Principles of Computer Systems

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POX

How do we construct systems that are

🗆 reliable

D portable

 \square efficient

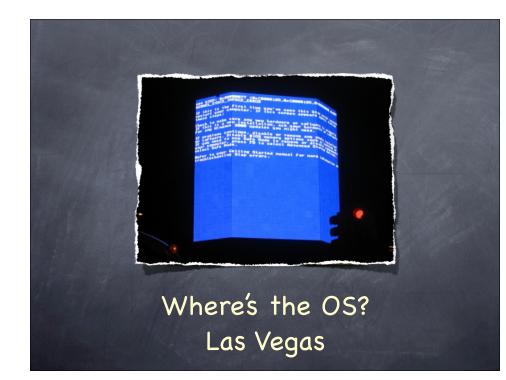


Meet the OS

- Software that manages a computer's resources
 - makes it easier to write the applications you want to write
 - makes you want to use the applications you wrote by running them efficiently

Why study Operating Systems?

- ◎ To learn how computers work
- To learn how to manage complexity through appropriate abstractions
 - □ infinite CPU, infinite memory, files, semaphores, etc.
- @ To learn about system design
 - □ performance vs. simplicity, HW vs. SW, etc.
- Because OSs are everywhere!

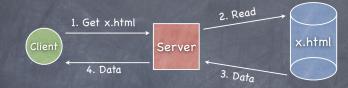




Why study Operating Systems?

Because you are worth it!

Running a Web Server



How does the OS

- allow multiple applications to communicate with each other?
- handle multiple concurrent requests?
- □ support access to shared data (such as the cache)?
- protect against malicious scripts?
- enable different apps to share the data they have produced?
- support consistent changes to complex data structures?
- handle clients and servers of different speed?
- transparently move to more powerful hardware?

Three steps to transmitting POX

- 1. How to approach problems
 - 🗆 fundamental issues
 - ▷ coordination, abstraction
 - 🗆 design space
 - 🗆 case studies
- Goal: Forever mutate your brain (Mwahahahahahahahahahaha)
- © Timescale: Big, long-term payoff

Three steps to transmitting POX

- 2. How to apply specific techniques
 - D Time-tested solutions to hard problems
 - □ Hacking will not succeed
 - ▷ concurrent programming, transactions, etc
- Goal: Be a good engineer (Mwahahaahhaha!)

Three steps to transmitting POX

3. How, in detail, current OSs work

- FS, network stack, internal data structures, VM... of
 - ▶ MacOS, Linux, iOS, Windows
- @ Goal: Well...now in detail how current OSs work!
- Timescale: Better be now, because all will change tomorrow

What is an OS?

An Operating System implements a virtual machine whose interface is more convenient* that the raw hardware interface



* easier to use, simpler to code, more reliable, more secure...

More than one hat

@ Referee

@ Illusionist

@ Glue

More than one hat

@ Referee

Manages shared resources such as CPU, memory, disks, networks, displays, cameras, etc.

Illusionist

□ Look! Infinite memory! Your own private processor!

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Offers a set of common services (e.g. U.I. routines)
 Separates apps from I/O devices

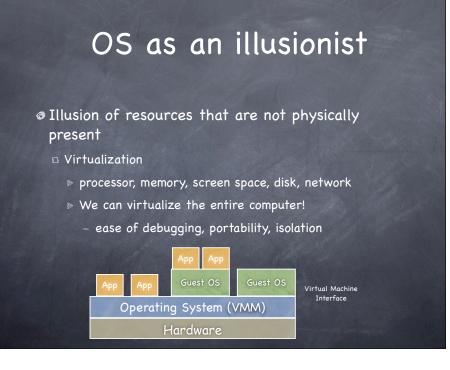
OS as a referee

Resource allocation

When multiple concurrent tasks, how does OS decide who gets how much?

Isolation

- $\hfill\square$ A faulty app should not disrupt other apps or OS
 - OS must export less than full power of underlying hardware
- Communication
 - Apps need to coordinate and share state
 - Web site: select ads, cache recent data, fetch/merge data from disk, etc



OS as an illusionist

- Illusion of resources that are not physically present
 - □ Atomic operations
 - ▶ hardware guarantees atomicity at the word level
 - what happens during concurrent updates to complex data structures?
 - what if computer crashes during a block write?

OS as a glue

- Offers standard services to simplify app design and facilitate sharing
 - □ send/receive of byte streams
 - 🗆 read/write files
 - 🗆 pass messages
 - □ share memory
- Decouples hardware and app development
 - ...but database may need to be aware of specific disk drive

What makes a good OS?

Reliability

- OS does exactly what is designed to do
- Security
 - OS cannot be compromised by a malicious attacker
- Portability
 - Ø OS does not change as hardware changes
- Performance
 - @ efficiency, overhead, fairness, latency, throughput, predictability
- Adoption
 - Are applications ported to the OS widely available?
 - Is hardware supported by the OS widely available?

Reliability

- The ability of a computer-related hardware or software component to consistently perform according to its specifications.
- In theory, a reliable product is totally free of technical errors (yeah, right)
- Availability: percentage of time system is useful
 - Depends on MTTF and MTTR

Security

- Includes privacy: data on the computer only accessible to authorized users
- Strong fault isolation helps, but not enough
 - D Email gives no strong assurance of sender's identity
 - Security mechanisms should not prevent legitimate sharing!

Enforcement mechanism

- 🗅 Ensures only permitted actions are allowed
- Security policy
 - Defines what is permitted

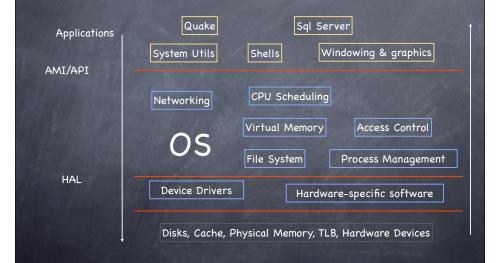
Portability

OSs can live more than your cat!
 must support applications not yet written
 must run on hardware not yet developed

Three interfaces

- Abstract Machine Interface (AMI)
 - ▷ between OS and apps: API + memory access model + legally executable instructions
- Application Programming Interface (API)
 - ▶ function calls provided to apps
- Hardware Abstraction Layer (HAL)
 - ▶ abstracts hardware internally to the OS

Logical OS Structure



Performance

Efficiency/Overhead

how much is lost by not running on bare hardware?

Fairness

how are resources divided?

Response time

how long does a task take to complete?

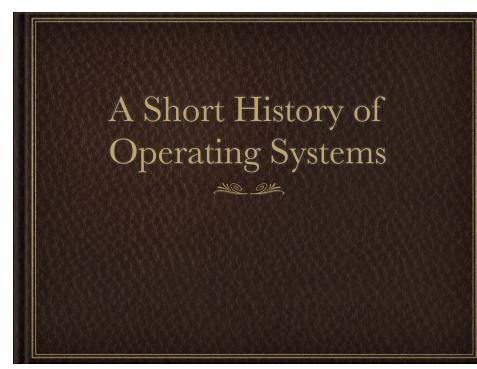
Throughput

how many tasks complete per unit of time?

Predictability

are performance metrics consistent over time?

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HISTORY OF OPERATING SYSTEMS: PHASES

• Phase 1: Hardware is expensive, humans are cheap

- User at console: single-user systems
- Batching systems
- Multi-programming systems

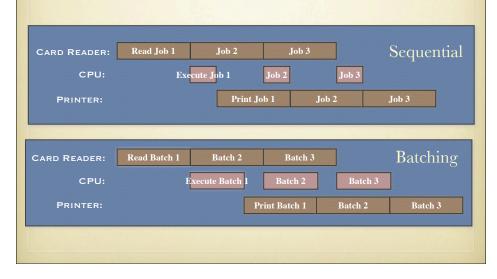
HAND PROGRAMMED MACHINES (1945-1955)

- Single user systems
- OS = loader + libraries of common subroutines
- Problem: low utilization of expensive components

 $\frac{\text{time device busy}}{\text{observation interval}} = \% \text{ utilization}$



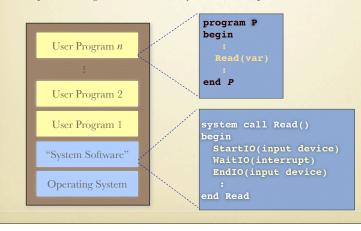
BATCH/OFF-LINE PROCESSING (1955-1965)



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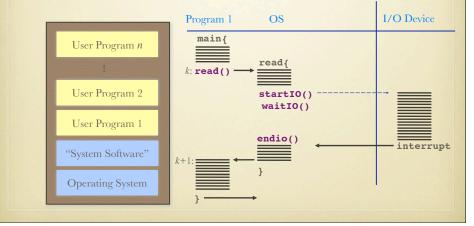
MULTIPROGRAMMING (1965-1980)

Keep several jobs in memory and multiplex CPU between jobs



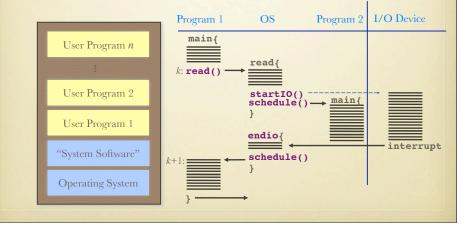
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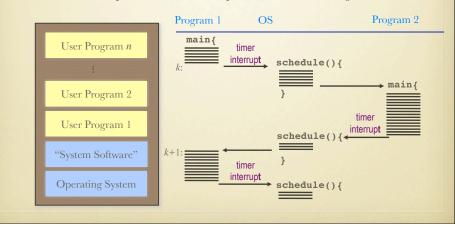


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 - Time sharing: Users use cheap terminals and share servers

TIMESHARING (1970-)

A timer interrupt is used to multiplex CPU between jobs



HISTORY OF OPERATING SYSTEMS: PHASES

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- Phase 2: Hardware is cheap, humans are expensive
 - Time sharing: Users use cheap terminals and share servers
- Phase 3: Hardware is very cheap, humans are very expensive
 - Personal computing: One system per user
 - Distributed computing: many systems per user
 - Ubiquitous computing: LOTS of systems per users

OPERATING SYSTEMS FOR PCS

Personal computing systems

□ Single user

- Utilization is no longer a concern
- Emphasis is on user interface and API
- □ Many services & features not present

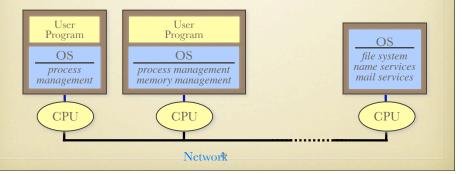
Evolution

- □ Initially: OS as a simple service provider (simple libraries)
- □ Now: Multi-application systems with support for coordination



DISTRIBUTED OPERATING SYSTEMS

- Abstraction: present a multi-processor system as a single processor one.
- New challenges in consistency, reliability, resource management, performance, etc.
- Examples: SANs, Oracle Parallel Server



UBIQUITOUS COMPUTING

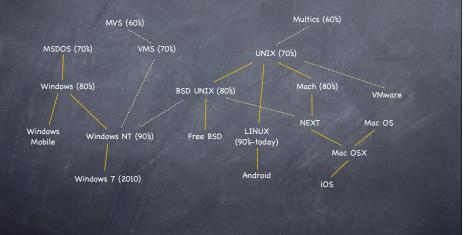
- PDAs, cellular phones, sensors
- Challenges
- □ Small memory size
- □ Slow processor
- □ Battery concerns
- □ Scale
- □ Security
- □ Naming



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- Richer Services
 - Real-time operating systems

Genealogy of modern Operating Systems



Cambia, Todo Cambia

Nothing wrong with batch systems

□ They were right for tradeoffs at the time

But tradeoffs change...

	1981	1996	2011	Factor
MIPS	1	300	10000	10K
\$/MIPS	\$100K	\$30	\$0.50	200K
DRAM	128KB	128MB	10GB	100K
Disk	10MB	4GB	1TB	100K
Home Internet	9.6Kbps	256 Kbps	5Mbps	500
LAN Network	3Mbps (shared)	10 Mbps	1Gbps	300
# Users	100	100 Mb/s	<<1	100+