

Algebraic Splines and Analysis - I : Lecture 2

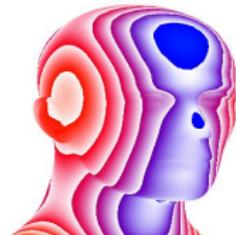
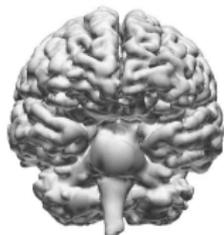
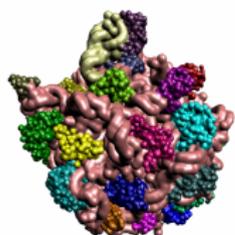
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Algebraic Curve, Surface Splines

We shall consider the modeling of domains and function fields using algebraic splines



Algebraic Splines are a complex of piecewise :

algebraic plane & space curves

algebraic surfaces



Algebraic Plane curves

- An algebraic plane curve in implicit form is a hyperelement of dimension 1 in R^2 :

$$f(x, y) = 0 \quad (1)$$

- An algebraic plane curve in parametric form is an algebraic variety of dimension 1 in R^3 . It is also a rational mapping from R^1 into R^2 .

$$x = f_1(s)/f_3(s) \quad (2)$$

$$y = f_2(s)/f_3(s) \quad (3)$$



Algebraic Space curves

- An algebraic space curve can be implicitly defined as the intersection of two surfaces given in implicit form:

$$f_1(x, y, z) = 0 \quad f_2(x, y, z) = 0 \quad (4)$$

- or alternatively as the intersection of two surfaces given in parametric form:

$$(x = f_{1,1}(s_1, t_1), y = f_{2,1}(s_1, t_1), z = f_{3,1}(s_1, t_1)) \quad (5)$$

$$(x = f_{1,2}(s_2, t_2), y = f_{2,2}(s_2, t_2), z = f_{3,2}(s_2, t_2)) \quad (6)$$

where all the $f_{i,j}$ are rational functions in s_i, t_j

- Rational algebraic space curves can also be represented as:

$$x = f_1(s), y = f_2(s), z = f_3(s) \quad (7)$$

where the f_i are rational functions in s .



Parameterization of algebraic curves

Theorem An algebraic curve P is rational iff the $\text{Genus}(P) = 0$.

The proof is classical, though non-trivial. See also, Abhyankar's *Algebraic Geometry for Scientists & Engineers AMS Publications, (1990)*

Constructive proof, genus computation, and parameterization algorithm is available from:

Automatic Parameterization of Rational Curves and Surfaces III :
Algebraic Plane Curves *Computer Aided Geometric Design, (1988)*



For algebraic space curves C given as intersection of two algebraic surfaces there exists a birational correspondence between points of C and points of a plane curve P .

The genus of C is same as the genus of P .

Hence C is rational iff $Genus(P) = 0$.

Algorithm :

- Construct a birationally equivalent plane curve P from C
- Generate a rational parametrization for P
- Construct a rational surface S containing C .

Automatic Parameterization of Rational Curves and Surfaces IV :
Algebraic Space Curves *ACM Transactions on Graphics*, (1989)



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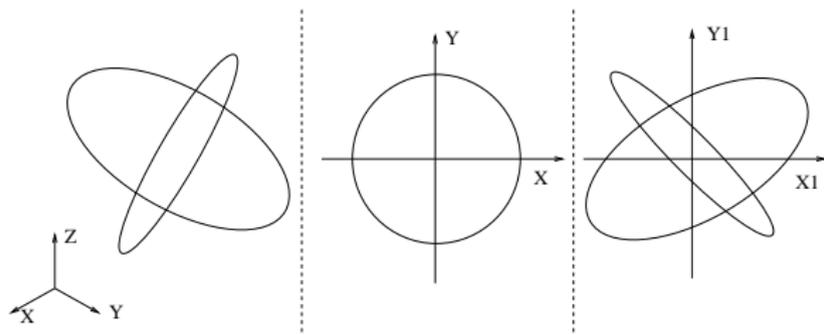
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Given: Irreducible space curve $C = (f = 0 \cap g = 0)$, and f, g not tangent along C .

Compute: Project C to an irreducible plane curve P , properly, to yield a birational map from P to C .



- 1 Space curve C as intersection of two axis aligned cylinders

$$C : (f = z^2 + x^2 - 1 \cap g = z^2 + y^2 - 1) \tag{8}$$

- 2 Badly chosen projection direction results in P not birationally related to C

$$P : (x^2 + z^2 - 1)^2 = 0 \tag{9}$$

- 3 Birationally equivalent plane curve P with properly chosen projection direction

$$P : (8y_1^2 - 4x_1y_1 + 5x_1^2 - 9)(8y_1^2 + 12x_1y_1 + 5x_1^2 - 1) = 0 \tag{10}$$

Projection can be computed using Elimination Theory. One way to eliminate a variable from two polynomials, is via Sylvester's polynomial resultant:

Given two polynomials

$$f(x) = a_m x^m + a_{m-1} x^{m-1} \dots a_0 \quad (11)$$

$$g(x) = b_n x^n + b_{n-1} x^{n-1} \dots b_0 \quad (12)$$

The Sylvester resultant matrix is constructed by rows of coefficients of f , shifted, followed by rows of coefficients of g , shifted.

To project along the z axis, write both equation as just polynomials in z , construct the matrix of coefficients in x, y , and the Sylvester resultant (projection) is the determinant.

Of course, the z axis may not be a proper projection direction. Hence first choose a valid transformation, to enable the projection to yield a rational (inverse) map.



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Choosing a valid projection direction:

Consider a general linear transformation to apply to f, g :

$$x = a_1x_1 + b_1y_1 + c_1z_1, \quad y = a_2x_1 + b_2y_1 + c_2z_1, \quad z = a_3x_1 + b_3y_1 + c_3z_1 \quad (13)$$

On substituting, we obtain the transformed equations

$$f_1(x_1, y_1, z_1) = 0, \quad g_1(x_1, y_1, z_1) = 0$$

Compute Resultant $h(x_1, y_1)$ eliminating z_1 to yield the projected plane curve $P: h = 0$.

To obtain a birational inverse map $z_1 = H(x_1, y_1)$, which exists when the projection degree is 1, we need to satisfy:

- Determinant of linear transformation to be nonzero
- Equation h of projected plane curve P is not a power of an irreducible polynomial.

A random choice of coefficients for the linear transformation, works with high probability.



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Concluding remarks

We require surfaces f, g are not tangent along C .

Birational map construction can be used for reducible space curves as well.

Irreducible space curves defined by more than two surfaces are difficult to handle outside of ideal-theoretic methods.



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Algebraic surfaces

- An algebraic surface in implicit form is a hyperplane of dimension 2 in R^3 :

$$f(x, y, z) = 0 \quad (14)$$

- An algebraic surface in parametric form is an algebraic variety of dimension 2 in R^3 . It is also a rational mapping from R^2 into R^3 .

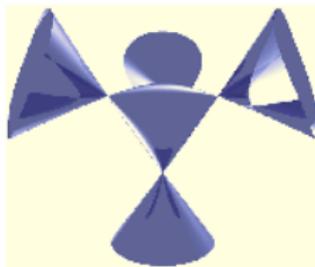
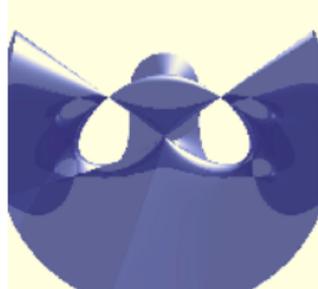
$$x = f_1(s, t)/f_4(s, t) \quad (15)$$

$$y = f_2(s, t)/f_4(s, t) \quad (16)$$

$$z = f_3(s, t)/f_4(s, t) \quad (17)$$



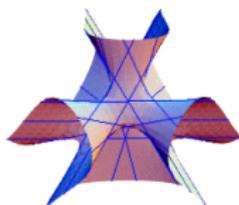
Example Algebraic Surfaces



The Clebsch Diagonal Cubic

$$81 * x^3 + 81 * y^3 + 81 * z^3 - 189 * x^2 * y - 189 * x^2 * z - 189 * y^2 * x - 189 * y^2 * z - 189 * z^2 * x - 189 * z^2 * y + 54 * x * y * z + 126 * x * y + 126 * x * z + 126 * y * z - 9 * x^2 - 9 * y^2 - 9 * z^2 - 9 * x - 9 * y - 9 * z + 1$$

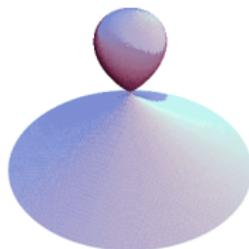
(27 real lines with 10 triple points)



The Cayley Cubic

$$-5 * x^2 * y - 5 * x^2 * z - 5 * y^2 * x - 5 * y^2 * z - 5 * z^2 * y - 5 * z^2 * x + 2 * x * y + 2 * x * z + 2 * y * z$$

(9 real lines = 6 connecting 4 double points, and 3 in a coplanar config)



The Ding-Dong Surface

$$x^2 + y^2 - (1 - z) * z^2$$



Cubic Algebraic Surfaces: Historical Gossip Column!

[1849 Cayley, Salmon] Exactly 27 straight lines on a general cubic surface

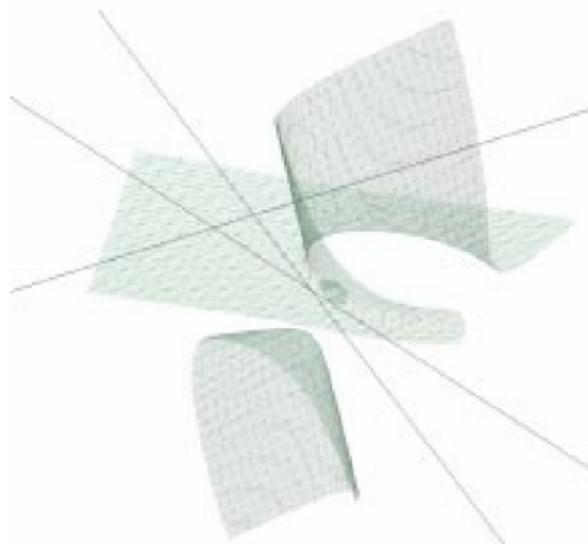
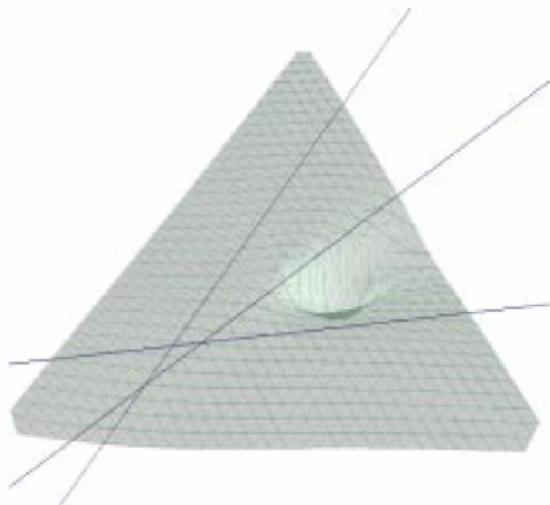
[1856 Steiner] The nine straight lines in which the surfaces of two arbitrarily given trihedra intersect each other determine together with one given point, a cubic surface.

[1858, 1863 Schläfli] classifies cubic surfaces into 23 species with respect to the number of real straight lines and tri-tangent planes on them

[1866 Cremona] establishes connections between the 27 lines on a cubic surface and Pascals Mystic hexagram:- If a hexagon is inscribed in any conic section, then the points where opposite sides meet are collinear.



45 Tri-Tangents on Smooth Cubic Surfaces



Why are the 27 lines useful to geometric modeling ?

Given two skew lines on the cubic surface $f(x, y, z) = 0$

$$l_1(u) = \begin{bmatrix} x_1(u) \\ y_1(u) \\ z_1(u) \end{bmatrix} \text{ and } l_2(u) = \begin{bmatrix} x_2(u) \\ y_2(u) \\ z_2(u) \end{bmatrix}$$

One can derive the following surface parameterization :

$$P(u, v) = \begin{bmatrix} x(u, v) \\ y(u, v) \\ z(u, v) \end{bmatrix} = \frac{al_1 + bl_2}{a + b} = \frac{a(u, v)l_1(u) + b(u, v)l_2(v)}{a(u, v) + b(u, v)}$$

where

$$a = a(u, v) = \nabla f(l_2(v)) \cdot [l_1(u) - l_2(v)]$$

$$b = b(u, v) = \nabla f(l_1(v)) \cdot [l_1(u) - l_2(v)]$$



Algorithm for Computing the 27 Lines

$$f(x, y, z) = Ax^3 + By^3 + Cz^3 + Dx^2y + Ex^2z + Fxy^2 + Gy^2z + Hxz^2 + Iyz^2 + Jxyz + kx^2 + Ly^2 + Mz^2 + Nxy + Oxz + Pyz + Qx + Ry + Sz + T = 0$$

Through intersection with tangent planes, one can reduce this to

$$\hat{f}_2(\hat{x}, \hat{y}) + \hat{g}_3(\hat{x}, \hat{y}) = 0$$

With a generic parameterization of the singular tangent cubics, one derives a polynomial $P_{81}(t)$ of degree 81.



Properties of the polynomial $P_{81}(t)$

Theorem The polynomial $P_{81}(t)$ obtained by taking the resultant of \hat{f}_2 and \hat{g}_3 factors as $P_{81}(t) = P_{27}(t)[P_3(t)]^6[P_6(t)]^6$, where $P_3(t)$, and $P_6(t)$ are degree 3 and 6 respectively.

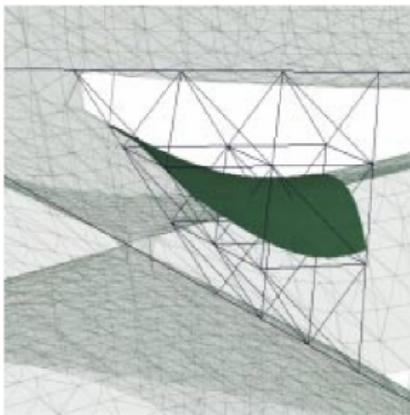
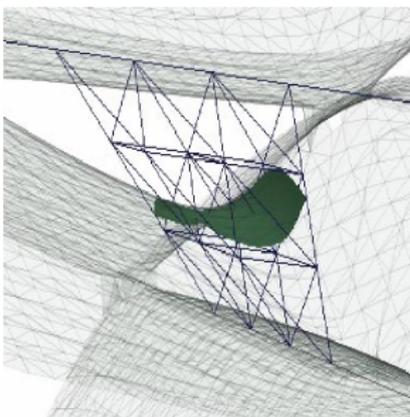
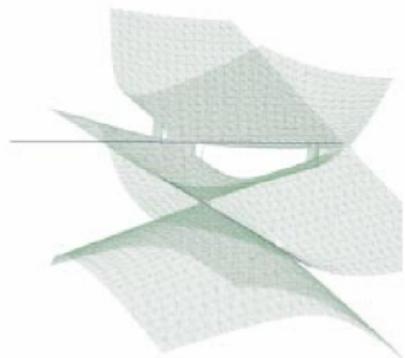
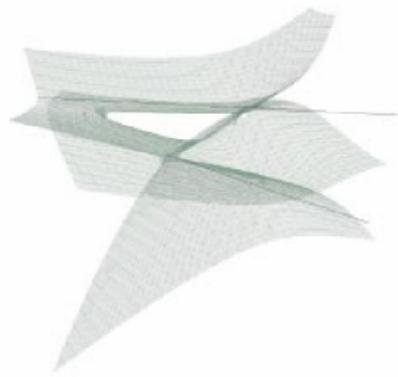
Theorem Simple real roots of $P_{27}(t) = 0$ correspond to real lines on the surface.

Proof and algorithm details available from

Rational parameterizations of non singular cubic surfaces *ACM Transactions on Graphics*, (1998)



Some Examples



Parameterization of algebraic surfaces

Theorem An algebraic surface S is rational iff the Arithmetic Genus(S)= Second Pluri-Genus (S) = 0.

The proof is attributed to Castelnuovo. See, Zariski's Algebraic Surfaces *Ergeb. Math. , Springer, (1935)*

Several examples of well known rational algebraic surfaces include: Cubic, Del Pezzo, Hirzebruch, Veronese, Steiner, etc.



What if the Algebraic Curve and/or Surface is Not Rational ?

Answer: Construct Rational Spline Approximations for a piecewise parameterization!



Rational Spline Approximation of Algebraic Plane Curves

Input: Given a real algebraic curve \mathbf{C} of degree d , a bounding box B , a finite precision real number ϵ and integers m, n with $m + n \leq d$. The curve \mathbf{C} within the bounding box B is denoted as \mathbf{C}_B .

Output: A C^{-1} , C^0 or C^1 continuous piecewise rational ϵ -approximation of all portions of \mathbf{C} within the given bounding box B , with each rational function $\frac{P_i}{Q_i}$ of degree $P_i \leq m$ and degree $Q_i \leq n$ and $m + n \leq d$.

Piecewise Rational Approximation of Real Algebraic Curves *Journal of Computational Mathematics*, (1997)



Rational Spline Approximation of $(x^2 + y^2)^3 - 4x^2y^2 = 0$ in Ganith

The screenshot shows the Ganith Algebraic Geometry Toolkit interface. The main window displays the equation $(x^2 + y^2)^3 - 4x^2y^2 = 0$ and its rational spline approximation. The approximation is shown in five windows, illustrating the process of approximating the curve with a series of line segments.

Command Window Content:

```

pieceRational2d(
(x^2+y^2)^3 -
4*x^2*y^2, 2, 3, 3,
-2.2,-2.2,0.05,1)
Interval:
[5^211,-0,058856672
9820937029+5^2+0,18
9875159296105019+5^
+5^2-0,095667091458
56986)
Created object 2...
Created object 3...

```



2. Algorithm

- Compute the intersections, the singular points S and the x -extreme points T of \mathbf{C}_B .
- Compute Newton factorization (via Hensel lifting) for each (x_i, y_i) in S and obtain a power series representation for each analytic branch of \mathbf{C} at (x_i, y_i) given by

$$\begin{cases} X(s) = x_i + s^{k_i} \\ Y(s) = \sum_{j=0}^{\infty} c_j^{(i)} s^j, \quad c_0^{(i)} = y_i \end{cases} \quad (18)$$

or

$$\begin{cases} Y(s) = y_i + s^{k_i} \\ X(s) = \sum_{j=0}^{\infty} \tilde{c}_j^{(i)} s^j, \quad \tilde{c}_0^{(i)} = x_i \end{cases} \quad (19)$$



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3. Algorithm Contd.

- Compute $\frac{P_{mn}(s)}{Q_{mn}(s)}$ the (m, n) rational Padé approximation of $Y(s)$.
- Compute $\beta > 0$ a real number, corresponding to points $(\tilde{x}_i = X(\beta), \tilde{y}_i = Y(\beta))$ and $(\hat{x}_i = X(-\beta), \hat{y}_i = Y(-\beta))$ on the analytic branch of the original curve \mathbf{C} , such that $\frac{P_{mn}(s)}{Q_{mn}(s)}$ is convergent for $s \in [-\beta, \beta]$.



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4. Algorithm Contd.

- Modify $P_{mn}(s)/Q_{mn}(s)$ to $\tilde{P}_{mn}(s)/\tilde{Q}_{mn}(s)$ such that $\tilde{P}_{mn}(s)/\tilde{Q}_{mn}(s)$ is C^1 continuous approximation of $Y(s)$ on $[0, \beta]$,
- Denote the set of all the points $(\tilde{x}_i, \tilde{y}_i)$, (\hat{x}_i, \hat{y}_i) , the set T and the boundary points of \mathbf{C}_B by V . Starting from each (simple) point (x_i, y_i) in V , \mathbf{C}_B is traced out by the Taylor approximation

$$X(s) = x_i + s$$

$$Y(s) = \sum_{j=0}^{\infty} c_j^{(i)} s^j, \quad c_0^{(i)} = y_i$$



4. Algorithm Contd.

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5. Results

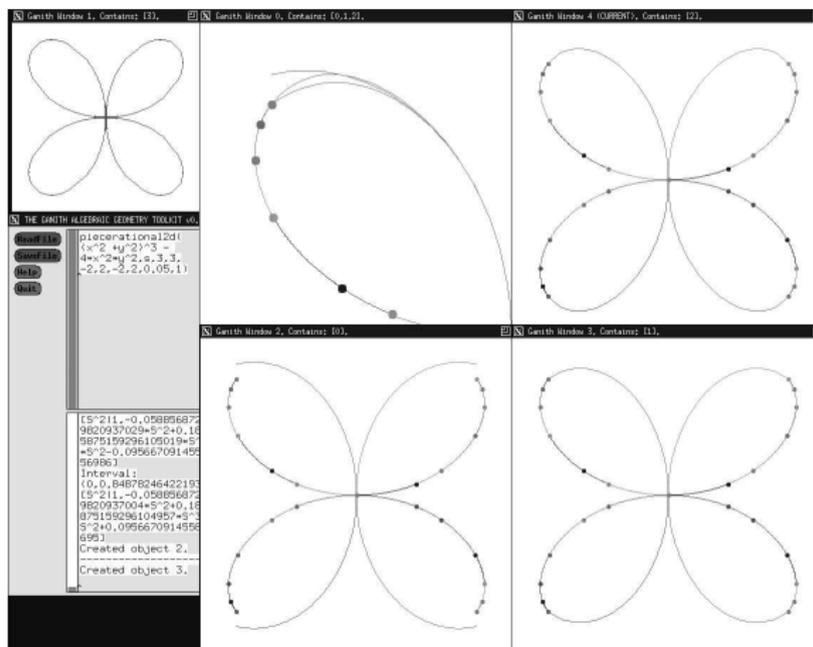


Figure: $(x^2 + y^2)^3 - 4x^2y^2 = 0$



Rational Spline Approximation of Space Curves

Given a real intersection space curve SC which is either the intersection of two implicitly defined surfaces $f_1(x, y, z) = 0$, $f_2(x, y, z) = 0$, or, the intersection of two parametric surfaces defined by

$$\begin{aligned} X_1(u_1, v_1) &= [G_{11}(u_1, v_1) \ G_{21}(u_1, v_1), \ G_{31}(u_1, v_1)]^T \\ X_2(u_2, v_2) &= [G_{12}(u_2, v_2) \ G_{22}(u_2, v_2), \ G_{32}(u_2, v_2)]^T \end{aligned}$$

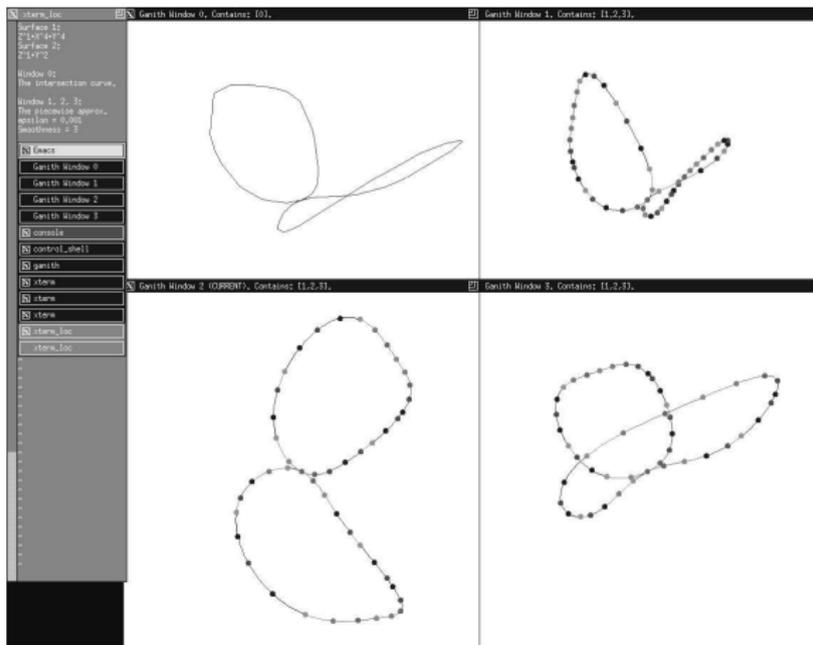
within a bounding box B and an error bound $\epsilon > 0$, a continuity index k , construct a C^k (or G^k) continuous piecewise parametric rational ϵ -approximation of all portions of SC within B .

NURBS Approximation of Surface/Surface Intersection Curves
Advances in Computational Mathematics, (1994)



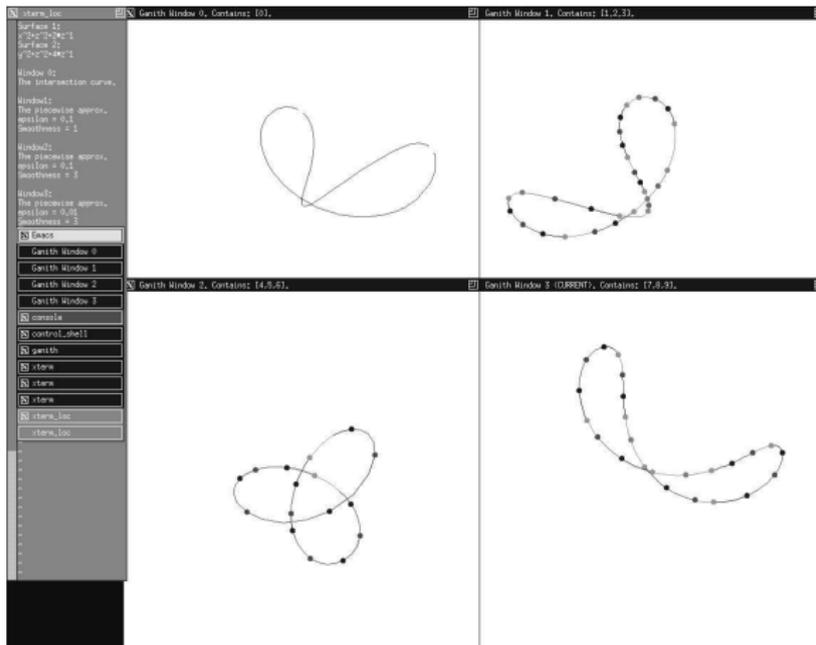
Results from Ganith - Intersection of Two implicit surfaces

Surfaces: $x^4 + y^4 + z = 0$ and $y^2 + z = 0$



Results from Ganith - Intersection of Implicit and Parametric Surfaces

Surfaces: $x^2 + z^2 + 2z = 0$ and $x = \frac{s+st^2}{1+t^2}, y = \frac{2-2t^2}{1+t^2}, z = \frac{4t-2-2t^2}{1+t^2}$



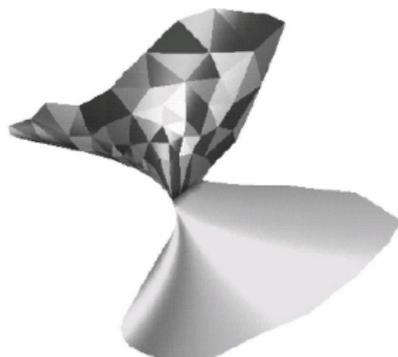
Rational Spline Approximation of Algebraic Surfaces

Given an implicit surface defined by a function $f(x, y, z) = 0$ and bounding box, create a piecewise rational spline approximation of the surface within the bounding box.

Spline Approximations of Real Algebraic Surfaces *Journal of Symbolic Computation, Special Issue on Parametric Algebraic Curves and Applications, (1997)*

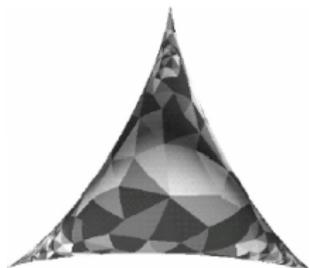


Cartan Surface: $f = x^2 - y * z^2 = 0$ has a singular point at $(0, 0, 0)$ and a singular line $(x = 0, z = 0)$.



Patch of a **Steiner Surface**:

$f = x^2 * y^2 + y^2 * z^2 + z^2 * x^2 - 4 * x * y * z = 0$ has a singular curve along x -axis, y -axis, z -axis and a triple point at the origin.



Lower Degree Spline Approximation of Rational Parametric Surfaces

For a rational parametric surface :

$$x(s, t) = \frac{X(s, t)}{W(s, t)}, y(s, t) = \frac{Y(s, t)}{W(s, t)}, z(s, t) = \frac{Z(s, t)}{W(s, t)}$$

Constructing lower degree rational spline approximations require solutions to sub-problems:

- 1 **Domain poles**
- 2 **Domain base points**
- 3 **Surface singularities**
- 4 **Complex parameter values**
- 5 **Infinite parameter values**

Triangulation and Display of Arbitrary Rational Parametric Surfaces,
Proceedings: IEEE Visualization '94 Conference

Finite Representations of Real Parametric Curves and Surfaces, *International Journal of Computational Geometry and Applications*, (1995)



Infinite parameter range

Consider the unit sphere:

implicit form: $f(x, y, z) = x^2 + y^2 + z^2 - 1 = 0$

parametric form:

$$x = 2s/(1 + s^2 + t^2) \quad (20)$$

$$y = 2t/(1 + s^2 + t^2) \quad (21)$$

$$z = 1 - s^2 - t^2/(1 + s^2 + t^2) \quad (22)$$

The point (0 , 0 , -1) can only be reached when both s and t tend to infinity.



Complex parameter values

We may need complex values to get real points

Consider the rational cubic curve:

$$\text{implicit form: } f(x, y) = x^3 + x^2 + y^2 = 0$$

$$\text{parametric form: } x(s) = -s^2 + 1, y(s) = -s(s^2 + 1)$$

The origin can only be reached with $s = \sqrt{-1}$.

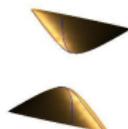


Poles

The denominator polynomial $f_4(s, t)$ may be 0, yielding a polynomial pole curve

Consider a hyperboloid of 2 sheets: implicit form:

$$f(x, y, z) = z^2 + yz + xz - y^2 - xy - x^2 - 1 = 0$$



parametric form:

$$x(s, t) = 4s/(5t^2 + 6st + 5s^2 - 1) \quad (23)$$

$$y(s, t) = 4t/(5t^2 + 6st + 5s^2 - 1) \quad (24)$$

$$z(s, t) = (5t^2 + 6st - 2t + 5s^2 - 2s + 1)/(5t^2 + 6st + 5s^2 - 1) \quad (25)$$

The problem arises from the polynomial pole curve

$5t^2 + 6st - 2t + 5s^2 - 2s + 1 = 0$ in the parameter domain.



Base points

All the polynomials may equal 0 for some values of s and t , thus causing curves (seam curves) to be missing from the parametric surface

Hyperboloid of 1 sheet with seam curve gaps caused by two base points :



Handling Base points

THEOREM : Let (a, b) be a base point of multiplicity q . Then for any $m \in R$, the image of a domain point approaching (a, b) along a line of slope m is given by $(X(m), Y(m), Z(m), W(m)) =$

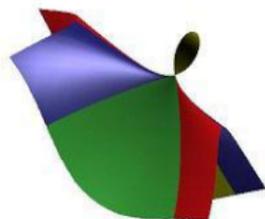
$$\sum_{i=0}^q \left(\frac{\partial^q X}{\partial s^{q-i} \partial t^i} (a, b) \right) m^i \dots \sum_{i=0}^q \left(\frac{\partial^q X}{\partial s^{q-i} \partial t^i} (a, b) \right) m^i$$

COROLLARY : If the curves $X(s, t) = 0, \dots, W(s, t) = 0$ share t tangent lines at (a, b) , then the seam curve $(X(m), Y(m), Z(m), W(m))$ has degree $q - t$. In particular, if $X(s, t) = 0$ have identical tangents at (a, b) , then for all $m \in R$ the coordinates $(X(m), \dots, W(m))$ represent a single point.

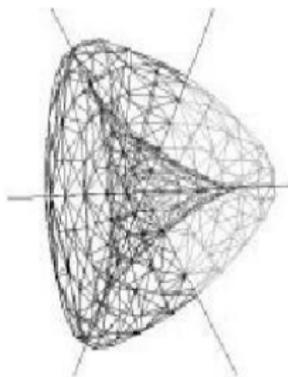


Parametric surfaces with a point, curve singularities

A Cubic Rational Surface:



The Steiner Rational Surface:



Algebraic Surface Blending, Joining, Least Squares Spline Approximations

Input: A collection of points, curves, derivative jets (scattered data) in 3D.

Output: A low degree, algebraic surface fit through the scattered set of points, curves, derivative jets, with prescribed higher order interpolation and least-squares approximation.

The mathematical model for this problem is a constrained minimization problem of the form :

$$\text{minimize } \mathbf{x}^T \mathbf{M}_A^T \mathbf{M}_A \mathbf{x} \quad \text{subject to } \mathbf{M}_I \mathbf{x} = \mathbf{0}, \quad \mathbf{x}^T \mathbf{x} = 1,$$

\mathbf{M}_I and \mathbf{M}_A are interpolation and least-square approximation matrices, and \mathbf{x} is a vector containing coefficients of an algebraic surface.



Algebraic Surface Blending, Joining, Least Squares Spline Approximations

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Definition

Two algebraic surfaces $f(x, y, z) = 0$ and $g(x, y, z) = 0$ meet with C^k *rescaling continuity* at a point p or along an irreducible algebraic curve C if and only if there exists two polynomials $a(x, y, z)$ and $b(x, y, z)$, not identically zero at p or along C , such that all derivatives of $af - bg$ up to order k vanish at p or along C .



Theorem

Let $g(x, y, z)$ and $h(x, y, z)$ be distinct, irreducible polynomials. If the surfaces $g(x, y, z) = 0$ and $h(x, y, z) = 0$ intersect transversally in a single irreducible curve C , then any algebraic surface $f(x, y, z) = 0$ that meets $g(x, y, z) = 0$ with C^k rescaling continuity along C must be of the form $f(x, y, z) = \alpha(x, y, z)g(x, y, z) + \beta(x, y, z)h^{k+1}(x, y, z)$. If $g(x, y, z) = 0$ and $h(x, y, z) = 0$ share no common components at infinity. Furthermore, the degree of $\alpha(x, y, z)g(x, y, z) \leq$ degree of $f(x, y, z)$ and the degree of $\beta(x, y, z)h^{k+1}(x, y, z) \leq$ degree of $f(x, y, z)$.

Higher-Order Interpolation and Least-Squares Approximation Using
Implicit Algebraic Surfaces *ACM Transactions on Graphics*, (1993)



Quartic Joining Surfaces

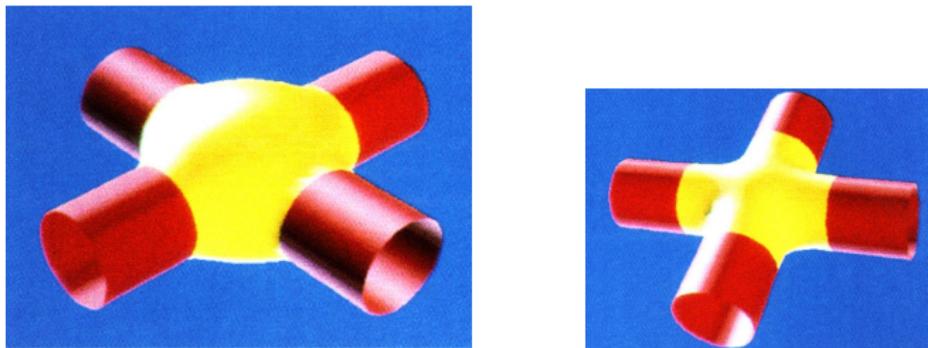


Figure: C^1 Interpolation at the Joins and Least-Squares Approximation in the Middle



Piecewise C^1 Cubic Fit

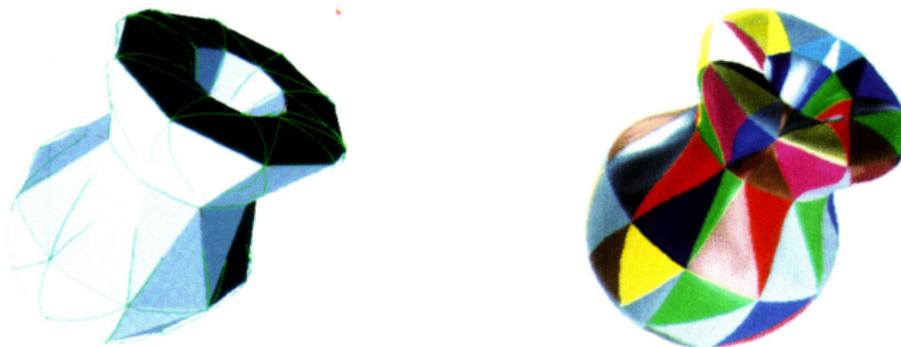
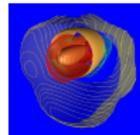
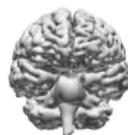
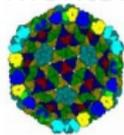
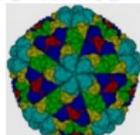


Figure: C^1 Cubic Rational Algebraic Spline



So what are Algebraic Splines, again ?

Collection (Complex) of smooth finite elements of polynomial (algebraic) curves and surfaces with prescribed order of continuity between the finite elements.



- 1 The splines are variously called Simplex, Box, Polyhedral depending on the support of the polynomial pieces.
- 2 The splines also can variously use the B-basis (B stands for Basis) or the BB-basis (BB stands for Bernstein-Bezier), or the C-basis (C for Chebyshev), etc. depending on the choice of polynomial basis
- 3 B-Splines (E.g. UBs or NUBs) or B-patches or Rational B-splines (e.g. NURBs) or T-Splines or X-splines etc. are just several examples of polynomial splines which are rational.



1 A-Splines:

- T-PACs, Cubics [Sederberg('98),Patterson-Paluzny('99)]
- C^k A-splines within triangles [Bajaj,Xu('99)]
- Regular A-splines over rectangular domains [Xu,Bajaj ('01)]
- A-splines in Data Fitting [Bajaj,Xu('03)]

2 A-Patches:

- C^1 piecewise quadric patches [Dahmen ('89)]
- Clough-Tocher split for C^1 cubic patches [Guo ('91)]
- Single valued cubic C^1 A-patches [Bajaj, Chen, Xu ('95)]
- Quintic C^2 A-patches [Bajaj, Xu ('97)]
- Rational C^1 A-patches [Xu, Bajaj ('01)]
- C^1 Prism A-patches and shell A-patches [Bajaj, Xu ('02,'03)]



C^k Triangular A-Splines

An A-spline element of degree d over the triangle $[p^1 p^2 p^3]$ is defined by

$$G_d(x, y) := F_d(\alpha) = F_d(\alpha_1, \alpha_2, \alpha_3) = 0$$

where

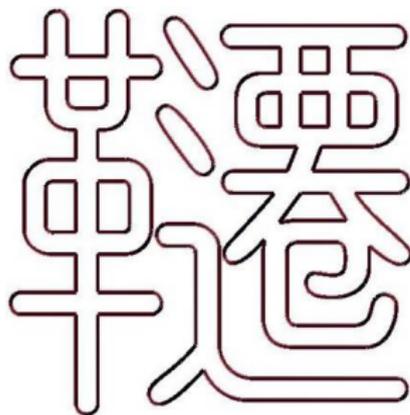
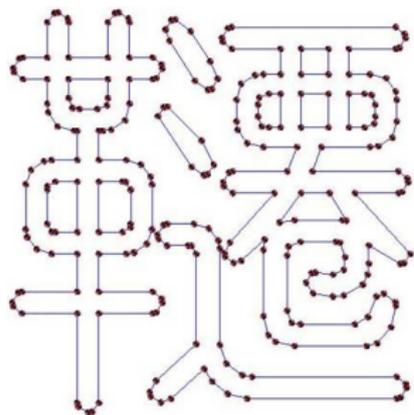
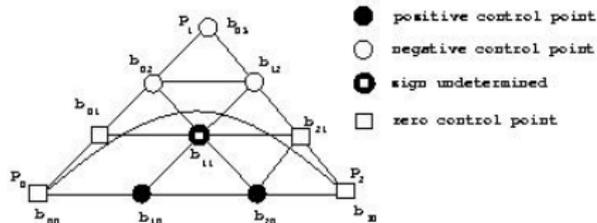
$$F_d(\alpha_1, \alpha_2, \alpha_3) = \sum_{i+j+k=d} b_{ijk} B_{ijk}^d(\alpha_1, \alpha_2, \alpha_3), B_{ijk}^d(\alpha_1, \alpha_2, \alpha_3) = \frac{d!}{i!j!k!} \alpha_1^i \alpha_2^j \alpha_3^k$$

and $(x, y)^T$ and $(\alpha_1, \alpha_2, \alpha_3)^T$ are related by

$$\begin{bmatrix} x \\ y \\ 1 \end{bmatrix} = \begin{bmatrix} p_1 & p_2 & p_3 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \alpha_3 \end{bmatrix}$$



C¹ Cubic Triangular A-Spline



Rational Parametric Form of A-Splines

and in parametric spline form

$$X(t) = \frac{\sum_{i=0}^d w_i B_i^d(t) b_i}{\sum_{i=0}^d w_i B_i^d(t)}, \quad t \in [0, 1]$$

where $b_i \in R^3$, $w_i \in R$ and $B_i^d(t) = \{d!/[i!(d-i)!]\} t^i (1-t)^{d-i}$

- A-Splines: Local Interpolation and Approximation Using G^k -Continuous Piecewise Real Algebraic Curves *Computer Aided Geometric Design*, (1999)



C^k A-Patches

A-Patches are surface finite elements.

A-Patch element of degree d over the tetrahedron p_1, p_2, p_3, p_4 is defined by

$$G_d(x, y, z) := F_d(\alpha) = F_d(\alpha_1, \alpha_2, \alpha_3, \alpha_4) = 0$$

where

$$F_d(\alpha_1, \alpha_2, \alpha_3, \alpha_4) = \sum_{i+j+k+l=d} \alpha_{ijkl} B_{ijkl}^d(\alpha_1, \alpha_2, \alpha_3, \alpha_4)$$

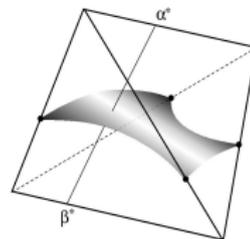
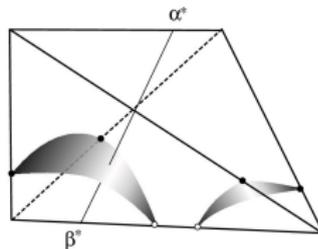
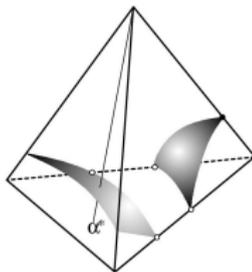
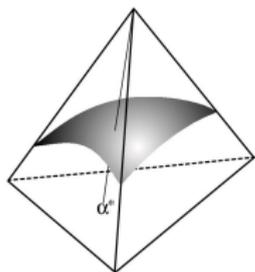
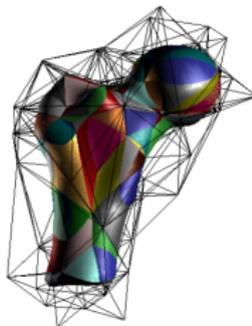
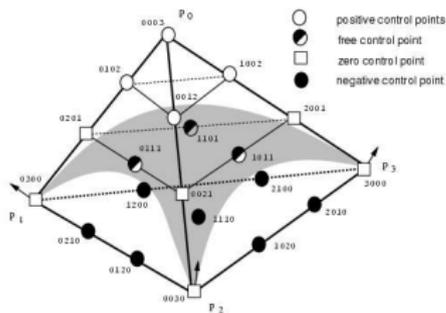
and

$(x, y, z)^T$ and $(\alpha_1, \alpha_2, \alpha_3, \alpha_4)^T$ are related by

$$\begin{bmatrix} x \\ y \\ z \\ 1 \end{bmatrix} = \begin{bmatrix} p_1 & p_2 & p_3 & p_4 \\ 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \alpha_3 \\ \alpha_4 \end{bmatrix}$$



Cubic A-patches on Tetrahedral Domains



- C^1 Modeling with Cubic A-patches *ACM Transactions on Graphics, 1995*
- C^1 Modeling with A-patches from Rational Trivariate Functions *Computer Aided Geometric Design, (2001)*



Prism C^1 A-patches

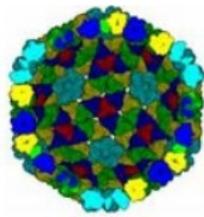
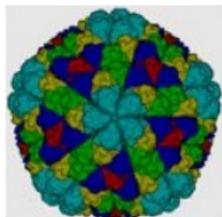
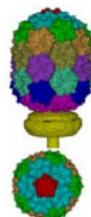
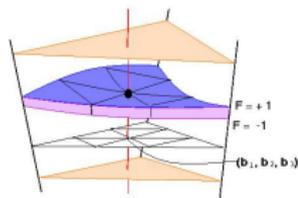
- Low degree algebraic surface finite element with dual implicit and rational parametric representations.
- The A-patch element is defined within a prism scaffold. For each triangle $v_i v_j v_k$ of a triangulation of the molecular surface, let

$$v_l(\lambda) = v_l + \lambda n_l, \quad l = i, j, k$$

Define the prism

$$D_{ijk} := \{p : p = b_1 v_i(\lambda) + b_2 v_j(\lambda) + b_3 v_k(\lambda), \lambda \in I_{ijk}\}$$

where (b_1, b_2, b_3) are the barycentric coordinates of points in $v_i v_j v_k$.



Can we convert between Algebraic Splines and Parametric Splines ?

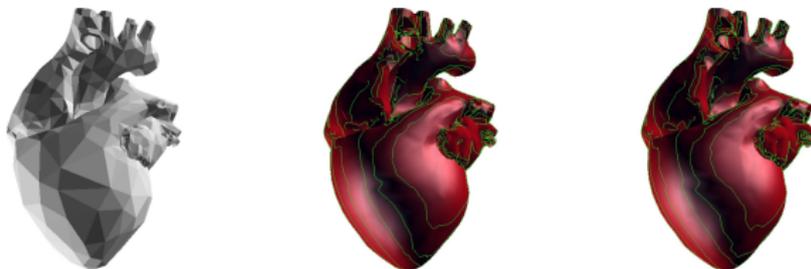


Figure: C^1 Rational Algebraic Splines

Answer: Since the algebraic plane/space curve and/or algebraic surface in general are not rational we need to construct rational parametric spline *approximations*. !

NURBs Approximation of A-splines and A-patches International Journal of Computational Geometry and Applications, (2003)

