Overview

What’s wrong with the sequential (SEQ) Y86?
- It’s slow!
- Each piece of hardware is used only a small fraction of the time.
- We would like to find a way to get more performance with only a little more hardware.

General Principles of Pipelining
- Express task as a collection of stages
- Move instructions through stages
- Process several instructions at any given moment

Creating a Pipelined Y86 Processor
- Rearrange SEQ
- Insert pipeline registers
- Deal with data and control hazards

Pipeline Correctness Axiom: A pipeline is correct only if the resulting machine satisfies the ISA (nonpipelined) semantics.

Pipelining: Laundry Example

Suppose you have four folks, each with a load of clothes to wash, dry, fold and stash away. There are four subtasks: wash, dry, fold, stash. Suppose each takes 30 minutes.

Time to do a load of laundry from start to finish: 2 hours. (That’s the latency.)
Sequential Laundry

- Sequential laundry takes 8 hours for 4 loads.
- If they learned pipelining, how long would laundry take?

Pipelined Laundry

Pipelined laundry takes 3.5 hours for 4 loads! But each load still takes 2 hours.

What’s the metric that improved? How would you measure the efficiency of the process if you were running a laundry service with loads (inputs) always ready to process?

Latency vs. Throughput

Latency is the time from start to finish for a given task.

Throughput is the number of tasks completed in a given time period.

Example: suppose that each laundry stage (wash, dry, fold, stash) takes 30 minutes. But you have a laundromat with 4 washers, 4 driers, 4 folding stations, 4 stashing stations.

- What is the latency?

Latency is 2 hours, because it still takes two hours to get any single load through the entire process.

- What is the highest possible throughput (per hour)?
Latency vs. Throughput

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*Throughput* is the number of tasks completed in a given time period.

**Example:** suppose that each laundry stage (wash, dry, fold, stash) takes 30 minutes. But you have a laundromat with 4 washers, 4 dryers, 4 folding stations, 4 stashing stations.

- **What is the latency?** Latency is 2 hours, because it still takes two hours to get any single load through the entire process.
- **What is the highest possible throughput (per hour)?** Throughput is (theoretically) 8 loads / hour since you can complete 8 loads every hour in steady state. How?

### Pipelining Lessons

- Pipelining doesn't help *latency* of a single task; it helps *throughput* of the entire workload.
- Multiple tasks operate simultaneously using different resources.
- Potential speedup = number of stages.
- Unbalanced lengths of pipe stages reduces speedup.
- Time to "fill" pipeline and time to "drain" it reduces speedup.
- May need to "stall" for dependencies.

### Computational Example

**System**

- Computation requires a total of 300 picoseconds.
- Needs an additional 20 picoseconds to save the result in the register.
- Must have a clock cycle of at least 320 ps. Why?

**3-Way Pipelined Version**

**System**

- Divide combinational logic into 3 blocks of 100 ps each.
- Can begin a new operation as soon as the previous one passes through stage A.
- Begin new operation every 120 ps. Why?
- Overall latency *increases!* It’s now 360 ps from start to finish.
Pipeline Diagrams

**Unpipelined**

<table>
<thead>
<tr>
<th>OP1</th>
<th>OP2</th>
<th>OP3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
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</table>

Cannot start new operation until the previous one completes.

**3-Way Pipelined**

<table>
<thead>
<tr>
<th>OP1</th>
<th>OP2</th>
<th>OP3</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>B</td>
<td>C</td>
</tr>
<tr>
<td>A</td>
<td>C</td>
<td></td>
</tr>
<tr>
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<td>B</td>
<td>C</td>
</tr>
</tbody>
</table>

Up to 3 operations in process simultaneously.

**Limitations: Non-uniform Delays**

<table>
<thead>
<tr>
<th>Comb Logic A</th>
<th>Reg</th>
<th>Comb Logic B</th>
<th>Reg</th>
<th>Comb Logic C</th>
<th>Reg</th>
<th>Clock</th>
</tr>
</thead>
<tbody>
<tr>
<td>50 ps</td>
<td>20</td>
<td>150 ps</td>
<td>20</td>
<td>100 ps</td>
<td>20</td>
<td></td>
</tr>
</tbody>
</table>

Delay = 510 ps
Throughput = 5.88 GIPS

**Limitations: Register Overhead**

<table>
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Delay = 420 ps
Throughput = 14.29 GIPS

As you try to deepen the pipeline, the overhead of loading registers becomes more significant.

**Percentage of clock cycle spent loading registers:**

- 1-stage pipeline: 6.25%
- 3-stage pipeline: 16.67%
- 6-stage pipeline: 28.57%

High speeds of modern processor designs are obtained through very deep pipelining. (Some models of x86 have a pipeline of 20-24 stages.)
The Performance Equation

CPU Time = \( \frac{\text{Seconds}}{\text{Program}} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}} \)

Clock Cycle Time
- Improves by a factor of almost N for N-deep pipeline.
- Not quite a factor of N due to pipeline overheads.

Cycles Per Instructions (CPI)
- In an ideal world, CPI would stay the same.
- An individual instruction takes N cycles.
- But we have N instructions in flight at a time.
- So, average CPI_{pipe} = (CPI_{no-pipe} \times N) / N
Thus, performance can improve by up to a factor of N.

Sequential System: Each operation may depend on the previous one. (It doesn’t matter for a sequential system. Why not?)

Data Hazards

Pipelined System:
- Result does not feed back around in time for the next operation.
- Pipelining has changed the behavior of the system. Alarm!!

Data Hazards in Processors

Result from one instruction is used as an operand for another; called read-after-write (RAW) dependency.
This is very common in actual programs.

Must make sure that our pipeline handles these properly and gets the right result.
Should minimize performance impact as much as possible.
A control hazard occurs if something interferes with the flow of control through the program. I.e., the PC is not determined quickly enough to allow fetching the next instruction.

```
xorq %rbx, %rbx
je Done
 irmovq $100, %rax
ret
Done: irmovq $200, %rax
ret
```

When the `je` instruction moves from the fetch to decode stage, what is the next instruction to fetch? When will you know?

**Pipeline Correctness Axiom:** A pipeline is correct only if the resulting machine satisfies the ISA (nonpipelined) semantics.

That is, the pipeline implementation must deal correctly with potential data and control hazards. *Any program that runs correctly on the sequential machine must run on the pipelined version with the exact same results.*

**SEQ Hardware**
- Stages occur in sequence.
- One operation in process at at time.
- One stage for each logical pipeline operation.
  - **Fetch:** get next instruction from memory.
  - **Decode:** figure out what to do, and get values from regfile.
  - **Execute:** compute.
  - **Memory:** access data memory if needed.
  - **Write back:** write results to regfile, if needed.

**SEQ+ Hardware**
- Still sequential implementation, but reorder PC stage to put at the beginning

**PC Stage**
- Task is to select PC for current instruction.
- Based on results computed by previous instruction.

**Processor State**
- PC is no longer stored in a register.
- But, can determine PC based on other stored information.