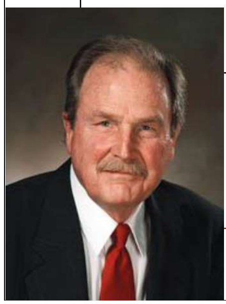
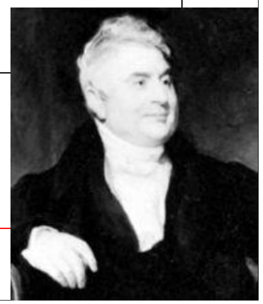
Topic 20: Huffman Coding

> *The author should gaze at Noah, and ... learn, as they did in the Ark, to crowd a great deal of matter into a very small compass.*
>
> Sydney Smith, Edinburgh Review





Agenda

- **Encoding**
- Compression
- Huffman Coding

Encoding

- UTCS
- 85 84 67 83
- 01010101 01010100 01000011 01010011

- What is stored in a jpg file? A text file? A Java file? A png file? A pdf file? An mp3 file? An mp4 file? An excel spreadsheet file? A zip file?
- open a bitmap in a text editor

Binary	Oct	Dec	Hex	Glyph	Binary	Oct	Dec	Hex	Glyph	Binary	Oct	Dec	Hex	Glyph
010 0000	040	32	20	(space)	100 0000	100	64	40	@	110 0000	140	96	60	`
010 0001	041	33	21	!	100 0001	101	65	41	A	110 0001	141	97	61	a
010 0010	042	34	22	"	100 0010	102	66	42	B	110 0010	142	98	62	b
010 0011	043	35	23	#	100 0011	103	67	43	C	110 0011	143	99	63	c
010 0100	044	36	24	$	100 0100	104	68	44	D	110 0100	144	100	64	d
010 0101	045	37	25	%	100 0101	105	69	45	E	110 0101	145	101	65	e
010 0110	046	38	26	&	100 0110	106	70	46	F	110 0110	146	102	66	f
010 0111	047	39	27	'	100 0111	107	71	47	G	110 0111	147	103	67	g
010 1000	050	40	28	(	100 1000	110	72	48	H	110 1000	150	104	68	h
010 1001	051	41	29	)	100 1001	111	73	49	I	110 1001	151	105	69	i
010 1010	052	42	2A	*	100 1010	112	74	4A	J	110 1010	152	106	6A	j
010 1011	053	43	2B	+	100 1011	113	75	4B	K	110 1011	153	107	6B	k
010 1100	054	44	2C	,	100 1100	114	76	4C	L	110 1100	154	108	6C	l
010 1101	055	45	2D	-	100 1101	115	77	4D	M	110 1101	155	109	6D	m
010 1110	056	46	2E	.	100 1110	116	78	4E	N	110 1110	156	110	6E	n
010 1111	057	47	2F	/	100 1111	117	79	4F	O	110 1111	157	111	6F	o
011 0000	060	48	30	0	101 0000	120	80	50	P	111 0000	160	112	70	p
011 0001	061	49	31	1	101 0001	121	81	51	Q	111 0001	161	113	71	q
011 0010	062	50	32	2	101 0010	122	82	52	R	111 0010	162	114	72	r
011 0011	063	51	33	3	101 0011	123	83	53	S	111 0011	163	115	73	s
011 0100	064	52	34	4	101 0100	124	84	54	T	111 0100	164	116	74	t
011 0101	065	53	35	5	101 0101	125	85	55	U	111 0101	165	117	75	u
011 0110	066	54	36	6	101 0110	126	86	56	V	111 0110	166	118	76	v
011 0111	067	55	37	7	101 0111	127	87	57	W	111 0111	167	119	77	w

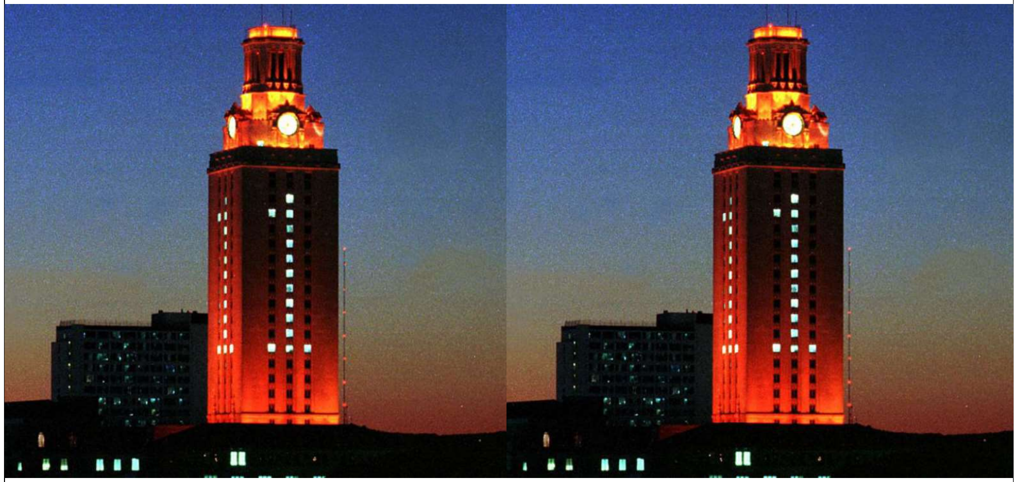
Text File

```
Project Gutenberg's The Adventures of Sherlock

This eBook is for the use of anyone anywhere at
almost no restrictions whatsoever.  You may cop
re-use it under the terms of the Project Gutenb
with this eBook or online at www.gutenberg.net


Title: The Adventures of Sherlock Holmes

Author: Arthur Conan Doyle

Posting Date: April 18, 2011 [EBook #1661]
First Posted: November 29, 2002
```


Text File???

```
50 72 6F 6A 65 63 74 20  47 75 74 65 6E 62 65 72  Project Gutenber
67 27 73 20 54 68 65 20  41 64 76 65 6E 74 75 72  g's The Adventur
65 73 20 6F 66 20 53 68  65 72 6C 6F 63 6B 20 48  es of Sherlock H
6F 6C 6D 65 73 2C 20 62  79 20 41 72 74 68 75 72  olmes, by Arthur
20 43 6F 6E 61 6E 20 44  6F 79 6C 65 0D 0A 0D 0A   Conan Doyle....
54 68 69 73 20 65 42 6F  6F 6B 20 69 73 20 66 6F  This eBook is fo
72 20 74 68 65 20 75 73  65 20 6F 66 20 61 6E 79  r the use of any
6F 6E 65 20 61 6E 79 77  68 65 72 65 20 61 74 20  one anywhere at
6E 6F 20 63 6F 73 74 20  61 6E 64 20 77 69 74 68  no cost and with
0D 0A 61 6C 6D 6F 73 74  20 6E 6F 20 72 65 73 74  ..almost no rest
72 69 63 74 69 6F 6E 73  20 77 68 61 74 73 6F 65  rictions whatsoe
```


Bitmap and JPEG File



Bitmap File????

```
07 01 06 05 00 04 03 01  04 02 02 05 03 03 07 02  ................
05 06 02 04 05 01 00 01  00 00 02 00 00 02 00 00  ................
03 01 04 07 05 09 0C 0A  07 0A 08 02 05 03 00 04  ................
05 00 04 03 00 04 03 01  03 03 02 05 03 04 05 03  ................
05 06 02 05 06 02 08 09  07 05 06 04 01 02 00 04  ................
04 04 0E 0E 0E 0E 0E 0E  0C 0B 0D 11 10 12 06 00  ................
0D 00 15 00 8F AD 3C FF  F8 A1 FB F9 BF F1 F8 B5  ....-<ÿø¡ûù¿ñøµ
FF FF C7 F3 FD A0 D2 EB  6B EE F9 7F FB FF 8F F7  ÿÿÇóý Òëkîù▮ûÿ÷
FF 94 FD FF A5 AA 9F 6B  02 03 00 00 00 16 42 81  ÿ▮ýÿ¥ª▮k......B
25 55 8D 3A 3D 68 25 2F  4C 1B 31 41 22 38 3F 30  %U:=h%/L.1A"8?0
0F 11 0B 0D 0B 0A 09 0E  11 00 04 07 00 01 02 06  ................
0C 0B 0C 13 10 06 0D 08  00 07 00 00 04 00 00 01  ................
00 02 07 06 04 09 08 00  04 03 00 03 02 05 0A 09  ................
06 0B 0A 01 06 05 12 06  12 03 05 0D 00 0E 00 28  ...............(
6C 17 E5 FA 7B EE FF AA  FB FB DD F5 F8 D8 FF FD  l.åú{îÿªûûÝõøØÿý
C5 FA F4 B3 F7 FA BC F9  FD D9 FF FF FA FC F5 F8  Åúô³÷ú¼ùýÙÿÿúüõø
FF FB E9 FF F8 CC C9 FC  82 F1 FA C8 FF F4 F5 FB  ÿûéÿøÌÉü▮ñúÈÿôõû
F5 E8 F9 FD DA FF FF E8  FF FF FB BC D7 CE 31 50  õèùýÚÿÿèÿÿû¼×Î1P
2B 01 16 00 00 02 00 01  03 04 09 08 04 0E 0D 03  +...............
03 04 00 09 08 04 0B 0C  0A 0A 0B 09 06 07 05 00  ................
01 00 00 01 00 06 07 05  0B 0C 0A 0D 0E 0C 19 11  ................
```

JPEG File

JPEG VS BITMAP

Tower-number1-1024x768.bmp	Bitmap image	2,305 KB
Tower-number1-1024x768.jpg	JPEG Image	283 KB

- JPEG File

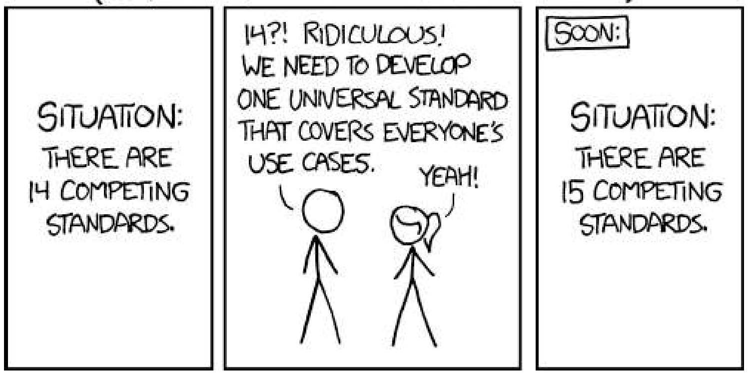
```
F D8 FF E0 00 10 4A 46  49 46 00 01 01 01 00 48  ÿØÿà..JFIF.....H
0 48 00 00 FF DB 00 43  00 02 01 01 02 01 01 02  .H..ÿÛ.C........
2 02 02 02 02 02 02 03  05 03 03 03 03 03 06 04  ................
4 03 05 07 06 07 07 07  06 07 07 08 09 0B 09 08  ................
8 0A 08 07 07 0A 0D 0A  0A 0B 0C 0C 0C 0C 07 09  ................
E 0F 0D 0C 0E 0B 0C 0C  0C FF DB 00 43 01 02 02  ..........ÿÛ.C...
2 03 03 03 06 03 03 06  0C 08 07 08 0C 0C 0C 0C  ................
C 0C 0C 0C 0C 0C 0C 0C  0C 0C 0C 0C 0C 0C 0C 0C  ................
C 0C 0C 0C 0C 0C 0C 0C  0C 0C 0C 0C 0C 0C 0C 0C  ................
C 0C 0C 0C 0C 0C 0C 0C  0C 0C 0C 0C 0C 0C FF C0  ..............ÿÀ
0 11 08 03 00 04 00 03  01 22 00 02 11 01 03 11  ........."......
1 FF C4 00 1F 00 00 01  05 01 01 01 01 01 01 00  .ÿÄ.............
0 00 00 00 00 00 00 01  02 03 04 05 06 07 08 09  ................
A 0B FF C4 00 B5 10 00  02 01 03 03 02 04 03 05  ..ÿÄ.µ..........
5 04 04 00 00 01 7D 01  02 03 00 04 11 05 12 21  .......}.......!
1 41 06 13 51 61 07 22  71 14 32 81 91 A1 08 23  1A..Qa."q.2´¡.#
2 B1 C1 15 52 D1 F0 24  33 62 72 82 09 0A 16 17  B±Á.RÑð$3br▮....
8 19 1A 25 26 27 28 29  2A 34 35 36 37 38 39 3A  ...%&'()*456789:
```


Encoding Schemes

- "It's all 1s and 0s"
- What do the 1s and 0s mean?
- 50 121 109
- ASCII -> 2ym
- Red Green Blue-> dark teal?

Why So Many Encoding / Decoding Schemes?

- Image file formats: bmp, png, jpg, gif, tiff, svg, cgm, pgm



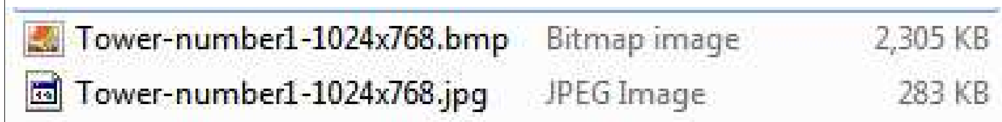
- XKCD, Standards: https://xkcd.com/927/

Agenda

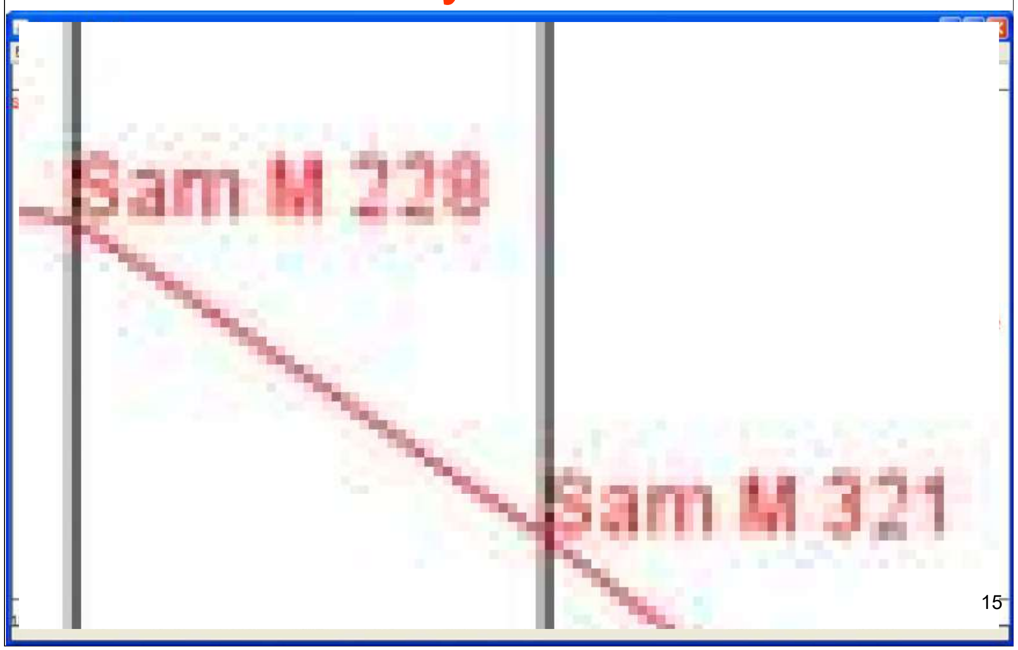
- Encoding
- **Compression**
- Huffman Coding

Compression

- Compression: Storing the same information but in a form that takes less memory
- lossless and lossy compression
- Recall:



Lossy Artifacts

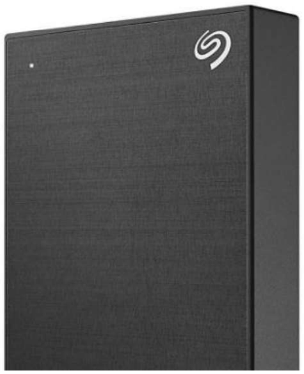


Compression

- 00000000000000000000000000000000 111111111111111111111111111111
- 0 00100000 1 00011110

Why Bother?

- Is compression really necessary?



5 Terabytes.
~5,000,000,000,0000 bytes

Clicker 1

- With computer storage so cheap, is compression really necessary?

A. No
B. Yes
C. It Depends

Little Pipes and Big Pumps

Home Internet Access

- 400 Mbps roughly $115 per month
- 12 months * 3 years * $115 =
- 400,000,000 *bits* /second = 5 * 10^7 bytes / sec

CPU Capability

- $2,000 for a good laptop or desktop
- **Intel® Core™ i9-7900X**
- Assume it lasts 3 years.
- Memory bandwidth 40 GB / sec = 4.0 * 10^{10} bytes / sec
- on the order of 6.4 * 10^{11} instructions / second

Mobile Devices?

Cellular Network

- Your mileage may vary …
- Mega bits per second
- AT&T

  17 mbps download, 7 mbps upload
- T-Mobile & Verizon

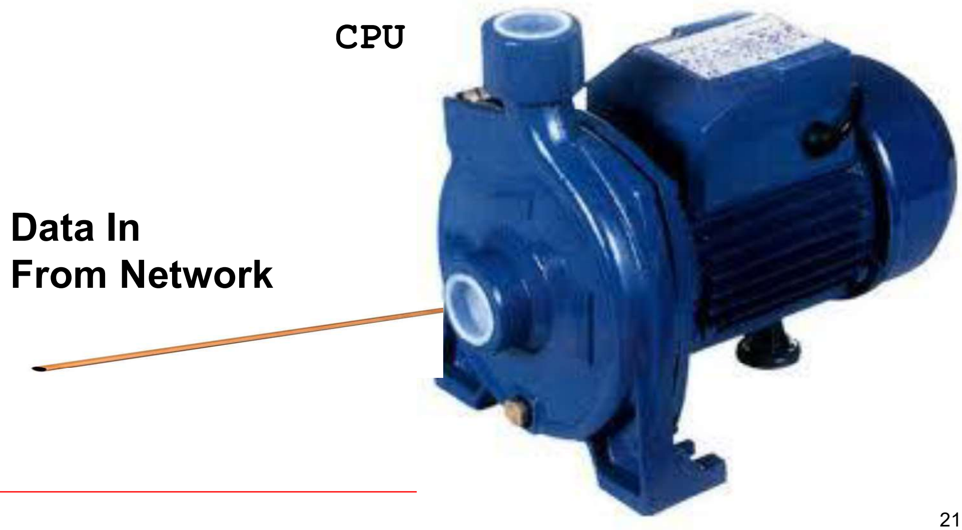
  12 mbps download, 7 mbps upload
- 17,000,000 bits per second = 2.125 x 10^6 bytes per second

http://tinyurl.com/q6o7wan

iPhone CPU

- Apple A6 System on a Chip
- Coy about IPS
- 2 cores
- Rough estimates: 1 x 10^{10} instructions per second

Little Pipes and Big Pumps



Agenda

- Encoding
- Compression
- **Huffman Coding**

Huffman Coding

- Proposed by Dr. David A. Huffman
  - Graduate class in 1951 at MIT with Robert Fano
  - term paper or final
  - term paper: prove min bits needed for binary coding of data
  - *A Method for the Construction of Minimum Redundancy Codes*
- Applicable to many forms of data transmission
  - Our example: text files
  - still used in fax machines, mp3 encoding, others

The Basic Algorithm

- Huffman coding is a form of statistical coding
- Not all characters occur with the same frequency, in typical text files. (can be true when reading raw bytes as well)
- Yet in ASCII all characters are allocated the same amount of space
  - 1 char = 1 byte, be it e or x
  - fixed width encoding

The Basic Algorithm

- Any savings in tailoring codes to frequency of character?
- Code word lengths are no longer fixed like ASCII or Unicode
- Code word lengths vary and will be shorter for the more frequently used characters
- Examples use characters for clarity, but in reality just read raw bytes from file.

The Basic Algorithm

1. **Scan file to be compressed and determine frequency of all values.**
2. **Sort or prioritize values based on frequency in file.**
3. **Build Huffman code tree based on prioritized values.**
4. **Perform a traversal of tree to determine new codes for values.**
5. **Scan file again to create new file using the new Huffman codes**

Building a Tree

Scan the original text

- Consider the following short text

  Eerie eyes seen near lake.

- Determine frequency of all numbers (values or in this case characters) in the text

Building a Tree

Scan the original text

Eerie eyes seen near lake.

- What characters are present?

E e r i space

y s n a r l k .

Building a Tree

Scan the original text

Eerie eyes seen near lake.

- What is the frequency of each character in the text?

Char	Freq.	Char	Freq.	Char	Freq.
E	1	y	1	k	1
e	8	s	2	.	1
r	2	n	2		
i	1	a	2		
space	4	l	1		

Building a Tree

Prioritize values from file

- Create binary tree nodes with a value and the frequency for each value
- Place nodes in a priority queue
  - The ***lower*** the frequency, the ***higher*** the priority in the queue

Building a Tree

- The queue after enqueueing all nodes

front ... back

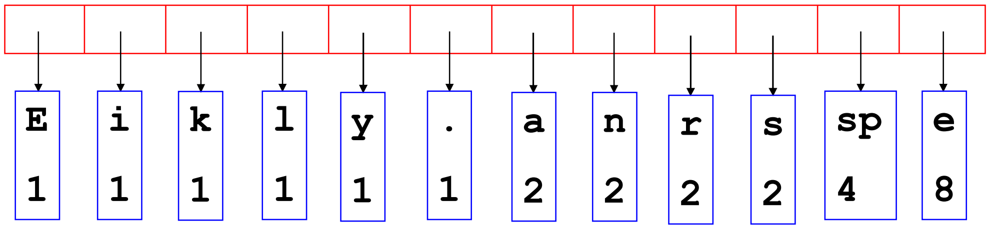
E	i	k	l	y	.	a	n	r	s	sp	e
1	1	1	1	1	1	2	2	2	2	4	8

- Null Pointers are not shown
- sp = space

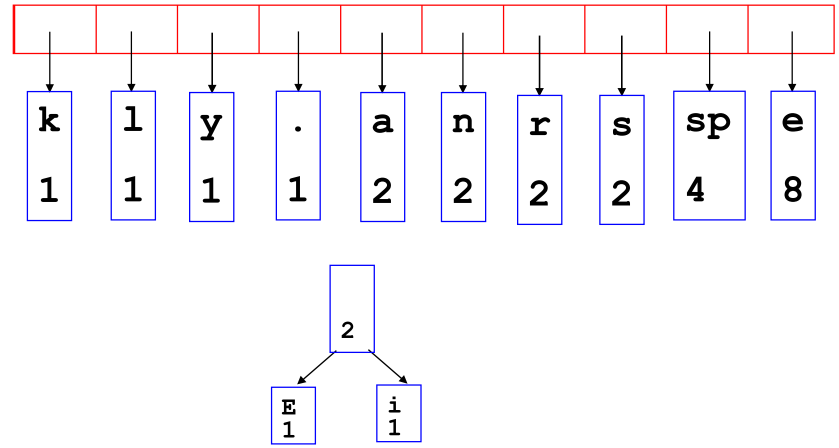
Building a Tree

- While priority queue contains two or more nodes
  - Create new node
  - Dequeue node and make it left child
  - Dequeue next node and make it right child
  - Frequency of new node equals sum of frequency of left and right children
    - New node does not contain value
  - Enqueue new node back into the priority queue

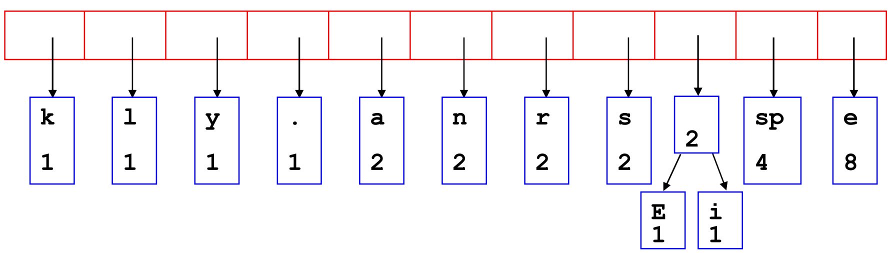
Building a Tree



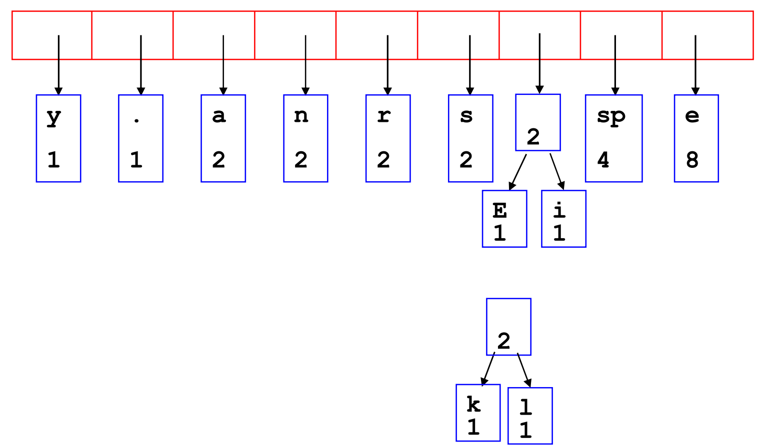
Building a Tree



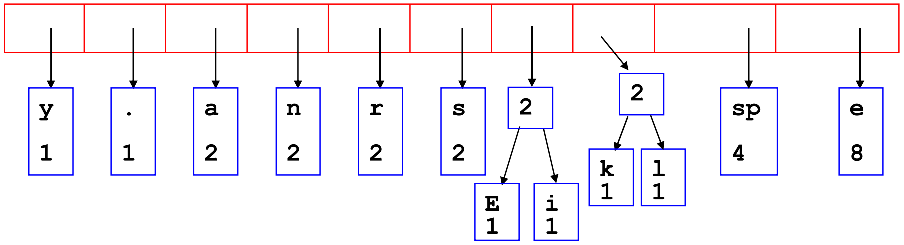
Building a Tree



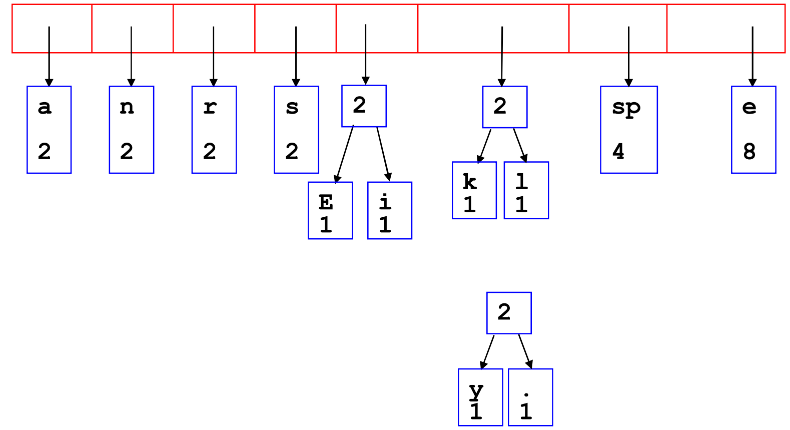
Building a Tree



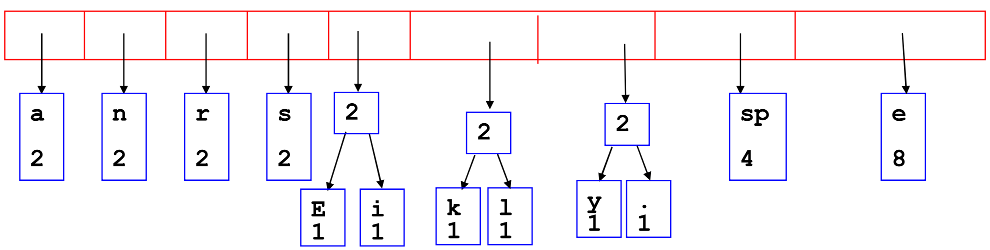
Building a Tree



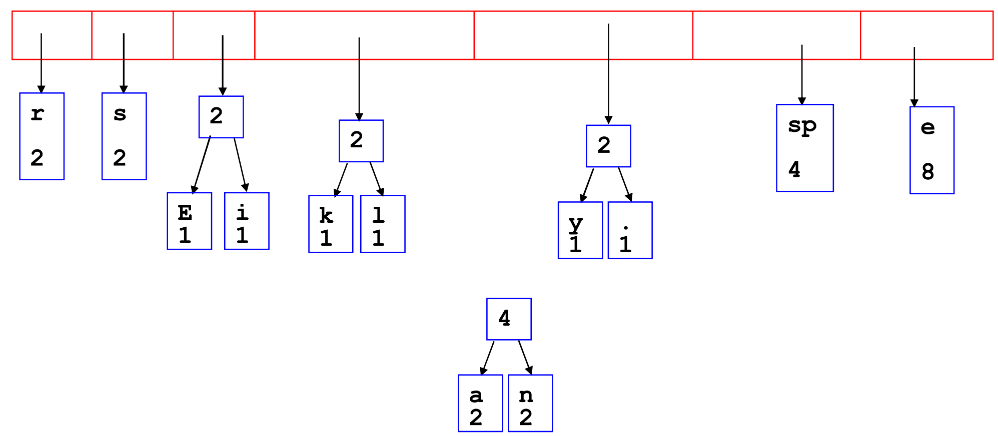
Building a Tree



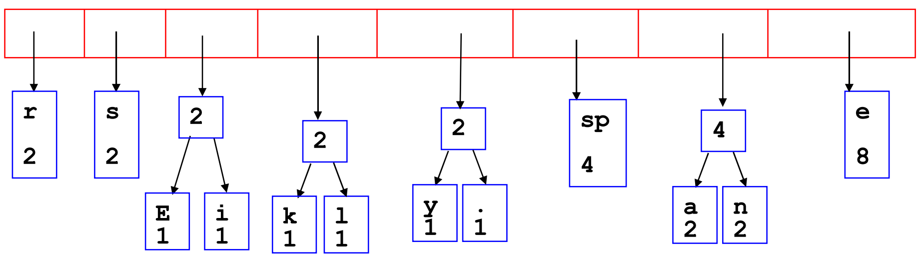
Building a Tree



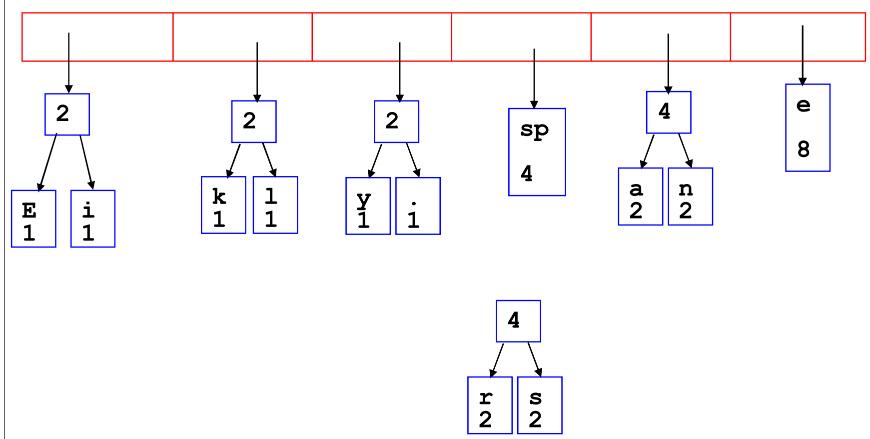
Building a Tree



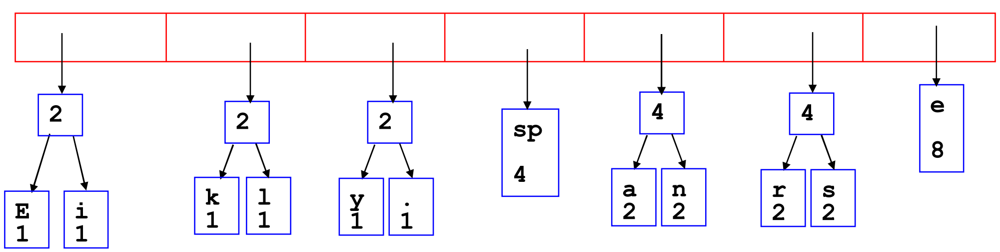
Building a Tree



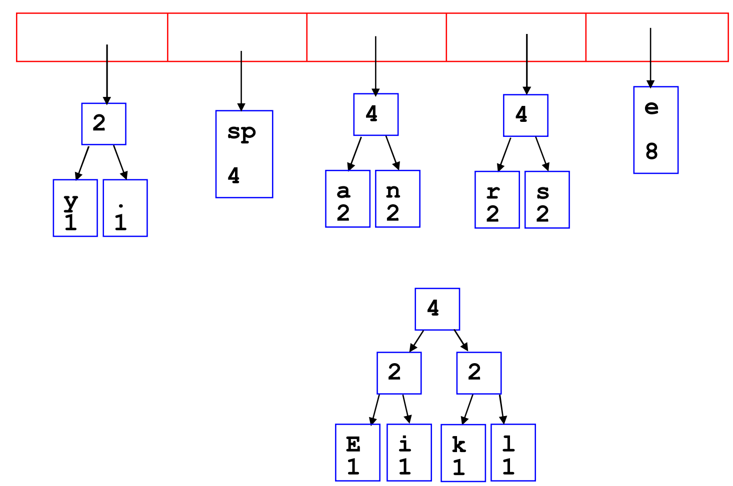
Building a Tree



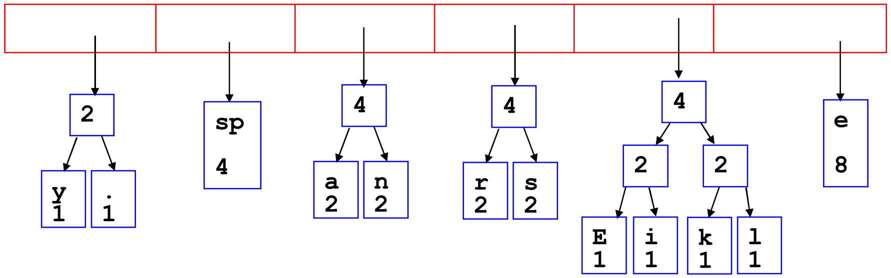
Building a Tree



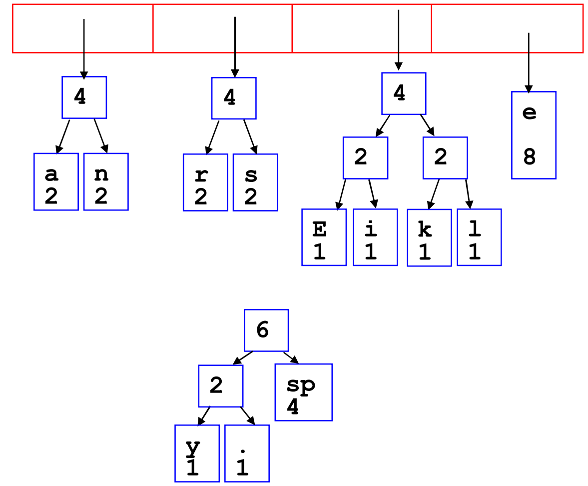
Building a Tree



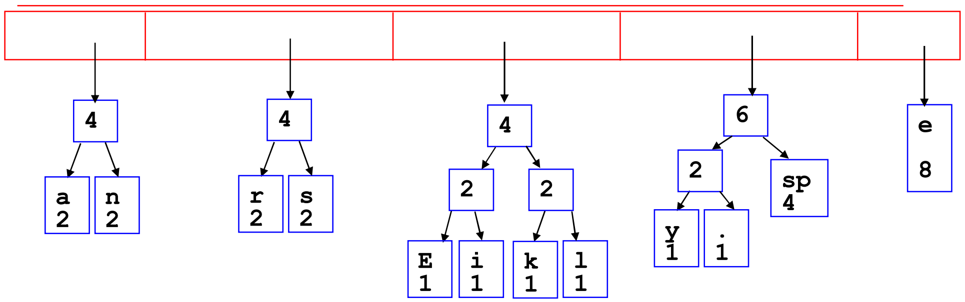
Building a Tree



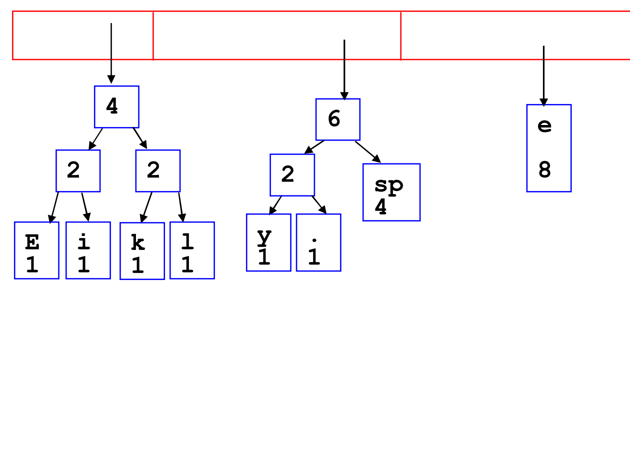
Building a Tree



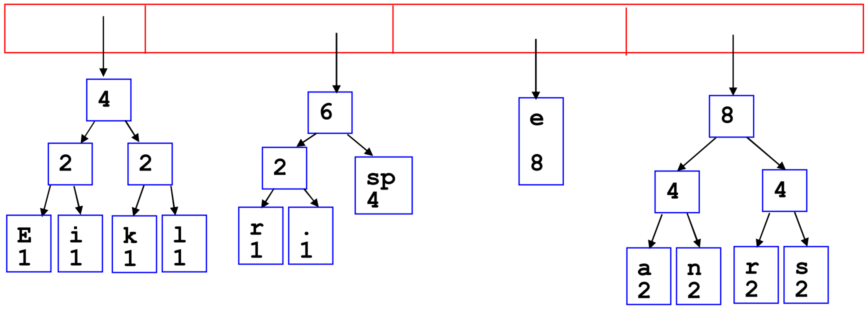
Building a Tree



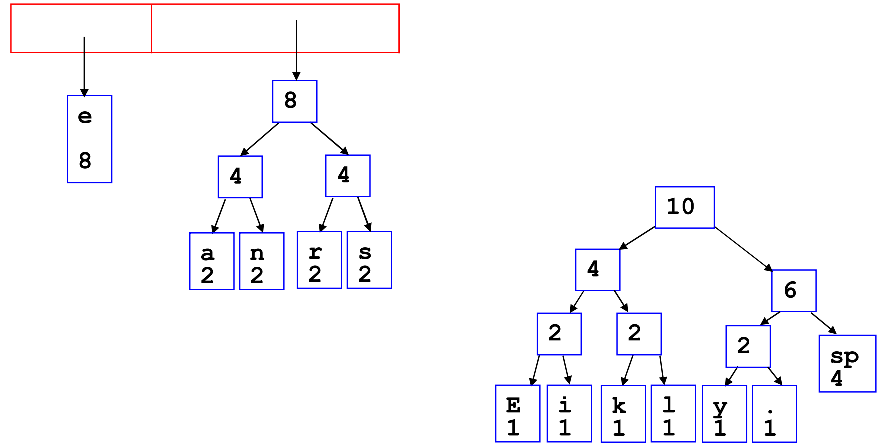
Building a Tree



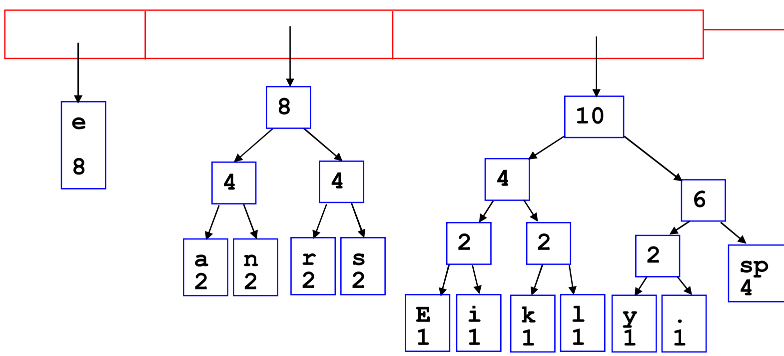
Building a Tree



Building a Tree



Building a Tree

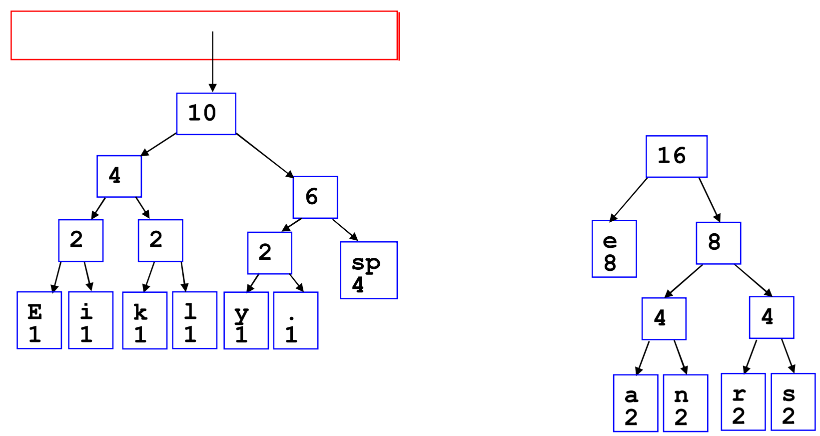


Clicker 2 - What is happening to the values with a low frequency compared to values with a high freq.?

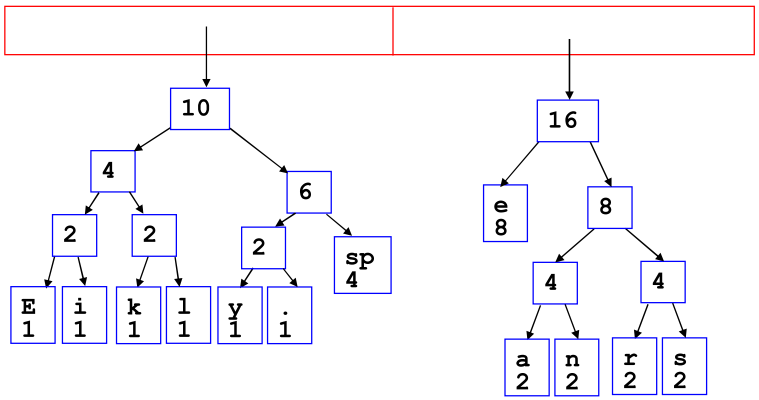
A. Smaller Depth B. Larger Depth

C. Something else

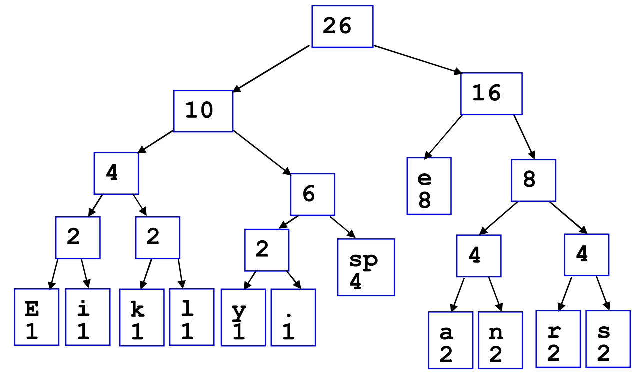
Building a Tree



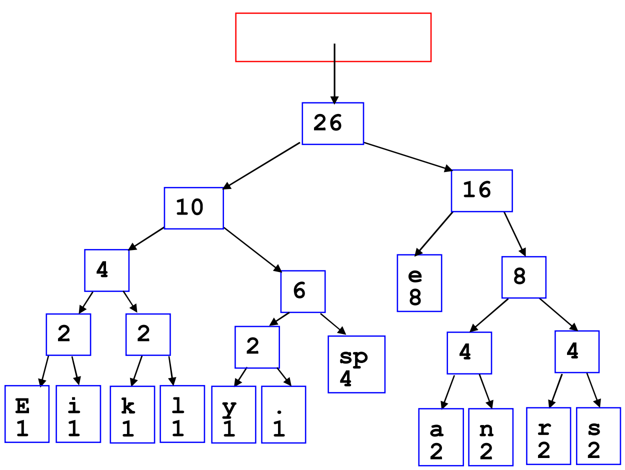
Building a Tree



Building a Tree



Building a Tree



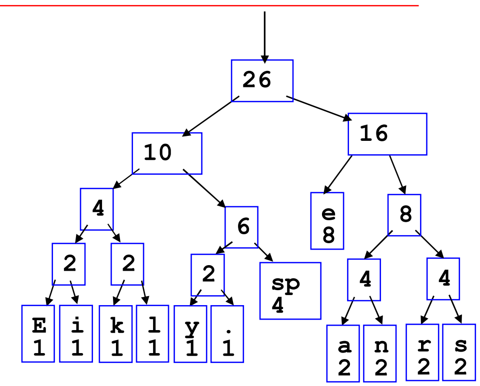
- After enqueueing this node there is only one node left in priority queue.

Building a Tree

Dequeue the single node left in the queue.

This tree contains the new code words for each character.

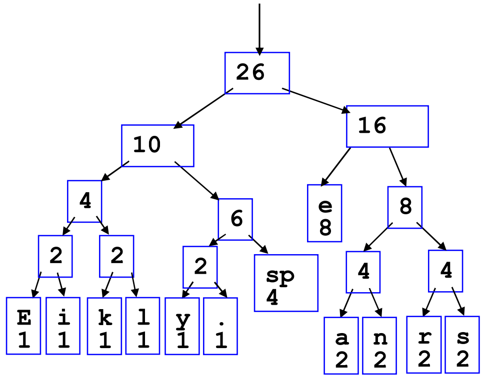
Frequency of root node should equal number of characters in text.



Eerie eyes seen near lake. 4 spaces, 26 characters total

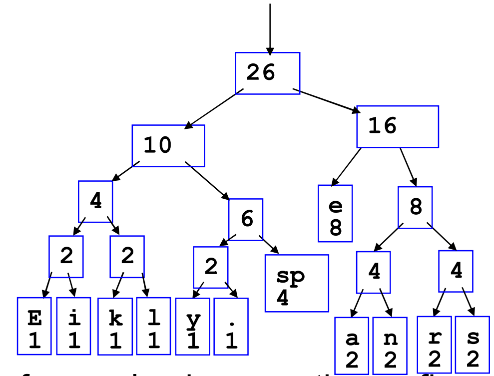
Encoding the File
Traverse Tree for Codes

- Perform a traversal of the tree to obtain new code words (sequence of 0's and 1's)
- left, append a 0 to code word
- right append a 1 to code word
- code word is only complete when a leaf node is reached



Encoding the File
Traverse Tree for Codes

Original Value	New Code
E (0100 0101)	0000
i (0110 1001)	0001
k (0110 1011)	0010
l (0110 1100)	0011
y (0111 1001)	0100
. (0010 1110)	0101
space (0010 0000)	011
e (0110 0101)	10
a (0110 0001)	1100
n (0110 1110)	1101
r (0111 0010)	1110
s (0111 0011)	1111



Prefix free codes. The code for a value in never the prefix of another code.

Encoding the File

- Rescan original file and encode file using new code words

Eerie eyes seen near lake.

```
0000101111000011001110
0100101111011111111010
1101011110110110011 10
0110011110000101001 01
```

Char	New Code
E	0000
i	0001
k	0010
l	0011
y	0100
.	0101
space	011
e	10
a	1100
n	1101
r	1110
s	1111

Encoding the File
Results

- Have we made things any better?
- 84 bits to encode the file
- ASCII would take 8 * 26 = 208 bits

```
0000101111000011001110
0100101111011111111010
1101011110110110011 10
0110011110000101001 01
```

- **If modified code used 4 bits per character are needed. Total bits 4 * 26 = 104. Savings not as great.**

Decoding the File

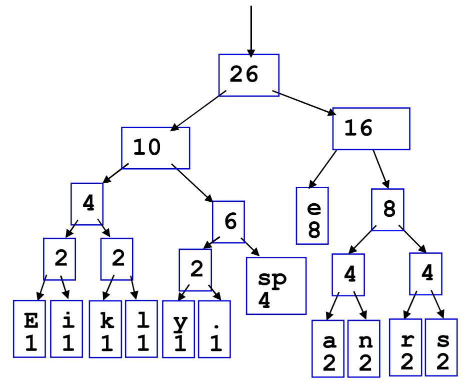
- How does receiver know what the codes are?
- Tree constructed for each file.
  - Considers frequency for each file
  - Big hit on compression, especially for smaller files
- Tree predetermined
  - based on statistical analysis of text files or other file types

Clicker 3 - Decoding the File

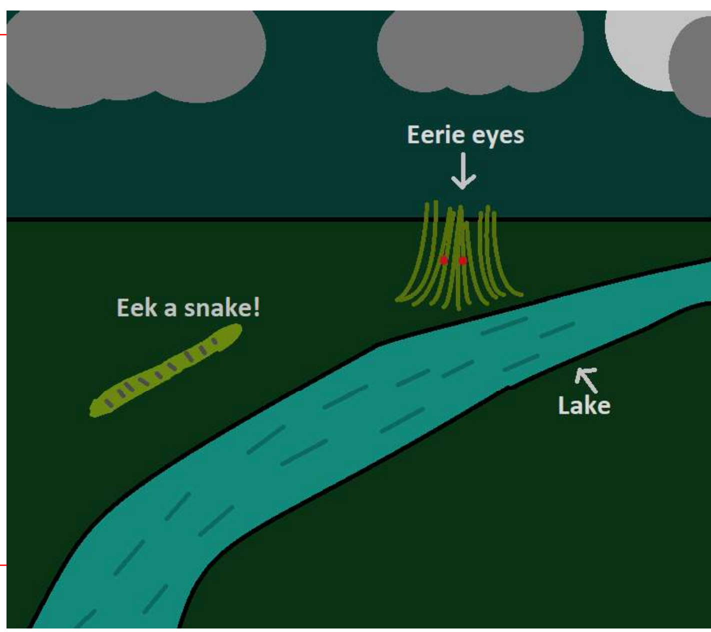
- Once receiver has tree it scans incoming bit stream
- 0 ⇒ go left
- 1 ⇒ go right

10100010011110001111111011100001010

A. elk nay sir
B. eek a snake
C. eek kin sly
D. eek snarl nil
E. eel a snarl



Alex Fall 2022



Assignment Hints

- reading chunks not chars
- header format
- the pseudo eof value
- the GUI

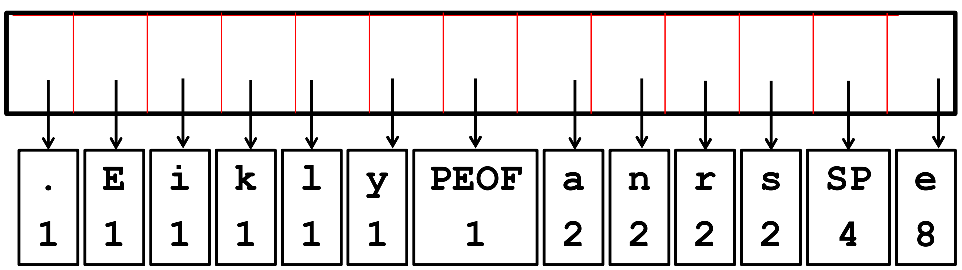
Assignment Example

- "Eerie eyes seen near lake." will result in different codes than those shown in slides due to:
  - adding elements in order to PriorityQueue
  - required pseudo eof value (PEOF)

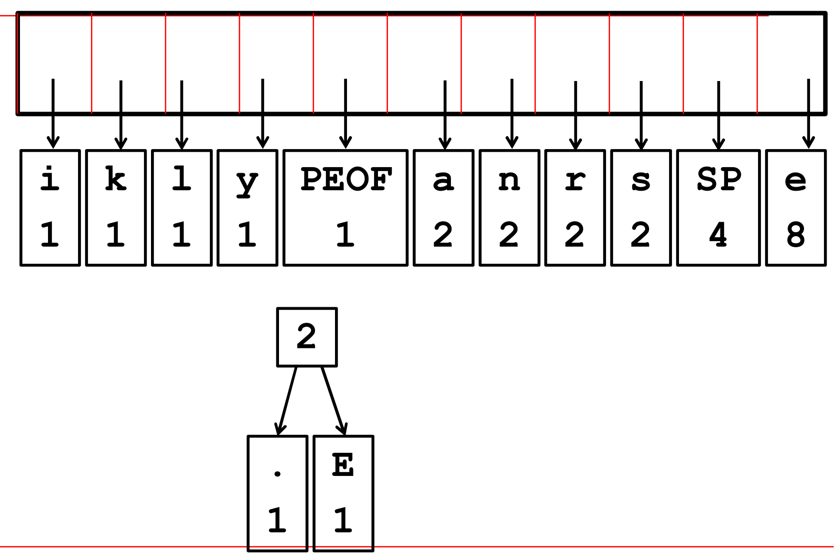
Assignment Example

Char	Freq.	Char	Freq.	Char	Freq.
E	1	y	1	k	1
e	8	s	2	.	1
r	2	n	2	PEOF	1
i	1	a	2		
space	4	l	1		

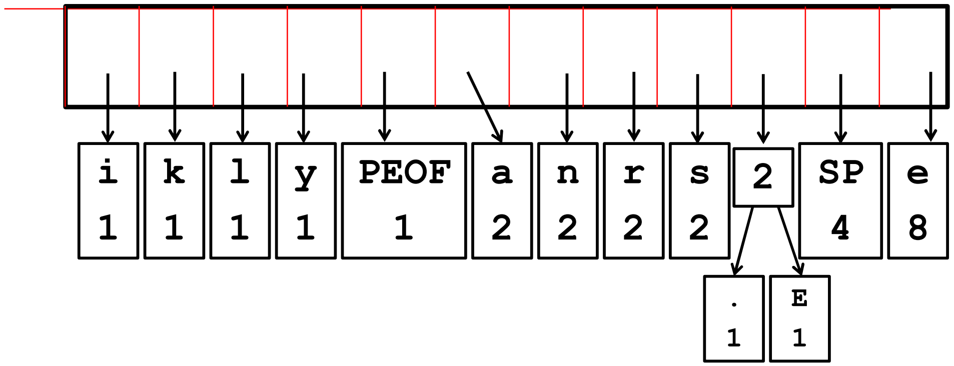
Assignment Example



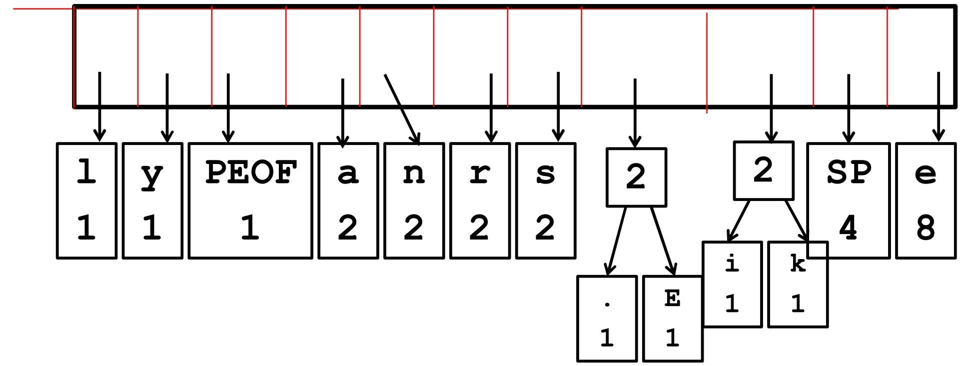
Assignment Example



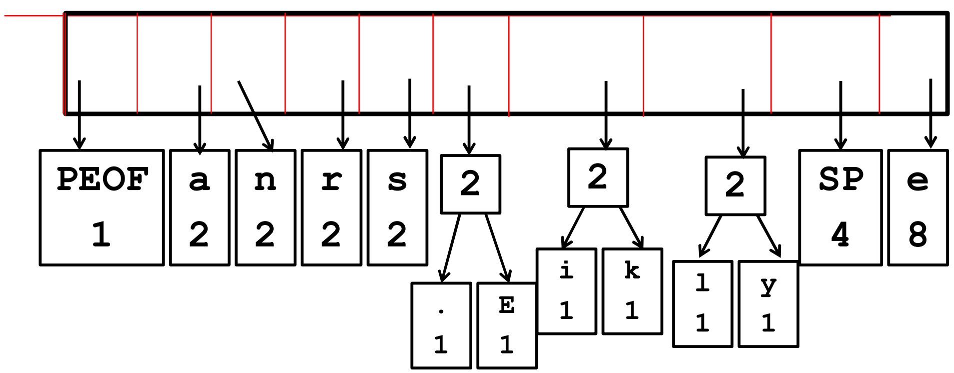
Assignment Example



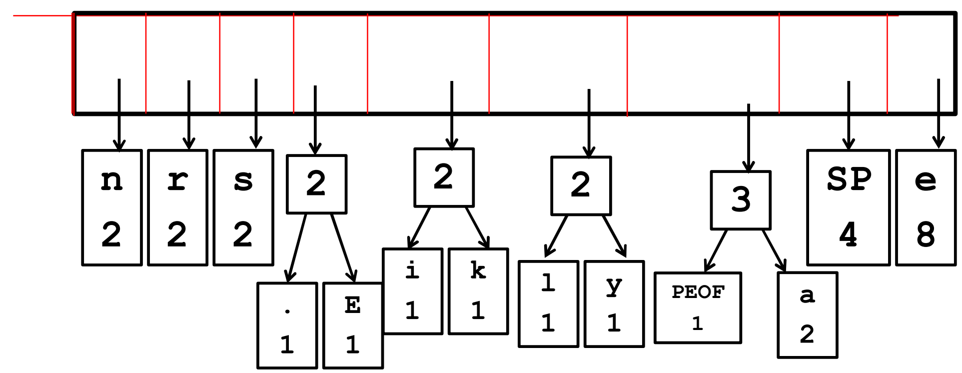
Assignment Example



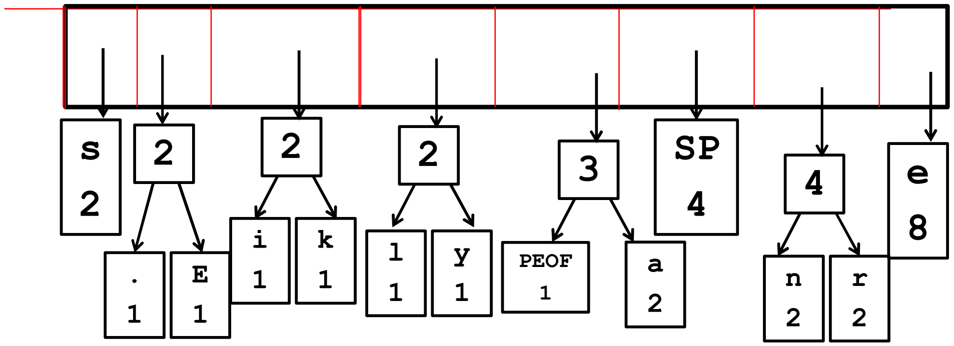
Assignment Example



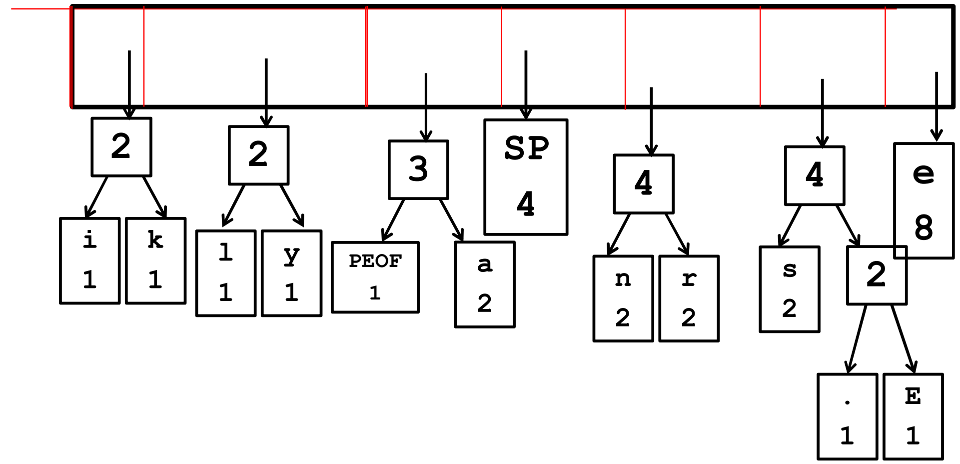
Assignment Example

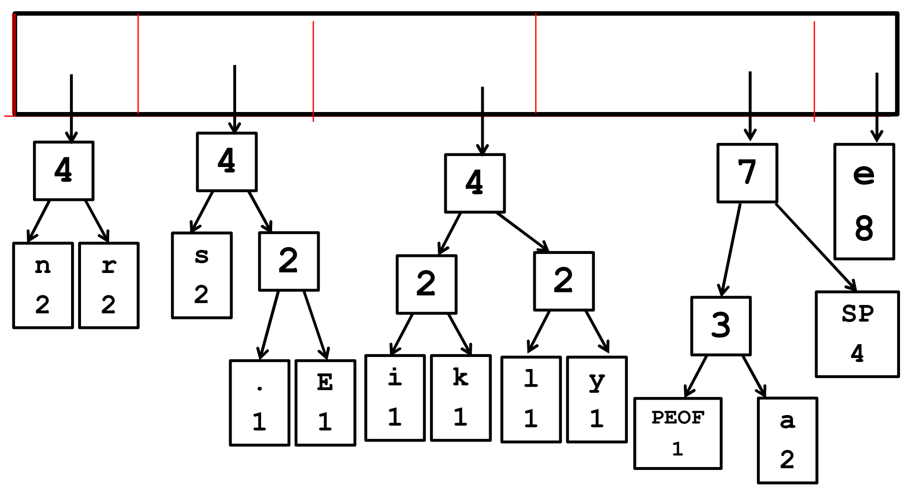


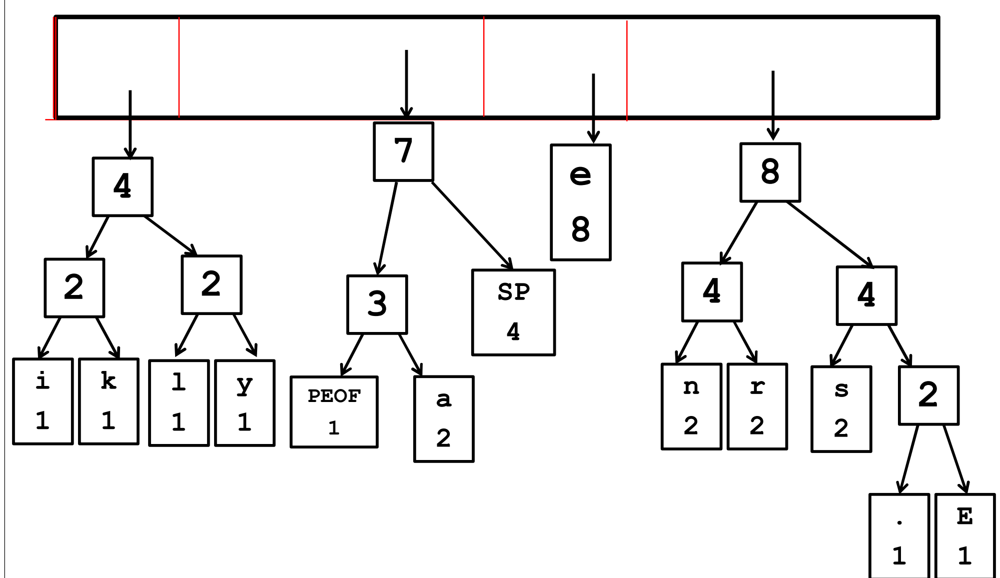
Assignment Example

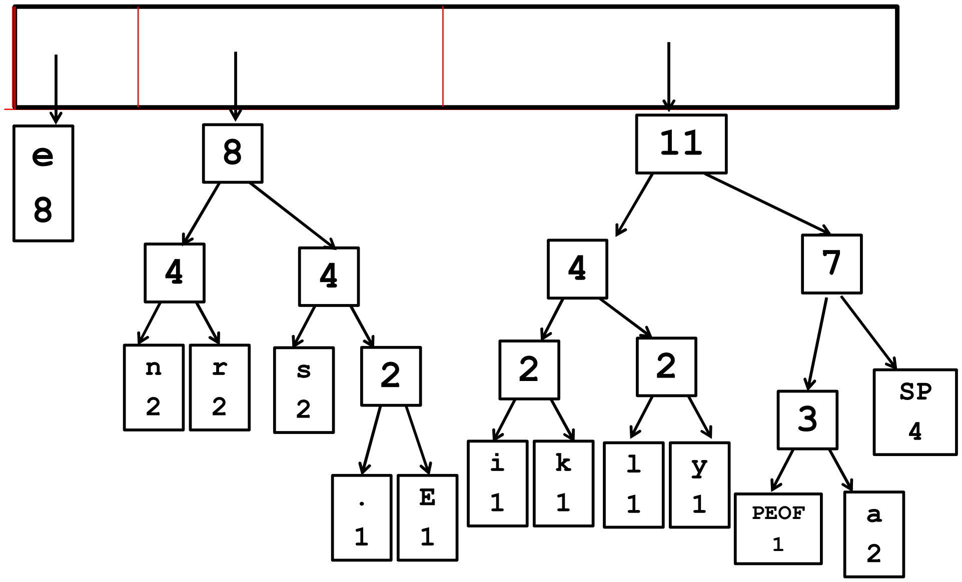


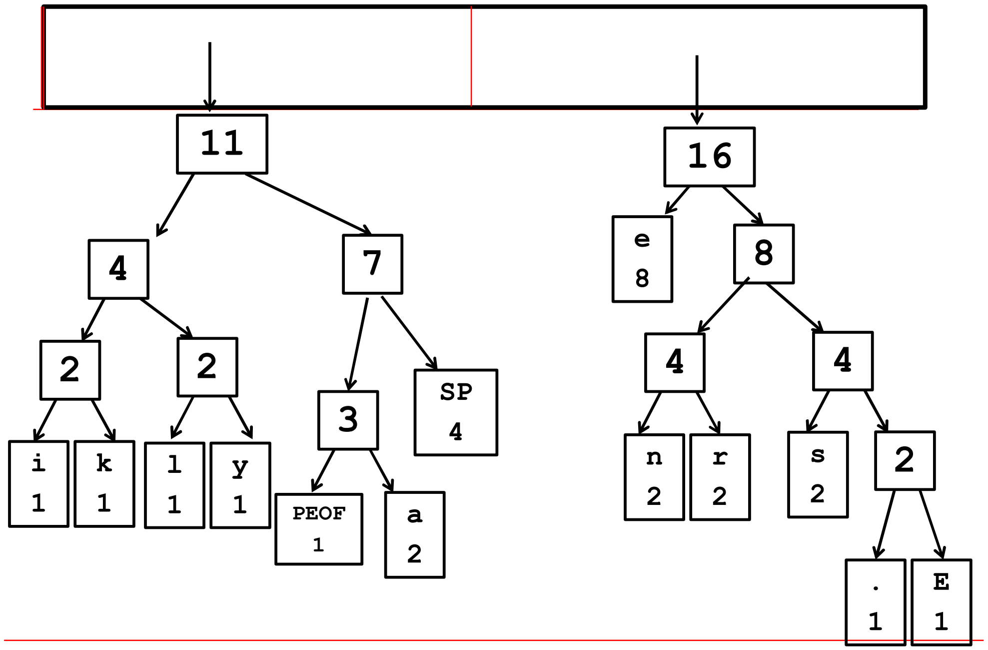
Assignment Example

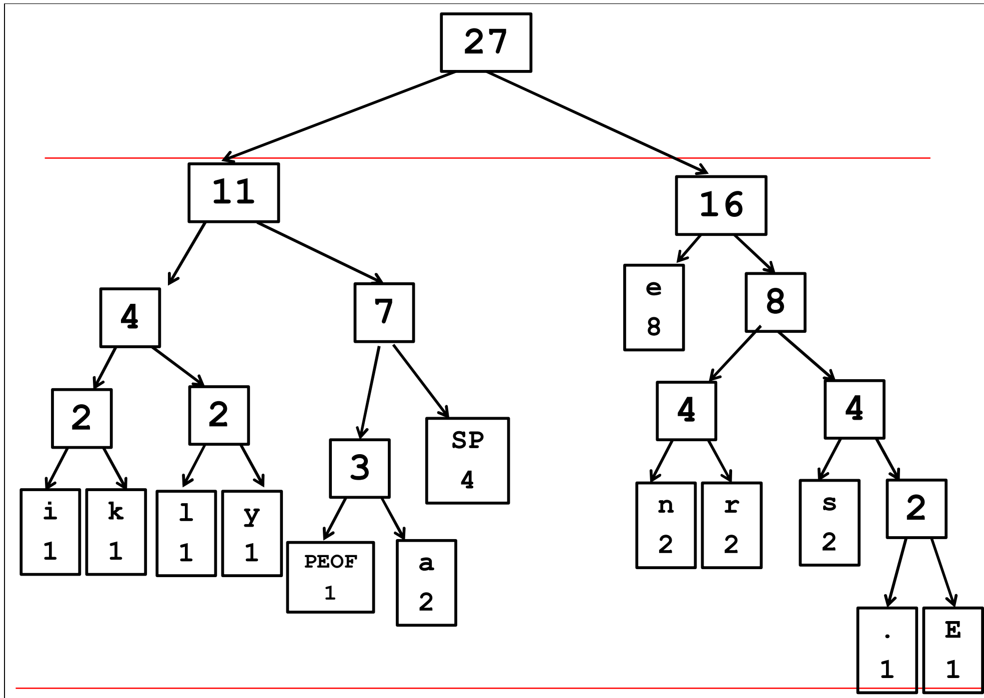












Codes

```
value: 32, equivalent char:  , frequency: 4, new code 011
value: 46, equivalent char: ., frequency: 1, new code 11110
value: 69, equivalent char: E, frequency: 1, new code 11111
value: 97, equivalent char: a, frequency: 2, new code 0101
value: 101, equivalent char: e, frequency: 8, new code 10
value: 105, equivalent char: i, frequency: 1, new code 0000
value: 107, equivalent char: k, frequency: 1, new code 0001
value: 108, equivalent char: l, frequency: 1, new code 0010
value: 110, equivalent char: n, frequency: 2, new code 1100
value: 114, equivalent char: r, frequency: 2, new code 1101
value: 115, equivalent char: s, frequency: 2, new code 1110
value: 121, equivalent char: y, frequency: 1, new code 0011
value: 256, equivalent char: ?, frequency: 1, new code 0100
```

Altering files

- Tower bit map (Eclipse/Huffman/Data). Alter the first 300 characters of line 16765 to this

```
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00~00~00~00~00
~00~00~00~00~00~00~00~00~00 xxx
```


Compression - Why Bother?

- Apostolos "Toli" Lerios
- Facebook Engineer
- Heads image storage group
- jpeg images already compressed
- look for ways to compress even more
- 1% less space = millions of dollars in savings