

ACL2 Support for Interactive Proof

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OUTLINE

[Introduction](#)

[Background](#)

[Demos \(I\)](#)

[Rewriting in ACL2](#)

[Demos \(II\)](#)

[Very Brief Survey of ACL2 Features](#)

[Conclusion](#)

OUTLINE

[Introduction](#)

[Background](#)

[Demos \(I\)](#)

[Rewriting in ACL2](#)

[Demos \(II\)](#)

[Very Brief Survey of ACL2 Features](#)

[Conclusion](#)

INTRODUCTION (1)

Quoting the [ACL2 home page](#):

*ACL2 is a **logic** and **programming language** in which you can model computer systems, together with a **tool** to help you prove properties of those models. “ACL2” denotes “**A Computational Logic for Applicative Common Lisp**”.*

Goal for this talk: Give a sense of the ACL2 system, especially how it supports **user interaction**.

Confession: there is considerable overlap with KeY invited talk given last month.

But I may skip some material. I hope to leave lots of time for discussion. **Please ask questions** during the talk!

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

BACKGROUND: OVERVIEW

- ▶ ACL2 is freely available, including libraries of *certifiable books*, from the [ACL2 home page](#).
- ▶ ACL2 is written mostly in itself (!).
 - ▶ About 10 MB of source code (Version 7.1).
- ▶ *Bleeding edge* for libraries (*community books*) and the ACL2 system are [available from Github](#).
 - ▶ Well over 400,000 *events* (theorems, definitions, other) are evaluated in the community books.
- ▶ [Workshop series: #13](#) is at UT, Oct. 1-2, 2015.

BACKGROUND: OVERVIEW (CONTINUED)

Development history:

- ▶ Bob Boyer and J Moore started ACL2 in 1989. I joined and Bob dropped out in 1993. J and I continue its development.
- ▶ *Boyer-Moore Theorem Provers* go back to the start of their collaboration in 1971.

Industrial usage: As far as I know, ACL2 is the only interactive theorem prover (ITP) used with some regularity at several companies:

- ▶ AMD, Centaur, IBM, Intel, Oracle, Rockwell Collins

There are also users in the **U.S. Government** and **universities**, including —

- ▶ UT Austin: x86 interpreter defined in ACL2, validation by co-simulation, proofs about x86 machine code

BACKGROUND: LOGIC

The ACL2 logic is a **first-order** logic with induction up to ε_0 .

But all ACL2 theories extend a given *ground-zero* theory, which axiomizes data types for:

- ▶ numbers (complex rationals), characters, strings, symbols;
- ▶ trees and lists, using a pairing operation (`cons`).

ACL2 extensions are *conservative* (a demo will discuss this).

- ▶ See M. Kaufmann and J Moore, “Structured Theory Development for a Mechanized Logic.” *Journal of Automated Reasoning* 26, no. 2 (2001) 161-203.

BACKGROUND: STRENGTHS

- ▶ Proof automation
- ▶ [Support for user interaction](#)
- ▶ Fast execution
- ▶ [Documentation](#) (about 100,000 lines for just the system)
- ▶ Interfaces include **Emacs**
(*Is that really an interface? A strength?*)
and the Eclipse-based [ACL2 Sedan](#).

A potential weakness: first-order logic with only basic quantifier support (but recursion helps).

BACKGROUND: MORE INFORMATION

NOTE: A longer variant of this talk, but targeted to CS grad students and with more focus on *using* ACL2, is here:

<http://www.cs.utexas.edu/users/kaufmann/talks/acl2-intro-2015-04/acl2-intro.pdf>

That talk mentions this link to several demos and their logs:

<http://www.cs.utexas.edu/users/kaufmann/talks/acl2-intro-2015-04/demos.tgz>

BACKGROUND: HIGHER-ORDER LOGIC?

Bob Boyer and I did some work in the mid-1980s on adapting the Boyer-Moore prover to the programming language, SASL, [which is a weakly typed ancestor of Haskell](#).

Our work is described at some length in a [Burroughs Technical Report](#), and is summarized briefly here:

Bob Boyer and Matt Kaufmann. A Prototype Theorem-Prover for a Higher-Order Functional Language (with R. Boyer). In *Proceedings of VERkshop III — a formal verification workshop*, Watsonville, CA, February 1985. ACM SIGSOFT Software Engineering Notes, Volume 10, Issue 4, August 1985, ACM, New York, NY, USA.

Not included above is a larger example (a SASL unification program).

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

DEMOS (I)

- ▶ All demos today, with logs, are available from [this link to a gzipped tarfile](#) `demo.s.tar.gz`.
- ▶ ACL2 programming and evaluation
[DEMO]: file `demo-1.lsp`
(log `demo-1-log.txt`)
- ▶ ACL2 as an automatic theorem prover
[DEMO]: file `demo-2.lsp`
(log `demo-2-log.txt`)
 - ▶ ACL2 provides **automation** for induction, linear arithmetic, Boolean reasoning, rule application, . . .
 - ▶ . . . but lemmas are usually needed (sometimes from libraries).

DEMOS (I) (CONTINUED)

ACL2 supports formally verified extensions.

In particular, [GL](#) is a *verified clause processor* defined and verified by an ACL2 user, Sol Swords.

GL does proofs about finite domains by *bit-blasting*.

The next demo illustrates GL. It also shows the use of [LOCAL](#), for “private” events (using *conservativity*).

```
[DEMO]: book demo-gl.lisp  
(log demo-gl-log.txt)
```


OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

REWRITING IN ACL2

ACL2 is typically controlled with [conditional rewrite](#) rules, although there are other [rule-classes](#) too.

The basic idea: the ACL2 rewriter automatically applies the rule

$$H \longrightarrow L = R$$

by replacing an instance L/s of L by R/s , when the rewriter can verify H/s .

The documentation topic for [rewrite](#) shows many ways to control the rewriter (needed only occasionally). I'll mention only a few:

REWRITING IN ACL2 (2)

- ▶ **backchain-limit**: limit effort to relieve hypotheses
- ▶ **force**: defer proving H if necessary
- ▶ **hide**: hide a term from the rewriter
- ▶ **syntaxp**: attach a heuristic filter on a rule

Example of `syntaxp`: consider $3 + (4 + x)$.

It's already in normal form: right-associated.

Our wish: $3 + (4 + x) = (3 + 4) + x = 7 + x$.

```
ACL2 !>:pe associativity-of-+
-997 (DEFAXIOM ASSOCIATIVITY-OF-+
      (EQUAL (+ (+ X Y) Z) (+ X (+ Y Z))))

ACL2 !>:pe fold-consts-in-+
-158 (DEFTHM FOLD-CONSTS-IN-+
      (IMPLIES (AND (SYNTAXP (QUOTE X))
                    (SYNTAXP (QUOTE Y)))
                (EQUAL (+ X (+ Y Z)) (+ (+ X Y) Z))))

ACL2 !>
```

REWRITING IN ACL2 (3)

$H \longrightarrow L = R$ (*ordinary rewrite rule*)

$H \longrightarrow L \sim R$ (*congruence-based rewrite rule*)

where $\sim$ is an equivalence relation.

Users prove [equivalence](#), [congruence](#), and [refinement](#) rules, to tell ACL2 when such rewrites are valid.

B. Brock, M. Kaufmann, and J S. Moore. Rewriting with Equivalence Relations in ACL2. *J. Aut. Reasoning* 40 (2008).

M. Kaufmann and J S. Moore. Double Rewriting for Equivalential Reasoning in ACL2. *Proc. ACL2 Workshop 2006*.

M. Kaufmann and J S. Moore. Rough Diamond: An Extension of Equivalence-based Rewriting. *Proc. ITP 2014*.

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

DEMOS (II)

Our final demo shows how ACL2 proof development often follows “*the method*”:

- ▶ ACL2 heuristically chooses and applies a destructor-style **induction scheme**.
- ▶ ACL2 **simplifies** the base and induction steps.
- ▶ The user looks at **key checkpoints**, which are unproved goals printed by ACL2.
- ▶ The user formulates **conditional rewrite rules** to simplify those checkpoints.
- ▶ Repeat.

DEMOS (II) (CONTINUED)

Quoting the documentation for [induction](#):

... the interested reader should see Chapter XIV of A Computational Logic (Boyer and Moore, Academic Press, 1979) which represents a fairly complete description of the induction heuristics of ACL2.

This demo should give a sense of how ACL2 chooses (and applies) induction schemes,
but the focus will be on [user interaction](#).

— [DEMO](#) (excerpted from [my TPHOLs 2008 talk](#)) —

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

VERY BRIEF SURVEY OF ACL2 FEATURES

Let's see a bit more about how ACL2 supports proof development . . .

- ▶ . . . by exploring briefly the [ACL2 documentation](#).

NOTE:

I would be very happy to elaborate on any of these topics!

VERY BRIEF SURVEY OF ACL2 FEATURES (2)

In particular, we might explore a few [debugging](#) features, as time and interest permit.

- ▶ [accumulated-persistence](#)
- ▶ [break-rewrite](#)
- ▶ [cgen](#)
- ▶ [cw-gstack](#)
- ▶ [disassemble\\$](#)
- ▶ [dmr](#)
- ▶ [failed-forcing](#)
- ▶ [failure](#)
- ▶ [find-lemmas](#)
- ▶ [forward-chaining-reports](#)
- ▶ [guard-debug](#)
- ▶ [measure-debug](#)
- ▶ [nil-goal](#)
- ▶ [print-gv](#)
- ▶ [profile-acl2](#)
- ▶ [profile-all](#)
- ▶ [proof-checker](#)
- ▶ [proof-tree](#)
- ▶ [pstack](#)
- ▶ [quick-and-dirty-subsumption-replacement-step](#)
- ▶ [redo-flat](#)
- ▶ [remove-hyps](#)
- ▶ [set-debugger-enable](#)
- ▶ [set-guard-msg](#)
- ▶ [sneaky](#)
- ▶ [spacewalk](#)
- ▶ [splitter](#)
- ▶ [time-tracker](#)
- ▶ [trace](#)
- ▶ [walkabout](#)
- ▶ [watch](#)

OUTLINE

Introduction

Background

Demos (I)

Rewriting in ACL2

Demos (II)

Very Brief Survey of ACL2 Features

Conclusion

OUTLINE

[Introduction](#)

[Background](#)

[Demos \(I\)](#)

[Rewriting in ACL2](#)

[Demos \(II\)](#)

[Very Brief Survey of ACL2 Features](#)

[Conclusion](#)

CONCLUSION

- ▶ ACL2 has a 25 (or 44) year history and is used in industry.
“Microprocessor design goes daily through numerous optimizations that affect thousands of lines of code. These optimizations must be proved correct.”
— Anna Slobodova, verification manager at Centaur Technology
- ▶ ACL2 provides automation but scales to large problems . . .
. . . with libraries and by supporting **user interaction**.
- ▶ For more information, see the [ACL2 home page](#).

THANK YOU!