

Corrigenda

Dong-Young Lee and Simon S. Lam, “Efficient and Accurate Protocols for Distributed Delaunay Triangulation under Churn,” in *Proceedings of IEEE ICNP*, October 2008.

1. The pseudocode of node u when it has received a NEIGHBOR_SET REQUEST from w is the following instead of the pseudocode shown in Figure 3 for the join protocol and in Figure 6 for the maintenance protocol.

On u 's receiving NEIGHBOR_SET REQUEST from w

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if  $w \notin C_u$  then
   $C_u \leftarrow C_u \cup \{w\}$ 
   $N_u \leftarrow$  neighbor nodes of  $u$  in  $DT(C_u)$ 
end if
 $N_w^u \leftarrow \{e \mid e \text{ is a neighbor of } w \text{ in } DT(C_u)\}$ 
Send( $w$ , NEIGHBOR_SET_REPLY( $N_w^u$ ))
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2. The proof of Theorem 2 requires the following unstated assumption: The joining node n is inside the convex hull of the set S of existing nodes. If this assumption is not satisfied, then the joining node should send neighbor-set requests to all new neighbors it discovers instead of just one neighbor in each simplex including n in $DT(C_n)$. That is, the join protocol in technical report TR-06-48 (Department of Computer Science, The University of Texas at Austin, October 2006) should be used.