

CS243: Discrete Structures

Introduction to Number Theory

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Announcements

- ▶ Midterms are graded – scores on Blackboard
 - ▶ Graded midterms handed back at the end of class
 - ▶ Make sure score on midterm matches grade on Blackboard
 - ▶ If not, let us know asap (within a week at latest)
- ▶ Sample solutions posted on course webpage – look over them!
- ▶ Since only 1-2 students answered Question (4c) correctly and since there was small typo, did not count (4c) when calculating grades – thus, grades out of 80 rather than 90

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Announcements, cont.

- ▶ Average on midterm 55 out of 80; standard deviation is 17
- ▶ Third homework also graded – scores on Blackboard
- ▶ Average on HW3 is 67 out of 100

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Introduction

- ▶ Number theory is the branch of mathematics that deals with integers and their properties
- ▶ Number theory has a number of applications in computer science, esp. in modern **cryptology**
- ▶ **This lecture:** Important results in number theory
- ▶ **Next lecture:** Continue discussion of number theory, look at applications of number theory in cryptology

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Divisibility

- ▶ Given two integers a and b where $a \neq 0$, we say a divides b if there is an integer c such that $b = ac$
- ▶ If a divides b , we write $a|b$; otherwise, $a \nmid b$
- ▶ **Example:** $2|6$, $2 \nmid 9$
- ▶ If $a|b$, a is called a **factor** of b
- ▶ b is called a **multiple** of a
- ▶ All integers divisible by a can be enumerated as:

$\dots, -3a, -2a, -a, 0, a, 2a, 3a, \dots$

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Example

- ▶ **Question:** If n and d are positive integers, how many integers not exceeding n are divisible by d ?
- ▶ **Recall:** All positive integers divisible by d are of the form dk
- ▶ We want to find how many numbers dk there are such that $0 < dk \leq n$.
- ▶ In other words, we want to know how many **integers** k there are such that $0 < k \leq \frac{n}{d}$
- ▶ How many integers are there between 1 and $\frac{n}{d}$?

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Properties of Divisibility

- ▶ **Theorem 1:** If $a|b$ and $a|c$, then $a|(b+c)$
- ▶ **Proof:** If $a|b$, then $\exists k_1$ such that $b = ak_1$
- ▶ Similarly, if $a|c$, then $\exists k_2$ such that $c = ak_2$
- ▶ Then, $b+c = a(k_1+k_2)$
- ▶ Hence, $a|(b+c)$ □

More Divisibility Properties

- ▶ **Theorem 2:** If $a|b$, then $a|bc$ for all integers c
- ▶ **Proof:** If $a|b$, then there exists k such that $b = ak$.
- ▶ Hence, $bc = a \cdot ck$
- ▶ Therefore, $a|bc$ □

Divisibility Properties, cont.

- ▶ **Theorem 3:** If $a|b$ and $b|c$, then $a|c$
- ▶ **Proof:** If $a|b$, there exists k_1 such that $b = ak_1$
- ▶ Since $b|c$, there exists k_2 such that $c = bk_2$.
- ▶ This implies $c = a \cdot k_1k_2$.
- ▶ Hence $a|c$. □

Divisibility Properties, cont.

- ▶ **Theorem 4** If $a|b$ and $a|c$, then $a|(mb+nc)$
- ▶ **Proof:** By Thm 2, if $a|b$, then $a|mb$
- ▶ By thm 2, if $a|c$, then $a|nc$.
- ▶ By Thm 1, if $a|mb$ and $a|nc$, then $a|(mb+nc)$

The Division Theorem

- ▶ **Division theorem:** Let a be an integer, and d a positive integer. Then, there are **unique** integers q, r with $0 \leq r < d$ such that $a = dq + r$
- ▶ Here, d is called **divisor**, and a is called **dividend**
- ▶ q is the **quotient**, and r is the **remainder**.
- ▶ We use the $r = a \bmod d$ notation to express the remainder
- ▶ The notation $q = a \operatorname{div} d$ expresses the quotient
- ▶ What is $101 \bmod 11$?
- ▶ What is $101 \operatorname{div} 11$?

Congruence Modulo

- ▶ Sometimes, we care if two integers a, b have the same remainder when divided by some number m .
- ▶ If so, a and b are **congruent modulo m** , $a \equiv b \pmod{m}$.
- ▶ More technically, if a and b are integers and m a positive integer, $a \equiv b \pmod{m}$ iff $m|(a-b)$
- ▶ **Example:** 7 and 13 are congruent modulo 3.
- ▶ **Example:** Find a number congruent to 7 modulo 4.

Congruence Modulo Theorem

- ▶ **Theorem:** $a \equiv b \pmod{m}$ iff $a \bmod m = b \bmod m$
- ▶ **Part 1, \Rightarrow :** Suppose $a \equiv b \pmod{m}$.
- ▶ Then, by definition of \equiv , $m|(a - b)$
- ▶ By definition of $|$, there exists k such that $a - b = mk$, i.e.,
 $a = b + mk$
- ▶ By division thm, $b = mp + r$ for some $0 \leq r < m$
- ▶ Then, $a = mp + r + mk = m(p + k) + r$
- ▶ Thus, $a \bmod m = r = b \bmod m$

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Congruence Modulo Theorem Proof, cont.

- ▶ **Theorem:** $a \equiv b \pmod{m}$ iff $a \bmod m = b \bmod m$
- ▶ **Part 2, \Leftarrow :** Suppose $a \bmod m = b \bmod m$
- ▶ Then, there exists some p_1, p_2, r such that $a = p_1 \cdot m + r$ and
 $b = p_2 \cdot m + r$ where $0 \leq r < m$
- ▶ Then, $a - b = p_1 \cdot m + r - p_2 \cdot m - r = m \cdot (p_1 - p_2)$
- ▶ Thus, $m|(a - b)$
- ▶ By definition of \equiv , $a \equiv b \pmod{m}$

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Example

- ▶ Prove that if $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then:

$$a + c \equiv b + d \pmod{m}$$

- ▶ **Proof:** Since $a \equiv b \pmod{m}$, $m|(a - b)$
- ▶ Since $c \equiv d \pmod{m}$, $m|(c - d)$
- ▶ By definition of $|$, there exists k_1, k_2 such that:

$$a - b = mk_1 \quad c - d = mk_2$$

- ▶ Adding these, we get: $a + c - (b + d) = m(k_1 + k_2)$
- ▶ Again, by definition of \equiv and $|$, this means
 $a + c \equiv b + d \pmod{m}$ □

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Applications of Congruence in Cryptography

- ▶ Congruences have many applications in cryptography
- ▶ For instance, Julius Caesar **encrypted** messages by shifting each letter three letters in the alphabet ("Caesar cipher")
- ▶ For example, the message "**I LIKE DISCRETE MATH**" would be encrypted as "**L OLNH GLYFUHVH PDVK**"
- ▶ Caesar's cipher example of **shift cipher**: shifts each letter by k
- ▶ For Caesar cipher, $k = 3$
- ▶ We can express shift ciphers using the modulo operator

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Shift Ciphers

- ▶ First, let's number letters A-Z with $0 - 25$
- ▶ Represent message with sequence of numbers
- ▶ **Example:** The sequence "25 0 2" represents "ZAC"
- ▶ To encrypt, apply **encryption function** f defined as:

$$f(x) = (x + k) \bmod 26$$

- ▶ Because f is bijective, its inverse yields decryption function:

$$f^{-1}(x) = (x - k) \bmod 26$$

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Ciphers and Congruence Modulo

- ▶ Shift cipher is a very primitive and insecure cipher because very easy to infer what k is
- ▶ But contains some useful ideas:
 - ▶ Encoding words as sequence of numbers
 - ▶ Use of modulo operator
- ▶ Modern encryption schemes much more sophisticated, but also share these principles
- ▶ More on this next lecture!

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Prime Numbers

- ▶ A positive integer p that is greater than 1 and divisible only by 1 and itself is called a **prime number**.
- ▶ **First few primes:** 2, 3, 5, 7, 11, ...
- ▶ A positive integer that is greater than 1 and that is not prime is called a **composite number**
- ▶ **Example:** 2, 4, 6, 8, 9, ...

Fundamental Theorem of Arithmetic

- ▶ **Fundamental Thm:** Every positive integer greater than 1 is either prime or can be written uniquely as a product of primes.
- ▶ This unique product of prime numbers for x is called the **prime factorization** of x
- ▶ **Examples:**
 - ▶ 12 =
 - ▶ 21 =
 - ▶ 99 =

Determining Prime-ness

- ▶ In many applications, such as crypto, important to determine if a number is prime – following thm is useful for this:
- ▶ **Theorem:** If n is composite, then it has a prime divisor less than or equal to \sqrt{n}
- ▶ **Proof:** Since n is composite, it can be written as $n = ab$ where $a > 1$ and $b > 1$.
- ▶ For contradiction, suppose neither a nor b are $\leq \sqrt{n}$, i.e., $a > \sqrt{n}$, $b > \sqrt{n}$
- ▶ Then, $n = ab > \sqrt{n}^2 = n$, a contradiction.
- ▶ Hence, either $a \leq \sqrt{n}$, or $b \leq \sqrt{n}$, and by the Fundamental Thm, is either itself a prime or has a factor less than itself. \square

Consequence of This Theorem

- Theorem:** If n is composite, then it has a prime divisor $\leq \sqrt{n}$
- ▶ Thus, to determine if n is prime, only need to check if it is divisible by primes $\leq \sqrt{n}$
 - ▶ **Example:** Show that 101 is prime
 - ▶ Since $\sqrt{101} < 11$, only need to check if it is divisible by 2, 3, 5, 7.
 - ▶ Since it is not divisible by any of these, we know it is prime.

Infinitely Many Primes

- ▶ **Theorem:** There are infinitely many prime numbers.
- ▶ **Proof:** (by contradiction) Suppose there are finitely many primes: p_1, p_2, \dots, p_n
- ▶ Now consider the number $Q = p_1 p_2 \dots p_n + 1$. Q is either prime or composite
- ▶ **Case 1:** Q is prime. We get a contradiction, because we assumed only prime numbers are p_1, \dots, p_n
- ▶ **Case 2:** Q is composite. In this case, Q can be written as product of primes.
- ▶ But Q is not divisible by any of p_1, p_2, \dots, p_n
- ▶ Hence, by Fundamental Thm, not composite $\Rightarrow \perp$ \square

A Word about Prime Numbers and Cryptography

- ▶ Prime numbers play a key role in modern cryptography
- ▶ Modern cryptography techniques rely on prime numbers to encrypt messages
- ▶ Security of encryption relies on prime factorization being intractable for sufficiently large numbers
- ▶ More on this later...

Greatest Common Divisors

- ▶ Suppose a and b are integers, not both 0.
- ▶ Then, the largest integer d such that $d|a$ and $d|b$ is called **greatest common divisor** of a and b , written $\gcd(a,b)$.
- ▶ Example: $\gcd(24, 36) =$
- ▶ Example: $\gcd(2^3 5, 2^2 3) =$
- ▶ Example: $\gcd(14, 25) =$
- ▶ Two numbers whose gcd is 1 are called **relatively prime**.
- ▶ Example: 14 and 25 are relatively prime

Least Common Multiple

- ▶ The **least common multiple** of a and b , written $\text{lcm}(a,b)$, is the smallest integer c such that $a|c$ and $b|c$.
- ▶ Example: $\text{lcm}(9, 12) =$
- ▶ Example: $\text{lcm}(2^3 3^5 7^2, 2^4 3^3) =$

Theorem about LCM and GCD

- ▶ **Theorem:** Let a and b be positive integers. Then, $ab = \gcd(a, b) \cdot \text{lcm}(a, b)$
- ▶ **Proof:** Let $a = p_1^{i_1} p_2^{i_2} \dots p_n^{i_n}$ and $b = p_1^{j_1} p_2^{j_2} \dots p_n^{j_n}$
- ▶ Then, $ab = p_1^{i_1+j_1} p_2^{i_2+j_2} \dots p_n^{i_n+j_n}$
- ▶ $\gcd(a, b) = p_1^{\min(i_1, j_1)} p_2^{\min(i_2, j_2)} \dots p_n^{\min(i_n, j_n)}$
- ▶ $\text{lcm}(a, b) = p_1^{\max(i_1, j_1)} p_2^{\max(i_2, j_2)} \dots p_n^{\max(i_n, j_n)}$
- ▶ Thus, we need to show $i_k + j_k = \min(i_k, j_k) + \max(i_k, j_k)$

Proof, cont.

- ▶ Show $i_k + j_k = \min(i_k, j_k) + \max(i_k, j_k)$
- ▶ Either (i) $i_k < j_k$ or (ii) $i_k \geq j_k$
- ▶ If (i), then $\min(i_k, j_k) = i_k$ and $\max(i_k, j_k) = j_k$
- ▶ Thus, $i_k + j_k = \min(i_k, j_k) + \max(i_k, j_k)$
- ▶ If (ii), then $\min(i_k, j_k) = j_k$ and $\max(i_k, j_k) = i_k$
- ▶ Hence $\min(i_k, j_k) + \max(i_k, j_k) = i_k + j_k$ □

Computing GCDs

- ▶ Simple algorithm to compute gcd of a, b :
 - ▶ Factorize a as $p_1^{i_1} p_2^{i_2} \dots p_n^{i_n}$
 - ▶ Factorize b as $p_1^{j_1} p_2^{j_2} \dots p_n^{j_n}$
 - ▶ $\gcd(a, b) = p_1^{\min(i_1, j_1)} p_2^{\min(i_2, j_2)} \dots p_n^{\min(i_n, j_n)}$
- ▶ But this algorithm is not very practical because prime factorization is **computationally expensive!**
- ▶ Much more efficient algorithm to compute gcd, called the **Euclidian algorithm**

Insight Behind Euclid's Algorithm

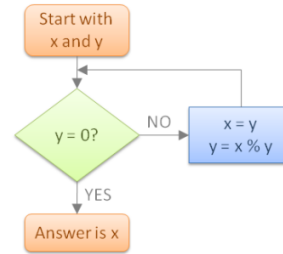
- ▶ **Theorem:** Let $a = bq + r$. Then, $\gcd(a, b) = \gcd(b, r)$
- ▶ **Proof:** We'll show that a, b and b, r have the same common divisors – implies they have the same gcd.
- ⇒ Suppose d is a common divisor of a, b , i.e., $d|a$ and $d|b$
 - ▶ By theorem we proved earlier, this implies $d|a - bq$
 - ▶ Since $a - bq = r$, $d|r$. Hence d is common divisor of b, r .
- ⇐ Now, suppose $d|b$ and $d|r$. Then, $d|bq + r$
 - ▶ Hence, $d|a$ and d is common divisor of a, b □

Using this Theorem

Theorem: Let $a = bq + r$. Then, $\gcd(a, b) = \gcd(b, r)$

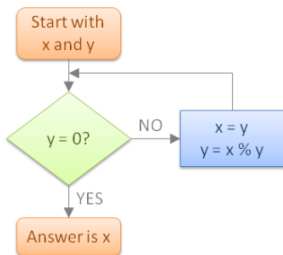
- ▶ Theorem suggests following strategy to compute $\gcd(a, b)$:
- ▶ Compute $r_1 = a \bmod b$ ($\gcd(a, b) = \gcd(b, r_1)$)
- ▶ Compute $r_2 = b \bmod r_1$ ($\gcd(a, b) = \gcd(r_1, r_2)$)
- ▶ Compute $r_3 = r_1 \bmod r_2$ ($\gcd(a, b) = \gcd(r_2, r_3)$)
- ▶ Repeat until remainder becomes 0
($\gcd(a, b) = \gcd(r_n, 0) = r_n$)
- ▶ The last non-zero remainder is the gcd of a and b !

Euclidian Algorithm



- ▶ Find gcd of 72 and 20
- ▶ $12 = 72\%20$
- ▶ $8 = 20\%12$
- ▶ $4 = 12\%8$
- ▶ $0 = 8\%4$
- ▶ gcd is 4!

Euclidian Algorithm Example



- ▶ Find gcd of 662 and 414
- ▶ $248 = 662\%414$
- ▶ $166 = 414\%248$
- ▶ $82 = 248\%166$
- ▶ $2 = 166\%82$
- ▶ $0 = 82\%2$
- ▶ gcd is 2!

GCD as Linear Combination

- ▶ $\gcd(a, b)$ can be expressed as a **linear combination** of a and b
- ▶ **Theorem:** If a and b are positive integers, then there exist integers s and t such that:

$$\gcd(a, b) = s \cdot a + t \cdot b$$

- ▶ Furthermore, Euclidian algorithm gives us a way to compute these integers s and t

Example

- ▶ Express $\gcd(252, 198)$ as a linear combination of 252 and 198
- ▶ First apply Euclid's algorithm (write $a = bq + r$ at each step):
 - $252 = 1 \cdot 198 + 54$
 - $198 = 3 \cdot 54 + 36$
 - $54 = 1 \cdot 36 + 18$
 - $36 = 2 \cdot 18 + 0 \Rightarrow$ gcd is 18
- ▶ Now, using (3), write 18 as $54 - 1 \cdot 36$
- ▶ Using (2), write 18 as $54 - 1 \cdot (198 - 3 \cdot 54)$
- ▶ Using (1), we have $54 = 252 - 1 \cdot 198$, thus:

$$18 = (252 - 1 \cdot 198) - 1(198 - 3 \cdot (252 - 1 \cdot 198))$$

Example, cont.

- $$18 = (252 - 1 \cdot 198) - 1(198 - 3 \cdot (252 - 1 \cdot 198))$$
- ▶ Now, let's simplify this:

$$18 = 252 - 1 \cdot 198 - 1 \cdot 198 + 3 \cdot 252 - 3 \cdot 198$$
 - ▶ Now, collect all 252 and 198 terms together:

$$18 = 4 \cdot 252 - 5 \cdot 198$$
 - ▶ Trace steps of Euclid's algorithm backwards to derive s, t :

$$\gcd(a, b) = s \cdot a + t \cdot b$$
 - ▶ This is known as the **extended Euclidian algorithm**

A Useful Result

- ▶ **Lemma:** If a, b are relatively prime and $a|bc$, then $a|c$.
- ▶ **Proof:** Since a, b are relatively prime $\gcd(a, b) = 1$
- ▶ By previous theorem, there exists s, t such that $1 = s \cdot a + t \cdot b$
- ▶ Multiply both sides by c : $c = csa + ctb$
- ▶ By earlier theorem, since $a|bc$, $a|ctb$
- ▶ Also, by earlier theorem, $a|csa$
- ▶ Therefore, $a|csa + ctb$, which implies $a|c$ since $c = csa + ctb$ \square

Example

Lemma: If a, b are relatively prime and $a|bc$, then $a|c$.

- ▶ Suppose $15 \mid 16 \cdot x$
- ▶ Here 15 and 16 are relatively prime
- ▶ Thus, previous theorem implies: $15 \mid x$

Question

- ▶ Suppose $ca \equiv cb \pmod{m}$. Does this imply $a \equiv b \pmod{m}$?
- ▶ **Counterexample:** Consider $14 \equiv 8 \pmod{6}$
- ▶ Thus, $2 \cdot 7 \equiv 2 \cdot 4 \pmod{6}$
- ▶ But $7 \not\equiv 4 \pmod{6}$
- ▶ Therefore, this implication does not hold in the general case!
- ▶ However, if c and m are relatively prime, it does hold

Another Useful Result

- ▶ **Theorem:** If $ca \equiv cb \pmod{m}$ and $\gcd(c, m) = 1$, then $a \equiv b \pmod{m}$
- ▶ **Proof:** Since $ca \equiv cb \pmod{m}$, we have $m \mid ca - cb$
- ▶ Rewriting, we get: $m \mid c(a - b)$
- ▶ Since m, c are relatively prime, previous thm implies $m \mid a - b$
- ▶ By definition of congruence, $a \equiv b \pmod{m}$

Examples

- ▶ If $15x \equiv 15y \pmod{4}$, is $x \equiv y \pmod{4}$?
- ▶ If $8x \equiv 8y \pmod{4}$, is $x \equiv y \pmod{4}$?
- ▶ **Counterexample:** $8 \cdot 2 \equiv 8 \cdot 3 \pmod{4}$, but $2 \not\equiv 3 \pmod{4}$