## Topic 15 <br> Implementing and Using Stacks

"stack n.
The set of things a person has to do in the future. "I haven't done it yet because every time I pop my stack something new gets pushed." If you are interrupted several times in the middle of a conversation, "My stack overflowed" means "I forget what we were talking about."

## -The Hacker's Dictionary

Friedrich L. Bauer
German computer scientist who proposed "stack method of expression evaluation" in 1955.


## Sharper Tools



## Stacks

Lists


## Stacks

- Access is allowed only at one point of the structure, normally termed the top of the stack
- access to the most recently added item only

Operations are limited:

- push (add item to stack)
- pop (remove top item from stack)
- top (get top item without removing it)
- isEmpty
- Described as a "Last In First Out" (LIFO) data structure


## Implementing a stack

- need an underlying collection to hold the elements of the stack
- 3 obvious choices?
- native array
- linked structure of nodes
- a list!!!
- Adding a layer of abstraction. A HUGE idea.
- array implementation
- linked list implementation

$\frac{\text { https://xkcd.com/2347/ }}{4}$


## Uses of Stacks

- The runtime stack used by $\mathrm{a}_{\text {ato }}$ process (running program) to keep track of methods in progress
- Search problems
- Undo, redo, back, forward





## Stack Operations

Assume a simple stack for integers.
Stack<Integer> s = new Stack<>();
s.push(12);
s.push(4);
s.push( s.top() + 2 );
s.pop();
s.push( s.top() );
//what are contents of stack?

## Clicker 1 - What is Output?

Stack<Integer> s = new Stack<>();
// put stuff in stack
for (int i $=0$; i $<5$; i++) s.push(i);
// Print out contents of stack. // Assume there is a size method. for (int $i=0 ; i<s . s i z e() ; i++)$ System.out.print(s.pop() + " ");
$\begin{array}{llllll}A & 0 & 1 & 2 & 3 & 4\end{array}$
$\begin{array}{llllll}\text { B } & 4 & 3 & 2 & 1 & 0\end{array}$
C 432

D 234
E No output due to runtime error

## Corrected Version

Stack<Integer> $s=$ new $S t a c k<I n t e g e r>() ;$ // put stuff in stack
for (int $i=0 ; i<5 ; i++$ ) s.push (i) ;
// print out contents of stack
// while emptying it
final int LIMIT = s.size();
for (int $i=0 ; i<L I M I T ; i++)$
System.out.print(s.pop() + " ");
/ /or
// while (!s.isEmpty())
/ /
System.out.println(s.pop ());

## Stack Operations

Write a method to print out contents of stack in reverse order.

## Applications of Stacks

## Mathematical Calculations

- What does $3+2$ * 4 equal?

$$
2 * 4+3 ? \quad 3 * 2+4 ?
$$

- The precedence of operators affects the order of operations.
- A mathematical expression cannot simply be evaluated left to right.
- A challenge when evaluating a program.
- Lexical analysis is the process of interpreting a program.

What about $1-2-4^{\wedge} 5{ }^{*} 3 * 6 / 7^{\wedge} 2^{\wedge} 3$

## Infix and Postfix Expressions

- The way we are use to writing expressions is known as infix notation
- Postfix expression does not
- require any precedence rules
- 32 * $1+$ is postfix of 3 * $2+1$
' evaluate the following postfix expressions and write out a
 corresponding infix expression:

$$
\begin{array}{ll}
2324^{*}+* & 1234^{\wedge} *+ \\
12-32^{\wedge} 3^{*} 6 /+ & 25 \wedge 1-
\end{array}
$$

## Clicker Question 2

- What does the following postfix expression evaluate to? 632 + *
A. 11
B. 18
C. 24
D. 30
E. 36


## Evaluation of Postfix Expressions

Easy to do with a stack
given a proper postfix expression:

- get the next token
- if it is an operand push it onto the stack
- else if it is an operator
- pop the stack for the right hand operand
- pop the stack for the left hand operand
- apply the operator to the two operands
- push the result onto the stack
- when the expression has been exhausted the result is the top (and only element) of the stack


## Infix to Postfix

- Convert the following equations from infix to postix:
$2^{\wedge} 3^{\wedge} 3+5^{*} 1$
$11+2-1^{*} 3 / 3+2^{\wedge} 2 / 3$
Problems:
Negative numbers?
parentheses in expression


## Infix to Postfix Conversion

- Requires operator precedence parsing algorithm
- parse v . To determine the syntactic structure of a sentence or other utterance
Operands: add to expression
Close parenthesis: pop stack symbols until an open parenthesis appears
Operators:
Have an on stack and off stack precedence Pop all stack symbols until a symbol of lower precedence appears. Then push the operator
End of input: Pop all remaining stack symbols and add to the expression


## Simple Example

Infix Expression: $\quad 3+2$ * 4
PostFix Expression:
Operator Stack:

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\iota$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

## Infix Expression: +2 * 4

PostFix Expression: 3
Operator Stack:

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\iota$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

## Infix Expression: 2 * 4

PostFix Expression: 3
Operator Stack: +

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $*$ | 2 | 2 |
| $/$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

Infix Expression: * 4
PostFix Expression: 32
Operator Stack: +

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\iota$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

## Infix Expression: 4

PostFix Expression: 32
Operator Stack: + *

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\iota$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

Infix Expression:
PostFix Expression: 324
Operator Stack: + *

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\Lambda$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

## Infix Expression:

PostFix Expression: 324 *
Operator Stack:
$+$

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\iota$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Simple Example

Infix Expression:
PostFix Expression: $324^{*}+$
Operator Stack:

## Precedence Table

| Symbol | Off Stack <br> Precedence | On Stack <br> Precedence |
| :--- | :--- | :--- |
| + | 1 | 1 |
| - | 1 | 1 |
| $\star$ | 2 | 2 |
| $\Lambda$ | 2 | 2 |
| $\wedge$ | 10 | 9 |
| $($ | 20 | 0 |

## Example

## $11+2^{\wedge} 4{ }^{\wedge} 3-\left((4+5)^{*} 6\right)^{\wedge} 2$ <br> Show algorithm in action on above equation

## Balanced Symbol Checking

- In processing programs and working with computer languages there are many instances when symbols must be balanced \{\},[],()

A stack is useful for checking symbol balance. When a closing symbol is found it must match the most recent opening symbol of the same type.

- Applicable to checking html and xml tags!


## Algorithm for Balanced Symbol Checking

- Make an empty stack
- read symbols until end of file
- if the symbol is an opening symbol push it onto the stack
- if it is a closing symbol do the following
- if the stack is empty report an error
- otherwise pop the stack. If the symbol popped does not match the closing symbol report an error
- At the end of the file if the stack is not empty report an error


## Algorithm in practice

- list[i] = 3 * ( $44-\operatorname{method}($ foo( list[ 2 * $(i+1)+\mathrm{foo}($ list[i-1] ) ) / 2 * ) - list[ method(list[0])];
- Complications
- when is it not an error to have non matching symbols?
- Processing a file
- Tokenization: the process of scanning an input stream. Each independent chunk is a token.
- Tokens may be made up of 1 or more characters

