

Exercises on Program Verification, Jayadev Misra

1. (3/3/98) Given a family of sets $in_j, 0 \leq j < N$ and $out_j, 0 \leq j < N$. Set a variable s to *true* iff
- $$(\forall i, j : 0 \leq i < N \wedge 0 \leq j < N \wedge i \neq j : in_i \subseteq out_j).$$

On first glance, a quadratic solution is obvious. We describe a linear solution below.

Invariant::

$$\begin{aligned} cin &= (\cup j : 0 \leq j < k : in_j), \\ cout &= (\cap j : 0 \leq j < k : out_j), \\ s &= (\forall i, j : 0 \leq i < k \wedge 0 \leq j < k \wedge i \neq j : in_i \subseteq out_j) \end{aligned}$$

Program::

```
cin := {}; cout := ALL; k := 0; s := true;
while k ≠ N do
  s := s ∧ (in_k ⊆ cout) ∧ (cin ⊆ out_k);
  cin := cin ∪ in_k; cout := cout ∩ out_k;
  k := k + 1
od
```

Note: It may be simpler to use integer arrays X, Y and compute the predicate $s = (\forall i, j : 0 \leq i < N \wedge 0 \leq j < N \wedge i \neq j : X_i \leq Y_j)$. Then \cup, \cap and can be replaced by *min, max*.

2. (3/3/98) Take Dijkstra's integer square root problem and solve it using binary search. From $s^2 \leq k \wedge (s+1)^2 > k$, use the invariant, $s^2 \leq k \wedge t^2 > k$ with the termination condition $t = s + 1$.