

A proof of Fermat's little theorem

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The following theorem, known as Fermat's little theorem, is a fundamental result in number theory. The theorem has many applications. Pratt [3] uses the theorem to certify that a number is prime. It is used in cryptographic protocols, such as the Diffie-Hellman key exchange [1].

Theorem 1 For any natural number n and prime number p , $n^p - n$ is a multiple of p .

There are several ways to prove this theorem, e.g. using induction on n . A proof using the pigeon-hole principle is as follows. For positive integers i and j , and prime p it can be shown that $i.n \equiv^{\text{mod } p} j.n$ if and only if $i \equiv^{\text{mod } p} j$. Then $\{i.n \bmod p \mid 1 < i < p\} = \{j \mid 1 < j < p\}$. The product of the elements of the sets in this equation are identical, so, $\Pi(\{i.n \mid 1 < i < p\}) \bmod p = \Pi(\{j \mid 1 < j < p\}) \bmod p$, or $n^{p-1} \times (p-1)! \equiv^{\text{mod } p} (p-1)!$. Since prime p does not divide $(p-1)!$, cancel $(p-1)!$ from both sides to get $n^{p-1} \equiv^{\text{mod } p} 1$. This is equivalent to $n^p \equiv^{\text{mod } p} n$, or $n^p - n$ is a multiple of p .

Dijkstra[2] gives a beautiful proof using elementary graph theory. The proof given here is based on Dijkstra's constructions though it does not use graph theory.

Proof of the theorem: Consider the set of words of length p over an alphabet of size n . Define an equivalence relation over the words, x and y are equivalent if and only if x is a rotation of y . We count the number and size of the equivalence classes.

Define q to be a *period* for x if q rotations of x , leftward for positive q and rightward for negative q , yields x . Clearly, 0 is a period for all x , 1 is a period for x if and only if all symbols in x are identical, and given periods q and q' for x , $a \times q + b \times q'$, for arbitrary integers a and b , are also periods for x . In particular, a multiple of a period is a period. A *simple period* is not a multiple of another period. For simple period q for x , all q rotations of x yield distinct words.

Let q be a simple period for a given x . We use Bézout's identity: for integers m and n , there exist integers a and b such that $a \times m + b \times n = \gcd(m, n)$, where \gcd is the greatest common divisor. Setting $m, n = p, q$ in Bézout's identity, $\gcd(p, q)$ is a period. Since p is prime, $\gcd(p, q)$ is either 1 or p , and since q is a

simple period, $q = 1$ or $q = p$. If $q = 1$, x consists of identical symbols. There are n such words so, $q = p$ for the remaining $n^p - n$ words. Therefore, each of these words belongs to an equivalence class of size p ; so, $n^p - n$ is a multiple of p .

Dijkstra's proof The following proof is a rewriting of the proof of Dijkstra [2]. For $n = 0$, $n^p - n$ is 0, hence a multiple of p . For positive integer n , take an alphabet of n symbols and construct a graph as follows: (1) each node of the graph is identified with a word of p symbols, and (2) there is an edge from x to y if rotating word x by one place to the left yields y . Observe:

1. No node is on two simple cycles because every node has a single successor and a single predecessor (which could be itself).
2. Each node is on a cycle of length p because successive p rotations of a word transforms it to itself.
3. Every simple cycle's length is a divisor of p , from (2). Since p is prime, the simple cycles are of length 1 or p .
4. A cycle of length 1 corresponds to a word of identical symbols. So, exactly n distinct nodes occur in cycles of length 1. The remaining $n^p - n$ nodes occur in simple cycles of length p .
5. A simple cycle of length p , from the definition of a simple cycle, has p distinct nodes. From (4), $n^p - n$ is a multiple of p .

References

- [1] W. Diffie and M. Hellman. New directions in cryptography. *IEEE Trans. Inform. Theory*, 22(6):644–654, 1976.
- [2] Edsger W. Dijkstra. A short proof of one of Fermat's theorems. EWD740: circulated privately, May 1980.
- [3] Vaughan R Pratt. Every prime has a succinct certificate. *SIAM Journal on Computing*, 4(3):214–220, 1975.