

1. (Compression)

$$\begin{aligned}
\text{(a)} \quad & \text{entropy} \\
& = -1/2 \log(1/2) - 1/8 \log(1/8) - 1/4 \log(1/4) - 1/8 \log(1/8) \\
& = -1/2(-1) - 1/8(-3) - 1/4(-2) - 1/8(-3) \\
& = 1/2 + 3/8 + 2/4 + 3/8 \\
& = 1.75
\end{aligned}$$

(b) Let T have a non-leaf node x whose only son is y ; we show that T is not optimal by displaying another tree T' , which is also a prefix code, whose weight is strictly lower.

If x is the root of T create T' by deleting x and letting y be the root. If x is non-root, change the parent of y to the parent of x , and delete x . In each case, the pathlengths to the descendants of y decrease and other pathlengths do not increase. So, the weight of T' is less than that of T .

(c) Let the weight of the optimal trees for S and R be s and r , respectively. We can get a tree for R from the one for S by expanding the leaf z into a non-leaf with children x and y . The weight of the resulting tree is $s + x + y$. This quantity is at least the weight of the optimal tree for R , r . That is,

$$r \leq s + x + y$$

2. (Powerlist)

(a) Proof is by induction on i .

• $i = 0$: We have to show $u_1 = u_0 \bowtie v_0$ and $v_1 = v_0 \bowtie u_0$.

$$\begin{aligned}
& u_1 \\
& = \{ \text{definition of } u_1 \} \\
& \quad u_0 \mid v_0 \\
& = \{ \text{laws of powerlist: } u_0 \text{ and } v_0 \text{ are singletons} \} \\
& \quad u_0 \bowtie v_0
\end{aligned}$$

• $i > 0$:

$$\begin{aligned}
& u_{i+1} \\
& = \{ \text{definition of } u_{i+1} \} \\
& \quad u_i \mid v_i \\
& = \{ \text{induction on both terms} \} \\
& \quad (u_{i-1} \bowtie v_{i-1}) \mid (v_{i-1} \bowtie u_{i-1}) \\
& = \{ \text{commutativity} \} \\
& \quad (u_{i-1} \mid v_{i-1}) \bowtie (v_{i-1} \mid u_{i-1}) \\
& = \{ \text{definitions of } u_i \text{ and } v_i \} \\
& \quad u_i \bowtie v_i
\end{aligned}$$

$$(b) \quad \begin{aligned} p \sqsubseteq \langle x \rangle &= (p == \langle x \rangle) \\ p \sqsubseteq r \mid s &= (p == r \mid s) \vee (p \sqsubseteq r) \end{aligned}$$

Another possible definition is

$$\begin{aligned} \langle x \rangle \sqsubseteq \langle y \rangle &= (x == y) \\ \langle x \rangle \sqsubseteq r \bowtie s &= (\langle x \rangle \sqsubseteq r) \\ p \bowtie q \sqsubseteq r \bowtie s &= (p \sqsubseteq r) \wedge (q \sqsubseteq s) \end{aligned}$$

3. (String Matching)

(a) Suppose $v[0..k]$ is the core. From the definition of core,

$$\begin{aligned} v[0..k] &= v[20 - k..20]. \text{ Hence,} \\ v[i] &= v[20 - k + i] \end{aligned}$$

Letting $i = 6$ and $20 - k + i = 11$, we get $k = 15$. That is, if $k = 15$, $v[6] = v[11]$. Since $v[6] \neq v[11]$, $v[0..15]$ is not the core.

$$(b) \quad \begin{aligned} &u \preceq v \\ \Rightarrow &\{c(u) \prec u\} \\ &c(u) \prec v \\ \Rightarrow &\{\text{definition of core: } w \prec v \equiv w \preceq c(v)\} \\ &c(u) \preceq c(v) \end{aligned}$$