Monitors

CS439: Principles of Computer Systems
February 10, 2016
Last Time

• Hardware Support for Mutual Exclusion and Synchronization
  – read-modify-write instructions
• Mutual Exclusion and Synchronization in Software
  – Abstractions built on top of hardware support
  – Locks
  – Semaphores
Semaphore Recap

• Two types: binary and counted
• Used for locking and synchronization
• Two atomic operations:
  – down(): wait for semaphore to be available, decrement value
  – up(): signal waiting threads semaphore is available, increment value
• Need separate semaphores for locking and synchronizing
Semaphore Operations

Semaphore::Down()
{
*if value is == 0, sleep
 *on signal, wake up
*decrement semaphore value
*return
}

Semaphore::Up{
 *increment semaphore value
*if any threads are sleeping
 on semaphore, wake one of them up
*return
}
Today’s Agenda

• More Synchronization in Software
  – Monitors
    • Locks and Condition Variables
    • Bounded Buffer problem

• Readers/Writers problem
Monitors
Introducing Monitors

• Monitors guarantee mutual exclusion
  – only one thread may execute a given monitor method at a time
• First introduced as a programming language construct (Mesa, Java)
• Now also define a programming convention and can be used in any language (C, C++, ...)
• Object-oriented style:
  – Collect related shared data into an object/module
    • Associate a lock with each one
    • All data is private
  – Define methods for accessing the shared data
    • These are critical sections
Monitors, Formally

A monitor defines a lock and zero or more condition variables for managing concurrent access to shared data.

- uses the lock to ensure that only a single thread is active in the monitor at any point
- the lock also provides mutual exclusion for shared data
- Condition variables enable threads to block waiting for an event inside of critical sections
  * release the lock at the same time the thread is put to sleep
Monitor Design

• Goal is to encapsulate shared data
  – Class in C++ or Java
  – Struct or File in C (logical encapsulation)

• Goal is to provide for mutual exclusion
  – Has a lock (exactly one!)

• Goal is to allow for synchronization
  – Has condition variables

• Goal is to allow operations on the shared data
  – Has functions or methods
Monitor Functions: Implementation

• Acquire the lock at the start of every function (first thing!)
  – Operate on the shared data
  – Temporarily release the lock if they can’t complete due to missing resource (use condition variable for this)
  – Reacquire the lock when they can continue (again, condition variable!)
  – Operate on the shared data

• Release the lock at the end
Semaphore Example:
Producers/Consumers (Recall)

Semaphore mutex = 1    //access to buffer
Semaphore empty = N    //count of empty slots
Semaphore full = 0     //count of full slots
int buffer[N]

BoundedBuffer::Producer()
    <produce item>
    empty->down()    //get empty spot
    mutex->down()    //get access to buffer
    <add item to buffer>
    mutex->up()      //release buffer
    full->up()       //another item in buffer
}

BoundedBuffer::Consumer()
    full->down()    //get item
    mutex->down()   //get access to buffer
    <remove item from buffer>
    mutex->up()     //release buffer
    empty->up()     //another empty slot
    <use item>
}
BoundedBuffer::Producer()
{
    //produce item
    empty->down() //get empty spot
    mutex->down() //get access to buffer

    //add item to buffer
    mutex->up() //release buffer
    full->up() //another item in buffer
}
Condition Variables

- Enable threads to wait efficiently for changes to shared state protected by a lock
- Each one is a queue of waiting threads (no state!)
- Enable the thread to sleep inside a critical section by *atomically* releasing the Lock at the same time the thread is put to sleep

**Rule:** A thread *must* hold the lock when doing condition variable operations
Condition Variable Operations

1. Wait(Lock lock)
   – atomic (release lock, move thread to waiting queue, suspend thread)
   – when the thread wakes up it re-acquires the lock (before returning from wait)
   – thread will always block

2. Signal(Lock lock)
   – wake up waiting thread, if one exists. Otherwise, it’s a no-op

3. Broadcast(Lock lock)
   – wake up all waiting threads, if any exist. Otherwise, it’s a no-op
Monitor Operations

Lock->Acquire()  //acquires lock, when returns, thread has lock
Lock->Release()  //releases lock

CondVar::Wait(lock){
    *move thread to waiting queue and suspend thread
    *release lock
    *on signal, wake up, re-acquire lock
    *return
}

CondVar::Signal(){
    *wake up a thread waiting on condVar
    *return
}

CondVar::Broadcast(){
    *wake up ALL threads waiting on condVar
    *return
}
Resource Variables

• Conditions variables (unlike semaphores) keep no state
• Each condition variable should have a resource variable that tracks the state of that resource
  – You must maintain this variable
• Check the resource variable before calling wait on the associated condition variable to ensure the resource really isn’t available
• Once the resource is available, claim it (subtract the amount you are using!)
• Before signaling that you are finished with a resource, indicate the resource is available by increasing the resource variable
Monitor Example: Items in a Queue

Lock lock;
Condition fullCV;
int queue_size;

void Add(item){
    lock->acquire()
    put item on queue;
    queue_size++;
    fullCV->signal(lock)
    lock->release()
}

void remove(){
    lock->acquire()
    while(queue_size <=0)
        fullCV->wait(lock);
    queue_size--;
    remove item;
    lock->release();
}
Signal() Semantics

Which thread executes once signal() is called?

– If there are no waiting threads, the signaler continues and the signal is effectively lost
– If there is a waiting thread (or two):
  • There are at least two ready threads: the one that called signal() and the one that was (or will be) awakened
  • Exactly one of the threads can execute or we will have more than one thread active in the monitor (violates mutual exclusion!)
  • So which thread gets to execute?
Whose turn is it?

**Mesa/Hansen Style**
- The thread that signals keeps the lock (and thus the processor)
- The waiting thread waits for the lock
  - Signal is only a hint that the condition may be true: shared state may have changed!
  - Adding signals affects performance, but never safety
- Implemented in Java and most real operating systems

**Hoare Style**
- The thread that signals gives up the lock and the waiting thread gets the lock
  - Signaling is atomic with the resumption of the waiting thread
  - Shared state cannot change before waiting thread is resumed
- When the thread that was waiting and is now executing exits or waits again, it releases the lock back to the signaling thread
- Implemented in most textbooks (not yours!)
More on Turns

• With Mesa-style, waiting thread may need to wait again after it is wakened (Why?), so while is important
• With Hoare-style, we can change while to if because a waiting thread runs immediately

```java
public void remove(){
    lock->Acquire()
    while (queue_size <= 0){
        full->wait(lock);
    }
    queue_size--;
    remove item;
    lock->Release();
}
```

Regardless, if you assume there is NO atomicity between signal() and the return from wait(), your code will always work.
Semaphore Example: Producers/Consumers (Recall)

Semaphore mutex = 1  //access to buffer
Semaphore empty = N  //count of empty slots
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int buffer[N]

BoundedBuffer::Producer()
{
  <produce item>
  empty->down() //get empty spot
  mutex->down() //get access to buffer

  <add item to buffer>

  mutex->up() //release buffer
  full->up() //another item in buffer
}

BoundedBuffer::Consumer()
{
  full->down()  //get item
  mutex->down() //get access to buffer

  <remove item from buffer>

  mutex->up() //release buffer
  empty->up() //another empty slot
  <use item>
}
Semaphore Example: Producers/Consumers (Recall)

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int buffer[N]

BoundedBuffer::Consumer()
{
    full->down()    //get item
    mutex->down()   //get access to buffer

    <remove item from buffer>

    mutex->up()     //release buffer
    empty->up()     //another empty slot
    <use item>
}


Monitor Example: Producers/Consumers

class BBMonitor{
    public: <methods>
    private: item buffer[N];
        Lock lock;
        Condition fullCV, emptyCV;
        int empty=N;
        int full=0;
    }

BoundedBuffer::Producer(){
    lock->Acquire()

    while(empty == 0)
        emptyCV->Wait(lock) //get empty spot
    empty--;

    <add item to buffer>

    full += 1
    fullCV->Signal(lock) //another item in buffer
    lock->Release()
}

BoundedBuffer::Consumer(){
    lock->Acquire()

    while(full == 0)
        fullCV->Wait(lock) //get item
    full--

    <remove item from buffer>

    empty++;
    emptyCV->Signal(lock) //another empty slot
    lock->Release()
}
Every monitor function should begin with what command?

A. Wait()
B. Signal()
C. Lock->Acquire()
D. Lock->Release()
E. Broadcast()
Comparing Monitors and Semaphores

• Condition variables do not have any history
  – on signal() if no one is waiting, the signal is a no-op
  – if thread then calls condition->wait(), it waits.

• Semaphores do have history
  – on up() if no one is waiting, the value of the semaphore is incremented
  – if a thread then calls semaphore->down(), the value is decremented and the thread continues
So... signal() and down()

• In semaphores, down() and up() are commutative
  – result is the same regardless of the order of execution

• Condition variables are not commutative
  – so they must be in a critical section to access state variables and do their job

You can implement Monitors with Semaphores.
## Monitors and Semaphores: Recap

<table>
<thead>
<tr>
<th>Both</th>
<th>Semaphores</th>
<th>Monitors</th>
</tr>
</thead>
<tbody>
<tr>
<td>• Provide mutual exclusion and synchronization</td>
<td>• Semaphores are basically generalized locks</td>
<td>• Consist of a lock and one or more condition variables</td>
</tr>
<tr>
<td>• Have signal() and wait()</td>
<td>• Binary and counting semaphores</td>
<td>• Encapsulate shared data</td>
</tr>
<tr>
<td>• Support a queue of threads that are waiting to access a critical section (e.g., to buy milk)</td>
<td>• Used for mutual exclusion and synchronization</td>
<td>• Use locks for mutual exclusion</td>
</tr>
<tr>
<td>• No busy waiting!</td>
<td></td>
<td>• Use condition variables for synchronization</td>
</tr>
</tbody>
</table>

No busy waiting!
A Different Type of Problem

• We’ve looked at problems where we protect shared data by only allowing one thread in the critical section at a time

• Is this always appropriate? When might we want to let more threads access shared data at once?
Readers/Writers Problem

• Data is shared among several threads
  – Some only read
  – Some only write
• To get correct results, we allow multiple readers at a time, but only one writer at a time
• How can we control access to the object to permit this protocol?
Correctness Criteria

• Each read or write of the shared data must happen within a critical section
• Guarantee mutual exclusion for writers
• Allow multiple readers to execute in the critical section at once
• Allow one writer (and no readers) to execute in the critical section at once
Readers and Writers: Monitor Solution

• What methods do we need?
• How many locks?
• How many condition variables?
• What should we name them?
• Any other variables?

• Assume we’re going to say <read> and <write> for accesses to the shared data.
Readers and Writers: Monitor Solution

Variables:

read()

write()

Is our solution fair?

A. Yes
B. No, favors readers
C. No, favors writers
Understanding Our Solution

It works, but it favors readers over writers

– Any reader blocks all writers
– All readers must finish before a writer can start
– Last reader will wake any writer, but a writer wakes all readers and writers
– If a writer exits and a reader goes next, then all readers that are waiting will get through
Readers and Writers: Monitor Solution

Variables:

write()

read()
Alternative Semantics

• It may be that you would like a writer to enter its critical section as soon as possible.
• How could we implement that?
Signal vs. Broadcast

• It is always safe to use broadcast() instead of signal()
  – only performance is affected

• signal() is preferable when
  – at most one waiting thread can make progress
  – any thread waiting on the condition variable can make progress

• broadcast() is preferable when
  – multiple waiting threads may be able to make progress
  – the same condition variable is used for multiple predicates
    • some waiting threads can make progress, others can’t
Summary

• A monitor wraps operations with a lock
• Condition variables release lock temporarily
• Monitors do not keep state---a call to wait() will *always* wait
• Monitors can be implemented by following the monitor rules for acquiring and releasing locks
Announcements

• Homework 3 due Friday in section
• Project 0 due 11:59p Friday
• Project 1 is due 2/26
  – Discussion Section this week will provide an introduction to Pintos. Read the documentation first!
• Exam 1 is in TWO weeks!
  – Wednesday, 2/24! 7p! WEL 2.224!