

Symbolically Synthesizing Small Circuits

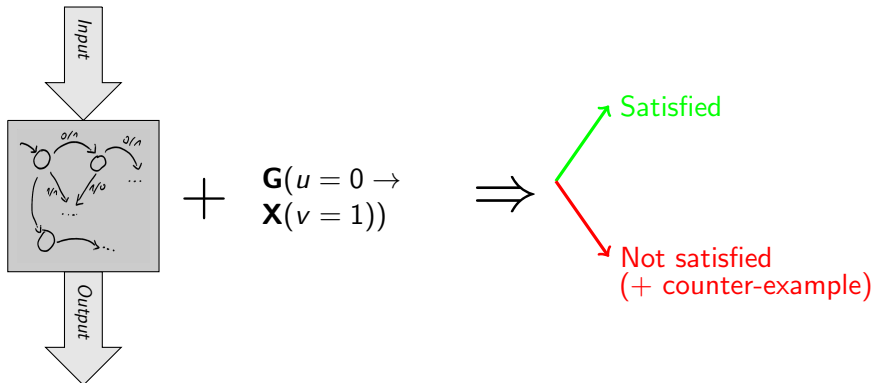
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Verification:



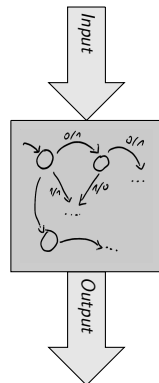
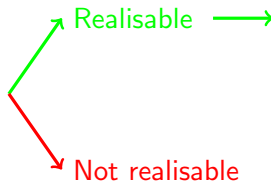
Synthesis:

$\mathbf{G}(u = 0 \rightarrow$
 $\mathbf{X}(v = 1))$

+

Input = $\{u, \dots\}$

Output = $\{v, \dots\}$



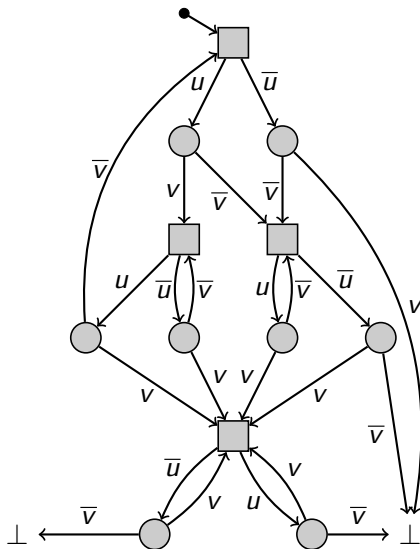
Synthesis in a nutshell

$$\mathbf{G}(u = 0 \\ \rightarrow \mathbf{X}(v = 1))$$

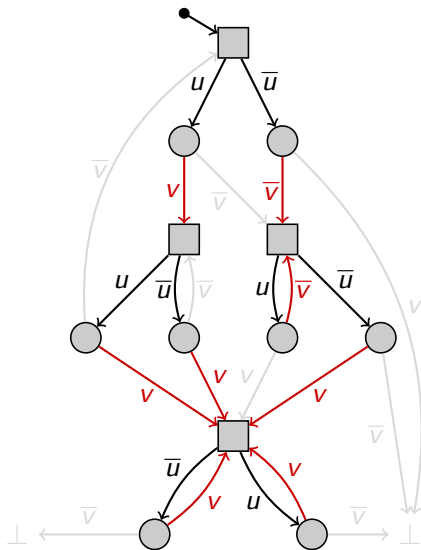
+

Input = $\{u, \dots\}$
Output = $\{v, \dots\}$

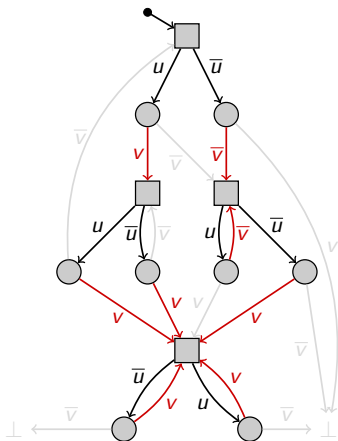
Synthesis in a nutshell



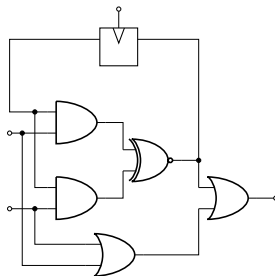
Synthesis in a nutshell



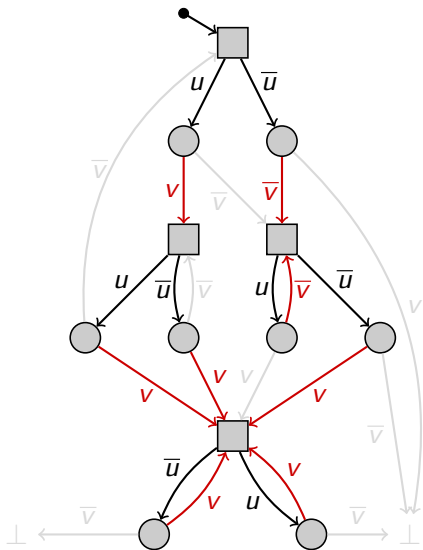
Synthesis in a nutshell



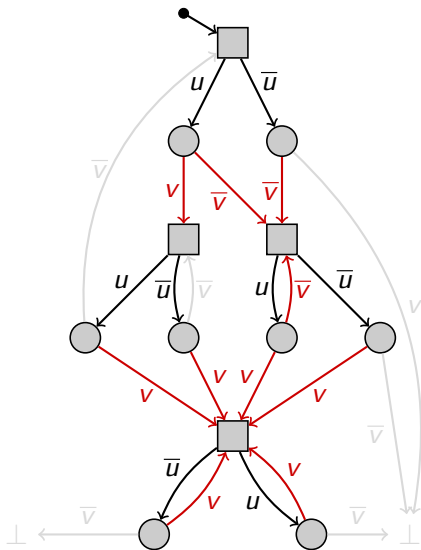
```
int main() {
    bool s = 1;
    while (true) {
        bool b = in-u();
        out-v(b | s);
        s = s  $\oplus$  u;
        wait();
    }
}
```



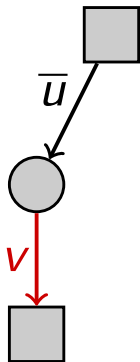
On general strategies



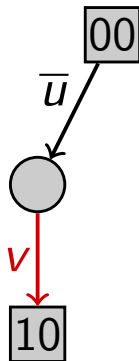
On general strategies



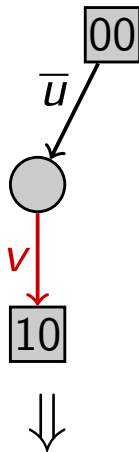
Encoding general strategies



Encoding general strategies

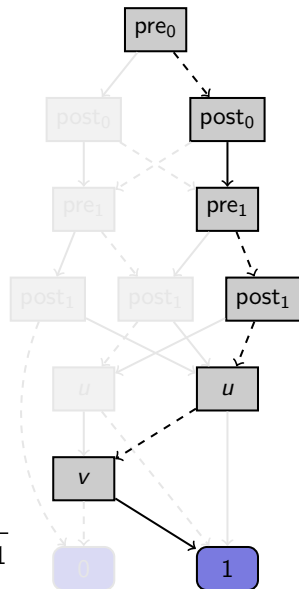
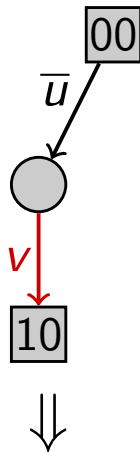


Encoding general strategies



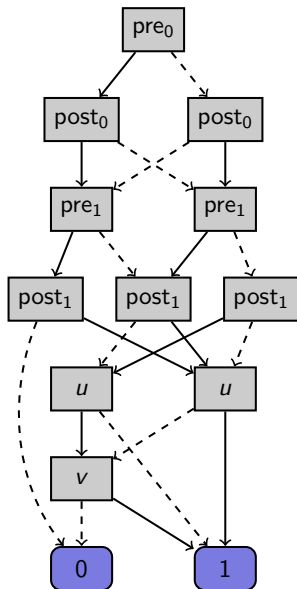
$$\overline{pre_0} \wedge \overline{pre_1} \wedge \bar{u} \wedge v \wedge pre_0 \wedge \overline{pre_1}$$

Encoding general strategies



$$\overline{pre_0} \wedge \overline{pre_1} \wedge \bar{u} \wedge v \wedge pre_0 \wedge \overline{pre_1}$$

Building circuits from general strategies



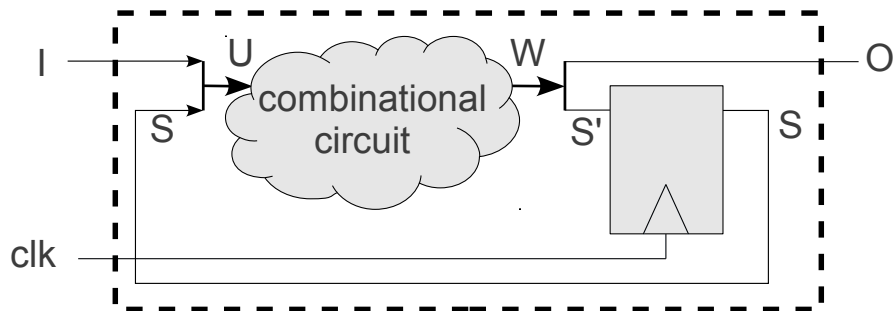
General strategies ...

- ... represent allowed behaviour of the system to synthesize
- ... are typically given symbolically
- ... are often **huge**. To have 64 state bits in the BDD is common.

What is a suitable implementation?

Any implementation that is a specialisation of the general strategy is fine.

Building circuits from general strategies

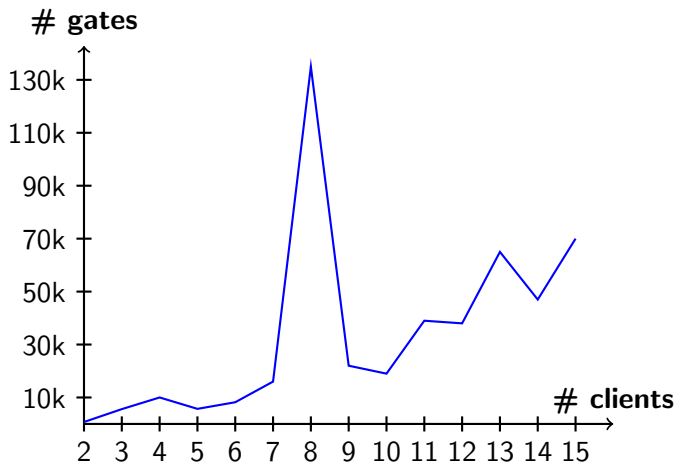


Previous approaches to building the combinational circuit

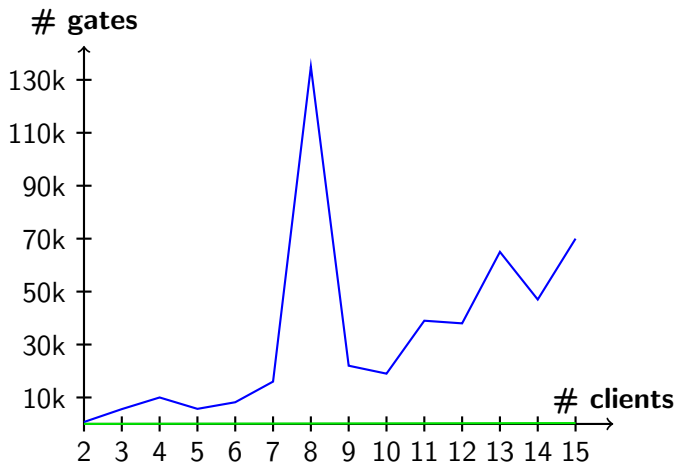
Experiences in the synthesis domain

- Kukula and Shiple (2000) - huge, deep circuits
- Bañeres et al. (2004) - extremely slow on synthesis problems
- Bloem et al. (2007b) - best approach so far, circuits are still large
- Jiang et al. (2009) - generates larger circuits than the approach by Bloem et al. (2007b)
- ...

Example: AMBA AHB bus arbiter (Bloem et al., 2007b)



Example: AMBA AHB bus arbiter (Bloem et al., 2007b)



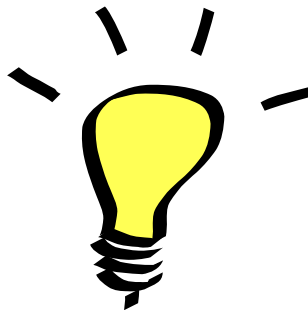
Our new approach

Main idea

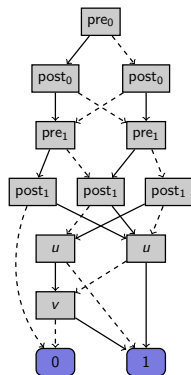
Use **computational learning of Boolean functions** as main tool for computing small circuit implementations.

Benefits

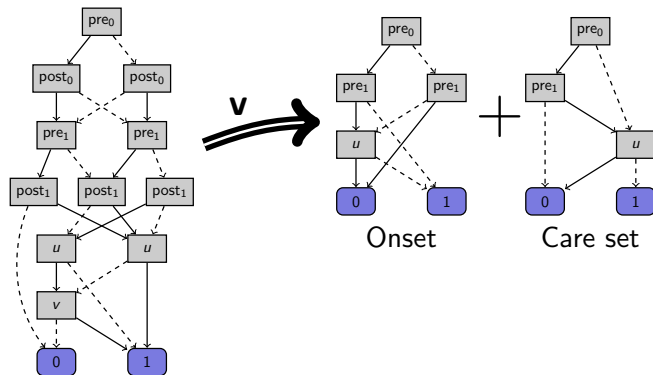
- We can efficiently benefit from a large number of *don't care* cases
- We obtain *shallow* circuits



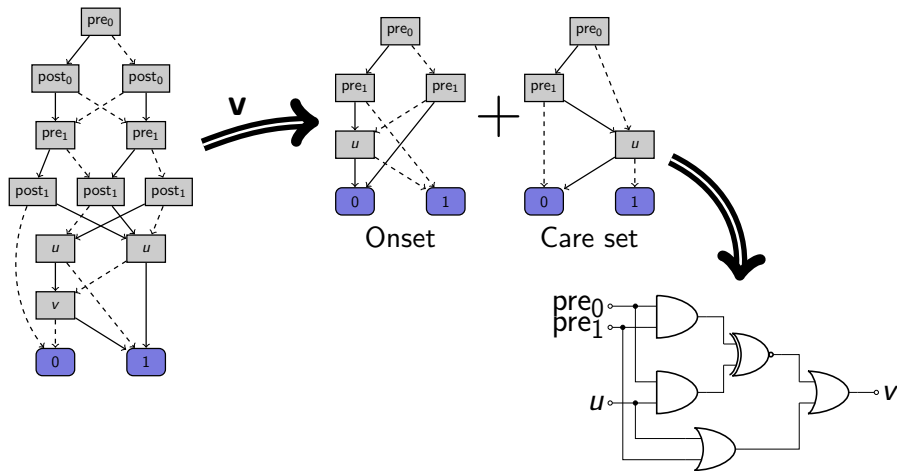
Bit-by-bit circuit extraction approach



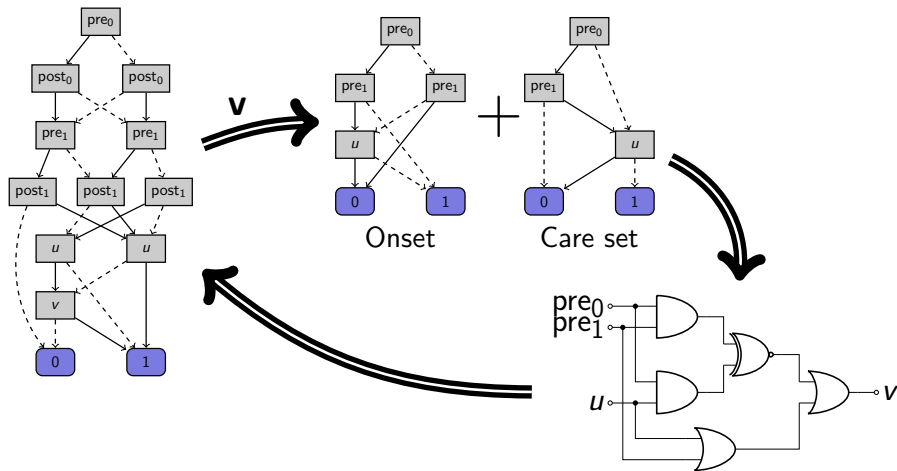
Bit-by-bit circuit extraction approach



Bit-by-bit circuit extraction approach

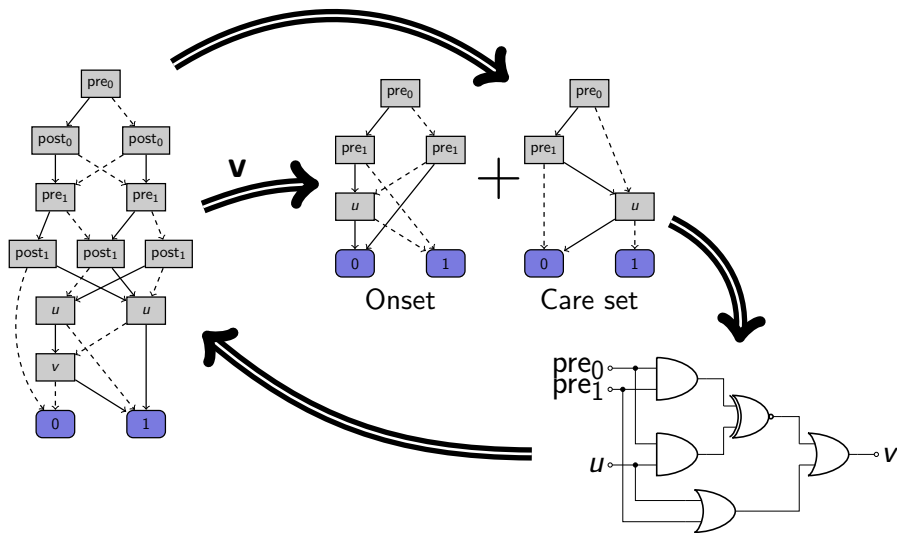


Bit-by-bit circuit extraction approach

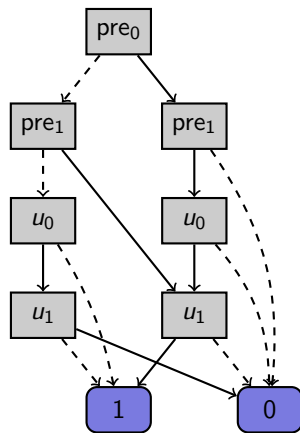


Bit-by-bit circuit extraction approach

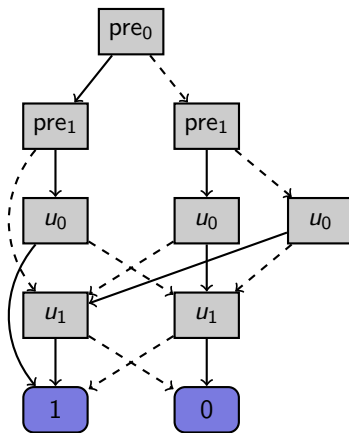
Reachable states/positions



Boolean formula learning (DNF)

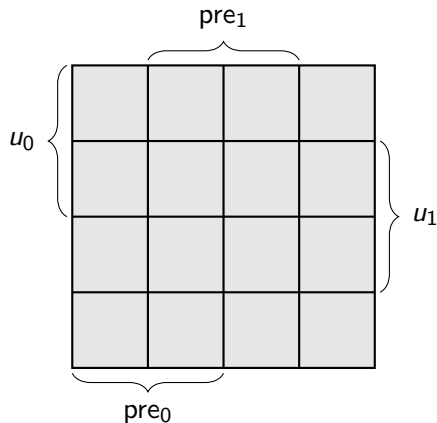


Onset



Care Set

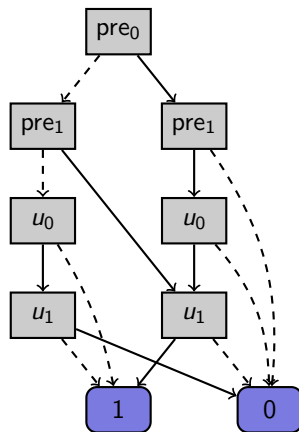
Boolean formula learning (DNF)



Boolean formula learning (DNF)

pre_1			
u_0	0	0	0
	0	1	1
	0	0	1
	0	0	0
pre_0			

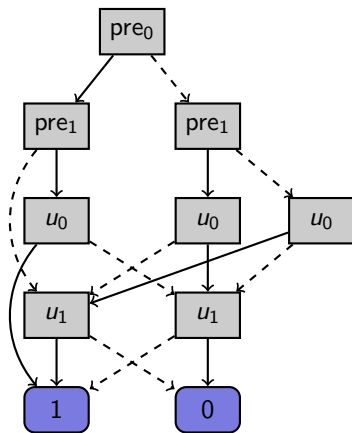
u_1



Onset

Boolean formula learning (DNF)

pre_1					
u_0	X	0	0	X	u_1
	0	1	X	0	
	0	X	1	X	
	X	0	X	1	
				pre_0	



Care Set

Boolean formula learning (DNF)

A 4x4 grid representing a truth table for Boolean formula learning. The grid is partitioned by brackets labeled u_0 , u_1 , pre_0 , and pre_1 .

	pre_1			
u_0	X	0	0	X
	0	1	X	0
	0	X	1	X
	X	0	X	1
	pre_0			u_1

The grid contains the following values:

- Row 1: X, 0, 0, X
- Row 2: 0, 1, X, 0
- Row 3: 0, X, 1, X
- Row 4: X, 0, X, 1

The cells are colored as follows:

- Red: (1,2), (1,3), (2,1), (2,4), (3,1), (4,2)
- Green: (2,2), (3,3), (4,4)
- Grey: (1,1), (1,4), (2,3), (3,2), (3,4), (4,1), (4,3)

Boolean formula learning (DNF)

pre ₁			
u ₀	X	0	0
	0	1	X
u ₁	0	X	1
	X	0	X
pre ₀			

Candidate solution:

Boolean formula learning (DNF)

u_0		pre_1			
		X	0	0	X
u_0		0	1	X	0
		0	X	1	X
		pre_0		u_1	
		X	0	X	1

Candidate solution:

$\psi = \text{false}$

Boolean formula learning (DNF)

		pre_1			
u_0	{	X	0	0	X
		0	1	X	0
u_0	{	0	X	1	X
		X	0	X	1
		pre_0		u_1	

Candidate solution:

$\psi = \text{false}$

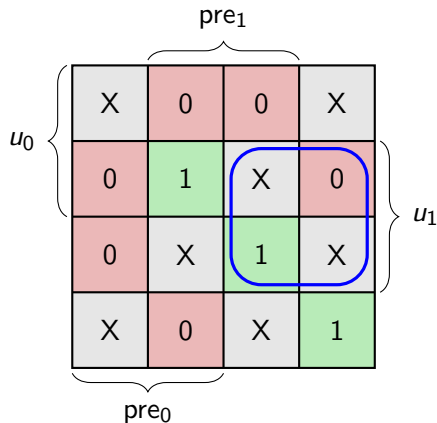
Boolean formula learning (DNF)

u_0		pre_1		
		X	0	
u_0		0	1	u_1
		X	X	
u_0		0	X	u_1
		1	X	
pre_0		X	0	
		X	1	

Candidate solution:

$\psi = \text{false}$

Boolean formula learning (DNF)



Candidate solution:

$\psi = \text{false}$

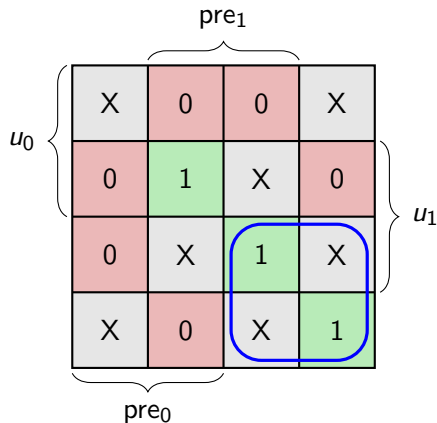
Boolean formula learning (DNF)

u_0				u_1			
pre_1						pre_0	
X	0	0	X	0	1	X	0
0	X	1	X	0	0	1	X
0	X	1	X	0	0	1	X
X	0	X	1	X	0	1	X

Candidate solution:

$\psi = \text{false}$

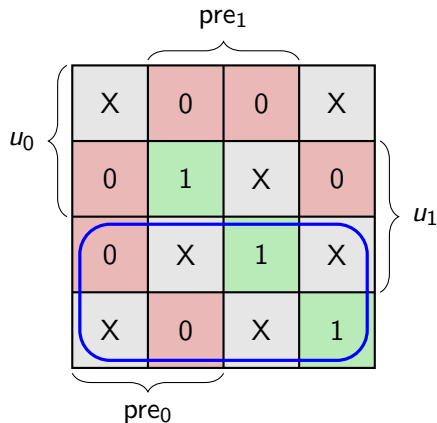
Boolean formula learning (DNF)



Candidate solution:

$\psi = \text{false}$

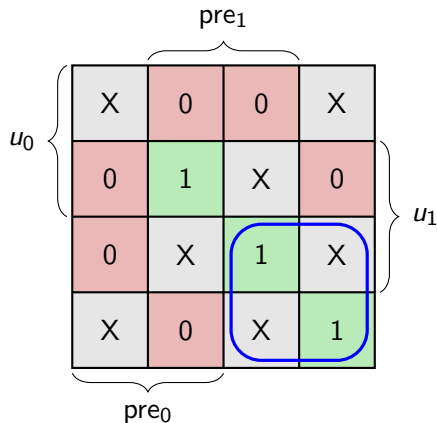
Boolean formula learning (DNF)



Candidate solution:

$\psi = \text{false}$

Boolean formula learning (DNF)



Candidate solution:

$\psi = \text{false}$

Boolean formula learning (DNF)

u_0				u_1			
pre_1							
X	0	0	X	0	1	X	0
0	X	1	X	X	0	1	X
X	0	X	1	X	0	X	1
		pre_0					

Candidate solution:

$$\psi = \neg pre_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

		pre ₁			
u ₀	X	0	0	X	u ₁
	0	1	X	0	
	0	X	X	X	
	X	0	X	X	
		pre ₀			

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

	pre ₁		
u ₀	X	0	0
	0	1	X
	0	X	X
	X	0	X
	pre ₀		
			u ₁

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

pre ₁			
X	0	0	X
0	1	X	0
0	X	X	X
X	0	X	X
pre ₀			

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

		pre ₁			
u ₀	X	0	0	X	u ₁
	0	1	X	0	
	0	X	X	X	
	X	0	X	X	
		pre ₀			

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

				pre ₁	
u ₀	X	0	0	X	u ₁
	0	1	X	0	
	0	X	X	X	
	X	0	X	X	
				pre ₀	

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

		pre ₁			
u ₀	X	0	0	X	u ₁
	0	1	X	0	
	0	X	X	X	
	X	0	X	X	
		pre ₀			

Candidate solution:

$$\psi = \neg \text{pre}_0 \wedge \neg u_0$$

Boolean formula learning (DNF)

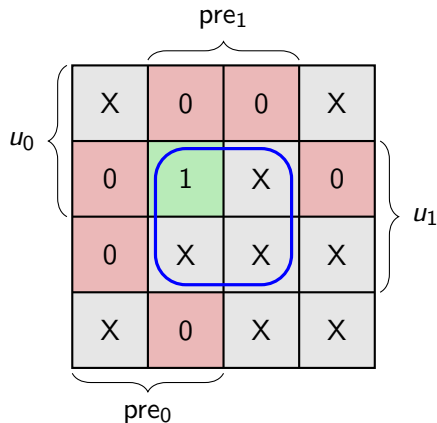
pre_1			
X	0	0	X
0	1	X	0
0	X	X	X
X	0	X	X
pre_0			

u_0 { } u_1 { }

Candidate solution:

$$\psi = \neg pre_0 \wedge \neg u_0$$

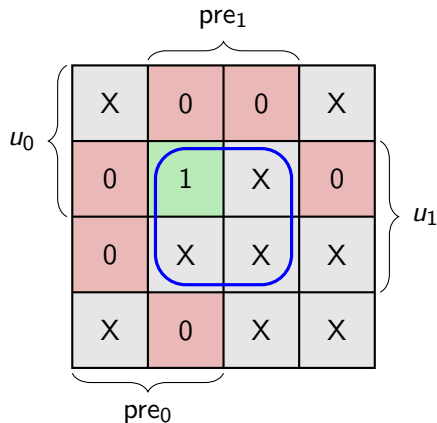
Boolean formula learning (DNF)



Candidate solution:

$$\psi = \neg pre_0 \wedge \neg u_0$$

Boolean formula learning (DNF)



Candidate solution:

$$\psi = (\neg \text{pre}_0 \wedge \neg u_0) \vee (\text{pre}_1 \wedge u_1)$$

Boolean formula learning (DNF)

		pre ₁		
u ₀	X	0	0	X
	0	X	X	0
	0	X	X	X
	X	0	X	X
		pre ₀		u ₁

Candidate solution:

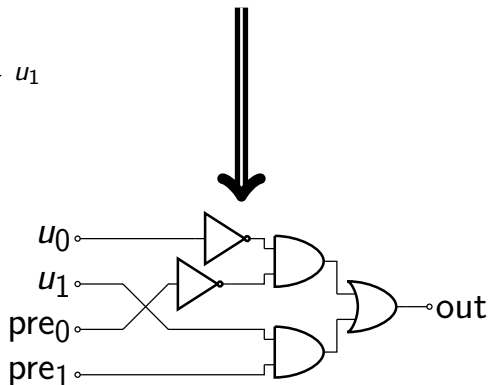
$$\psi = (\neg \text{pre}_0 \wedge \neg u_0) \vee (\text{pre}_1 \wedge u_1)$$

Boolean formula learning (DNF)

				pre_1			
u_0	X	0	0	X	u_1		
	0	X	X	0			
	0	X	X	X			
	X	0	X	X			
				pre_0			

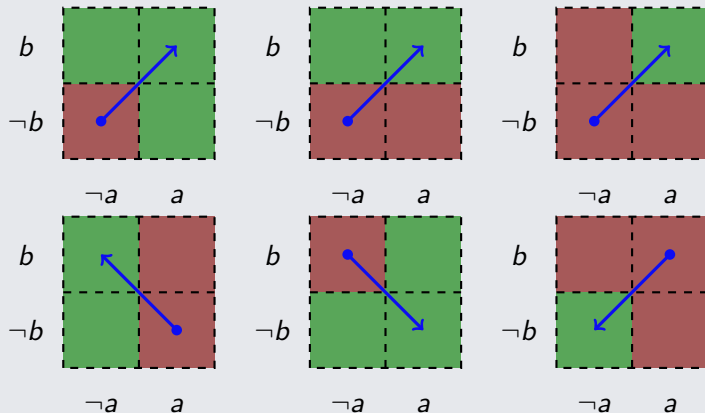
Candidate solution:

$$\psi = (\neg pre_0 \wedge \neg u_0) \vee (pre_1 \wedge u_1)$$



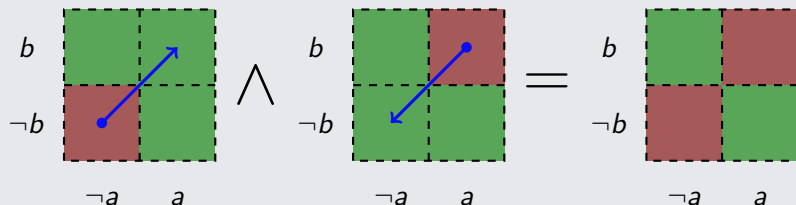
CDNF learning: Bshouty's algorithm

Z-monotone functions



Overlay example: decomposing a XOR into CDNF

Idea



$$(a \vee b) \wedge (\neg a \vee \neg b) = \text{Target function}$$

Algorithm

While target function not found:

- Search for **false-positive** z (as in CNF)
If found, add a new z -based DNF
- Search for **false-negative** z (as in DNF)
If found, update all DNFs to accept z

Example: learning $a \text{ XOR } b$

Candidate solution

(**true**)

Operation

Check for false-positive

Example: learning $a \text{ XOR } b$

Candidate solution

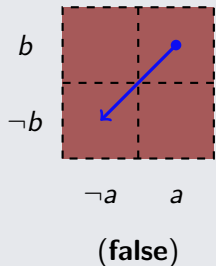
(true)

Operation

Check for false-positive $\rightarrow \mathbf{x}$: ab

Example: learning $a \text{ XOR } b$

Candidate solution

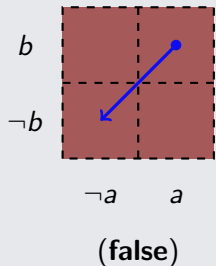


Operation

Check for false-negative

Example: learning $a \text{ XOR } b$

Candidate solution

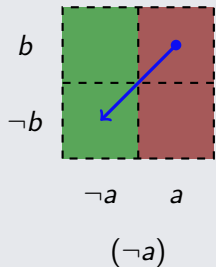


Operation

Check for false-negative $\rightarrow \mathbf{X}$: $\bar{a}b$

Example: learning $a \text{ XOR } b$

Candidate solution

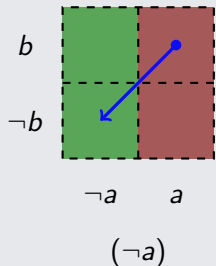


Operation

Check for false-positive

Example: learning $a \text{ XOR } b$

Candidate solution

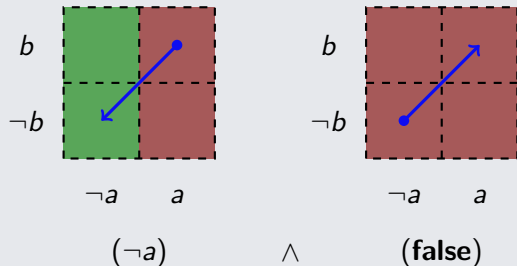


Operation

Check for false-positive $\rightarrow \mathbf{x}: \overline{ab}$

Example: learning $a \text{ XOR } b$

Candidate solution

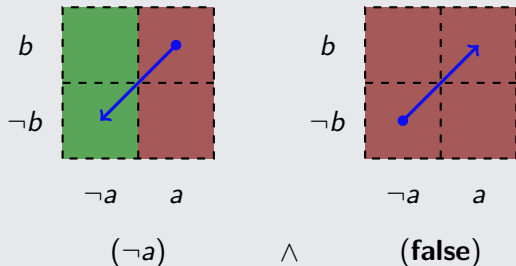


Operation

Check for false-negative

Example: learning $a \text{ XOR } b$

Candidate solution

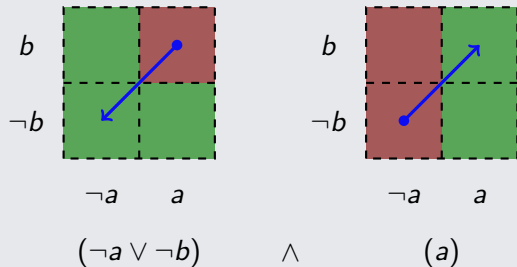


Operation

Check for false-negative $\rightarrow \mathbf{x}: \bar{b}a$

Example: learning $a \text{ XOR } b$

Candidate solution

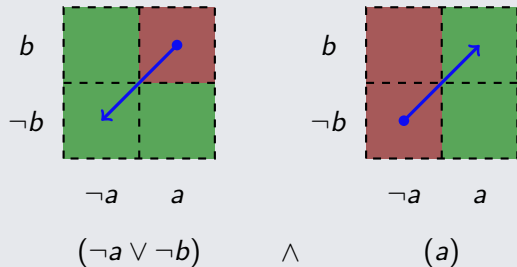


Operation

Check for false-positive

Example: learning $a \text{ XOR } b$

Candidate solution

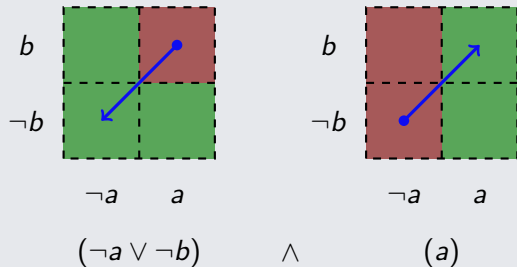


Operation

Check for false-positive $\rightarrow \checkmark$

Example: learning $a \text{ XOR } b$

Candidate solution

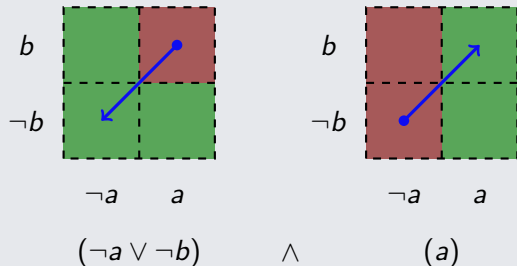


Operation

Check for false-negative

Example: learning $a \text{ XOR } b$

Candidate solution

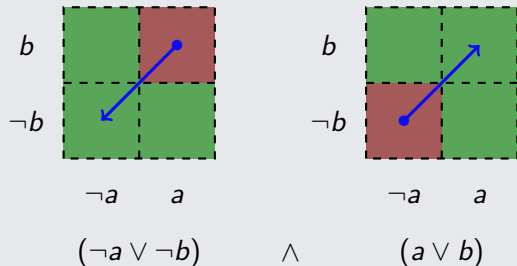


Operation

Check for false-negative $\rightarrow \mathbf{X}: \bar{a}b$

Example: learning $a \text{ XOR } b$

Candidate solution

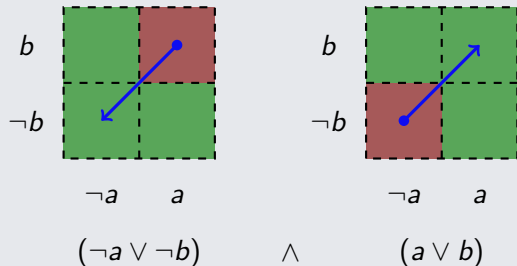


Operation

Check for false-positive

Example: learning $a \text{ XOR } b$

Candidate solution

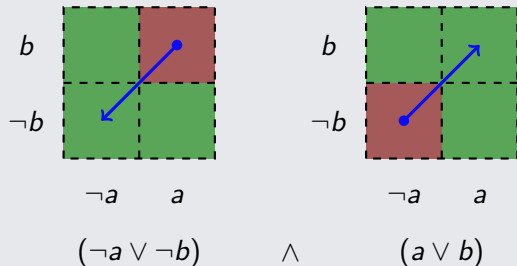


Operation

Check for false-positive $\rightarrow \checkmark$

Example: learning $a \text{ XOR } b$

Candidate solution

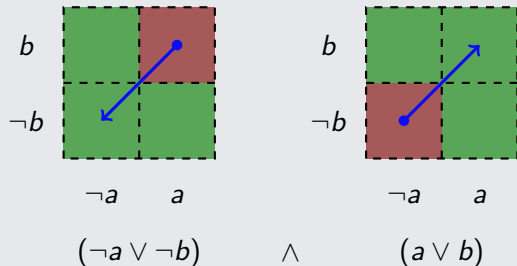


Operation

Check for false-negative

Example: learning $a \text{ XOR } b$

Candidate solution

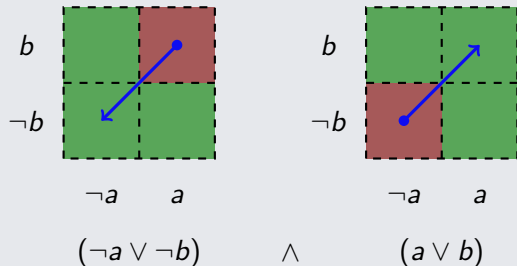


Operation

Check for false-negative $\rightarrow \checkmark$

Example: learning $a \text{ XOR } b$

Candidate solution



Operation

No counter-examples? Done!

Benchmarks

- RATS_Y – GR(1) synthesis
 - IBM generalised buffer (Bloem et al., 2007b)
 - AMBA AHB Arbiter (Bloem et al., 2007a,b)
- UNBEAST – bounded synthesis
 - Load balancer (Ehlers, 2010)
 - Various benchmarks by Jobstmann and Bloem (2006)

Implementation

- Uses CuDD as BDD library
- Written in C++
- Imports a general strategy as CuDD BDD (written to a file)
- Postprocessing/verification: ABC and NuSMV (BMC)

Modes of operation

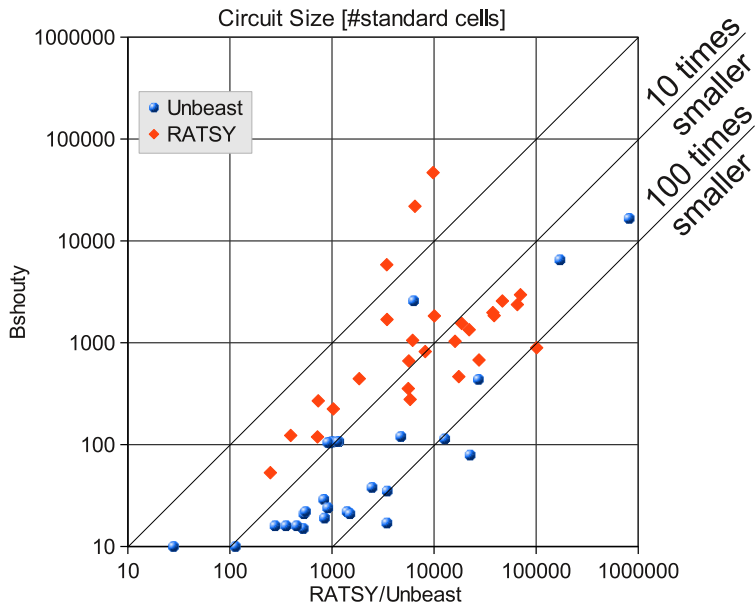
Any combination of ...

- ... DNF/CDNF learning or their dual case, CNF/DCNF learning
- ... Smart variable order selection on/off

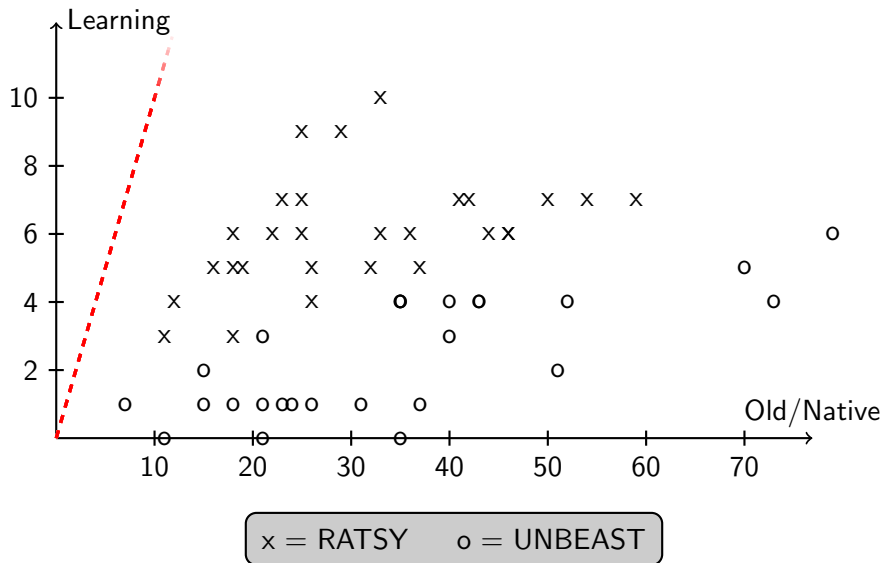
Best modes found

- **CDNF** learning + smart variable order selection **on**
- **DNF** learning + smart variable order selection **on**

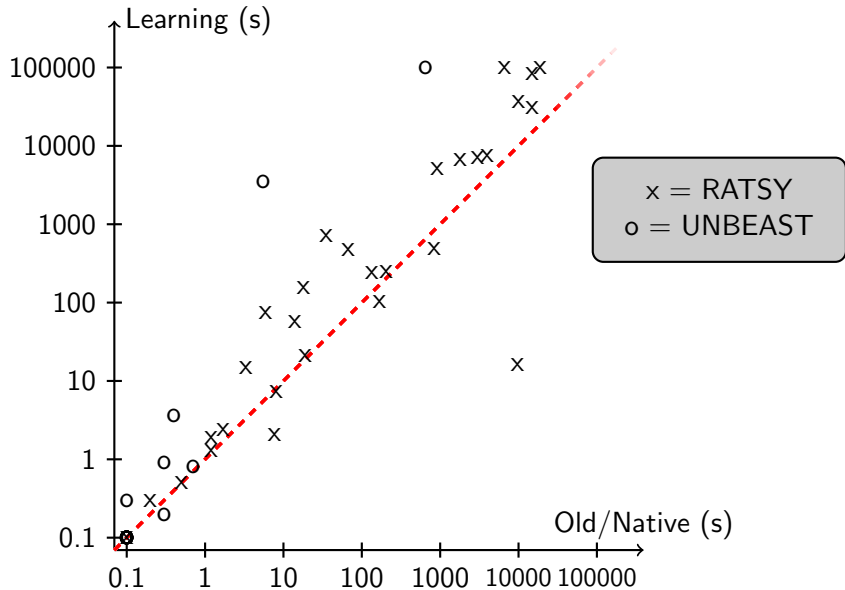
Circuit sizes (CDNF)



Circuit depths (CDNF)



Computation times

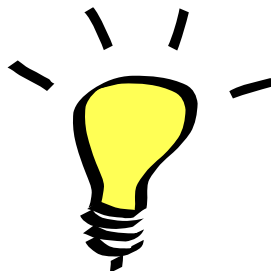


Learning-based implementation extraction:

- an *efficient* and *effective* way to compute circuits in reactive synthesis
- *not restricted to BDDs*, but can be used with *any* symbolic data structure
- better suited for the task than all other techniques that we tried

Implementation available (for CuDD)

The URL is in the paper!



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