

# RESPONSE TIME DISTRIBUTIONS FOR A MULTI-CLASS QUEUE WITH FEEDBACK

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## Abstract

A single server queue with feedback and multiple customer classes is analyzed. Arrival processes are independent Poisson processes. Each round of service is exponentially distributed. After receiving a round of service, a customer may depart or rejoin the end of the queue for more service. The number of rounds of service required by a customer is a random variable with a general distribution. Our main contribution is characterization of response time distributions for the customer classes. Our results generalize in some respects previous analyses of processor-sharing models. They also represent initial efforts to understand response time behavior along paths with loops in local balanced queueing networks.

## 1. INTRODUCTION

Many service facilities can be modeled as a feedback queue such as shown in Figure 1. Of interest in this paper is a single-server queue with infinite waiting room and  $Q$  types of customers. The arrival process of type  $q$  customers is an independent Poisson process ( $q = 1, 2, \dots, Q$ ). Each new arrival joins the end of the queue. The customer at the head of the queue receives from the server a round of service, which is an independent exponentially distributed random variable with mean  $1/\mu$  seconds. After receiving a round of service, a customer may depart or rejoin the end of the queue for more service. The number of rounds of service required by a type  $q$  customer is a random variable with a general probability distribution  $\{a_r^{(q)}, r = 1, 2, \dots, R\}$  where  $a_r^{(q)}$  is the probability of a type  $q$  customer requiring exactly  $r$  rounds of service.

The queue length distribution of the above model is readily available since the feedback queue described is an open queueing network satisfying local balance [1]. The contribution of this paper is to characterize response time distributions of the different types of customers; specifically, we solved for the conditional response time distributions of an arbitrary customer requiring  $r$  rounds of service for  $r = 1, 2, \dots, R$ .

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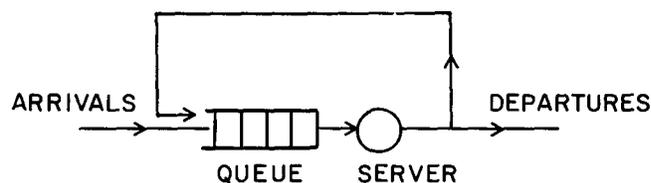


Fig. 1. A feedback queue model.

### Relationship to prior work

Our feedback queue model is like a time-sharing model with exponentially distributed service "quantums." Time-sharing models were first studied by Kleinrock [2] who solved for the mean response time of a customer conditioning on his (total) service requirement. He considered two cases: (a) constant quantum size  $\Delta$ , and (b) the limiting case of  $\Delta \rightarrow 0$  called processor-sharing. Customers are assumed to arrive according to a Poisson process. In case (a), the number of service quanta required by a customer is geometrically distributed. In case (b), the service requirements are characterized by an exponential distribution. (This is called the processor-sharing M/M/1 queue.) Kleinrock's conditional mean response time result was later shown to hold for a processor-sharing M/G/1 queue (i.e. service requirements characterized by a general distribution) as well by Sakata, Noguchi and Oizumi [3]. Higher order response time statistics are much harder to get. The response time distribution for the processor-sharing M/M/1 queue was obtained by Coffman, Muntz, and Trotter [4]. The response time distribution for the constant quantum size case was obtained by Muntz [5] assuming exponentially distributed service requirements.

Our feedback queue model is different from the time-sharing models in several respects. A round of service in our model, corresponding to a service quantum in time-sharing models, is exponentially distributed. Our model can be used, however, to approximate processor-sharing by making  $1/\mu$  very small relative to the mean service requirement.

Aside from the quantum size assumption, our model is more general than those of [4,5] in two respects: (i) multiple types of customers, and (ii) the number of rounds of service for each customer type has a general probability distribution. Specifically, distributions of service requirements that are admissible in our model are those with moment generating functions of the form

$$B_q^*(s) = \sum_{r=1}^R a_r^{(q)} \left( \frac{\mu}{s+\mu} \right)^r \quad (1)$$

Our model is also different from the feedback queue model of Takács [6]. In his model, each round of service can have a general distribution. However, he considered one type of customers only and the number of rounds of service required by a customer is geometrically distributed; in other words, after each round of service, a customer always departs with probability  $(1-p)$  and rejoins the end of the queue with probability  $p$  (memoryless behavior).

The original motivation of this work stems from our efforts to characterize the response time in a network of queues. For a network of FCFS queues that satisfies local balance, J. Wong [7] found the response time distribution of customers traversing loop-free paths. Our results in this paper represent efforts to understand the response time behavior along paths with loops in the simplest form of queueing networks satisfying local balance.

### Assumptions and definitions

Consider the following example of 2 types of customers. Type 1 customers arrive according to a Poisson process with rate  $\alpha_1$  customers per second. The number of rounds of service required by a type 1

customer has the probability distribution

$$a_r^{(1)} = \begin{cases} 1/100 & r = 1, 2, \dots, 100 \\ 0 & \text{otherwise} \end{cases}$$

Type 2 customers arrive according to a Poisson process with rate  $\alpha_2$  customers per second. The number of rounds of service required by a type 2 customer has the probability distribution

$$a_r^{(2)} = \begin{cases} 1/10 & r = 1, 2, \dots, 10 \\ 0 & \text{otherwise} \end{cases}$$

Using the properties of Poisson processes, the above model is equivalent to the following model with 100 types of customers. Type  $r$  customers ( $r = 1, 2, \dots, 100$ ) require exactly  $r$  rounds of service and arrive according to a Poisson process with rate

$$\gamma_r = \begin{cases} 0.01\alpha_1 + 0.1\alpha_2 & r = 1, 2, \dots, 10 \\ 0.01\alpha_1 & r = 11, 12, \dots, 100 \\ 0 & \text{otherwise} \end{cases}$$

We shall, without any loss of generality, consider the following model. There are  $R$  types of customers. The arrival process of the  $r^{\text{th}}$  type is Poisson at rate  $\gamma_r$  customers per second. A type  $r$  customer requires exactly  $r$  rounds of service. It should be obvious that if we can derive response time distributions for this model, response time distributions for any model with  $Q$  customer types and service time requirements characterized by Eq. (1) can be easily obtained.

Let  $t_r$  be the response time of attaining exactly  $r$  rounds of service;  $r = 1, 2, \dots, R$  and obviously  $t_0 = 0$ . We shall solve for its moment generating function

$$T_r^*(s) = E[e^{-st_r}]$$

where  $E[\cdot]$  denotes the expectation of the function of random variable(s) inside the brackets.

We shall only consider steady-state results. For a single-server queue, stationarity is assured if the traffic intensity  $\rho < 1$  where  $\rho = \sum_{r=1}^R \gamma_r(r/\mu)$ ; see Cohen [8].

Customers in the queue are differentiated into  $R$  different classes; class  $k$  consists of all those customers in the queue with exactly  $k$  more rounds of service to go, where  $k = 1, 2, \dots, R$ .

Let us follow the progress of a "tagged" customer and introduce some more notation. Upon his initial arrival, the tagged customer finds  $n_k$  class  $k$  customers in the queue ( $k = 1, 2, \dots, R$ ). The system state thus found at an arrival instant is denoted by  $\underline{n} = (n_1, n_2, \dots, n_R)$  and is described by the moment generating function

$$P^*(\underline{z}) = E[z_1^{n_1} z_2^{n_2} \dots z_R^{n_R}]$$

where  $\underline{z}$  is the shorthand notation for  $(z_1, z_2, \dots, z_R)$ .

At the end of the tagged customer's  $r^{\text{th}}$  round of service (given that he requires at least  $r$  rounds), let the system state at that instant be denoted by  $\underline{m}^{(r)} = (m_1^{(r)}, m_2^{(r)}, \dots, m_R^{(r)})$  where  $m_k^{(r)}$  is the number of customers who have exactly  $k$  more rounds of service to go. Define  $M^{(r)} = \sum_{k=1}^R m_k^{(r)}$ .

In order to characterize  $T_r^*(s)$ , we shall need to first characterize the joint distribution of  $t_r$  and  $\underline{m}^{(r)}$ , which is described by

$$U_r^*(s, \underline{z}) = E[e^{-st} z_1^{m_1^{(r)}} z_2^{m_2^{(r)}} \dots z_R^{m_R^{(r)}}]$$

### Summary of results

We derived a recursive equation relating  $U_{r+1}^*(s, \underline{z})$  to  $U_r^*(s, \underline{z})$  [Lemma 2]. An explicit solution of  $U_r^*(s, \underline{z})$  was found, from which  $T_r^*(s)$  was obtained [Theorem 1]. We then proved that the stationary distribution of  $\underline{m}^{(r)}$ ,  $r = 1, 2, \dots, R$ , is the same as that of  $\underline{n}$  [Theorem 2]. With this result, we solved for the mean value of  $t_r$  [Theorem 3]; this last result is similar to the mean response time result of processor-sharing models. We also obtained an efficient recursive algorithm to calculate the second order statistics of  $t_r$  [Theorem 4]. Some numerical results are shown in Section 3.

## 2. THE ANALYSIS

Consider the system state  $\underline{n} = (n_1, n_2, \dots, n_R)$  at arrival instants. Recall that  $n_k$  is the number of class  $k$  customers with exactly  $k$  more rounds of service to go. The aggregate arrival rate of customers to the  $k^{\text{th}}$  class is

$$\lambda_k = \sum_{i=k}^R \gamma_i \quad (2)$$

since any new arrival who requires at least  $k$  rounds of service must enter and leave the  $k^{\text{th}}$  class exactly once.

Lemma 1. The moment generating function of  $\underline{n}$  is

$$P^*(\underline{z}) = \frac{1-\rho}{1 - \sum_{k=1}^R \rho_k z_k} \quad (3)$$

where  $\rho_k = \lambda_k/\mu$  and  $\rho = \sum_{k=1}^R \rho_k$ .

Proof. Given Poisson arrival processes, the system state probabilities at an arrival instant are the same as system state probabilities at a random time instant [9]. With each round of service being exponentially distributed with the same mean  $(1/\mu)$ , we have an open queueing network that satisfies local balance [1].

Eq. (3) has been obtained by Reiser and Kobayashi [10]. (Q. E. D.)

Since each round of service is exponentially distributed, it has the moment generating function

$$B^*(s) = \frac{\mu}{s + \mu} \quad (4)$$

A recursive solution of  $U_r^*(s, \underline{z})$  is next given.

Lemma 2.

$$U_0^*(s, \underline{z}) = P^*(\underline{z}) \quad (5)$$

$$U_{r+1}^*(s, \underline{z}) = \gamma_1(s, \underline{z}) U_r^*(s, \underline{y}(s, \underline{z})) \quad r \geq 0 \quad (6)$$

where

$$y(s, \underline{z}) = (y_1(s, \underline{z}), y_2(s, \underline{z}), \dots, y_R(s, \underline{z})),$$

$$y_1(s, \underline{z}) = B^*(s + \sum_{i=1}^R \gamma_i(1 - z_i)),$$

and

$$y_k(s, \underline{z}) = z_{k-1} y_1(s, \underline{z}) \quad \text{for } 2 \leq k \leq R$$

Proof. For  $r = 0$ ,  $t_0 = 0$  and  $\underline{m}^{(0)} = \underline{n}$ . This and the definition of  $U_r^*(s, \underline{z})$  yield (5) at once.

To show (6), consider the time period between  $t_r$  and  $t_{r+1}$  during which the server served  $M^{(r)} + 1$  customers, where  $M^{(r)} = \sum_{k=1}^R m_k^{(r)}$  and the extra one is for the tagged customer's  $(r + 1)^{st}$  round. During the same time period, each class  $k$  customer became a class  $(k - 1)$  customer where  $k = 2, 3, \dots, R$ . Furthermore, let  $A_k(t)$  be the number of external new arrivals to class  $k$  during time  $t (= t_{r+1} - t_r)$  according to a Poisson process of rate  $\gamma_k$  customers per second. We note that class  $R$  is an exception in that its  $m_R^{(r+1)}$  customers are all new arrivals. Thus, conditioning on  $t_r$  and  $\underline{m}^{(r)}$ , we have

$$U_{r+1}^*(s, \underline{z}/t_r, \underline{m}^{(r)}) = E[e^{-s(t+t_r)} z_1^{m_2^{(r)} + A_1(t)} z_2^{m_3^{(r)} + A_2(t)} \dots z_R^{A_R(t)} / t_r, \underline{m}^{(r)}]$$

$$= e^{-st} \prod_{k=2}^R z_{k-1}^{m_k^{(r)}} E[e^{-st} z_1^{A_1(t)} z_2^{A_2(t)} \dots z_R^{A_R(t)} / M^{(r)}].$$

The last quantity on the right hand side is  $(y_1(s, \underline{z}))^{M^{(r)} + 1}$  because  $t$  is the sum of  $M^{(r)} + 1$  independent identically distributed random variables with the moment generating function  $B^*(s)$ . The above equation can be rewritten as

$$U_{r+1}^*(s, \underline{z}/t_r, \underline{m}^{(r)}) = y_1(s, \underline{z}) \{ e^{-st} y_1(s, \underline{z}) \prod_{k=2}^R [z_{k-1} y_1(s, \underline{z})]^{m_k^{(r)}} \}$$

Unconditioning on  $t_r$  and  $\underline{m}^{(r)}$ , (6) follows. (Q. E. D.)

Explicit solutions for  $U_r^*(s, \underline{z})$  and  $T_r^*(s)$  can now be shown.

Theorem 1. (i)  $U_r^*(s, \underline{z}) = \frac{1 - \rho}{P_r(s) - \sum_{k=1}^R Q_{k,r}(s) z_k}$  (7)  $r \geq 0$

where  $P_r(s)$  and  $Q_{k,r}(s)$  are polynomials in  $s$  given by

$$\begin{bmatrix} P_r(s) \\ Q_{1,r}(s) \\ Q_{2,r}(s) \\ \vdots \\ Q_{R-1,r}(s) \\ Q_{R,r}(s) \end{bmatrix} = \begin{bmatrix} (1 + \frac{\Sigma}{\mu} + \rho_1) & -1 & 0 & 0 & \dots & 0 \\ \gamma_1/\mu & 0 & 1 & 0 & & 0 \\ \gamma_2/\mu & 0 & 0 & 1 & & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & 1 \\ \gamma_{R-1}/\mu & 0 & 0 & \dots & 0 & 1 \\ \gamma_R/\mu & 0 & 0 & \dots & 0 & 0 \end{bmatrix}^r \begin{bmatrix} 1 \\ \rho_1 \\ \rho_2 \\ \vdots \\ \rho_{R-1} \\ \rho_R \end{bmatrix} \quad (8)$$

$$(ii) \quad T_r^*(s) = \frac{1 - \rho}{P_r(s) - \sum_{k=1}^R Q_{k,r}(s)} \quad (9)$$

Proof. (i) Because of (3) and (5), (7) holds for  $r = 0$  with  $P_0(s) = 1$  and

$Q_{k,0}(s) = \rho_k$  for  $1 \leq k \leq R$ . Assuming that (7) holds for  $r$ , we use (6) and (4) to express  $U_{r+1}^*(s, \underline{z})$  as follows.

$$\begin{aligned} U_{r+1}^*(s, \underline{z}) &= \frac{1}{1 + \frac{s}{\mu} + \sum_{i=1}^R \frac{\gamma_i}{\mu}(1-z_i)} \cdot \frac{1 - \rho}{P_r(s) - \frac{Q_{1,r}(s) - \sum_{k=1}^{R-1} Q_{k+1,r}(s) z_k}{1 + (s/\mu) + \sum_{i=1}^R (\gamma_i/\mu)(1-z_i)}} \\ &= \frac{1 - \rho}{\left\{ \left(1 + \frac{s}{\mu} + \sum_{i=1}^R \frac{\gamma_i}{\mu}\right) P_r(s) - Q_{1,r}(s) \right\} - \sum_{k=1}^{R-1} \left[ \frac{\gamma_k}{\mu} P_r(s) + Q_{k+1,r}(s) \right] z_k - \frac{\gamma_R}{\mu} z_R P_r(s)} \end{aligned}$$

Thus, the form of (7) is maintained, and it is evident from the above that

$$\begin{bmatrix} P_{r+1}(s) \\ Q_{1,r+1}(s) \\ \vdots \\ Q_{R,r+1}(s) \end{bmatrix} = \begin{bmatrix} \left(1 + \frac{s}{\mu} + \rho_1\right) & -1 & 0 & \dots & 0 \\ \gamma_1/\mu & 0 & 1 & & \cdot \\ \vdots & \vdots & \ddots & \ddots & 0 \\ \cdot & \cdot & & 0 & 1 \\ \gamma_R/\mu & \dots & \dots & 0 & 0 \end{bmatrix} \begin{bmatrix} P_r(s) \\ Q_{1,r}(s) \\ \vdots \\ Q_{R,r}(s) \end{bmatrix} \quad (10)$$

The recursion in (10) started at  $r = 0$  clearly yields (8).

(ii) (9) follows from (7) and  $T_r^*(s) = U_r^*(s, \underline{1})$ . (Q. E. D.)

For  $r = 1, 2$  and  $3$ , we show  $U_r^*(s, \underline{z})$  below.

$$U_1^*(s, \underline{z}) = \frac{1 - \rho}{1 + \frac{s}{\mu} - \sum_{k=1}^R \rho_k z_k}$$

$$U_2^*(s, \underline{z}) = \frac{1 - \rho}{\left(1 + \frac{s}{\mu}\right)^2 + \frac{s}{\mu} \rho_1 - \sum_{k=1}^R \left(\rho_k + \frac{s}{\mu} \gamma_k\right) z_k}$$

$$\begin{aligned} U_3^*(s, \underline{z}) &= (1-\rho) / \left\{ \left(1 + \frac{s}{\mu}\right)^3 + 2\left(\frac{s}{\mu}\right)^2 \rho_1 + \frac{s}{\mu} (\rho_2 + 2\rho_1 + \rho_1^2) - \sum_{i=1}^{R-1} \left\{ \frac{\gamma_i}{\mu} \left[ \left(1 + \frac{s}{\mu}\right)^2 + \frac{s}{\mu} \rho_1 \right] + \left[ \rho_{i+1} + \frac{\gamma_{i+1}}{\mu} \frac{s}{\mu} \right] \right\} z_i \right. \\ &\quad \left. - \frac{\gamma_R}{\mu} \left[ \left(1 + \frac{s}{\mu}\right)^2 + \frac{s}{\mu} \rho_1 \right] z_R \right\} \end{aligned}$$

From the above, we obtain  $T_r^*(s)$  for  $r = 1, 2$  and  $3$  by letting  $\underline{z} = \underline{1}$  in  $U_r^*(s, \underline{z})$ .

$$T_1^*(s) = \frac{1 - \rho}{\left(1 + \frac{s}{\mu}\right) - \rho}$$

$$T_2^*(s) = \frac{1 - \rho}{\left(1 + \frac{s}{\mu}\right)^2 - \rho}$$

$$T_3^*(s) = \frac{1 - \rho}{(1 + \frac{s}{\mu})^3 + \rho_1 (\frac{s}{\mu})^2 - \rho}$$

We note that the solutions for  $U_r^*(s, \underline{z})$  and  $T_r^*(s)$  become quite complex if one tries to solve for  $P_r(s)$  and  $Q_{k,r}(s)$  explicitly using the matrix equation (8) when  $r \geq 4$ . In what follows, we turn our attention to finding the moments of  $t_r$ . To do so, we need the following result concerning the distribution of  $\underline{m}^{(r)}$ .

**Theorem 2.** For any  $r \geq 0$ ,  $\underline{m}^{(r)}$  and  $\underline{n}$  have the same stationary distribution.

That is

$$U_r^*(0, \underline{z}) = E[z_1^{m_1^{(r)}} z_2^{m_2^{(r)}} \dots z_R^{m_R^{(r)}}] = P^*(\underline{z}) \quad (11)$$

**Proof.** By (5), (11) holds true for  $r = 0$ . Assume that (11) holds true for some  $r$  so that

$U_r^*(0, \underline{z}) = P^*(\underline{z})$ . By (3), (6) and the induction hypothesis,

$$\begin{aligned} U_{r+1}^*(0, \underline{z}) &= y_1(0, \underline{z}) \cdot \frac{1 - \rho}{1 - \sum_{k=1}^R \rho_k y_k(0, \underline{z})} = \frac{1 - \rho}{y_1(0, \underline{z}) - (\rho_1 + \sum_{k=1}^{R-1} \rho_{k+1} z_k)} \\ &= \frac{1 - \rho}{1 + \sum_{i=1}^R \frac{\gamma_i}{\mu} (1 - z_i) - \rho_1 - \sum_{k=1}^{R-1} \rho_{k+1} z_k} = \frac{1 - \rho}{1 - \sum_{k=1}^R \rho_k z_k} \end{aligned}$$

which is  $P^*(\underline{z})$ . The last equality is obtained using the following relationships:

$$\rho_1 = \frac{\lambda_1}{\mu} = \sum_{i=1}^R \frac{\gamma_i}{\mu} \quad \text{and} \quad \rho_k = \frac{\lambda_k}{\mu} + \rho_{k+1} \quad \text{for } 1 \leq k \leq R-1. \quad (\text{Q. E. D.})$$

The moments of  $t_r$  can be obtained from the moment generating function of  $t_r$  and  $\underline{m}^{(r)}$  as follows.

$$\begin{aligned} E[t_r^n] &= (-1)^n \frac{\partial^n}{\partial s^n} U_r^*(s, \underline{z}) \Big|_{s=0, \underline{z}=\underline{1}} \\ &= (-1)^n \frac{\partial^n}{\partial s^n} U_r^*(s, z, z, \dots, z) \Big|_{s=0, z=1} \end{aligned} \quad (12)$$

**Theorem 3.** The conditional mean response time is

$$E[t_r] = \frac{r/\mu}{1 - \rho} \quad (13)$$

The above theorem is proved by first expressing  $E[t_{r+1}]$  in terms of  $E[t_r]$  using (6), (11) and (12). (13) is then obtained by induction starting with  $E[t_0] = 0$ . (See [11] for details of Proof.)

**Theorem 4.** The second order statistics of the conditional response time can be found recursively using

$$\text{Var}(t_{r+1}) = \text{Var}(t_r) + \frac{1 - 2\rho r}{\mu^2(1-\rho)^2} + \frac{2}{\mu} E[t_r M^{(r)}] \quad (14)$$

$$E[t_{r+1} M^{(r+1)}] = \sum_{i=1}^R E[t_{r+1} m_i^{(r+1)}] \quad (15)$$

and

$$E[t_{r+1} m_i^{(r+1)}] = \frac{2\rho_i}{\mu(1-\rho)^2} + \frac{r\gamma_i}{\mu^2(1-\rho)} + \frac{\gamma_i}{\mu} E[t_r M^{(r)}] + E[t_r m_{i+1}^{(r)}] \quad 1 \leq i \leq R \quad (16)$$

where  $\text{Var}(t_r)$  is the variance of  $t_r$  and  $E[t_r m_{R+1}^{(r)}]$  is zero, with the initial condition

$$\text{Var}(t_0) = 0$$

$$E[t_0 m_i^{(0)}] = 0 \quad \text{for } 1 \leq i \leq R.$$

The above theorem is proved by taking derivatives of (6), using the moment generating properties of transforms; (11) and (13) are used to simplify the resulting expressions. (See [11] for details of proof.)

### 3. DISCUSSIONS AND NUMERICAL EXAMPLES

The conditional mean response time result in Theorem 3 is analogous to results from analyses of a processor-sharing queue [2,3]. The mean response time

$$E[t_r] = \frac{r/\mu}{1-\rho}$$

of a type  $r$  job varies linearly as its (expected) service requirement  $r/\mu$ .

The contribution of this paper is the derivation of higher order statistics for the response times of different types of jobs; also the service requirements (in number of rounds of service) of each type of jobs can have a general probability distribution.

By assuming that each round of service is exponentially distributed, the multi-class feedback queue considered is an open queueing network satisfying local balance. Each type of jobs corresponds to customers following a fixed path. The key idea in our solution approach is to develop a recursive relationship between the response time of a path and the response time of the same path extended by one more transition.

To illustrate our results, we apply the recursive algorithm in Theorem 4 to calculate the standard deviation of  $t_r$  for the following two examples.

Example 1. The service requirements of customers have a small coefficient of variation. The probability of a customer requiring  $r$  rounds of service is

$$a_r = \begin{cases} 1/3 & r = 19, 20, 21 \\ 0 & \text{otherwise} \end{cases}$$

Example 2. The service requirements of customers have a large coefficient of variation. The probability of a customer requiring  $r$  rounds of service is

$$a_r = \begin{cases} 80/99 & r = 1 \\ 19/99 & r = 100 \\ 0 & \text{otherwise} \end{cases}$$

The standard deviation of  $t_r$  is plotted versus  $\rho$  in Figures 2 and 3 for Examples 1 and 2 respectively, for different values of  $r$ .

For comparison, we have plotted two additional curves in Figures 2 and 3. One is the standard deviation of the response time versus  $\rho$  of an arbitrary job for our feedback queue (a round-robin system). The other is the standard deviation of the response time of an arbitrary job in a FCFS system (no feedback; a customer requiring  $r$  rounds of service gets all of them at the same time).

In both examples, the FCFS system gives rise to a smaller standard deviation for the response time of an arbitrary customer than the round-robin system.

In Figure 2, note that all customers require  $r=19, 20$  or  $21$  rounds of service. (The  $r=1$  and  $r=10$  curves correspond to no customers.) Therefore, FCFS gives rise to a smaller standard deviation for the response time of all customers than round-robin.

In Figure 3, the standard deviation of  $t_r$  for small values of  $r$  is smaller than the FCFS standard deviation at the same  $\rho$ . The exact crossover point depends upon  $\rho$ .

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Fig. 2. Standard deviation of response time versus  $\rho$  for service requirements with a small coefficient of variation.

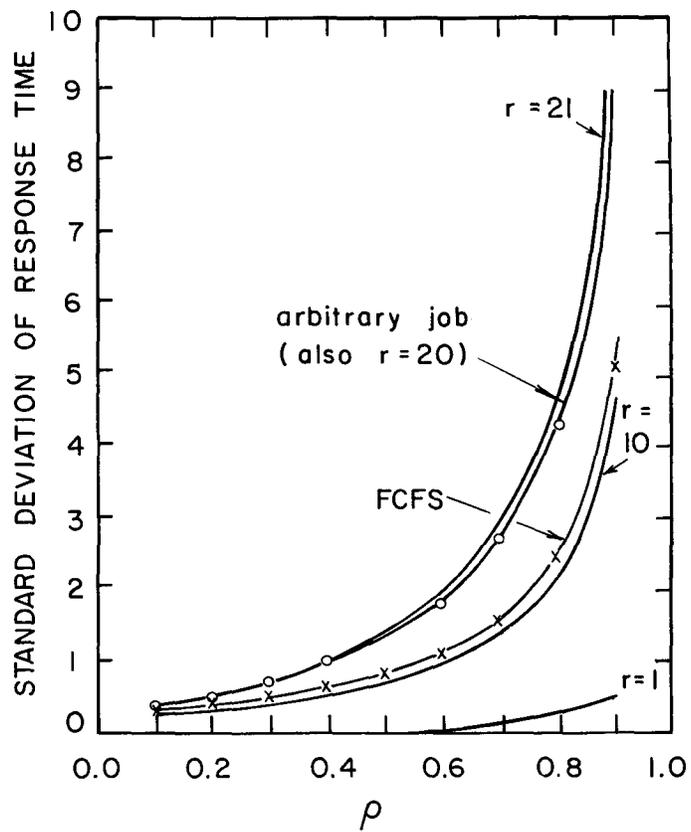


Fig. 3. Standard deviation of response time versus  $\rho$  for service requirements with a large coefficient of variation.

