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|#

; Proof of the Correctness of the ISQRT Program
;

EVENT: Start with the library "mc20-2" using the compiled version.

#|

The following C function ISQRT computes the integer square root of a given nonnegative integer i.

```
/* computes the integer square root of a given nonnegative integer */  
isqrt(int i)  
{  
    int j;  
  
    j = (i / 2);  
    while ((i / j) < j)  
        j = (j + (i / j)) / 2;  
    return (j);  
}
```

Here is the MC68020 assembly code of the above ISQRT program. The code is generated by "gcc -O".

The program expects an nonnegative integer input which is on the sp stack at address (sp)+4. The local variable j is assigned to d2. The program exits with the answer in register d0 and the program counter equals the return address saved on the sp stack at address (sp).

```
0x2318 <isqrt>: linkw a6,#0
0x231c <isqrt+4>: movel d2,sp@-
0x231e <isqrt+6>: movel a6@(8),d1
0x2322 <isqrt+10>: movel d1,d2
0x2324 <isqrt+12>: bra 0x232e <isqrt+22>
0x2326 <isqrt+14>: movel d2,d0
0x2328 <isqrt+16>: divsll d1,d0,d0
0x232c <isqrt+20>: addl d0,d1
0x232e <isqrt+22>: tstl d1
0x2330 <isqrt+24>: bge 0x2334 <isqrt+28>
0x2332 <isqrt+26>: addql #1,d1
0x2334 <isqrt+28>: asrl #1,d1
0x2336 <isqrt+30>: movel d2,d0
0x2338 <isqrt+32>: divsll d1,d0,d0
0x233c <isqrt+36>: cmpl d0,d1
0x233e <isqrt+38>: bgt 0x2326 <isqrt+14>
0x2340 <isqrt+40>: movel d1,d0
0x2342 <isqrt+42>: movel a6@(-4),d2
0x2346 <isqrt+46>: unlk a6
0x2348 <isqrt+48>: rts
```

The machine code of the above program is:

```
<isqrt>: 0x4e56 0x0000 0x2f02 0x222e 0x0008 0x2401 0x6008 0x2002
<isqrt+16>: 0x4c41 0x0800 0xd280 0x4a81 0x6c02 0x5281 0xe281 0x2002
<isqrt+32>: 0x4c41 0x0800 0xb280 0x6ee6 0x2001 0x242e 0xffffc 0x4e5e
<isqrt+48>: 0x4e75
```

In the logic, it is going to be a list like:

'(78	86	0	0	47	2	34	46
0	8	36	1	96	8	32	2
76	65	8	0	210	128	74	129
108	2	82	129	226	129	32	2
76	65	8	0	178	128	110	230

```

32      1      36      46      255      252      78      94
78      117)
|#
; in the logic, the program is defined by (isqrt-code).

DEFINITION:
ISQRT-CODE
= '(78 86 0 0 47 2 34 46 0 8 36 1 96 8 32 2 76 65 8 0
  210 128 74 129 108 2 82 129 226 129 32 2 76 65 8 0
  178 128 110 230 32 1 36 46 255 252 78 94 78 117)

THEOREM: mean-lessp-1
(((b + a) ÷ 2) < b) = (a < b)

; isqrt1 is a function in the Logic simulating the loop of the above code.

DEFINITION:
isqrt1(i, j)
= if  $j \simeq 0$  then fix(i)
  elseif  $(i \div j) < j$  then isqrt1(i, (j + (i ÷ j)) ÷ 2)
  else fix(j) endif

; isqrt is a function in the Logic simulating the above code.

DEFINITION: isqrt(i) = isqrt1(i, i ÷ 2)

; the computation time of the loop.

DEFINITION:
isqrt1-t(i, j)
= if  $j \simeq 0$  then 0
  elseif  $(i \div j) < j$ 
  then splus(10, isqrt1-t(i, (j + (i ÷ j)) ÷ 2))
  else 8 endif

DEFINITION: isqrt-t(i) = splus(8, isqrt1-t(i, i ÷ 2))

EVENT: Enable iplus.

EVENT: Enable integerp.

EVENT: Enable iquotient.

EVENT: Enable ilessp.

```

EVENT: Disable remainder.

EVENT: Disable quotient.

THEOREM: isqrt-no-overflow

$$\begin{aligned} & (\text{int-rangep}(2 * j, n) \wedge ((i \div j) < j)) \\ \rightarrow & \quad \text{int-rangep}(j + (i \div j), n) \end{aligned}$$

THEOREM: j-nonzerop

$$((1 < i) \wedge (0 < j)) \rightarrow (((j + (i \div j)) \div 2) \neq 0)$$

THEOREM: j-int-rangep

$$\begin{aligned} & (\text{int-rangep}(2 * j, n) \wedge ((i \div j) < j)) \\ \rightarrow & \quad \text{int-rangep}(2 * ((j + (i \div j)) \div 2), n) \end{aligned}$$

; an induction hint.

DEFINITION:

$$\begin{aligned} & \text{isqrt-induct}(s, i, j) \\ = & \quad \text{if } j \simeq 0 \text{ then t} \\ & \quad \text{elseif } (i \div j) < j \\ & \quad \text{then isqrt-induct}(\text{stepn}(s, 10), i, (j + (i \div j)) \div 2) \\ & \quad \text{else t endif} \end{aligned}$$

DEFINITION:

$$\begin{aligned} & \text{isqrt-statep}(s, i) \\ = & \quad ((\text{mc-status}(s) = \text{'running}) \\ & \quad \wedge \text{evenp}(\text{mc-pc}(s)) \\ & \quad \wedge \text{rom-addrp}(\text{mc-pc}(s), \text{mc-mem}(s), 50) \\ & \quad \wedge \text{mcode-addrp}(\text{mc-pc}(s), \text{mc-mem}(s), \text{ISQRT-CODE}) \\ & \quad \wedge \text{ram-addrp}(\text{sub}(32, 8, \text{read-sp}(s)), \text{mc-mem}(s), 16) \\ & \quad \wedge (i = \text{iread-mem}(\text{add}(32, \text{read-sp}(s), 4), \text{mc-mem}(s), 4)) \\ & \quad \wedge \text{ilessp}(1, i)) \end{aligned}$$

DEFINITION:

$$\begin{aligned} & \text{isqrt-s0p}(s, i, j) \\ = & \quad ((\text{mc-status}(s) = \text{'running}) \\ & \quad \wedge \text{evenp}(\text{mc-pc}(s)) \\ & \quad \wedge \text{rom-addrp}(\text{sub}(32, 30, \text{mc-pc}(s)), \text{mc-mem}(s), 50) \\ & \quad \wedge \text{mcode-addrp}(\text{sub}(32, 30, \text{mc-pc}(s)), \text{mc-mem}(s), \text{ISQRT-CODE}) \\ & \quad \wedge \text{ram-addrp}(\text{sub}(32, 4, \text{read-an}(32, 6, s)), \text{mc-mem}(s), 16) \\ & \quad \wedge (i = \text{iread-dn}(32, 2, s)) \\ & \quad \wedge (j = \text{iread-dn}(32, 1, s)) \\ & \quad \wedge \text{int-rangep}(2 * j, 32) \\ & \quad \wedge \text{ilessp}(1, i) \\ & \quad \wedge \text{ilessp}(0, j)) \end{aligned}$$

; from the initial state to s0.

THEOREM: initial-j-int-rangep
 $\text{int-rangep}(i, n) \rightarrow \text{int-rangep}(2 * (i \div 2), n)$

THEOREM: isqrt-s-s0

$\text{isqrt-statep}(s, i)$
 $\rightarrow (\text{isqrt-s0p}(\text{stepn}(s, 8)), i, i \div 2)$
 $\wedge (\text{linked-rts-addr}(\text{stepn}(s, 8)) = \text{rts-addr}(s))$
 $\wedge (\text{linked-a6}(\text{stepn}(s, 8)) = \text{read-an}(32, 6, s))$
 $\wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, 8)))$
 $= \text{sub}(32, 4, \text{read-sp}(s)))$
 $\wedge (\text{rn-saved}(\text{stepn}(s, 8)) = \text{read-dn}(32, 2, s)))$

THEOREM: isqrt-s-s0-rfile

$(\text{isqrt-statep}(s, i) \wedge \text{d3-7a2-5p}(rn))$
 $\rightarrow (\text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(\text{stepn}(s, 8))))$
 $= \text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(s)))$

THEOREM: isqrt-s-s0-mem

$(\text{isqrt-statep}(s, i) \wedge \text{disjoint}(x, k, \text{sub}(32, 8, \text{read-sp}(s)), 16))$
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 8)), k) = \text{read-mem}(x, \text{mc-mem}(s), k))$

; from s0 to s0 (induction csae), or from s0 to exit (base case).

; base case:

; the basics we need to prove:

; 0. the machine is still running.
; 1. the pc is at the right position.
; 2. d0 contains the local variable j, the computing square root.
; 3. a6 should have the right value.
; 4. sp(a7) should have the right value.

THEOREM: isqrt-s0-sn

$(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) \not< j))$
 $\rightarrow ((\text{mc-status}(\text{stepn}(s, 8)) = \text{'running})$
 $\wedge (\text{mc-pc}(\text{stepn}(s, 8)) = \text{linked-rts-addr}(s))$
 $\wedge (\text{iread-dn}(32, 0, \text{stepn}(s, 8)) = j)$
 $\wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, 8))) = \text{linked-a6}(s))$
 $\wedge (\text{read-rn}(32, 15, \text{mc-rfile}(\text{stepn}(s, 8)))$
 $= \text{add}(32, \text{read-an}(32, 6, s), 8)))$

; 5. d2, which is used by this program, should be restored.

THEOREM: isqrt-s0-sn-d2

$(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) \not< j) \wedge (\text{oplen} \leq 32))$
 $\rightarrow (\text{read-rn}(\text{oplen}, 2, \text{mc-rfile}(\text{stepn}(s, 8))) = \text{head}(\text{rn-saved}(s), \text{oplen}))$

; and 6. those registers untouched by this program should still
; have their original contents.

THEOREM: isqrt-s0-sn-rfile
 $(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) \not< j) \wedge \text{d3-7a2-5p}(rn))$
 $\rightarrow (\text{read-rn}(oplen, rn, \text{mc-rfile}(\text{stepn}(s, 8))))$
 $= \text{read-rn}(oplen, rn, \text{mc-rfile}(s)))$

; and 7. the memory is correctly changed.

THEOREM: isqrt1-s0-sn-mem
 $(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) \not< j))$
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 8))), k) = \text{read-mem}(x, \text{mc-mem}(s), k))$

THEOREM: isqrt-s0p-j._nonzerop
 $\text{isqrt-s0p}(s, i, j) \rightarrow ((j \neq 0) \wedge (j \in \mathbb{N}))$

; induction case: s0 --> s0.
; we need to prove:
; 0. the state predicate at s0 is still satisfied.
; 1. a6 is unchanged.
; 2. the content of original a6 on the sp stack is not changed.
; 3. the return address on the sp attack is not changed.
; 4. the content of original d2 on the sp stack is not changed.

THEOREM: isqrt-s0-s0
 $(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) < j))$
 $\rightarrow (\text{isqrt-s0p}(\text{stepn}(s, 10), i, (j + (i \div j)) \div 2)$
 $\wedge (\text{read-rn}(oplen, 14, \text{mc-rfile}(\text{stepn}(s, 10))))$
 $= \text{read-rn}(oplen, 14, \text{mc-rfile}(s)))$
 $\wedge (\text{linked-a6}(\text{stepn}(s, 10)) = \text{linked-a6}(s))$
 $\wedge (\text{linked-rts-addr}(\text{stepn}(s, 10)) = \text{linked-rts-addr}(s))$
 $\wedge (\text{rn-saved}(\text{stepn}(s, 10)) = \text{rn-saved}(s)))$

; and 5. the registers that are not modified by this program still have
; their previous values.

THEOREM: isqrt-s0-s0-rfile
 $(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) < j) \wedge \text{d3-7a2-5p}(rn))$
 $\rightarrow (\text{read-rn}(oplen, rn, \text{mc-rfile}(\text{stepn}(s, 10))))$
 $= \text{read-rn}(oplen, rn, \text{mc-rfile}(s)))$

; and 6. the memory is correctly changed.

THEOREM: isqrt-s0-s0-mem
 $(\text{isqrt-s0p}(s, i, j) \wedge ((i \div j) < j))$
 $\rightarrow (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, 10))), k) = \text{read-mem}(x, \text{mc-mem}(s), k))$

EVENT: Disable isqrt-statep.

EVENT: Disable isqrt-s0p.

THEOREM: isqrt1-correctness

$$\begin{aligned} & \text{isqrt-s0p}(s, i, j) \\ \rightarrow & ((\text{mc-status}(\text{stepn}(s, \text{isqrt1-t}(i, j))) = \text{'running}) \\ & \wedge (\text{mc-pc}(\text{stepn}(s, \text{isqrt1-t}(i, j))) = \text{linked-rts-addr}(s)) \\ & \wedge (\text{iread-dn}(32, 0, \text{stepn}(s, \text{isqrt1-t}(i, j))) = \text{isqrt1}(i, j)) \\ & \wedge (\text{read-rn}(32, 14, \text{mc-rfile}(\text{stepn}(s, \text{isqrt1-t}(i, j)))) \\ & \quad = \text{linked-a6}(s)) \\ & \wedge (\text{read-rn}(32, 15, \text{mc-rfile}(\text{stepn}(s, \text{isqrt1-t}(i, j)))) \\ & \quad = \text{add}(32, \text{read-an}(32, 6, s), 8)) \\ & \wedge (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, \text{isqrt1-t}(i, j))), k) \\ & \quad = \text{read-mem}(x, \text{mc-mem}(s), k))) \end{aligned}$$

THEOREM: isqrt1-d2

$$\begin{aligned} & (\text{isqrt-s0p}(s, i, j) \wedge (\text{oplen} \leq 32)) \\ \rightarrow & (\text{read-rn}(\text{oplen}, 2, \text{mc-rfile}(\text{stepn}(s, \text{isqrt1-t}(i, j))))) \\ & = \text{head}(\text{rn-saved}(s), \text{oplen})) \end{aligned}$$

THEOREM: isqrt1-rfile

$$\begin{aligned} & (\text{isqrt-s0p}(s, i, j) \wedge \text{d3-7a2-5p}(rn)) \\ \rightarrow & (\text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(\text{stepn}(s, \text{isqrt1-t}(i, j))))) \\ & = \text{read-rn}(\text{oplen}, rn, \text{mc-rfile}(s))) \end{aligned}$$

; after an execution of this program, the machine state satisfies:
; 0. normal exit.
; 1. the pc is returned to the next instruction of the caller.
; 2. the result -- ISQRT(i), is stored in D0.
; 3. a6, used by LINK, is restored to its original content.
; 4. the stack pointer sp(a7) is updated correctly to pop off one frame.

THEOREM: isqrt-correctness

$$\begin{aligned} & \text{isqrt-statep}(s, i) \\ \rightarrow & ((\text{mc-status}(\text{stepn}(s, \text{isqrt-t}(i))) = \text{'running}) \\ & \wedge (\text{mc-pc}(\text{stepn}(s, \text{isqrt-t}(i))) = \text{rts-addr}(s)) \\ & \wedge (\text{iread-dn}(32, 0, \text{stepn}(s, \text{isqrt-t}(i))) = \text{isqrt}(i)) \\ & \wedge (\text{read-an}(32, 6, \text{stepn}(s, \text{isqrt-t}(i))) = \text{read-an}(32, 6, s)) \\ & \wedge (\text{read-an}(32, 7, \text{stepn}(s, \text{isqrt-t}(i))) \\ & \quad = \text{add}(32, \text{read-an}(32, 7, s), 4))) \\ ; & 5. \text{ d2, used for local variable } j, \text{ is restored to its original value.} \end{aligned}$$

THEOREM: isqrt-d2

$$\begin{aligned} & (\text{isqrt-statep}(s, i) \wedge (oplen \leq 32)) \\ \rightarrow & \quad (\text{read-rn}(oplen, 2, \text{mc-rfile}(\text{stepn}(s, \text{isqrt-t}(i))))) \\ = & \quad \text{read-dn}(oplen, 2, s) \end{aligned}$$

; and 6. the registers that are not modified by this program still have
; their original values.

THEOREM: isqrt-rfile

$$\begin{aligned} & (\text{isqrt-statep}(s, i) \wedge \text{d3-7a2-5p}(rn)) \\ \rightarrow & \quad (\text{read-rn}(oplen, rn, \text{mc-rfile}(\text{stepn}(s, \text{isqrt-t}(i))))) \\ = & \quad \text{read-rn}(oplen, rn, \text{mc-rfile}(s)) \end{aligned}$$

; and 7. the memory is correctly changed. i.e. specify clearly the portions
; of the memory were changed and prove that is a truth.

THEOREM: isqrt-read-mem

$$\begin{aligned} & (\text{isqrt-statep}(s, i) \wedge \text{disjoint}(x, k, \text{sub}(32, 8, \text{read-sp}(s)), 16)) \\ \rightarrow & \quad (\text{read-mem}(x, \text{mc-mem}(\text{stepn}(s, \text{isqrt-t}(i))), k) \\ = & \quad \text{read-mem}(x, \text{mc-mem}(s), k)) \end{aligned}$$

EVENT: Disable isqrt-t.

```
#|
; we next need to prove that isqrt does compute the integer square root.
; the following is the same as the proof in the file isqrt-ada.events.
(defn sq (j)
  (times j j))

(prove-lemma isqrt1-lower-bound (rewrite)
  (implies (not (zerop j))
            (not (lessp i (sq (isqrt1 i j))))))

(prove-lemma quotient-by-2 (rewrite)
  (not (lessp (plus (quotient x 2) (quotient x 2))
              (sub1 x)))))

(prove-lemma main-trick (rewrite)
  (not (lessp (sq (add1 (quotient (plus j k) 2)))
              (plus (times j k) j)))
       ((induct (difference j k)))))

(prove-lemma sq-add1-non-zero (rewrite)
  (not (equal (sq (add1 x)) 0)))
```

```

(prove-lemma main (rewrite)
  (implies (not (zerop j))
    (lessp i
      (sq (add1 (quotient (plus j (quotient i j)) 2))))))
  ((disable sq)))

(prove-lemma isqrt1-upper-bound (rewrite)
  (implies (lessp i (sq (add1 j)))
    (lessp i (sq (add1 (isqrt1 i j))))))
  ((disable sq)))

(prove-lemma isqrt->isqrt1 (rewrite)
  (implies (lessp 1 i)
    (lessp i (sq (add1 (quotient i 2))))))

; (isqrt i) is the square root of i: (isqrt i)^2 <= i < [(isqrt i)+1]^2.
(prove-lemma isqrt-logic-correctness ()
  (implies (lessp 1 i)
    (and (lessp i (sq (add1 (isqrt i)))))
      (not (lessp i (sq (isqrt i)))))))
  ((disable sq isqrt1)))

|#

```

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