

## CS311H: Discrete Mathematics

### More Number Theory

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## Review

- ▶ What is the fundamental theorem of arithmetic?
- ▶ Complete the following:
  - ▶  $\gcd(a, b) \cdot \text{lcm}(a, b) = ?$
  - ▶  $\gcd(a, b) = \gcd(b, ?)$
- ▶ Why is the previous equation useful?
- ▶  $\gcd(a, b)$  can be expressed as  $\dots$  ?
- ▶ What does it mean for  $a, b$  to be relatively prime?

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## A Useful Result

- ▶ Suppose  $a|bc$ . Does that mean  $a|c$ ?
- ▶ **Lemma:** If  $a, b$  are relatively prime and  $a|bc$ , then  $a|c$ .
- ▶ **Proof:** Since  $a, b$  are relatively prime  $\gcd(a, b) = 1$
- ▶ By previous theorem, there exists  $s, t$  such that  $1 = s \cdot a + t \cdot b$
- ▶ Multiply both sides by  $c$ :  $c = csa + ctb$
- ▶ By earlier theorem, since  $a|bc$ ,  $a|ctb$
- ▶ Also, by earlier theorem,  $a|csa$
- ▶ Therefore,  $a|csa + ctb$ , which implies  $a|c$  since  $c = csa + ctb$   $\square$

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## Example

**Lemma:** If  $a, b$  are relatively prime and  $a|bc$ , then  $a|c$ .

- ▶ Suppose  $15 \mid 16 \cdot x$
- ▶ Here 15 and 16 are relatively prime
- ▶ Thus, previous theorem implies:  $15 \mid x$

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## Question

- ▶ Suppose  $ca \equiv cb \pmod{m}$ . Does this imply  $a \equiv b \pmod{m}$ ?
- ▶
- ▶
- ▶
- ▶
- ▶

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## Another Useful Result

- ▶ **Theorem:** If  $ca \equiv cb \pmod{m}$  and  $\gcd(c, m) = 1$ , then  $a \equiv b \pmod{m}$
- ▶
- ▶
- ▶
- ▶

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## Examples

- ▶ If  $15x \equiv 15y \pmod{4}$ , is  $x \equiv y \pmod{4}$ ?
- ▶ If  $8x \equiv 8y \pmod{4}$ , is  $x \equiv y \pmod{4}$ ?
- ▶

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## Linear Congruences

- ▶ A congruence of the form  $ax \equiv b \pmod{m}$  where  $a, b, m$  are integers and  $x$  a variable is called a **linear congruence**.
- ▶ Given such a linear congruence, often need to answer:
  1. Are there any solutions?
  2. What are the solutions?
- ▶ **Example:** Does  $8x \equiv 2 \pmod{4}$  have any solutions?
- ▶ **Example:** Does  $8x \equiv 2 \pmod{7}$  have any solutions?
- ▶ **Question:** Is there a systematic way to solve linear congruences?

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## Determining Existence of Solutions

- ▶ **Theorem:** The linear congruence  $ax \equiv b \pmod{m}$  has solutions iff  $\gcd(a, m) \mid b$ .
- ▶ Proof involves two steps:
  1. If  $ax \equiv b \pmod{m}$  has solutions, then  $\gcd(a, m) \mid b$ .
  2. If  $\gcd(a, m) \mid b$ , then  $ax \equiv b \pmod{m}$  has solutions.
- ▶ First prove (1), then (2).

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## Proof, Part I

If  $ax \equiv b \pmod{m}$  has solutions, then  $\gcd(a, m) \mid b$ .

- ▶
- ▶
- ▶
- ▶
- ▶
- ▶

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## Proof, Part II

If  $\gcd(a, m) \mid b$ , then  $ax \equiv b \pmod{m}$  has solutions.

- ▶ Let  $d = \gcd(a, m)$  and suppose  $d \mid b$
- ▶ Then, there is a  $k$  such that  $b = dk$
- ▶ By earlier theorem, there exist  $s, t$  such that  $d = s \cdot a + t \cdot m$
- ▶ Multiply both sides by  $k$ :  $dk = a \cdot (sk) + m \cdot (tk)$
- ▶ Since  $b = dk$ , we have  $b = a \cdot (sk) + m \cdot (tk)$
- ▶ Thus,  $b \equiv a \cdot (sk) \pmod{m}$
- ▶ Hence,  $sk$  is a solution.  $\square$

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## Examples

- ▶ Does  $5x \equiv 7 \pmod{15}$  have any solutions?
- ▶ Does  $3x \equiv 4 \pmod{7}$  have any solutions?

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## Finding Solutions

- ▶ Can determine existence of solutions, but how to find them?
- ▶ **Theorem:** Let  $d = \gcd(a, m) = sa + tm$ . If  $d|b$ , then the solutions to  $ax \equiv b \pmod{m}$  are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

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## Example

Let  $d = \gcd(a, m) = sa + tm$ . If  $d|b$ , then the solutions to  $ax \equiv b \pmod{m}$  are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

- ▶ What are the solutions to the linear congruence  $3x \equiv 4 \pmod{7}$ ?

▶

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## Another Example

Let  $d = \gcd(a, m) = sa + tm$ . If  $d|b$ , then the solutions to  $ax \equiv b \pmod{m}$  are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

- ▶ What are the solutions to the linear congruence  $3x \equiv 1 \pmod{7}$ ?

▶

▶

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## Inverse Modulo $m$

- ▶ The **inverse of  $a$  modulo  $m$** , written  $\bar{a}$  has the property:

$$a\bar{a} \equiv 1 \pmod{m}$$

- ▶ **Theorem:** Inverse of  $a$  modulo  $m$  exists if and only if  $a$  and  $m$  are relatively prime.
- ▶ **Proof:** Inverse must satisfy  $ax \equiv 1 \pmod{m}$
- ▶
- ▶
- ▶ Does 3 have an inverse modulo 7?

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## Example

- ▶ Find an inverse of 3 modulo 7.
- ▶ An inverse is any solution to  $3x \equiv 1 \pmod{7}$
- ▶ Earlier, we already computed solutions for this equation as:

$$x = -2 + 7u$$

- ▶ Thus,  $-2$  is an inverse of 3 modulo 7
- ▶  $5, 12, -9, \dots$  are also inverses

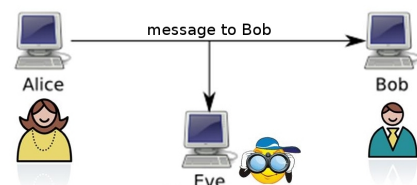
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## Cryptography

- ▶ Cryptography is the study of techniques for secure transmission of information in the presence of adversaries



- ▶ How can Alice send secret messages to Bob without Eve being able to read them?

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## Private vs. Public Crypto Systems

- ▶ Two different kinds of cryptography systems:
  1. Private key cryptography (also known as **symmetric**)
  2. Public key cryptography (**asymmetric**)
- ▶ In private key cryptography, sender and receiver agree on **secret key** that both use to encrypt/decrypt the message
- ▶ In public key cryptography, a **public key** is used to encrypt the message, and **private key** is used to decrypt the message

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## Private Key Cryptography

- ▶ Private key crypto is classical method, used since antiquity
- ▶ Caesar's cipher is an example of private key cryptography
- ▶ Caesar's cipher is **shift cipher** where  $f(p) = (p + k) \pmod{26}$
- ▶ Both receiver and sender need to know  $k$  to encrypt/decrypt
- ▶ Modern symmetric algorithms: RC4, DES, AES, ...
- ▶ **Main problem:** How do you exchange secret key in a secure way?

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## Public Key Cryptography

- ▶ Public key cryptography is the modern method: different keys are used to encrypt vs. decrypt message
- ▶ Most commonly used public key system is **RSA**
- ▶ Great application of number theory and things we've learned

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## RSA History



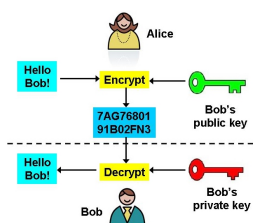
- ▶ Named after its inventors Rivest, Shamir, and Adleman, all researchers at MIT (1978)
- ▶ Actually, similar system invented earlier by British researcher Clifford Cocks, but classified – unknown until 90's

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## RSA Overview



- ▶ Bob has two keys: public and private
- ▶ Everyone knows Bob's public key, but only he knows his private key
- ▶ Alice encrypts message using Bob's public key
- ▶ Bob decrypts message using private key
- ▶ Since public key cannot decrypt, no one can read message except Bob

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## High Level Math Behind RSA

- ▶ In the RSA system, **private key** consists of two **very large prime numbers**  $p, q$
- ▶ **Public key** consists of a number  $n$ , which is the product of  $p, q$  and another number  $e$ , which is relatively prime with  $(p - 1)(q - 1)$
- ▶ Encrypt messages using  $n, e$ , but to decrypt, must know  $p, q$
- ▶ In theory, can extract  $p, q$  from  $n$  using **prime factorization**, but this is intractable for very large numbers
- ▶ **Security of RSA relies on inherent computational difficulty of prime factorization**

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## Encryption in RSA

- ▶ To send message to Bob, Alice first represents message as a sequence of numbers
- ▶ Call this number representing message  $M$
- ▶ Alice then uses Bob's public key  $n, e$  to perform encryption as:

$$C = M^e \pmod{n}$$

- ▶  $C$  is called the **ciphertext**

## RSA Decryption

- ▶ **Decryption key**  $d$  is the inverse of  $e$  modulo  $(p-1)(q-1)$ :

$$d \cdot e \equiv 1 \pmod{(p-1)(q-1)}$$

- ▶ **Decryption function:**  $C^d \pmod{n}$
- ▶ As we saw earlier,  $d$  can be computed reasonably efficiently if we know  $(p-1)(q-1)$
- ▶ However, since adversaries do not know  $p, q$ , they cannot compute  $d$  with reasonable computational effort!

## Security of RSA

- ▶ The encryption function used in RSA is a **trapdoor function**
- ▶ Trapdoor function is easy to compute in one direction, but very difficult in reverse direction without additional knowledge
- ▶ Decryption without private key is very hard because requires prime factorization (which is intractable for large enough numbers)
- ▶ **Interesting fact:** There are efficient (poly-time) prime factorization algorithms for quantum computers (e.g., Shor's algorithm)
- ▶ If we could build quantum computers with sufficient "qubits", RSA would no longer be secure!