

# **Machine-Level Programming IV: x86-64 Procedures, Data**

# Today

- **Procedures (x86-64)**
- **Arrays**
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
- **Structures**
  - Allocation
  - Access

# x86-64 Integer Registers: Usage Conventions

|             |               |
|-------------|---------------|
| <b>%rax</b> | Return value  |
| <b>%rbx</b> | Callee saved  |
| <b>%rcx</b> | Argument #4   |
| <b>%rdx</b> | Argument #3   |
| <b>%rsi</b> | Argument #2   |
| <b>%rdi</b> | Argument #1   |
| <b>%rsp</b> | Stack pointer |
| <b>%rbp</b> | Callee saved  |

|             |              |
|-------------|--------------|
| <b>%r8</b>  | Argument #5  |
| <b>%r9</b>  | Argument #6  |
| <b>%r10</b> | Caller saved |
| <b>%r11</b> | Caller Saved |
| <b>%r12</b> | Callee saved |
| <b>%r13</b> | Callee saved |
| <b>%r14</b> | Callee saved |
| <b>%r15</b> | Callee saved |

# x86-64 Registers

- **Arguments passed to functions via registers**
  - If more than 6 integral parameters, then pass rest on stack
  - These registers can be used as caller-saved as well
- **All references to stack frame via stack pointer**
  - Eliminates need to update `%ebp/%rbp`
- **Other Registers**
  - 6 callee saved
  - 2 caller saved
  - 1 return value (also usable as caller saved)
  - 1 special (stack pointer)

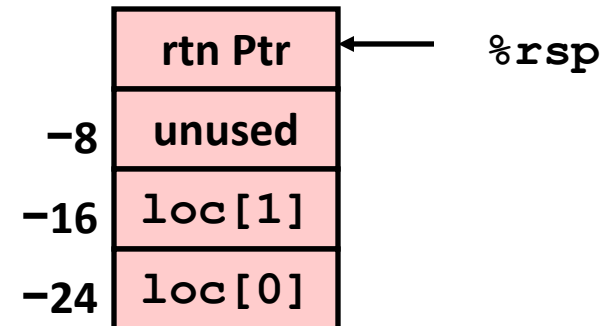
# x86-64 Locals in the Red Zone

```
/* Swap, using local array */  
void swap_a(long *xp, long *yp)  
{  
    volatile long loc[2];  
    loc[0] = *xp;  
    loc[1] = *yp;  
    *xp = loc[1];  
    *yp = loc[0];  
}
```

```
swap_a:  
    movq    (%rdi), %rax  
    movq    %rax, -24(%rsp)  
    movq    (%rsi), %rax  
    movq    %rax, -16(%rsp)  
    movq    -16(%rsp), %rax  
    movq    %rax, (%rdi)  
    movq    -24(%rsp), %rax  
    movq    %rax, (%rsi)  
    ret
```

## ■ Avoiding Stack Pointer Change

- Can hold all information within small window beyond stack pointer



# Interesting Features of Stack Frame

## ■ Allocate entire frame at once

- All stack accesses can be relative to `%rsp`
- Do by decrementing stack pointer
- Can delay allocation, since safe to temporarily use red zone

## ■ Simple deallocation

- Increment stack pointer
- No base/frame pointer needed

# Today

- Procedures (x86-64)
- **Arrays**
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
- Structures

# Basic Data Types

## ■ Integral

- Stored & operated on in general (integer) registers
- Signed vs. unsigned depends on instructions used

| <b>Intel</b> | <b>ASM</b> | <b>Bytes</b> | <b>C</b>                            |
|--------------|------------|--------------|-------------------------------------|
| byte         | <b>b</b>   | 1            | <b>[unsigned] char</b>              |
| word         | <b>w</b>   | 2            | <b>[unsigned] short</b>             |
| double word  | <b>d</b>   | 4            | <b>[unsigned] int</b>               |
| quad word    | <b>q</b>   | 8            | <b>[unsigned] long int (x86-64)</b> |

## ■ Floating Point

- Stored & operated on in floating point registers

| <b>Intel</b> | <b>ASM</b> | <b>Bytes</b> | <b>C</b>           |
|--------------|------------|--------------|--------------------|
| Single       | <b>s</b>   | 4            | <b>float</b>       |
| Double       | <b>d</b>   | 8            | <b>double</b>      |
| Extended     | <b>t</b>   | 10/12/16     | <b>long double</b> |

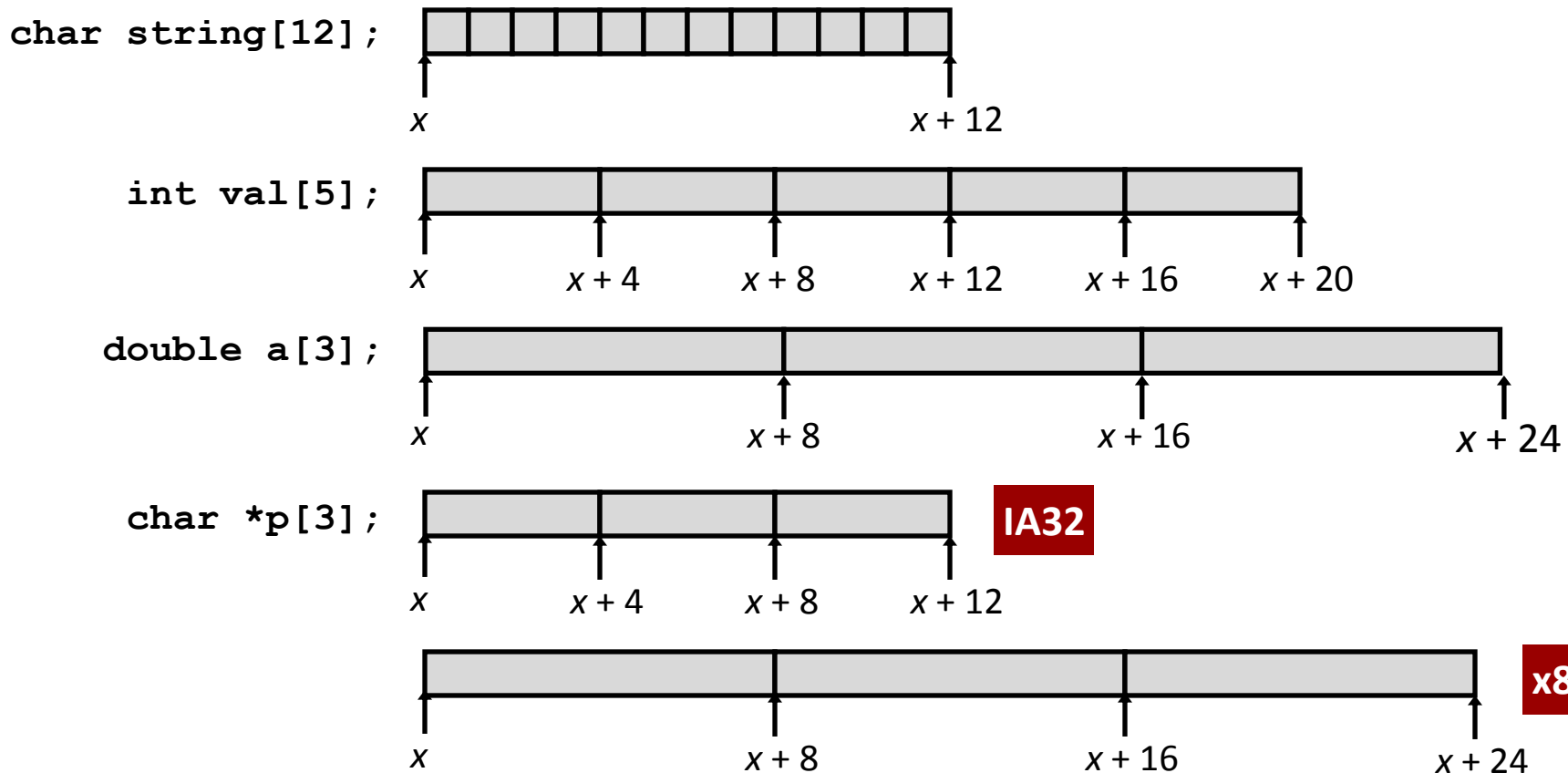


# Array Allocation

## ■ Basic Principle

$T$   $A[L]$  ;

- Array of data type  $T$  and length  $L$
- Contiguously allocated region of  $L * \text{sizeof}(T)$  bytes

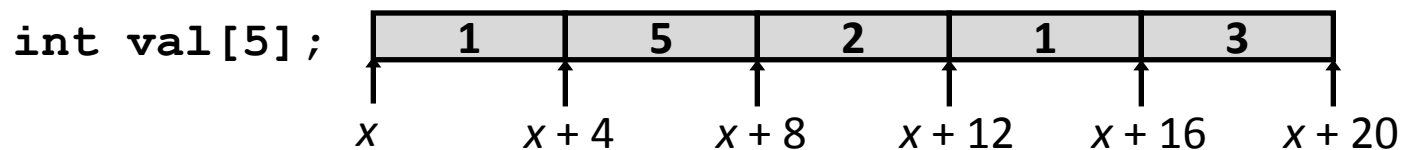


# Array Access

## ■ Basic Principle

$T$   $\mathbf{A}[L]$  ;

- Array of data type  $T$  and length  $L$
- Identifier  $\mathbf{A}$  can be used as a pointer to array element 0: Type  $T^*$

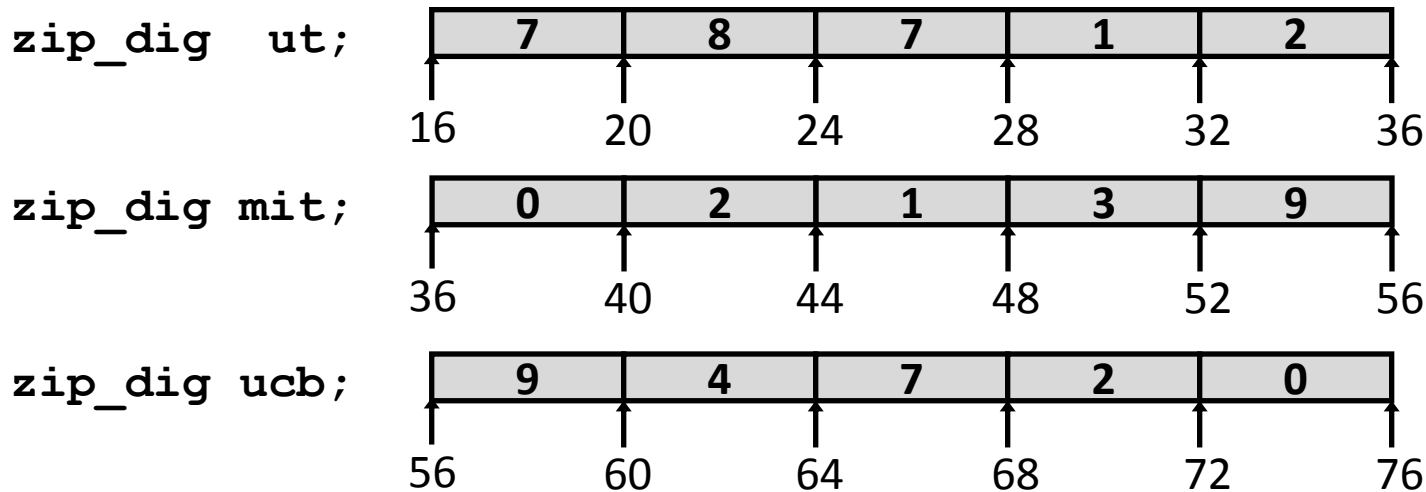


| ■ Reference              | Type               | Value  |
|--------------------------|--------------------|--------|
| <code>val[4]</code>      | <code>int</code>   | 3      |
| <code>val</code>         | <code>int *</code> | $x$    |
| <code>val+1</code>       | <code>int *</code> | $x+4$  |
| <code>&amp;val[2]</code> | <code>int *</code> | $x+8$  |
| <code>val[5]</code>      | <code>int</code>   | ??     |
| <code>*(val+1)</code>    | <code>int</code>   | 5      |
| <code>val + i</code>     | <code>int *</code> | $x+4i$ |

# Array Example

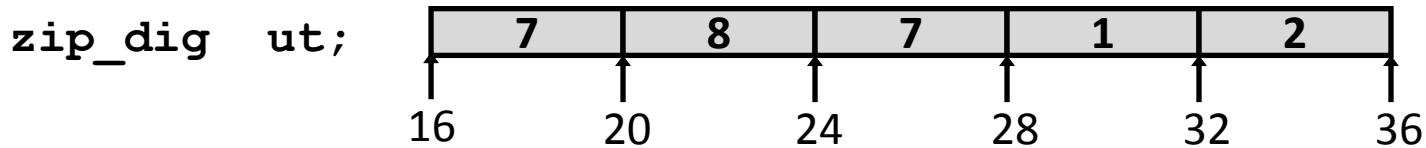
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig ut = { 7, 8, 7, 1, 2 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “`zip_dig cmu`” equivalent to “`int ut[5]`”
- Example arrays were allocated in successive 20 byte blocks
  - Not guaranteed to happen in general

# Array Accessing Example



```
int get_digit
  (zip_dig z, int dig)
{
  return z[dig];
}
```

## IA32

```
# %edx = z
# %eax = dig
movl (%edx,%eax,4),%eax # z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at  $4 * \%eax + \%edx$
- Use memory reference  $(\%edx, \%eax, 4)$

# Array Loop Example (IA32)

```
void zincr(zip_dig z) {
    int i;
    for (i = 0; i < ZLEN; i++)
        z[i]++;
}
```

```
# edx = z
movl    $0, %eax          # %eax = i
.L4:                                # loop:
addl    $1, (%edx,%eax,4)  # z[i]++
addl    $1, %eax          # i++
cmpl    $5, %eax          # i:5
jne     .L4                # if !=, goto loop
```

# Pointer Loop Example (IA32)

```
void zincr_p(zip_dig z) {
    int *zend = z+ZLEN;
    do {
        (*z)++;
        z++;
    } while (z != zend);
}
```

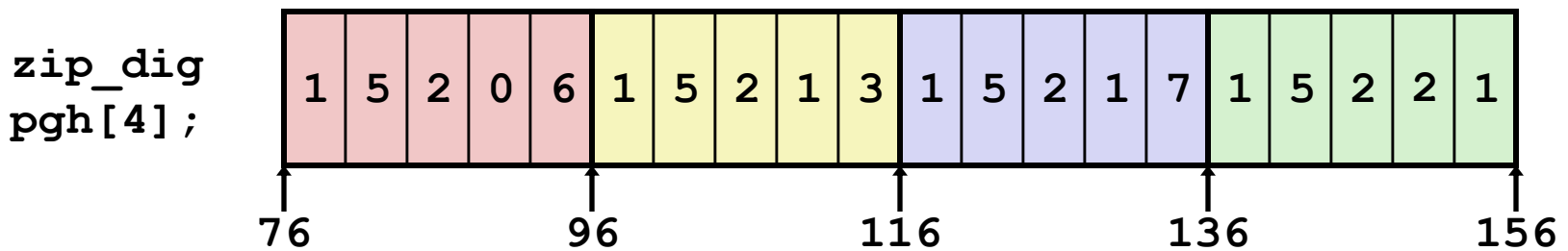


```
void zincr_v(zip_dig z) {
    void *vz = z;
    int i = 0;
    do {
        (*((int *) (vz+i)))++;
        i += ISIZE;
    } while (i != ISIZE*ZLEN);
}
```

```
# edx = z = vz
movl  $0, %eax
.L8:
addl  $1, (%edx,%eax)
addl  $4, %eax
cmpl  $20, %eax
jne   .L8
# i = 0
# loop:
# Increment vz+i
# i += 4
# Compare i:20
# if !=, goto loop
```

# Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```



- **“zip\_dig pgh[4]” equivalent to “int pgh[4][5]”**
  - Variable **pgh**: array of 4 elements, allocated contiguously
  - Each element is an array of 5 **int**'s, allocated contiguously
- **“Row-Major” ordering of all elements guaranteed**

# Multidimensional (Nested) Arrays

## ■ Declaration

$T \ A[R][C];$

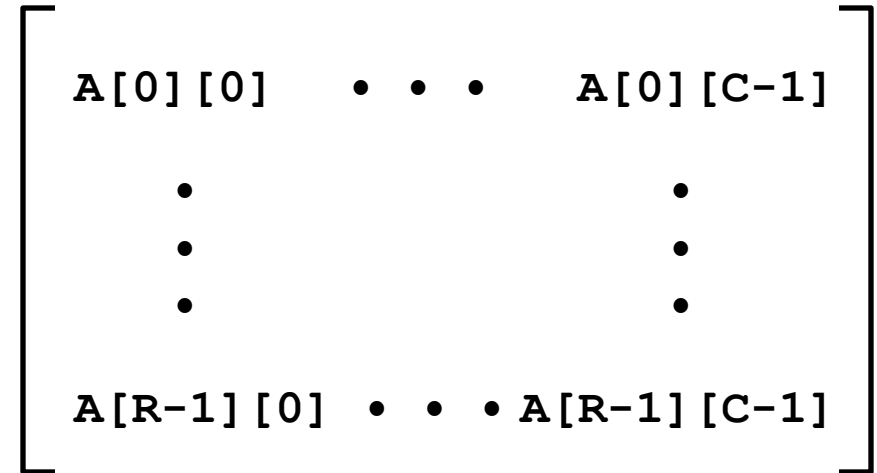
- 2D array of data type  $T$
- $R$  rows,  $C$  columns
- Type  $T$  element requires  $K$  bytes

## ■ Array Size

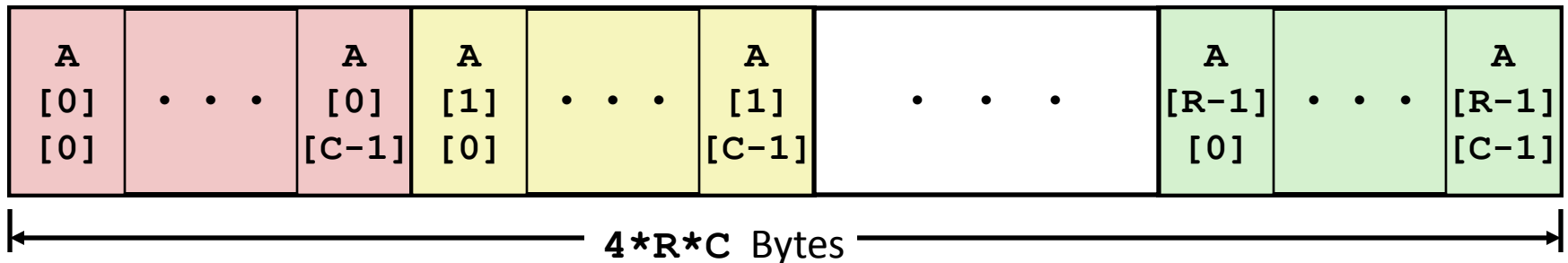
- $R * C * K$  bytes

## ■ Arrangement

- Row-Major Ordering



`int A[R][C];`



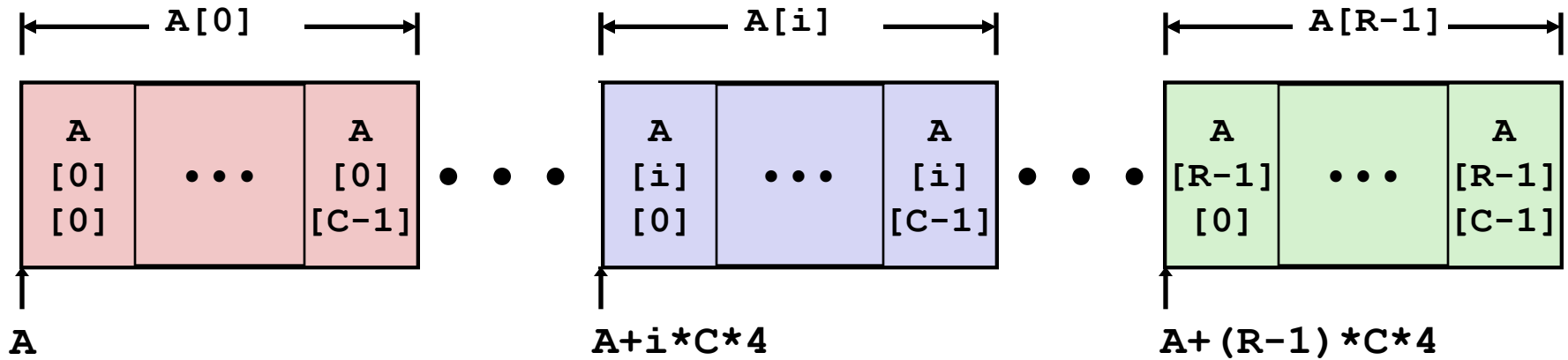


# Nested Array Row Access

## ■ Row Vectors

- $\mathbf{A}[i]$  is array of  $C$  elements
- Each element of type  $T$  requires  $K$  bytes
- Starting address  $\mathbf{A} + i * (C * K)$

```
int A[R][C];
```



# Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

```
# %eax = index
leal (%eax,%eax,4),%eax # 5 * index
leal pgh(,%eax,4),%eax # pgh + (20 * index)
```

## ■ Row Vector

- `pgh[index]` is array of 5 `int`'s
- Starting address `pgh+20*index`

## ■ IA32 Code

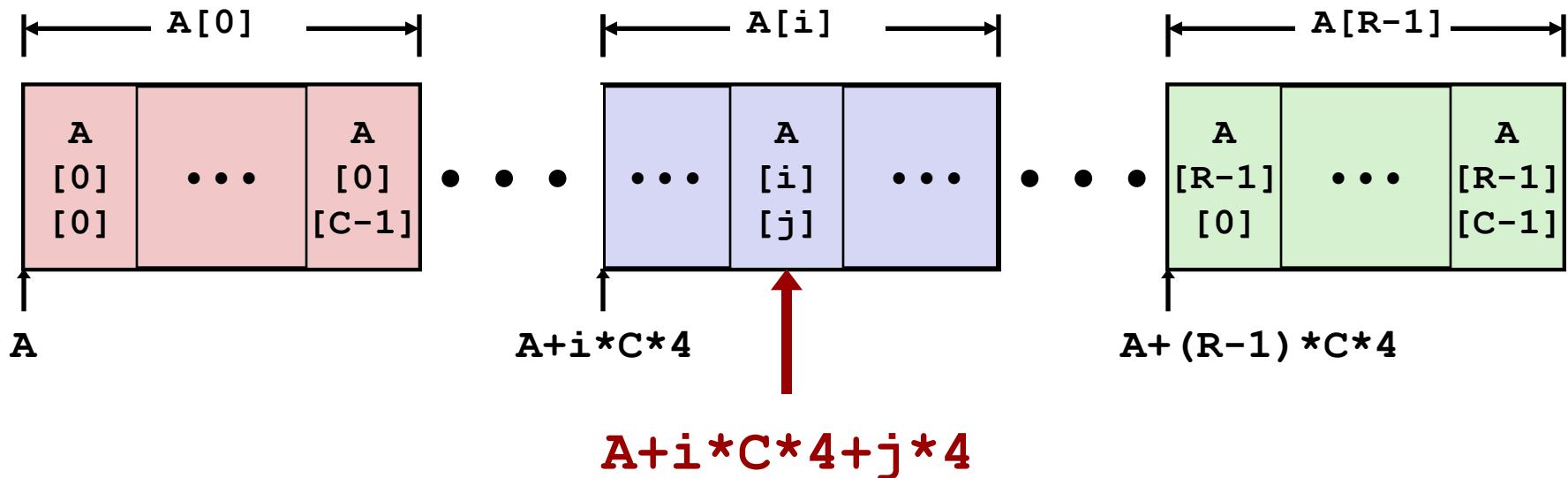
- Computes and returns address
- Compute as `pgh + 4*(index+4*index)`

# Nested Array Row Access

## ■ Array Elements

- $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
- Address  $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



# Nested Array Element Access Code

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```

```
movl    8(%ebp), %eax    # index
leal   (%eax,%eax,4), %eax    # 5*index
addl   12(%ebp), %eax    # 5*index+dig
movl   pgh(,%eax,4), %eax    # offset 4*(5*index+dig)
```

## ■ Array Elements

- `pgh[index][dig]` is `int`
- Address: `pgh + 20*index + 4*dig`
  - = `pgh + 4*(5*index + dig)`

## ■ IA32 Code

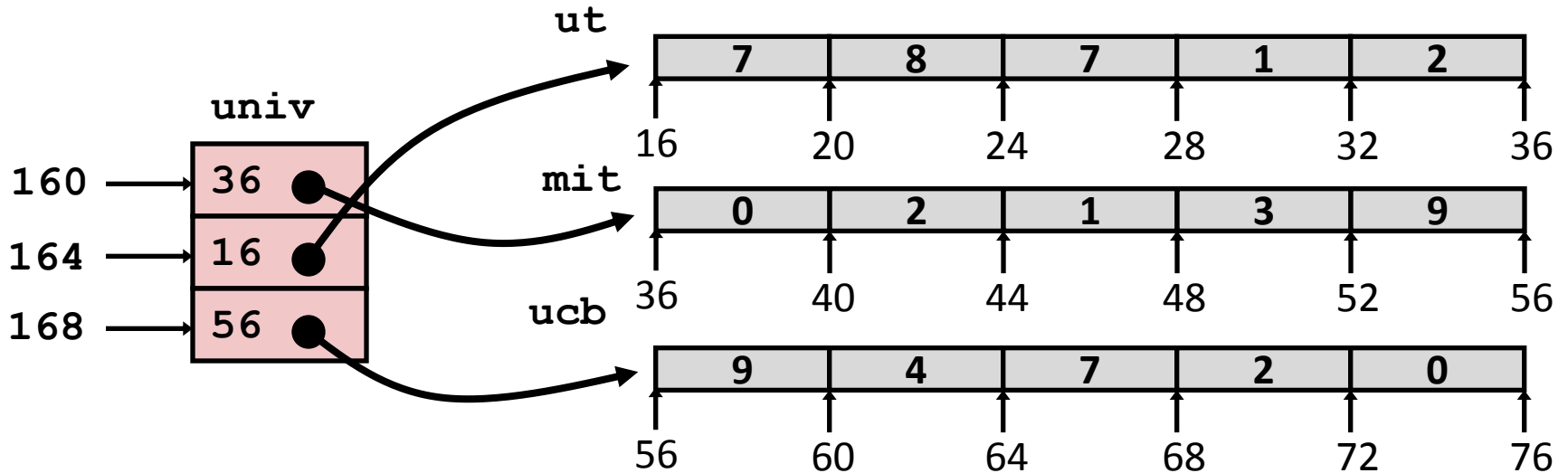
- Computes address `pgh + 4*((index+4*index)+dig)`

# Multi-Level Array Example

```
zip_dig ut = { 7, 8, 7, 1, 2 };  
zip_dig mit = { 0, 2, 1, 3, 9 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3  
int *univ[UCOUNT] = {mit, ut, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
  - 4 bytes
- Each pointer points to array of `int`'s



# Element Access in Multi-Level Array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

```
movl    8(%ebp), %eax           # index
movl    univ(,%eax,4), %edx     # p = univ[index]
movl    12(%ebp), %eax          # dig
movl    (%edx,%eax,4), %eax     # p[dig]
```

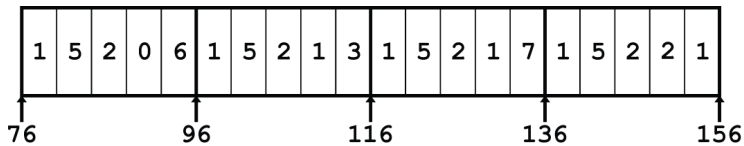
## ■ Computation (IA32)

- Element access  $\text{Mem}[\text{Mem}[\text{univ}+4*\text{index}]+4*\text{dig}]$
- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

# Array Element Accesses

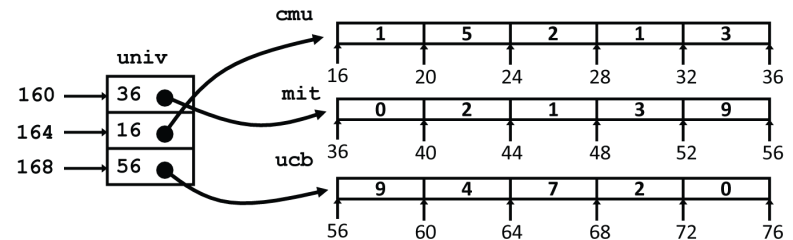
## Nested array

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```



## Multi-level array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```



Accesses looks similar in C, but addresses very different:

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{dig}]$

$\text{Mem}[\text{Mem}[\text{univ} + 4 * \text{index}] + 4 * \text{dig}]$

# N X N Matrix Code

## ■ Fixed dimensions

- Know value of N at compile time

```
#define N 16
typedef int fix_matrix[N][N];
/* Get element a[i][j] */
int fix_ele
    (fix_matrix a, int i, int j)
{
    return a[i][j];
}
```

## ■ Variable dimensions, explicit indexing

- Traditional way to implement dynamic arrays

```
#define IDX(n, i, j) ((i)*(n)+(j))
/* Get element a[i][j] */
int vec_ele
    (int n, int *a, int i, int j)
{
    return a[IDX(n,i,j)];
}
```

## ■ Variable dimensions, implicit indexing

- Now supported by gcc

```
/* Get element a[i][j] */
int var_ele
    (int n, int a[n][n], int i, int j)
{
    return a[i][j];
}
```



# 16 X 16 Matrix Access

## ■ Array Elements

- Address  $\mathbf{A} + i * (\mathbf{C} * \mathbf{K}) + j * \mathbf{K}$
- $\mathbf{C} = 16, \mathbf{K} = 4$

```
/* Get element a[i][j] */  
int fix_ele(fix_matrix a, int i, int j) {  
    return a[i][j];  
}
```

```
movl    12(%ebp), %edx    # i  
sall    $6, %edx         # i*64  
movl    16(%ebp), %eax    # j  
sall    $2, %eax         # j*4  
addl    8(%ebp), %eax     # a + j*4  
movl    (%eax,%edx), %eax # *(a + j*4 + i*64)
```

# n X n Matrix Access

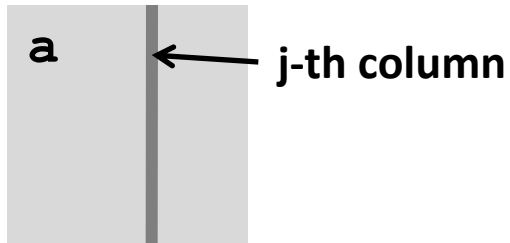
## ■ Array Elements

- Address  $\mathbf{A} + i * (\mathbf{C} * \mathbf{K}) + j * \mathbf{K}$
- $\mathbf{C} = \mathbf{n}, \mathbf{K} = 4$

```
/* Get element a[i][j] */  
int var_ele(int n, int a[n][n], int i, int j) {  
    return a[i][j];  
}
```

```
movl    8(%ebp), %eax    # n  
sall    $2, %eax        # n*4  
movl    %eax, %edx      # n*4  
imull   16(%ebp), %edx   # i*n*4  
movl    20(%ebp), %eax   # j  
sall    $2, %eax        # j*4  
addl    12(%ebp), %eax   # a + j*4  
movl    (%eax,%edx), %eax # *(a + j*4 + i*n*4)
```

# Optimizing Fixed Array Access



## ■ Computation

- Step through all elements in column j

## ■ Optimization

- Retrieving successive elements from single column

```
#define N 16
typedef int fix_matrix[N][N];
```

```
/* Retrieve column j from array */
void fix_column
(fix_matrix a, int j, int *dest)
{
    int i;
    for (i = 0; i < N; i++)
        dest[i] = a[i][j];
}
```

# Optimizing Fixed Array Access

## ■ Optimization

- Compute  $ajp = \&a[i][j]$ 
  - Initially =  $a + 4*j$
  - Increment by  $4*N$

| Register          | Value             |
|-------------------|-------------------|
| <code>%ecx</code> | <code>ajp</code>  |
| <code>%ebx</code> | <code>dest</code> |
| <code>%edx</code> | <code>i</code>    |

```
/* Retrieve column j from array */  
void fix_column  
    (fix_matrix a, int j, int *dest)  
{  
    int i;  
    for (i = 0; i < N; i++)  
        dest[i] = a[i][j];  
}
```

```
.L8:                                # loop:  
    movl    (%ecx), %eax             #   Read *ajp  
    movl    %eax, (%ebx,%edx,4)     #   Save in dest[i]  
    addl    $1, %edx                #   i++  
    addl    $64, %ecx               #   ajp += 4*N  
    cmpl   $16, %edx               #   i:N  
    jne    .L8                     #   if !=, goto loop
```

# Optimizing Variable Array Access

- Compute  $ajp = \&a[i][j]$ 
  - Initially  $= a + 4*j$
  - Increment by  $4*n$

| Register | Value |
|----------|-------|
| %ecx     | ajp   |
| %edi     | dest  |
| %edx     | i     |
| %ebx     | 4*n   |
| %esi     | n     |

```
/* Retrieve column j from array */  
void var_column  
    (int n, int a[n][n],  
     int j, int *dest)  
{  
    int i;  
    for (i = 0; i < n; i++)  
        dest[i] = a[i][j];  
}
```

```
.L18:                                # loop:  
    movl    (%ecx), %eax              #   Read *ajp  
    movl    %eax, (%edi,%edx,4)      #   Save in dest[i]  
    addl    $1, %edx                 #   i++  
    addl    $ebx, %ecx               #   ajp += 4*n  
    cmpl   $edx, %esi               #   n:i  
    jg     .L18                      #   if >, goto loop
```

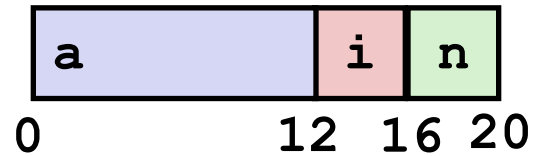
# Today

- Procedures (x86-64)
- Arrays
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
- **Structures**
  - Allocation
  - Access

# Structure Allocation

```
struct rec {  
    int a[3];  
    int i;  
    struct rec *n;  
};
```

## Memory Layout

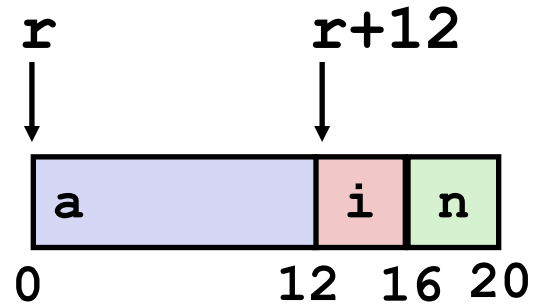


## ■ Concept

- Contiguously-allocated region of memory
- Refer to members within structure by names
- Members may be of different types

# Structure Access

```
struct rec {  
    int a[3];  
    int i;  
    struct rec *n;  
};
```



## ■ Accessing Structure Member

- Pointer indicates first byte of structure
- Access elements with offsets

```
void  
set_i(struct rec *r,  
      int val)  
{  
    r->i = val;  
}
```

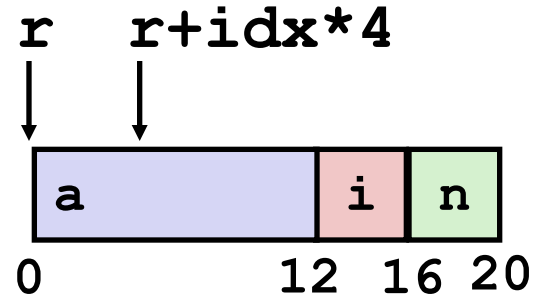
## IA32 Assembly

```
# %edx = val  
# %eax = r  
movl %edx, 12(%eax) # Mem[r+12] = val
```



# Generating Pointer to Structure Member

```
struct rec {  
    int a[3];  
    int i;  
    struct rec *n;  
};
```



## ■ Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Arguments
  - Mem[%ebp+8]: **r**
  - Mem[%ebp+12]: **idx**

```
int *get_ap  
(struct rec *r, int idx)  
{  
    return &r->a[idx];  
}
```

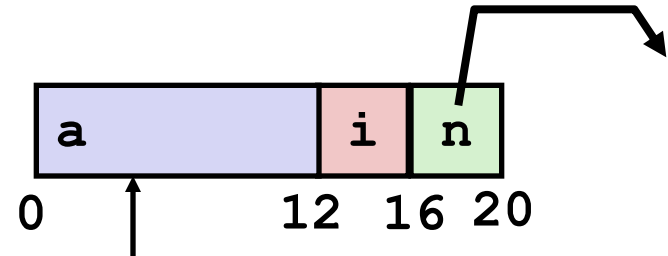
```
movl    12(%ebp), %eax    # Get idx  
sall    $2, %eax         # idx*4  
addl    8(%ebp), %eax     # r+idx*4
```

# Following Linked List

## ■ C Code

```
void set_val
(struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->n;
    }
}
```

```
struct rec {
    int a[3];
    int i;
    struct rec *n;
};
```



Element i

| Register | Value |
|----------|-------|
| %edx     | r     |
| %ecx     | val   |

```
.L17:                # loop:
    movl    12(%edx), %eax    # r->i
    movl    %ecx, (%edx,%eax,4) # r->a[i] = val
    movl    16(%edx), %edx    # r = r->n
    testl   %edx, %edx       # Test r
    jne     .L17             # If != 0 goto loop
```

# Summary

## ■ Procedures in x86-64

- Stack frame is relative to stack pointer
- Parameters passed in registers

## ■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

## ■ Structures

- Allocation
- Access