CS429: Computer Organization and Architecture

Bits and Bytes

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There are 10 kinds of people in the world: those who understand binary, and those who don’t!

- Why bits?
- Representing information as bits
  - Binary and hexadecimal
  - Byte representations: numbers, characters, strings, instructions, etc.
- Bit level manipulations
  - Boolean algebra
  - C constructs
Great Reality 7: Whatever you plan to store on a computer ultimately has to be represented as a finite collection of bits.

That’s true whether it’s integers, reals, characters, strings, data structures, instructions, programs, pictures, videos, etc.

That really means that only discrete quantities can be represented exactly. Non-discrete (continuous) quantities have to be approximated.
Base 10 Number Representation.

- Fingers *are* called as “digits.”
- Natural representation for financial transactions. Floating point number cannot exactly represent $1.20.
- Even carries through in scientific notation: $1.5213 \times 10^4$

*If we lived in Homer Simpson’s world, we’d all use octal!*
Implementing Electronically

- 10 different values are hard to store. ENIAC (First electronic computer) used 10 vacuum tubes / digits
- They’re hard to transmit. Need high precision to encode 10 signal levels on single wire.
- Messy to implement digital logic functions: addition, multiplication, etc.
Alternative: Binary Representation

**Base 2 Number Representation**
- Represent $15213_{10}$ as $11101101101101_2$
- Represent $1.20_{10}$ as $1.0011001100110011[0011]...2$
- Represent $1.5213 \times 10^4$ as $1.1101101101101_2 \times 2^{13}$

**Electronic Implementation**
- Easy to store bits with bistable elements.
- Reliably transmitted on noisy and inaccurate wires.
To store data of type $X$, someone had to invent a mapping from items of type $X$ to bit strings. That’s the representation mapping.

In a sense the representation is arbitrary. The representation is just a mapping from the domain onto a finite set of bit strings.

But some representations are better than others. Why would that be? Hint: what operations do you want to support?
Suppose you want to represent the finite set of natural numbers $[0\ldots 7]$ as 3-bit strings.

<table>
<thead>
<tr>
<th>Dec.</th>
<th>Rep1</th>
<th>Rep2</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>101</td>
<td>000</td>
</tr>
<tr>
<td>1</td>
<td>011</td>
<td>001</td>
</tr>
<tr>
<td>2</td>
<td>111</td>
<td>010</td>
</tr>
<tr>
<td>3</td>
<td>000</td>
<td>011</td>
</tr>
<tr>
<td>4</td>
<td>110</td>
<td>100</td>
</tr>
<tr>
<td>5</td>
<td>010</td>
<td>101</td>
</tr>
<tr>
<td>6</td>
<td>001</td>
<td>110</td>
</tr>
<tr>
<td>7</td>
<td>100</td>
<td>111</td>
</tr>
</tbody>
</table>

Can you see why one of these representations is “better” than the other? Hint: How would you do addition using Rep1?
A “good” mapping will map X data onto bit strings (B) in a way that makes it easy to compute common operations on that data. I.e., the following diagram must commute, for a reasonable choice of conc-op.
Representing Data: Integer Addition

```
int x;
int y;
...
t = x + y;
```

To carry out any operation at the C level means converting the data into bit strings, and implementing an operation on the bit strings that has the “intended effect” under the mapping.
Important Fact 1: If you are going to represent any type in $k$ bits, you can only represent $2^k$ different values. There are exactly as many ints as floats on x86.

Important Fact 2: The same bit string can represent an integer (signed or unsigned), float, character string, list of instructions, address, etc. depending on the context.
Since it’s tedious always to think in terms of bits, we group them together into larger units. *Sizes of these units depends on the architecture/language.*
Byte = 8 bits
Which can be represented in various forms:

- Binary: 00000000₂ to 11111111₂
- Decimal: 0₁₀ to 255₁₀
- Hexadecimal: 00₁₆ to FF₁₆
  - Base 16 number representation
  - Use characters '0' to '9' and 'A' to 'F'
  - Write FA1D37B₁₆ in C as 0xFA1D37B or 0xfa1d37b

BTW: one hexadecimal digit represents 4 bits (one nybble).
### Memory as Array

<table>
<thead>
<tr>
<th>Computer Address</th>
<th>Content</th>
<th>Programmers Name</th>
<th>Type</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>900000000</td>
<td>00</td>
<td>sum</td>
<td>int</td>
<td>000000FF(255₁₀)</td>
</tr>
<tr>
<td>900000001</td>
<td>00</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000002</td>
<td>00</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000003</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000004</td>
<td>FF</td>
<td>age</td>
<td>short</td>
<td>FFFF(-1₁₀)</td>
</tr>
<tr>
<td>900000005</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000006</td>
<td>1F</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000007</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000008</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>900000009</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000000A</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000000B</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000000C</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000000D</td>
<td>FF</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000000E</td>
<td>90</td>
<td>average</td>
<td>double</td>
<td>1FFFFFFFFFFFFFFF(4.45015E-308₁₀)</td>
</tr>
<tr>
<td>90000000F</td>
<td>00</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>90000010</td>
<td>00</td>
<td>ptrSum</td>
<td>int*</td>
<td>900000000</td>
</tr>
<tr>
<td>90000011</td>
<td>00</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Note: All numbers in hexadecimal

**Note:** this picture is appropriate for a 32-bit, big endian machine. How did I know that?
Conceptually, memory is a very large array of bytes.
Actually, it’s implemented with hierarchy of different memory types.
- SRAM, DRAM, disk.
- The OS only allocates storage for regions actually used by program.

In Unix and Windows, address space private to particular “process.”
- Encapsulates the program being executed.
- Program can clobber its own data, but not that of others.
Byte-Oriented Memory Organization

Compiler and Run-Time System Control Allocation

- Where different program objects should be stored.
- Multiple storage mechanisms: static, stack, and heap.
- In any case, all allocation within single virtual address space.
Machines generally have a specific “word size.”

- It’s the nominal size of addresses on the machine.
- Most current machines run 64-bit software (8 bytes).
  - 32-bit software limits addresses to 4GB.
  - Becoming too small for memory-intensive applications.
- All x86 current hardware systems are 64 bits (8 bytes).
  Potentially address around $1.8 \times 10^{19}$ bytes.
- Machines support multiple data formats.
  - Fractions or multiples of word size.
  - Always integral number of bytes.
- X86-hardware systems operate in 16, 32, and 64-bit modes.
  - Initially starts in 286 mode, which is 16-bit.
  - Under programmer control, 32- and 64-bit modes are enabled.
Word-Oriented Memory Organization

Addresses Specify Byte Locations

- Which is the address of the *first* byte in word.
- Addresses of successive words differ by 4 (32-bit) or 8 (64-bit).
- Addresses of multi-byte data items are typically *aligned* according to the size of the data.

<table>
<thead>
<tr>
<th>32-bit words</th>
<th>64-bit words</th>
<th>bytes</th>
<th>addr.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Addr: 0000</td>
<td>Addr: 0000</td>
<td></td>
<td>0000</td>
</tr>
<tr>
<td>Addr: 0004</td>
<td></td>
<td></td>
<td>0001</td>
</tr>
<tr>
<td>Addr: 0008</td>
<td></td>
<td></td>
<td>0002</td>
</tr>
<tr>
<td>Addr: 0012</td>
<td></td>
<td></td>
<td>0003</td>
</tr>
<tr>
<td></td>
<td>Addr: 0008</td>
<td></td>
<td>0004</td>
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<tr>
<td></td>
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<td></td>
<td>0005</td>
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<td>0013</td>
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<td></td>
<td>0014</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>0015</td>
</tr>
</tbody>
</table>
The *integer data types* (int, long int, short, char) can all be either signed or unsigned.
Byte Ordering

How should bytes within a multi-byte data item be ordered in memory?

Given 64-bit hex value 0x0001020304050607, it is common to store this in memory in one of two formats: big endian or little endian.

Note that “endian-ness” only applies to multi-byte primitive data items, not to strings, arrays, or structs.
**Byte Ordering Examples**

**Big Endian:** Most significant byte has lowest (first) address.

**Little Endian:** Least significant byte has lowest address.

**Example:**
- Int variable \( x \) has 4-byte representation \( 0x01234567 \).
- Address given by \&x is \( 0x100 \)

### Big Endian:

<table>
<thead>
<tr>
<th>Address:</th>
<th>0x100</th>
<th>0x101</th>
<th>0x102</th>
<th>0x103</th>
</tr>
</thead>
<tbody>
<tr>
<td>Value:</td>
<td>01</td>
<td>23</td>
<td>45</td>
<td>67</td>
</tr>
</tbody>
</table>

### Little Endian:

<table>
<thead>
<tr>
<th>Address:</th>
<th>0x100</th>
<th>0x101</th>
<th>0x102</th>
<th>0x103</th>
</tr>
</thead>
<tbody>
<tr>
<td>Value:</td>
<td>67</td>
<td>45</td>
<td>23</td>
<td>01</td>
</tr>
</tbody>
</table>
## Conventions

- Sun, PowerPC Macintosh computers are “big endian” machines: most significant byte has lowest (first) address.
- Alpha, Intel MacIntosh, x86s are “little endian” machines: least significant byte has lowest address.
- ARM processor offers support for big endian, but mainly they are used in their default, little endian configuration.
- There are many (hundreds) of microcontrollers so check before you start programming!
Disassembly

- Yields textual representation of binary machine code.
- Generated by program that reads the machine code.

Example Fragment (IA32):

<table>
<thead>
<tr>
<th>Address</th>
<th>Instruction</th>
<th>Code</th>
<th>Assembly Rendition</th>
</tr>
</thead>
<tbody>
<tr>
<td>8048365:</td>
<td>5b</td>
<td>pop %ebx</td>
<td></td>
</tr>
<tr>
<td>8048366:</td>
<td>81 c3 ab 12 00 00</td>
<td>add $0x12ab,%ebx</td>
<td></td>
</tr>
<tr>
<td>804836c:</td>
<td>83 bb 28 00 00 00</td>
<td>cmpl $0x0,0x28(%ebx)</td>
<td></td>
</tr>
</tbody>
</table>

Deciphering Numbers: Consider the value $0x12ab$ in the second line of code:

- Pad to 4 bytes: $0x000012ab$
- Split into bytes: 00 00 12 ab
- Make little endian: ab 12 00 00
**Examining Data Representations**

**Code to Print Byte Representations of Data**
Casting a pointer to unsigned `char *` creates a byte array.

```c
typedef unsigned char *pointer;

void show_bytes(pointer start, int len)
{
    int i;
    for (i = 0; i < len; i++)
        printf(“%p\t0x%.2x\n”, start+i, start[i]);
    printf("\n");
}
```

**Printf directives:**
- `%p`: print pointer
- `%x`: print hexadecimal
```c
int a = 15213;
printf("int a = 15213;\n");
show_bytes((pointer) &a, sizeof(int));
```

**Result (Linux):**

```c
int a = 15213;
0x7fff90c56c7c 0x6d
0x7fff90c56c7d 0x3b
0x7fff90c56c7e 0x00
0x7fff90c56c7f 0x00
```
int A = 15213;
int B = −15213;
long int C = 15213;

\[15213_{10} = 0011101101101101_2 = 3B6D_{16}\]

<table>
<thead>
<tr>
<th></th>
<th>Linux (little endian)</th>
<th>Alpha (little endian)</th>
<th>Sun (big endian)</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>6D 3B 00 00</td>
<td>6D 3B 00 00</td>
<td>00 00 3B 6D</td>
</tr>
<tr>
<td>B</td>
<td>93 C4 FF FF</td>
<td>93 C4 FF FF</td>
<td>FF FF C4 93</td>
</tr>
<tr>
<td>C</td>
<td>6D 3B 00 00 00 00 00 00 00</td>
<td>6D 3B 00 00 00 00 00 00 00</td>
<td>00 00 00 00 00 00 3B 6D</td>
</tr>
</tbody>
</table>

We’ll cover the representation of negatives later.
Representing Pointers

```c
int B = -15213;
int *P = &B;
```

**Linux Address:**
Hex: BFFFF8D4AFBB4CD0
In memory: D0 4C BB AF D4 F8 FF BF

**Sun Address:**
Hex: EFFFFFB2CAA2C15C0
In Memory: EF FF FB 2C AA 2C 15 C0

*Pointer values generally are not predictable. Different compilers and machines assign different locations.*
All modern machines implement the IEEE Floating Point standard. This means that it is consistent across all machines.

```plaintext
float F = 15213.0;
```

Binary: 01000110011011011011010000000000
Hex: 466DB400
In Memory (Linux/Alpha): 00 B4 6D 46
In Memory (Sun): 46 6D B4 00

Note that it’s not the same as the `int` representation, but you can see that the `int` is in there, if you know where to look.
Representing Strings

Strings are represented by a sequence of characters.
Each character is encoded in ASCII format.
  - Standard 7-bit encoding of character set.
  - Other encodings exist, but are less common.
Strings should be null-terminated. That is, the final character has ASCII code 0. I.e., a string of $k$ chars requires $k + 1$ bytes.

Compatibility

- *Byte ordering (endian-ness) is not an issue* since the data are single byte quantities.
- Text files are generally platform independent, except for different conventions of line break character(s).
Encode Program as Sequence of Instructions

- Each simple operation
  - Arithmetic operation
  - Read or write memory
  - Conditional branch

- Instructions are encoded as sequences of bytes.
  - Alpha, Sun, PowerPC Mac use 4 byte instructions (Reduced Instruction Set Computer” (RISC)).
  - PC’s and Intel Mac’s use variable length instructions (Complex Instruction Set Computer (CISC)).

- Different instruction types and encodings for different machines.

- Most code is not binary compatible.

**Remember:** Programs are byte sequences too!
For this example, Alpha and Sun use two 4-byte instructions. They use differing numbers of instructions in other cases.

PC uses 7 instructions with lengths 1, 2, and 3 bytes. Windows and Linux are not fully compatible.

Different machines typically use different instructions and encodings.

**Instruction sequence for sum program:**

**Alpha:** 00 00 30 42 01 80 FA 68  
**Sun:** 81 C3 E0 08 90 02 00 09  
**PC:** 55 89 E5 8B 45 OC 03 45 08 89 EC 5D C3
Assembly vs. Machine Code

Machine code bytes

<table>
<thead>
<tr>
<th>Machine code bytes</th>
<th>Assembly language statements</th>
</tr>
</thead>
<tbody>
<tr>
<td>B8 22 11 00 FF</td>
<td>movl $0xFFFF001122, %eax</td>
</tr>
<tr>
<td>01 CA</td>
<td>addl %ecx, %edx</td>
</tr>
<tr>
<td>31 F6</td>
<td>xorl %esi, %esi</td>
</tr>
<tr>
<td>53</td>
<td>pushl %ebx</td>
</tr>
<tr>
<td>8B 5C 24 04</td>
<td>movl 4(%esp), %ebx</td>
</tr>
<tr>
<td>8D 34 48</td>
<td>leal (%eax, %ecx, 2), %esi</td>
</tr>
<tr>
<td>39 C3</td>
<td>cmpl %eax, %ebx</td>
</tr>
<tr>
<td>72 EB</td>
<td>jnae foo</td>
</tr>
<tr>
<td>C3</td>
<td>retl</td>
</tr>
</tbody>
</table>

Instruction stream

B8 22 11 00 FF 01 CA 31 F6 53 8B 5C 24
04 8D 34 48 39 C3 72 EB C3
Recall **Great Reality 7**: Whatever you plan to store on a computer ultimately has to be represented as a finite collection of bits.

That’s true whether it’s integers, reals, characters, strings, data structures, instructions, programs, pictures, videos, audio files, etc. Anything!
Developed by George Boole in the 19th century, Boolean algebra is the algebraic representation of logic. We encode “True” as 1 and “False” as 0.
Boolean Algebra

**And:** $A \ & \ B = 1$ when both $A = 1$ and $B = 1$.

<table>
<thead>
<tr>
<th>$A$</th>
<th>$B$</th>
<th>$&amp;$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>0</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
</tbody>
</table>

**Or:** $A \ | \ B = 1$ when either $A = 1$ or $B = 1$.

| $A$ | $B$ | $|$ |
|-----|-----|----|
| 0   | 0   | 0  |
| 0   | 1   | 1  |
| 1   | 0   | 1  |
| 1   | 1   | 1  |

**Not:** $\sim A = 1$ when $A = 0$.

<table>
<thead>
<tr>
<th>$A$</th>
<th>$\sim$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
</tr>
</tbody>
</table>

**Xor:** $A \ ^\wedge \ B = 1$ when either $A = 1$ or $B = 1$, but not both.

<table>
<thead>
<tr>
<th>$A$</th>
<th>$B$</th>
<th>$\ ^\wedge$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>0</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
</tbody>
</table>
In a 1937 MIT Master’s Thesis, Claude Shannon showed that Boolean algebra would be a great way to model digital networks.

At that time, the networks were relay switches. But today, all combinational circuits can be described in terms of Boolean “gates.”
\[ \langle \{0, 1\}, \lor, \land, \sim, 0, 1 \rangle \text{ forms a Boolean algebra.} \]

- Or is the sum operation.
- And is the product operation.
- \( \sim \) is the “complement” operation (not additive inverse).
- 0 is the identity for sum.
- 1 is the identity for product.
Some boolean algebra properties are similar to integer arithmetic, some are not.

**Commutativity:**

\[
A | B = B | A \\
A \& B = B \& A \\
A + B = B + A \\
A \times B = B \times A
\]

**Associativity:**

\[
(A | B) | C = A | (B | C) \\
(A \& B) \& C = A \& (B \& C) \\
(A + B) + C = A + (B + C) \\
(A \times B) \times C = A \times (B \times C)
\]

**Product Distributes over Sum:**

\[
A \& (B | C) = (A \& B) | (A \& C) \\
A \times (B + C) = (A \times B) + (A \times C)
\]

**Sum and Product Identities:**

\[
A | 0 = A \\
A \& 1 = A \\
A + 0 = A \\
A \times 1 = A
\]
Boolean Algebra Properties

**Zero is product annihilator:**

\[ A \& 0 = 0 \quad \quad A \ast 0 = 0 \]

**Cancellation of negation:**

\[ \sim (\sim A)) = A \quad \quad \neg(\neg A) = A \]

The following boolean algebra rules don’t have analogs in integer arithmetic.

**Boolean:** Sum distributes over product

\[ A|(B \& C) = (A|B) \& (A|C) \quad A + (B \ast C) \neq (A + B) \ast (A + C) \]

**Boolean:** Idempotency

\[ A|A = A \quad \quad A + A \neq A \]

\[ A \& A = A \quad \quad A \ast A \neq A \]
**Boolean: Absorption**

\[ A | (A \& B) = A \quad A + (A \ast B) \neq A \]
\[ A \& (A | B) = A \quad A \ast (A + B) \neq A \]

**Boolean: Laws of Complements**

\[ A | \sim A = 1 \quad A + -A \neq 1 \]

**Ring: Every element has additive inverse**

\[ A | \sim A \neq 0 \quad A + -A = 0 \]
Properties of & and ^

Commutative sum: \( A \hat{\lor} B = B \hat{\lor} A \)

Commutative product: \( A \land B = B \land A \)

Associative sum: \( (A \hat{\lor} B) \hat{\lor} C = A \hat{\lor} (B \hat{\lor} C) \)

Associative product: \( (A \land B) \land C = A \land (B \land C) \)

Prod. over sum: \( A \land (B \hat{\lor} C) = (A \land B) \hat{\lor} (A \land C) \)

0 is sum identity: \( A \hat{\lor} 0 = A \)

1 is prod. identity: \( A \land 1 = A \)

0 is product annihilator: \( A \land 0 = 0 \)

Additive inverse: \( A \hat{\lor} A = 0 \)
DeMorgan’s Laws
Express & in terms of |, and vice-versa:

\[ A \& B = \sim (\sim A | \sim B) \]
\[ A | B = \sim (\sim A \& \sim B) \]

Exclusive-Or using Inclusive Or:

\[ A ^ B = (\sim A \& B) | (A \& \sim B) \]
\[ A ^ B = (A | B) \& \sim (A \& B) \]
We can also operate on bit vectors (bitwise). All of the properties of Boolean algebra apply:

\[
\begin{array}{c}
\text{01101001} & \text{01101001} & \text{01101001} \\
\& \text{01010101} & | \text{01010101} & \sim \text{01010101} & \sim \text{01010101} \\
\hline
\text{01000001} & \text{01111101} & \text{00111100} & \text{10101010}
\end{array}
\]
The operations &, |, ~, ^ are all available in C.

- Apply to any integral data type: long, int, short, char.
- View the arguments as bit vectors.
- Operations are applied bit-wise to the argument(s).

Examples: (char data type)

\[
\begin{align*}
\sim 0x41 & \rightarrow 0xBE \\
\sim 01000001_2 & \rightarrow 10111110_2 \\
\sim 0x00 & \rightarrow 0xFF \\
\sim 00000000_2 & \rightarrow 11111111_2 \\
0x69 \& 0x55 & \rightarrow 0x41 \\
01101001_2 \& 01010101_2 & \rightarrow 01000001_2 \\
0x69|0x55 & \rightarrow 0x7D \\
01101001_2|01010101_2 & \rightarrow 01111101_2 \\
01101001_2^01010101_2 & \rightarrow 00111100_2 \\
\end{align*}
\]
Logical Operators in C

There is another set of operators in C, called the logical operators, (&&, ||, !). These treat inputs as booleans, not as strings of booleans.

- View 0 as “False.”
- View anything nonzero as “True.”
- Always return 0 or 1.
- Always do short-circuit evaluation (early termination)
- There isn’t a “logical” xor, but != works if you know the inputs are boolean.

Examples:

```
!0x41 → 0x00
!0x00 → 0x01
!!0x41 → 0x01
!!0x69 && 0x55 → 0x01
!!0x69 || 0x55 → 0x01
```
Representing Sets with Masks

Representation
A bit vector $a$ may represent a subset $S$ of some “reference set” (actually list) $L$: $a_j = 1$ iff $L[j] \in S$

Bit vector A:
\[
\begin{array}{c}
01101001 \\
76543210
\end{array}
\]
represents $\{0, 3, 5, 6\}$

Bit vector B:
\[
\begin{array}{c}
01010101 \\
76543210
\end{array}
\]
represents $\{0, 2, 4, 6\}$

What bit operations on these set representations correspond to: intersection, union, complement?
Bit vector A: 01101001 = \{0, 3, 5, 6\}
Bit vector B: 01010101 = \{0, 2, 4, 6\}

**Operations:**

Given the two sets above, perform these bitwise ops to obtain:

<table>
<thead>
<tr>
<th>Set operation</th>
<th>Boolean op</th>
<th>Result</th>
<th>Set</th>
</tr>
</thead>
<tbody>
<tr>
<td>Intersection</td>
<td>A &amp; B</td>
<td>01000001</td>
<td>{0, 6}</td>
</tr>
<tr>
<td>Union</td>
<td>A</td>
<td>B</td>
<td>01111101</td>
</tr>
<tr>
<td>Symmetric difference</td>
<td>A ^ B</td>
<td>00111100</td>
<td>{2, 3, 4, 5}</td>
</tr>
<tr>
<td>Complement</td>
<td>^A</td>
<td>10010110</td>
<td>{1, 2, 4, 7}</td>
</tr>
</tbody>
</table>
Given such a set of switches on floors A and B, how could you store the following information conveniently?

1. Which lights are on on each floor?
2. Which lights are on both floors?
3. Which lights are on on either floor?
4. Which lights are on on floor A but not floor B?
**Left Shift:** \( x \ll y \)
Shift bit vector \( x \) left by \( y \) positions
- Throw away extra bits on the left.
- Fill with 0’s on the right.

**Right Shift:** \( x \gg y \)
Shift bit vector \( x \) right by \( y \) positions.
- Throw away extra bits on the right.
- **Logical shift:** Fill with 0’s on the left.
- **Arithmetic shift:** Replicate with most significant bit on the left.

Unlike Java, C uses the same operator for logical and arithmetic right shift; the compiler “guesses” which one you meant according to the type of the operand.
Shift Examples

<table>
<thead>
<tr>
<th>Argument x</th>
<th>01100010</th>
</tr>
</thead>
<tbody>
<tr>
<td>x &lt;&lt; 3</td>
<td>00010000</td>
</tr>
<tr>
<td>x &gt;&gt; 2 (logical)</td>
<td>00011000</td>
</tr>
<tr>
<td>x &gt;&gt; 2 (arithmetic)</td>
<td>00011000</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Argument x</th>
<th>10100010</th>
</tr>
</thead>
<tbody>
<tr>
<td>x &lt;&lt; 3</td>
<td>00010000</td>
</tr>
<tr>
<td>x &gt;&gt; 2 (logical)</td>
<td>00101000</td>
</tr>
<tr>
<td>x &gt;&gt; 2 (arithmetic)</td>
<td>11101000</td>
</tr>
</tbody>
</table>

For right shift, the compiler will choose arithmetic shift if the argument is signed, and logical shift if unsigned.
Bitwise XOR is a form of addition, with the extra property that each value is its own additive inverse: \( A \oplus A = 0 \).

```c
void funny_swap(int *x, int *y)
{
    *x = *x ^ *y; /* #1 */
    *y = *x ^ *y; /* #2 */
    *x = *x ^ *y; /* #3 */
}
```

<table>
<thead>
<tr>
<th></th>
<th>*x</th>
<th>*y</th>
</tr>
</thead>
<tbody>
<tr>
<td>Begin</td>
<td>A</td>
<td>B</td>
</tr>
<tr>
<td>1</td>
<td>( A ) ( \oplus B )</td>
<td>( B )</td>
</tr>
<tr>
<td>2</td>
<td>( A ) ( \oplus B )</td>
<td>((A \oplus B) \oplus B = A)</td>
</tr>
<tr>
<td>3</td>
<td>((A \oplus B) \oplus A = B)</td>
<td>( A )</td>
</tr>
<tr>
<td>End</td>
<td>( B )</td>
<td>( A )</td>
</tr>
</tbody>
</table>

Is there ever a case where this code fails?
Main Points

It’s all about bits and bytes.

- Numbers
- Programs
- Text

Different machines follow different conventions.

- Word size
- Byte ordering
- Representations

Boolean algebra is the mathematical basis.

- Basic form encodes “False” as 0 and “True” as 1.
- General form is like bit-level operations in C; good for representing and manipulating sets.