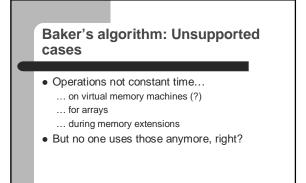


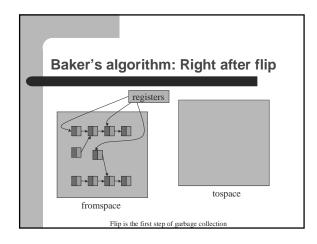
Baker's algorithm: How it's different

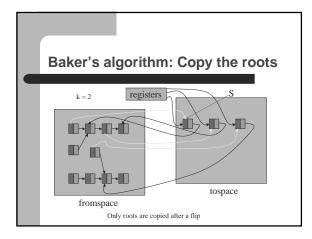
- When *tospace* fills up, do a flip and copy only roots
- Every *cons* does a few iterations of garbage collection
- Each *car* and *cdr* checks for a forwarding address updates pointer if found

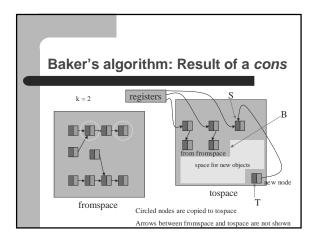
Baker's algorithm: Proven guarantees

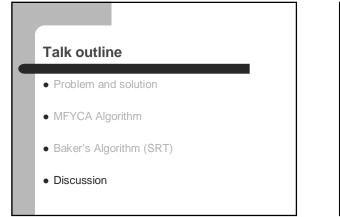
- Serial Real-Time system
 - one thread
 - all operations O(1) time
- Won't run out of *tospace* (unless we really run out of space)
 - (1 + 1/k) times as much storage needed as for MFYCA
 - Tradeoff between time and space

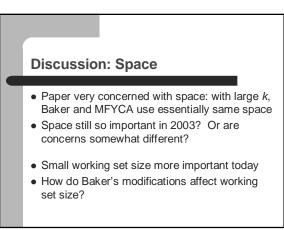












Discussion: Breadth-first order

- The "graph" of objects is traversed in breadthfirst order
 - True for both MFYCA and Baker
- Why?
- Is this beneficial? Consider locality.

Discussion: T pointer • Baker's algorithm adds T pointer to MFYCA • New objects allocated at the end of the free space; garbage-collected accessible objects copied to beginning of free space • Maximum Conservation of the space • Stammed Index Tree space • Why do this? • Does this have other effects? Could MFYCA have used a T pointer?

Discussion: Read barrier overhead

- The (potential) copying done by each *car* and *cdr* has to be done by MFYCA, too
- However, each *car* and *cdr* checks if the node is in *fromspace* vs. *tospace*
- Is this bad?
- How bad?
- Consider database application (from paper)

Discussion: More operations

- How well does Baker's algorithm apply to a larger set of operations than *cons*, *car*, *cdr*, *rplaca*, *rplacd*, *eq*, and *atom*?
- Consider imperative languages

Conclusion

- Baker's algorithm: Modification to MFYCA
- Contributions:
 - Real-time: all operations constant time
 - Space efficiency and flexibility: can choose k for space-time tradeoff
 - Proof: Correct and doesn't run out of space when it shouldn't