

## Insertion Sort

The previous example is an **insertion sort**, which is a simple sorting algorithm that requires  $O(n^2)$  comparisons to sort an array of  $N$  elements. How do we know that  $T(n)$  is  $O(n^2)$ ?

The algorithm starts with the pointer at position  $a_1$  in the array, and makes 1 comparison (with  $a_0$ ). Then the pointer moves to position  $a_2$  and makes 2 comparisons, and so on to  $a_n$ . So the number of comparisons is:

$$\sum_{i=1}^{n-1} i = 1 + 2 + 3 + \dots + n - 1$$

What is the *closed form* for this arithmetic series? Well, Gauss observed that you can add this series to itself, using the sum in reverse, and divide by 2 to get the closed form.

$$\begin{array}{r} 1 + 2 + 3 + \dots + n - 1 \\ n - 1 + n - 2 + n - 3 + \dots + n - n - 1 \\ \hline n + n + n + \dots + n \end{array}$$

Where there are  $n-1$  terms. Since we added the same sum twice, we must divide this expression by 2, giving:

$$\sum_{i=1}^{n-1} i = \frac{n(n-1)}{2} = \frac{n^2 - n}{2} \cong O\left(\frac{1}{2}n^2 - \frac{1}{2}n\right) = O(n^2)$$