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EVENT: Start with the initial **nqthm** theory.

; Talk about journal level proof-checkers.

; This sequence of lemmas describes the relationship between Ackermann's original
; function and R. Peter's version of it. The statement of the theorem
; appears in a manuscript by George D. Herbert, Department of
; Computer and Information Sciences, University of North Florida, Jacksonville,
; Florida. I was asked to review this result by the Monthly. -- rsb

; Ackermann's original function.

DEFINITION:

$a(m, n, a)$
= if $m \simeq 0$ then $n + a$
elseif $n \simeq 0$ then 0

```

    elseif  $n = 1$  then  $a$ 
    else  $a(m - 1, a(m, n - 1, a), a)$  endif

; Peter's simplification of it.

DEFINITION:
 $p(m, n)$ 
= if  $m \simeq 0$  then  $1 + n$ 
  elseif  $n \simeq 0$  then  $p(m - 1, 1)$ 
  else  $p(m - 1, p(m, n - 1))$  endif

; A stupid lemma forcing the theorem-prover to see the obvious.

THEOREM: a-open
 $((m \not\simeq 0) \wedge (1 < n)) \rightarrow (a(m, n, a) = a(m - 1, a(m, n - 1, a), a))$ 

; A nice little lemma used in Herbert's proof.

THEOREM: a-2-2->4
 $a(m, 2, 2) = 4$ 

; Another stupid little lemma.

THEOREM: plus-2
 $(2 + x) = (1 + (1 + x))$ 

; The main result.

THEOREM: a->p
 $(p(0, n) = (1 + n))$ 
 $\wedge ((0 < m) \rightarrow ((3 + p(m, n)) = a(m - 1, 1 + (1 + (1 + n)), 2)))$ 

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From cowles%UWYO.BITNET@CUNYVM.CUNY.EDU Wed Sep 7 09:16:01 1988
 X-St-Return-Receipt-Requested:
 Date: Tue, 6 Sep 88 22:38:14 MDT
 From: cowles%UWYO.BITNET@CUNYVM.CUNY.EDU (John Cowles)
 Subject: Ackermann's Original Function.
 To: boyer@cli.com, moore@cli.com, kaufmann@cli.com

In the file basic.events, it is claimed that Ackermann's original function may be recursively defined on the nonnegative integers by

```

ack-h(x,y,z) <= if x=0 then y+z
                  else if z=0 then 0
                  else if z=1 then y

```

else ack-h(x-1, y, ack-h(x,y,z-1)).

However a look at the English translation, in van Heijenoort's source book in mathematical logic, of Ackermann's 1928 paper shows that the original function may be recursively defined on the nonnegative integers by

```
ack(x,y,z) <== if x=0 then y+z
                else if x<=2 and z=0 then x-1
                else if x>2 and z=0 then y
                else ack( x-1, y, ack(x,y,z-1) ).
```

For all nonnegative y and z, $\text{ack}(0,y,z) = y+z = \text{ack-h}(0,y,z)$.

For all nonnegative y and z, $\text{ack}(1,y,z) = y*z = \text{ack-h}(1,y,z)$.

For all nonnegative y and positive z, $\text{ack}(2,y,z) = y^z = \text{ack-h}(2,y,z)$.

For all nonnegative y and positive z, $\text{ack}(3,y,z) = \text{ack-h}(3,y,z+1)$.

For $x>3$, the exact relationship between $\text{ack}(x,y,z)$ and $\text{ack-h}(x,y,z)$ seems to be difficult to determine. In any case, ack and ack-h are not the same function!

From cowles@CORRAL.UWyo.Edu Fri Sep 30 17:17:18 1988

X-St-Return-Receipt-Requested:

Date: Fri, 30 Sep 88 15:09:58 MDT

From: cowles@CORRAL.UWyo.Edu (John Cowles)

Subject: more on Ackermann

To: boyer@cli.com

1. Please feel free to add my note on Ackermann's function as a comment. But see 4 below for more on Ackermann's function.

4.

SEVERAL VERSIONS OF ACKERMANN'S FUNCTION

by J. Cowles and T. Bailey
Dept. of Computer Science
University of Wyoming
Laramie, WY

We have located several versions of Ackermann's function which are listed below. The claims made below about these functions should be carefully

verified (by machine if possible).

The original version (1928) of Ackermann's function may be defined recursively on the nonnegative integers by

```
ack( x, y, z ) <== if x=0 then y+z
                    else if x<=2 and z=0 then x-1
                    else if x>2 and z=0 then y
                    else ack( x-1, y, ack(x,y,z-1) ).
```

Then for all nonnegative integers x, y, and z,

```
ack( x, y, 0 ) equals if x=1 then 0
                      else if x=2 then 1
                      else y ;
```

```
ack( 0, y, z ) equals y+z ;
```

```
ack( 1, y, z ) equals y*z ;
```

```
ack( 2, y, z ) equals y^z ;
```

```
ack( 3, y, z ) equals iter-exp(y,z+1) ;
```

```
ack( x, 1, z ) equals if x=0 then z+1
                      else if x=1 then z
                      else 1 .
```

Here the function iter-exp is defined recursively on the nonnegative integers by

```
iter-exp( y, z ) <== if z=0 then 1
                     else y ^ iter-exp(y,z-1).
```

=====

The Herbert version of Ackermann's function may be defined recursively on the nonnegative integers by

```
ack-h( x, y, z ) <== if x=0 then y+z
                     else if z=0 then 0
                     else if z=1 then y
                     else ack-h( x-1, y, ack-h(x,y,z-1) )
```

Then for all nonnegative integers x , y , and z ,

$\text{ack-h}(x, y, 0)$ equals if $x=0$ then y
 else 0 ;

$\text{ack-h}(x, y, 1)$ equals if $x=0$ then $y+1$
 else y ;

$\text{ack-h}(0, y, z)$ equals $y+z$;

$\text{ack-h}(1, y, z)$ equals $y*z$;

$\text{ack-h}(2, y, z)$ equals if $z=0$ then 0
 else y^z ;

$\text{ack-h}(3, y, z)$ equals if $z=0$ then 0
 else $\text{iter-exp}(y,z)$;

$\text{ack-h}(x, 2, 2)$ equals 4 .

=====

The Meyer and Ritchie version (1967) of Ackermann's function may be defined recursively on the nonnegative integers by

$\text{ack-mr}(x, z) \leq$ if $z=0$ then 1
 else if $x=0$ and $z=1$ then 2
 else if $x=0$ and $z>1$ then $z+2$
 else $\text{ack-mr}(x-1, \text{ack-mr}(z,z-1))$.

Then for all nonnegative integers x and z ,

$\text{ack-mr}(x, 0)$ equals 1 ;

$\text{ack-mr}(0, z)$ equals if $z=0$ then 1
 else if $z=1$ then 2
 else $2+z$;

$\text{ack-mr}(1, z)$ equals if $z=0$ then 1
 else $2*z$;

$\text{ack-mr}(2, z)$ equals 2^z ;

$\text{ack-mr}(3, z)$ equals $\text{iter-exp}(2,z)$;

ack-mr(x, 1) equals 2 ;

ack-mr(x, 2) equals 4 .

=====

The R. Peter version (1935) of Ackermann's function may be defined recursively on the nonnegative integers by

ack-p(x, z) <== if x=0 then z+1
 else if z=0 then ack-p(x-1,1)
 else ack-p(x-1, ack-p(x,z-1)).

Then for all nonnegative integers z,

ack-p(0, z) equals z+1 ;

ack-p(1, z) equals z+2 ;

ack-p(2, z) equals $2*z + 3$;

ack-p(3, z) equals $2^{(z+3)} - 3$;

ack-p(4, z) equals iter-exp(2,z+3) - 3 .

=====

The Z. Manna version (1974) of Ackermann's function may be defined recursively on the nonnegative integers by

ack-m(x, z) <== if x*z = 0 then x+z+1
 else ack-m(x-1, ack-m(x,z-1)).

Then for all nonnegative integers z,

ack-m(0, z) equals z+1 ;

ack-m(1, z) equals z+2 ;

ack-m(2, z) equals $2*z + 3$;

ack-m(3, z) equals $7*(2^z) - 3$;

ack-m(4, z) equals $h(2,z) - 3$.

Here the function h is defined recursively on
the nonnegative integers by

$h(x, z) \leq$ if $z=0$ then y^3
 else $7*(y^{[h(y,z-1)-3]})$.

=====

Since everyone else has a version of Ackermann's function,
it should cause little or no harm if we also define a version.

ack-cb(x, y, z) $\leq$ if $x=0$ then $y+z$
 else if $x=1$ and $z=0$ then 0
 else if $x>1$ and $z=0$ then 1
 else ack-cb($x-1$, y, ack-cb($x,y,z-1$)).

Then for all nonnegative integers x, y, and z,

ack-cb(x, y, 0) equals if $x=0$ then y
 else if $x=1$ then 0
 else 1 ;

ack-cb(0, y, z) equals $y+z$;

ack-cb(1, y, z) equals $y*z$;

ack-cb(2, y, z) equals y^z ;

ack-cb(3, y, z) equals iter-exp(y,z) ;

ack-cb(x, y, 1) equals if $x=0$ then $y+1$
 else y ;

ack-cb(x, 2, 2) equals 4 .

=====

Then for all nonnegative integers x, y, and z,

ack-h(x, y, $z+1$) equals ack-cb(x, y, $z+1$) ;

ack-mr(x, $z+2$) equals ack-cb(x, 2, $z+2$)

```

                                equals  ack-h ( x, 2, z+2 ) ;
ack-p( x+1, z )  equals  ack-mr( x, z+3 ) - 3
                                equals  ack-h ( x, 2, z+3 ) - 3
                                equals  ack-cb( x, 2, z+3 ) - 3 .

```

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