

A Functional Specification and Validation Model for Networks on Chip in the ACL2 Logic

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* Part of this work was done while visiting the CS department of UT

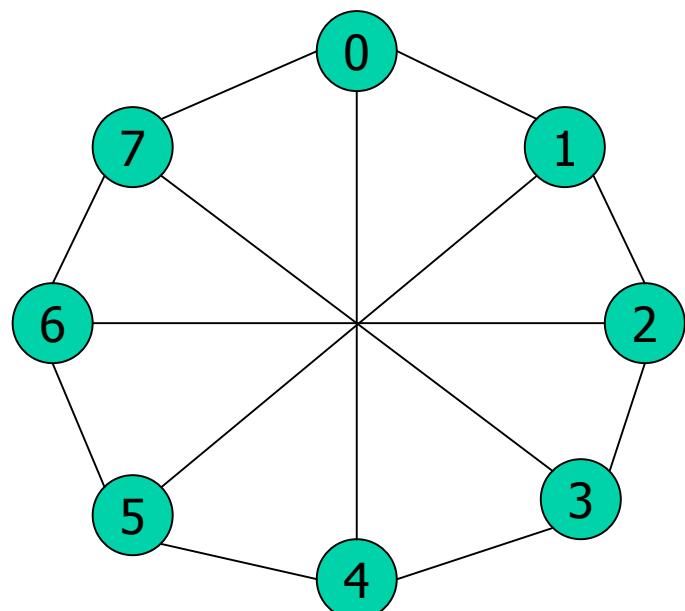


VDS

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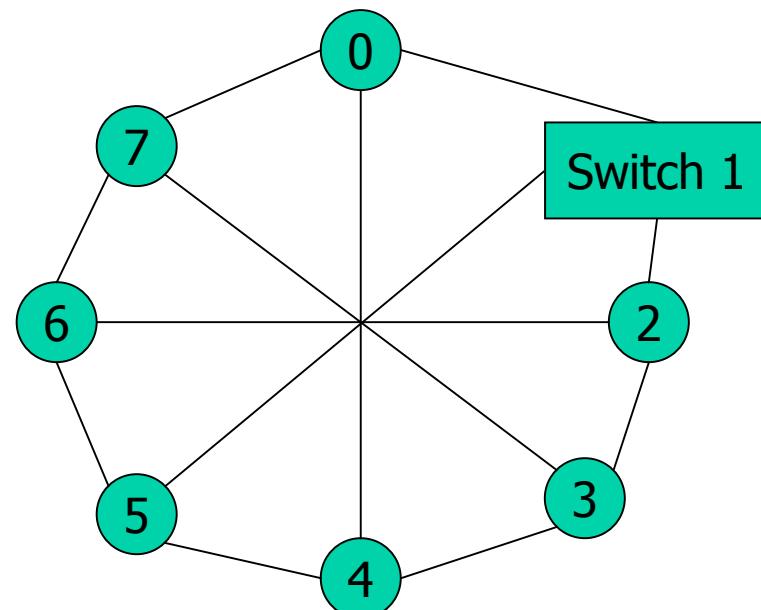
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Initial Motivation: Octagon Network on Chip

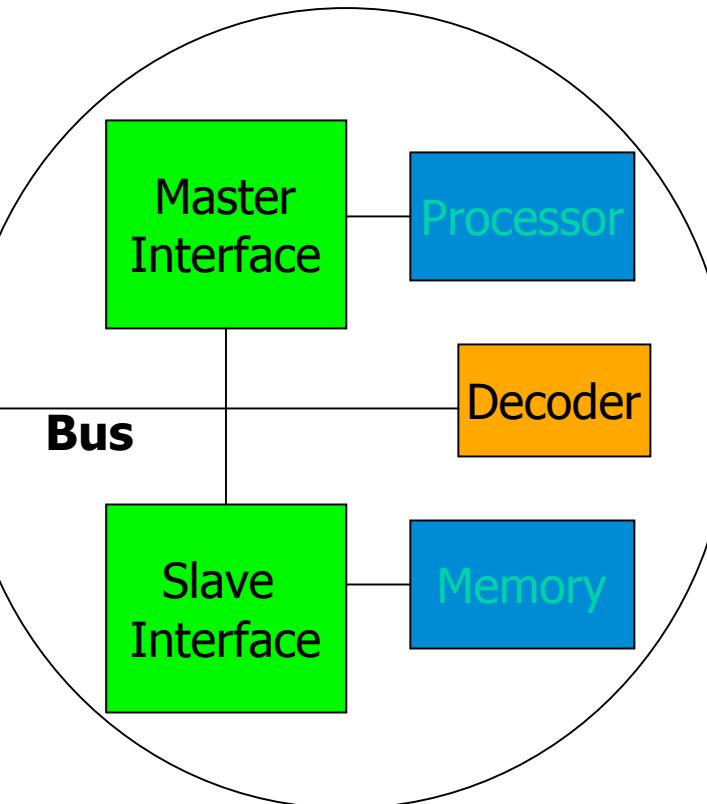


- Design by STmicroelectronics ref: DAC 2001 and IEEE MICRO 2002 by F. Karim *et al.*
- 8 Nodes
- Extensible to $N = 4*i$
- Bidirectional Links
- Simple Shortest Path Routing Algorithm

Octagon System



Parameterized Octagon Unit

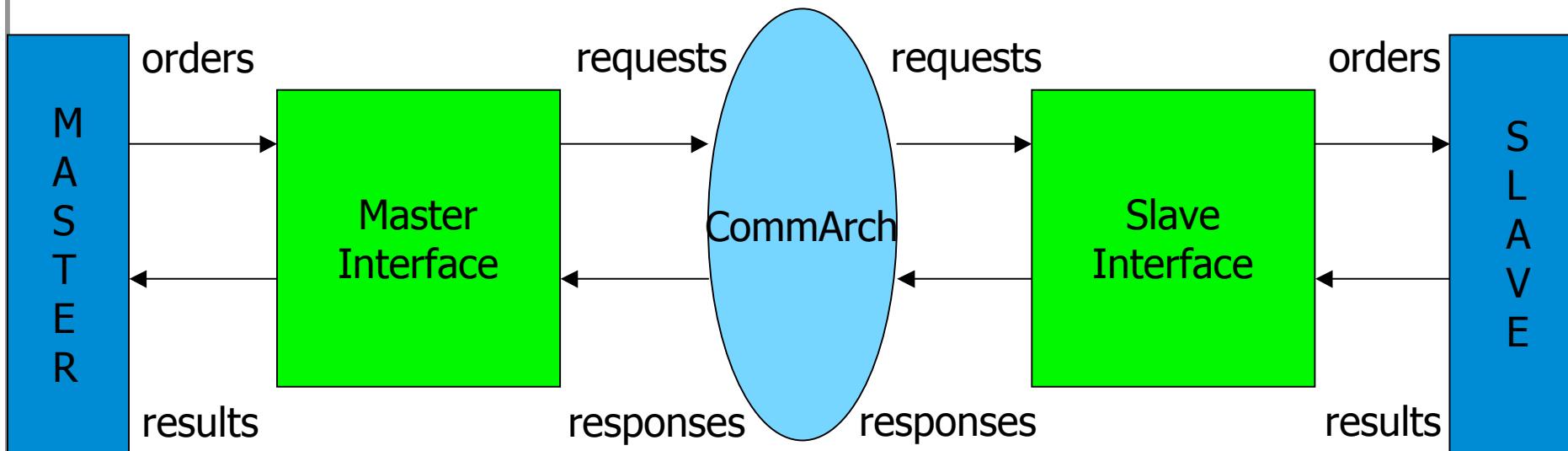


Node System

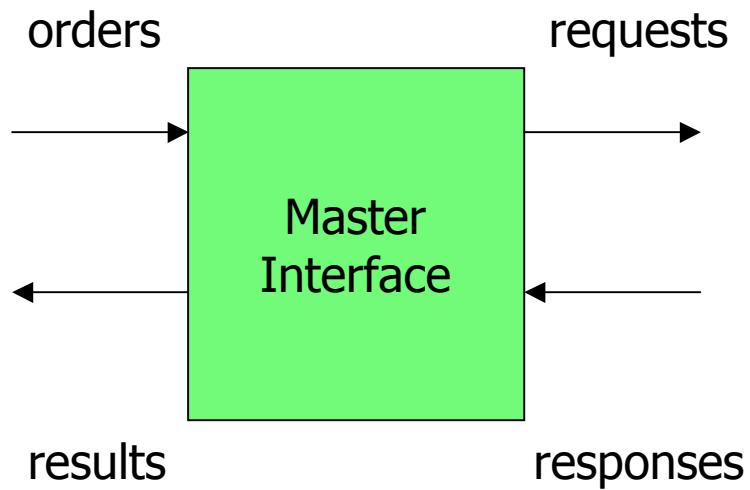
Outline

- Functional modeling of communications
 - Definition of transfers
 - Correctness Criteria
- Informal presentation of the Octagon
- ACL2 model of the Octagon

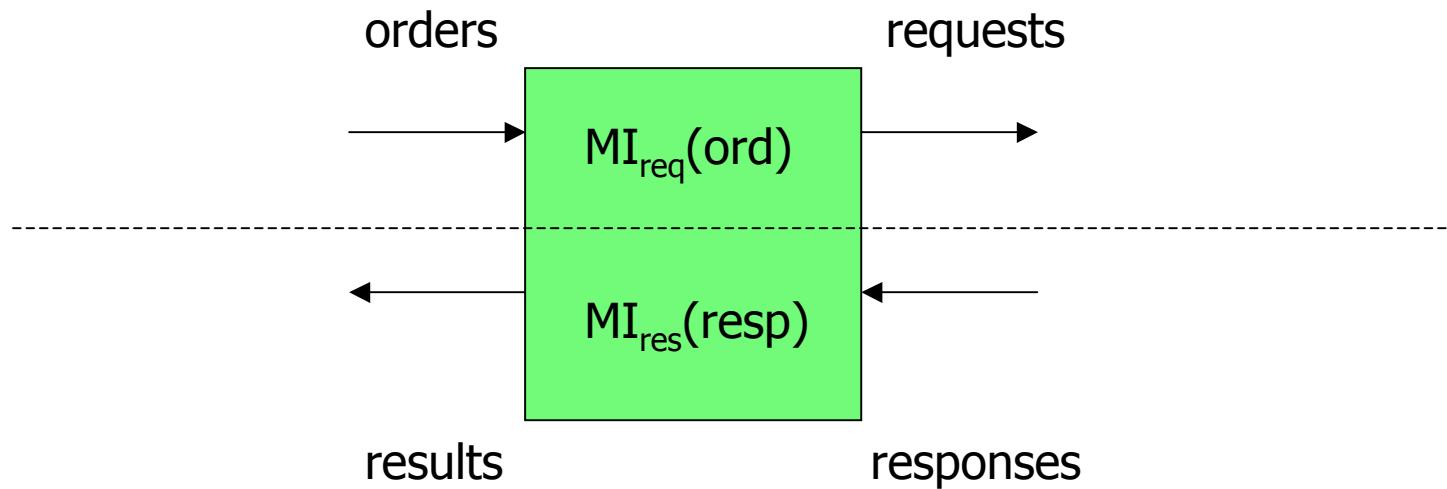
Communication Scheme



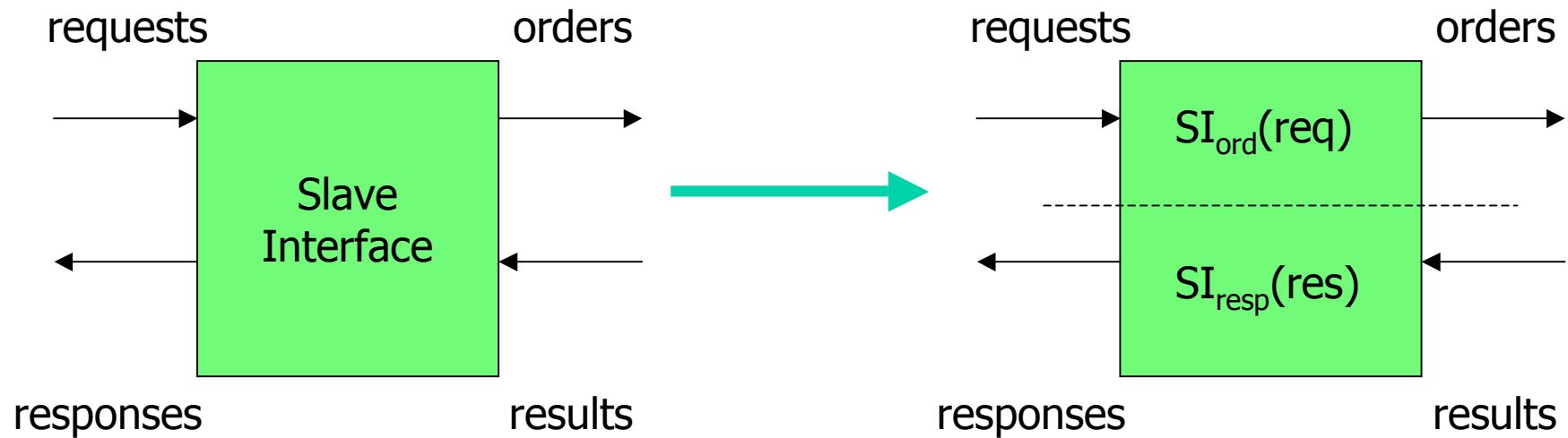
Functional Modeling



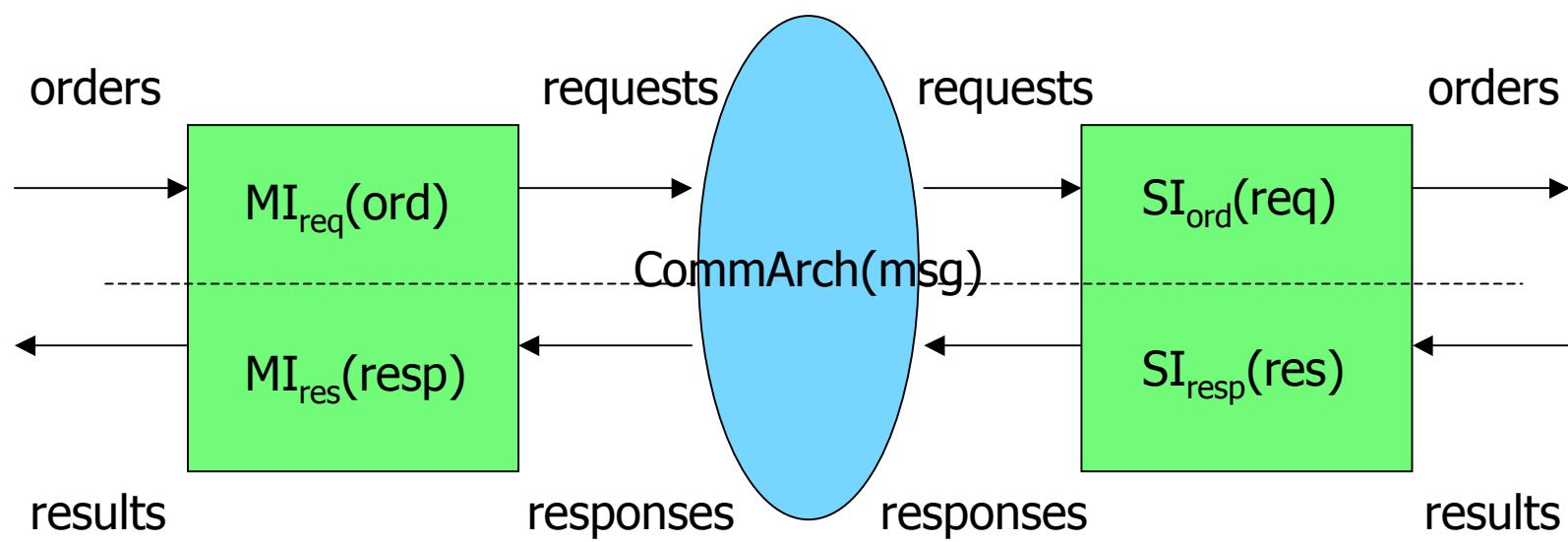
Functional Modeling



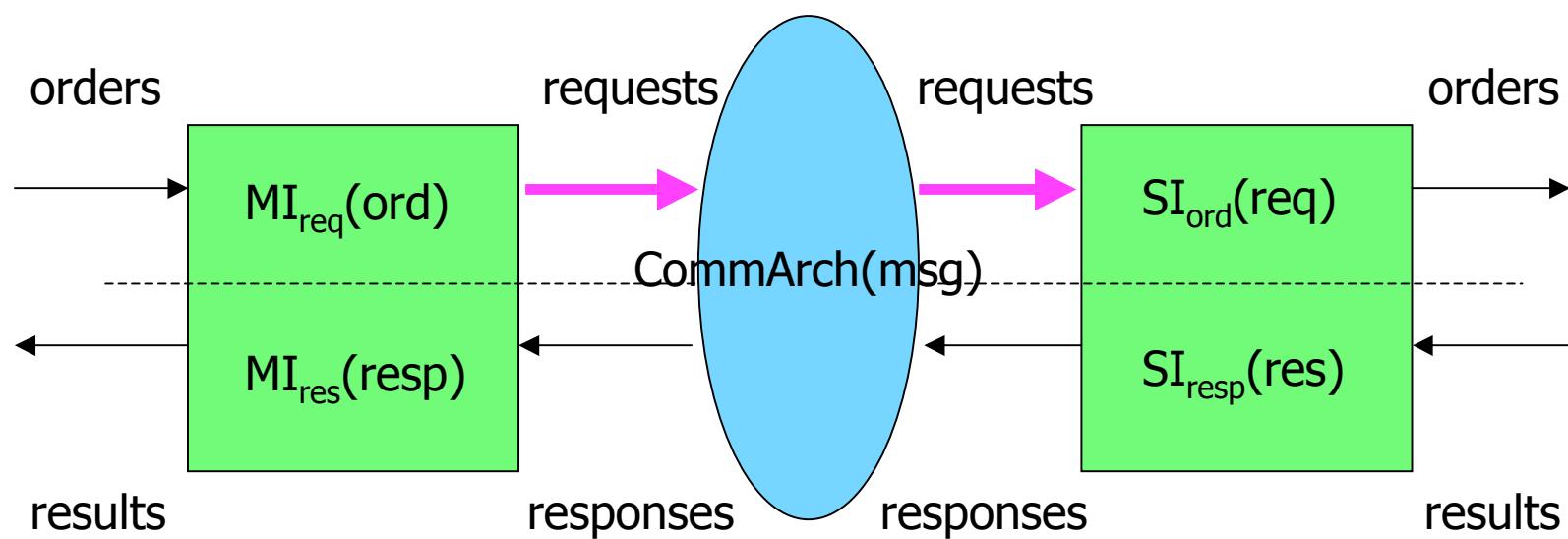
Functional Modeling



Definition of Transfers

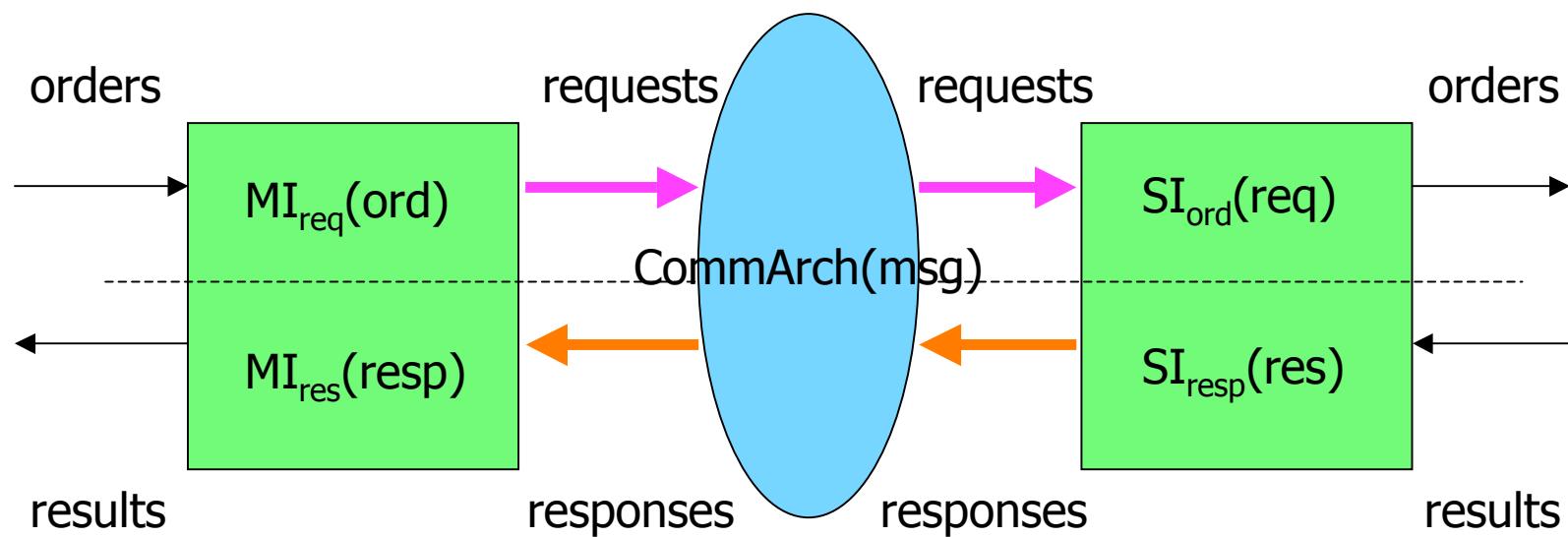


Definition of Transfers



$$\mathbf{Trans_ord(ord)} = \mathbf{SI_{ord} \circ CommArch \circ MI_{req}(ord)}$$

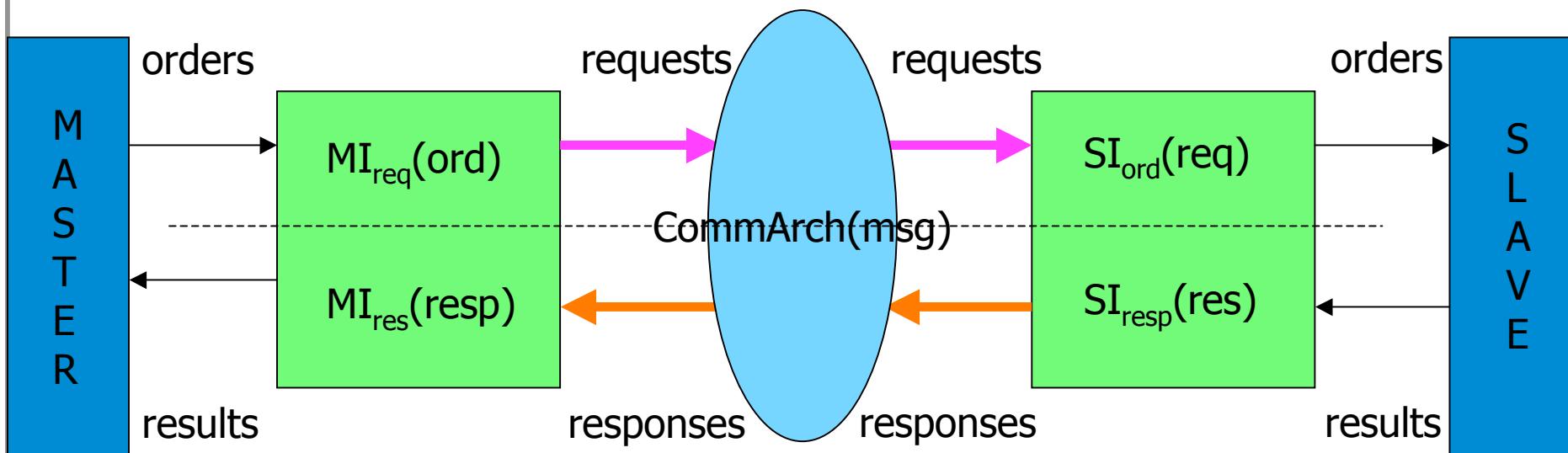
Definition of Transfers



$$Trans_ord(ord) = SI_{ord} \circ \text{CommArch} \circ MI_{req}(ord)$$

$$\mathbf{Trans_res(res) = MI_{res} \circ CommArch \circ SI_{resp}(res)}$$

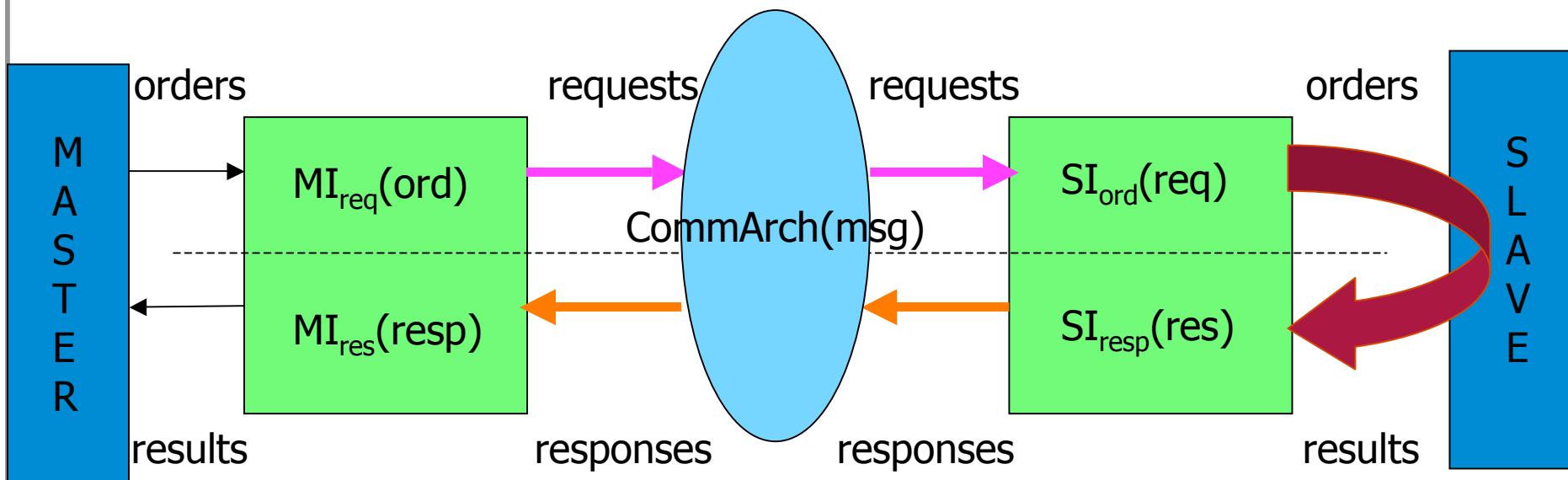
Correctness Criteria



THM
 $Trans_ord(ord) = SI_{ord} \circ CommArch \circ MI_{req}(ord) = ord$

$Trans_res(res) = MI_{res} \circ CommArch \circ SI_{resp}(res) = res$

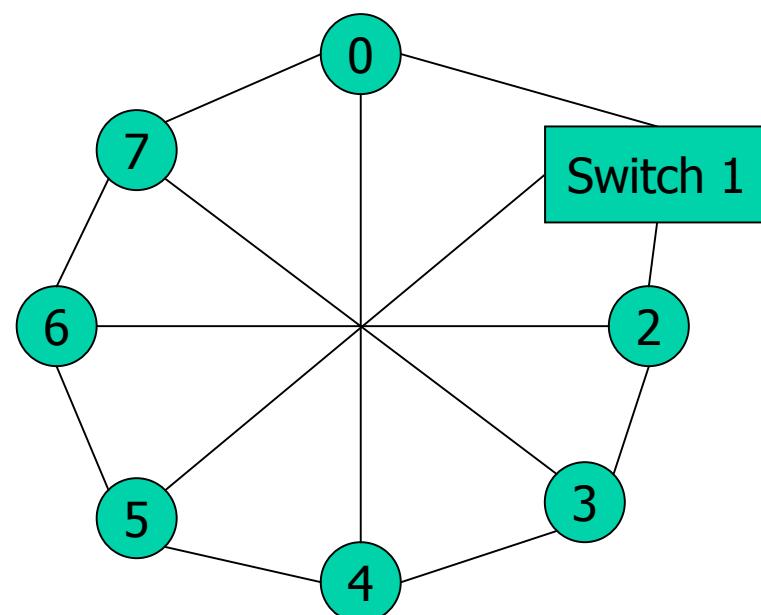
Correctness Criteria



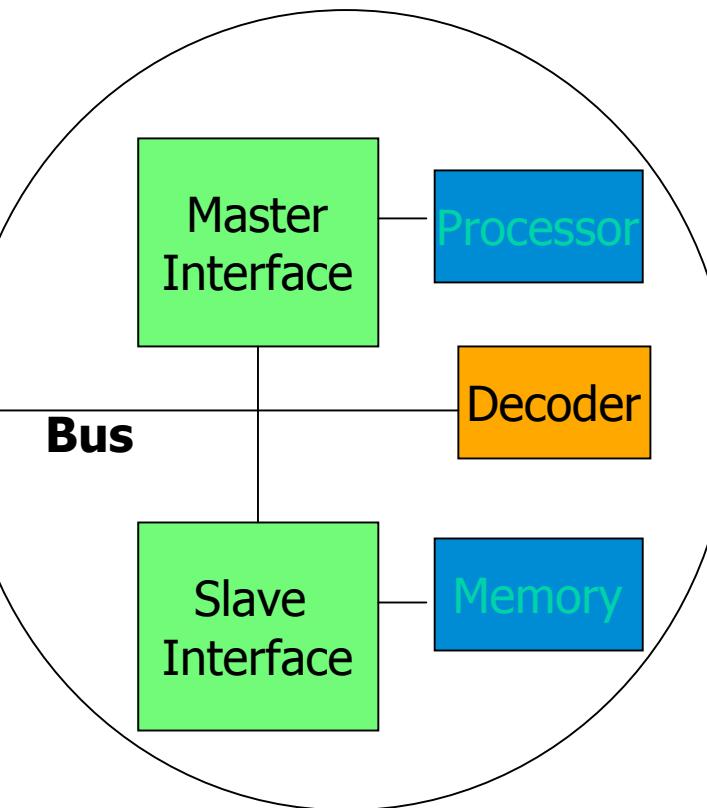
THM: Transfer(ord) = Slave(ord)

Transfer(ord) = Trans_res o Slave o Trans_ord(ord)

Octagon System

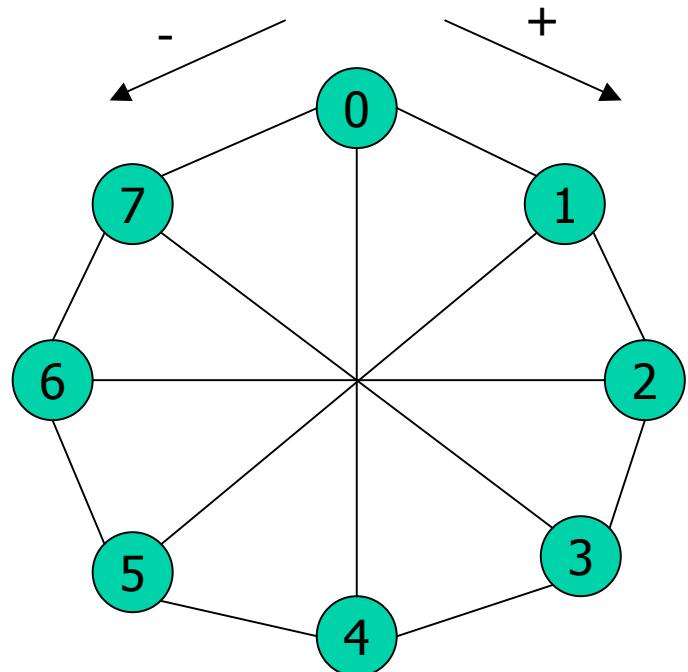


Parameterized Octagon Unit



Node System

Routing Algorithm



If $(\text{dest} - \text{from}) \bmod n \leq n/4$

then go clockwise (+)

If $3n/4 \leq (\text{dest}-\text{from}) \bmod n$

then go counter clockwise (-)

Else go across

Routing Definition

- Formals
 - Origin and destination
 - Number of nodes
- Result
 - A path which is a list of nodes
- Role
 - Compute the path between origin and destination

Routing Definition

```
(defun route (from dest n)
  (declare (xargs :measure (min (nfix (mod (- dest from) (* 4 n)))
                                (nfix (mod (- from dest) (* 4 n)))))))
  (cond ((or
           (not (integerp dest)) (< dest 0) (< (- (* 4 n) 1) dest)
           (not (integerp from)) (< from 0) (< (- (* 4 n) 1) from)
           (not (integerp n)) (<= n 0))
           nil)
         ((equal (- dest from) 0)
          (cons from nil))
         ((and (< 0 (mod (- dest from) (* 4 n)))
               (<= (mod (- dest from) (* 4 n)) n))
          (cons from (route (n_clockwise from n) dest n)))
         ((<= (* 3 n) (mod (- dest from) (* 4 n)))
          (cons from (route (n_counter_clockwise from n) dest n)))
         (t
          (cons from (route (n_across from n) dest n)))))))
```



Routing Theorem

```
(defthm CORRECTNESS_OF_ROUTE
  (implies (and (integerp from) (<= 0 from) (< from (* 4 n))
                 (integerp to) (<= 0 to) (< to (* 4 n)) (integerp n)
                 (< 0 n))
              (and ;; a route is a consp
                   (consp (route from to n))
                   ;; every node is an integer
                   (all_intp (route from to n))
                   ;; every node number is positive
                   (all_pos_intp (route from to n))
                   ;; every route contains no duplicate
                   (no-duplicatesp (route from to n)))
              ;; every node is less than the maximum of nodes
              (all_inf_np (route from to n) (* 4 n))
              ;; a route is made of available moves
              (AvailableMovep (route from to n) n)
              ;; the first node is the starting node
              (equal (car (route from to n)) from)
              ;; the last node is the final node
              (equal (car (last (route from to n))) to))))
```

Messages not lost in the medium {

Connection of a well-formed Master/Slave pair {



Scheduling Policy

- Circuit switched mode
 - A path is allocated during the complete transfer
- Formals
 - Travel list and two “working” variables
- Result
 - A set of non overlapping communications

Scheduling Policy

```
(defun scheduler (tl non_ovlp_r prev)
  ;; extracts non overlapping communications of tl
  (if (endp tl)
      (rev non_ovlp_r)
      (let ((route_i (cdr (car tl))))
        (if (no_intersectx route_i prev)
            (scheduler (cdr tl)
                      (cons (car tl) non_ovlp_r)
                      (append route_i prev))
            (scheduler (cdr tl) non_ovlp_r prev)))))

(defthm all_no_duplicatesp_scheduler
  ;; if tl contains routes with no duplicate, then grab_nodes
  ;; of scheduler contains no duplicate
  (implies (all_no_duplicatesp tl)
           (no-duplicatesp (grab_nodes (scheduler tl nil prev)))))
```

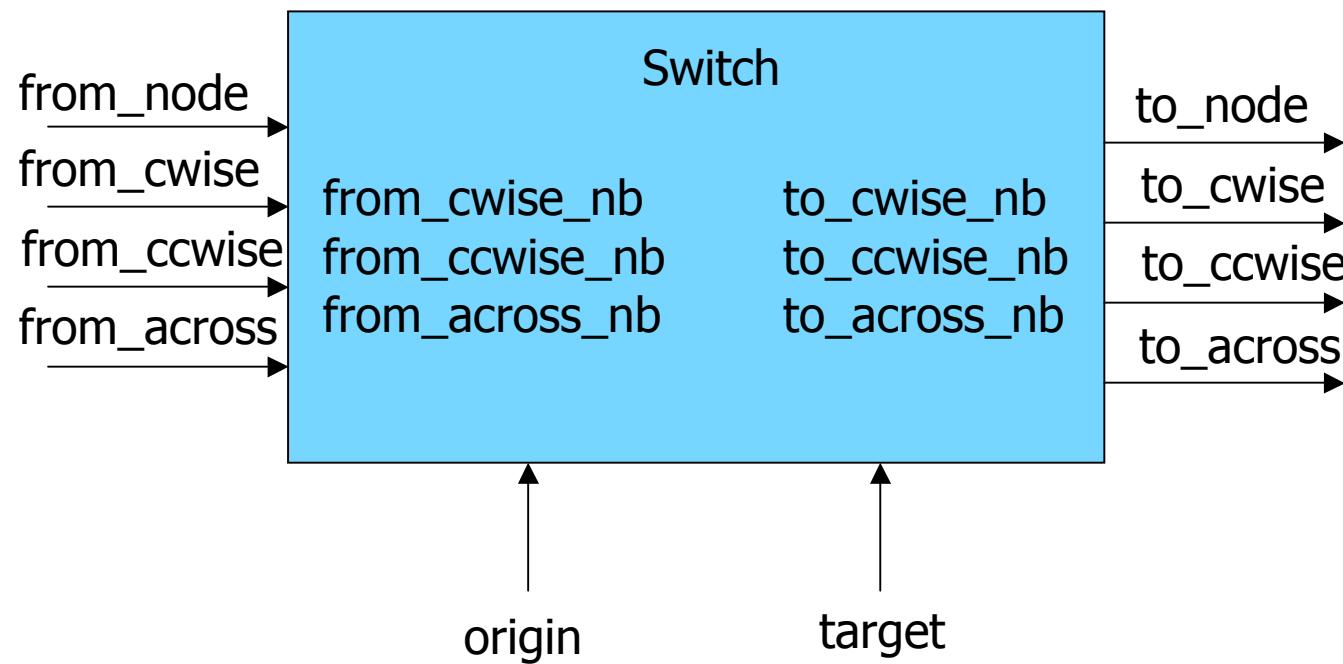
Scheduler correctness



Routing Correctness

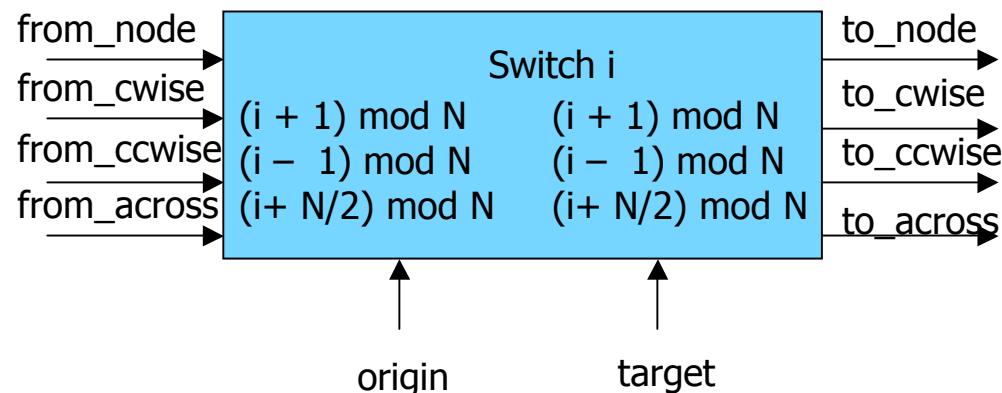
- } □ Every scheduled transfer is a well-formed non-overlapping Master/Slave communication

Switch



Traveling Functions

- Model the interconnection structure
 - **Trip: Travel List x NumNode ↳ Travel List**
 - Connection of Switches
 - Iterative calls to the switch function



Traveling Correctness

- We do not lose a message if its route is made of correct moves
 - $\text{Trip}(\text{tl}, \text{Num_Node}) = \text{tl}$
- We do lose a message if its route is not made of correct moves
 - $\text{Trip}(\text{tl}, \text{Num_Node}) \neq \text{tl}$

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(AvailableMovep (route from to n) n) \rightarrow $\text{Trip}(\text{tl}, \text{Num_Node}) = \text{tl}$



Memory Model

Local Memory = $(d_0 \ d_1 \ \dots \ d_{ms-1})$;; manipulated using nth and put-nth

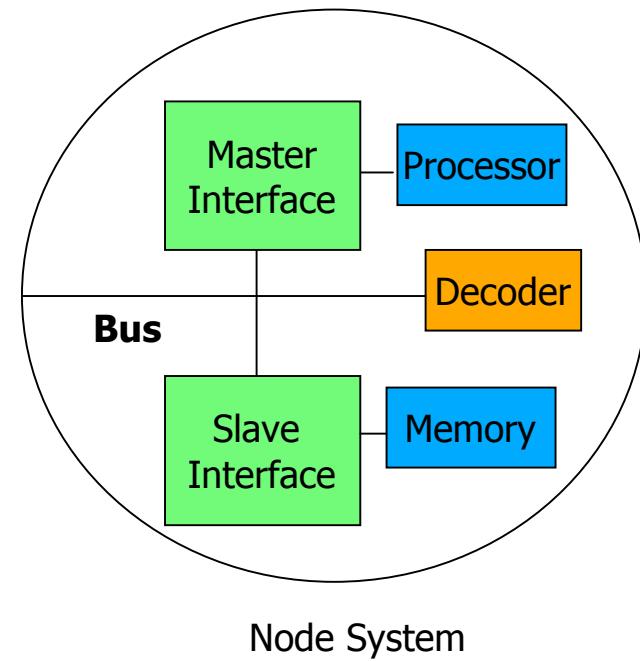
Local memory of node i

Global Memory = $(d_0 \ d_1 \ d_2 \ \dots \ d_{i.ms-1} \ d_{i.ms} \ \dots \ d_{(i+1).ms-1} \ d_{(i+1).ms} \ \dots \ d_{n.ms-1})$



Node System

- If a request arrives at a node
 - Function **Node** $\approx \text{SI_resp} \circ \text{Memory} \circ \text{SI_ord}$
- If a response is coming
 - Function **Node** $\approx \text{MI_res}$
- If a request is emitted
 - Function **Node** $\approx \text{MI_req}$



Node System

```
(defun node (op loc dat Glob_Mem
  nw_stat nw_r_dat ;; stat and dat coming from the network
  nw_r/w nw_addr nw_dat ;; msg coming from the network
  IncomingResponse IncomingRequest ;; scheduler commands
  node_nb ;; node number
  ms) ;; size of local memory
(if (equal op 'NO_OP) ;; node master is doing nothing
  ;; if there is a command from the scheduler, we do it
  (if (equal IncomingRequest 1) ;; a request is coming
      (let ((dec (decoder nw_addr ms node_nb)))
        (mv-let (st dat memo)
          (nw_transfer nw_r/w nw_addr nw_dat Glob_Mem
                      dec node_nb ms)
          (mv 'NO_OP 'NO_DATA ;; to the master
              st dat 'NO_MSG_DATA ;; res sent to th nw
              memo))) ;; new memory
    (if (equal IncomingResponse 1)
        ;; a response is coming -> call MI_res
```



Node System

```
;; else the node master is doing a write or read operation
(let ((dec (decoder loc ms node_nb)))
  (if (equal dec 1)
      ;; local transfer
      ;; else the node sends a msg to the nw
      (mv-let (r/w addr data)
              (mi_req op loc dat)
              (mv 'NO_OP 'NO_DATA ;; nothing to the master
                  r/w addr data ;; msg sent to the nw
                  Glob_Mem))))))
```



Node System

```
(defun nw_transfer (r/w addr data Glob_Mem sl_select node_nb ms)
  (mv-let (op loc dat)
    (si_ord r/w addr data sl_select ms)
    (mv-let (stat dat memo)
      (memory op loc dat
              (get_local_mem Glob_Mem node_nb ms))
      (mv-let (st d)
        (si_resp stat dat sl_select)
        (mv st d
            (update_local_mem Glob_Mem memo
                               node_nb ms)))))))
```



Octagon Definition

```
(defun Octagon (op_lst N ms Glob_Mem)
  ;; model of the complete network, i.e. nodes connected to Octagon
  ;; runs the Network once and returns loc_done nw_done and the memory
  (mv-let (loc_done nw_op Glob_Mem1)
    ;; first we collect the messages and execute the local ops
    ;; we also get the non local requests
    (collect_msg op_lst nil nil Glob_MEMORY ms) ;; node function
    ;; then we compute the travel list
    (let* ((tl (make_travel_list nw_op nil ms N)) ; Routing
           ;; we extract the set of non-overlapping communications
           (novlp (scheduler tl nil nil)))
      ;; we move every request to their destination
      (tl_at_dest (trip novlp N)))
  (mv-let (cr_lst Glob_Mem2)
    ;; we compute the response of every request (node)
    (ComputeResponses tl_at_dest Glob_Mem1 ms nil)
    ;; we move the responses back to their source node
    (let ((tl_back (trip cr_lst N)))
      ;; we compute the final result
      (mv-let (nw_done Glob_Mem3)
        (ComputeRes tl_back Glob_Mem2 ms nil)
        (mv loc_done nw_done Glob_Mem3)))))))
```



Final Theorems

```
(defthm mem_ok_read_Octagon
  ; we prove that in case of read transfers the memory is not
  ; changed
  (implies (and (all_read_op_lstp op_lst)
                (all_non_loc_op_lstp op_lst ms)
                (all_node_nb_validp op_lst (* 4 N))
                (all_address_validp op_lst (* 4 N) ms)
                (equal (len Glob_Mem) (* (* 4 N) ms)))
          (integerp N) (< 0 N)
          (true-listp Glob_Mem)
          (NODE_MEM_SIZEp ms))
    (equal (mv-nth 2 (Octagon op_lst N ms Glob_Mem))
           Glob_Mem)))
```



Final Theorems

```
(defthm mem_ok_write_octagon
  (implies (and (all_write_op_lstp op_lst)
                (all_non_loc_op_lstp op_lst ms)
                (all_node_nb_validp op_lst (* 4 N))
                (all_address_validp op_lst (* 4 N) ms)
                (equal (len Glob_Mem) (* (* 4 N) ms)))
            (integerp N) (< 0 N)
            (true-listp Glob_Mem)
            (NODE_MEM_SIZEp ms))
    (equal (mv-nth 2 (Octagon op_lst N MS Glob_Mem))
           (good_mem_write
             (scheduler
               (make_travel_list
                 (mv-nth 1 (collect_msg op_lst nil nil
                                         Glob_Mem ms))
                 NIL ms N)
               nil nil)
             Glob_Mem ms))))
```

(put-nth (addr req_lst_i)
 (data req_lst_i)
 Glob_Mem)

req_lst



Conclusions

- Specification and validation of a parameterized Octagon system
 - Industrial case study
 - Functional framework
- First sketch of a methodology
 - 1st design steps of generic NoC
- Original work
 - New application of ACL2
 - New application of theorem proving

Prospective Applications

- Packet and Wormhole Routing
- Protocols
 - Ethernet
- Buses
 - AMBA AHB (extension of ACL2'03)
 - Crossbars

Thank you !!

