

CS380C Compilers

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Introduction

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Today's Plan

- ◆ Motivation
 - ◆ Why study compilers?
- ◆ Let's get started
 - ◆ Look at some sample optimizations and assorted issues
- ◆ A few administrative matters
 - ◆ Course details

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Motivation

- ◆ Q: Why study compilers?

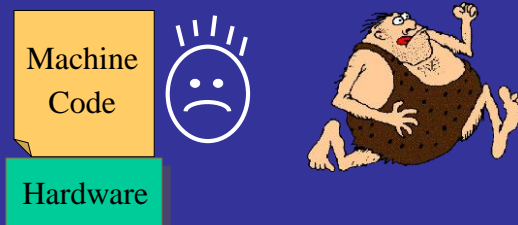
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Life B.C.

- ◆ Before compilers



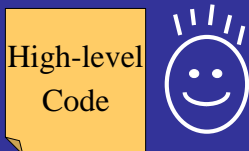
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Liberation

Along came Backus



Hardware

Compilers liberate the programmer from the machine

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Traditional View of Compilers

- ◆ Translate high-level language to machine code
- ◆ High-level programming languages
 - ◆ Increase programmer productivity
 - ◆ Improve program maintenance
 - ◆ Improve portability
- ◆ Low-level architectural details
 - ◆ Instruction set
 - ◆ Addressing modes
 - ◆ Registers, cache, and the rest of the memory hierarchy
 - ◆ Pipelines, instruction-level parallelism

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Optimization

- ◆ Translation is not enough
 - ◆ Backus recognized the importance of obtaining good performance
- ◆ Can perform tedious optimizations that programmers won't do

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Consider Matrix Multiplication

- ◆ Obvious code

```
for i = 1 to n
  for j = 1 to n
    for k = 1 to n
      c[i,j] = c[i, j] + a[i, k]* b[k,j]
```

- ◆ Tiled code– can be significantly faster

```
for it = 1 to n by t
  for jt = 1 to n by t
    for kt = 1 to n by t
      for i = it to it+t-1
        for j = jt to jt+t-1
          for k = kt to kt+t-1
            c[i,j] = c[i, j] + a[i, k]* b[k,j]
```

Why don't we want
programmers to write this code?

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Translation + Optimization

- ◆ Enable language design to flourish
 - ◆ Functional languages
 - ◆ Object oriented languages
 - ◆ . . .

Compilers liberate language designers

- ◆ Logic languages

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Isn't Compilation A Solved Problem?

- ◆ “Optimization for scalar machines is a problem that was solved ten years ago”
-- David Kuck, 1990
- ◆ Languages keep changing
 - ◆ Wacky ideas (*e.g.*, OOP and GC) have gone mainstream
- ◆ Machines keep changing
 - ◆ New features present new problems (*e.g.*, MMX, IA64, trace caches)
 - ◆ Changing costs lead to different concerns (*e.g.*, loads)
- ◆ Applications keep changing
 - ◆ Interactive, real-time, mobile

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Isn't Compilation A Solved Problem? (cont)

- ◆ Values keep changing
- ◆ We used to just care about run-time performance
- ◆ Now?
 - ◆ Compile-time performance
 - ◆ Code size
 - ◆ Correctness
 - ◆ Energy consumption
 - ◆ Security
 - ◆ Fault tolerance

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Value-Added Compilation

- ◆ The more we rely on software, the more we demand more of it
- ◆ Compilers can help— **treat code as data**
 - ◆ Analyze the code
- ◆ Correctness
- ◆ Security

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Correctness and Security

- ◆ Can we check whether pointers and addresses are valid?
- ◆ Can we detect when untrusted code accesses a sensitive part of a system?
- ◆ Can we detect whether locks are used properly?
- ◆ Can we use compilers to certify that code is correct?
- ◆ Can we use compilers to verify that a given compiler transformation is correct?

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Value-Added Compilation

- ◆ The more we rely on software, the more we demand more of it
- ◆ Compilers can help— **treat code as data**
 - ◆ Analyze the code
- ◆ **Correctness**
- ◆ **Security**
- ◆ **Reliability**
- ◆ **Program understanding**
- ◆ **Program evolution**
- ◆ **Software testing**
- ◆ **Reverse engineering**
- ◆ **Program obfuscation**
- ◆ **Code compaction**
- ◆ **Energy efficiency**

Computation important ⇒ understanding **computation important**

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Freedom Cuts Both Ways

- ◆ Just as compilers liberate the language designer, they also liberate the computer architect
- ◆ Can we change the ISA from one generation to the next?
 - ◆ Yes, if we trust our compilers
- ◆ Enables richer design space
 - ◆ VLIW
 - ◆ IA64
 - ◆ TRIPS
 - ◆ Multicore
 - ◆ Heterogeneous multi-core
 - ◆ Reconfigurable architectures

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Benefits to the Architect (cont)

- ◆ Two benefits of the compiler
 - ◆ Can simplify the hardware by shifting burden to the compiler
 - ◆ VLIW, IA64, TRIPS, software controlled caches, Cell
 - ◆ Can let the compiler inform the hardware
 - ◆ Bias bits
 - ◆ Prefetch instructions

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Virtualization is a Virtue

- ◆ High-level languages provide virtualization
 - ◆ Why is virtualization good?
- ◆ We can virtualize at many levels
 - ◆ Transmeta: dynamically compile x86 to VLIW
 - ◆ GPUs rely on dynamic compilation
 - ◆ JVMs and JITs

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The Point

- ◆ Compilers are a fundamental building block of modern systems
- ◆ We need to understand their power and limitations
 - ◆ Computer architects
 - ◆ Language designers
 - ◆ Software engineers
 - ◆ OS/Runtime system researchers
 - ◆ Security researchers
 - ◆ Formal methods researchers (model checking, automated theorem proving)

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Plan For Today

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Types of Optimizations

- ◆ Definition
 - ◆ An *optimization* is a transformation that is expected to improve the program in some way; often consists of *analysis* and *transformation*
e.g., decreasing the running time or decreasing memory requirements
- ◆ **Machine-independent optimizations**
 - ◆ Eliminate redundant computation
 - ◆ Move computation to less frequently executed place
 - ◆ Specialize some general purpose code
 - ◆ Remove useless code

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Types of Optimizations (cont)

- ◆ **Machine-dependent optimizations**
 - ◆ Replace a costly operation with a cheaper one
 - ◆ Replace a sequence of operations with a cheaper one
 - ◆ Hide latency
 - ◆ Improve locality
 - ◆ Reduce power consumption
- ◆ **Enabling transformations**
 - ◆ Expose opportunities for other optimizations
 - ◆ Help structure optimizations

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Sample Optimizations

- ◆ Arithmetic simplification
 - ◆ **Constant folding**
e.g., $x = 8/2;$

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Sample Optimizations (cont)

- ◆ Constant propagation

- ◆ *e.g.*, $\mathbf{x} = 3;$
 $\mathbf{y} = 4 + \mathbf{x};$

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Sample Optimizations (cont)

- ◆ Common subexpression elimination (CSE)

- ◆ *e.g.*, $\mathbf{x} = \mathbf{a} + \mathbf{b};$
 $\mathbf{y} = \mathbf{a} + \mathbf{b};$

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Sample Optimizations (cont)

- ◆ Dead (unused) assignment elimination

- ◆ e.g., `x = 3;`
... `x` not used...
`x = 4;`

This assignment is dead

- ◆ Dead (unreachable) code elimination

- ◆ e.g., `if (false == true) {`
 `printf("debugging...");`
}

This statement is dead

Sample Optimizations (cont)

- ◆ Loop-invariant code motion

- ◆ e.g., `for i = 1 to 10 do`
 `x = 3;`
 ...



- ◆ e.g., `x = 3;`
`for i = 1 to 10 do`
 ...

Sample Optimizations (cont)

- ◆ Induction variable elimination

- ◆ *e.g.*,

```
for i = 1 to 10 do
  a[i] = a[i] + 1;
```



```
for p = &a[1] to &a[10] do
  *p = *p + 1
```

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Sample Optimizations (cont)

- ◆ Loop unrolling

- ◆ *e.g.*,

```
for i = 1 to 10 do
  a[i] = a[i] + 1;
```



```
for i = 1 to 10 by 2 do
  a[i]    = a[i] + 1;
  a[i+1] = a[i+1] + 1;
```

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Is an Optimization Worthwhile?

- ◆ Criteria for evaluating optimizations
 - ◆ **Safety**: Does it preserve behavior?
 - ◆ **Profitability**: Does it actually improve the code?
 - ◆ **Opportunity**: Is it widely applicable?
 - ◆ **Cost (compilation time)**: Can it be practically performed?
 - ◆ **Cost (complexity)**: Can it be practically implemented?

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Scope of Analysis/Optimizations

- | | |
|---|---|
| ◆ Peephole | ◆ Local |
| ◆ Consider a small window of instructions | ◆ Consider blocks of straight line code (no control flow) |
| ◆ Usually machine-specific | ◆ Simple to analyze |

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Scope of Analysis/Optimizations (cont)

- ◆ **Global (intraprocedural)**
 - ◆ Consider entire procedures
 - ◆ Must consider branches, loops, merging of control flow
 - ◆ Use data-flow analysis
 - ◆ Make simplifying assumptions at procedure calls
- ◆ **Whole program (interprocedural)**
 - ◆ Consider multiple procedures
 - ◆ Analysis even more complex (calls, returns)
 - ◆ Hard with separate compilation

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Time of Optimization

- ◆ Compile time
- ◆ Link time
- ◆ Configuration time
- ◆ Runtime

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Optimization Dimensions: A Rich Space

- ◆ Abstraction level
 - ◆ Machine-dependent, machine-independent
- ◆ Goal
 - ◆ Performance, correctness, etc
 - ◆ Enabling transformation
- ◆ Scope
 - ◆ Peephole, local, global, interprocedural
- ◆ Timing
 - ◆ Compile time, link time, configuration time, run time

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Limits of Compiler Optimizations

- ◆ Fully Optimizing Compiler (FOC)
 - ◆ $FOC(P) = P_{opt}$
 - ◆ P_{opt} is the *smallest* program with same I/O behavior as P
- ◆ Observe
 - ◆ If program Q produces no output and never halts, $FOC(Q) = L: \text{goto } L$
- ◆ Aha! We've solved the halting problem?!
- ◆ Moral
 - ◆ Cannot build FOC
 - ◆ Can always build a better optimizing compiler

January 21, 2015 (full employment theorem for compiler writers!)

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Optimizations Don't Always Help

◆ Common Sub-expression Elimination

$x = a + b$	→	$t = a + b$
$y = a + b$		$x = t$
		$y = t$
$\underbrace{\hspace{10em}}$		$\underbrace{\hspace{10em}}$
2 adds		1 add
4 variables		5 variables

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Optimizations Don't Always Help (cont)

◆ Fusion and Contraction

<pre>for i = 1 to n T[i] = A[i] + B[i] for i = 1 to n C[i] = D[i] + T[i]</pre>	→	<pre>for i = 1 to n t = A[i] + B[i] C[i] = D[i] + t</pre>
--	---	---

t fits in a register, so no loads or stores in this loop.

Huge win on most machines.

Degrades performance on machines with hardware managed stream buffers.

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Optimizations Don't Always Help

- ◆ **Backpatching**

`o.foo();` } In Java, the address of `foo()` is often not known until runtime (due to dynamic class loading), so the method call requires a **table lookup**.

After the first execution of this statement, **backpatching** replaces the table lookup with a direct call to the proper function.

Q: How could this optimization ever hurt?

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Phase Ordering Problem

- ◆ In what order should optimizations be performed?

- ◆ **Simple dependences**

- ◆ One optimization creates opportunity for another
e.g., copy propagation and dead code elimination

- ◆ **Cyclic dependences**

- ◆ *e.g.*, constant folding and constant propagation

- ◆ **Adverse interactions**

- ◆ *e.g.*, common sub-expression elimination and register allocation

- ◆ *e.g.*, register allocation and instruction scheduling

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Engineering Issues

- ◆ Building a compiler is an engineering activity
- ◆ Balance multiple goals
 - ◆ Benefit for *typical* programs
 - ◆ Complexity of implementation
 - ◆ Compilation speed
- ◆ Overall Goal
 - ◆ Identify a small set of general analyses and optimization
 - ◆ Easier said than done: just one more...

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Two Approaches— Which is Better?

- ◆ Build a compiler from scratch

- ◆ Extend an existing compiler

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Administrative Matters

- ◆ Turn to your syllabus

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Next Time

- ◆ Reading
 - ◆ Syllabus
- ◆ Lecture
 - ◆ Undergraduate compilers in a day!

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