Advanced Topics

Last Time

• Experimental Methodology

Today

- What's a managed language?
- Alias Analysis dealing with pointers
 - Focus on statically typed managed languages
 - Method invocation resolution

Alias Analysis

What is pointer analysis?

- For umanaged languages like C, C++, and managed languages like C#, Java, Modula-3
- Goal: given two memory references, *m* and *n*, and a program point, *p*, can *m* and *n* access the same memory location?

Pointer analysis is an important problem

- Compiler optimizations require alias information
- Error detection
- Program slicing/debugging
- Data structure *shape* analysis: is it a tree?
- Object inlining, object colocation, synchronization elimination, etc.

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Pointer analysis is hard

- Undecidable in general [Landi 92]
- Many approximations
 - Near linear, but imprecise [Steensgaard 97]
 - Double exponential shape analysis [Sagiv 98]
- Why haven't we solved this problem? [Hind 01]
- \Rightarrow No one algorithm satisfies all client needs

Sample Client driven results:

- Error detection: needs high precision, but a demand driven precision can provide precision at low cost [Guyer/Lin 03]
- Scalar optimizations: do not need high precision [Chowdhury et al. 03, Diwan et al. 98/01]
- Object colocation: a fast, simple, and *unsound* analysis is effective [Guyer/McKinley 04]

Taxonomy

Interprocedural data-flow analysis

- Interprocedural or intraprocedural? Do we compute the effects of calls?
- Context-sensitive or insensitive? Do we take into account the calling context?
- Flow-sensitive or insensitive? Do we take into account control flow

More precise means more expensive.

- Address-taken: Flow and context insensitive. Intraprocedural or interprocedural. Every variable that may be aliased goes in a single set. Single pass.
- **Steensgaard:** Flow and context insensitive. Interprocedural. Uses union-find to create sets of aliases. Single pass.
- Anderson: Flow and context insensitive. Interprocedural. Creates sets of aliases. Iterative on the CFG.
- **Burke:** Flow-insensitive and context sensitive. Interprocedural. Iterative on the CFG and Call graph.
- **Choi et al.:** Flow and context sensitive. Interprocedural. Iterative on the CFG and Call graph.

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Alias Analysis

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Example: Simple Alias Analysis

Language context

- Modula-3: a statically-typed, object-oriented programming language.
- Techniques apply directly to Java and C#, and may apply to C++.

Clients:	scalar	optimizations	and	method	resolution

Alias analysis: Type-based alias analysis

Using types is

• Fast and surprisingly effective for these clients

Evaluation: A limit analysis computes a runtime upper bound on effectiveness of these compiler optimizations.

A. Diwan, K. S. McKinley, and J. E. B. Moss, Using Types to Analyze and Optimize Object-Oriented Programs, ACM Transactions on Programming Languages and Systems, 23(1): 30-72, January 2001.

-anguages and Systems, 25(1). 50-72, January 2001.					
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Type-Based Alias Analysis

Use language types to disambiguate memory references



Strengths and Weakness of Analysis

	Types	Instructions
Speed	O(Instructions)	O(Instructions ^d)
Unsafe programs	No	Yes
Precision	Good	Better

Type-based alias analysis is fast. Is it effective?

Evaluation

	Static	Run time	Upper-bound
Coverage	\checkmark		
Run-time impact		\checkmark	
Compare analyses	\checkmark	\checkmark	~
Compare to best			\checkmark
Deficiencies			\checkmark

Lots of alias analysis work uses static metrics and does not consider the clients.

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Client 1: Redundant Load Elimination (RLE)

Modula-3 Benchmarks

Name	Lines	Instructions	% Heap loads	% Other loads
format	395	1,879,195	10	17
dformat	602	1,442,541	9	19
write-pickle	654	1,614,437	13	16
k-tree	726	50,297,517	10	21
slisp	1,645	11,462,791	27	9
рр	2,328	45,779,402	11	19
m2tom3	10,574	50,894,990	8	28
m3cg	16,475	5,636,004	8	21

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Static Evaluation

Average number of potential aliases for each reference.

	Declarations	Fields	Statements
Interprocedural	4.7	3.4	3.4
			[0.3 to 21]
Intraprocedural	54.1	12.7	12.7
			[2 to 28]

- Bad news!
- Much higher accuracy with interprocedural analysis.
- Flow-insensitive data-flow analysis contributes very little to the precision of a type-based alias analysis.

Run-Time Impact of Alias Analysis

Measure the performance improvement due to RLE.



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Comparing Alias Analysis to an Upper-bound

What's Left to Disambiguate?



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Bottom Line

- All metrics are useful.
- In this case, static measures are misleading.
- Alias analysis & RLE yield modest improvements for our benchmarks.
- There is not much room for improvement in the alias analysis for RLE.
- RLE acts locally.

Client 2: Resolving Method Invocations



- A polymorphic method invocation may call more than one procedure at run time.
- A monomorphic method invocation always calls the same procedure at run time.

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Resolving Method Invocations

Method invocations

- are more expensive than direct calls
- inhibit compiler optimizations

We want to resolve monomorphic method invocations: replacing method invocations with direct calls.

Full resolution is undecidable so any resolution technique must be conservative and assume polymorphic.

Technique	Complexity	Source
Type hierarchy	O(Types * Methods)	borrowed
(Fernandez, Dean et al.)		
Intraprocedural		
type propagation	$O(\sum_p n_p * v_p)$	simplified
Aggregate	$O(\sum_{p} n_p * v_p)$	new
Interprocedural		
type propagation	$O(p\sum_p n_p * v_p)$	simplified

Type Hierarchy Analysis

Bounds the procedures a method invocation may call by examining the type hierarchy declarations for method overrides.



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Type Hierarchy Analysis



Declarations set types. Analysis ignores control flow.

Intraprocedural Type Propagation: Data-Flow Analysis



Allocation, type discrimination, and assignment affect types

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Aggregate Analysis

Use type-based alias analysis to disambiguate pointer references. Type-based alias analysis

- merges all instances of an object or record type, and
- ignores all control flow.

Aggregate analysis finds monomorphic uses of general data structures.

Aggregate Analysis Example



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Interprocedural Type Propagation

- Merge parameter information for all callers (context insensitive)
- Merge return information for all callers at a call site
- Reanalyze when more refined information becomes available
- Update call graph as necessary



Modula-3 Benchmarks

		Static	Metho	d Inv.	
Name	Lines	Total	Mono	Rslv	Dyn MI
format	395	37	26	26	47,064
dformat	602	95	77	77	30,775
k-tree	726	13	5	4	714,619
slisp	1645	223	94	94	67,253
рр	2328	24	1	1	458
dom	6186	222	136	128	12,377
postcard	8214	293	129	87	3,076
m2tom3	10574	1821	1110	1110	19,886,862
m3cg	16475	1808	295	213	32,850
trestle	28977	430	22	9	10,756
Total		4966	1895	1749	(92%)

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Effectiveness



Upper-Bound Comparison

Compare results with monomorphic method invocations at run time



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Cause Evaluation

Find the source of loss of type information



Source	Solution
Control merge	Context sensitive analysis
Data merge	More powerful alias analysis
Unavailable	Analyze libraries

Cause of analysis failure



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Run-Time Impact of Resolving Method Invocations

Bottom Line

- Method resolution improves performance up to 11%.
- Inlining & method resolution improves performance up to 28% for these benchmarks.
- Different analyses are effective under different situations; thus we need a range of analyses.
- There is not much room for improvement here.
- \Longrightarrow We can eliminate method invocation overhead for monomorphic call sites.

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Cumulative Impact of Optimizations



Summary

Type-based analyses are effective at

- reducing the overhead of linked structures, and
- reducing the overhead of method invocations,

Compiler optimizations need static, dynamic, and limit evaluation.

Open Questions

- Further evaluation of type-based alias analysis.
- Applying limit analysis to other optimizations.

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Next Time

Interprocedural optimization

Grove & Torczon, Interprocedural Constant Propagation: A Study of Jump Function Implementations, ACM Conference on Programming Language Design and Implementation, pp. 90-99, Albuquerque NM, June 1993.

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