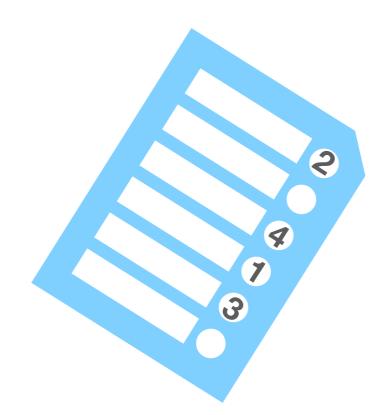
A Formalization of an Instant Run-Off Voting Scheme



Shilpi Goel & Mayank Manjrekar

ACL2-2018

15 min. Rump Session Talk

Instant Run-Off Voting

Instant Run-Off Voting a.k.a. Single-Winner Ranked Choice Voting

Instant Run-Off Voting a.k.a. Single-Winner Ranked Choice Voting a.k.a. Single Transferable Voting

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IRV is a *preferential voting* scheme: voters rank candidates in order of preference to elect one winner.

	Preference 1	Preference 2	Preference 3
Voter 1	Α	В	С
Voter 2	Α	С	В
Voter 3	С	В	Α
Voter 4	В	С	

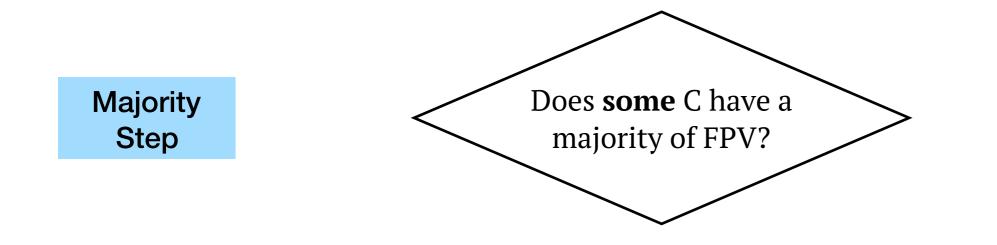
Where is IRV Used?

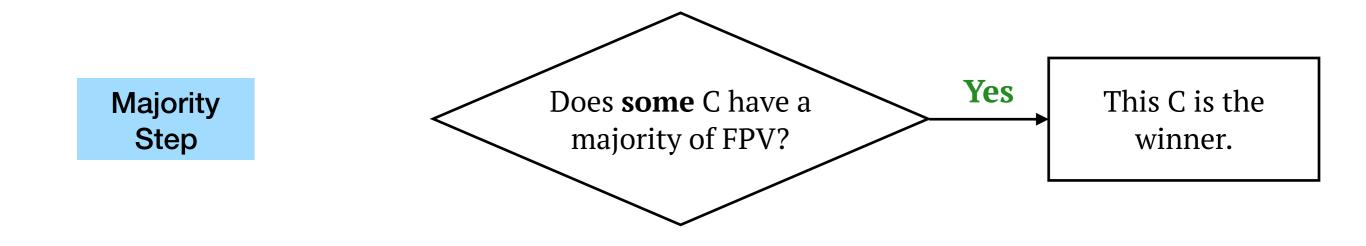
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 - Mayor of London
 - Members of the Australian Parliament's lower House
- Entertainment
 - Oscar's Best Picture Award
- Computer Science
 - Planning
 - Rank Aggregation Engines

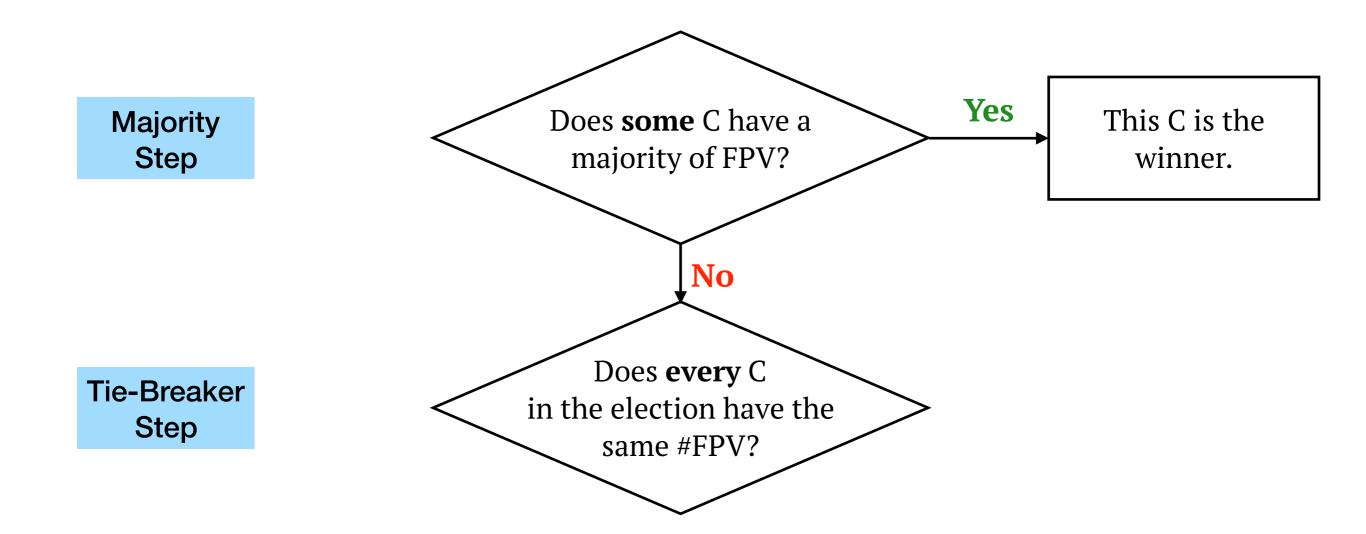
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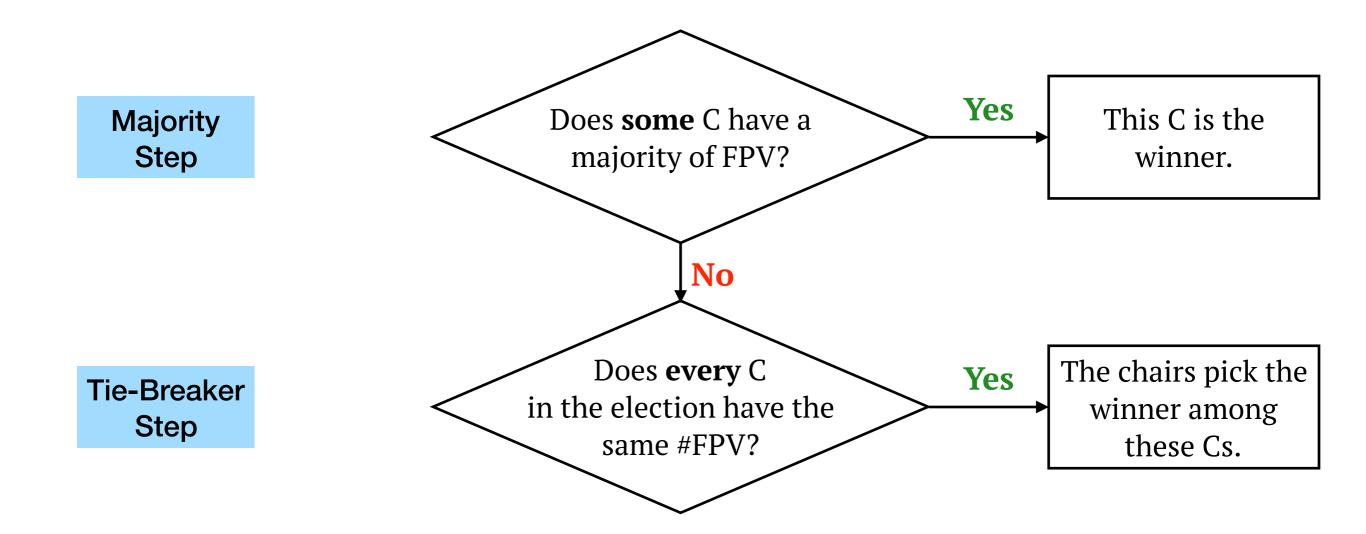
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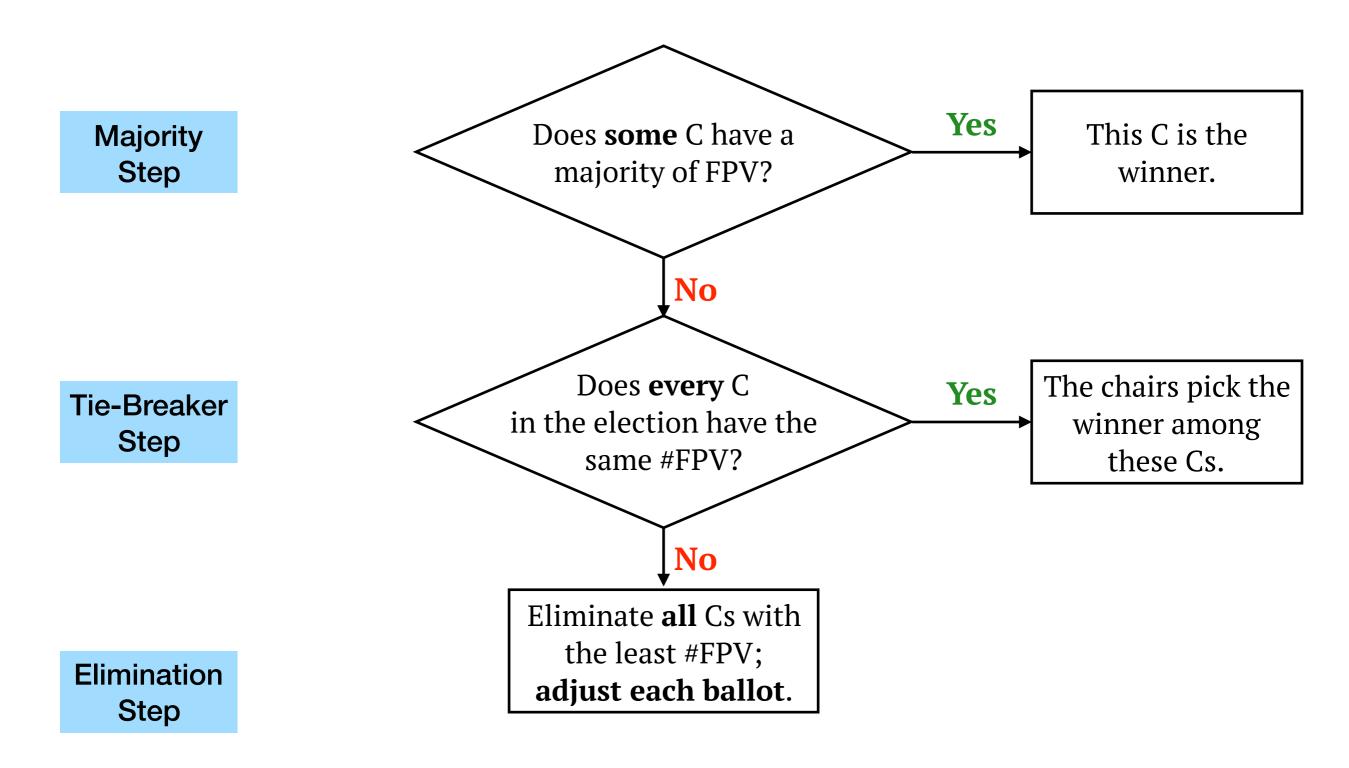
We got interested because of ACL2-2018's slogan election.

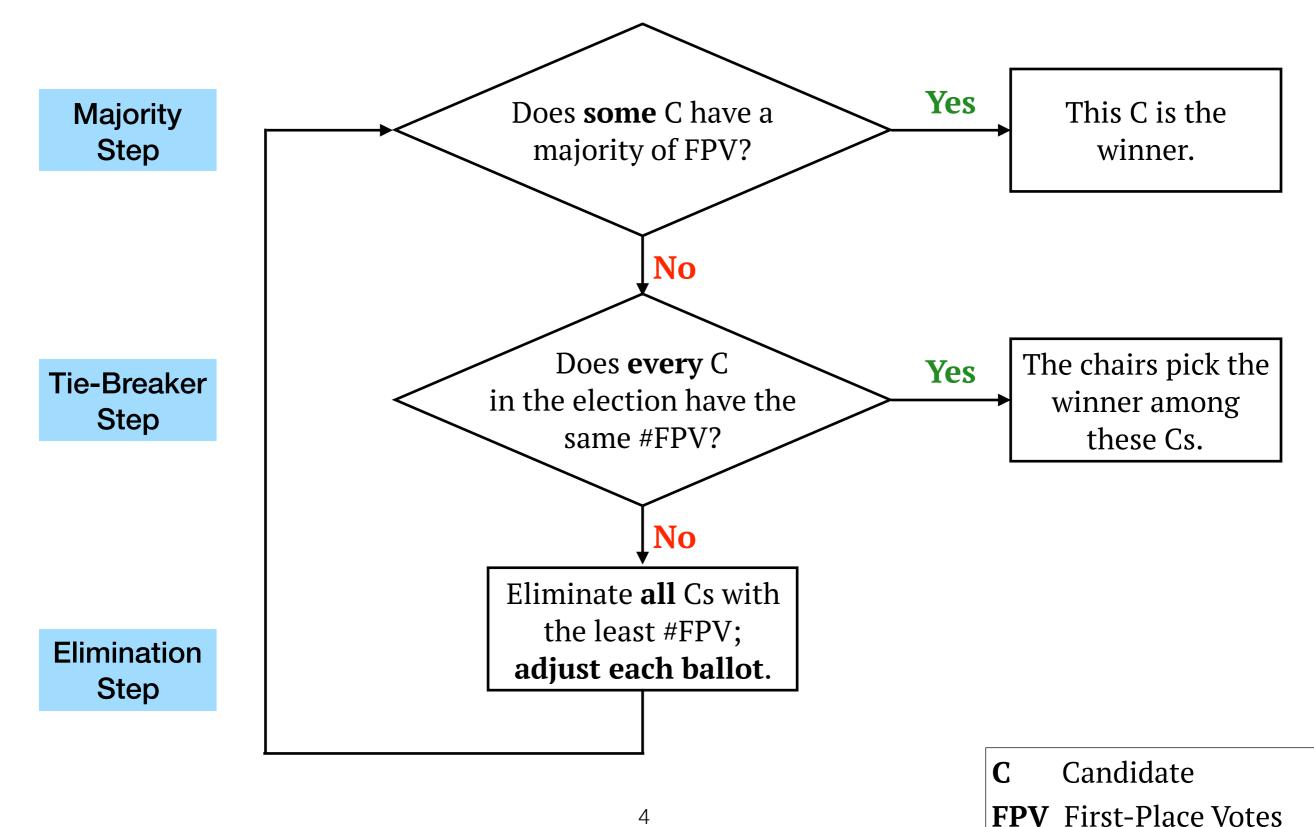












	Preference 1	Preference 2	Preference 3
Voter 1	Α	В	С
Voter 2	Α	С	В
Voter 3	С	В	Α
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Voter 4	В	С	
	Preference 1	Preference 2	Preference 3
Voter 1	Preference 1 A	Preference 2	Preference 3
Voter 1 Voter 2		Preference 2 -	Preference 3 -
	Α	Preference 2 - -	Preference 3 - - A

	Preference 1	Preference 2	Preference 3
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Voter 3	С	В	Α
Voter 4	В	С	
	Preference 1	Preference 2	Preference 3
Voter 1	Α	_	-
Voter 2	Α	-	-
Voter 3	-	-	Α
Voter 4	-	-	
	Preference 1		
Voter 1	Α	-	
Voter 2	Α		
Voter 3	Α		

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Voter 1	Α	_	-
Voter 2	Α	-	-
Voter 3	-	-	Α
Voter 4	-	-	
	Preference 1		
Voter 1	Α		
Voter 2	Α		
Voter 3	Α		

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Voter 1	Α	_	_
Voter 2	Α	_	-
Voter 3	-	_	Α
Voter 4	-	_	
	Preference 1		
Voter 1	Α		
Voter 2	Α	Α	wins
Voter 3	Α		

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Voter 1	Α	_	-
Voter 2	Α	-	_
Voter 3	-	-	Α
Voter 4	-	-	
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 Vote counting was done manually for ACL2-2018 slogans.

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Voter 1	Α	_	_
Voter 2	Α	_	-
Voter 3	-	_	Α
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- Matt Kaufmann:

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Voter 2	Α	-	-
Voter 3	-	_	Α
Voter 4	-	-	
	Preference 1		
Voter 1	Α		
Voter 2	Α	Α	wins
Voter 3	Α		

- Vote counting was done manually for ACL2-2018 slogans.
- This scheme seems a little unfair...

...different notions of fairness

Matt Kaufmann: *"if I were to do this again ... if there's a tie for least first-place votes, then it's broken by which of those has the least second-place votes, etc., before deleting candidates."*

	Preference 1	Preference 2	Preference 3
Voter 1	Α	В	С
Voter 2	Α	С	В
Voter 3	С	В	Α
Voter 4	В	С	

	Preference 1	Preference 2	Preference 3
Voter 1	Α	В	С
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Voter 2	Α	С	В
Voter 3	С	В	Α
Voter 4	В	С	

Pick-Candidate(B,C) = B

	Preference 1	Preference 2	Preference 3	
Voter 1	Α	B	С	
Voter 2	Α	С	B	Pick-Candidate(B,C) = I
Voter 3	С	B	Α	
Voter 4	В	С		

	Preference 1	Preference 2
Voter 1	Α	С
Voter 2	Α	С
Voter 3	С	Α
Voter 4	С	

	Preference 1	Preference 2	Preference 3	
Voter 1	Α	B	С	
Voter 2	Α	С	В	Pick-Candidate(B,C) = B
Voter 3	С	В	Α	
Voter 4	В	С		

	Preference 1	Preference 2
Voter 1	Α	С
Voter 2	Α	С
Voter 3	С	Α
Voter 4	С	

	Preference 1	Preference 2	Preference 3	
Voter 1	Α	B	С	
Voter 2	Α	С	B	Pick-Candidate(B,C) =
Voter 3	С	B	Α	
Voter 4	В	С		

	Preference 1	Preference 2	
Voter 1	Α	С	
Voter 2	Α	С	Eliminate A
Voter 3	С	Α	
Voter 4	С		

	Preference 1	Preference 2	Preference 3	
Voter 1	Α	B	С	
Voter 2	Α	С	B	Pick-Ca
Voter 3	С	B	Α	
Voter 4	В	С		
	Preference 1	Preference 2		
Voter 1	Α	С		
Voter 2	Α	С	Eliminate A	
Voter 3	С	Α		
Voter 4	С			
	Preference 1			
Voter 1	С			
Voter 2	С			
Voter 3	С			
Voter 4	С	lbo	oks]/	'pro

Lbooks]/projects/irv

	Preference 1	Preference 2	Preference 3
Voter 1	Α	B	С
Voter 2	Α	С	В
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Voter 4	В	С	
	Preference 1	Preference 2	
Voter 1	Α	С	•
Voter 2	Α	С	Eliminate A
Voter 3	С	Α	-
Voter 4	С		
	Preference 1		
Voter 1	С	C wins	S
Voter 2	С		
Voter 3	C		
Voter 4	С		oks]/

[books]/projects/irv

```
(defun irv (xs)
  (cond
   ((or (not (irv-ballot-p xs))
        (endp xs))
    nil)
  (t
   (b* ((cids (candidate-ids xs))
        (maj-winner? (first-choice-of-majority-p cids xs))
        ((when (natp maj-winner?)) maj-winner?)
        (weak-candidate
        (candidate-with-least-nth-place-votes 0 cids xs))
        (new-xs (eliminate-candidate weak-candidate xs)))
        (irv new-xs)))))
```

[books]/projects/irv

(defun candidate-with-least-nth-place-votes (n cids xs)
 (cond ((endp cids) nil)

```
((< n (number-of-candidates xs))
 (let* ((relevant-candidates
                (candidates-with-min-votes n cids xs)))
 (if (equal (len relevant-candidates) 1)
        (car relevant-candidates)
        (candidate-with-least-nth-place-votes
                (1+ n) relevant-candidates xs))))</pre>
```

```
(t
;; Tie persisted throughout all the preference
;; levels. Use a tie-breaker function.
 (pick-candidate cids))))
```

[books]/projects/irv

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      (car relevant-candidates)
    (candidate-with-least-nth-place-votes
      (1+ n) relevant-candidates xs))))
(t
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 (pick-candidate cids))))
                                       Constrained function:
```

returns a member of its input

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- Social Choice Theory
 - Very rich
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- Computational Choice Theory
 - Additional sets of concerns: parallelizability, importance of tie-breaking, etc.

Fairness Criteria

Our formalization meets the following criteria that should be satisfied by IRV schemes:

- Majority Winner Criterion
- Condorcet Loser Criterion
- Majority Loser Criterion

Majority Winner Criterion

If a candidate is preferred by an absolute majority of voters, then that candidate must win.

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Straightforward; needed lemmas like:

- If e gets the majority of first-place votes, then there cannot be a tie for the maximum number of first-place votes.

If a candidate L loses a head-to-head competition against every other candidate, then L must not win the overall election.

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Proof Sketch: Let w = (irv xs).

- If w won by a majority, then w would still have majority in every head-to-head competition; therefore, w != l.
- Otherwise, induct on xs.

- *Base Case*: xs has two candidates.

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 - Note that $w \neq id$.
 - -If l == id, then $w \neq l$.
 - Otherwise, l is still the head-to-head loser after id is eliminated. So by the induction hypothesis, w ≠ l.

Majority Loser Criterion

If a majority of voters prefers every other candidate over a given candidate l, then l must not win.

Note that *l* has a majority of last-place votes.

```
(defthm irv-satisfies-the-majority-loser-criterion
  (implies
   (and (< (majority (number-of-voters xs))
        (count l (make-nth-choice-list (last-place xs) xs)))
        (< 1 (number-of-candidates xs))
        (irv-ballot-p xs))
   (not (equal (irv xs) l))))
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   (not (equal (irv xs) l))))
```

But, a candidate who gets the majority of last-place votes must be the Condorcet Loser.

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- Possible future work: include other properties & schemes

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:doc irv::instant-runoff-voting

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Thanks!

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