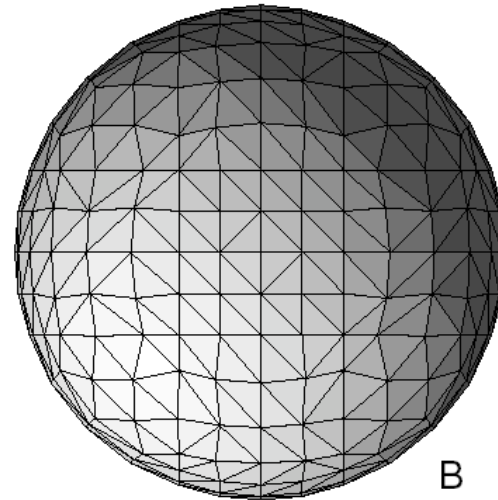
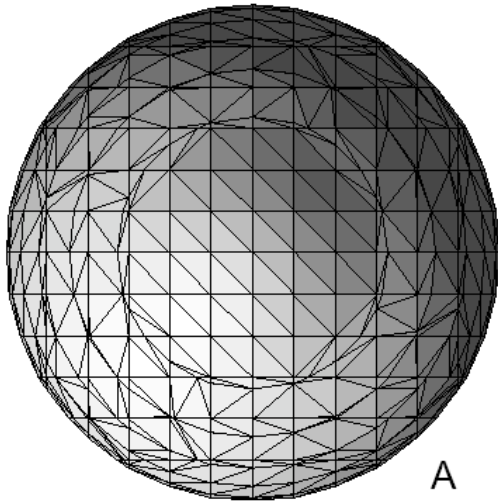


# Mesh Displacement

- Remove small triangles + good aspect ratio



# Mesh Displacement



Marching Cubes



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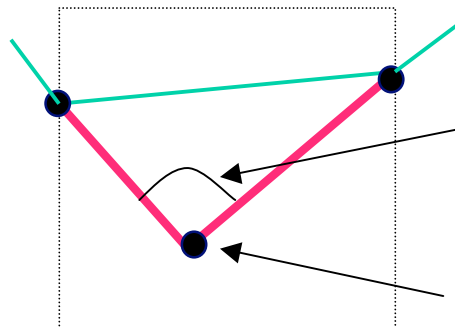
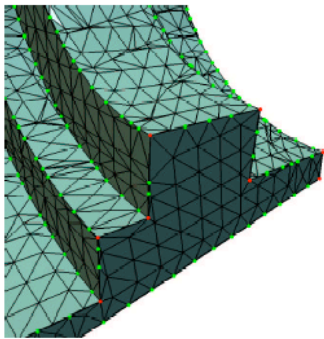
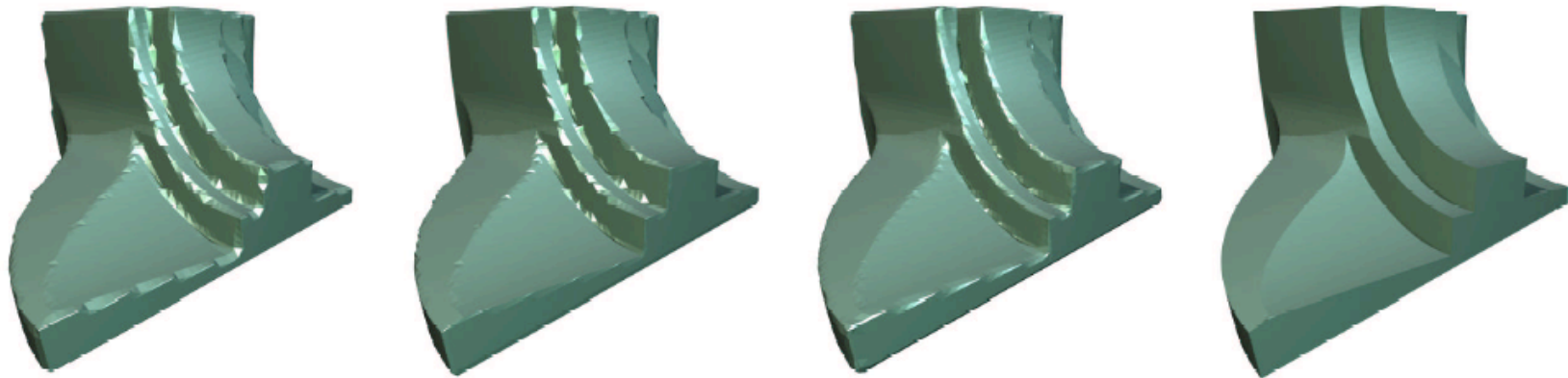
Mesh Displacement

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# Feature Sensitive Surface Extraction

- Extended Marching Cubes



compare the angle with threshold

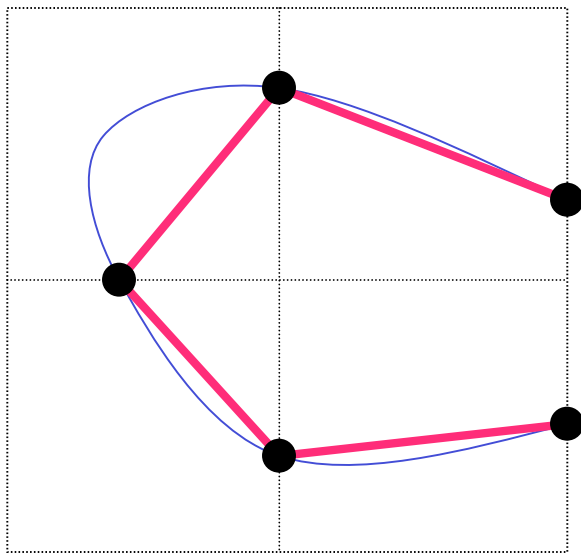
Adding a (edge, corner) feature point



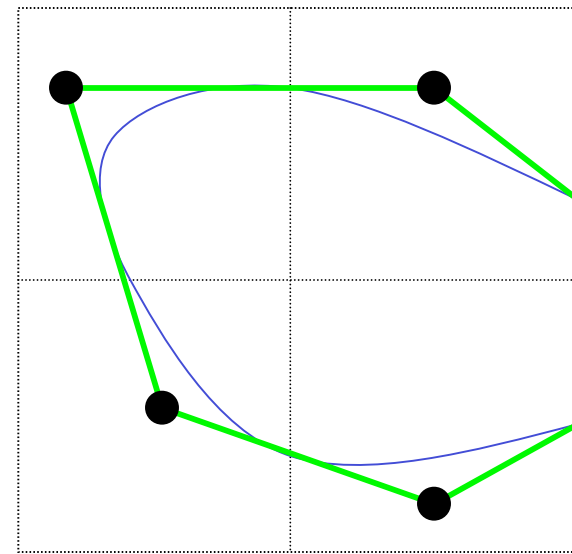


# Dual Contouring

- Primal Contouring vs Dual Contouring



Primal contour

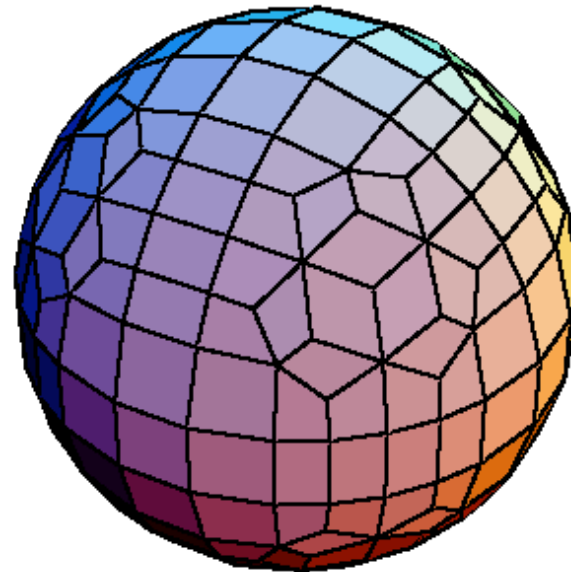
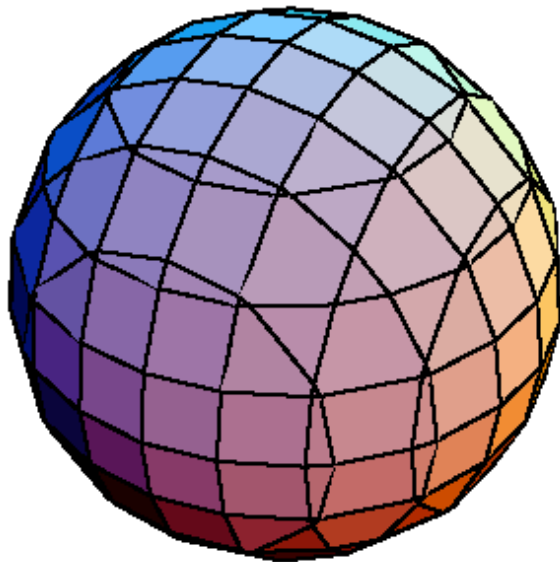


Dual Contour



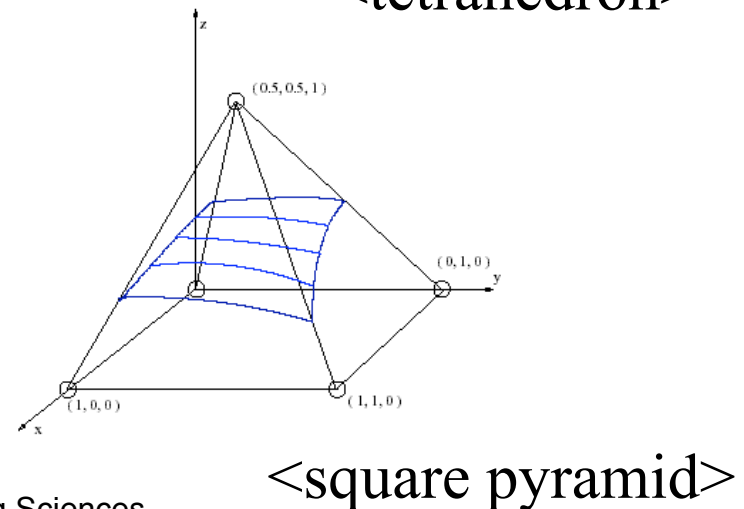
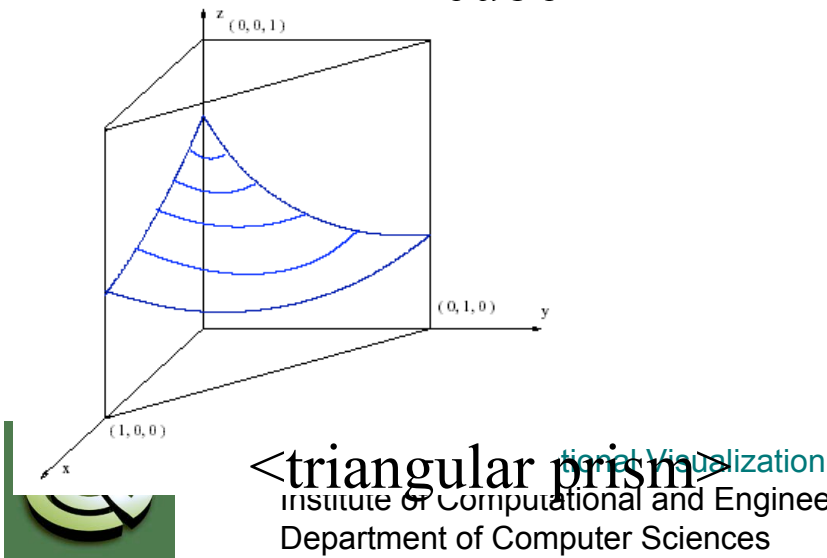
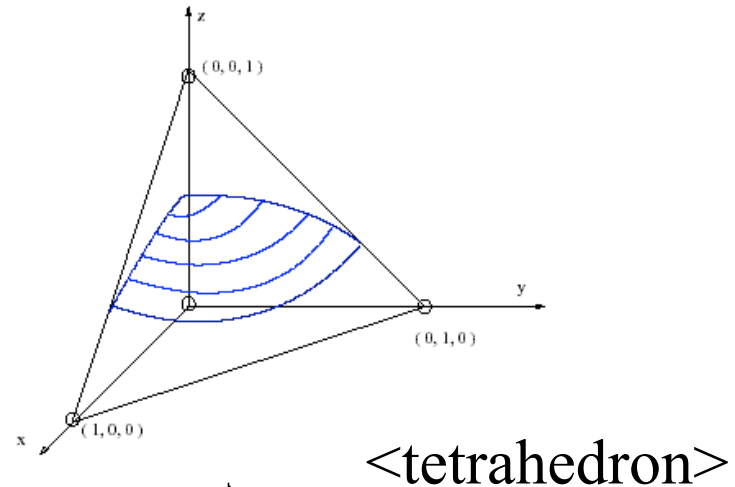
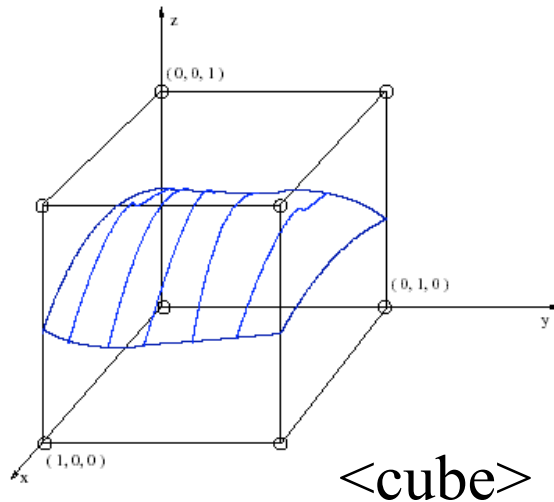
# Dual Contouring

- Polygons with better aspect ratio



# Algebraic Patches: Smooth Boundary Elements

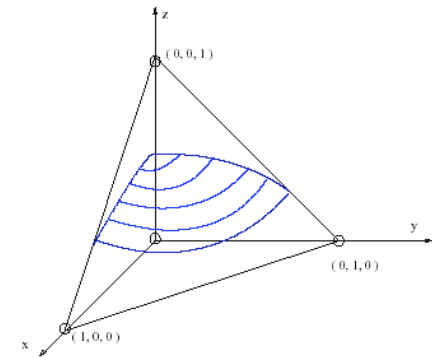
- Implicit form of Isocontour :  $f(x,y,z) = w$



# A-Patches

- Given tetrahedron vertices  $p_i=(x_i,y_i,z_i)$ ,  $i=1,2,3,4$ ,  $\alpha$  is barycentric coordinates of  $p=(x,y,z)$  :

$$\begin{bmatrix} x \\ y \\ z \\ 1 \end{bmatrix} = \begin{bmatrix} x_1 & x_2 & x_3 & x_4 \\ y_1 & y_2 & y_3 & y_4 \\ z_1 & z_2 & z_3 & z_4 \\ 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \alpha_3 \\ \alpha_4 \end{bmatrix}$$



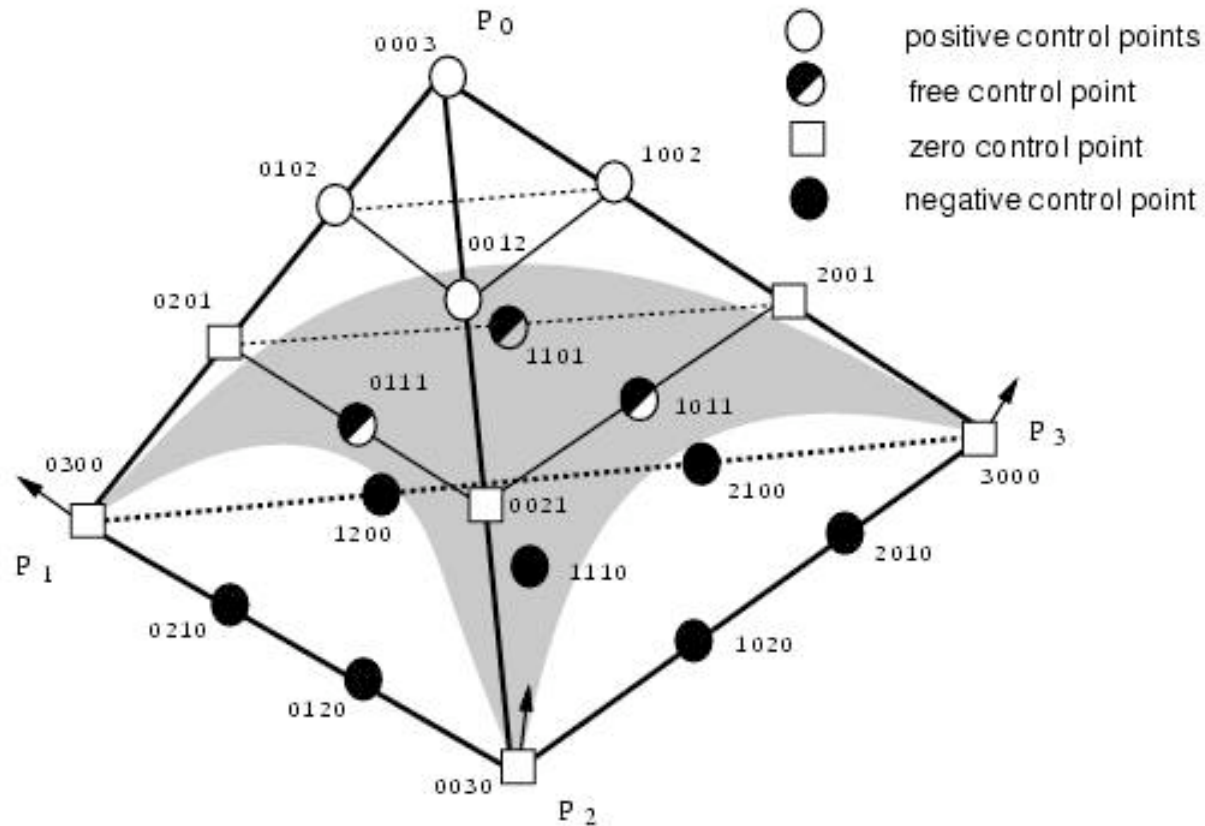
- function  $f(p)$  of degree  $n$  can be expressed in Bernstein-Bezier form :

$$f(p) = \sum_{|\lambda|=n} b_\lambda B_\lambda^n(\alpha), \quad \lambda \in \mathbb{Z}_+^4 \quad B_\lambda^n(\alpha) = \frac{n!}{\lambda_1! \lambda_2! \lambda_3! \lambda_4!} \alpha_1^{\lambda_1} \alpha_2^{\lambda_2} \alpha_3^{\lambda_3} \alpha_4^{\lambda_4}$$

- Algebraic surface patch(A-patch) within the tet is defined as  $f(p)=0$ .



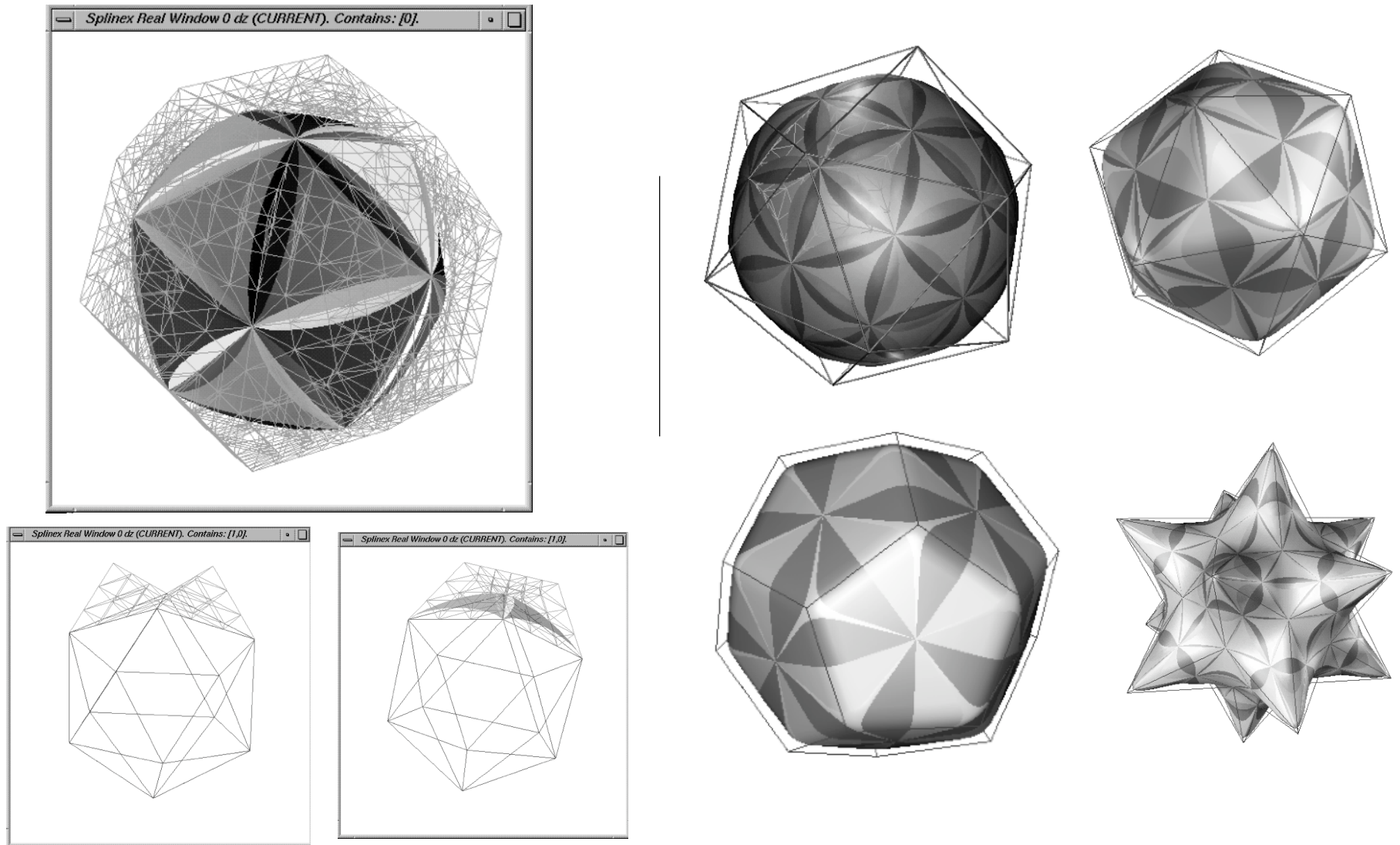
## A-patch Surface ( $C^1$ ) Interpolant



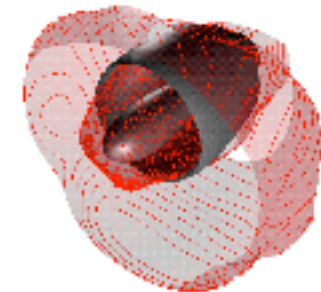
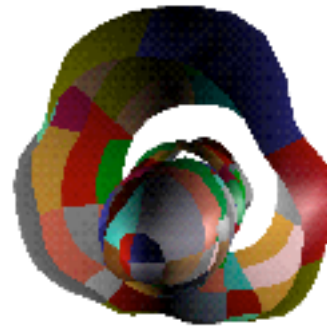
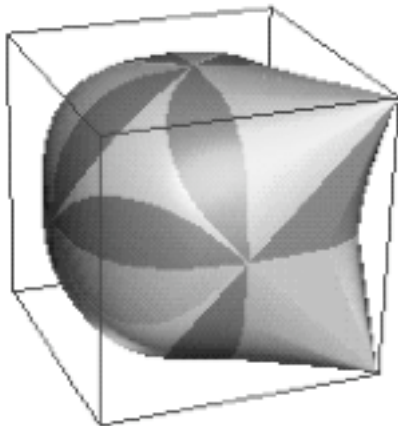
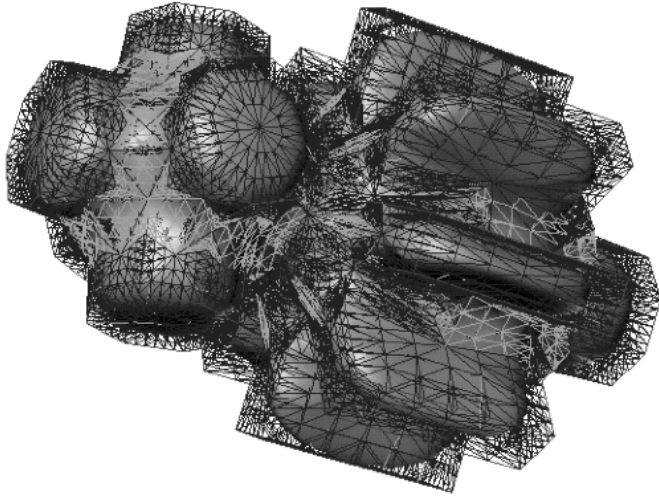
- An implicit single-sheeted interpolant over a tetrahedron



# A-patch Contouring

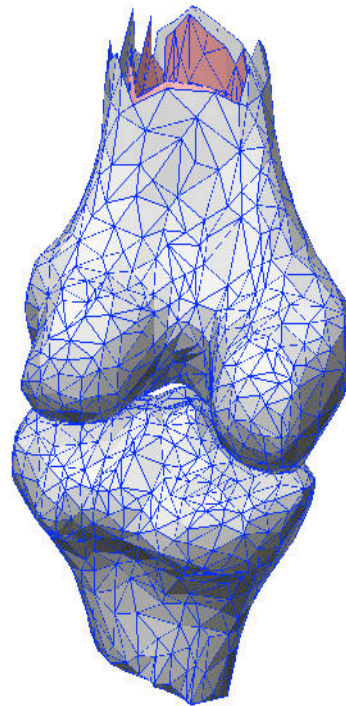


# A-patch Contouring



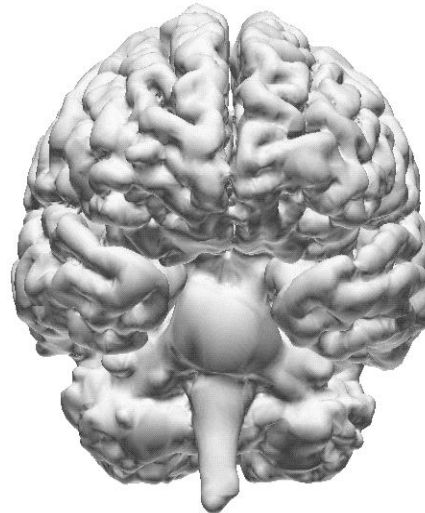
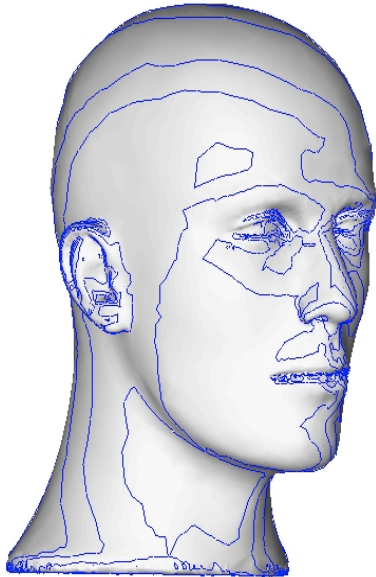
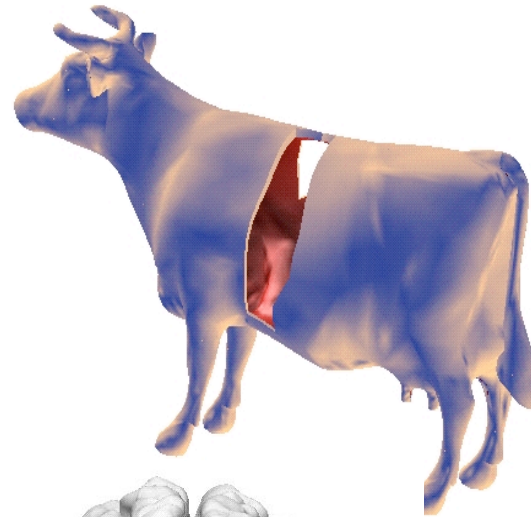
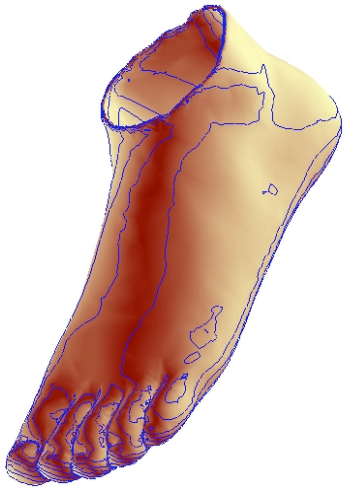


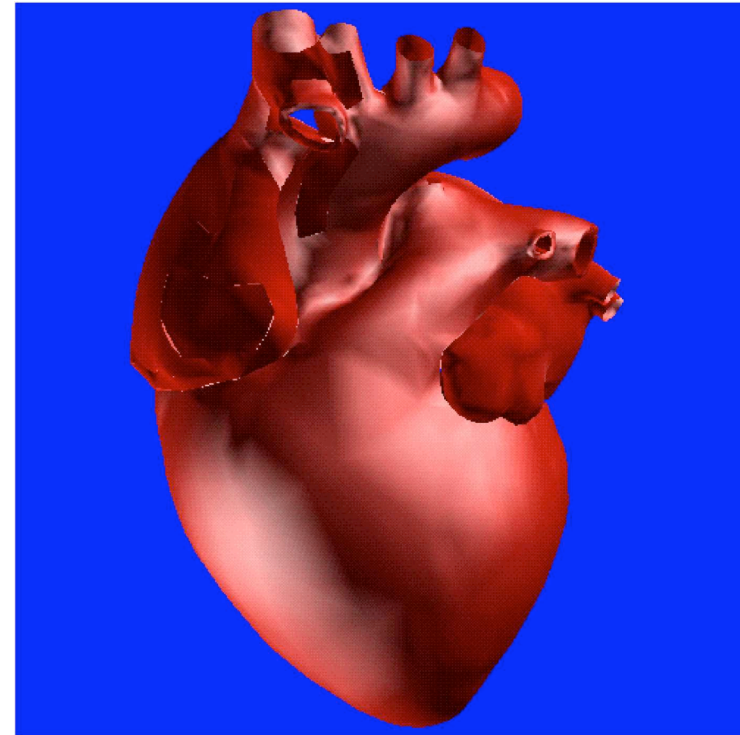
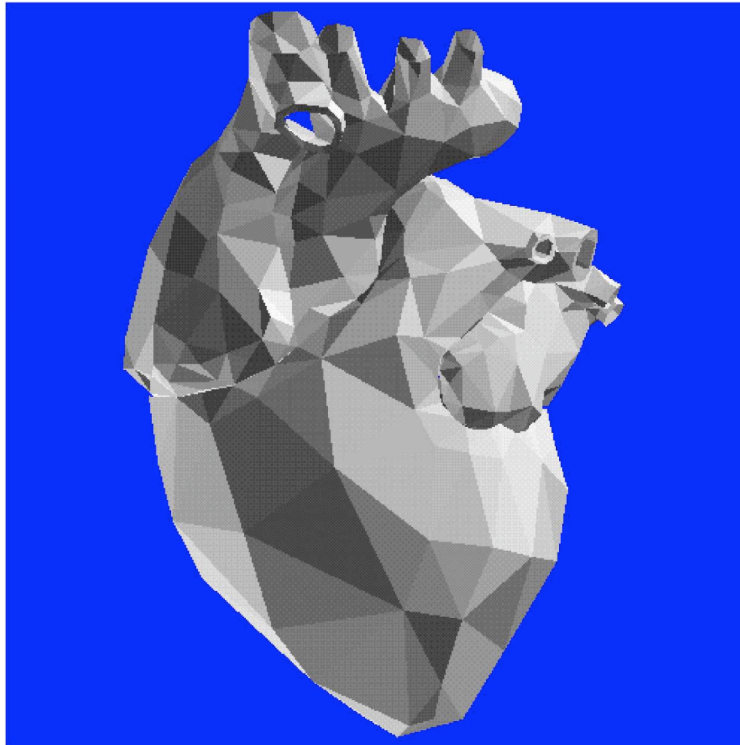
## Finite Elements from Images





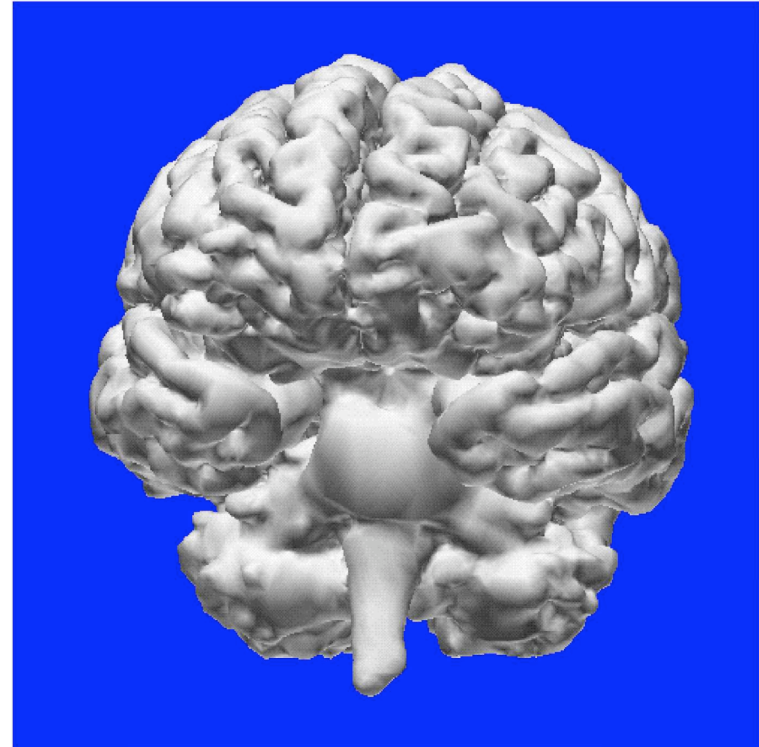
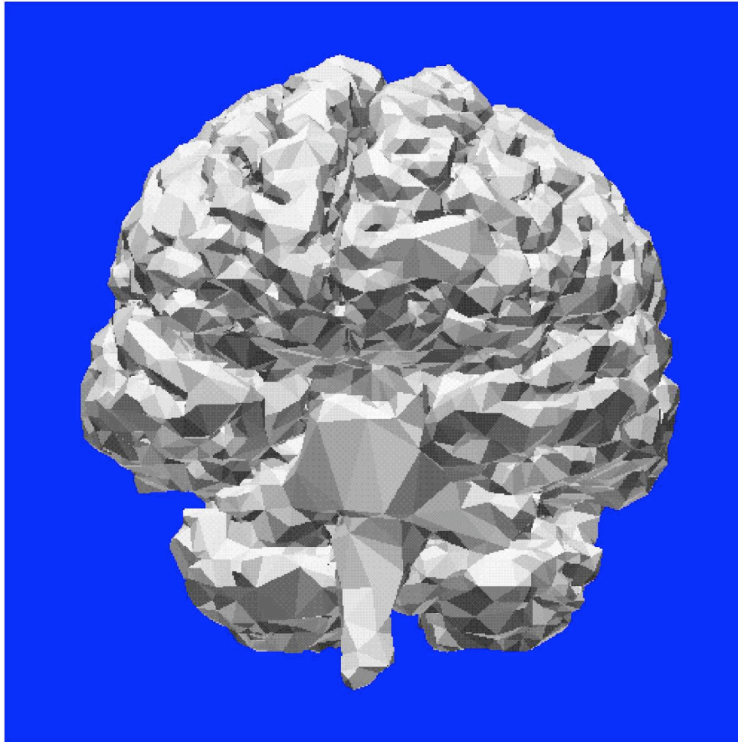
# Examples with Shell Finite Elements





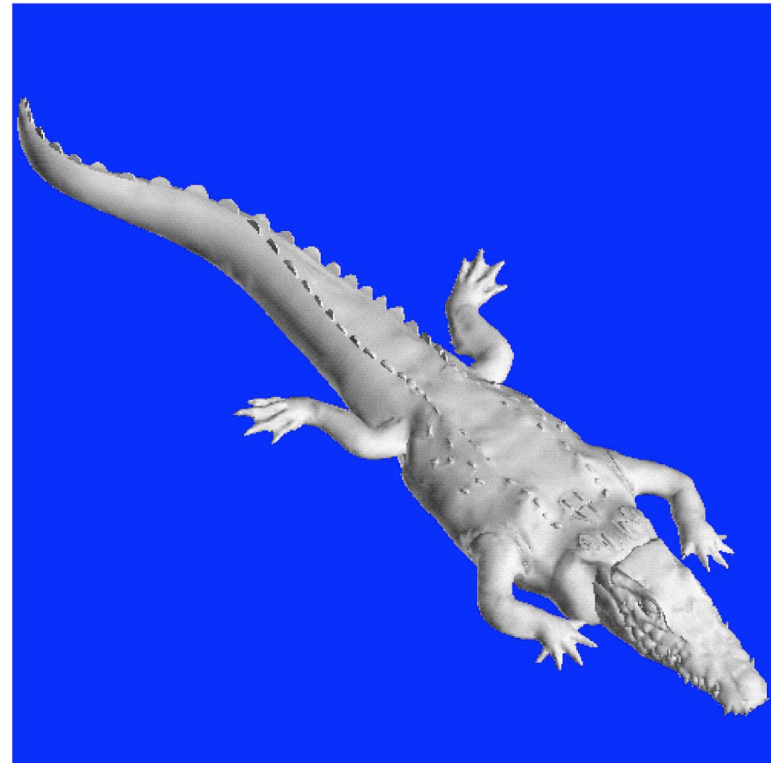
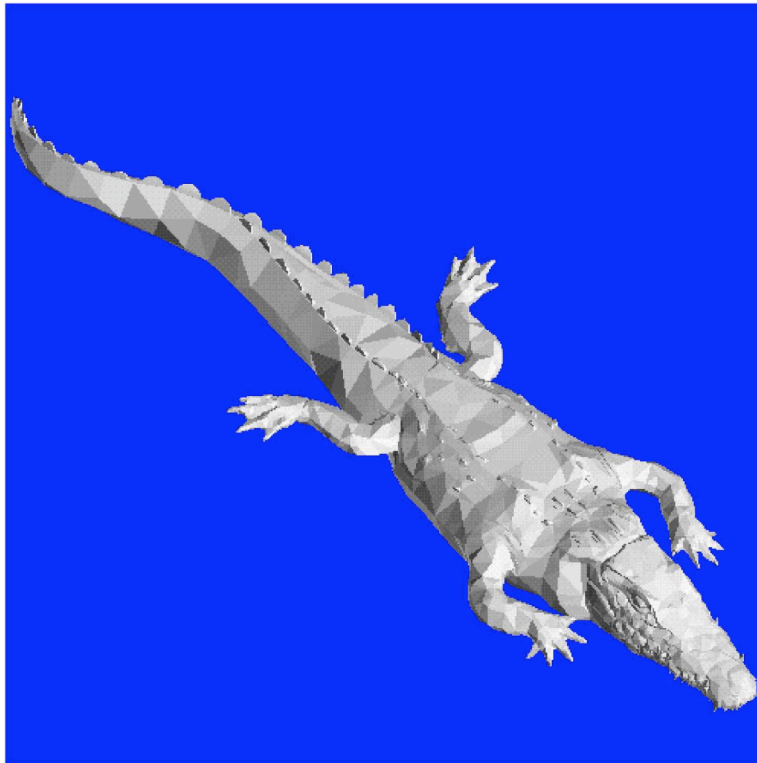
Adaptive feature of the reconstruction: The flat parts use less patches than the curved parts





Adaptive feature of the reconstruction: The flat parts use less patches than the curved parts





Capturing detail structures.



Center for Computational Visualization  
Institute of Computational and Engineering Sciences  
Department of Computer Sciences

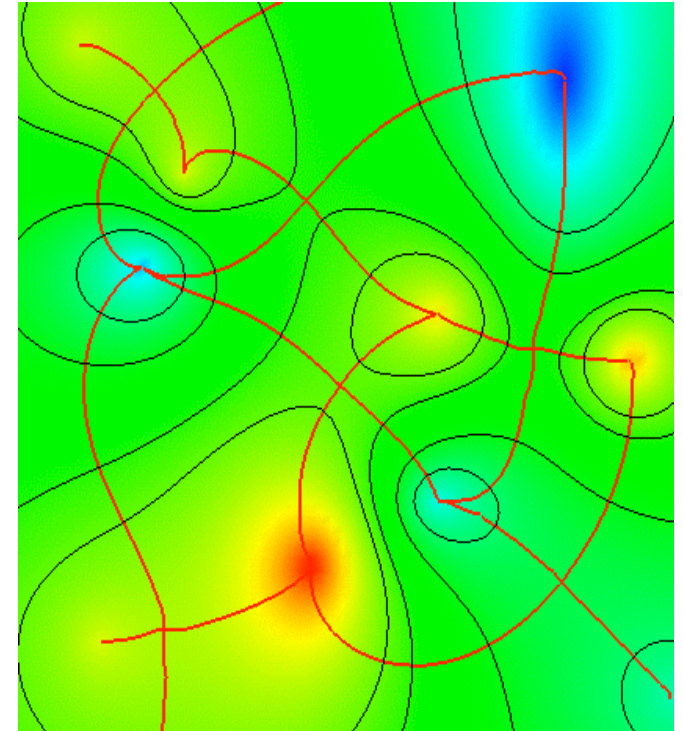
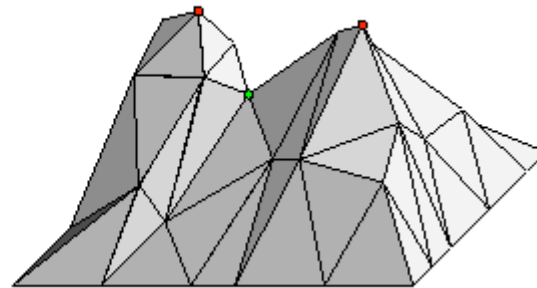
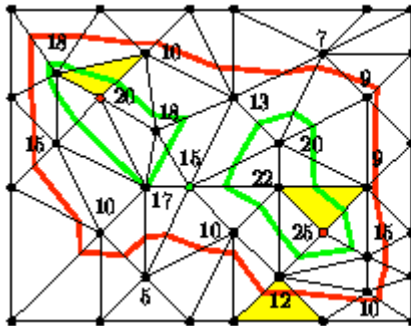
University of Texas at Austin

October 2007



# Interactive Isocontour Queries

- Input:
  - Scalar Field  $F$  defined on a mesh
  - Multiple Isovalues  $w$  in unpredictable order
- Output (for each isovalue  $w$ ):  
Contour  $C(w) = \{x \mid F(x) = w\}$



# Related Work

		Search Space	
		Geometric	Value
<b>Contouring Strategy</b>	<b>Cell by Cell</b>	<b>Loreson/Cline (Marching Cubes)</b> <b>Wilhelms/Van Gelder (octree)</b>	<b>Giles/Haimes (min-sorted ranges)</b> <b>Shen/Livnat/Johnson/Hansen (LxL lattice)</b> <b>Gallagher (span decomposed into buckets)</b> <b>Shen/Johnson (hierachical min-max ranges)</b> <b>Cignoni/Montani/Puppo/Scopigno</b> <b>Livnat/Shen/Johnson (kd-tree)</b>
	<b>Mesh Propagation</b>	<b>Howie/Blake (propagation)</b> <b>Itoh/Koyamada (extrema graph)</b> <b>Itoh/Yamaguchi/Koyamada (volume thinnig)</b>	<b>van Kreveld</b>  <b>Bajaj/Pascucci/Schikore</b>  <b>van Kreveld /van Oostrum/Bajaj/ Pascucci/Schikore</b>



# Isocontour Query Problem

Lower Bound

Input size  $n$   
Output size  $m$

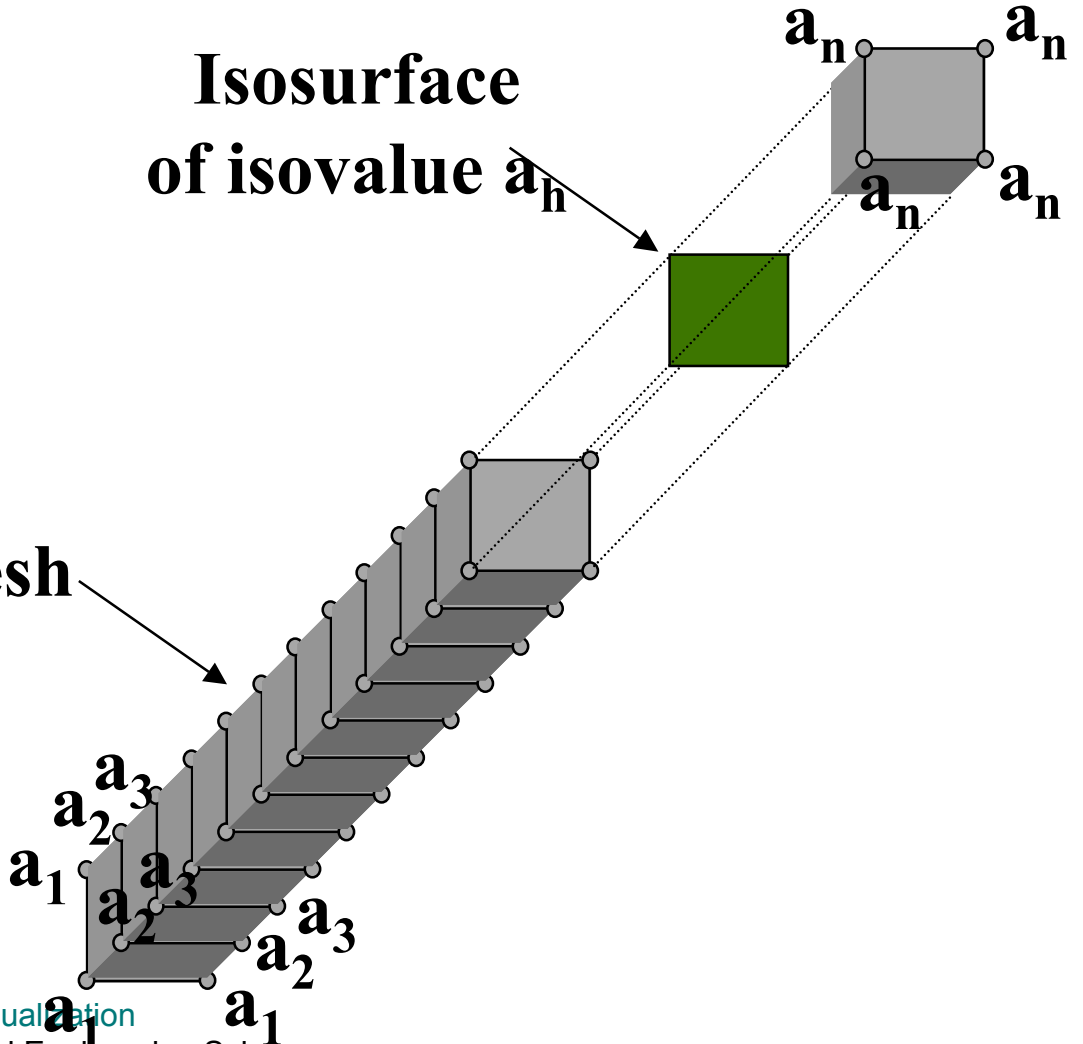
$m + \log(n)$

Isosurface  
of isovalue  $a_h$

Mesh

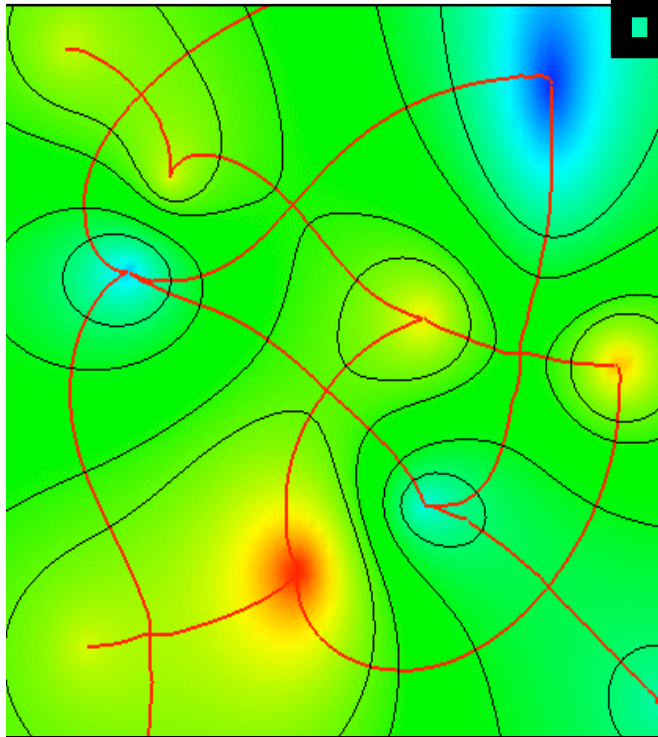
The search for  $a_h$   
takes at least  $\log(n)$

$a_1, a_2, a_3, a_4, \dots, a_n$



# Optimal Single-Resolution Isocontouring

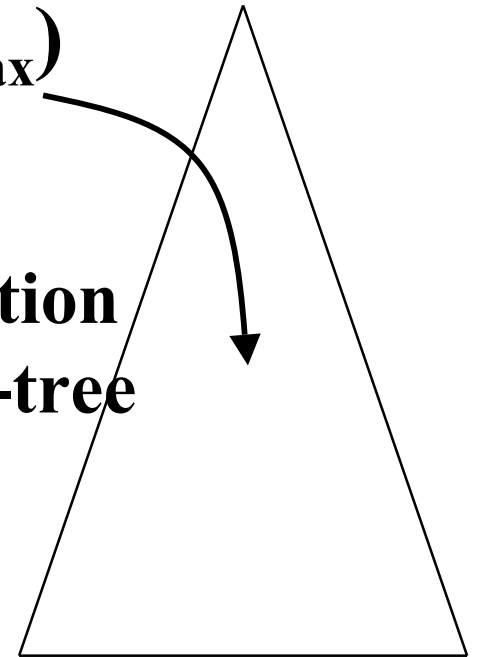
## The basic scheme



Preprocessing:  $(f_{\min}, f_{\max})$

For each cell  $c$  in  $M$

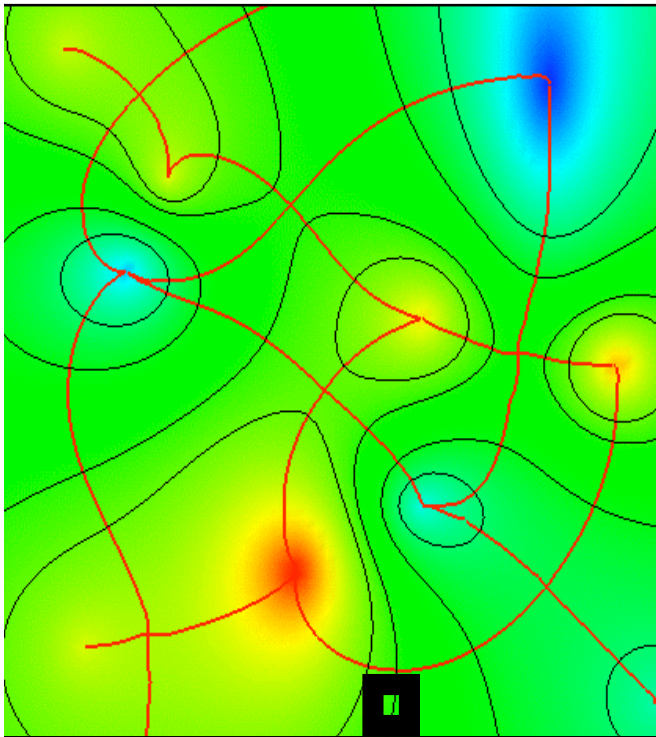
Enter its range of function values into an interval-tree





# Optimal Single-Resolution Isocontouring

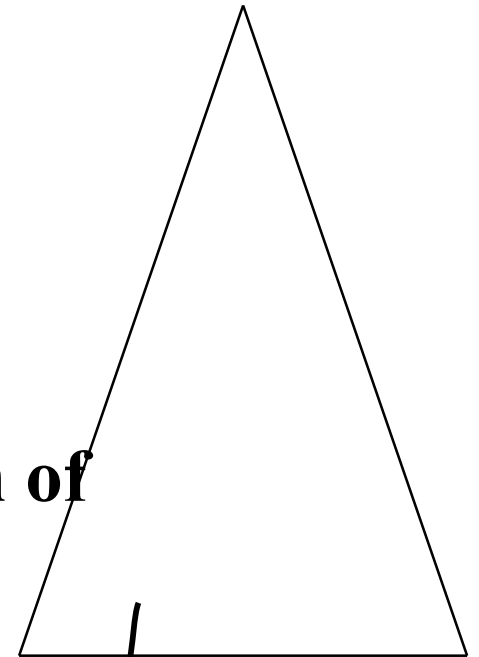
## The basic scheme



**Isocontour query  $W$**

**For each interval  
containing  $W$**

**Compute the portion of  
isocontour in the  
corresponding cell**

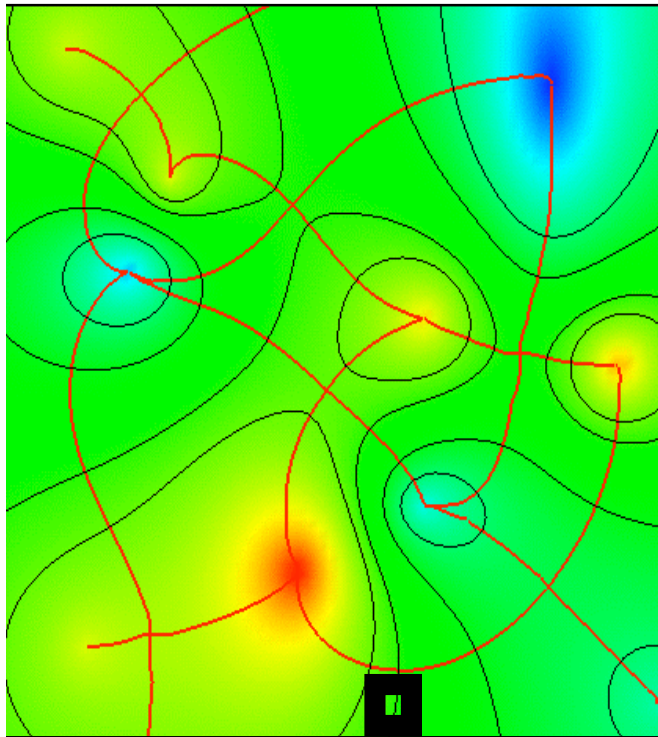


**$(f_{\min}, f_{\max})$**



# Optimal Single-Resolution Isocontouring

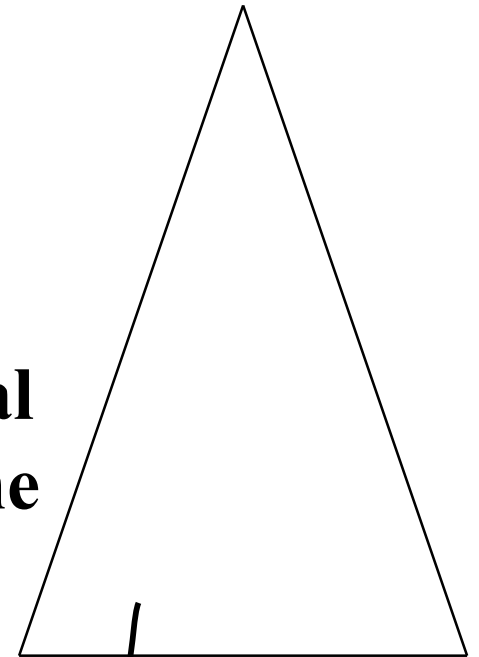
## The basic scheme



Isocontour query  $W$

Complexity:  $m + \log(n)$

Optimal but impractical  
because of the size of the  
interval-tree

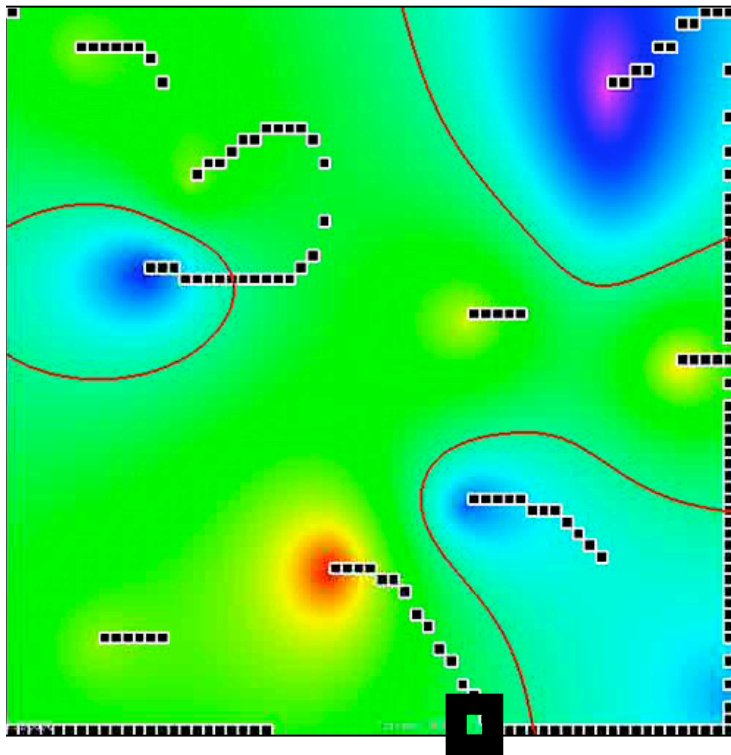


$(f_{\min}, f_{\max})$



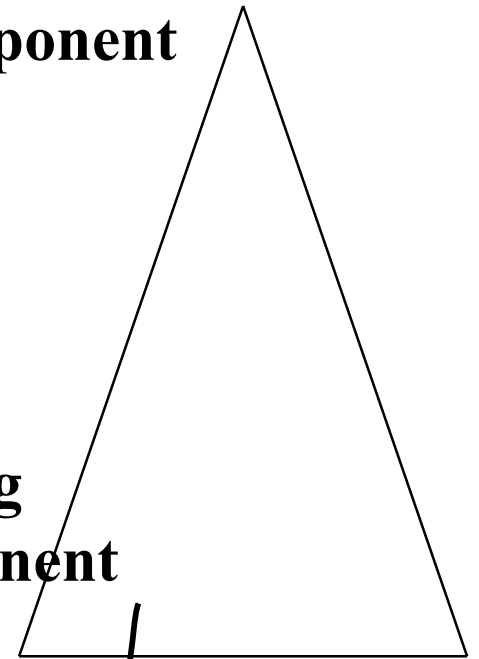
# Optimal Single-Resolution Isocontouring

## Seed Set Optimization



**For each connected component we need only one cell (and then propagate by adjacency in the mesh)**

**Seed Set:  
a set of cells intersecting every connected component of every isocontour**

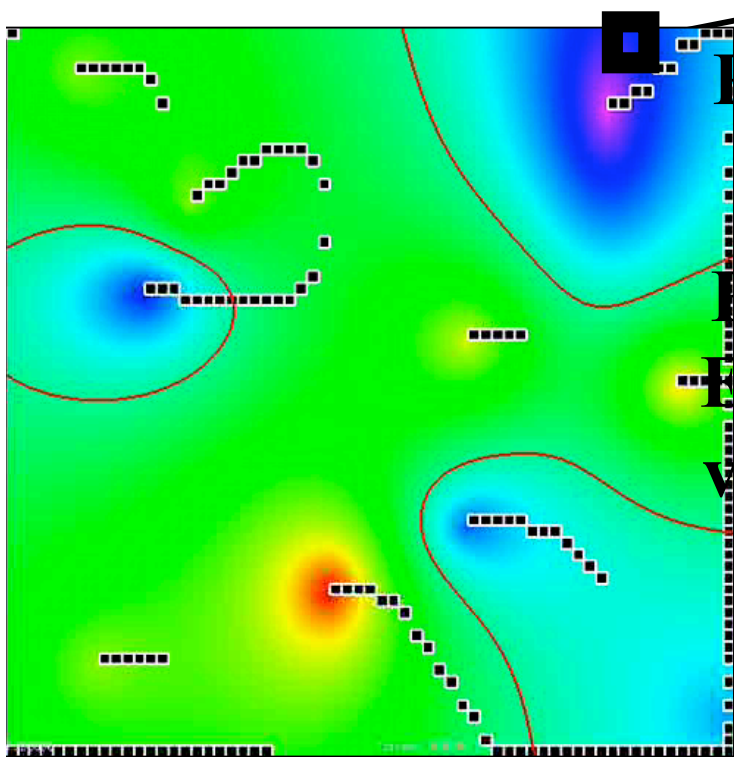


$(f_{\min}, f_{\max})$



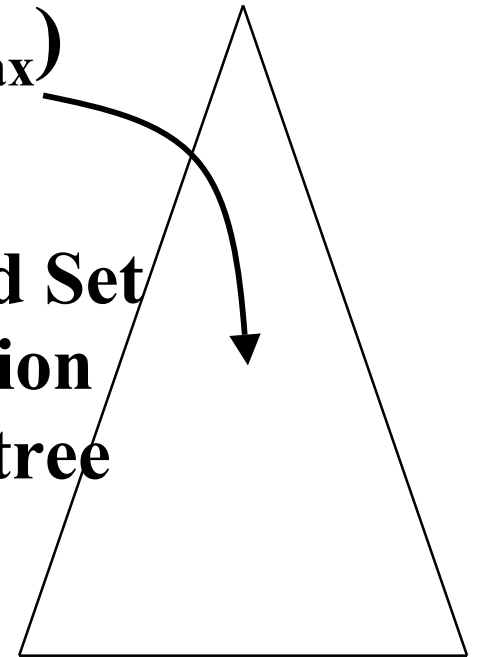
# Optimal Single-Resolution Isocontouring

## The basic scheme



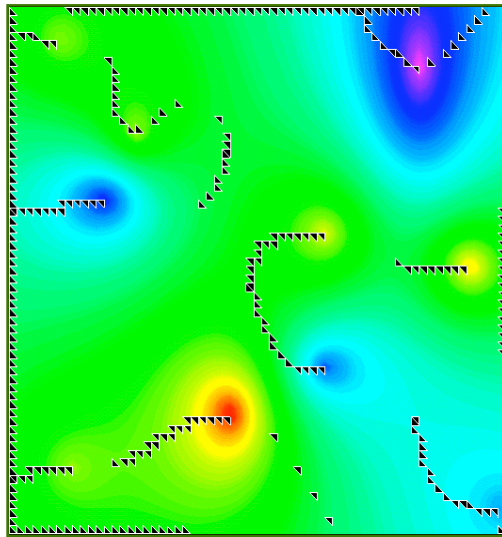
Preprocessing ( $f_{\min}, f_{\max}$ )  
(revised):

For each cell  $c$  in a Seed Set  
Enter its range of function  
values into an interval-tree

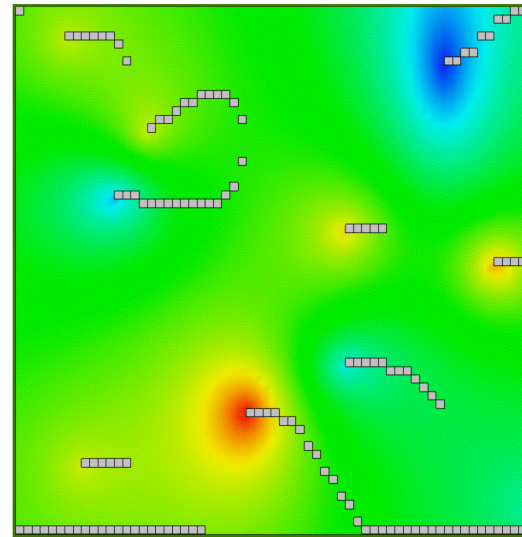


# Optimal Single-Resolution Isocontouring

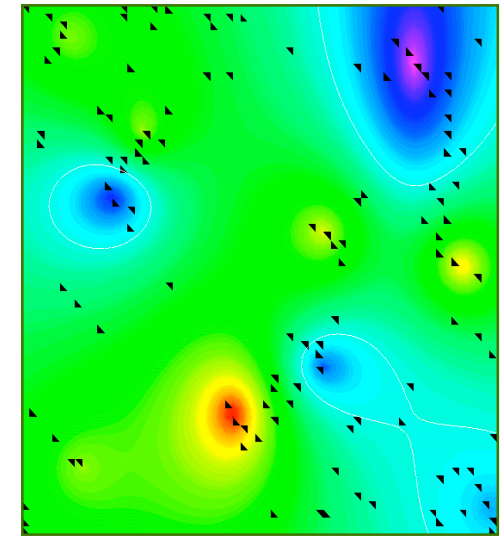
## Seed Set Generation ( $k$ seeds from $n$ cells)




Domain Sweep



Responsibility Propagation



Range Sweep

Time	$O(n)$	$O(n)$	$O(n \log n)$
Space	$O(k)$	$O(k)$	$O(n)$
$k =$	?	?	$2 k_{\min}$
Test	 238 seed cells 0.04 seconds	177 seed cells 0.05 seconds	59 seed cells 1.02 seconds

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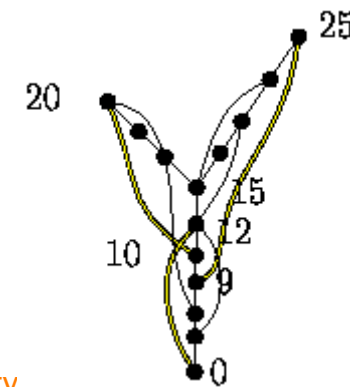
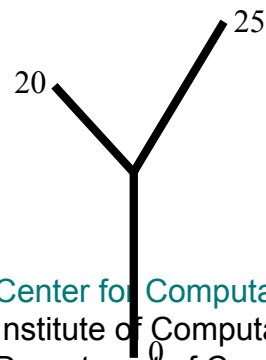
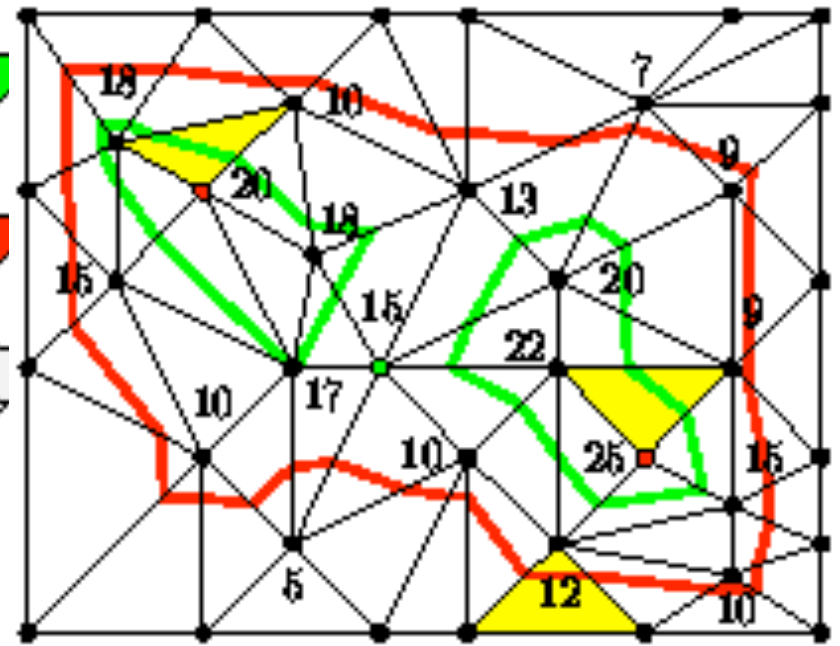
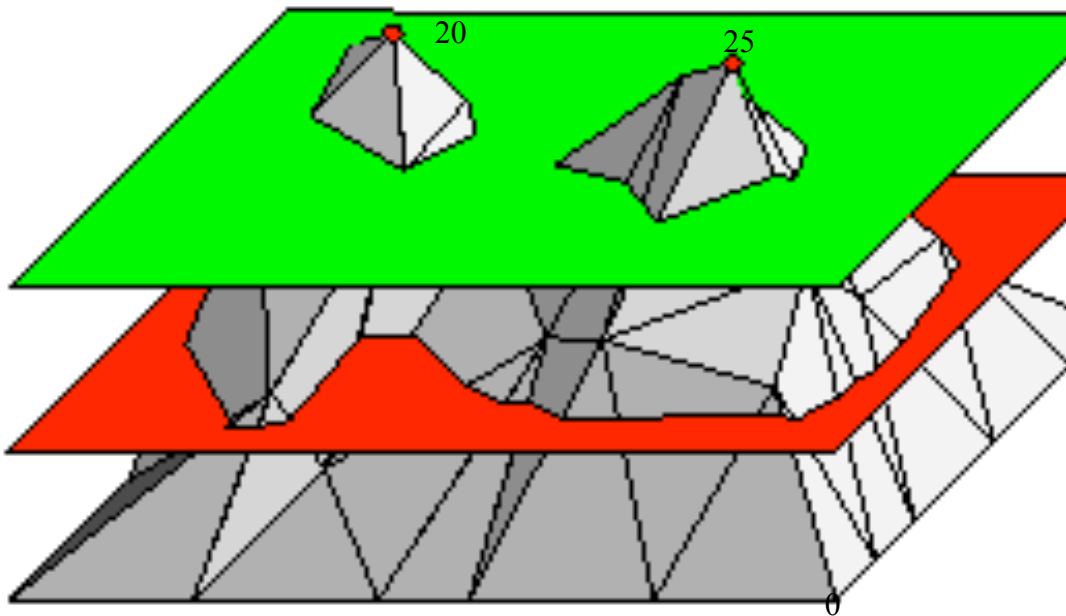
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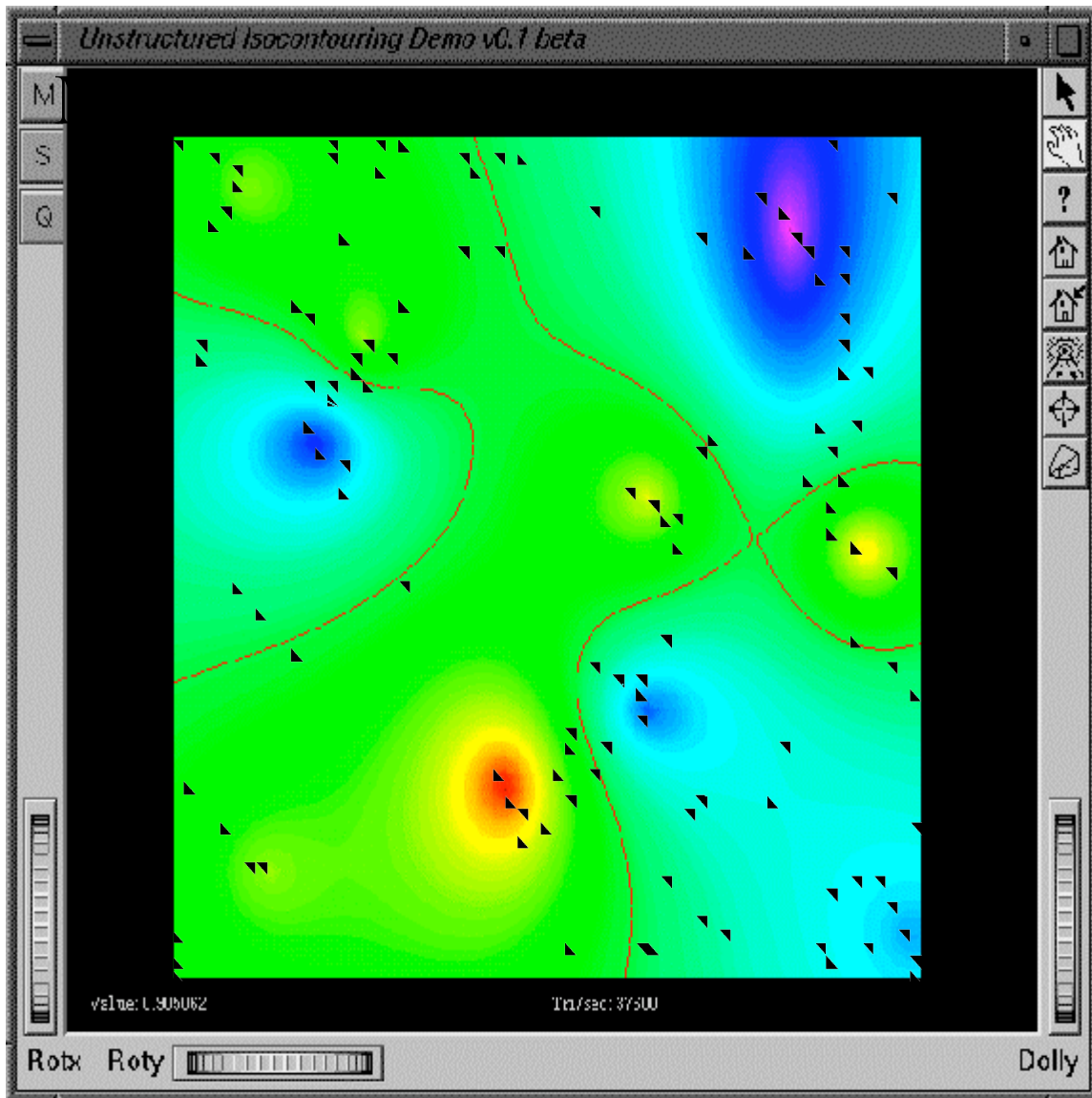


# Optimal Single-Resolution Isocontouring

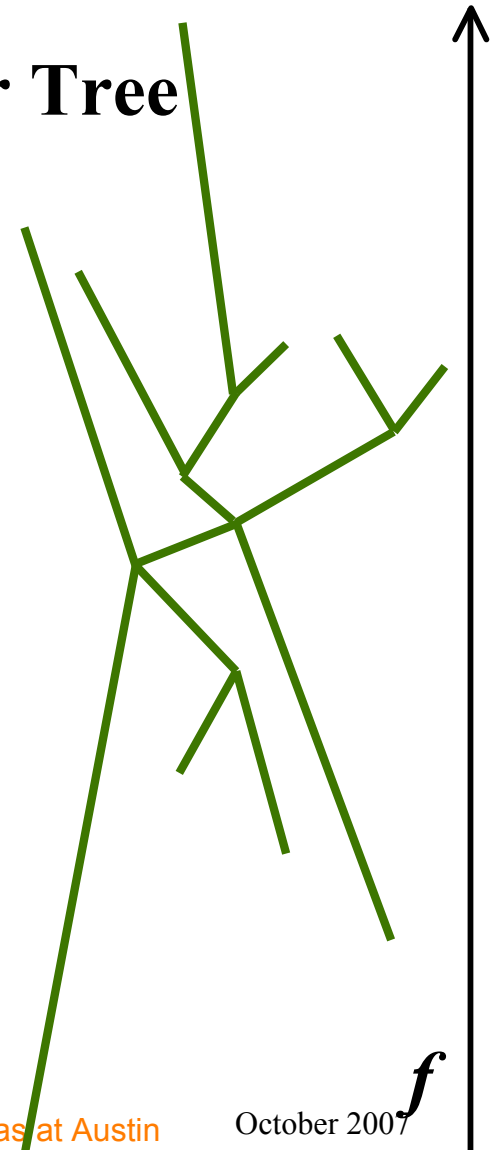
## Contour tree



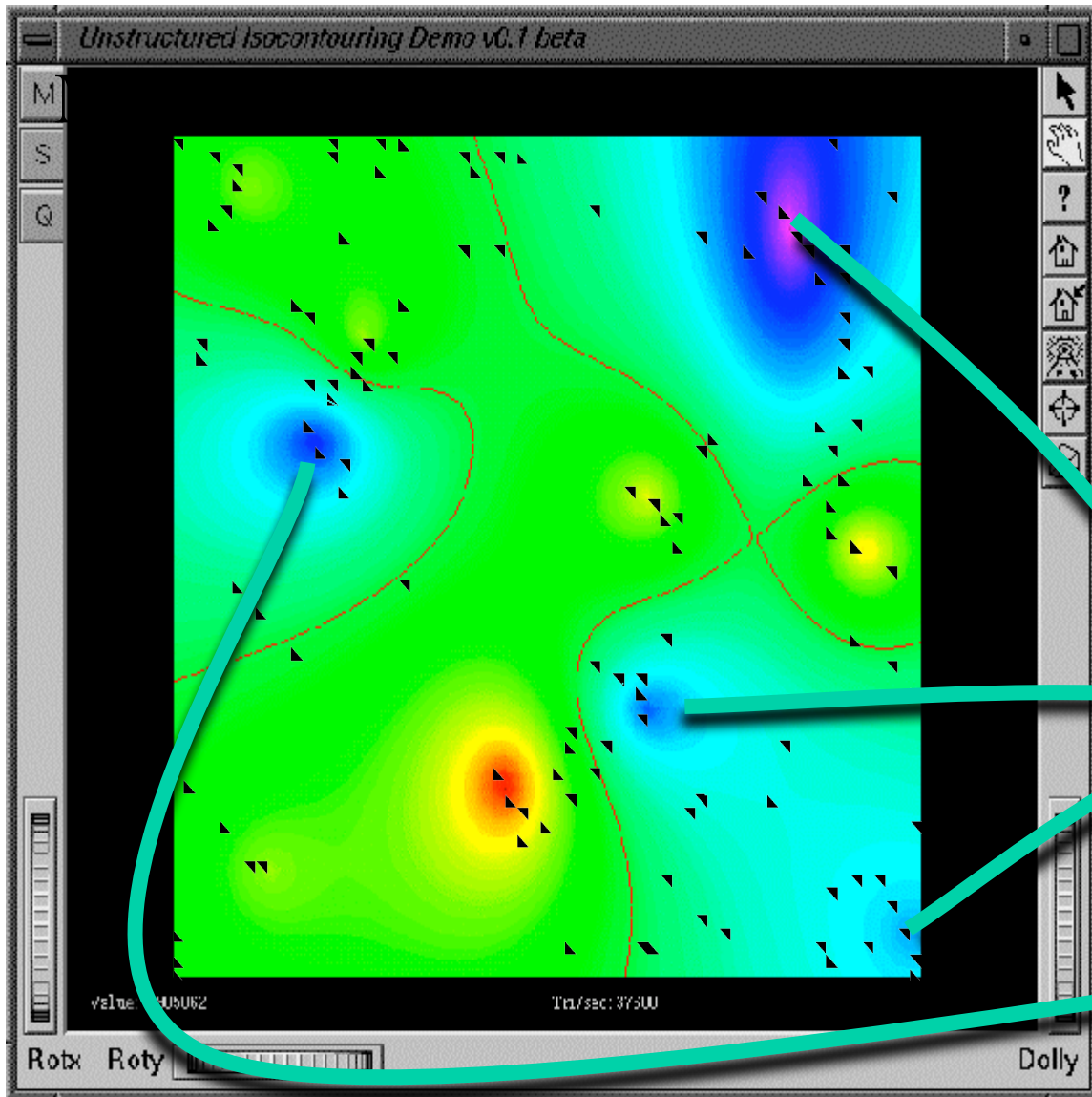
# Optimal Single-Resolution Isocontouring



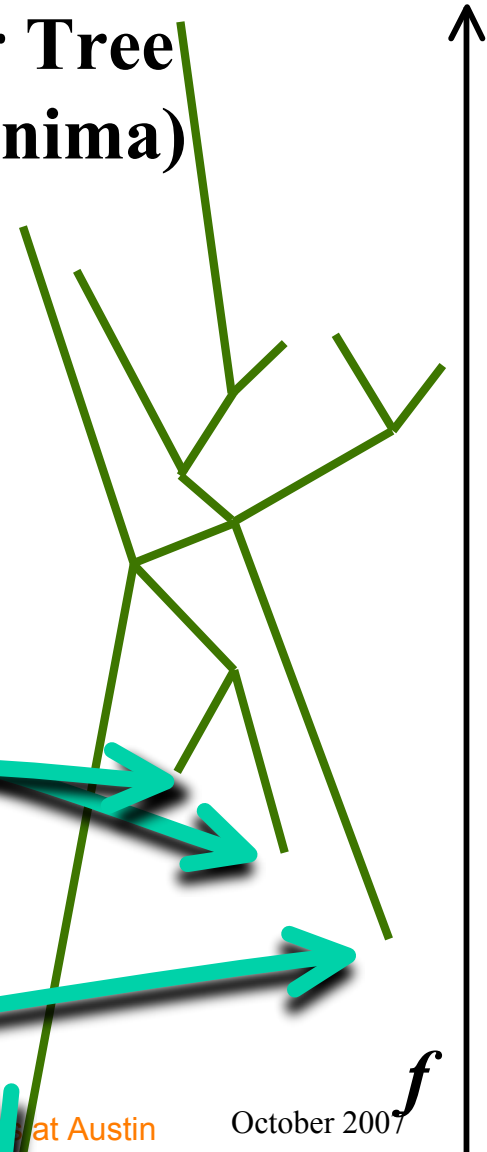
## Contour Tree



# Optimal Single-Resolution Isocontouring

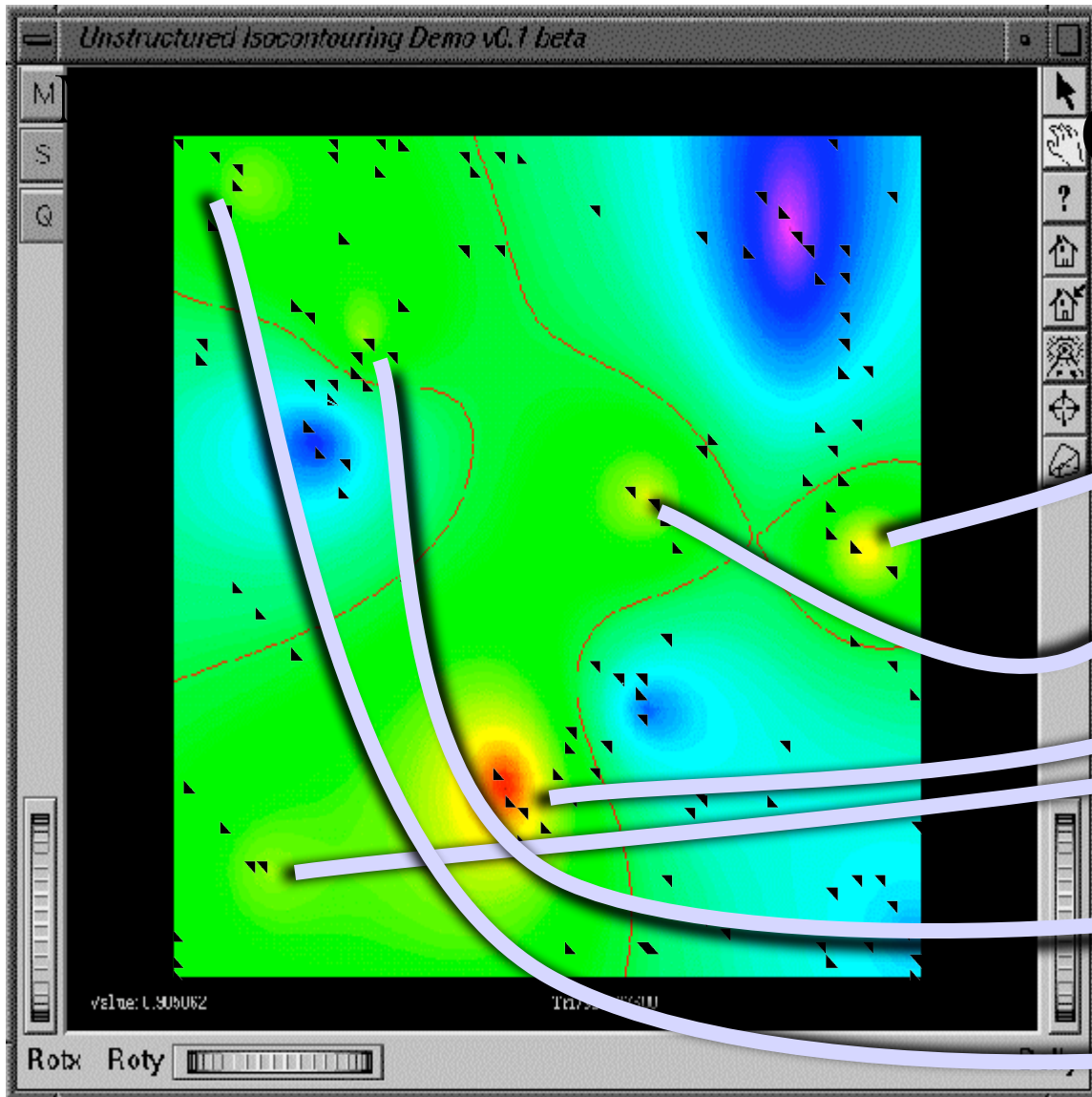


**Contour Tree  
(local minima)**

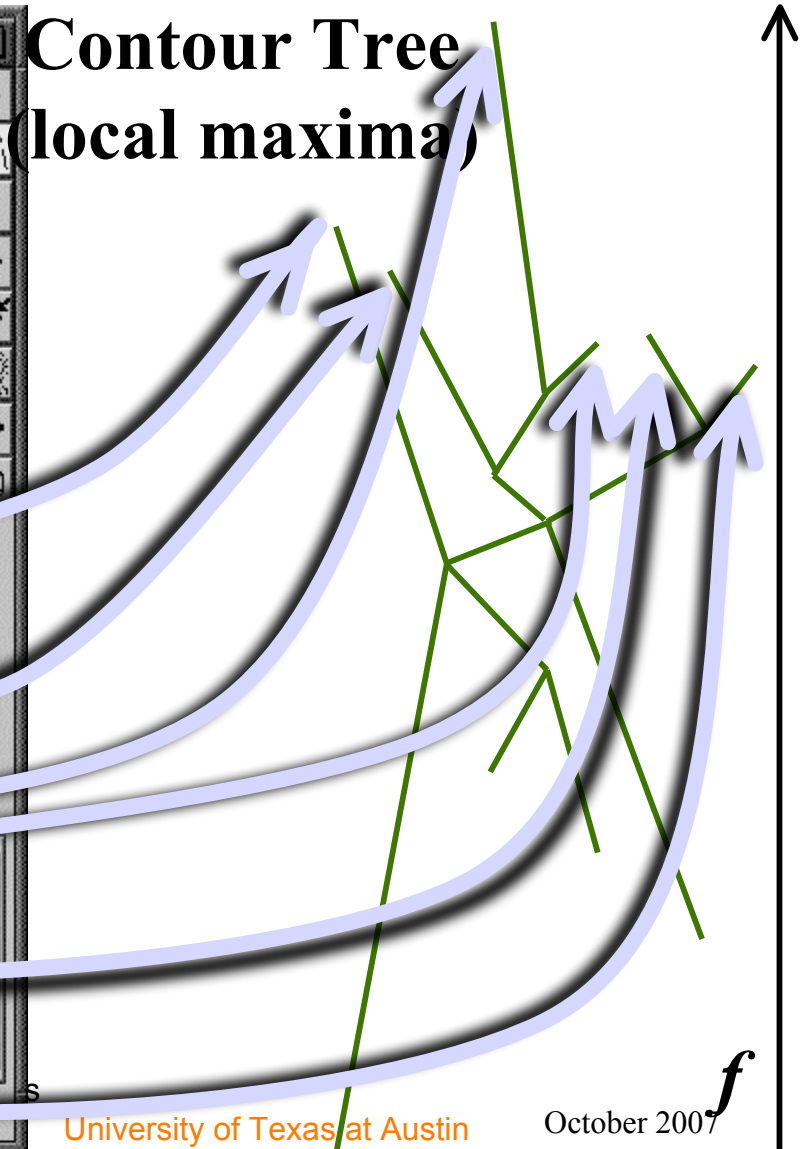




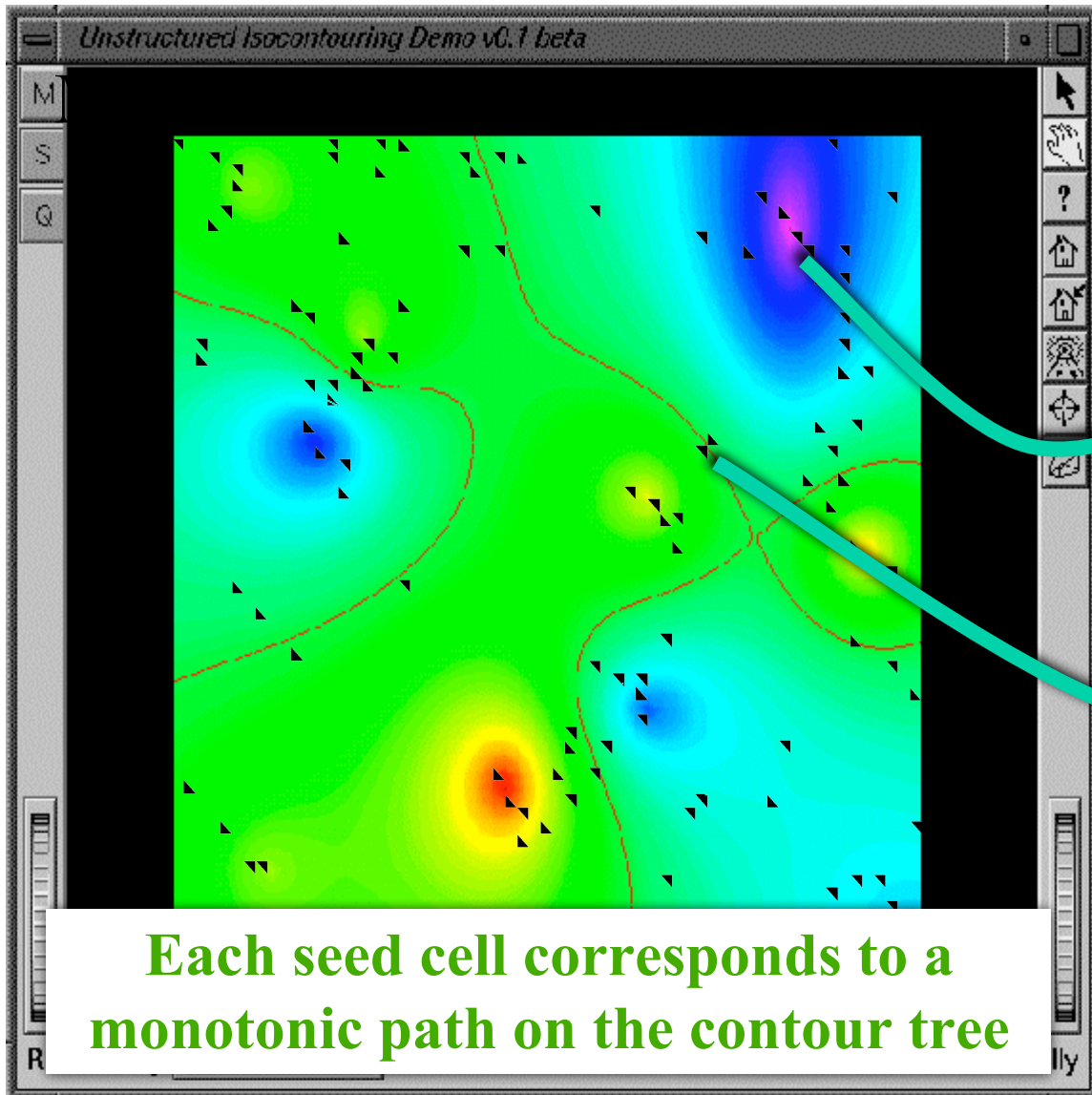
# Optimal Single-Resolution Isocontouring



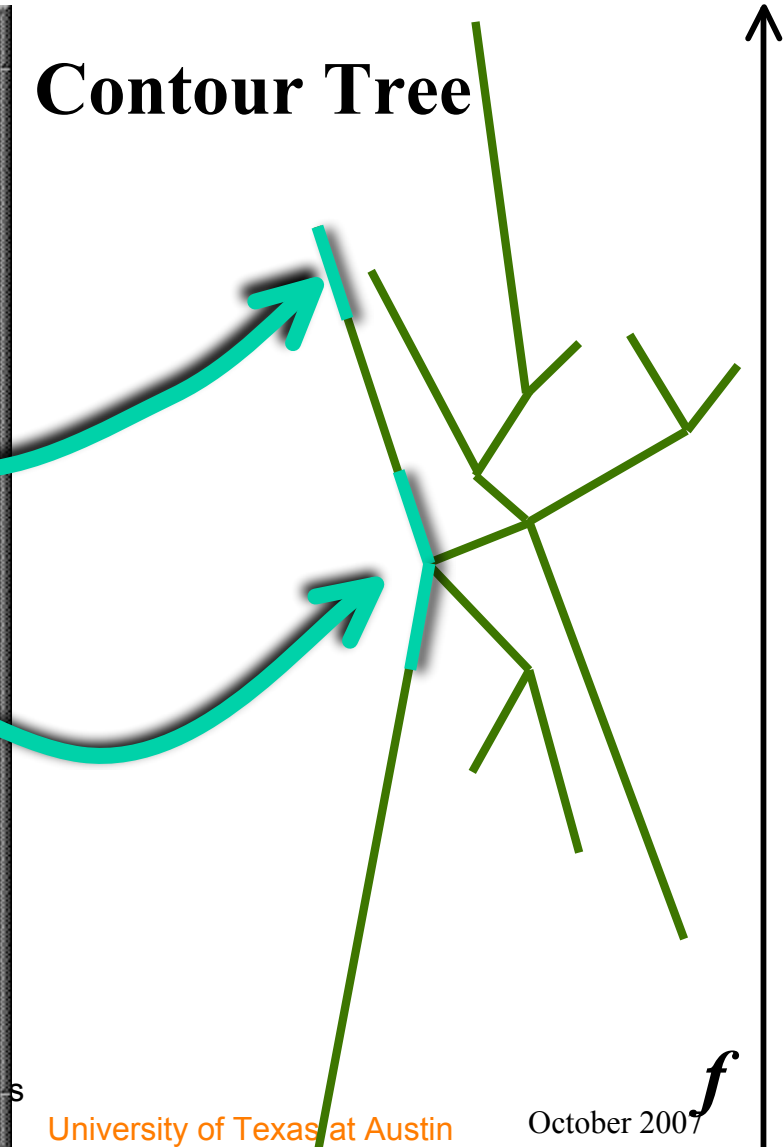
Contour Tree  
(local maxima)



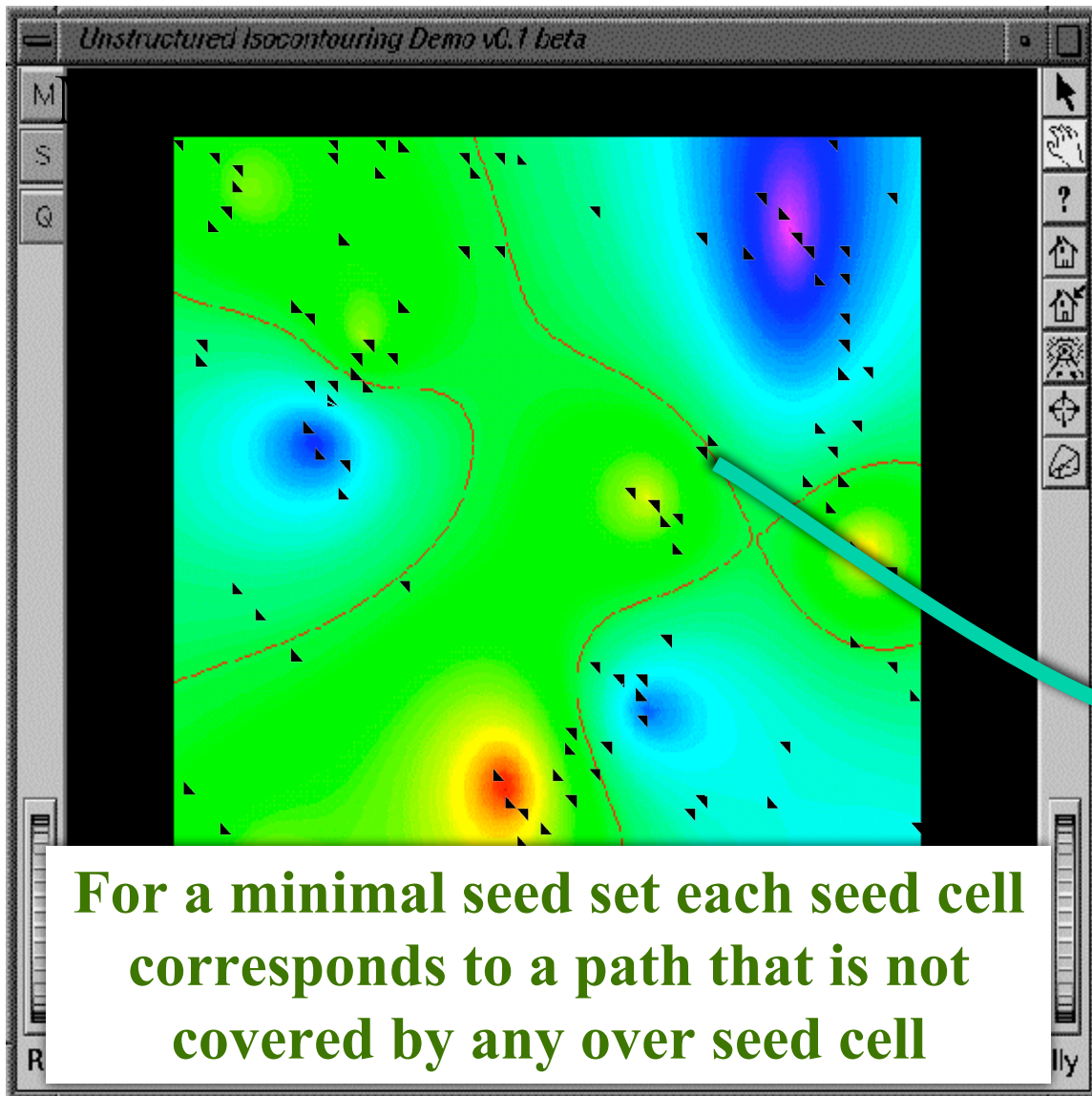
# Optimal Single-Resolution Isocontouring



## Contour Tree



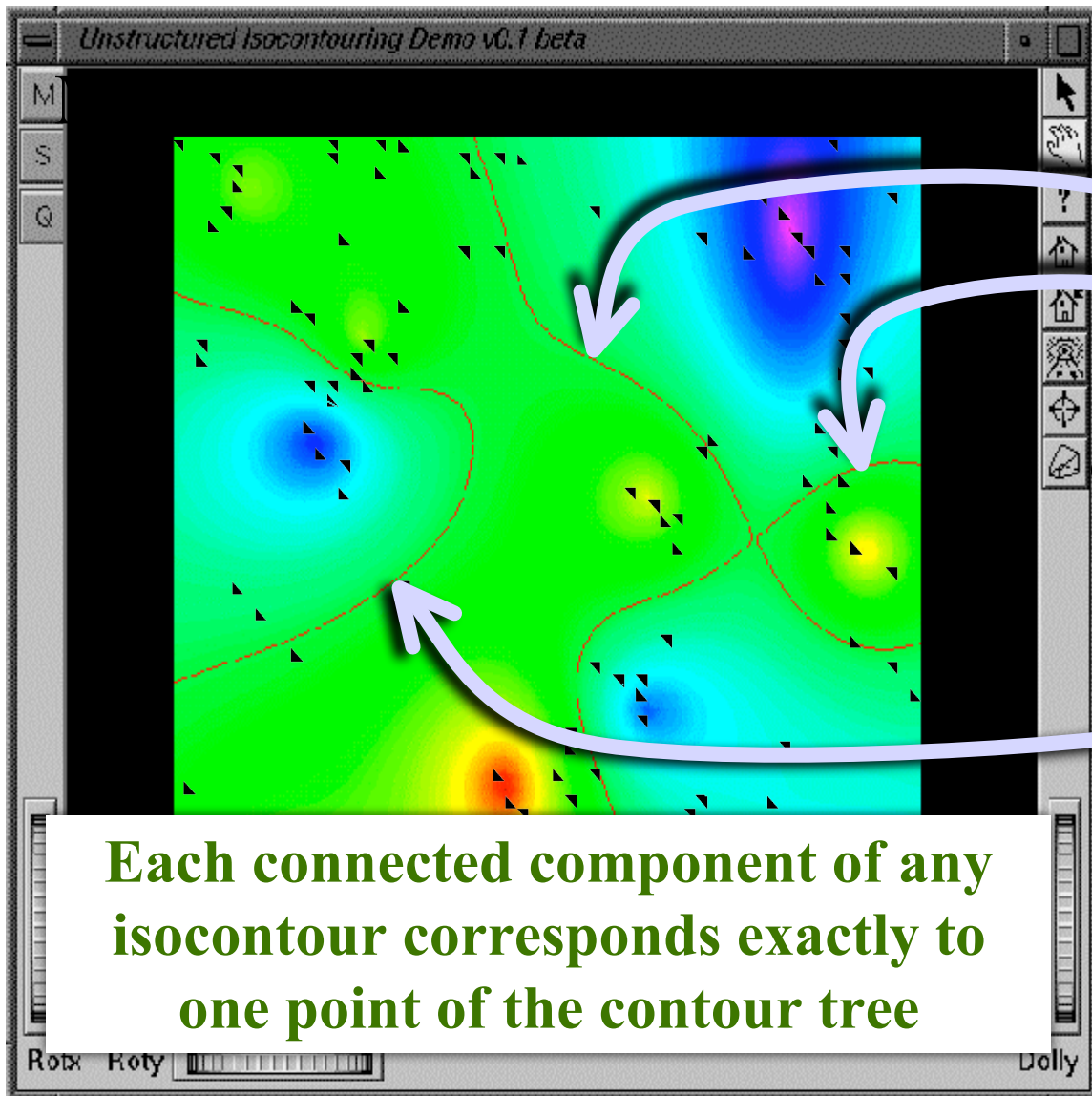
# Optimal Single-Resolution Isocontouring



## Contour Tree

**For a minimal seed set each seed cell corresponds to a path that is not covered by any over seed cell**

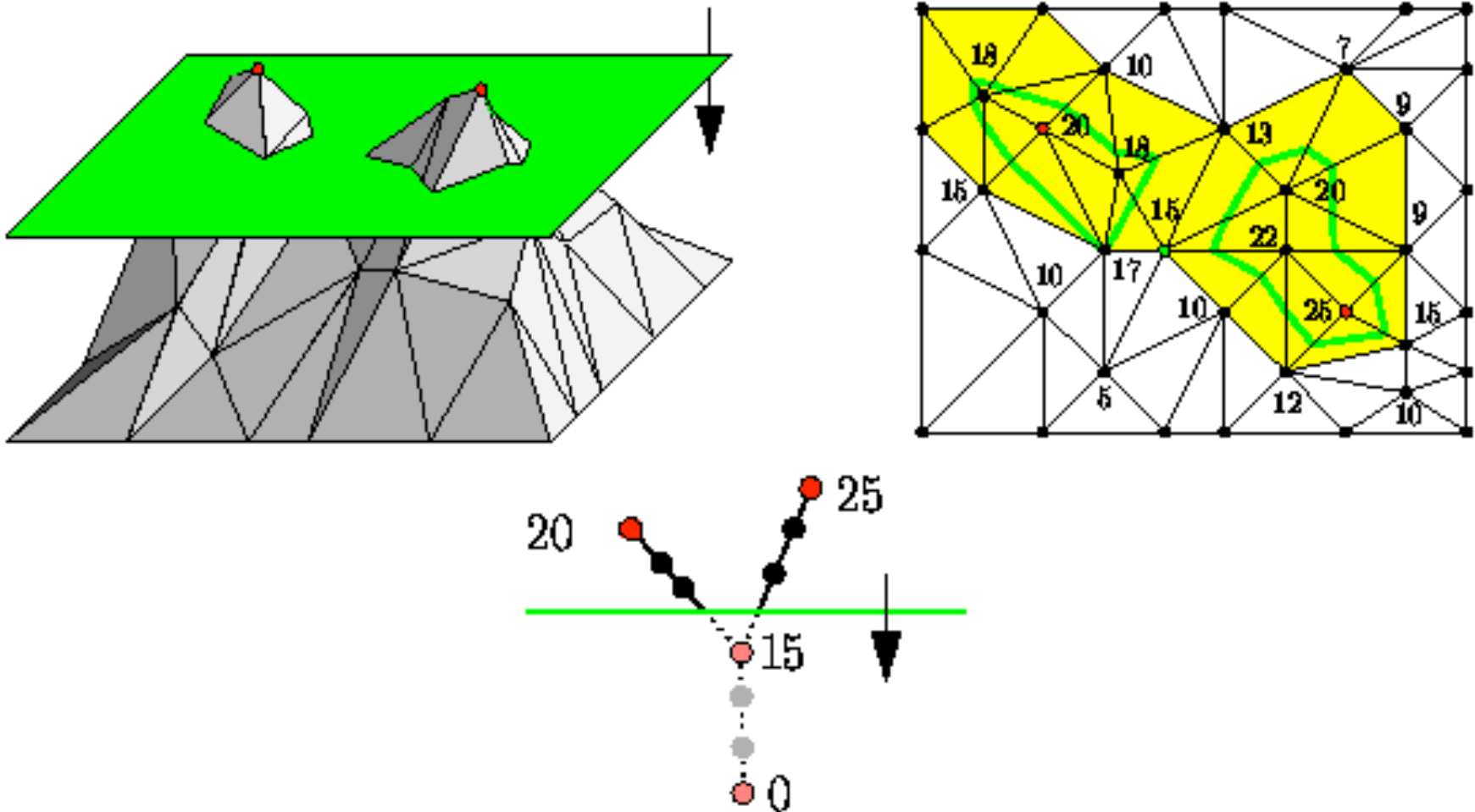
# Optimal Single-Resolution Isocontouring



Contour Tree

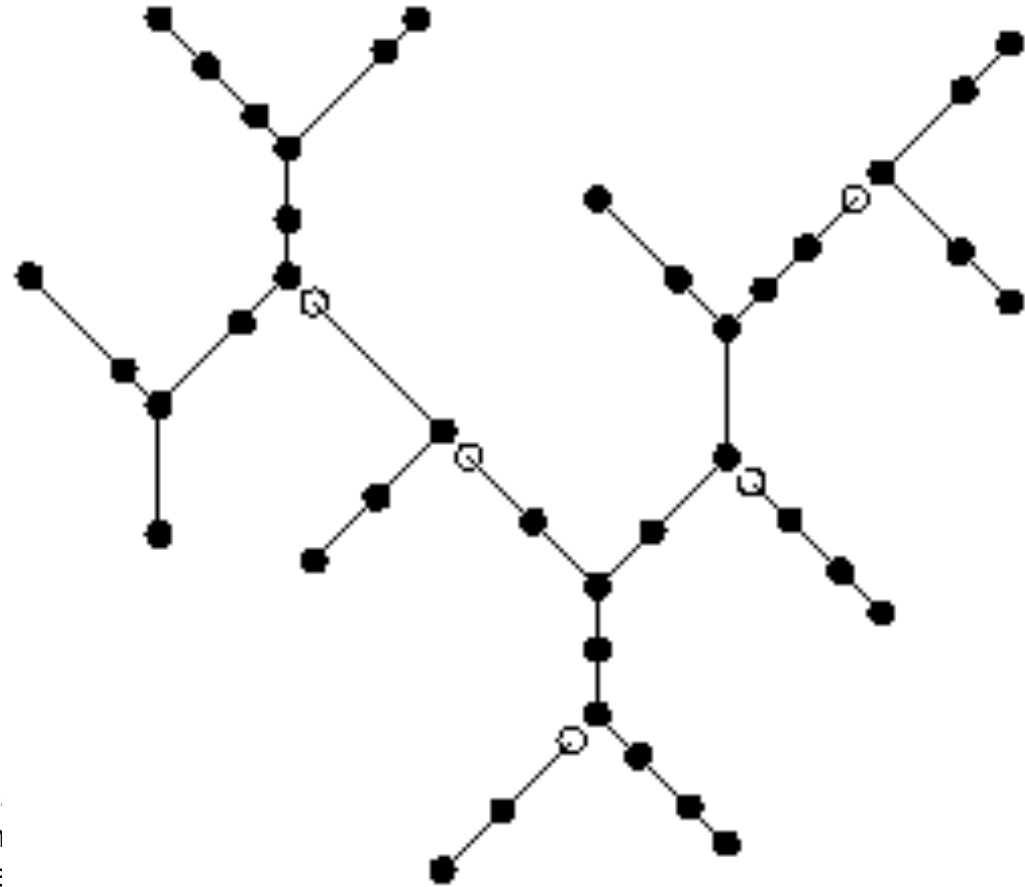
Current  
isovalue

# Optimal Single-Resolution Isocontouring



## Optimal Single-Resolution Isocontouring

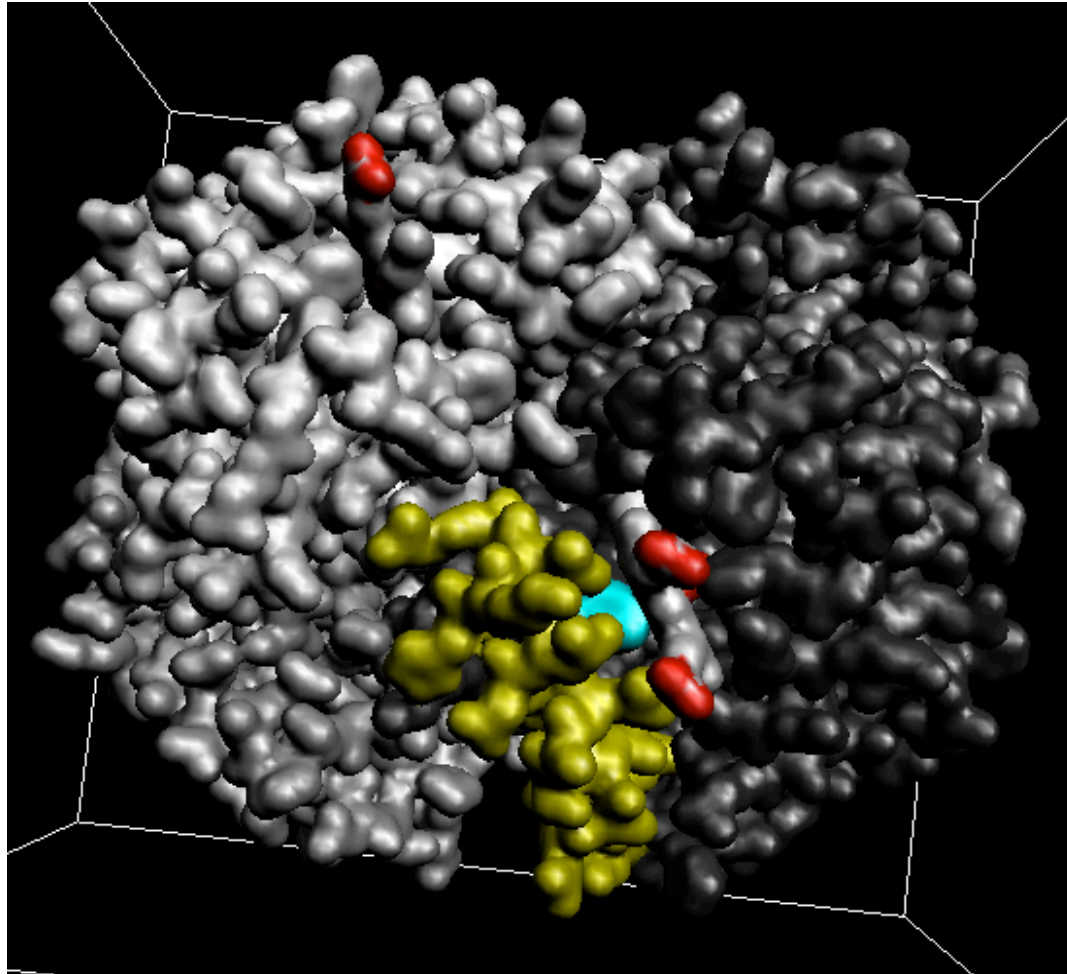
- The number of seeds selected is the minimum plus the number of local minima.

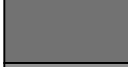


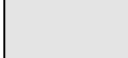







# Structural Analysis

## Contour Spectrum and Contour Tree on Hemoglobin Dynamics



	Subunit A
	Subunit B
	Subunit C
	Subunit D

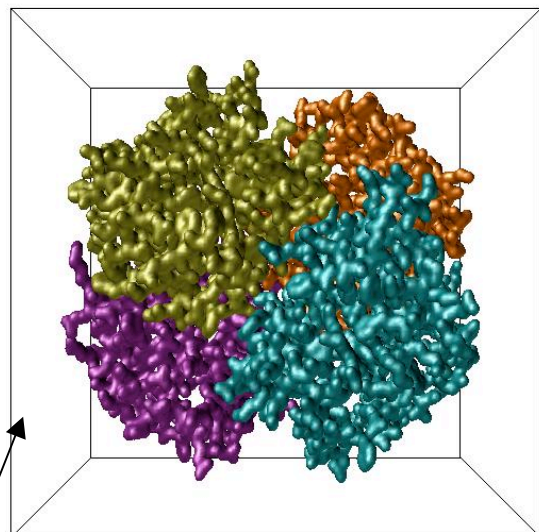
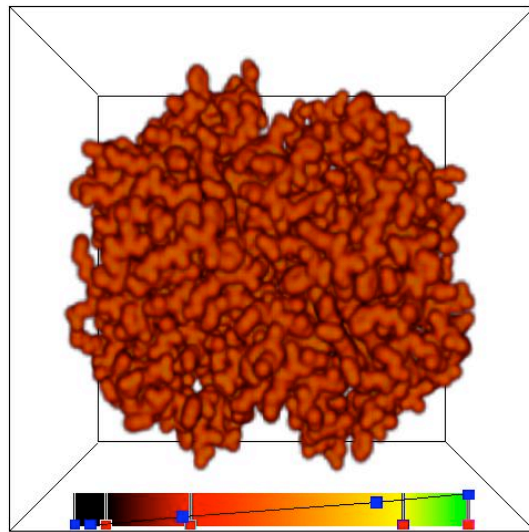
Within Subunit A

	F helix
	Histidine Ligand(HIS87)
	O <sub>2</sub>

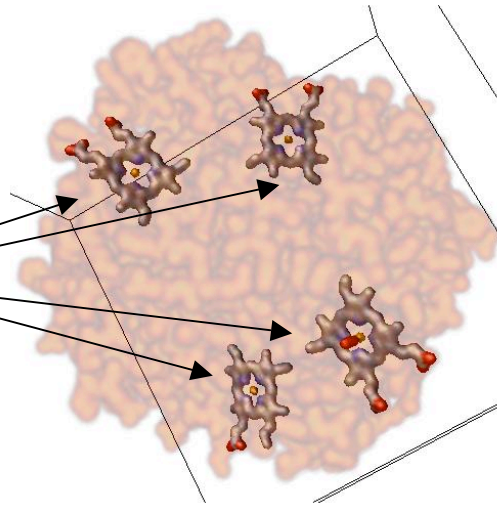
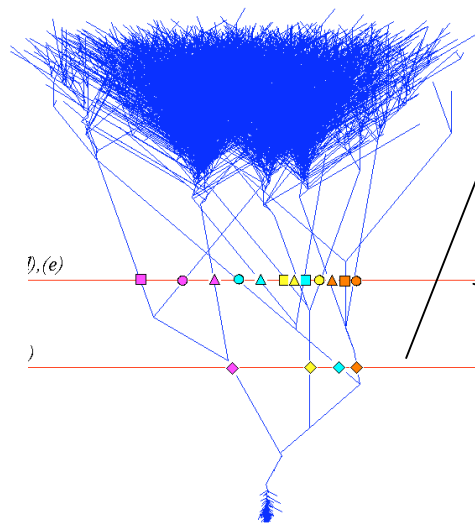
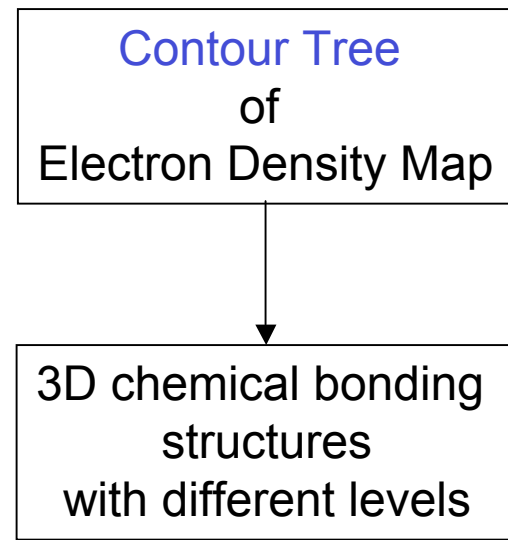
- Oxy process : O<sub>2</sub> binds to the Fe<sup>2+</sup> ion on the opposite side of the histidine ligand. F helix shifts position through the oxy-deoxy cycle.



# Topological Analysis & Visualization



Four Polypeptide chains



**Functional groups**  
Atoms belonging to the same  
contour have stronger linkage

Each chain consists of  
heme, iron, and globin

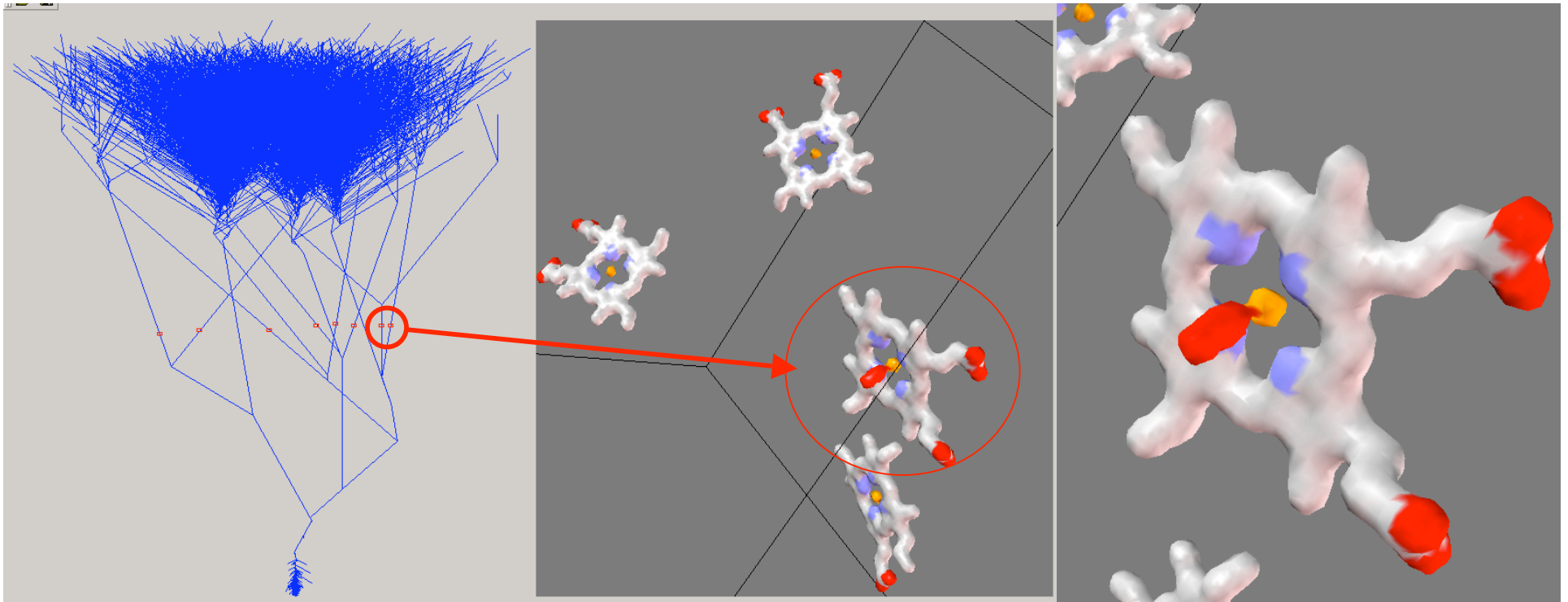
•M. van Kreveld, R. van Oostrum, C. Bajaj, V. Pascucci, and D. Schikore, *Chap5, pg 71 - 86, 2004 ed. by S. Rana, John Wiley & Sons, Ltd, 2004*  
•C. Bajaj, V. Pascucci, and D. Schikore, *Proceedings of the 1997 IEEE Visualization Conference, 167-173, October 1997 Phoenix, Arizona*





# Topological Analysis using the CONTOUR TREE

- Oxygenated Hemoglobin ( T=1 )

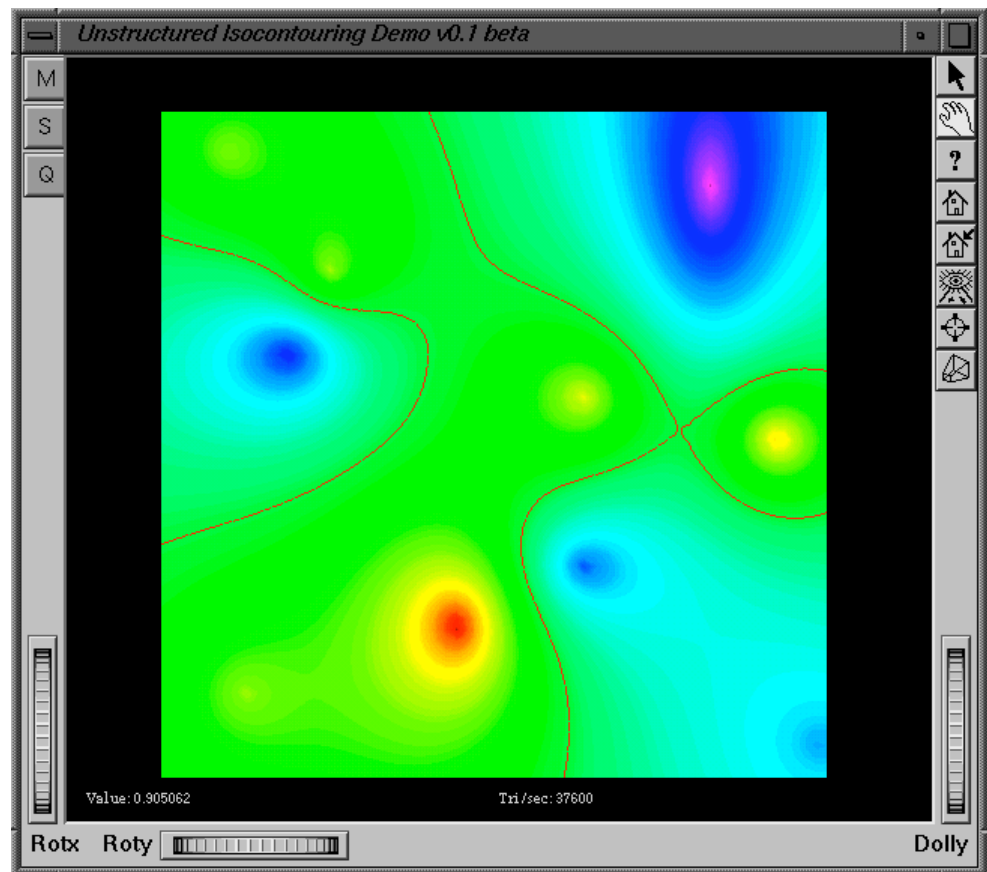


<isovalue = 31>



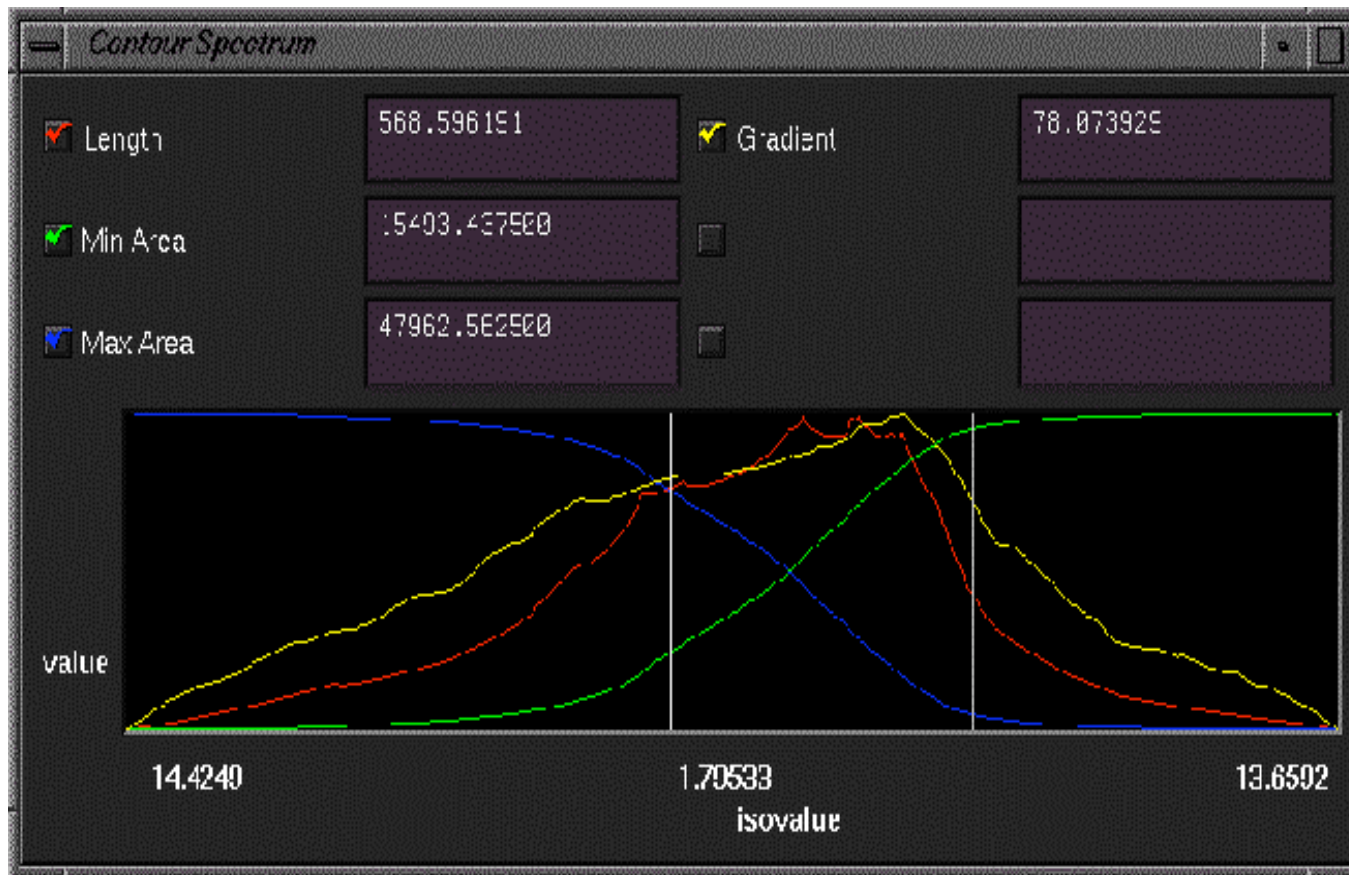
# Spectral Analysis

- Consider a terrain of which you want to compute the length of each isocontour and the area contained inside each isocontour.



# Spectral Analysis

## Graphical User Interface for Static Data



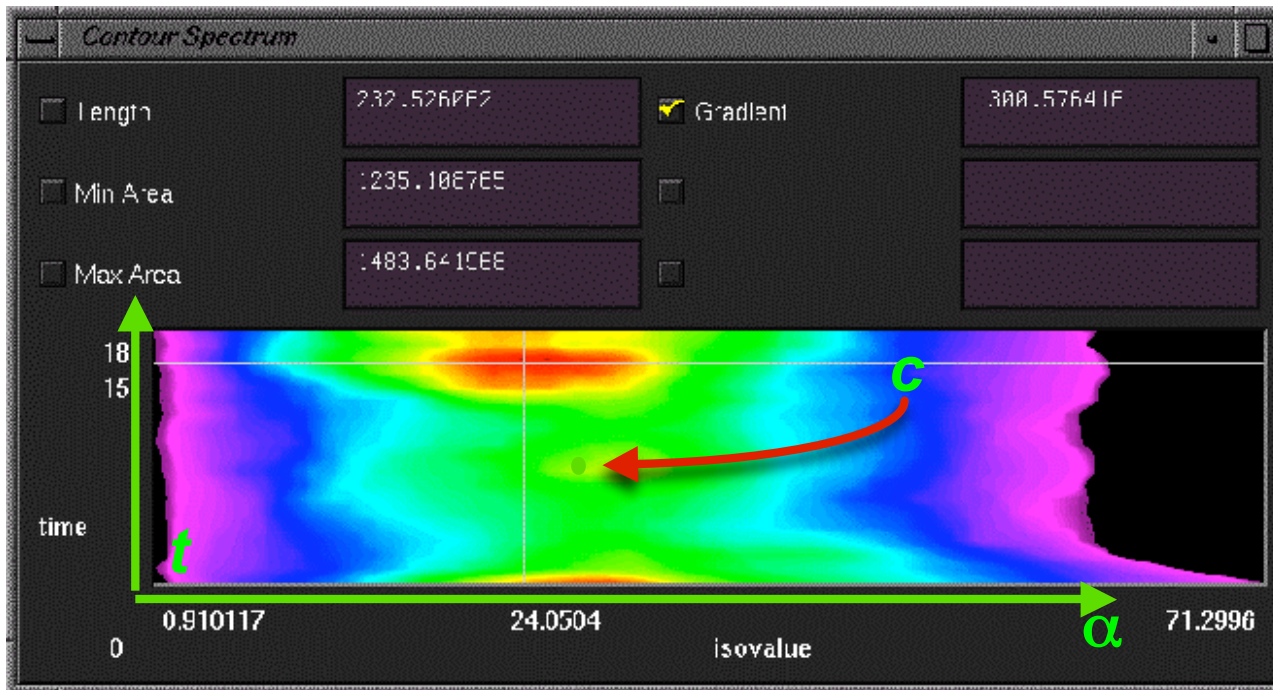
- The horizontal axis spans the scalar values  $\alpha$ .
- Plot of a set of signatures (length, area, gradient ...) as functions of the scalar value  $\alpha$ .

- Vertical axis spans normalized ranges of each



# Spectral Analysis

## Graphical User Interface for time varying data



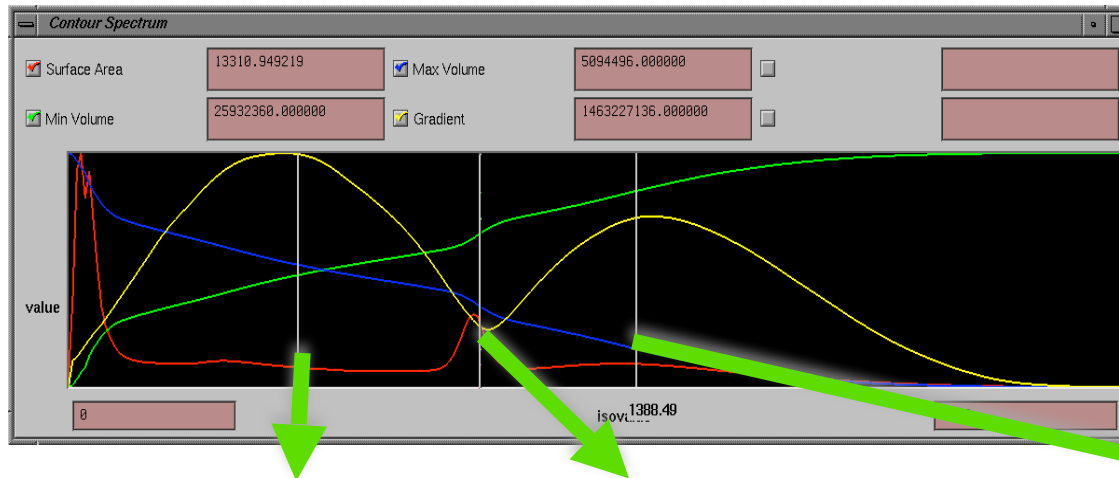
The horizontal axis spans the scalar value dimension  $\alpha$   
 The vertical axis spans the time dimension  $t$

high  
 $(\alpha, t) \rightarrow c$   
 The color  $c$   
 is mapped  
 to the  
 magnitude  
 of a  
 signature  
 function of  
 time  $t$  and  
 isovalue  $\alpha$   
 low

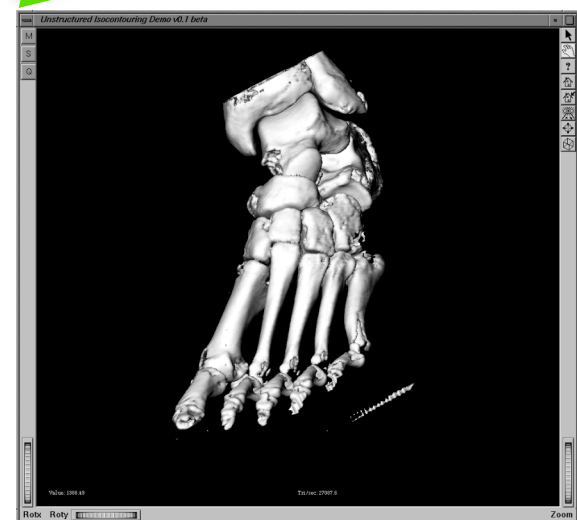
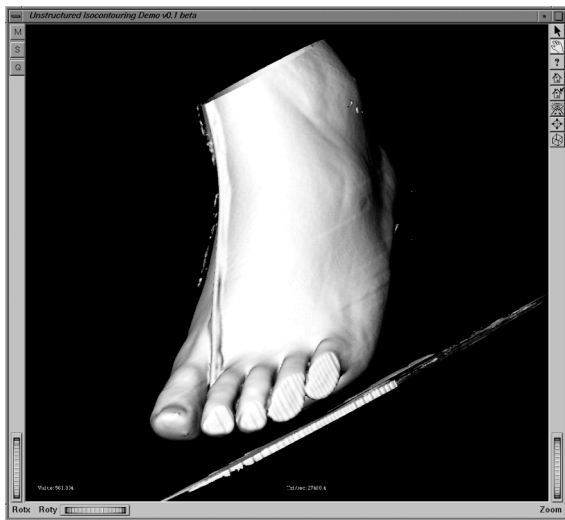
magnitude



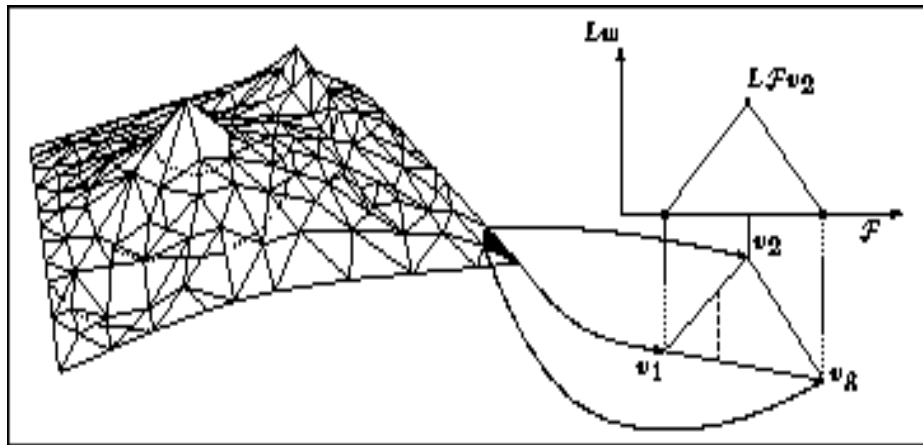
# Contouring based Selection



- *The contour spectrum allows the development of an adaptive ability to separate interesting isovalues from the others.*

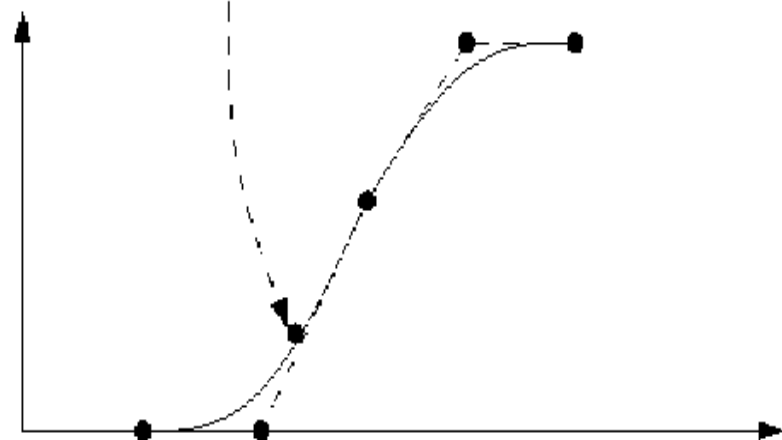
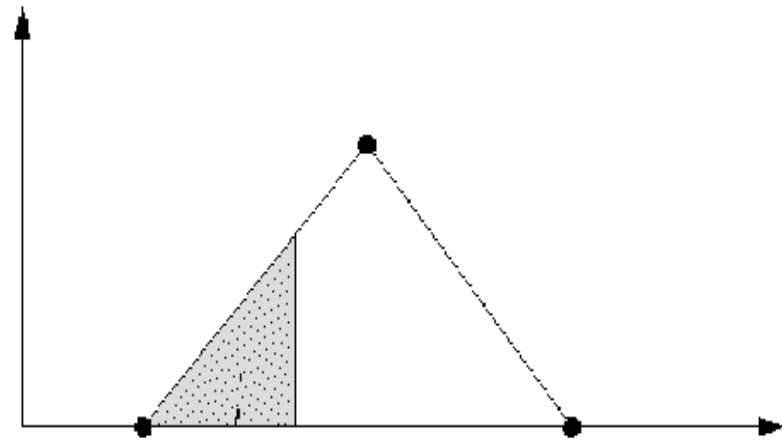


# Spectral Analysis (signature computation)



- The length of each contour is a  $C^0$  spline function.

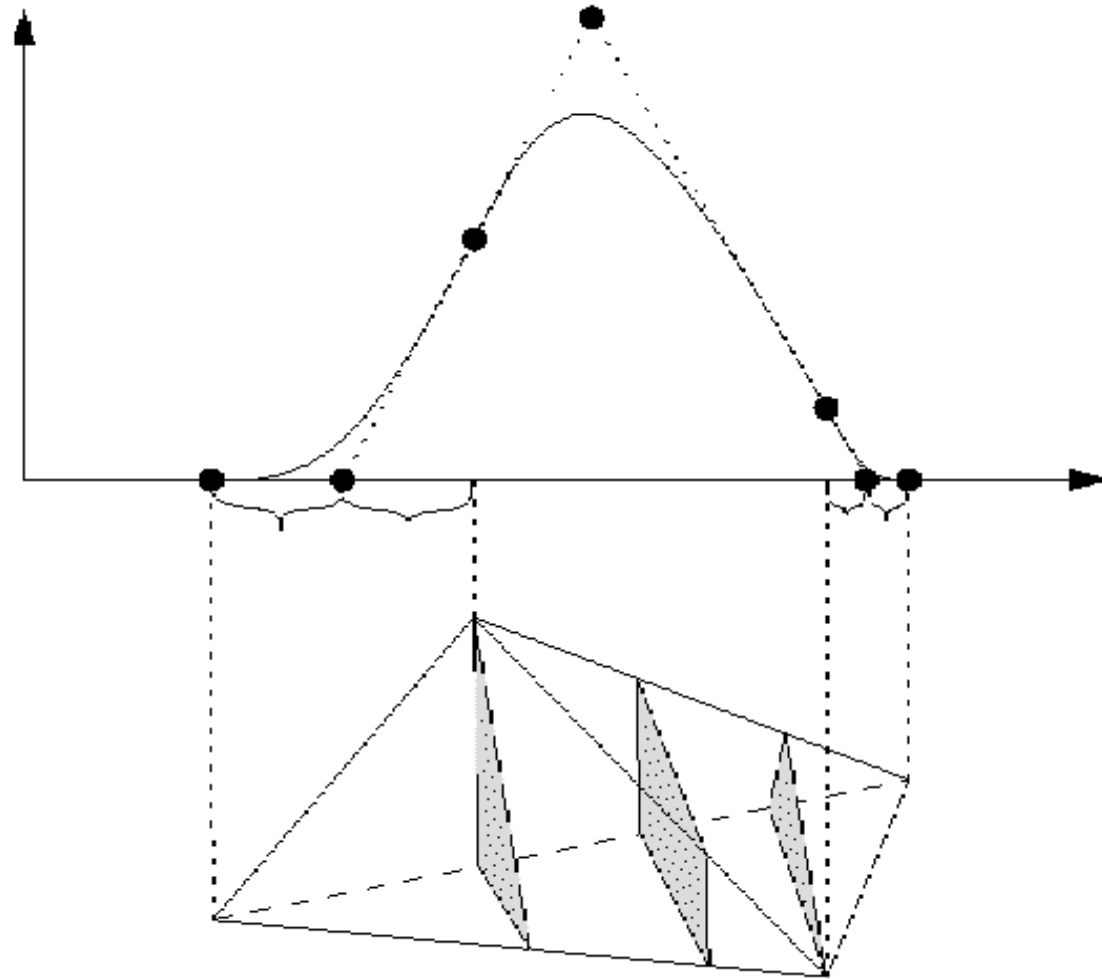
The area inside/outside each isocontour is a  $C^1$  spline function.





# Spectral Analysis (signature computation)

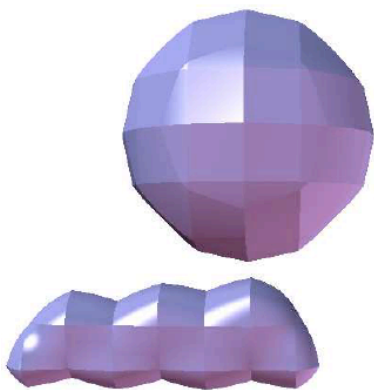
- In general the size of each isocontour of a scalar field of dimension  $d$  is a spline function of  $d-2$  continuity.
- The size of the region inside/outside is given by a spline function of  $d-1$  continuity



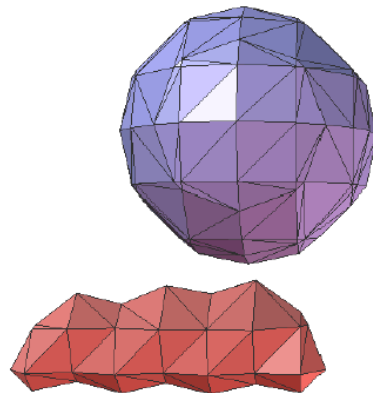
# Applications

## [ Contour Tree Based Visualization ]

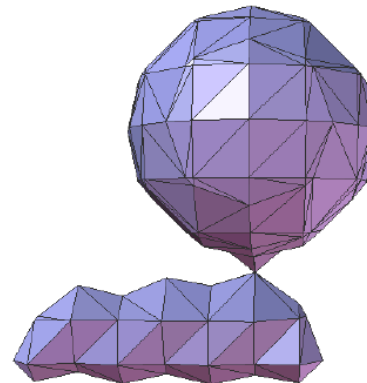
- Perform Tetrahedral Decomposition of Rectilinear Data (Trilinear Isosurface Topology is preserved)
- Apply Contour Tree and Seed Set computation, and Contour Propagation for Isosurface Component Segmentation



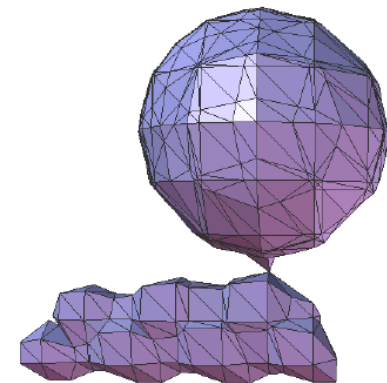
(a) Trilinear



(b) Cell Decomposition



(c) Marching Cubes



(d) Marching Tetra

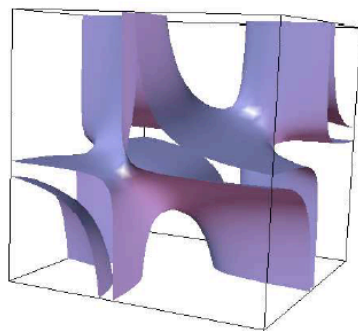




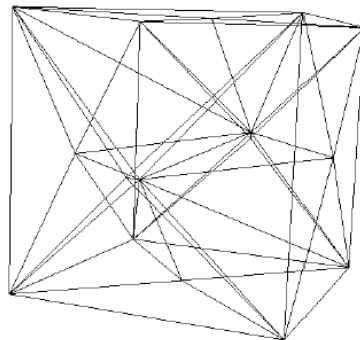
# Applications

## [ Trilinear Interval Volume Tetrahedrization ]

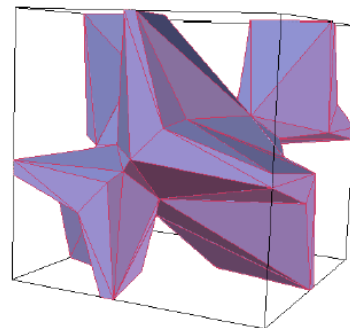
- Perform topology preserving tetrahedral decomposition method
- Apply interval volume tetrahedrization to each tetrahedra generated from our method



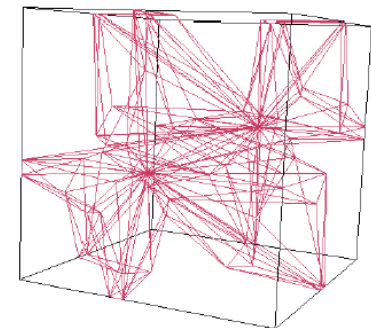
(a) boundary isosurfaces



(b) tetrahedral decomposition



(c) interval volume



(d) wireframe



# Further Reading

- C. Bajaj (ed) “DataVisualization Techniques”, John Wiley & Sons 1998
- C. Bajaj, V. Pascucci, D. Schikore, “Contour Spectrum” IEEE Viz, 1997
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- B. Sohn, C. Bajaj. “Topology Preserving Tetrahedral Decomposition of Trilinear Cell”, CS/ICES Tech. Rep. TR2004.
- S. Goswami, A. Gillette, C. Bajaj “Efficient Delaunay Mesh Generation from Sampled Scalar Functions”, 16h IMR, 2007
- J. Bloomenthal, C. Bajaj, J. Blinn, M. Gascuel, A. Rockwood, B. Wyvill, G. Wyvill **Introduction to Implicit Surfaces** Morgan Kaufman Publishers Inc., (1997).

