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|#

EVENT: Start with the library "gcd" using the compiled version.

#|

The following C program computes the greatest common divisor of three nonnegative integers a, b and c. We investigate the machine code of this program generated by a widely used C compiler gcc, and verify the correctness of the code. The aim here is to see how to handle subroutine calls.

```
gcd3(a, b, c)
long int a, b, c;
{
    gcd(gcd(a, b), c);
}
```

Here is the MC68020 assembly code of the above GCD program. The code is generated by gcc.

```

0x2324 <gcd3>:      linkw a6,#0
0x2328 <gcd3+4>:    movel a2,sp@-
0x232a <gcd3+6>:    movel a6@(16),sp@-
0x232e <gcd3+10>:   movel a6@(12),sp@-
0x2332 <gcd3+14>:   movel a6@(8),sp@-
0x2336 <gcd3+18>:   lea @#0x2350 <gcd>,a2
0x233c <gcd3+24>:   jsr a2@
0x233e <gcd3+26>:   addqw #8,sp
0x2340 <gcd3+28>:   movel d0,sp@-
0x2342 <gcd3+30>:   jsr a2@
0x2344 <gcd3+32>:   moveal a6@(-4),a2
0x2348 <gcd3+36>:   unlk a6
0x234a <gcd3+38>:   rts

```

The machine code of the above program is:

```

<gcd3>:      0x4e56  0x0000  0x2f0a  0x2f2e  0x0010  0x2f2e  0x000c  0x2f2e
<gcd3+16>:   0x0008  0x45f9  0x0000  0x2350  0x4e92  0x504f  0x2f00  0x4e92
<gcd3+32>:   0x246e  0xffff  0x4e5e  0x4e75

```

```

'(78      86      0      0      47      10      47      46
  0      16      47      46      0      12      47      46
  0       8      69     249     0       0      35     80
  78     146     80     79     47      0      78     146
  36     110     255     252     78     94      78     117)
|#

```

; now we start to verify this GCD3 program, defined by (gcd3-code).

DEFINITION:

GCD3-CODE

```

= '(78 86 0 0 47 10 47 46 0 16 47 46 0 12 47 46 0 8 69
    249 -1 -1 -1 -1 78 146 80 79 47 0 78 146 36 110 255
    252 78 94 78 117)

```

CONSERVATIVE AXIOM: gcd3-load

gcd3-loadp(s)

```

= (evenp (GCD3-ADDR)
   ∧ (GCD3-ADDR ∈ ℕ)
   ∧ nat-range (GCD3-ADDR, 32)
   ∧ rom-addrp (GCD3-ADDR, mc-mem(s), 40)
   ∧ mcode-addrp (GCD3-ADDR, mc-mem(s), GCD3-CODE)
   ∧ gcd-loadp(s)
   ∧ (pc-read-mem (add (32, GCD3-ADDR, 20), mc-mem(s), 4) = GCD-ADDR))

```

Simultaneously, we introduce the new function symbols *gcd3-loadp* and *gcd3-addr*.

THEOREM: *stepn-gcd3-loadp*

$$\text{gcd3-loadp}(\text{stepn}(s, n)) = \text{gcd3-loadp}(s)$$

DEFINITION:  $\text{gcd3}(a, b, c) = \text{gcd}(\text{gcd}(a, b), c)$

DEFINITION:  $\text{gcd3-t0}(a, b, c) = 7$

DEFINITION:  $\text{gcd3-t1}(a, b, c) = \text{gcd-t}(a, b)$

DEFINITION:  $\text{gcd3-t2}(a, b, c) = 3$

DEFINITION:  $\text{gcd3-t3}(a, b, c) = \text{gcd-t}(\text{gcd}(a, b), c)$

DEFINITION:  $\text{gcd3-t4}(a, b, c) = 3$

DEFINITION:

$$\begin{aligned} &\text{gcd3-t}(a, b, c) \\ = &\text{splus}(\text{gcd3-t0}(a, b, c), \\ &\quad \text{splus}(\text{gcd3-t1}(a, b, c), \\ &\quad \quad \text{splus}(\text{gcd3-t2}(a, b, c), \text{splus}(\text{gcd3-t3}(a, b, c), \text{gcd3-t4}(a, b, c)))))) \end{aligned}$$

; the initial state.

DEFINITION:

$$\begin{aligned} &\text{gcd3-statep}(s, a, b, c) \\ = &((\text{mc-status}(s) = \text{'running'}) \\ &\quad \wedge \text{gcd3-loadp}(s) \\ &\quad \wedge (\text{mc-pc}(s) = \text{GCD3-ADDR}) \\ &\quad \wedge \text{ram-addrp}(\text{sub}(32, 36, \text{read-sp}(s)), \text{mc-mem}(s), 52) \\ &\quad \wedge (a = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 4), \text{mc-mem}(s), 4)) \\ &\quad \wedge (b = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 8), \text{mc-mem}(s), 4)) \\ &\quad \wedge (c = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 12), \text{mc-mem}(s), 4)) \\ &\quad \wedge (0 < a) \\ &\quad \wedge (0 < b) \\ &\quad \wedge (0 < c)) \end{aligned}$$

; the state after the execution of the first JSR instruction, but before  
; the execution of the subroutine GCD.

DEFINITION:

$$\begin{aligned} &\text{gcd3-s0p}(s, a, b, c) \\ = &((\text{mc-status}(s) = \text{'running'}) \\ &\quad \wedge \text{gcd3-loadp}(s)) \end{aligned}$$

$\wedge$  (mc-pc( $s$ ) = GCD-ADDR)  
 $\wedge$  (read-an(32, 2,  $s$ ) = GCD-ADDR)  
 $\wedge$  (rts-addr( $s$ ) = add(32, GCD3-ADDR, 26))  
 $\wedge$  ram-addrp(sub(32, 12, read-sp( $s$ )), mc-mem( $s$ ), 52)  
 $\wedge$  equal\*(read-an(32, 6,  $s$ ), add(32, read-sp( $s$ ), 20))  
 $\wedge$  ( $a$  = iread-mem(add(32, read-sp( $s$ ), 4), mc-mem( $s$ ), 4))  
 $\wedge$  ( $b$  = iread-mem(add(32, read-sp( $s$ ), 8), mc-mem( $s$ ), 4))  
 $\wedge$  ( $c$  = iread-mem(add(32, read-sp( $s$ ), 12), mc-mem( $s$ ), 4))  
 $\wedge$  ( $0 < a$ )  
 $\wedge$  ( $0 < b$ )  
 $\wedge$  ( $0 < c$ )

; the state right after return from the first call to subroutine GCD.

DEFINITION:

gcd3-s1p( $s, a, b, c$ )

= ((mc-status( $s$ ) = 'running')  
 $\wedge$  gcd3-loadp( $s$ )  
 $\wedge$  (read-an(32, 2,  $s$ ) = GCD-ADDR)  
 $\wedge$  (mc-pc( $s$ ) = add(32, GCD3-ADDR, 26))  
 $\wedge$  ram-addrp(sub(32, 16, read-sp( $s$ )), mc-mem( $s$ ), 52)  
 $\wedge$  equal\*(read-an(32, 6,  $s$ ), add(32, read-sp( $s$ ), 16))  
 $\wedge$  (iread-dn(32, 0,  $s$ ) = gcd( $a, b$ ))  
 $\wedge$  ( $c$  = iread-mem(add(32, read-sp( $s$ ), 8), mc-mem( $s$ ), 4))  
 $\wedge$  ( $0 < a$ )  
 $\wedge$  ( $0 < b$ )  
 $\wedge$  ( $0 < c$ )

; the state after the execution of the second JSR, but before the  
 ; execution of the subroutine GCD.

DEFINITION:

gcd3-s2p( $s, a, b, c$ )

= ((mc-status( $s$ ) = 'running')  
 $\wedge$  gcd3-loadp( $s$ )  
 $\wedge$  (mc-pc( $s$ ) = GCD-ADDR)  
 $\wedge$  (rts-addr( $s$ ) = add(32, GCD3-ADDR, 32))  
 $\wedge$  ram-addrp(sub(32, 16, read-sp( $s$ )), mc-mem( $s$ ), 52)  
 $\wedge$  equal\*(read-an(32, 6,  $s$ ), add(32, read-sp( $s$ ), 16))  
 $\wedge$  (gcd( $a, b$ ) = iread-mem(add(32, read-sp( $s$ ), 4), mc-mem( $s$ ), 4))  
 $\wedge$  ( $c$  = iread-mem(add(32, read-sp( $s$ ), 8), mc-mem( $s$ ), 4))  
 $\wedge$  ( $0 < a$ )  
 $\wedge$  ( $0 < b$ )  
 $\wedge$  ( $0 < c$ )

; the state returned from the second call to the subroutine GCD.

DEFINITION:

gcd3-s3p( $s, a, b, c$ )  
 $=$  ((mc-status( $s$ ) = 'running')  
 $\wedge$  gcd3-loadp( $s$ )  
 $\wedge$  (mc-pc( $s$ ) = add(32, GCD3-ADDR, 32))  
 $\wedge$  ram-addrp(sub(32, 12, read-sp( $s$ )), mc-mem( $s$ ), 44)  
 $\wedge$  equal\*(read-an(32, 6,  $s$ ), add(32, read-sp( $s$ ), 12))  
 $\wedge$  (gcd(gcd( $a, b$ ),  $c$ ) = iread-dn(32, 0,  $s$ ))  
 $\wedge$  ( $0 < a$ )  
 $\wedge$  ( $0 < b$ )  
 $\wedge$  ( $0 < c$ ))

; from the initial state to s0.

THEOREM: gcd3-s-s0

**let**  $s0$  **be** stepn( $s$ , gcd3-t0( $a, b, c$ ))  
**in**  
gcd3-statep( $s, a, b, c$ )  
 $\rightarrow$  (gcd3-s0p( $s0, a, b, c$ )  
 $\wedge$  (linked-rtss-addr( $s0$ ) = rts-addr( $s$ ))  
 $\wedge$  (linked-a6( $s0$ ) = read-an(32, 6,  $s$ ))  
 $\wedge$  (read-rn(32, 14, mc-rfile( $s0$ ))  
 $=$  sub(32, 4, read-sp( $s$ )))  
 $\wedge$  (rn-saved( $s0$ ) = read-an(32, 2,  $s$ ))) **endlet**

THEOREM: gcd3-s-s0-rfile

(gcd3-statep( $s, a, b, c$ )  $\wedge$  d2-7a3-5p( $rn$ ))  
 $\rightarrow$  (read-rn( $oplen, rn, mc-rfile(stepn(s, gcd3-t0(a, b, c)))$ )  
 $=$  read-rn( $oplen, rn, mc-rfile(s)$ ))

THEOREM: gcd3-s-s0-mem

(gcd3-statep( $s, a, b, c$ )  $\wedge$  disjoint( $x, k$ , sub(32, 36, read-sp( $s$ )), 52))  
 $\rightarrow$  (read-mem( $x, mc-mem(stepn(s, gcd3-t0(a, b, c)))$ ,  $k$ )  
 $=$  read-mem( $x, mc-mem(s)$ ,  $k$ ))

; from s0 to s1.

THEOREM: gcd3-s0-s1

**let**  $s1$  **be** stepn( $s$ , gcd3-t1( $a, b, c$ ))  
**in**  
gcd3-s0p( $s, a, b, c$ )  
 $\rightarrow$  (gcd3-s1p( $s1, a, b, c$ )  
 $\wedge$  (linked-rtss-addr( $s1$ ) = linked-rtss-addr( $s$ )))

$\wedge$  (read-rn (32, 14, mc-rfile (s1))  
 $=$  read-rn (32, 14, mc-rfile (s)))  
 $\wedge$  (linked-a6 (s1) = linked-a6 (s))  
 $\wedge$  (rn-saved (s1) = rn-saved (s))) **endlet**

THEOREM: gcd3-s0-s1-rfile  
 $(\text{gcd3-s0p}(s, a, b, c) \wedge \text{d2-7a2-5p}(rn) \wedge (\text{oplen} \leq 32))$   
 $\rightarrow$  (read-rn (oplen, rn, mc-rfile (stepn (s, gcd3-t1 (a, b, c))))  
 $=$  read-rn (oplen, rn, mc-rfile (s)))

THEOREM: gcd3-s0-s1-mem  
 $(\text{gcd3-s0p}(s, a, b, c) \wedge \text{disjoint}(x, k, \text{sub}(32, 32, \text{read-an}(32, 6, s)), 52))$   
 $\rightarrow$  (read-mem (x, mc-mem (stepn (s, gcd3-t1 (a, b, c))), k)  
 $=$  read-mem (x, mc-mem (s), k))

; from s1 to s2.

THEOREM: gcd3-s1-s2  
**let** s2 **be** stepn (s, gcd3-t2 (a, b, c))  
**in**  
 gcd3-s1p (s, a, b, c)  
 $\rightarrow$  (gcd3-s2p (s2, a, b, c)  
 $\wedge$  (linked-rtts-addr (s2) = linked-rtts-addr (s))  
 $\wedge$  (read-rn (32, 14, mc-rfile (s2))  
 $=$  read-rn (32, 14, mc-rfile (s)))  
 $\wedge$  (linked-a6 (s2) = linked-a6 (s))  
 $\wedge$  (rn-saved (s2) = rn-saved (s))) **endlet**

THEOREM: gcd3-s1-s2-rfile  
 $(\text{gcd3-s1p}(s, a, b, c) \wedge \text{d2-7a3-5p}(rn))$   
 $\rightarrow$  (read-rn (oplen, rn, mc-rfile (stepn (s, gcd3-t2 (a, b, c))))  
 $=$  read-rn (oplen, rn, mc-rfile (s)))

THEOREM: gcd3-s1-s2-mem  
 $(\text{gcd3-s1p}(s, a, b, c) \wedge \text{disjoint}(x, k, \text{sub}(32, 32, \text{read-an}(32, 6, s)), 52))$   
 $\rightarrow$  (read-mem (x, mc-mem (stepn (s, gcd3-t2 (a, b, c))), k)  
 $=$  read-mem (x, mc-mem (s), k))

; from s2 to s3.

THEOREM: gcd-nonzero  
 $((a \neq 0) \wedge (b \neq 0)) \rightarrow (\text{gcd}(a, b) \neq 0)$

THEOREM: gcd3-s2-s3  
**let** s3 **be** stepn (s, gcd3-t3 (a, b, c))  
**in**

```

gcd3-s2p(s, a, b, c)
→ (gcd3-s3p(s3, a, b, c)
   ∧ (linked-rtts-addr(s3) = linked-rtts-addr(s))
   ∧ (read-rn(32, 14, mc-rfile(s3))
      = read-rn(32, 14, mc-rfile(s)))
   ∧ (linked-a6(s3) = linked-a6(s))
   ∧ (rn-saved(s3) = rn-saved(s))) endlet

```

THEOREM: gcd3-s2-s3-rfile  
 $(\text{gcd3-s2p}(s, a, b, c) \wedge \text{d2-7a2-5p}(rn) \wedge (\text{oplen} \leq 32))$   
 $\rightarrow (\text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(\text{stepn}(s, \text{gcd3-t3}(a, b, c))))$   
 $= \text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(s)))$

THEOREM: gcd3-s2-s3-mem  
 $(\text{gcd3-s2p}(s, a, b, c) \wedge \text{disjoint}(x, k, \text{sub}(32, 32, \text{read-an}(32, 6, s)), 52))$   
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, \text{gcd3-t3}(a, b, c))), k)$   
 $= \text{read-mem}(x, \text{mc-mem}(s), k))$

; from s3 to exit.

THEOREM: gcd3-s3-sn  
**let** *sn* **be**  $\text{stepn}(s, \text{gcd3-t4}(a, b, c))$   
**in**  
gcd3-s3p(*s*, *a*, *b*, *c*)  
 $\rightarrow ((\text{mc-status}(sn) = \text{'running'})$   
 $\wedge (\text{mc-pc}(sn) = \text{linked-rtts-addr}(s))$   
 $\wedge (\text{iread-dn}(32, 0, sn) = \text{gcd}(\text{gcd}(a, b), c))$   
 $\wedge (\text{read-rn}(32, 14, \text{mc-rfile}(sn)) = \text{linked-a6}(s))$   
 $\wedge (\text{read-rn}(32, 15, \text{mc-rfile}(sn))$   
 $= \text{add}(32, \text{read-an}(32, 6, s), 8)))$  **endlet**

THEOREM: gcd3-s3-sn-rfile  
 $(\text{gcd3-s3p}(s, a, b, c) \wedge (\text{oplen} \leq 32) \wedge \text{d2-7a2-5p}(rn))$   
 $\rightarrow (\text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(\text{stepn}(s, \text{gcd3-t4}(a, b, c))))$   
 $= \text{if } rn = 10 \text{ then head}(\text{rn-saved}(s), \text{oplen})$   
 $\text{else read-rn}(\text{oplen}, rn, \text{mc-rfile}(s)) \text{ endif})$

THEOREM: gcd3-s3-sn-mem  
 $(\text{gcd3-s3p}(s, a, b, c) \wedge \text{disjoint}(x, k, \text{sub}(32, 32, \text{read-an}(32, 6, s)), 52))$   
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, \text{gcd3-t4}(a, b, c))), k)$   
 $= \text{read-mem}(x, \text{mc-mem}(s), k))$

EVENT: Disable gcd3-t0.

EVENT: Disable gcd3-t1.

EVENT: Disable gcd3-t2.

EVENT: Disable gcd3-t3.

EVENT: Disable gcd3-t4.

; the correctness of the program GCD3.

THEOREM: gcd3-correctness

```
let sn be stepn(s, gcd3-t(a, b, c))
in
gcd3-statep(s, a, b, c)
→ ((mc-status(sn) = 'running)
   ∧ (mc-pc(sn) = rts-addr(s))
   ∧ (read-rn(32, 14, mc-rfile(sn))
      = read-rn(32, 14, mc-rfile(s)))
   ∧ (read-rn(32, 15, mc-rfile(sn))
      = add(32, read-rn(32, 15, mc-rfile(s)), 4))
   ∧ (((oplen ≤ 32) ∧ d2-7a2-5p(rn))
      → (read-rn(oplen, rn, mc-rfile(sn))
          = read-rn(oplen, rn, mc-rfile(s))))
   ∧ (disjoint(x, k, sub(32, 36, read-sp(s)), 52)
      → (read-mem(x, mc-mem(sn), k)
          = read-mem(x, mc-mem(s), k)))
   ∧ (iread-dn(32, 0, sn) = gcd(gcd(a, b), c))) endlet
```

EVENT: Disable gcd3-t.

; in the logic, the function gcd3 does computes the greatest common divisor  
; of its three arguments.

THEOREM: remainder-trans

$$(((a \bmod b) = 0) \wedge ((b \bmod c) = 0)) \rightarrow ((a \bmod c) = 0)$$

THEOREM: gcd3-is-cd

$$\begin{aligned} & ((a \bmod \text{gcd3}(a, b, c)) = 0) \\ \wedge & ((b \bmod \text{gcd3}(a, b, c)) = 0) \\ \wedge & ((c \bmod \text{gcd3}(a, b, c)) = 0) \end{aligned}$$

THEOREM: cd-divides-gcd

$$(((a \bmod x) = 0) \wedge ((b \bmod x) = 0)) \rightarrow ((\text{gcd}(a, b) \bmod x) = 0)$$



THEOREM: gcd3-the-greatest

$$\begin{aligned} & ((a \neq 0) \\ & \wedge (b \neq 0) \\ & \wedge (c \neq 0) \\ & \wedge ((a \bmod x) = 0) \\ & \wedge ((b \bmod x) = 0) \\ & \wedge ((c \bmod x) = 0)) \\ & \rightarrow (\text{gcd3}(a, b, c) \neq x) \end{aligned}$$

EVENT: Disable remainder-trans.

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