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|#

EVENT: Start with the library "unsolv" using the compiled version.

EVENT: For efficiency, compile those definitions not yet compiled.

DEFINITION: $\text{symbol}(x) = (x \in '(0\ 1))$

DEFINITION:

half-tape(x)

= **if** $x \simeq \text{nil}$ **then** $x = 0$

else $\text{symbol}(\text{car}(x)) \wedge \text{half-tape}(\text{cdr}(x))$ **endif**

DEFINITION:

tape(x) = $(\text{listp}(x) \wedge \text{half-tape}(\text{car}(x)) \wedge \text{half-tape}(\text{cdr}(x)))$

DEFINITION: $\text{operation}(x) = (x \in '(1\ \text{r}\ 0\ 1))$

DEFINITION: $\text{state}(x) = \text{litatom}(x)$

DEFINITION:

$\text{turing-4tuple}(x)$
 $= \text{listp}(x)$
 $\quad \wedge \text{state}(\text{car}(x))$
 $\quad \wedge \text{symbol}(\text{cadr}(x))$
 $\quad \wedge \text{operation}(\text{caddr}(x))$
 $\quad \wedge \text{state}(\text{cadddr}(x))$
 $\quad \wedge (\text{cddddr}(x) = \mathbf{nil})$

DEFINITION:

$\text{turing-machine}(x)$
 $= \text{if } x \simeq \mathbf{nil} \text{ then } x = \mathbf{nil}$
 $\quad \text{else } \text{turing-4tuple}(\text{car}(x)) \wedge \text{turing-machine}(\text{cdr}(x)) \text{ endif}$

DEFINITION:

$\text{instr}(st, sym, tm)$
 $= \text{if listp}(tm)$
 $\quad \text{then if } st = \text{car}(\text{car}(tm))$
 $\quad \quad \text{then if } sym = \text{car}(\text{cdr}(\text{car}(tm))) \text{ then } \text{cdr}(\text{cdr}(\text{car}(tm)))$
 $\quad \quad \quad \text{else } \text{instr}(st, sym, \text{cdr}(tm)) \text{ endif}$
 $\quad \quad \text{else } \text{instr}(st, sym, \text{cdr}(tm)) \text{ endif}$
 $\quad \text{else f endif}$

DEFINITION:

$\text{new-tape}(op, tape)$
 $= \text{if } op = 'l$
 $\quad \text{then } \text{cons}(\text{cdr}(\text{car}(tape)), \text{cons}(\text{car}(\text{car}(tape)), \text{cdr}(tape)))$
 $\quad \text{elseif } op = 'r$
 $\quad \text{then } \text{cons}(\text{cons}(\text{car}(\text{cdr}(tape)), \text{car}(tape)), \text{cdr}(\text{cdr}(tape)))$
 $\quad \text{else } \text{cons}(\text{car}(tape), \text{cons}(op, \text{cdr}(\text{cdr}(tape)))) \text{ endif}$

DEFINITION:

$\text{tmi}(st, tape, tm, n)$
 $= \text{if } n \simeq 0 \text{ then BTM}$
 $\quad \text{elseif } \text{instr}(st, \text{car}(\text{cdr}(tape)), tm)$
 $\quad \text{then } \text{tmi}(\text{car}(\text{cdr}(\text{instr}(st, \text{car}(\text{cdr}(tape)), tm))),$
 $\quad \quad \text{new-tape}(\text{car}(\text{instr}(st, \text{car}(\text{cdr}(tape)), tm)), tape),$
 $\quad \quad tm,$
 $\quad \quad n - 1)$
 $\quad \text{else } tape \text{ endif}$

DEFINITION:

INSTR-DEFN

```

= '((st sym tm)
  (if
    (listp tm)
    (if
      (equal st (car (car tm)))
      (if
        (equal sym (car (cdr (car tm))))
        (cdr (cdr (car tm)))
        (instr st sym (cdr tm)))
      (instr st sym (cdr tm)))
    f))

```

DEFINITION:

NEW-TAPE-DEFN

```

= '((op tape)
  (if
    (equal op 'l)
    (cons
      (cdr (car tape))
      (cons (car (car tape)) (cdr tape)))
    (if
      (equal op 'r)
      (cons
        (cons (car (cdr tape)) (car tape))
        (cdr (cdr tape)))
      (cons (car tape) (cons op (cdr (cdr tape)))))))

```

DEFINITION:

TMI-DEFN

```

= '((st tape tm)
  (if
    (instr st (car (cdr tape)) tm)
    (tmi
      (car (cdr (instr st (car (cdr tape)) tm)))
      (new-tape
        (car (instr st (car (cdr tape)) tm))
        tape)
      tm)
    tape))

```

DEFINITION: $\text{kwote}(x) = \text{list}('quote, x)$

DEFINITION:

$\text{tmi-fa}(tm)$
 $= \text{list}(\text{list}('tm, \text{nil}, \text{kwote}(tm)),$

```

cons('instr, INSTR-DEFN),
cons('new-tape, NEW-TAPE-DEFN),
cons('tmi, TMI-DEFN))

```

DEFINITION:

$\text{tmi-x}(st, \text{tape}) = \text{list}('tmi, \text{kwote}(st), \text{kwote}(\text{tape}), '(\text{tm}))$

DEFINITION:

```

length(x)
=  if x  $\simeq$  nil then 0
   else 1 + length(cdr(x)) endif

```

DEFINITION: $\text{tmi-k}(st, \text{tape}, tm, n) = (n - (1 + \text{length}(tm)))$

DEFINITION: $\text{tmi-n}(st, \text{tape}, tm, k) = (k + (1 + \text{length}(tm)))$

THEOREM: length-0

$(\text{length}(x) = 0) = (x \simeq \text{nil})$

THEOREM: plus-equal-0

$((i + j) = 0) = ((i \simeq 0) \wedge (j \simeq 0))$

THEOREM: plus-difference

```

((i - j) + j)
=  if i  $\leq$  j then fix(j)
   else i endif

```

EVENT: Enable difference; name this event 'difference-off'.

THEOREM: pr-eval-fn-0

```

((n  $\simeq$  0)
 $\wedge$  (fn  $\neq$  'quote)
 $\wedge$  (fn  $\neq$  'if)
 $\wedge$  ( $\neg$  unsolv-subrp(fn))
 $\wedge$  (vargs = ev('list, args, va, fa, n)))
 $\rightarrow$  (ev('al, cons(fn, args), va, fa, n) = BTM)

```

THEOREM: pr-eval-fn-1

```

((n  $\neq$  0)
 $\wedge$  (fn  $\neq$  'quote)
 $\wedge$  (fn  $\neq$  'if)
 $\wedge$  ( $\neg$  unsolv-subrp(fn))
 $\wedge$  (vargs = ev('list, args, va, fa, n)))
 $\rightarrow$  (ev('al, cons(fn, args), va, fa, n)
    =  if btmp(vargs) then BTM

```

```

elseif btmp (get (fn, fa)) then BTM
else ev ('a1,
        cadr (get (fn, fa)),
        pairlist (car (get (fn, fa)), vars),
        fa,
        n - 1) endif)

```

THEOREM: pr-eval-unsolv-subrp
 $(\text{unsolv-subrp}(fn) \wedge (vars = \text{ev}('list, args, va, fa, n)))$
 $\rightarrow (\text{ev}('a1, \text{cons}(fn, args), va, fa, n)$
 $= \text{if btmp}(vars) \text{ then BTM}$
 $\text{else unsolv-apply-subr}(fn, vars) \text{ endif})$

THEOREM: pr-eval-if
 $(vx1 = \text{ev}('a1, x1, va, fa, n))$
 $\rightarrow (\text{ev}('a1, \text{list}('if, x1, x2, x3), va, fa, n)$
 $= \text{if btmp}(vx1) \text{ then BTM}$
 $\text{elseif } vx1 \text{ then ev}('a1, x2, va, fa, n)$
 $\text{else ev}('a1, x3, va, fa, n) \text{ endif})$

THEOREM: pr-eval-quote
 $\text{ev}('a1, \text{list}('quote, x), va, fa, n) = x$

THEOREM: pr-eval-nlistp
 $(\text{ev}('a1, 0, va, fa, n) = 0)$
 $\wedge (\text{ev}('a1, 1 + n, va, fa, n) = (1 + n))$
 $\wedge (\text{ev}('a1, \text{pack}(x), va, fa, n)$
 $= \text{if pack}(x) = 't \text{ then } t$
 $\text{elseif pack}(x) = 'f \text{ then } f$
 $\text{elseif pack}(x) = 'nil \text{ then nil}$
 $\text{else get}(\text{pack}(x), va) \text{ endif})$

THEOREM: evlist-nil
 $\text{ev}('list, nil, va, fa, n) = nil$

THEOREM: evlist-cons
 $((vx = \text{ev}('a1, x, va, fa, n)) \wedge (vl = \text{ev}('list, l, va, fa, n)))$
 $\rightarrow (\text{ev}('list, \text{cons}(x, l), va, fa, n)$
 $= \text{if btmp}(vx) \text{ then BTM}$
 $\text{elseif btmp}(vl) \text{ then BTM}$
 $\text{else cons}(vx, vl) \text{ endif})$

EVENT: Enable unsolv-subrp; name this event 'unsolv-subrp-off'.

EVENT: Enable ev; name this event 'ev-off'.

DEFINITION:

cnb(x)
 = **if** listp(x) **then** cnb(car(x)) $\wedge$ cnb(cdr(x))
 else $\neg$ btmp(x) **endif**

THEOREM: cnb-nbtm

cnb(x) $\rightarrow$ (btmp(x) = **f**)

THEOREM: cnb-cons

(cnb(cons(a , b)) = (cnb(a) $\wedge$ cnb(b)))
 $\wedge$ (cnb(x) $\rightarrow$ cnb(car(x)))
 $\wedge$ (cnb(x) $\rightarrow$ cnb(cdr(x)))

THEOREM: cnb-litatom

litatom(x) $\rightarrow$ cnb(x)

THEOREM: cnb-numberp

($x \in \mathbf{N}$) $\rightarrow$ cnb(x)

EVENT: Enable cnb; name this event 'cnb-off'.

THEOREM: get-tmi-fa

(get('tm, tmi-fa(tm)) = list(**nil**, kwote(tm)))
 $\wedge$ (get('instr, tmi-fa(tm)) = INSTR-DEFN)
 $\wedge$ (get('new-tape, tmi-fa(tm)) = NEW-TAPE-DEFN)
 $\wedge$ (get('tmi, tmi-fa(tm)) = TMI-DEFN)

EVENT: Enable tmi-fa; name this event 'tmi-fa-off'.

DEFINITION:

instrn(st , sym , tm , n)
 = **if** $n \simeq 0$ **then** BTM
 elseif listp(tm)
 then if $st = \text{car}(\text{car}(tm))$
 then if $sym = \text{car}(\text{cdr}(\text{car}(tm)))$ **then** cdr(cdr(car(tm)))
 else instrn(st , sym , cdr(tm), $n - 1$) **endif**
 else instrn(st , sym , cdr(tm), $n - 1$) **endif**
 else f endif

DEFINITION:

pr-eval-instr-induction-scheme(st , sym , tm -, va , tm , n)
 = **if** $n \simeq 0$ **then** **t**
 else pr-eval-instr-induction-scheme('st,
 'sym,

```

      '(cdr tm),
      list (cons ('st,
                  pr-eval (st,
                          va,
                          tmi-fa (tm),
                          n)),
            cons ('sym,
                  pr-eval (sym,
                          va,
                          tmi-fa (tm),
                          n)),
            cons ('tm,
                  pr-eval (tm-,
                          va,
                          tmi-fa (tm),
                          n))),
      tm,
      n - 1) endif

```

THEOREM: pr-eval-instr

```

(cnb (ev ('al, st, va, tmi-fa (tm), n))
  ∧ cnb (ev ('al, sym, va, tmi-fa (tm), n))
  ∧ cnb (ev ('al, tm-, va, tmi-fa (tm), n)))
→ (ev ('al, list ('instr, st, sym, tm-), va, tmi-fa (tm), n)
    = instrn (ev ('al, st, va, tmi-fa (tm), n),
               ev ('al, sym, va, tmi-fa (tm), n),
               ev ('al, tm-, va, tmi-fa (tm), n),
               n))

```

THEOREM: pr-eval-new-tape

```

(cnb (ev ('al, op, va, tmi-fa (tm), n)) ∧ cnb (ev ('al, tape, va, tmi-fa (tm), n)))
→ (ev ('al, list ('new-tape, op, tape), va, tmi-fa (tm), n)
    = if n ≃ 0 then BTM
      else new-tape (ev ('al, op, va, tmi-fa (tm), n),
                      ev ('al, tape, va, tmi-fa (tm), n)) endif)

```

THEOREM: cnb-instrn

```

((¬ btmp (instrn (st, sym, tm, n))) ∧ cnb (tm)) → cnb (instrn (st, sym, tm, n))

```

THEOREM: cnb-new-tape

```

(cnb (op) ∧ cnb (tape)) → cnb (new-tape (op, tape))

```

EVENT: Enable new-tape; name this event 'new-tape-off'.

DEFINITION:

```

tmin(st, tape, tm, n)
=  if  $n \simeq 0$  then BTM
    elseif btmp(instrn(st, car(cdr(tape)), tm,  $n - 1$ )) then BTM
    elseif instrn(st, car(cdr(tape)), tm,  $n - 1$ )
    then tmin(car(cdr(instrn(st, car(cdr(tape)), tm,  $n - 1$ ))),
              new-tape(car(instrn(st, car(cdr(tape)), tm,  $n - 1$ )), tape),
              tm,
               $n - 1$ )
    else tape endif

```

DEFINITION:

```

pr-eval-tmi-induction-scheme(st, tape, tm-, va, tm, n)
=  if  $n \simeq 0$  then t
    else pr-eval-tmi-induction-scheme(' (car
                                          (cdr
                                           (instr st
                                            (car (cdr tape))
                                            tm))),
                                          '(new-tape
                                           (car
                                            (instr st
                                             (car (cdr tape))
                                             tm))
                                           tape),
                                          'tm,
                                          list (cons ('st,
                                                       ev ('al,
                                                            st,
                                                            va,
                                                            tmi-fa(tm),
                                                            n)),
                                                  cons ('tape,
                                                       ev ('al,
                                                            tape,
                                                            va,
                                                            tmi-fa(tm),
                                                            n)),
                                                  cons ('tm,
                                                       ev ('al,
                                                            tm-,
                                                            va,
                                                            tmi-fa(tm),
                                                            n))),
                                          tm,

```


$n - 1$) **endif**

THEOREM: pr-eval-tmi

$$\begin{aligned} & (\text{cnb}(\text{ev}('a1, st, va, \text{tmi-fa}(tm), n)) \\ & \quad \wedge \text{cnb}(\text{ev}('a1, tape, va, \text{tmi-fa}(tm), n)) \\ & \quad \wedge \text{cnb}(\text{ev}('a1, tm-, va, \text{tmi-fa}(tm), n))) \\ \rightarrow & (\text{ev}('a1, \text{list}('tmi, st, tape, tm-), va, \text{tmi-fa}(tm), n) \\ & \quad = \text{tmin}(\text{ev}('a1, st, va, \text{tmi-fa}(tm), n), \\ & \quad \quad \text{ev}('a1, tape, va, \text{tmi-fa}(tm), n), \\ & \quad \quad \text{ev}('a1, tm-, va, \text{tmi-fa}(tm), n), \\ & \quad \quad n)) \end{aligned}$$

THEOREM: pr-eval-tmi-x

$$\begin{aligned} & (\text{cnb}(st) \wedge \text{cnb}(tape) \wedge \text{cnb}(tm)) \\ \rightarrow & (\text{ev}('a1, \text{tmi-x}(st, tape), \text{nil}, \text{tmi-fa}(tm), n) \\ & \quad = \text{if } n \simeq 0 \text{ then BTM} \\ & \quad \text{else tmin}(st, tape, tm, n) \text{ endif}) \end{aligned}$$

EVENT: Enable tmi-x; name this event 'tmi-x-off'.

THEOREM: instrn-instr

$$(\text{length}(tm) < n) \rightarrow (\text{instrn}(st, sym, tm, n) = \text{instr}(st, sym, tm))$$

THEOREM: nbtmp-instr

$$\text{turing-machine}(tm) \rightarrow (\neg \text{btm}(\text{instr}(st, sym, tm)))$$

THEOREM: instrn-non-f

$$(\text{turing-machine}(tm) \wedge (n \leq \text{length}(tm))) \rightarrow \text{instrn}(st, sym, tm, n)$$

THEOREM: tmin-btm

$$\begin{aligned} & (\text{turing-machine}(tm) \wedge (n \leq \text{length}(tm))) \\ \rightarrow & (\text{tmin}(st, tape, tm, n) = \text{BTM}) \end{aligned}$$

THEOREM: tmin-tmi

$$\begin{aligned} & \text{turing-machine}(tm) \\ \rightarrow & (\text{tmi}(st, tape, tm, k) = \text{tmin}(st, tape, tm, k + (1 + \text{length}(tm)))) \end{aligned}$$

THEOREM: symbol-cnb

$$\text{symbol}(sym) \rightarrow \text{cnb}(sym)$$

EVENT: Enable symbol; name this event 'symbol-off'.

THEOREM: half-tape-cnb

$$\text{half-tape}(x) \rightarrow \text{cnb}(x)$$

THEOREM: tape-cnb

$\text{tape}(x) \rightarrow \text{cnb}(x)$

EVENT: Enable tape; name this event ‘tape-off’.

THEOREM: operation-cnb

$\text{operation}(op) \rightarrow \text{cnb}(op)$

EVENT: Enable operation; name this event ‘operation-off’.

THEOREM: turing-machine-cnb

$\text{turing-machine}(tm) \rightarrow \text{cnb}(tm)$

EVENT: Enable turing-machine; name this event ‘turing-machine-off’.

THEOREM: turing-completeness-of-lisp

$(\text{state}(st) \wedge \text{tape}(tape) \wedge \text{turing-machine}(tm))$
 $\rightarrow (((\neg \text{btmp}(\text{pr-eval}(\text{tmi-x}(st, tape), \mathbf{nil}, \text{tmi-fa}(tm), n)))$
 $\rightarrow (\neg \text{btmp}(\text{tmi}(st, tape, tm, \text{tmi-k}(st, tape, tm, n)))))$
 $\wedge ((\neg \text{btmp}(\text{tmi}(st, tape, tm, k)))$
 $\rightarrow (\text{tmi}(st, tape, tm, k)$
 $= \text{pr-eval}(\text{tmi-x}(st, tape),$
 $\mathbf{nil},$
 $\text{tmi-fa}(tm),$
 $\text{tmi-n}(st, tape, tm, k)))))$

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