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|#

```
;          Proof of the Correctness of a Binary Search Program
;
```

EVENT: Start with the library "mc20-2" using the compiled version.

#|

The following C function BSEARCH determines if a given value x occurs in the sorted array a.

```
/* from K&R */
/* bsearch: find x in a[0] <= a[1] <= ... <= a[n-1] */
int bsearch (int x, int a[], int n)
{
    int low, high, mid;

    low = 0;
    high = n;
    while (low < high) {
        mid = (low + high) / 2;
```

```

    if (x < a[mid])
        high = mid;
    else if (x > a[mid])
        low = mid + 1;
    else return mid;
}
return -1;
}

```

Here is the MC68020 assembly code of the above BSEARCH program. The code is generated by "gcc -O".

```

0x22f0 <bsearch>:      linkw a6,#0
0x22f4 <bsearch+4>:    moveml d2-d3,sp@-
0x22f8 <bsearch+8>:    move1 a6@ (8),d3
0x22fc <bsearch+12>:   moveal a6@ (12),a0
0x2300 <bsearch+16>:   clrl d1
0x2302 <bsearch+18>:   move1 a6@ (16),d2
0x2306 <bsearch+22>:   cmpl d1,d2
0x2308 <bsearch+24>:   ble 0x232a <bsearch+58>
0x230a <bsearch+26>:   move1 d1,d0
0x230c <bsearch+28>:   addl d2,d0
0x230e <bsearch+30>:   bpl 0x2312 <bsearch+34>
0x2310 <bsearch+32>:   addql #1,d0
0x2312 <bsearch+34>:   asrl #1,d0
0x2314 <bsearch+36>:   cmpl 0(a0)[d0.l*4],d3
0x2318 <bsearch+40>:   bge 0x231e <bsearch+46>
0x231a <bsearch+42>:   move1 d0,d2
0x231c <bsearch+44>:   bra 0x2306 <bsearch+22>
0x231e <bsearch+46>:   cmpl 0(a0)[d0.l*4],d3
0x2322 <bsearch+50>:   ble 0x232c <bsearch+60>
0x2324 <bsearch+52>:   move1 d0,d1
0x2326 <bsearch+54>:   addql #1,d1
0x2328 <bsearch+56>:   bra 0x2306 <bsearch+22>
0x232a <bsearch+58>:   move1 #-1,d0
0x232c <bsearch+60>:   moveml a6@ (-8),d2-d3
0x2332 <bsearch+66>:   unlk a6
0x2334 <bsearch+68>:   rts

```

The machine code of the above program is:

```

<bsearch>:      0x4e56  0x0000  0x48e7  0x3000  0x262e  0x0008  0x206e  0x000c
<bsearch+16>:   0x4281  0x242e  0x0010  0xb481  0x6f20  0x2001  0xd082  0x6a02
<bsearch+32>:   0x5280  0xe280  0xb6b0  0x0c00  0x6c04  0x2400  0x60e8  0xb6b0

```

```

<bsearch+48>:  0x0c00  0x6f08  0x2200  0x5281  0x60dc  0x70ff  0x4cee  0x000c
<bsearch+64>:  0xff8  0x4e5e  0x4e75

```

In the Logic, this is:

```

'(78      86      0      0      72      231      48      0
   38      46      0      8      32      110      0      12
   66      129     36      46      0      16      180     129
  111      32      32      1     208      130      106      2
   82      128     226     128     182      176      12      0
  108      4       36      0      96      232      182     176
   12      0      111      8      34      0       82     129
   96      220     112     255     76      238      0      12
  255     248      78      94      78      117)

```

|#

; in the logic, the above program is defined by (bsearch-code).

DEFINITION:

BSEARCH-CODE

```

=  '(78 86 0 0 72 231 48 0 38 46 0 8 32 110 0 12 66 129
    36 46 0 16 180 129 111 32 32 1 208 130 106 2 82 128
    226 128 182 176 12 0 108 4 36 0 96 232 182 176 12 0
    111 8 34 0 82 129 96 220 112 255 76 238 0 12 255 248
    78 94 78 117)

```

THEOREM: mean-bounds

$(i < j) \rightarrow (((i + j) \div 2) < j) \wedge (((i + j) \div 2) \not< i)$

THEOREM: ilessp-lessp

$((x \in \mathbf{N}) \wedge (y \in \mathbf{N})) \rightarrow (\text{ilessp}(x, y) = (x < y))$

EVENT: Disable ilessp.

; bsearch1 is a function in the logic to simulate the loop of the
; above code.

DEFINITION:

bsearch1(x, lst, i, j)

= **let** k **be** $(i + j) \div 2$

in

if $i < j$

then if ilessp($x, \text{get-nth}(k, lst)$) **then** bsearch1(x, lst, i, k)

```

    elseif ilessp (get-nth (k, lst), x)
    then bsearch1 (x, lst, 1 + k, j)
    else k endif
else -1 endif endlet

```

; bsearch is a function in the logic to simulate the above code.

DEFINITION: $\text{bsearch}(x, n, \text{lst}) = \text{bsearch1}(x, \text{lst}, 0, n)$

; the computation time of the loop.

DEFINITION:

```

bsearch1-t (x, lst, i, j)
= let k be (i + j) ÷ 2
  in
  if i < j
  then if ilessp (x, get-nth (k, lst))
    then splus (10, bsearch1-t (x, lst, i, k))
    elseif ilessp (get-nth (k, lst), x)
    then splus (13, bsearch1-t (x, lst, 1 + k, j))
    else 13 endif
  else 6 endif endlet

```

; the computation time of the code.

DEFINITION:

$\text{bsearch-t}(x, n, \text{lst}) = \text{splus}(6, \text{bsearch1-t}(x, \text{lst}, 0, n))$

; an induction hint.

DEFINITION:

```

bsearch-induct (s, x, lst, i, j)
= let k be (i + j) ÷ 2
  in
  if i < j
  then if ilessp (x, get-nth (k, lst))
    then bsearch-induct (stepn (s, 10), x, lst, i, k)
    elseif ilessp (get-nth (k, lst), x)
    then bsearch-induct (stepn (s, 13), x, lst, 1 + k, j)
    else t endif
  else t endif endlet

```

; the preconditions of the initial state.

DEFINITION:

bsearch-statep(s, x, a, n, lst)
 $= ((\text{mc-status}(s) = \text{'running'})$
 $\wedge \text{evenp}(\text{mc-pc}(s))$
 $\wedge \text{rom-addrp}(\text{mc-pc}(s), \text{mc-mem}(s), 70)$
 $\wedge \text{mcode-addrp}(\text{mc-pc}(s), \text{mc-mem}(s), \text{BSEARCH-CODE})$
 $\wedge \text{ram-addrp}(\text{sub}(32, 12, \text{read-sp}(s)), \text{mc-mem}(s), 28)$
 $\wedge \text{ram-addrp}(a, \text{mc-mem}(s), 4 * n)$
 $\wedge \text{mem-ilst}(4, a, \text{mc-mem}(s), n, lst)$
 $\wedge \text{disjoint}(\text{sub}(32, 12, \text{read-sp}(s)), 28, a, 4 * n)$
 $\wedge (a = \text{read-mem}(\text{add}(32, \text{read-sp}(s), 8), \text{mc-mem}(s), 4))$
 $\wedge (n = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 12), \text{mc-mem}(s), 4))$
 $\wedge (x = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 4), \text{mc-mem}(s), 4))$
 $\wedge \text{int-rangep}(2 * n, 32)$
 $\wedge (n \in \mathbf{N}))$

; the conditions of an intermediate state.

DEFINITION:

bsearch-s0p(s, x, a, n, lst, i, j)
 $= ((\text{mc-status}(s) = \text{'running'})$
 $\wedge \text{evenp}(\text{mc-pc}(s))$
 $\wedge \text{rom-addrp}(\text{sub}(32, 22, \text{mc-pc}(s)), \text{mc-mem}(s), 70)$
 $\wedge \text{mcode-addrp}(\text{sub}(32, 22, \text{mc-pc}(s)), \text{mc-mem}(s), \text{BSEARCH-CODE})$
 $\wedge \text{ram-addrp}(\text{sub}(32, 8, \text{read-an}(32, 6, s)), \text{mc-mem}(s), 28)$
 $\wedge \text{ram-addrp}(a, \text{mc-mem}(s), 4 * n)$
 $\wedge \text{mem-ilst}(4, a, \text{mc-mem}(s), n, lst)$
 $\wedge \text{disjoint}(\text{sub}(32, 8, \text{read-an}(32, 6, s)), 28, a, 4 * n)$
 $\wedge (a = \text{read-an}(32, 0, s))$
 $\wedge (i = \text{nat-to-int}(\text{read-dn}(32, 1, s), 32))$
 $\wedge (j = \text{nat-to-int}(\text{read-dn}(32, 2, s), 32))$
 $\wedge (x = \text{nat-to-int}(\text{read-dn}(32, 3, s), 32))$
 $\wedge \text{int-rangep}(2 * j, 32)$
 $\wedge (i \in \mathbf{N})$
 $\wedge (j \in \mathbf{N})$
 $\wedge (n \in \mathbf{N})$
 $\wedge (i \leq n)$
 $\wedge (j \leq n))$

; the initial segment. From the initial state to s0.

THEOREM: bsearch-s-s0p

bsearch-statep(s, x, a, n, lst) $\rightarrow$ bsearch-s0p(stepn($s, 6$), $x, a, n, lst, 0, n$)

THEOREM: bsearch-s-s0

$\text{bsearch-statep}(s, x, a, n, lst)$
 $\rightarrow ((\text{linked-rtts-addr}(\text{stepn}(s, 6)) = \text{rts-addr}(s))$
 $\quad \wedge (\text{linked-a6}(\text{stepn}(s, 6)) = \text{read-an}(32, 6, s))$
 $\quad \wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, 6)))$
 $\quad \quad = \text{sub}(32, 4, \text{read-sp}(s)))$
 $\quad \wedge (\text{movem-saved}(\text{stepn}(s, 6), 4, 8, 2)$
 $\quad \quad = \text{readm-rn}(32, '(2\ 3), \text{mc-rfile}(s))))$

THEOREM: $\text{bsearch-s-s0-rfile}$
 $(\text{bsearch-statep}(s, x, a, n, lst) \wedge \text{d4-7a2-5p}(rn))$
 $\rightarrow (\text{read-rn}(oplen, rn, \text{mc-rfile}(\text{stepn}(s, 6)))$
 $\quad = \text{read-rn}(oplen, rn, \text{mc-rfile}(s)))$

THEOREM: bsearch-s-s0-mem
 $(\text{bsearch-statep}(s, x, a, n, lst) \wedge \text{disjoint}(x, k, \text{sub}(32, 12, \text{read-sp}(s)), 28))$
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 6)), k) = \text{read-mem}(x, \text{mc-mem}(s), k))$

; from s0 to s0 (induction case), from s0 to exit (base case).
 ; base case: s0 --> sn.

THEOREM: $\text{bsearch-s0-sn-base1}$
 $(\text{bsearch-s0p}(s, x, a, n, lst, i, j) \wedge (i \not\leq j))$
 $\rightarrow ((\text{mc-status}(\text{stepn}(s, 6)) = \text{'running'})$
 $\quad \wedge (\text{mc-pc}(\text{stepn}(s, 6)) = \text{linked-rtts-addr}(s))$
 $\quad \wedge (\text{iread-dn}(32, 0, \text{stepn}(s, 6)) = -1)$
 $\quad \wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, 6))) = \text{linked-a6}(s))$
 $\quad \wedge (\text{read-rn}(32, 15, \text{mc-rfile}(\text{stepn}(s, 6)))$
 $\quad \quad = \text{add}(32, \text{read-an}(32, 6, s), 8))$
 $\quad \wedge (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 6)), l)$
 $\quad \quad = \text{read-mem}(x, \text{mc-mem}(s), l)))$

THEOREM: $\text{bsearch-s0-sn-rfile-base1}$
 $(\text{bsearch-s0p}(s, x, a, n, lst, i, j)$
 $\quad \wedge (i \not\leq j)$
 $\quad \wedge \text{d2-7a2-5p}(rn)$
 $\quad \wedge (oplen \leq 32))$
 $\rightarrow (\text{read-rn}(oplen, rn, \text{mc-rfile}(\text{stepn}(s, 6)))$
 $\quad = \text{if } \text{d4-7a2-5p}(rn) \text{ then } \text{read-rn}(oplen, rn, \text{mc-rfile}(s))$
 $\quad \quad \text{else } \text{get-vlst}(oplen, 0, rn, '(2\ 3), \text{movem-saved}(s, 4, 8, 2)) \text{ endif})$

EVENT: Enable iplus.

EVENT: Enable iquotient.

EVENT: Disable quotient.

EVENT: Disable remainder.

EVENT: Disable times.

THEOREM: bsearch-crock
 $(\text{int-range}(2 * j, n) \wedge (i < j)) \rightarrow \text{int-range}(i + j, n)$

THEOREM: bsearch-s0-sn-base2
let k **be** $(i + j) \div 2$
in
 $(\text{bsearch-s0p}(s, x, a, n, lst, i, j)$
 $\wedge (i < j)$
 $\wedge (\neg \text{ilessp}(x, \text{get-nth}(k, lst)))$
 $\wedge (\neg \text{ilessp}(\text{get-nth}(k, lst), x))$
 $\rightarrow ((\text{mc-status}(\text{stepn}(s, 13)) = \text{'running'})$
 $\wedge (\text{mc-pc}(\text{stepn}(s, 13)) = \text{linked-rtts-addr}(s))$
 $\wedge (\text{iread-dn}(32, 0, \text{stepn}(s, 13)) = ((i + j) \div 2))$
 $\wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, 13)))$
 $\quad = \text{linked-a6}(s))$
 $\wedge (\text{read-rn}(32, 15, \text{mc-rfile}(\text{stepn}(s, 13)))$
 $\quad = \text{add}(32, \text{read-an}(32, 6, s), 8))$
 $\wedge (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 13)), l)$
 $\quad = \text{read-mem}(x, \text{mc-mem}(s, l)))$ **endlet**

THEOREM: bsearch1-s0-sn-rfile-base2
let k **be** $(i + j) \div 2$
in
 $(\text{bsearch-s0p}(s, x, a, n, lst, i, j)$
 $\wedge (i < j)$
 $\wedge (\neg \text{ilessp}(x, \text{get-nth}(k, lst)))$
 $\wedge (\neg \text{ilessp}(\text{get-nth}(k, lst), x))$
 $\wedge \text{d2-7a2-5p}(rn)$
 $\wedge (\text{oplen} \leq 32))$
 $\rightarrow (\text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(\text{stepn}(s, 13)))$
 $\quad = \text{if } \text{d4-7a2-5p}(rn) \text{ then } \text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(s))$
 $\quad \text{else } \text{get-vlst}(\text{oplen},$
 $\quad \quad 0,$
 $\quad \quad rn,$
 $\quad \quad \text{'(2 3)},$
 $\quad \text{movem-saved}(s, 4, 8, 2))$ **endif** **endlet**

; from s0 to s0 (induction case).

THEOREM: bsearch-s0-s0-1

```

let  $k$  be  $(i + j) \div 2$ 
in
  (bsearch-s0p ( $s, x, a, n, lst, i, j$ )
     $\wedge$  ( $i < j$ )
     $\wedge$  ilessp ( $x, \text{get-nth}(k, lst)$ )
   $\rightarrow$  (bsearch-s0p (stepn ( $s, 10$ ),  $x, a, n, lst, i, k$ )
     $\wedge$  (read-rn ( $oplen, 14, \text{mc-rfile}(\text{stepn}(s, 10))$ )
      = read-rn ( $oplen, 14, \text{mc-rfile}(s)$ ))
     $\wedge$  (linked-a6 (stepn ( $s, 10$ )) = linked-a6 ( $s$ ))
     $\wedge$  (linked-rt-s-addr (stepn ( $s, 10$ )) = linked-rt-s-addr ( $s$ ))
     $\wedge$  (movem-saved (stepn ( $s, 10$ ), 4, 8, 2)
      = movem-saved ( $s, 4, 8, 2$ ))
     $\wedge$  (read-mem ( $x, \text{mc-mem}(\text{stepn}(s, 10)), l$ )
      = read-mem ( $x, \text{mc-mem}(s), l$ ))) endlet

```

THEOREM: bsearch-s0-s0-rfile1

```

let  $k$  be  $(i + j) \div 2$ 
in
  (bsearch-s0p ( $s, x, a, n, lst, i, j$ )
     $\wedge$  ( $i < j$ )
     $\wedge$  ilessp ( $x, \text{get-nth}(k, lst)$ )
     $\wedge$  d4-7a2-5p ( $rn$ )
   $\rightarrow$  (read-rn ( $oplen, rn, \text{mc-rfile}(\text{stepn}(s, 10))$ )
    = read-rn ( $oplen, rn, \text{mc-rfile}(s)$ )) endlet

```

THEOREM: bsearch-s0-s0-2

```

let  $k$  be  $(i + j) \div 2$ 
in
  (bsearch-s0p ( $s, x, a, n, lst, i, j$ )
     $\wedge$  ( $i < j$ )
     $\wedge$  ( $\neg$  ilessp ( $x, \text{get-nth}(k, lst)$ ))
     $\wedge$  ilessp ( $\text{get-nth}(k, lst), x$ )
   $\rightarrow$  (bsearch-s0p (stepn ( $s, 13$ ),  $x, a, n, lst, 1 + k, j$ )
     $\wedge$  (read-rn ( $oplen, 14, \text{mc-rfile}(\text{stepn}(s, 13))$ )
      = read-rn ( $oplen, 14, \text{mc-rfile}(s)$ ))
     $\wedge$  (linked-a6 (stepn ( $s, 13$ )) = linked-a6 ( $s$ ))
     $\wedge$  (linked-rt-s-addr (stepn ( $s, 13$ )) = linked-rt-s-addr ( $s$ ))
     $\wedge$  (movem-saved (stepn ( $s, 13$ ), 4, 8, 2)
      = movem-saved ( $s, 4, 8, 2$ ))
     $\wedge$  (read-mem ( $x, \text{mc-mem}(\text{stepn}(s, 13)), l$ )
      = read-mem ( $x, \text{mc-mem}(s), l$ ))) endlet

```

THEOREM: bsearch1-s0-s0-rfile2

```

let  $k$  be  $(i + j) \div 2$ 
in

```



```

(bsearch-s0p (s, x, a, n, lst, i, j)
  ∧ (i < j)
  ∧ (¬ ilessp (x, get-nth (k, lst)))
  ∧ ilessp (get-nth (k, lst), x)
  ∧ d4-7a2-5p (rn))
→ (read-rn (oplen, rn, mc-rfile (stepn (s, 13)))
    = read-rn (oplen, rn, mc-rfile (s))) endlet

```

EVENT: Disable bsearch-statep.

EVENT: Disable bsearch-s0p.

```

; put together. s0 --> sn

```

THEOREM: bsearch-s0-sn

```

let sn be stepn (s, bsearch1-t (x, lst, i, j))
in
  bsearch-s0p (s, x, a, n, lst, i, j)
→ ((mc-status (sn) = 'running)
    ∧ (mc-pc (sn) = linked-rtts-addr (s))
    ∧ (iread-dn (32, 0, sn) = bsearch1 (x, lst, i, j))
    ∧ (read-rn (32, 14, mc-rfile (sn)) = linked-a6 (s))
    ∧ (read-rn (32, 15, mc-rfile (sn))
        = add (32, read-an (32, 6, s), 8))
    ∧ (read-mem (x, mc-mem (sn), k) = read-mem (x, mc-mem (s), k))) endlet

```

THEOREM: bsearch-s0-sn-rfile

```

(bsearch-s0p (s, x, a, n, lst, i, j) ∧ d2-7a2-5p (rn) ∧ (oplen ≤ 32))
→ (read-rn (oplen, rn, mc-rfile (stepn (s, bsearch1-t (x, lst, i, j))))
    = if d4-7a2-5p (rn) then read-rn (oplen, rn, mc-rfile (s))
      else get-vlst (oplen, 0, rn, ' (2 3), movem-saved (s, 4, 8, 2)) endif)

```

```

; the correctness statement of this BSEARCH program.
; after an execution of this program, the machine state satisfies:
; 0. normal exit.
; 1. the program counter is returned to the next instruction of the caller.
; 2. the result -- (bsearch x n lst), is stored in the register D0.
; 3. a6, used by LINK, is restored to its original value.
; 4. a7, the stack pointer, is updated correctly to pop off one frame.
; 5. the registers d2-d7 and a2-a5 maintain their original values.
; 6. the memory is only locally changed wrt this program.

```

THEOREM: bsearch-correctness

```

let sn be stepn (s, bsearch-t (x, n, lst))

```

```

in
bsearch-statep(s, x, a, n, lst)
→ ((mc-status(sn) = 'running)
   ∧ (mc-pc(sn) = rts-addr(s))
   ∧ (read-rn(32, 14, mc-rfile(sn))
      = read-rn(32, 14, mc-rfile(s)))
   ∧ (read-rn(32, 15, mc-rfile(sn))
      = add(32, read-sp(s), 4))
   ∧ ((d2-7a2-5p(rn) ∧ (oplen ≤ 32))
      → (read-rn(oplen, rn, mc-rfile(sn))
          = read-rn(oplen, rn, mc-rfile(s))))
   ∧ (disjoint(x, k, sub(32, 12, read-sp(s)), 28)
      → (read-mem(x, mc-mem(sn), k)
          = read-mem(x, mc-mem(s), k)))
   ∧ (iread-dn(32, 0, sn) = bsearch(x, n, lst))) endlet

```

EVENT: Disable bsearch-t.

```

; in the logic, bsearch has these properties:
;   1. if it returns a nonnegative integer i, then lst[i] = x.
;   2. if it returns -1, then x is not in lst.

```

DEFINITION:

```

member1(x, lst, i, j)
= if i < j
  then if x = get-nth(i, lst) then t
    else member1(x, lst, 1 + i, j) endif
  else f endif

```

DEFINITION:

```

orderedp(lst)
= if lst ≈ nil then t
  elseif cdr(lst) ≈ nil then t
  else ileq(car(lst), cadr(lst)) ∧ orderedp(cdr(lst)) endif

```

THEOREM: leq-trans

$$((\neg \text{ilessp}(x, y)) \wedge (\neg \text{ilessp}(y, z))) \rightarrow (\neg \text{ilessp}(x, z))$$

THEOREM: int-equal

$$(\text{integerp}(x) \wedge \text{integerp}(y) \wedge (\neg \text{ilessp}(x, y)) \wedge (\neg \text{ilessp}(y, x))) \\ \rightarrow ((x = y) = \mathbf{t})$$

THEOREM: orderedp-ordered

$$(\text{orderedp}(lst) \wedge (i \leq j) \wedge (j < \text{len}(lst))) \\ \rightarrow (\text{ilessp}(\text{get-nth}(j, lst), \text{get-nth}(i, lst)) = \mathbf{f})$$

THEOREM: bsearch1-found
 $((\text{bsearch1}(x, \text{lst}, i, j) \neq -1) \wedge \text{lst-integerp}(\text{lst}) \wedge \text{integerp}(x))$
 $\rightarrow (\text{get-nth}(\text{bsearch1}(x, \text{lst}, i, j), \text{lst}) = x)$

THEOREM: bsearch1-not-found-1
 $(\text{orderedp}(\text{lst})$
 $\wedge \text{ilessp}(\text{get-nth}(k, \text{lst}), x)$
 $\wedge (i \leq k)$
 $\wedge (k < j)$
 $\wedge (j \leq \text{len}(\text{lst})))$
 $\rightarrow (\text{member1}(x, \text{lst}, i, j) = \text{member1}(x, \text{lst}, 1 + k, j))$

THEOREM: bsearch1-not-found-2-lemma
 $(\text{orderedp}(\text{lst})$
 $\wedge \text{ilessp}(x, \text{get-nth}(k, \text{lst}))$
 $\wedge (k \leq i)$
 $\wedge (j < \text{len}(\text{lst})))$
 $\rightarrow (\neg \text{member1}(x, \text{lst}, i, j))$

THEOREM: bsearch1-not-found-2
 $(\text{orderedp}(\text{lst})$
 $\wedge \text{ilessp}(x, \text{get-nth}(k, \text{lst}))$
 $\wedge (k < j)$
 $\wedge (j \leq \text{len}(\text{lst})))$
 $\rightarrow (\text{member1}(x, \text{lst}, i, j) = \text{member1}(x, \text{lst}, i, k))$

EVENT: Disable bsearch1-not-found-2-lemma.

THEOREM: bsearch1-not-found
 $((\text{bsearch1}(x, \text{lst}, i, j) = -1)$
 $\wedge \text{orderedp}(\text{lst})$
 $\wedge \text{lst-integerp}(\text{lst})$
 $\wedge \text{integerp}(x)$
 $\wedge (j \leq \text{len}(\text{lst})))$
 $\rightarrow (\neg \text{member1}(x, \text{lst}, i, j))$

DEFINITION:
 $\text{member2}(x, \text{lst}, i, j)$
 $=$ **if** $i < j$
 then if $x = \text{get-nth}(i, \text{lst})$ **then t**
 else $\text{member2}(x, \text{cdr}(\text{lst}), i, j - 1)$ **endif**
 else f endif

THEOREM: member2-member
 $\text{member2}(x, \text{lst}, 0, \text{len}(\text{lst})) = (x \in \text{lst})$

THEOREM: member2-lemma

$$(x \neq \text{get-nth}(i, lst)) \rightarrow (\text{member2}(x, lst, 1 + i, j) = \text{member2}(x, lst, i, j))$$

THEOREM: member1-member2

$$\text{member1}(x, lst, i, j) = \text{member2}(x, lst, i, j)$$

THEOREM: bsearch-found

$$\begin{aligned} & ((\text{bsearch}(x, n, lst) \neq -1) \wedge \text{lst-integerp}(lst) \wedge \text{integerp}(x)) \\ \rightarrow & (\text{get-nth}(\text{bsearch}(x, n, lst), lst) = x) \end{aligned}$$

THEOREM: bsearch-not-found

$$\begin{aligned} & ((\text{bsearch}(x, \text{len}(lst), lst) = -1) \\ & \wedge \text{orderedp}(lst) \\ & \wedge \text{lst-integerp}(lst) \\ & \wedge \text{integerp}(x)) \\ \rightarrow & (x \notin lst) \end{aligned}$$

; an upper bound.

THEOREM: bsearch1-t-0

$$\text{bsearch1-t}(x, lst, i, i) = 6$$

THEOREM: bsearch1-t-crock

$$\begin{aligned} & ((i < j) \\ & \wedge ((26 + (13 * \log(2, (j - 1) - ((i + j) \div 2)))) \\ & \quad \not\leq \text{bsearch1-t}(x, lst, 1 + ((i + j) \div 2), j))) \\ \rightarrow & (((26 + (13 * \log(2, j - i))) \\ & \quad < (13 + \text{bsearch1-t}(x, lst, 1 + ((i + j) \div 2), j))) \\ & \quad = \text{f}) \end{aligned}$$

THEOREM: bsearch1-t-ubound

$$(26 + (13 * \log(2, j - i))) \not\leq \text{bsearch1-t}(x, lst, i, j)$$

EVENT: Disable bsearch1-t-crock.

THEOREM: bsearch-t-ubound-la

$$\text{bsearch-t}(x, n, lst) \leq (32 + (13 * \log(2, n)))$$

THEOREM: bsearch-t-ubound

$$(n < \exp(2, 31)) \rightarrow (\text{bsearch-t}(x, n, lst) \leq 435)$$

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