

GSS GSKNN

BLIS-Based High Performance Computing Kernels
in N-body Problems

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The 3rd BLIS Retreat
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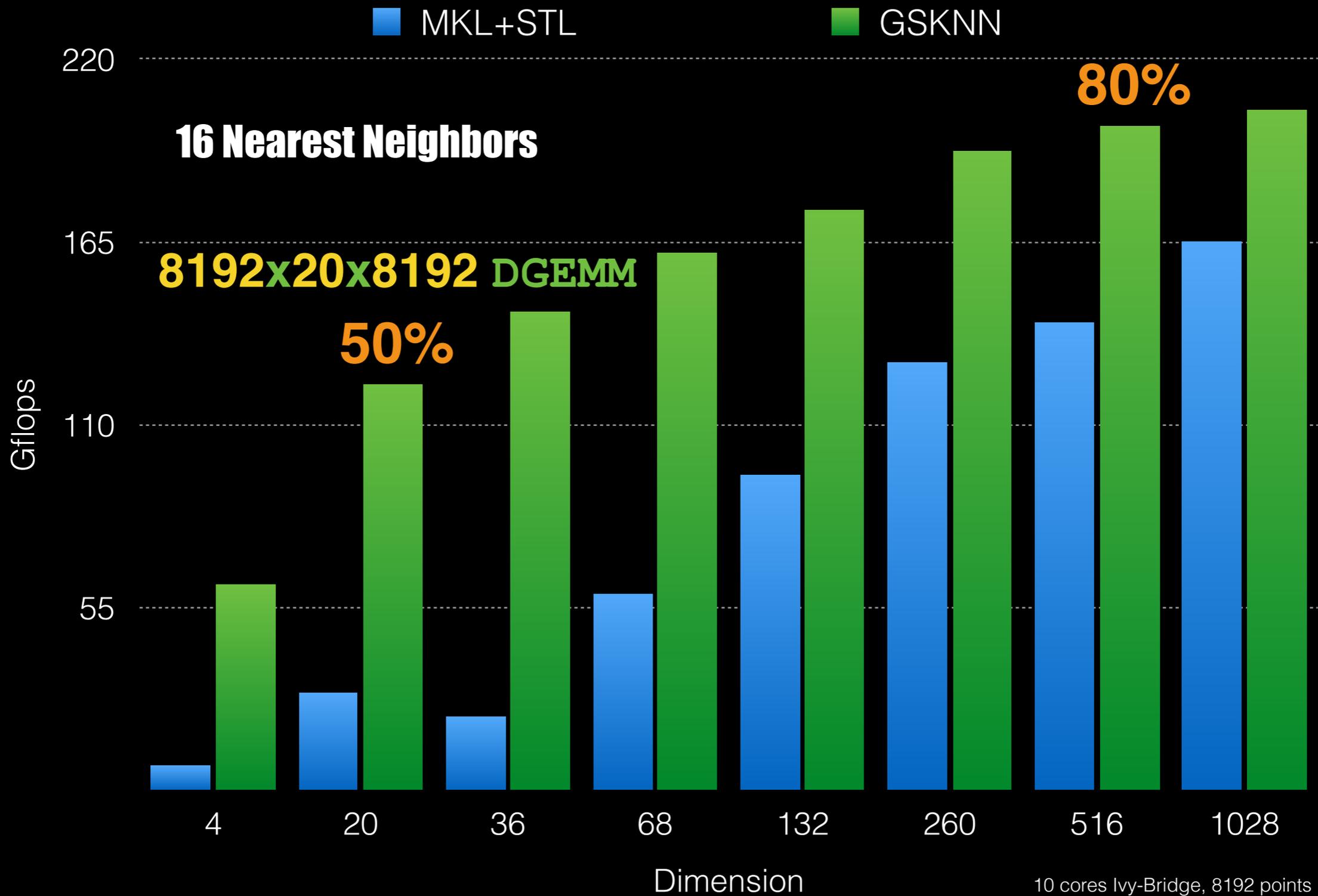


N-body Problems



N-body Problems

- N-body problems aim to describe the interaction (relation) of N points $\{ \mathbf{X} \}$ in a d dimensional space.
- $\mathcal{K}(\mathbf{x}_i, \mathbf{x}_j) = \mathbf{K}_{ij}$ describes the interaction between \mathbf{x}_i and \mathbf{x}_j .
- 3 operations: Kernel Summation $\mathbf{u} = \mathbf{K}\mathbf{w}$, Kernel Inversion $\mathbf{w} = (\mathbf{K} + \lambda \mathbf{I})^{-1} \mathbf{u}$ and Nearest-Neighbors.
- 2D and 3D applications can be found in computational physics, geophysical exploration and medical imaging.
- High dimension applications in computational statistic include clustering, classification and regression.

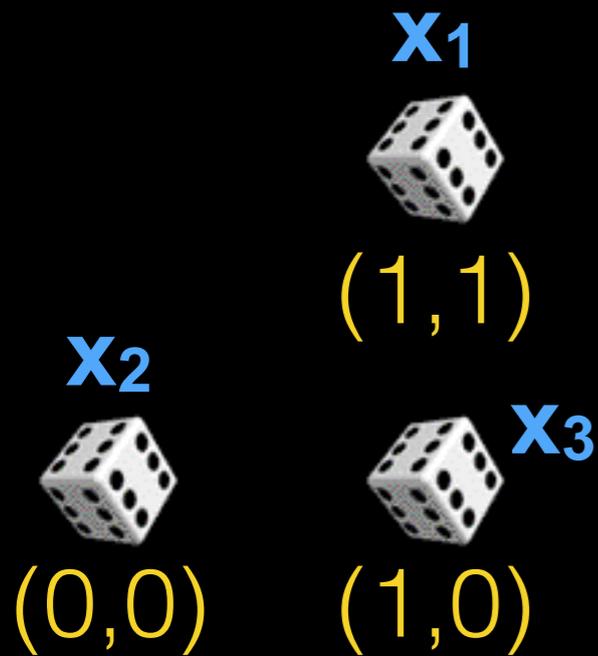


10 cores Ivy-Bridge, 8192 points

Outline

- Kernel Summation ($\mathbf{u}=\mathbf{Kw}$) and Nearest-Neighbors.
- How GEMM is applied in the conventional approach?
- Why GEMM can be memory bound in these operations?
- What insight is required to design an algorithm that avoids redundant memory operations but still preserves the efficiency?
- How GSKS and GSKNN are inspired by the BLIS framework in their design?

Linear Kernel: $\mathcal{K}(\mathbf{x}_i, \mathbf{x}_j) = \mathbf{x}_i^T \mathbf{x}_j$



R

1	0	1
1	0	0

\mathbf{x}_1 \mathbf{x}_2 \mathbf{x}_3

\mathbf{x}_1^T

1	1
0	0
1	0

\mathbf{x}_2^T

\mathbf{x}_3^T

$Q^T = R$

2	0	1
0	0	0
1	0	1

$K = Q^T R$

Kernel Summation **Kw**

2	0	1	\times	=	1	3
0	0	0			1	0
1	0	1			1	2

Nearest-Neighbors

2	0	1
0	0	0
1	0	1

\mathbf{x}_2	\mathbf{x}_3
\mathbf{x}_1	\mathbf{x}_2
\mathbf{x}_2	\mathbf{x}_1

Other Kernels

$\mathcal{K}(\mathbf{x}_i, \mathbf{x}_j) = f(\mathbf{x}_i^T \mathbf{x}_j)$, e.g. Gaussian kernel

$$\mathcal{K}(x_i, x_j) = \exp(-\|x_i - x_j\|_2^2 / (2h^2))$$

The expansion exposes **GEMM** operations:

$$\|x_i - x_j\|_2^2 = \|x_i\|_2^2 + \|x_j\|_2^2 - 2x_i^T x_j$$

GEMM

$\mathbf{x}_1^T \mathbf{x}_1$	2
$\mathbf{x}_2^T \mathbf{x}_2$	0
$\mathbf{x}_3^T \mathbf{x}_3$	1

\mathbf{x}_1^T	1	1
\mathbf{x}_2^T	0	0
\mathbf{x}_3^T	1	0

2	0	1
0	0	0
1	0	1

Q^T

$K = Q^T R$

$1 + 2 - 2 * 1$

The Big Picture

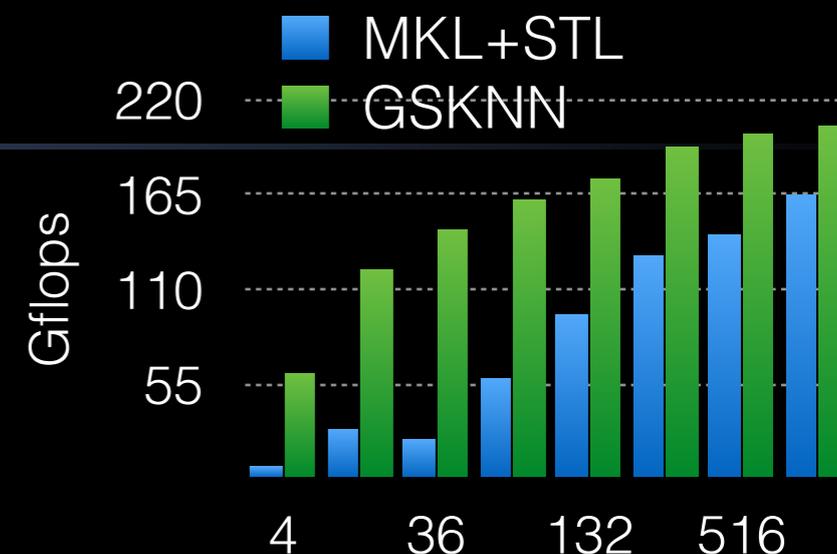
- Kw takes $O(N^2)$ if K is precomputed, otherwise $O(dN^2)$. The cost is too expensive when N is large.
- Exhaustive search requires $O(N^2 \log(k))$ if K is precomputed, otherwise $O(dN^2 + N^2 \log(k))$.
- Divide-and-conquer approximation: Barnes-Hut or FMM for kernel summation, and randomized KD-tree or locality sensitive hashing for kNN.
- Still the subproblem of all these algorithms is to solve several smaller dense kernel summation or kNN.
- Solving the subproblem fast benefits all these methods.

Subproblem

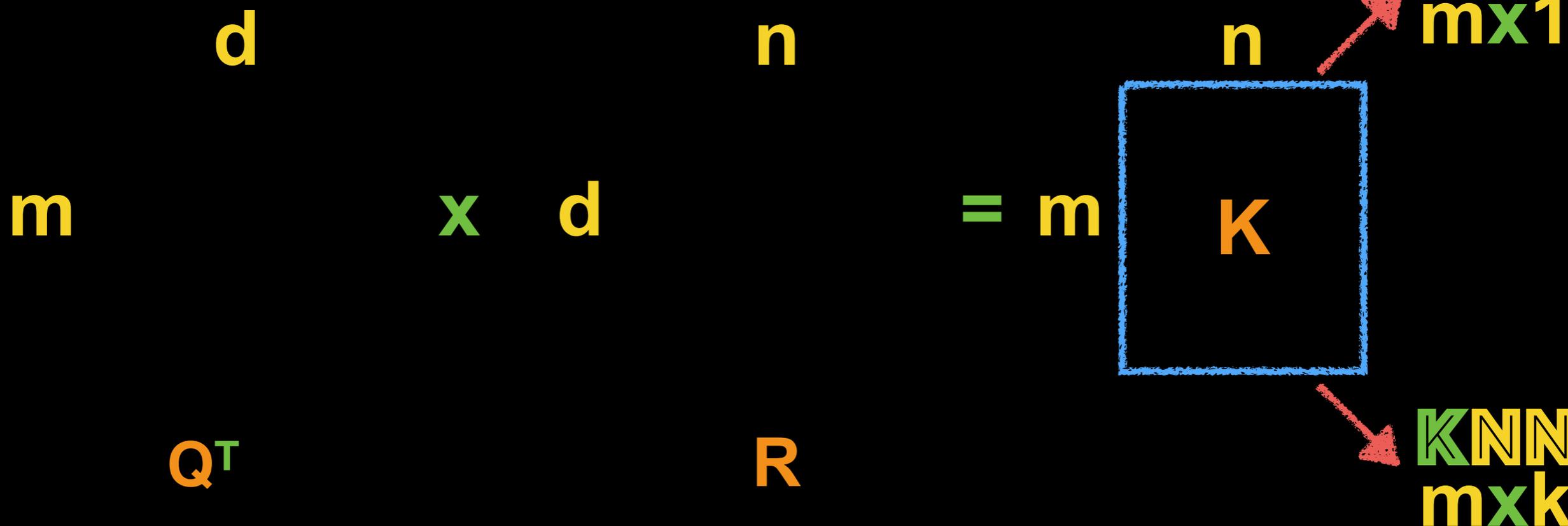
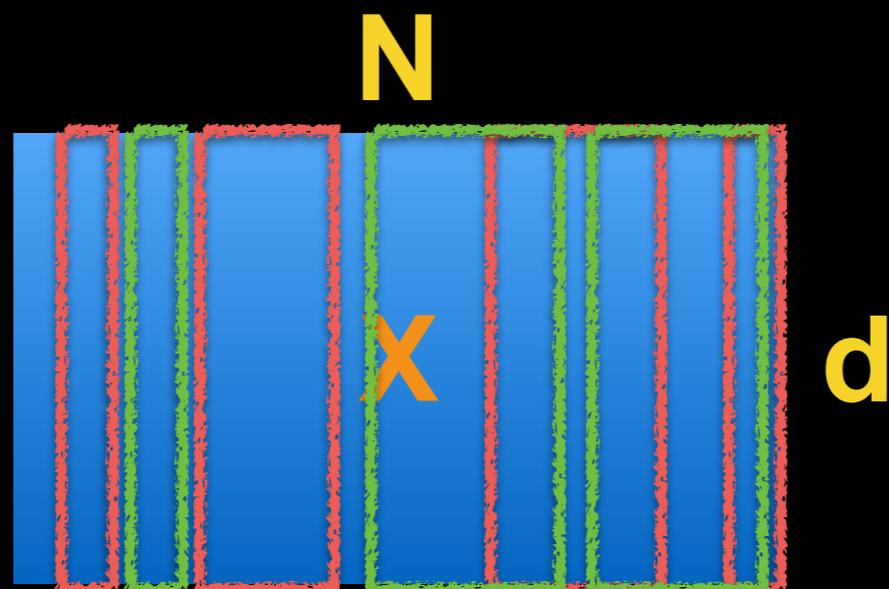
- Take two subsets **Q** and **R** from **X**.
- Compute $\mathcal{K}(Q,R)$ with **GEMM** using:

$$\|x_i - x_j\|_2^2 = \|x_i\|_2^2 + \|x_j\|_2^2 - 2x_i^T x_j$$

- Compute **Kw** with **GEMV** or select **k** entries in each row.
- Rely on BLAS, VML (Vectorized Math Library) and STL.
- What can possibly go wrong?



Visualization



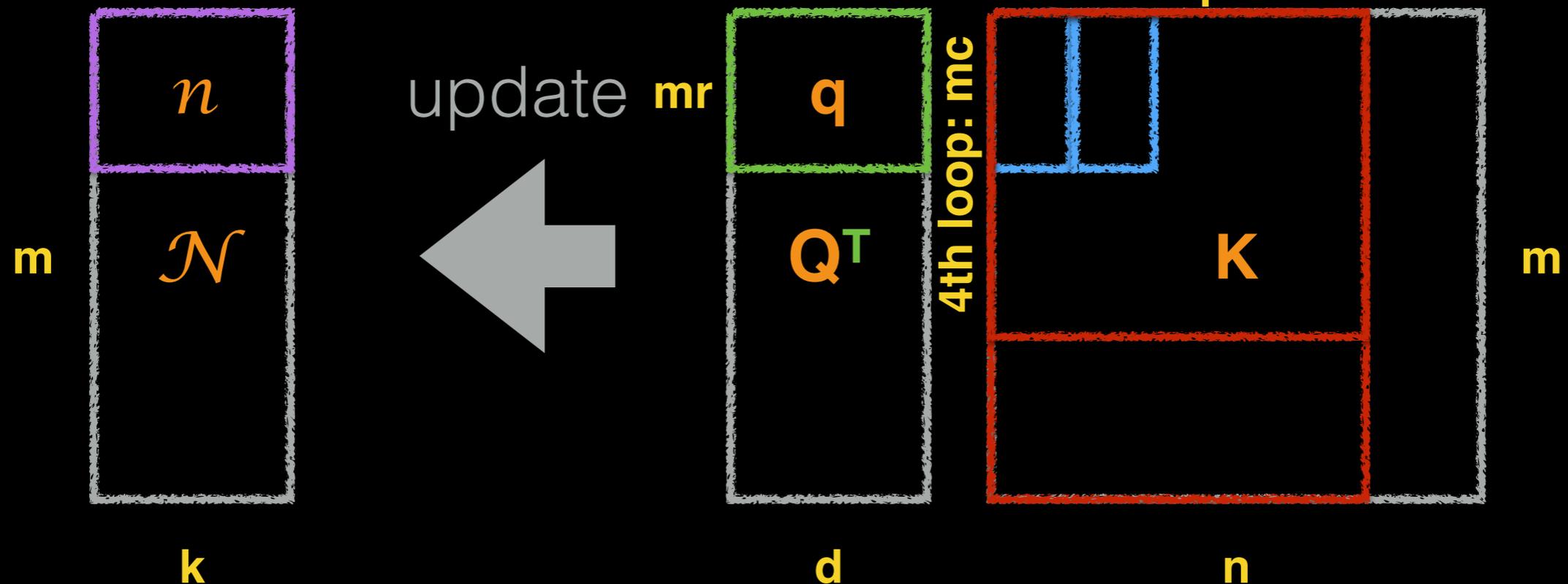
Insights

- **Q**, **R** and **K** can't be stored.
 - Collect **Q** and **R** from **X** during packing.
 - $\mathcal{K}(\mathbf{x}_i, \mathbf{x}_j) = \mathbf{K}_{ij}$ must be computed in registers.
 - **Kw** or **k-select** must be completed in registers.
 - Only store the output. i.e.
-
- We need a special packing routine.
 - Fuse **GEMM** with distance calculations, special function evaluations, **Kw** or **k-select**.

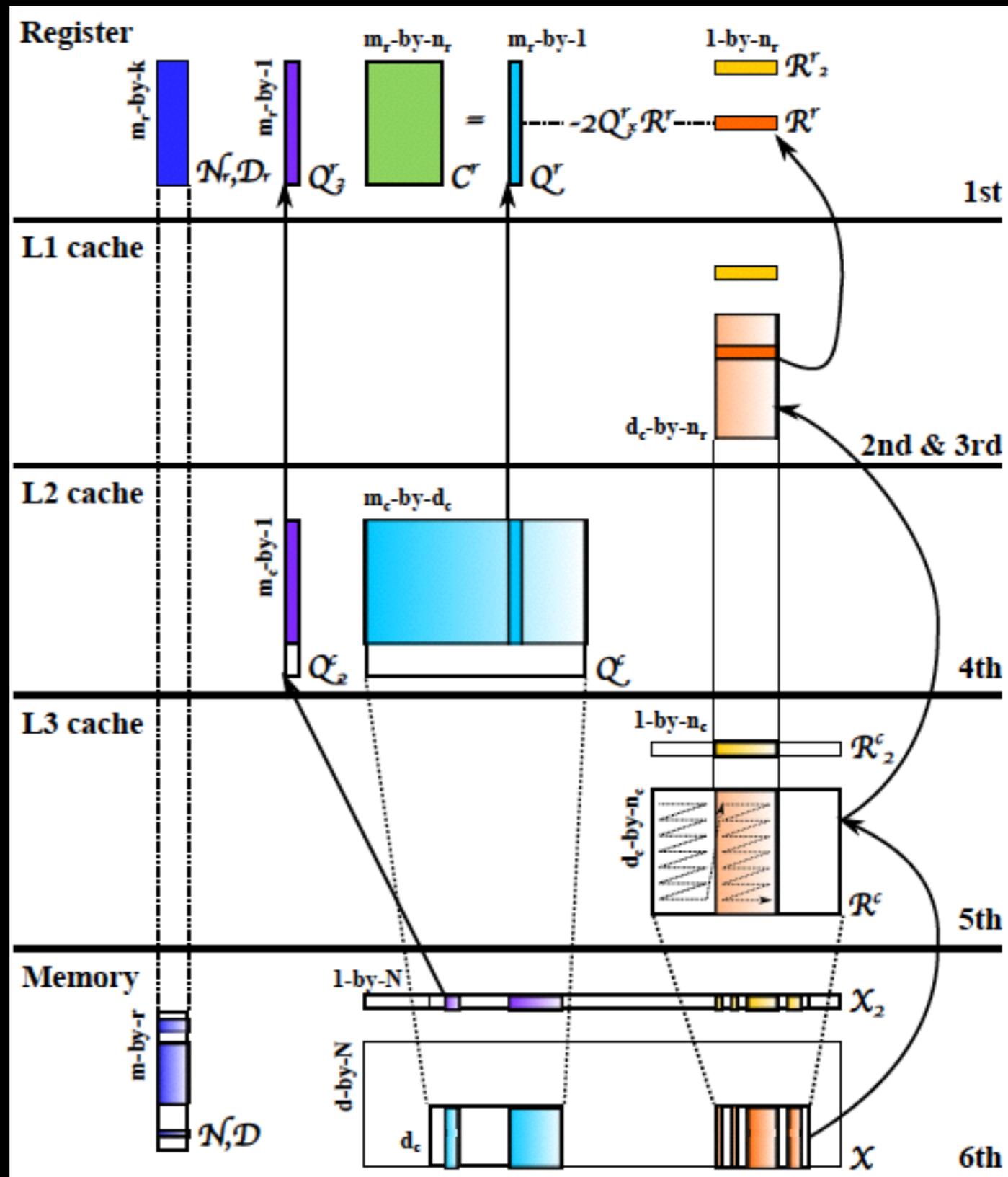
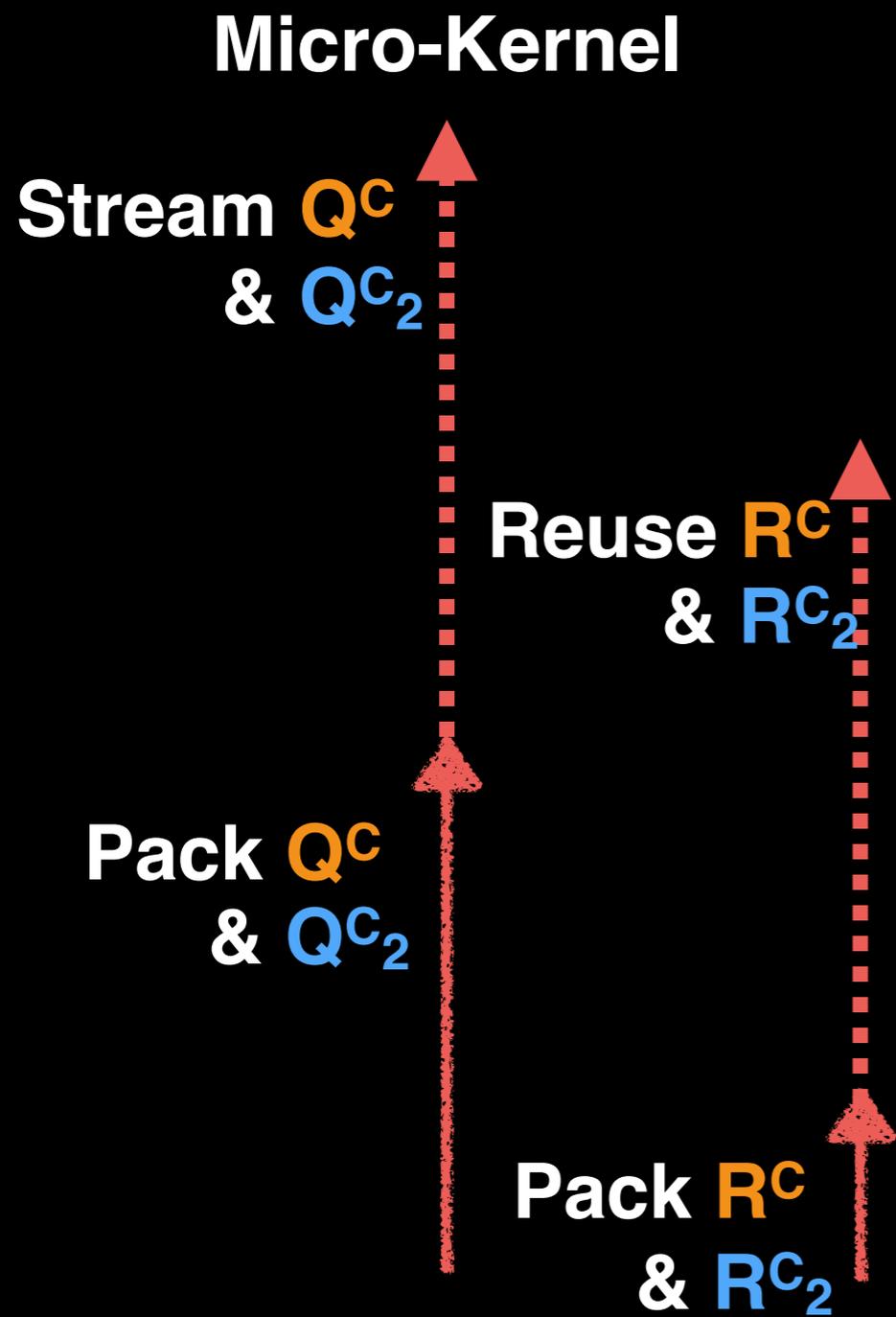
Code Fusion in BLIS *Slice and Dice!*

Code fusion is done in micro-kernel,
and the BLIS framework is maintained.

neighbor lists

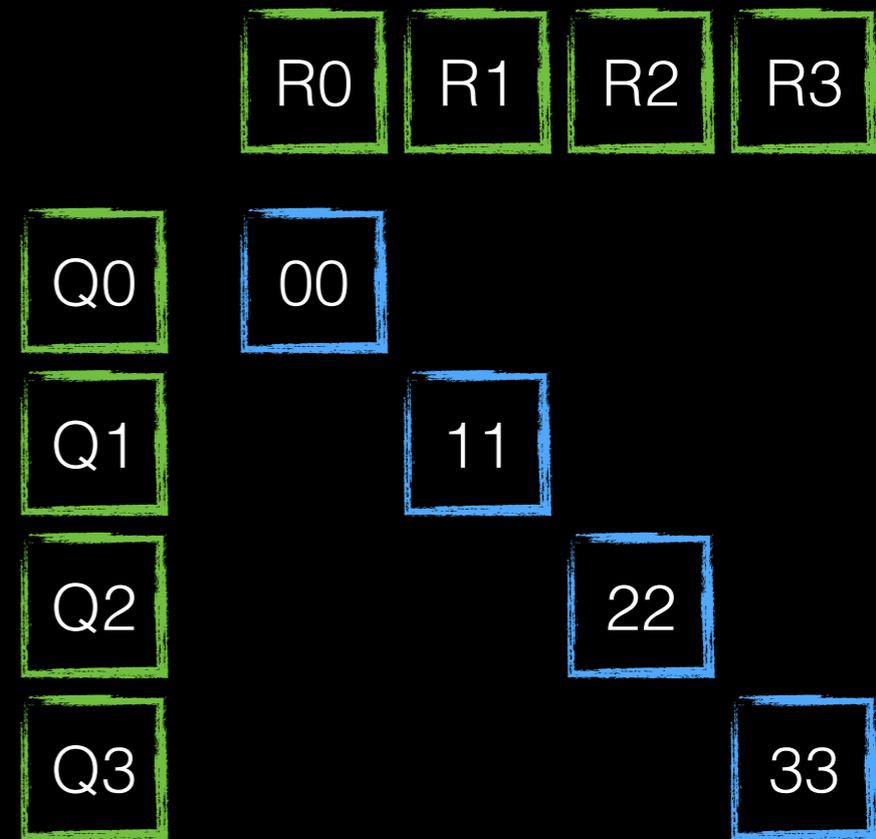


GSKNN and BLIS ($K=Q^T R$)



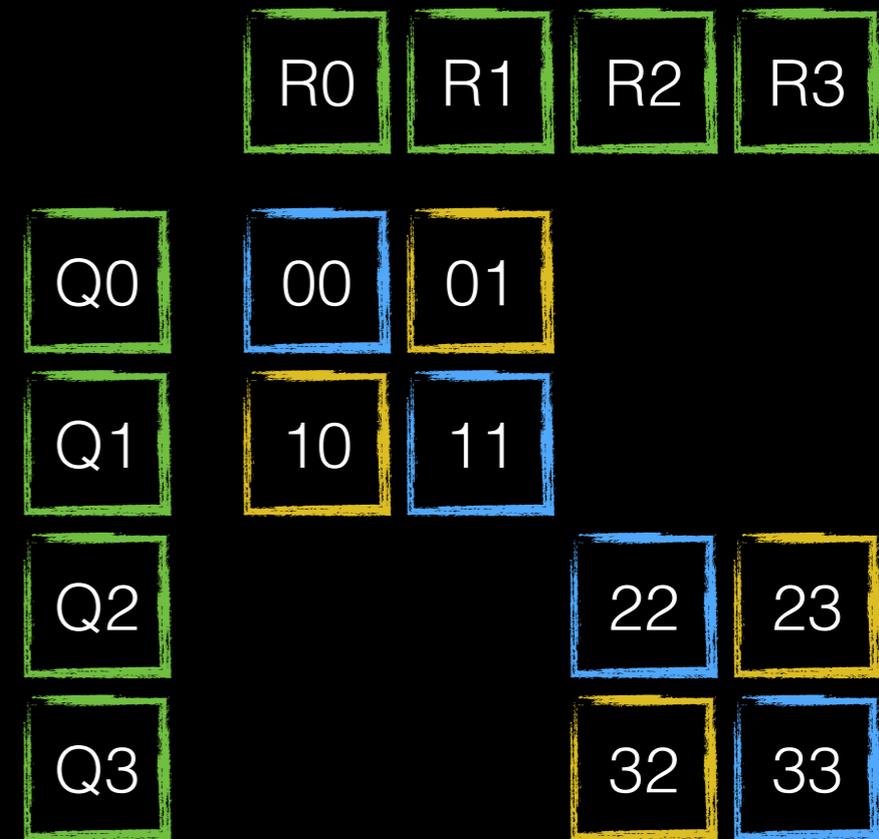
Micro-Kernel

LOAD Q
LOAD R
FMA Q, R, C03_0
SHUFFLE
FMA Q, R, C03_1
PERMUTE2F128
FMA Q, R, C03_2
SHUFFLE
FMA Q, R, C03_3



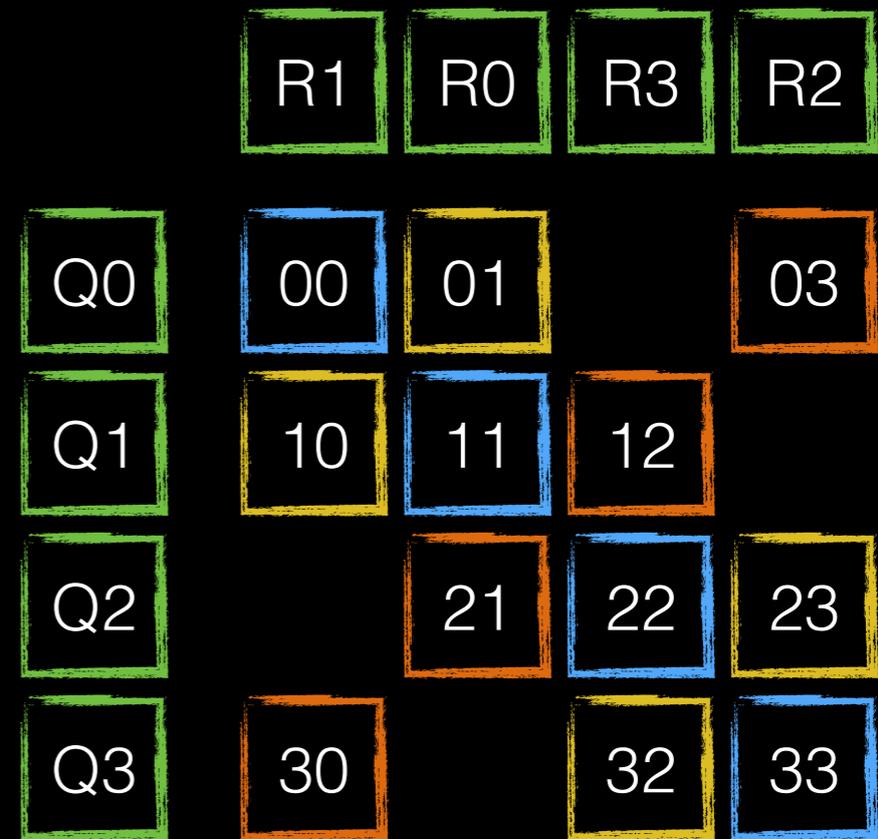
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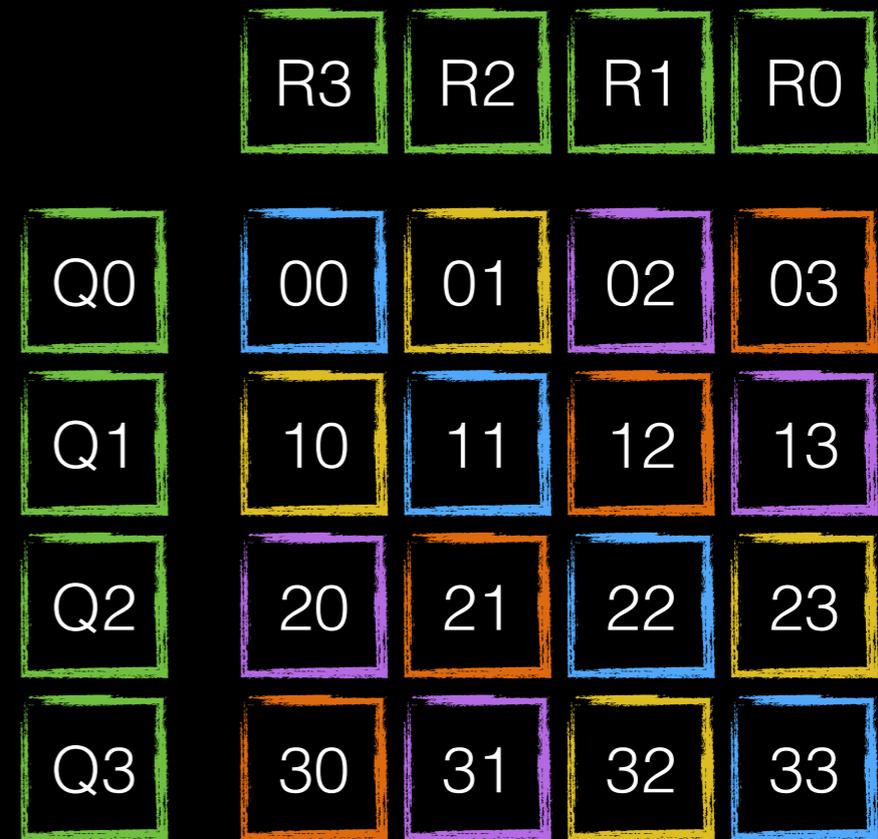
Micro-Kernel

LOAD Q
LOAD R
FMA Q, R, C03_0
SHUFFLE
FMA Q, R, C03_1
PERMUTE2F128
FMA Q, R, C03_2
SHUFFLE
FMA Q, R, C03_3



Micro-Kernel

LOAD Q
LOAD R
FMA Q, R, C03_0
SHUFFLE
FMA Q, R, C03_1
PERMUTE2F128
FMA Q, R, C03_2
SHUFFLE
FMA Q, R, C03_3



Micro-Kernel with p-norm

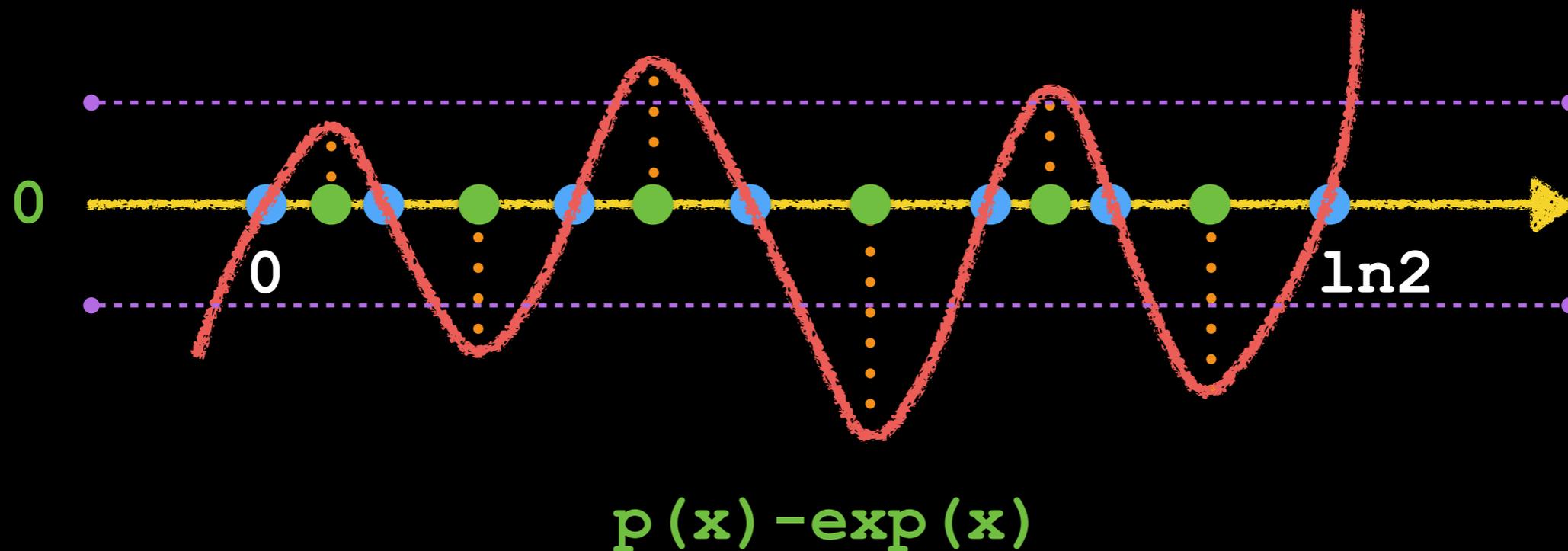
LOAD	Q			1-norm
LOAD	R			SUB
FMA	Q, R, C03_0			AND (flip signed bit)
SHUFFLE				ADD
FMA	Q, R, C03_1			inf-norm
PERMUTE2F128				SUB
FMA	Q, R, C03_2			AND (flip signed bit)
SHUFFLE				MAX
FMA	Q, R, C03_3			p-norm
				SUB
				POW (SVML)
				ADD

Vectorized Math Functions

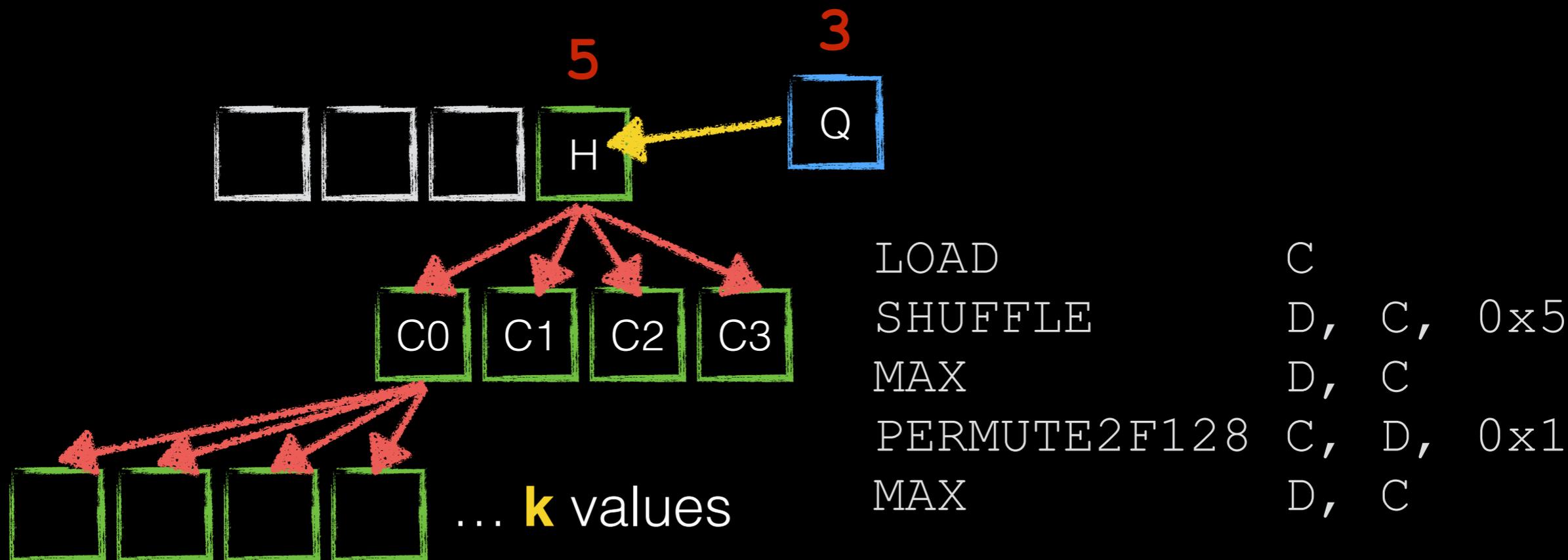
- With a high precision (**20** digits in decimal), **Remez exchange algorithm** can generate an 11 order near minimax polynomial with **1E-18** relative error.

$$P_{11}(x) = c_{11} + (\dots + (c_5 + (c_4 + (c_3 + (c_2 + (c_1 + c_0x)x)x)x)x)x\dots)x$$

$$= 1\text{ADD} + 11\text{FMA}$$



Vectorized Max Heap



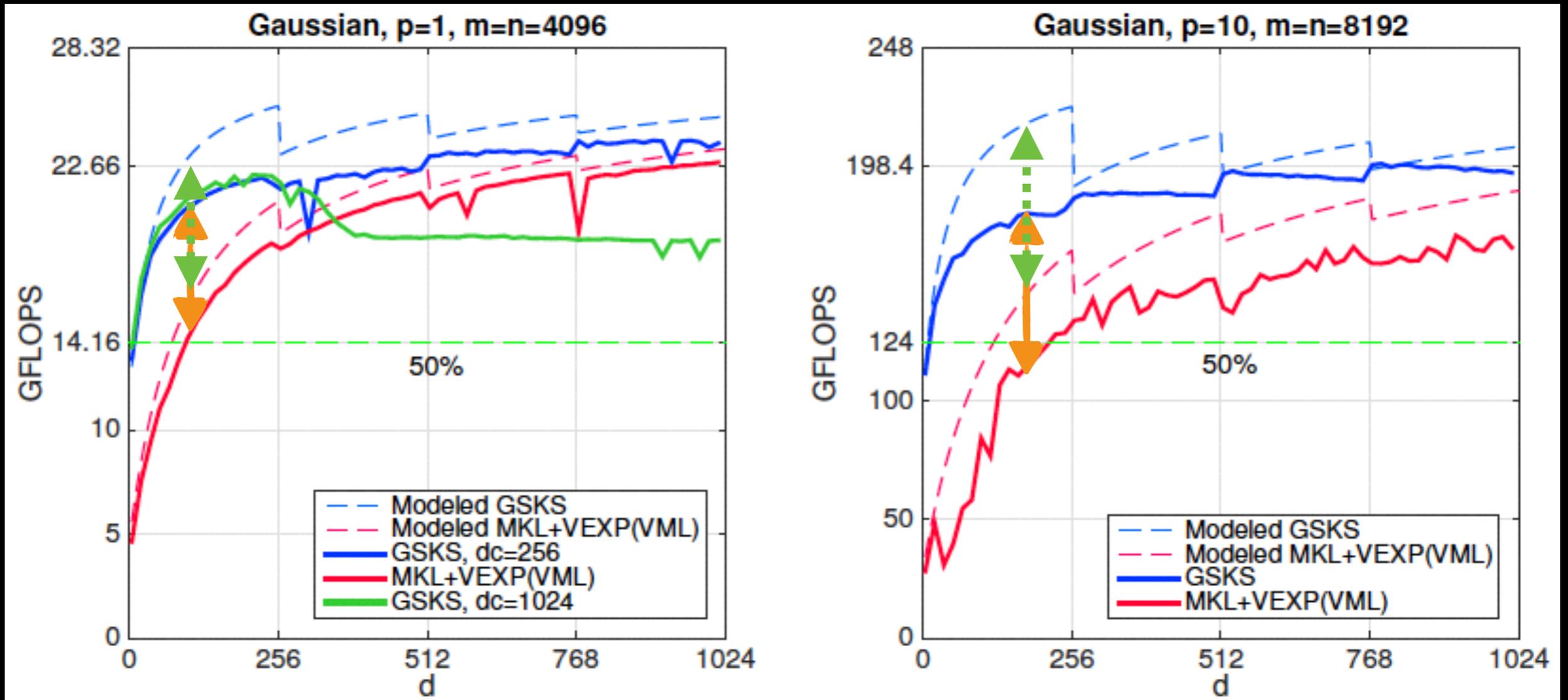
Find the max child

C: [1, 3, 4, 2] → [4, 3, 4, 3] → [4, 4, 4, 4]
 D: [4, 2, 1, 3] [3, 4, 3, 4]



Efficiency Analysis

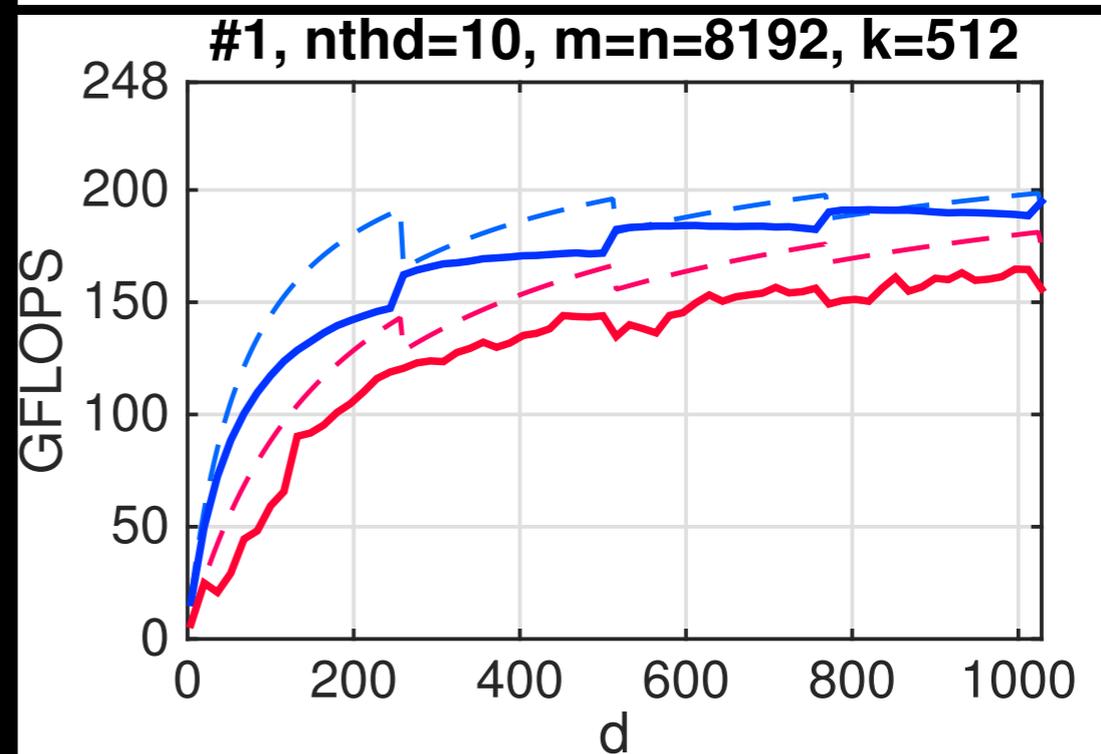
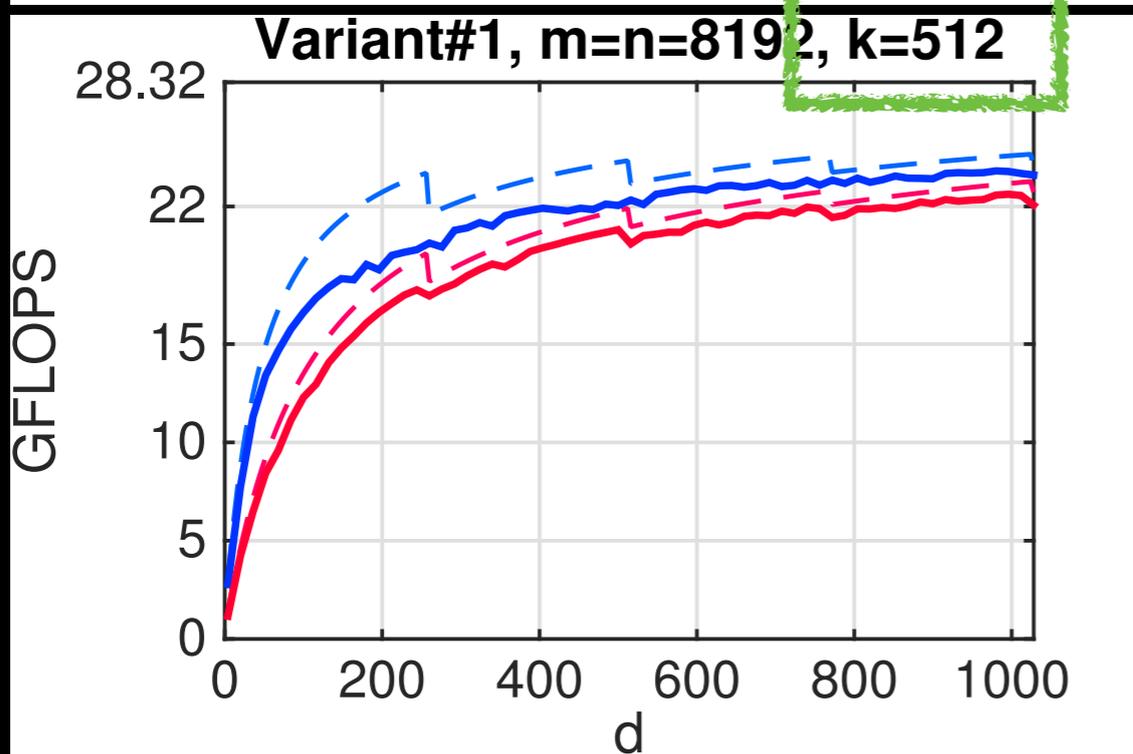
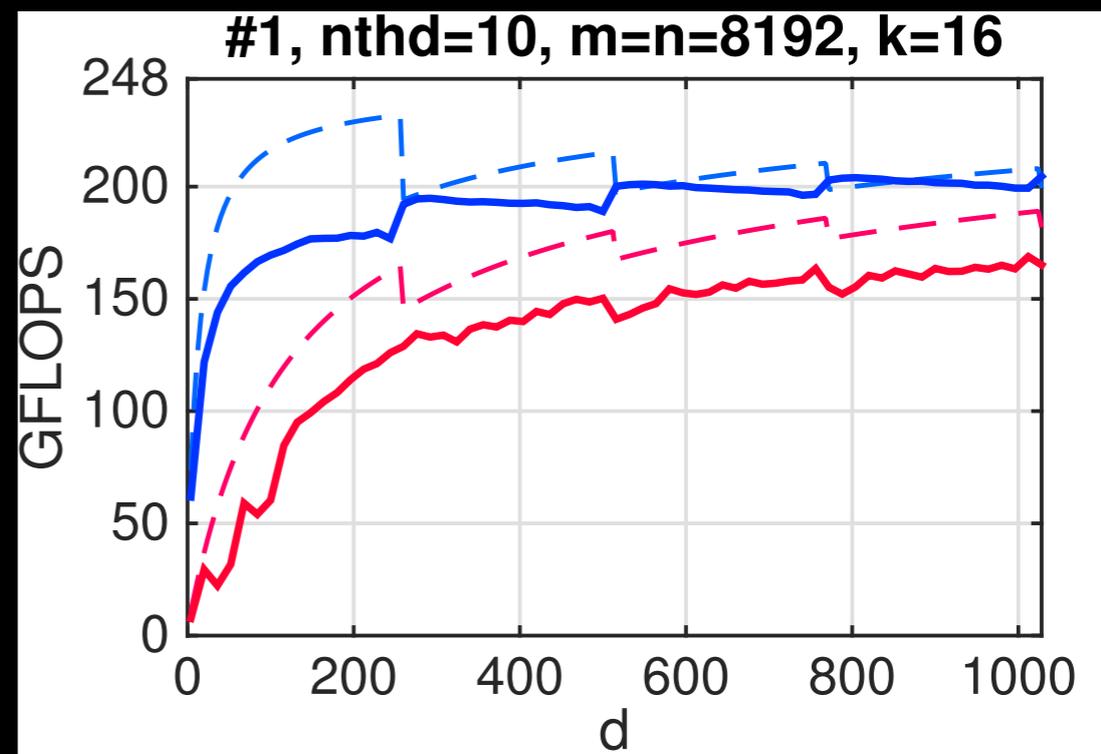
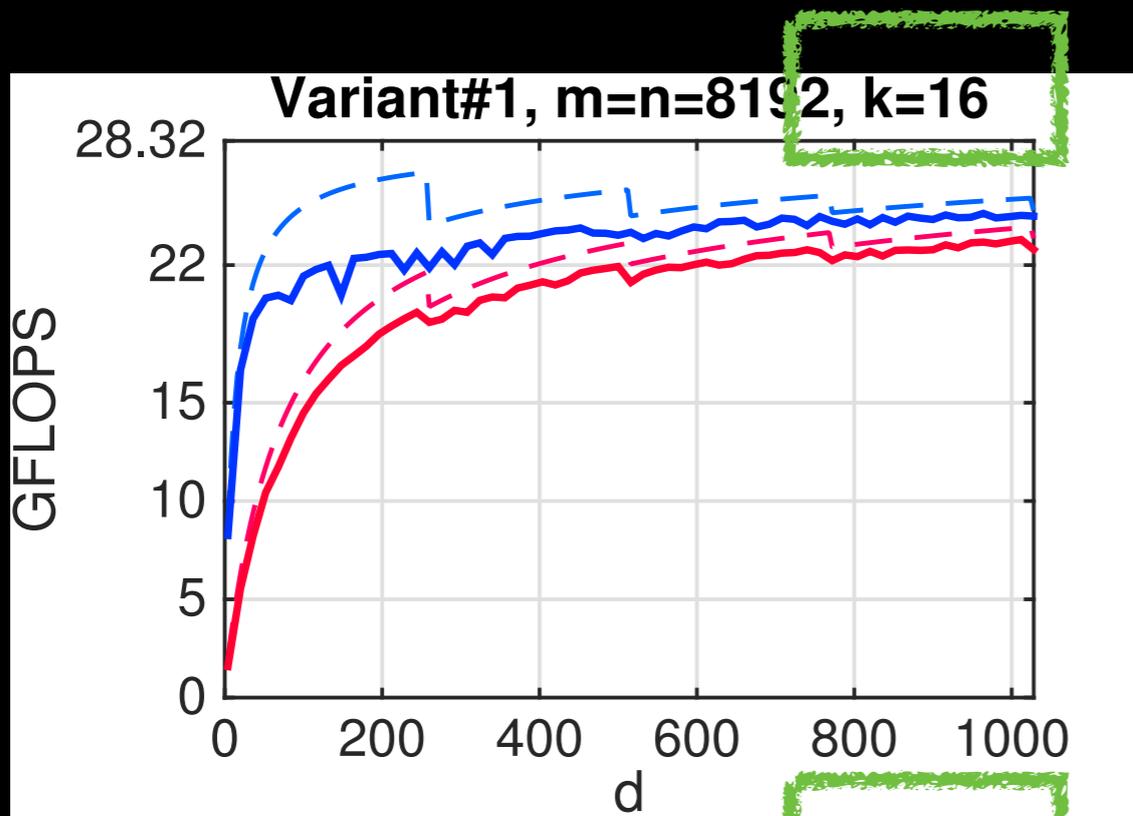
$$T_{BLAS} = T_{GSKS} + T_R + T_Q + T_K$$



$$\frac{mn(2d+36)}{T_{GSKS}} - \frac{mn(2d+36)}{T_{BLAS}} = ?$$

GSKNN Efficiency Graphs

Memory Bound



Conclusion

- The **GEMM** approach in N-body problems is a good example to show the current BLAS library is lacking flexibility for lower level integration.
- The algorithmic innovation of **GSKS** and **GSKNN** is to break through the interface, seeking for the lowest memory complexity.
- We exploit these observations with the help of the **BLIS** framework.
- Ongoing work includes other operations. e.g. kernel inversion, **k**-meaning clustering. Port to GPU and other accelerators.

Question?

GSKS

GSKNN

github.com/ChenhanYu/ks

github.com/ChenhanYu/rnn