

CS341 Automata Theory
Homework Assignment #6

Do not forget to write your name and EID.

1. Give state diagram for TM that recognizes the following language:
 $\{w2w^R | w \in (0 \cup 1)^*\}$
2. Design a TM that computes the function
 - (a) $f(n) = 3n$
 - (b) $f(n) = n*(n-1)$
 - (c) $f(n, m) = n+m+3$
3. A Turing machine with left reset is similar to an ordinary TM except that the transition function has the form
 $\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{R, RESET\}$.
If $\delta(q, a) = (r, b, RESET)$, when the machine is in state q and reads an a , then the machine's head jumps to the left end of the tape after writing b on the tape and entering state r . These machines cannot move the tape head one symbol to the left. Show that Turing machines with left reset recognize the class of Turing-recognizable languages (ie, show this model is equivalent to the original TM model).
4. For the following languages/decision problems, prove whether the language is (I) decidable, (II) semi-decidable but not decidable, or (III) not semi-decidable.
 - (a) $L = \{ \langle M \rangle \mid \text{there is at least one input on which TM } M \text{ halts within 1000 steps} \}$
 - (b) $L = \{ \langle M \rangle \mid M \text{ accepts at least one string that ends with } 0 \}$
 - (c) $L = \{ \langle M_1, M_2 \rangle \mid \varepsilon \in L(M_1) \cup L(M_2) \}$
5. problem 5.9 in Sipser
6. Show the decidable languages are closed under complement.
7. Show the semi-decidable languages are closed under concatenation.
8. problem 4.15 in Sipser
9. Consider the following question: Given a DFA D , is $L(D)$ infinite? Show that this language is decidable by defining the corresponding language and proving the language is decidable.