EVENT: Start with the library "c-proc-call2".

```
;; I know that in the absence of resource-errors the two MG interpreters
;; and the two clocks are the same. Therefore, after I have proved everything
;; with the resource-error versions, I automatically gain the non-resource-error
;; versions.
;;
                                                                       ;;
                             EXACT TIME LEMMA
;;
                                                                       ;;
;;
                                                                       ;;
THEOREM: not-resource-errorp-not-zerop-n
(ok-mg-statement (stmt, r-cond-list, name-alist, proc-list)
 \land (\neg resource-errorp (mg-meaning-r (stmt,
                               proc-list,
                               mg-state,
                               n,
                               list (length (temp-stk)),
                                  p-ctrl-stk-size(ctrl-stk))))))
\rightarrow (n \not\simeq 0)
(defn exact-time-induction-hint (cinfo r-cond-list t-cond-list stmt
    proc-list mg-state n code2 subr
                 ctrl-stk temp-stk name-alist)
      (if (zerop n)
          t
      (if (resources-inadequatep stmt proc-list
  (list (length temp-stk) (p-ctrl-stk-size ctrl-stk)))
  t
      (if (equal (car stmt) 'no-op-mg)
          t.
      (if (equal (car stmt) 'signal-mg)
          t.
      (if (equal (car stmt) 'prog2-mg)
          (and (exact-time-induction-hint
 cinfo
 r-cond-list
 t-cond-list
```

```
(prog2-left-branch stmt)
 proc-list
 mg-state
  (sub1 n)
  (append (code (translate (nullify (translate (nullify cinfo)
       t-cond-list
       (prog2-left-branch stmt)
      proc-list))
  t-cond-list
   (prog2-right-branch stmt)
  proc-list))
 code2)
 subr ctrl-stk temp-stk name-alist)
                (exact-time-induction-hint
  (translate cinfo t-cond-list (prog2-left-branch stmt) proc-list)
                  r-cond-list
 t-cond-list
  (prog2-right-branch stmt)
 proc-list
  (mg-meaning-r (prog2-left-branch stmt) proc-list mg-state (sub1 n)
(list (length temp-stk) (p-ctrl-stk-size ctrl-stk)))
  (sub1 n)
 code2
 subr ctrl-stk temp-stk name-alist))
       (if (equal (car stmt) 'loop-mg)
           (and (exact-time-induction-hint
                   (make-cinfo (append (code cinfo)
       (list (list 'dl (label-cnt cinfo) nil '(no-op))))
       (cons (cons 'leave (add1 (label-cnt cinfo)))
     (label-alist cinfo))
       (add1 (add1 (label-cnt cinfo))))
                  (cons 'leave r-cond-list) t-cond-list
  (loop-body stmt)
 proc-list
 mg-state
 (sub1 n)
  (cons (list 'jump (label-cnt cinfo))
(cons (list 'dl (add1 (label-cnt cinfo)) nil '(push-constant (nat 2)))
      (cons '(pop-global c-c)
   code2)))
 subr ctrl-stk temp-stk name-alist)
                (exact-time-induction-hint
 cinfo
                  (cons 'leave r-cond-list) t-cond-list
```

```
stmt
 proc-list
  (mg-meaning-r (loop-body stmt) proc-list mg-state (sub1 n)
(list (length temp-stk) (p-ctrl-stk-size ctrl-stk)))
  (sub1 n)
 code2
 subr ctrl-stk temp-stk name-alist))
       (if (equal (car stmt) 'if-mg)
            (and (exact-time-induction-hint
   (make-cinfo
     (append (code cinfo)
    (list (list 'push-local (if-condition stmt))
  '(fetch-temp-stk)
                                   (list 'test-bool-and-jump 'false (label-cnt cinfo))))
     (label-alist cinfo)
     (add1 (add1 (label-cnt cinfo))))
                   r-cond-list t-cond-list
   (if-true-branch stmt)
                   proc-list
  mg-state
  (sub1 n)
   (cons (list 'jump (add1 (label-cnt cinfo)))
 (cons (list 'dl (label-cnt cinfo) nil '(no-op))
       (append (code (translate
       (nullify
 (translate
   (make-cinfo
    nil
     (label-alist cinfo)
     (add1 (add1 (label-cnt cinfo))))
                                                    t-cond-list
   (if-true-branch stmt)
  proc-list))
       t-cond-list
       (if-false-branch stmt)
      proc-list))
       (cons (list 'dl (add1 (label-cnt cinfo)) nil '(no-op))
     code2))))
  subr ctrl-stk temp-stk name-alist)
                (exact-time-induction-hint
  (add-code
    (translate
      (make-cinfo
(append (code cinfo)
```

```
(list (list 'push-local (if-condition stmt))
  '(fetch-temp-stk)
                                   (list 'test-bool-and-jump 'false (label-cnt cinfo))))
(label-alist cinfo)
(add1 (add1 (label-cnt cinfo))))
     t-cond-list
      (if-true-branch stmt)
     proc-list)
    (list (list 'jump (add1 (label-cnt cinfo)))
  (list 'dl (label-cnt cinfo) nil '(no-op))))
                  r-cond-list t-cond-list
                  (if-false-branch stmt)
 proc-list
                  mg-state
                  (sub1 n)
                  (cons (list 'dl (add1 (label-cnt cinfo)) nil '(no-op))
                        code2)
 subr ctrl-stk temp-stk name-alist))
        (if (equal (car stmt) 'begin-mg)
            (and
              (exact-time-induction-hint
(make-cinfo (code cinfo)
    (append (make-label-alist (when-labels stmt)
      (label-cnt cinfo))
    (label-alist cinfo))
    (add1 (add1 (label-cnt cinfo))))
                (append (when-labels stmt) r-cond-list)
                t-cond-list
(begin-body stmt)
proc-list
mg-state
(sub1 n)
(cons (list 'jump (add1 (label-cnt cinfo)))
      (cons (list 'dl (label-cnt cinfo) nil '(push-constant (nat 2)))
    (cons '(pop-global c-c)
  (append
    (code (translate
    (nullify
                                               (set-label-alist
(translate
 (make-cinfo (code cinfo)
     (append (make-label-alist (when-labels stmt)
       (label-cnt cinfo))
     (label-alist cinfo))
```

```
(add1 (add1 (label-cnt cinfo))))
  t-cond-list
  (begin-body stmt)
 proc-list)
(label-alist cinfo)))
                                             t-cond-list
    (when-handler stmt)
    proc-list))
    (cons (list 'dl (add1 (label-cnt cinfo)) nil '(no-op))
  code2)))))
subr ctrl-stk temp-stk name-alist)
              (exact-time-induction-hint
                 (add-code
  (set-label-alist
   (translate
                    (make-cinfo (code cinfo)
(append (make-label-alist (when-labels stmt)
  (label-cnt cinfo))
(label-alist cinfo))
(add1 (add1 (label-cnt cinfo))))
   t-cond-list
    (begin-body stmt)
   proc-list)
   (label-alist cinfo))
  (list (list 'jump (add1 (label-cnt cinfo)))
(list 'dl (label-cnt cinfo) nil '(push-constant (nat 2)))
'(pop-global c-c)))
                r-cond-list
t-cond-list
(when-handler stmt) proc-list
(set-condition (mg-meaning-r (begin-body stmt) proc-list mg-state (sub1 n)
     (list (length temp-stk) (p-ctrl-stk-size ctrl-stk)))
       'normal)
(sub1 n)
(cons (list 'dl (add1 (label-cnt cinfo)) nil '(no-op))
 code2)
                subr ctrl-stk temp-stk name-alist))
       (if (equal (car stmt) 'proc-call-mg)
           (exact-time-induction-hint
             (make-cinfo nil
     (cons (cons 'routineerror 0)
   (make-label-alist (make-cond-list (fetch-called-def stmt proc-list)) 0))
     1)
 (make-cond-list (fetch-called-def stmt proc-list))
```

```
(make-cond-list (fetch-called-def stmt proc-list))
                 (def-body (fetch-called-def stmt proc-list))
proc-list
 (make-call-environment mg-state stmt (fetch-called-def stmt proc-list))
 (sub1 n)
 (cons '(dl 0 nil (no-op))
       (cons (list 'pop* (data-length (def-locals (fetch-called-def stmt proc-list))))
             '((ret))))
 (call-name stmt)
                 (cons (p-frame
(make-frame-alist (fetch-called-def stmt proc-list) stmt ctrl-stk temp-stk)
(tag 'pc (cons subr
       (add1 (PLUS (LENGTH (CODE CINFO))
   (data-length (DEF-LOCALS (FETCH-CALLED-DEF
       STMT PROC-LIST)))
   (length (DEF-LOCALS (FETCH-CALLED-DEF
STMT PROC-LIST)))
   (LENGTH (CALL-ACTUALS STMT)))))))
       ctrl-stk)
                 (append (reverse (mg-to-p-local-values (def-locals (fetch-called-def stmt)
 (map-down-values (mg-alist mg-state)
  (bindings (top ctrl-stk))
  temp-stk))
 (make-name-alist (fetch-called-def stmt proc-list)))
        (if (equal (car stmt) 'predefined-proc-call-mg)
    t.
     f)))))))))))
    ((lessp (COUNT n)))
    (INSTRUCTIONS (BASH (ENABLE WHEN-HANDLER BEGIN-BODY IF-TRUE-BRANCH
IF-FALSE-BRANCH LOOP-BODY PROG2-RIGHT-BRANCH
PROG2-LEFT-BRANCH))))
;; The cond-list is required in the translation so that
;; I can convert between MG and Piton conditions; in the recognizer it is only
;; required that I have some set of conditions which could be signalled. It had better
   be that any signalled are on the translator cond-list or I won't be able to do the
::
;; mapping. However, I can sometimes signal 'leave and 'routineerror even though
;; these are not on the translator cond-list. This is because their map functions are
;; computed independently of the list. Therefore the appropriate relation between the
   recognizer-cond-alist and translator-cond-alist is that (cond-subsetp rec-list trans-li
;;
```

```
THEOREM: exact-time-lemma
(ok-mg-statement (stmt, r-cond-list, name-alist, proc-list)
```

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-translation-parameters (*cinfo*, *t-cond-list*, *stmt*, *proc-list*, *code2*)
- \land ok-mg-statep (*mg-state*, *r-cond-list*)
- $\land \quad \text{cond-subsetp}\left(\textit{r-cond-list}, \textit{ t-cond-list}\right)$
- \land (code (translate-def-body (assoc (*subr*, *proc-list*), *proc-list*))
 - = append (code (translate (*cinfo*, *t-cond-list*, *stmt*, *proc-list*)), $code_2$))
- \land user-defined-procp (*subr*, *proc-list*)
- \wedge plistp (*temp-stk*)
- \wedge listp (*ctrl-stk*)
- \wedge mg-vars-list-ok-in-p-state (mg-alist (*mg-state*),

bindings
$$(top (ctrl-stk)),$$

temp-stk)

- \land no-p-aliasing (bindings (top (*ctrl-stk*)), mg-alist (*mg-state*))
- \land signatures-match (mg-alist (*mg-state*), *name-alist*)
- $\wedge \quad \text{normal} (mg\text{-state})$
- \wedge all-cars-unique (mg-alist (*mg-state*))
- \land (\neg resource-errorp (mg-meaning-r (*stmt*,

```
proc-list,
```

mg-state, n,

```
list (length (temp-stk),
```

```
p-ctrl-stk-size(ctrl-stk))))))
```

- \rightarrow (p (map-down (*mg-state*,
 - $proc{-list},$
 - $\mathit{ctrl-stk},$
 - temp-stk,

tag('pc, cons(*subr*, length(code(*cinfo*)))),

t-cond-list),

```
clock(stmt, proc-list, mg-state, n))
```

= p-state (tag ('pc,

```
\cos\left(subr\right)
```

if normal (mg-meaning-r (stmt,

proc-list,

mg-state, n,

list (length (*temp-stk*),

p-ctrl-stk-size(*ctrl-stk*))))

then length (code (translate (*cinfo*,

t-cond-list,

stmt,

proc-list)))

else find-label (fetch-label (cc (mg-meaning-r (stmt,

proc-list,

```
mg-state,
                                                                              n,
                                                                              list (length (temp-stk),
                                                                                   p-ctrl-stk-size(ctrl-stk)))),
                                                           label-alist (translate (cinfo,
                                                                                 t-cond-list,
                                                                                 stmt,
                                                                                 proc-list))),
                                              append (code (translate (cinfo,
                                                                        t-cond-list,
                                                                        stmt,
                                                                        proc-list)),
                                                        code2)) endif)),
                   ctrl-stk,
                   map-down-values (mg-alist (mg-meaning-r (stmt,
                                                                 proc-list,
                                                                 mg-state,
                                                                 n,
                                                                list (length (temp-stk),
                                                                     p-ctrl-stk-size(ctrl-stk)))),
                                       bindings (top (ctrl-stk)),
                                       temp-stk),
                   translate-proc-list (proc-list),
                   list (list ('c-c,
                            mg-cond-to-p-nat (cc (mg-meaning-r (stmt,
                                                                   proc-list,
                                                                    mg-state,
                                                                   n,
                                                                   list (length (temp-stk),
                                                                        p-ctrl-stk-size(ctrl-stk)))),
                                                t-cond-list))),
                   MG-MAX-CTRL-STK-SIZE,
                   MG-MAX-TEMP-STK-SIZE,
                   MG-WORD-SIZE,
                    'run))
EVENT: Disable exact-time-lemma.
```

THEOREM: exact-time-lemma2

(ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-translation-parameters (*cinfo*, *t-cond-list*, *stmt*, *proc-list*, *code2*)
- \land ok-mg-statep (*mg-state*, *r-cond-list*)
- \land cond-subsetp (*r-cond-list*, *t-cond-list*)

(code (translate-def-body (assoc (*subr*, *proc-list*), *proc-list*)) \wedge append (code (translate (*cinfo*, *t-cond-list*, *stmt*, *proc-list*)), = code2))user-defined-procp (*subr*, *proc-list*) Λ \wedge plistp (temp-stk) \wedge listp(ctrl-stk)mg-vars-list-ok-in-p-state (mg-alist (mg-state), \wedge bindings (top (*ctrl-stk*)), temp-stk) no-p-aliasing (bindings (top (ctrl-stk)), mg-alist (mq-state)) \wedge \land signatures-match (mg-alist (*mg-state*), *name-alist*) \wedge normal (*mg-state*) \land all-cars-unique (mg-alist (*mg-state*)) $(\neg$ resource-errorp (mg-meaning-r (*stmt*, \wedge proc-list, mg-state, n,list (length (temp-stk)),p-ctrl-stk-size(*ctrl-stk*))))) (offset = length(code(cinfo)))) \wedge (p (map-down (*mg-state*, proc-list, ctrl-stk, temp-stk, tag('pc, cons(*subr*, *offset*)), *t*-cond-list), clock(stmt, proc-list, mq-state, n))= p-state (tag ('pc, $\cos(subr,$ if normal (mg-meaning-r (*stmt*, proc-list, mg-state, n,list (length (temp-stk), p-ctrl-stk-size(*ctrl-stk*)))) then length (code (translate (*cinfo*, t-cond-list, stmt, proc-list))) else find-label (fetch-label (cc (mg-meaning-r (*stmt*, proc-list, mg-state, n,

```
p-ctrl-stk-size(ctrl-stk)))),
                                       label-alist (translate (cinfo,
                                                              t-cond-list,
                                                              stmt,
                                                              proc-list))),
                           append (code (translate (cinfo,
                                                     t-cond-list,
                                                     stmt,
                                                     proc-list)),
                                    code2)) endif)),
ctrl-stk,
map-down-values (mg-alist (mg-meaning-r (stmt,
                                             proc-list,
                                             mg-state,
                                             n,
                                             list (length (temp-stk)),
                                                 p-ctrl-stk-size(ctrl-stk)))),
                   bindings (top (ctrl-stk)),
                   temp-stk),
translate-proc-list (proc-list),
list(list('c-c,
         mg-cond-to-p-nat (cc (mg-meaning-r (stmt,
                                                proc-list,
                                                mg-state,
                                                n,
                                                list (length (temp-stk)),
                                                     p-ctrl-stk-size(ctrl-stk)))),
                             t-cond-list))),
MG-MAX-CTRL-STK-SIZE,
MG-MAX-TEMP-STK-SIZE,
MG-WORD-SIZE,
'run))
```

```
EVENT: Disable exact-time-lemma2.
```

```
;; induction work. The user supplies only a single list with his input statement.
;; That list must be an identifier-plistp and hence may not contain 'leave.
;;
;; The hypotheses of this theorem guarantee the following:
      1. mg-state is of the form <alist current-condition>
;;
         a. the alist is of the form < <name type value> ... >
;;
         b. the condition is 'normal, 'routineerror, 'timed-out, or
;;
               is in cond-list
;;
     2. proc-list is a syntactically legal list of micro-gypsy procedures (which may
;;
               be mutually recursive)
;;
     3. stmt is a syntactically legal micro-Gypsy statement with respect to proc-list
;;
               with conditions from cond-list and variables from (mg-alist mg-state)
;;
      4. the translation environment described by cinfo is legitimate
::
         a. cinfo is of the form <code label-alist label-cnt cond-list>, where
;;
                 cond-list and code are proper lists
;;
         b. cond-list is a proper list containing only legal identifiers, 'leave, or
;;
                 'routineerror
;;
         c. translation stmt with cinfo does not generate any labels which would
;;
                 duplicate labels in (code cinfo) or code2
;;
     5. subr is the name of some procedure in proc-list
::
     6. the body of procedure subr is the context in which stmt "lives"; that is
;;
             the translation of the body of subr is equal to the append of the
;;
             the translation of stmt with code2
;;
     7. execution of stmt does not yield the condition timed-out.
;;
```

THEOREM: cons-preserves-cond-subsetp cond-subsetp $(y, z) \rightarrow$ cond-subsetp $(y, \cos(x, z))$

THEOREM: cond-subsetp-reflexive ok-cond-list (cond-list) \rightarrow cond-subsetp (cond-list, cond-list)

THEOREM: mg-meaning-preserves-signatures-match3

(ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-mg-statep (*mg-state*, cond-list)
- \land signatures-match (mg-alist (*mg-state*), *name-alist*))
- \rightarrow signatures-match (mg-alist (*mg-state*),

mg-alist (mg-meaning (stmt, proc-list, mg-state, n)))

THEOREM: mg-state-decomposition

 $(\neg \text{ resource-errorp}(state))$

 \rightarrow (mg-state (cc (*state*), mg-alist (*state*), 'run) = *state*)

THEOREM: signatures-match-implies-signatures-equal signatures-match (*alist1*, *alist2*)

 \rightarrow (signature (*alist1*) = signature (*alist2*))

EVENT: Disable signatures-match-implies-signatures-equal.

Theorem: translation-is-correct 2

(ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-translation-parameters (*cinfo*, *cond-list*, *stmt*, *proc-list*, *code2*)
- \land ok-mg-statep (*mg-state*, cond-list)
- \land (code (translate-def-body (assoc (*subr*, *proc-list*), *proc-list*))
 - = append (code (translate (*cinfo*, *cond-list*, *stmt*, *proc-list*)), *code2*))
- \land user-defined-procp (*subr*, *proc-list*)
- \land plistp(*temp-stk*)
- $\wedge \quad \text{listp}(ctrl-stk)$
- \land mg-vars-list-ok-in-p-state (mg-alist (*mg-state*),

bindings (top (ctrl-stk)),

temp-stk)

 \land no-p-aliasing (bindings (top (*ctrl-stk*)), mg-alist (*mg-state*))

- \land signatures-match (mg-alist (*mg-state*), *name-alist*)
- \land all-cars-unique (mg-alist (*mg-state*))
- \land (\neg resource-errorp (mg-meaning-r (*stmt*,

proc-list,

```
mg-state,
```

```
n, list (length (temp-stk),
```

```
p-ctrl-stk-size(ctrl-stk)))))
```

- $\land \quad (pc\text{-}offset = \text{length}(\text{code}(cinfo)))$
- $\land \quad (\operatorname{cc}(mg\text{-state}) \neq \texttt{'leave}))$
- \rightarrow (map-up (p (map-down (*mg-state*,

proc-list,

```
ctrl-stk,
```

temp-stk, tag('pc, cons(subr, pc-offset)),

```
cond-list),
```

 $\operatorname{clock}(stmt, proc-list, mg-state, n)),$

signature (mg-alist (*mg-state*)),

cond-list)

= mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))

;; The following few functions are for making the "initial story" about how you invoke the ;; MG compiler at the highest level. That is, we want to be able to compute the meaning of

;; an MG statement. To do so, I build a procedure around it and show by the lemma

```
;; translation-is-correct2 that the meaning of the statement is the execution of that
;; procedure.
;;
;; This takes an alist from the mg-state and turns it into
;; a list of local-var-decls as might appear in a Micro-Gypsy
;; procedure.
;; Does this do anything; the alist is already in the right form.
DEFINITION:
make-mg-locals-list (mg-alist)
= if mg-alist \simeq nil then nil
    else cons (list (name (car (mg-alist))),
                 m-type (car (mg-alist)),
                  m-value (car (mq-alist))),
              make-mg-locals-list (cdr (mg-alist))) endif
THEOREM: make-mg-local-list-preserves-listcars
listcars (make-mg-locals-list (lst)) = listcars (lst)
```

```
THEOREM: make-mg-locals-list-ok-mg-local-data-plistp
mg-alistp(mg-alist)
→ ok-mg-local-data-plistp(make-mg-locals-list(mg-alist))
;; The point of this is to cons up a new user-defined-procedure
;; definition from the following components:
;; alist: the Micro-Gypsy variable alist in which the statement is
;; to be interpreted.
;; subr: a user-supplied name;
;; stmt: the statement which we are interpreting
```

```
;; cond-list: a list of conditions which we will allow
```

```
;; to be raised.
```

DEFINITION:

```
make-mg-proc(alist, subr, stmt, cond-list)
= list(subr, nil, cond-list, make-mg-locals-list(alist), nil, stmt)
;; Initial-temp-stk places the values of the mg-alist variables onto the temp-stk;
;; initial-bindings creates the Piton bindings corresponding to that initial
;; temp-stk.
```

DEFINITION:

initial-temp-stk-reversed (*mg-alist*)

```
= if mg-alist \simeq nil then nil 
elseif simple-mg-type-refp (cadr (car (mg-alist))) 
then cons (mg-to-p-simple-literal (caddr (car (mg-alist))), 
initial-temp-stk-reversed (cdr (mg-alist))) 
else append (mg-to-p-simple-literal-list (caddr (car (mg-alist))), 
initial-temp-stk-reversed (cdr (mg-alist))) endif
```

THEOREM: initial-temp-stk-reversed-plistp plistp (initial-temp-stk-reversed (x))

```
DEFINITION:
```

initial-temp-stk(mg-alist) = reverse(initial-temp-stk-reversed(mg-alist))

```
DEFINITION:
```

```
 \begin{array}{ll} \operatorname{initial-bindings}\left(mg\text{-}alist, n\right) \\ = & \operatorname{if} mg\text{-}alist \simeq \operatorname{nil} \operatorname{then} \operatorname{nil} \\ & \operatorname{elseif} \operatorname{simple-mg-type-refp}\left(\operatorname{cadr}\left(\operatorname{car}\left(mg\text{-}alist\right)\right)\right) \\ & \operatorname{then} \operatorname{cons}\left(\operatorname{cons}\left(\operatorname{caar}\left(mg\text{-}alist\right), \operatorname{tag}\left(\texttt{'nat}, n\right)\right), \\ & \operatorname{initial-bindings}\left(\operatorname{cdr}\left(mg\text{-}alist\right), 1+n\right)\right) \\ & \operatorname{else} \operatorname{cons}\left(\operatorname{cons}\left(\operatorname{caar}\left(mg\text{-}alist\right), \operatorname{tag}\left(\texttt{'nat}, n\right)\right), \\ & \operatorname{initial-bindings}\left(\operatorname{cdr}\left(mg\text{-}alist\right), \\ & n + \operatorname{array-length}\left(\operatorname{cadr}\left(\operatorname{car}\left(mg\text{-}alist\right)\right)\right))) \text{ endif} \end{array} \right)
```

THEOREM: length-initial-bindings length (initial-bindings (alist, n)) = length (alist)

```
;; 10/4/88 changing the final return pc from nil to (pc (subr . 0))
;; It is ignored but must be a legal pc value.
```

DEFINITION:

```
;; the following and guarantee the following:
```

```
;; proc-list: must be a legitimate MG procedure list;
```

```
a list of conditions you will allow to be raised;
      cond-list:
;;
                   a legitimate MG state with current condition constrained by cond-list,
      mg-state:
;;
                      names on the alist must all be unique;
;;
                   must be a legal statement with respect to proc-list and cond-list;
;;
      stmt:
     n:
                   the clock parameter saying how long to let the thing run;
;;
          an integer telling me the size of the implementation ;
      :
;;
              a litatom which is not the name of any procedure on proc-list;
      subr:
;;
;;
;; Supplied with these things, I can show you the meaning of a statement (provided that
;; the computation does not run out of time or space.
```

```
;;(defn ok-cond-list1 (cond-list)
;; (and (ok-cond-list cond-list)
;; (not (member 'leave cond-list))))
```

DEFINITION:

new-proc-name (x, proc-list)= $(ok-mg-namep(x) \land (\neg defined-procp(x, proc-list)))$

THEOREM: new-proc-doesnt-affect-fetch-called-def

((car(stmt) = 'proc-call-mg))

 \land new-proc-name (car (*new-proc*), *proc-list*)

 $\wedge \quad \text{ok-mg-statement} \left(\textit{stmt}, \textit{ cond-list}, \textit{ name-alist}, \textit{ proc-list} \right) \right)$

- $\rightarrow \quad (\text{fetch-called-def}\left(\textit{stmt}, \ \text{cons}\left(\textit{new-proc}, \ \textit{proc-list}\right)\right)$
 - = fetch-called-def(*stmt*, *proc-list*))

THEOREM: new-proc-doesnt-affect-mg-meaning-proc-call-case $((n \not\simeq 0)$

 $\wedge \quad \text{normal}(mg\text{-}state)$

 $\wedge \quad (\operatorname{car}(stmt) = \texttt{'proc-call-mg})$

 \land ((ok-mg-def-plistp(*proc-list*))

 $\land \quad \text{ok-mg-statement} (\text{def-body} (\text{fetch-called-def} (stmt, proc-list)), \\ \\ \qquad \qquad \text{make-cond-list} (\text{fetch-called-def} (stmt, stmt)), \\ \end{aligned}$

proc-list)),

make-name-alist (fetch-called-def (stmt,

proc-list)),

proc-list)

 \land new-proc-name (car (*new-proc*), *proc-list*))

 \rightarrow (mg-meaning (def-body (fetch-called-def (*stmt*, *proc-list*)),

proc-list,

make-call-environment (mg-state,

stmt, fetch-called-def(*stmt*,

proc-list)),

n - 1)

 $= \quad \text{mg-meaning} \, (\text{def-body} \, (\text{fetch-called-def} \, (\textit{stmt}, \textit{proc-list})),$

 $\cos(new-proc, proc-list),$

 ${\it make-call-environment}\,({\it mg-state},$

stmt,

fetch-called-def (stmt,

proc-list)),

n - 1)))

 \land ok-mg-def-plistp (*proc-list*)

 \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)

 \land new-proc-name (car (*new-proc*), *proc-list*))

- \rightarrow (mg-meaning (stmt, proc-list, mg-state, n)
 - = mg-meaning (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*))

DEFINITION:

meaning-induction-hint3 (stmt,

proc-list, mg-state, n, name-alist, cond-list, new-proc)

 $= if n \simeq 0 then t$ elseif \neg normal(*mg-state*) then t

Λ

elseif car(stmt) ='no-op-mg then t

 $\mathbf{elseif} \, \mathrm{car} \, (stmt) =$ 'signal-mg then t

```
\mathbf{elseif} \, \mathrm{car} \, (stmt) = \texttt{'prog2-mg}
```

then meaning-induction-hint3 (prog2-left-branch (stmt),

proc-list, mg-state,

n - 1,

name-alist,

cond-list, new-proc)

meaning-induction-hint3 (prog2-right-branch (stmt),

proc-list,

mg-meaning (prog2-left-branch (stmt),

```
cons(new-proc, proc-list),
```

```
mg-state,
n-1),
n = 1,
name-alist,
```

cond-list,

new-proc)

```
elseif car(stmt) = \text{'loop-mg}
then meaning-induction-hint3 (loop-body (stmt),
                                proc-list,
                                mg-state,
                                n - 1,
                                name-alist,
                                cons('leave, cond-list),
                                new-proc)
      \wedge
         meaning-induction-hint 3(stmt,
                                     proc-list,
                                    mg-meaning (loop-body (stmt)),
                                                  cons (new-proc,
                                                        proc-list),
                                                   mg-state,
                                                   n - 1),
                                     n - 1,
                                     name-alist,
                                     cond-list,
                                     new-proc)
elseif car(stmt) = 'if-mg
then meaning-induction-hint3 (if-false-branch (stmt),
                                proc-list,
                                mg-state,
                                n - 1,
                                name-alist,
                                 cond-list,
                                new-proc)
         meaning-induction-hint3 (if-true-branch (stmt)),
      \wedge
                                    proc-list,
                                     mg-state,
                                     n - 1,
                                     name-alist,
                                     cond-list,
                                     new-proc)
elseif car(stmt) = 'begin-mg
then meaning-induction-hint3 (begin-body (stmt),
                                proc-list,
                                mg-state,
                                n - 1,
                                name-alist,
                                append (when-labels (stmt), cond-list),
                                new-proc)
         meaning-induction-hint3 (when-handler (stmt)),
      \wedge
```

proc-list, set-condition (mg-meaning (begin-body (stmt), cons (new-proc, proc-list), mg-state, n - 1),'normal), n - 1, name-alist, cond-list, new-proc) elseif car(stmt) = 'proc-call-mgthen meaning-induction-hint3 (def-body (fetch-called-def (*stmt*, *proc-list*)), proc-list, make-call-environment (mg-state, stmt, fetch-called-def (stmt,proc-list)),n - 1, make-name-alist (fetch-called-def(stmt, proc-list)), make-cond-list (fetch-called-def (stmt, proc-list)), *new-proc*) elseif car(stmt) ='predefined-proc-call-mg then t else f endif ;; >>> For an automatic proof, these rules should be oriented in the other direction. they could cause looping. ;;

Othe

THEOREM: new-proc-doesnt-affect-mg-meaning

(ok-mg-def-plistp(proc-list)

- $\wedge \quad \text{ok-mg-statement} \left(\textit{stmt}, \textit{ cond-list}, \textit{ name-alist}, \textit{ proc-list} \right)$
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)
 - = mg-meaning (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*))

DEFINITION:

meaning-induction-hint4 (stmt,

proc-list,mg-state,n,name-alist,cond-list,

new-proc, sizes) = if $n \simeq 0$ then t elseif \neg normal(*mg-state*) then t elseif resources-inadequatep (stmt, proc-list, sizes) then t elseif car(stmt) = 'no-op-mg then t elseif car(stmt) ='signal-mg then t **elseif** car(stmt) = 'prog2-mgthen meaning-induction-hint4 (prog2-left-branch (*stmt*), proc-list, mg-state, n - 1, name-alist, cond-list, new-proc, sizes) meaning-induction-hint4 (prog2-right-branch (*stmt*), \wedge proc-list, mg-meaning-r (prog2-left-branch (stmt), cons (new-proc, proc-list), mg-state, n - 1, sizes), n - 1, name-alist,cond-list, new-proc, sizes) elseif car(stmt) = 'loop-mgthen meaning-induction-hint4 (loop-body (stmt), proc-list, mg-state, n - 1, name-alist, cons('leave, cond-list), new-proc, sizes) \wedge meaning-induction-hint4 (*stmt*, proc-list, mg-meaning-r (loop-body (stmt), cons (new-proc, proc-list),

mg-state,

n - 1, sizes), n - 1, name-alist, cond-list, new-proc, sizes) elseif car(stmt) = 'if-mgthen meaning-induction-hint4 (if-false-branch (*stmt*), proc-list, mg-state, n - 1, name-alist, cond-list, new-proc, sizes) meaning-induction-hint4 (if-true-branch (*stmt*), \wedge proc-list, mg-state, n - 1, name-alist, cond-list, new-proc, sizes) **elseif** car(stmt) = 'begin-mgthen meaning-induction-hint4 (begin-body (*stmt*), proc-list, mg-state, n - 1, name-alist, append (when-labels (*stmt*), *cond-list*), new-proc, sizes) meaning-induction-hint4 (when-handler (stmt)), \wedge proc-list, set-condition (mg-meaning-r (begin-body (stmt), cons (new-proc, proc-list), mg-state, n - 1, sizes), 'normal), n - 1, name-alist,

```
cond-list,
                                     new-proc,
                                     sizes)
elseif car(stmt) = 'proc-call-mg
then meaning-induction-hint4 (def-body (fetch-called-def (stmt, proc-list)),
                                proc-list,
                                make-call-environment (mg-state,
                                                         stmt,
                                                         fetch-called-def(stmt,
                                                                          proc-list)),
                                 n - 1,
                                 make-name-alist (fetch-called-def(stmt,
                                                                    proc-list)),
                                 make-cond-list (fetch-called-def(stmt,
                                                                  proc-list)),
                                 new-proc,
                                list (t-size (sizes)
                                     + data-length (def-locals (fetch-called-def (stmt,
                                                                                   proc-list))),
                                     c-size (sizes)
                                     + (2
                                              length (def-locals (fetch-called-def (stmt,
                                          +
                                                                                   proc-list)))
                                              length (def-formals (fetch-called-def (stmt,
                                          +
                                                                                    proc-list)))))))
elseif car(stmt) = 'predefined-proc-call-mg then t
else f endif
```

THEOREM: new-proc-doesnt-affect-resources-inadequatep (ok-mg-def-plistp (proc-list)

- \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (resources-inadequatep (*stmt*, cons (*new-proc*, *proc-list*), *sizes*)
 - = resources-inadequatep (*stmt*, *proc-list*, *sizes*))

THEOREM: new-proc-doesnt-affect-mg-meaning-r (ok-mg-def-plistp (*proc-list*)

- \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)
 - = mg-meaning-r (stmt, cons (new-proc, proc-list), mg-state, n, sizes))

THEOREM: new-proc-doesnt-affect-mg-meaning-r-2

(ok-mg-def-plistp (*proc-list*)

 \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)

- new-proc-name (car (*new-proc*), *proc-list*)) Λ
- (mg-meaning-r (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*, *sizes*) = mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*))

THEOREM: new-proc-doesnt-affect-clock-prog2-case

 $((n \not\simeq 0)$

 $\wedge \quad \text{normal}(mq\text{-state})$ (car(stmt) = 'prog2-mg) \wedge ((ok-mg-def-plistp(*proc-list*)) \wedge \wedge ok-mg-statement (prog2-right-branch (*stmt*), cond-list, name-alist, proc-list) Λ

- new-proc-name (car (*new-proc*), *proc-list*))
- $(\operatorname{clock}(\operatorname{prog2-right-branch}(stmt)),$

proc-list,

mg-meaning (prog2-left-branch (stmt), cons (new-proc, proc-list),

mg-state,

n - 1),

- n 1)
- clock (prog2-right-branch (*stmt*), =
 - cons (new-proc, proc-list),

```
mg-meaning (prog2-left-branch (stmt),
```

```
\cos(new-proc, proc-list),
```

```
mg-state,
```

```
n - 1),
```

```
(n - 1)))
```

- ((ok-mg-def-plistp(*proc-list*) \wedge
 - \wedge ok-mg-statement (prog2-left-branch (*stmt*),
 - cond-list,
 - name-alist.
 - proc-list)
 - new-proc-name (car (new-proc), proc-list)) Λ

$$\rightarrow$$
 (clock (prog2-left-branch (*stmt*), *proc-list*, *mg-state*, $n-1$)

- clock (prog2-left-branch (stmt)),=
 - $\cos(new-proc, proc-list),$

```
mg-state,
```

- (n 1)))
- \land ok-mg-def-plistp (*proc-list*)
- ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*) \wedge
- \land new-proc-name (car (*new-proc*), *proc-list*))
- $(\operatorname{clock}(stmt, proc-list, mq-state, n))$ \rightarrow
 - = $\operatorname{clock}(stmt, \operatorname{cons}(new-proc, proc-list), mg-state, n))$

THEOREM: new-proc-doesnt-affect-clock-loop-case $((n \not\simeq 0)$ \wedge normal (*mq-state*) $\land \quad (\operatorname{car}(stmt) = \text{'loop-mg})$ \wedge ((ok-mg-def-plistp(*proc-list*)) \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*) new-proc-name (car (*new-proc*), *proc-list*)) \wedge $(\operatorname{clock}(stmt,$ \rightarrow proc-list, mg-meaning (loop-body (stmt)), $\cos(new-proc, proc-list),$ mg-state, n - 1),n - 1)= clock (*stmt*, $\cos(new-proc, proc-list),$ mg-meaning (loop-body (stmt)), cons (new-proc, proc-list), mg-state, n - 1), (n - 1)))((ok-mg-def-plistp(*proc-list*)) Λ \wedge ok-mg-statement (loop-body (*stmt*), cons('leave, cond-list), name-alist, proc-list) new-proc-name (car (*new-proc*), *proc-list*)) Λ $(\operatorname{clock}(\operatorname{loop-body}(stmt), proc-list, mg-state, n-1))$ \rightarrow = clock (loop-body (*stmt*), cons (new-proc, proc-list), mg-state, (n-1))) \wedge ok-mg-def-plistp (*proc-list*) \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*) \land new-proc-name (car (*new-proc*), *proc-list*)) $(\operatorname{clock}(stmt, proc-list, mg-state, n))$ \rightarrow $\operatorname{clock}(stmt, \operatorname{cons}(new-proc, proc-list), mg-state, n))$ = THEOREM: new-proc-doesnt-affect-clock-if-case $((n \not\simeq 0)$ \wedge normal (*mg-state*) \wedge (car(*stmt*) = 'if-mg)

- \wedge ((ok-mg-def-plistp (*proc-list*))
 - \wedge ok-mg-statement (if-true-branch (*stmt*),

```
cond-list,
```

```
name-alist,
```

proc-list)

- \land new-proc-name (car (*new-proc*), *proc-list*))
- $\rightarrow \quad ({\rm clock}\,({\rm if-true-branch}\,(stmt),\,proc\text{-}list,\,mg\text{-}state,\,n\,-\,1)$
 - = clock (if-true-branch (*stmt*),
 - $\cos(new-proc, proc-list),$
 - mg-state,
 - (n 1)))
- \land ((ok-mg-def-plistp (*proc-list*))
 - \land ok-mg-statement (if-false-branch (*stmt*),
 - cond-list,

name-alist,

proc-list)

- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (clock (if-false-branch (*stmt*), *proc-list*, *mg-state*, n 1)
 - = clock (if-false-branch (*stmt*),

cons (new-proc, proc-list), mg-state,

- n 1)))
- \land ok-mg-def-plistp (*proc-list*)
- \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (clock (*stmt*, *proc-list*, *mg-state*, *n*)
 - = clock (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*))

 $T{\tt Heorem: new-proc-doesnt-affect-clock-begin-case}$

 $((n \not\simeq 0)$

- $\land \quad \text{normal}(mg\text{-state})$
- $\land \quad (\operatorname{car}(stmt) = \texttt{'begin-mg})$
- \land ((ok-mg-def-plistp(*proc-list*))
 - \wedge ok-mg-statement (when-handler (*stmt*),

cond-list,

name-alist,

```
proc-list)
```

- $\wedge \quad \text{new-proc-name}\left(\text{car}\left(\textit{new-proc} \right), \textit{proc-list} \right) \right)$
- \rightarrow (clock (when-handler (*stmt*),

proc-list,

set-condition (mg-meaning (begin-body (*stmt*),

cons (new-proc, proc-list), mg-state,

$$(n - 1)$$

'normal),

n - 1)

 $\operatorname{clock}(\operatorname{when-handler}(stmt)),$ = cons (new-proc, proc-list), set-condition (mg-meaning (begin-body (stmt), cons (new-proc, proc-list), mg-state, n - 1), 'normal), (n - 1)))Λ ((ok-mg-def-plistp(*proc-list*)) Λ ok-mg-statement (begin-body (stmt)), append (when-labels (*stmt*), *cond-list*), name-alist, proc-list) new-proc-name (car (*new-proc*), *proc-list*)) \wedge $(\operatorname{clock}(\operatorname{begin-body}(stmt), proc-list, mg-state, n-1))$ \rightarrow $\operatorname{clock}(\operatorname{begin-body}(stmt)),$ = cons (new-proc, proc-list), mg-state, (n-1))) \wedge ok-mg-def-plistp (*proc-list*) ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*) Λ new-proc-name (car (*new-proc*), *proc-list*)) \wedge $(\operatorname{clock}(stmt, proc-list, mg-state, n))$ \rightarrow $\operatorname{clock}(stmt, \operatorname{cons}(new-proc, proc-list), mg-state, n))$ = THEOREM: new-proc-doesnt-affect-clock-proc-call-case $((n \not\simeq 0)$ \wedge normal (*mg-state*) $\wedge \quad (\operatorname{car}(stmt) = \mathsf{'proc-call-mg})$ Λ ((ok-mg-def-plistp(*proc-list*)) \land ok-mg-statement (def-body (fetch-called-def (*stmt*, *proc-list*)), make-cond-list (fetch-called-def (stmt,

proc-list)),

make-name-alist (fetch-called-def (*stmt*,

proc-list)),

proc-list)

 \land new-proc-name (car (*new-proc*), *proc-list*))

 \rightarrow (clock (def-body (fetch-called-def (*stmt*, *proc-list*)),

proc-list,

make-call-environment (mg-state,

stmt,

fetch-called-def(*stmt*, *proc-list*)),

n - 1)

 $= \ \ \, {\rm clock} \, ({\rm def}{\rm -body} \, ({\rm fetch-called-def} \, (stmt, \, proc-list)),$

 $\cos(new-proc, proc-list),$

make-call-environment (mg-state,

stmt,

fetch-called-def (stmt,

proc-list)),

(n - 1)))

 \land ok-mg-def-plistp (*proc-list*)

 \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)

 \land new-proc-name (car (*new-proc*), *proc-list*))

 \rightarrow (clock (*stmt*, *proc-list*, *mg-state*, *n*)

= clock (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*))

 $T{\tt HEOREM:}\ new-proc-doesnt-affect-clock$

(ok-mg-def-plistp (proc-list)

- \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow (clock (*stmt*, *proc-list*, *mg-state*, *n*)
 - = clock (*stmt*, cons (*new-proc*, *proc-list*), *mg-state*, *n*))

THEOREM: new-proc-doesnt-affect-ok-mg-statement-proc-call-case ((car(stmt) = 'proc-call-mg)

- \wedge ok-mg-def-plistp (*proc-list*)
- \land ok-mg-statement (*stmt*, *cond-list*, *name-alist*, *proc-list*)
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow ok-mg-statement (*stmt*, *cond-list*, *name-alist*, cons (*new-proc*, *proc-list*))

EVENT: Enable ok-mg-statement.

THEOREM: new-proc-doesnt-affect-ok-mg-statement

(ok-mg-def-plistp(proc-list))

- $\wedge \quad \text{ok-mg-statement} \left(\textit{stmt}, \textit{ cond-list}, \textit{ name-alist}, \textit{ proc-list} \right)$
- \land new-proc-name (car (*new-proc*), *proc-list*))
- \rightarrow ok-mg-statement (*stmt*, *cond-list*, *name-alist*, cons (*new-proc*, *proc-list*))

THEOREM: new-proc-preserves-ok-mg-def

- $(\text{new-proc-name}\left(\operatorname{car}\left(\mathit{new-proc}\right), \, \mathit{pl}\right)$
- \land ok-mg-def-plistp (pl)
- \land ok-mg-def (*def*, *pl*))
- \rightarrow ok-mg-def (*def*, cons (*new-proc*, *pl*))

THEOREM: new-proc-preserves-ok-mg-def-plistp1 (new-proc-name (car (*new-proc*), *pl2*)

- (new-proc-name (car (*new-proc*), p
- $\wedge \quad \text{ok-mg-def-plistp} \left(\textit{pl2} \right)$
- $\wedge \quad \text{ok-mg-def-plistp1} \left(\textit{pl1}, \textit{pl2} \right) \right)$
- \rightarrow ok-mg-def-plistp1 (*pl1*, cons (*new-proc*, *pl2*))

THEOREM: make-alist-make-locals-list-preserves-signatures-match mg-alistp (mg-alist)

 \rightarrow signatures-match (*mg-alist*,

make-alist-from-formals (make-mg-locals-list (*mg-alist*)))

;; To do this one I needed to change the hyp the cond-list is a cond-identifierp-plistp ;; to an identifier plistp.

THEOREM: new-proc-preserves-ok-mg-def-plistp

(ok-mg-statement (*stmt*, *cond-list*, mg-alist (*mg-state*), *proc-list*)

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-mg-statep (*mg-state*, *cond-list*)
- \wedge identifier-plistp(cond-list)
- \land all-cars-unique (mg-alist (*mg-state*))
- \land new-proc-name (*subr*, *proc-list*)
- $\land \quad (\text{length}(cond-list) < (((\exp(2, \text{MG-WORD-SIZE}) 1) 1))))$
- \rightarrow ok-mg-def-plistp (cons (make-mg-proc (mg-alist (*mg-state*),

 $subr, \\ stmt, \\ cond-list),$

proc-list))

THEOREM: mg-to-p-simple-literal-list-listp listp (mg-to-p-simple-literal-list (x)) = listp (x)

THEOREM: initial-temp-stk-reversed-listp (mg-alistp $(x) \land$ listp $(x)) \rightarrow$ listp (initial-temp-stk-reversed (x))

THEOREM: length-initial-temp-stk-reversed mg-alistp (*alist*) \rightarrow (length (initial-temp-stk-reversed (*alist*)) = data-length (*alist*))

DEFINITION:

initial-bindings-induction-hint (mg-alist, n, lst)

= if *mg-alist* \simeq nil then t

elseif simple-mg-type-refp (cadar(mg-alist))

then initial-bindings-induction-hint (cdr (*mg-alist*),

$$1 + n$$
,

 $\cos(\text{mg-to-p-simple-literal}(\text{caddar}(mg-alist))),$

lst))

else initial-bindings-induction-hint (cdr (mg-alist),

n + array-length (cadar (mg-alist)),append (reverse (mg-to-p-simple-literal-list (caddar (mg-alist))), lst)) endif THEOREM: initial-bindings-ok-in-initial-temp-stk1 $(\text{mg-alistp}(mg-alist) \land (n = \text{length}(lst)) \land \text{all-cars-unique}(mg-alist))$ \rightarrow mg-vars-list-ok-in-p-state (*mg-alist*, initial-bindings (mq-alist, n), append (initial-temp-stk (*mq-alist*), *lst*)) THEOREM: initial-bindings-ok-in-initial-temp-stk $(ok-mg-state, cond-list) \land all-cars-unique (mg-alist (mg-state)))$ \rightarrow mg-vars-list-ok-in-p-state (mg-alist (*mg-state*), initial-bindings (mg-alist (*mg-state*), 0), initial-temp-stk (mg-alist (mg-state))) THEOREM: initial-bindings-all-pointers-bigger all-cars-unique (*alist*) \rightarrow all-pointers-bigger (collect-pointers (initial-bindings (*alist*, *n*), *alist*), n) THEOREM: no-p-aliasing-in-initial-bindings (all-cars-unique $(mq-alist) \land mg-alistp (mq-alist) \land (n \in \mathbf{N})$) \rightarrow no-p-aliasing (initial-bindings (*mg-alist*, *n*), *mg-alist*) THEOREM: leave-not-state-cc $(ok-mg-statep(mg-state, cond-list) \land identifier-plistp(cond-list))$ \rightarrow (cc(*mg-state*) \neq 'leave) **THEOREM:** translation-is-correct3 (ok-mg-statement (*stmt*, *cond-list*, mg-alist (*mg-state*), *proc-list*) \wedge ok-mg-def-plistp (*proc-list*)

- \land ok-mg-statep (*mg-state*, *cond-list*)
- \land identifier-plistp (*cond-list*)
- \wedge all-cars-unique (mg-alist (*mg-state*))
- \wedge new-proc-name (*subr*, *proc-list*)
- $\land \quad (\text{length}(cond-list) < (((\exp(2, \text{MG-WORD-SIZE}) 1) 1)))$
- \land (\neg resource-errorp (mg-meaning-r (*stmt*,

proc-list,

mg-state, n,

list (length (initial-temp-stk (mg-alist (mg-state))),

p-ctrl-stk-size (list (cons (initial-bindings (mg-alist (mg-state),

0),

list (tag('pc,

 $\cos\left(subr,\right)$

 $\rightarrow (\text{map-up}(p(\text{map-down1}(\textit{mg-state},\textit{proc-list},\textit{cond-list},\textit{subr},\textit{stmt}), \\ \text{clock}(\textit{stmt},\textit{proc-list},\textit{mg-state},\textit{n})),$

signature (mg-alist (mg-state)), cond-list)

= mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))

;; This is just a slightly cleaned up version of the previous lemma.

THEOREM: translation-is-correct4

(ok-mg-statement (*stmt*, *cond-list*, mg-alist (*mg-state*), *proc-list*)

- \land ok-mg-def-plistp (*proc-list*)
- \land ok-mg-statep (*mg-state*, cond-list)
- \wedge identifier-plistp (*cond-list*)
- \land all-cars-unique (mg-alist (*mg-state*))
- \land new-proc-name (*subr*, *proc-list*)
- $\land \quad (\text{length}(\textit{cond-list}) < (((\exp(2, \text{MG-WORD-SIZE}) 1) 1))$
- \land (\neg resource-errorp (mg-meaning-r (*stmt*,

proc-list,

$$mg$$
-state,

n,

list (data-length (mg-alist (mg-state)),

2 + length(mg-alist(mg-state)))))))

 \rightarrow (map-up (p (map-down1 (*mg-state*, *proc-list*, *cond-list*, *subr*, *stmt*),

 $\operatorname{clock}(stmt, proc-list, mg-state, n)),$

signature (mg-alist (mg-state)),

cond-list)

= mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))

DEFINITION:

ok-execution-environment (stmt, cond-list, proc-list, mg-state, subr, n)

- = (ok-mg-statement (stmt, cond-list, mg-alist (mg-state), proc-list)
 - $\wedge \quad \text{ok-mg-def-plistp} \left(\textit{proc-list} \right)$
 - \land ok-mg-statep (*mg-state*, cond-list)
 - \land identifier-plistp (*cond-list*)
 - \land all-cars-unique (mg-alist (*mg-state*))
 - \land new-proc-name (*subr*, *proc-list*)
 - $\land \quad (\text{length}(cond-list) < (((\exp(2, \text{MG-WORD-SIZE}) 1) 1))))$

THEOREM: translation-is-correct5

(ok-execution-environment(stmt, cond-list, proc-list, mg-state, subr, n)

 \land (\neg resource-errorp (mg-meaning-r (*stmt*,

 $\begin{array}{l} proc-list,\\ mg\text{-}state,\\ n,\\ list (data-length (mg\text{-}alist (mg\text{-}state)),\\ \mathbf{2}+ \text{length (mg\text{-}alist (mg\text{-}state))))))))\end{array}$

 $\begin{array}{ll} \rightarrow & (\text{map-up} \left(\text{p} \left(\text{map-down1} \left(\textit{mg-state}, \textit{proc-list}, \textit{cond-list}, \textit{subr}, \textit{stmt} \right), \\ & \quad \text{clock} \left(\textit{stmt}, \textit{proc-list}, \textit{mg-state}, \textit{n} \right) \right), \\ & \quad \text{signature} \left(\text{mg-alist} \left(\textit{mg-state} \right) \right), \\ & \quad \textit{cond-list} \right) \\ & = & \quad \text{mg-meaning} \left(\textit{stmt}, \textit{proc-list}, \textit{mg-state}, \textit{n} \right) \right) \end{array}$

EVENT: Make the library "ca10".

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