Open book and notes.

Max points = 75

Time = 75 min

Do all questions.

- 1. (Recursion and Induction; 45 points)
  - (a) (4 points) Define function take where take n xs is a list containing the first n items of list xs in order, where n is a natural number. If n exceeds the length of xs then the entire list xs is returned.
  - (b) (4 points) Define function drop where

$$xs = (take n xs) ++ (drop n xs)$$

Recall: ++ concatenates two lists.

(c) (9 points) A matrix of integers is represented as a list of lists where each component list is a row of the matrix. Thus, the matrix in Table 1

> 1 6

Table 1: Matrix

is represented by [[3,7,4],[2,1,9],[1,2,4],[0,6,2]]. Assume that the matrix is not empty, so no row or column is empty.

Write functions to compute: (1) a list of row sums, [14, 12, 7, 8] for this example, (2) the sum of all the matrix elements, 41 in this case, and (3) a list of column sums, [6, 16, 19] in this case.

You may have to write auxiliary functions. Use map and foldr lib-

(d) (10 points) You are asked to find the  $i^{th}$  value in rank order in a given list of n integers, n > 0. Assume that  $1 \le i \le n$ . For example, given the list [31,34,22,3,19,5,20,6], its sorted version is [3,5,6,19,20,22,31,34]. For i=1, the value is 3, for i=2, the value is 5, for i = 5, the value is 20, etc.

To find the  $i^{th}$  value, you don't need to sort the list. You can use function partition, given in Page 151 in the class notes. Given a list (x:xs), partition xs into lh and rh using x for v in function partition. If the length of lh is greater than or equal to i, look for the  $i^{th}$  value in lh, otherwise, find the suitable value as x or in rh.

Define function rank i xs. Use function partition (you don't have to rewrite it). You can get the length of a list using Haskell function length. Assume that the list elements are distinct.

- (e) (9 points) You are given a list xs of strings. Write a function to produce a list of strings ys, of the same length as xs, where the i<sup>th</sup> string of ys is the concatenation of all strings of xs up to and including the i<sup>th</sup> one. Thus, given that xs is [''this'', ''is'', ''cute''], ys is [''this'', ''thisiscute'']. Use ++ for list (string) concatenation.
- (f) (9 points) You are given a list xs of bit strings, all of the same length. Write a function to output True iff all adjacent pairs of strings have Hamming distance of 1. Consider the first and last strings of xs to be adjacent. Thus, ["00", "01", "11", "10"] should produce True and ["00", "01", "10", "11"] produces False. For input [], output True.

Hint: You may use the function right\_rotate that rotates a list to the right by one position. This function is defined in Page 145 of the notes. Just use the function, don't reproduce it.

## 2. (String Searching; 20 points)

- (a) (Core Computation; 8 points) Let  $Z_i$  be the alternating binary string  $\overbrace{0101\cdots01}^{2i}$  of length 2i (i.e., i pairs of 01). Thus,  $Z_4$  is 01010101. Show that  $c^j(Z_n0) = Z_{n-j}0$ , where c denotes the core function, and  $1 \le j \le n$ . That is,  $c^1(Z_40) = c(010101010) = 0101010 = Z_30$ , and  $c^3(Z_40) = c^2(c(010101010)) = c^2(0101010) = c(c(0101010)) = c(01010) = 010 = Z_10$ .
- (b) (Core Computation; 6 points) Consider string p[0..12] where p[3], p[5] and p[9] are all different symbols. What can you say about the length of the core?
- (c) (Rabin-Karp algorithm; 6 points) We are given a very long genome sequence. This string will be repeatedly searched with patterns of small lengths, say of lengths about 20. How can we apply the Rabin-Karp algorithm effectively? There are several possible approaches; any good one will do. Explain your algorithm in no more than 5 sentences.

## 3. (Relational Algebra; 10 points)

- (a) (5 points) Let the Tables 6.1 (page 170), 6.4 (page 172) and 6.5 (page 173) in your notes be denoted by relations R, S and T, respectively. Write a query to find the names of theatres which are showing PG movies in which Will Smith is acting. You don't have to simplify or compute the value of the query.
- (b) (5 points) Show that for arbitrary relations R, S, and attribute a,  $\pi_a(R \cap S) = \pi_a(R) \cap \pi_a(S)$  does not necessarily hold.