

# A Little Blocked Literal Goes a Long Way

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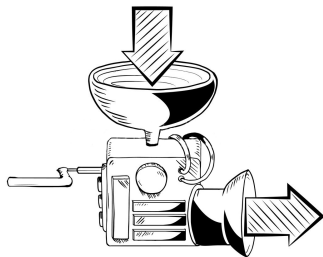
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- Brief overview of QBF and corresponding proof systems.
- **Main result: QRAT** simulates **long-distance resolution**.
  - QRAT is the QBF generalization of DRAT.
  - Simulation is **polynomial**.
- We have an **implementation** and **evaluation** of the simulation.

# Satisfiability of Quantified Boolean Formulas (QSAT)

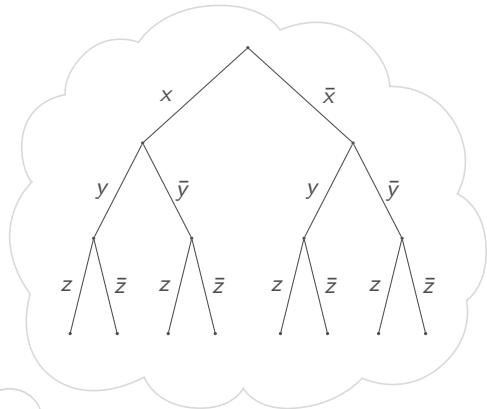
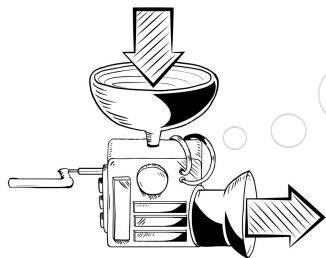
“For every truth value of  $x$ ,  
does there exist a truth value of  $y$ ,  
such that ...”

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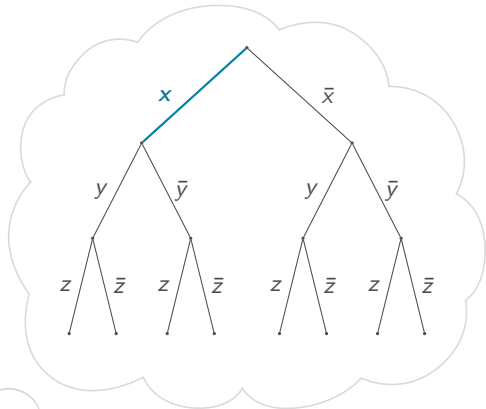
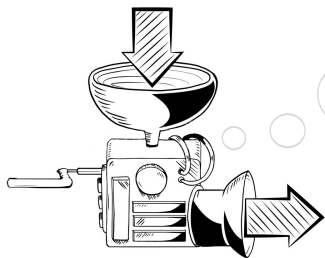
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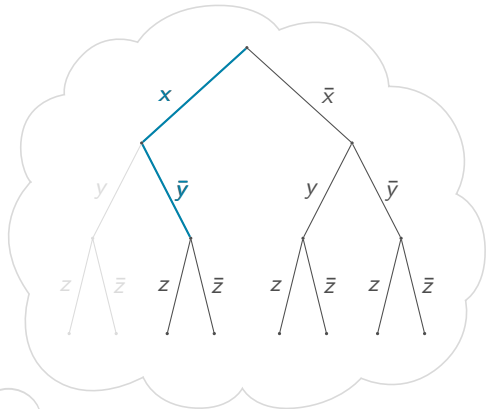
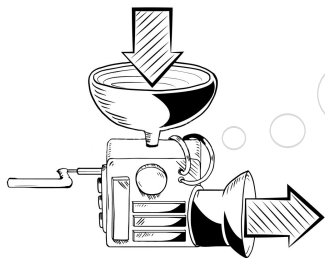
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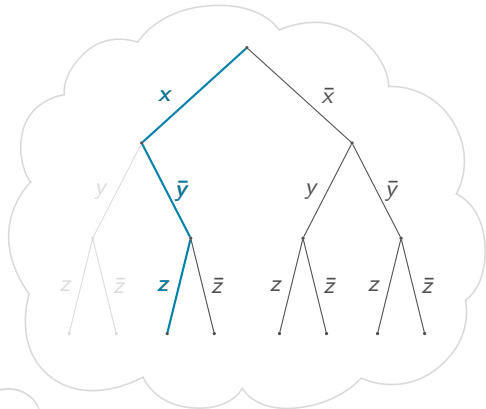
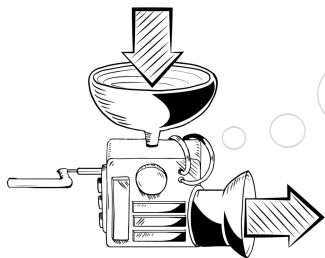
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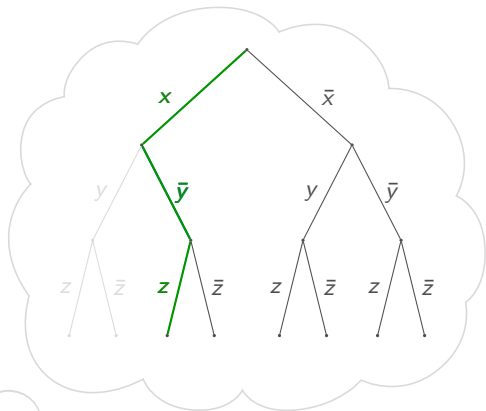
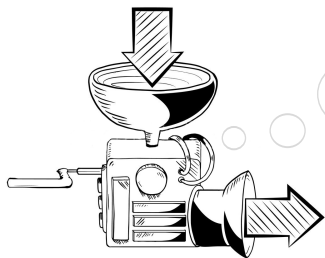
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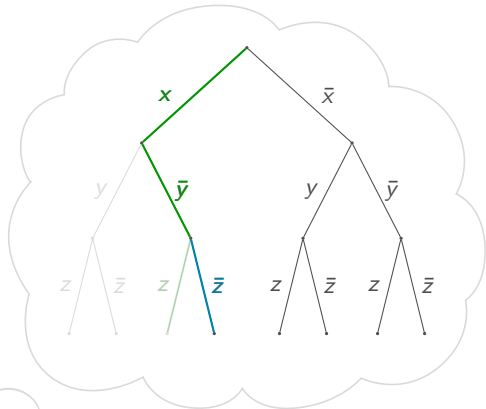
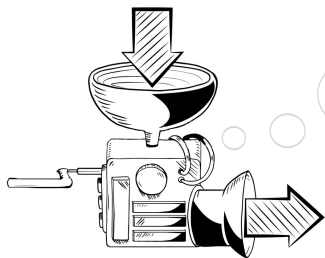
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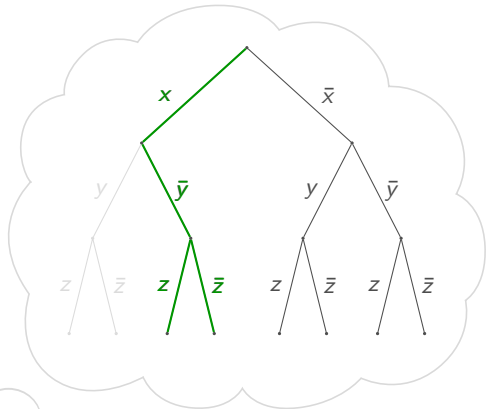
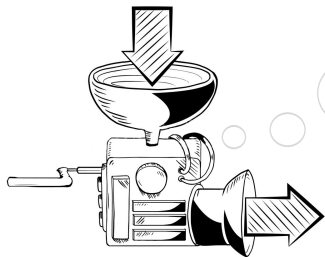
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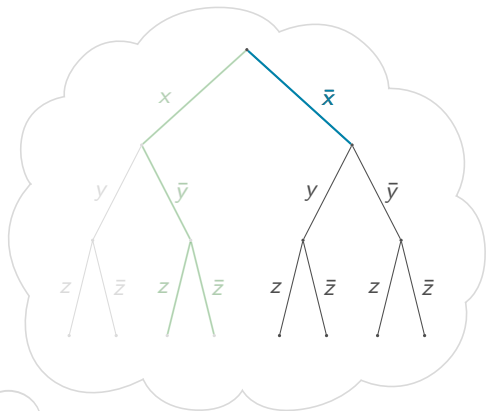
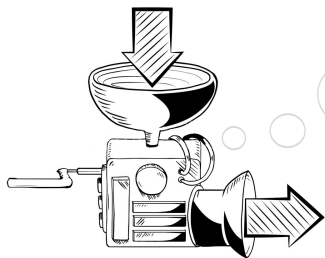
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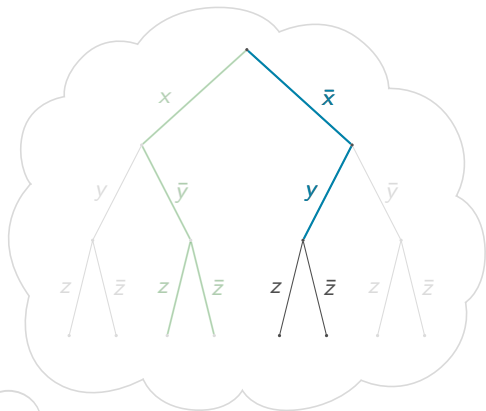
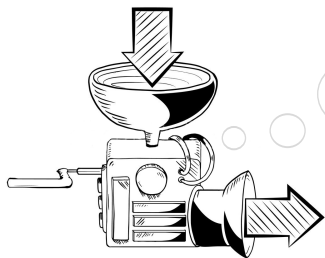
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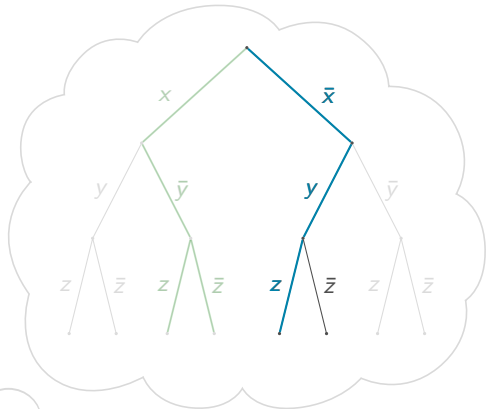
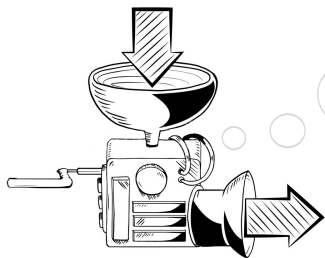
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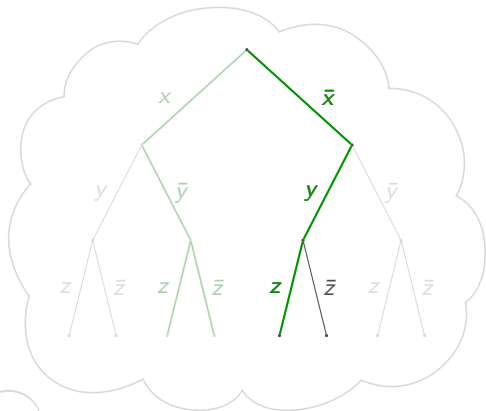
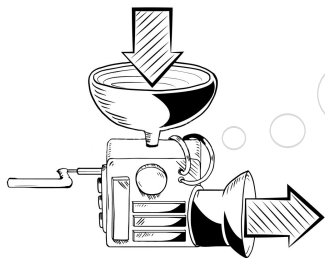
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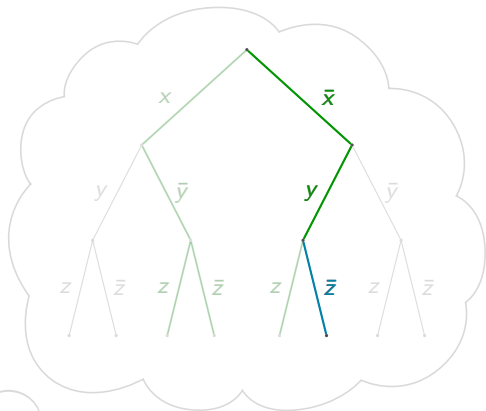
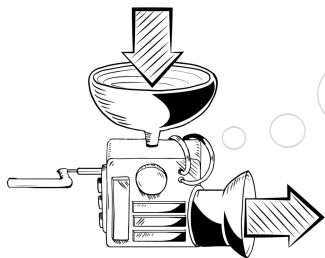
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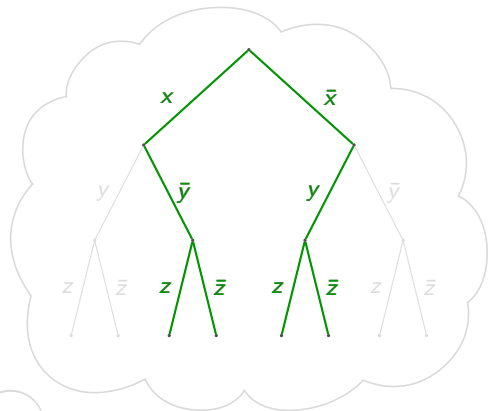
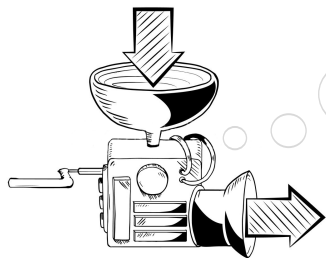
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Satisfiable

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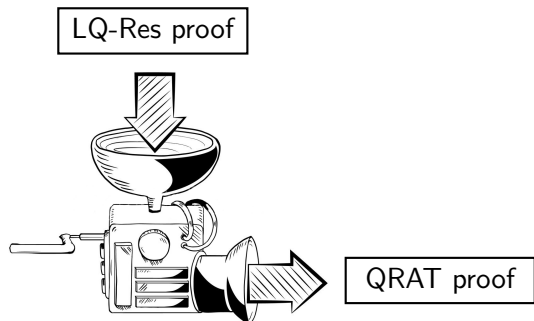
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- **Open question**: If there is a **short LQ-Res proof** of a QBF, is there also a **short QRAT proof**?
  - **Short** = **polynomial** with respect to the size of the formula.
  - Our answer: **Yes!**

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# Simulating LQ-Res With QRAT

- How to show that there is a short QRAT proof for every short LQ-Res proof?
- ➔ Answer: With a simulation procedure.
  - Takes as input an LQ-Res proof and transforms it into a short QRAT proof.



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- Each clause in an LQ-Res proof is either contained in the formula or derived via one of the following two rules:

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# Proving Unsatisfiability of QBFs: Long-Distance Resolution

- Example proof with long-distance resolution:

$$\phi = \exists e_1 \forall u_1 \exists e_2 \exists e_3. (\bar{e}_1 \vee \bar{u}_1 \vee e_3) \wedge (\bar{u}_1 \vee e_2 \vee \bar{e}_3) \wedge (e_1 \vee u_1 \vee e_2) \wedge (\bar{e}_2)$$

$$\begin{array}{c}
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  - Blocked-literal elimination
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  - $\forall$ -reduction of non-complementary literals
  - Blocked-literal elimination (QRAT-literal elimination)
  - Blocked-literal addition (QRAT-literal addition)

## Example: QRAT proof

1.  $a_n \vee \bar{x}_n \vee \bar{c}_1 \vee \cdots \vee \bar{c}_{n-1}$  (Q-res)
2.  $b_n \vee x_n \vee \bar{c}_1 \vee \cdots \vee \bar{c}_{n-1}$  (Q-res)
3.  $a_{n-1} \vee \bar{x}_{n-1} \vee \bar{b}_n \vee \bar{x}_n \vee \bar{c}_1 \vee \cdots \vee \bar{c}_{n-1}$  (Q-res)
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- ⋮



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- The blocked-literal definition is based on **outer resolvents**:
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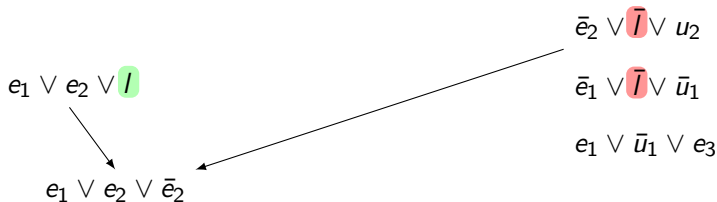
$$\bar{e}_2 \vee \bar{I} \vee u_2$$

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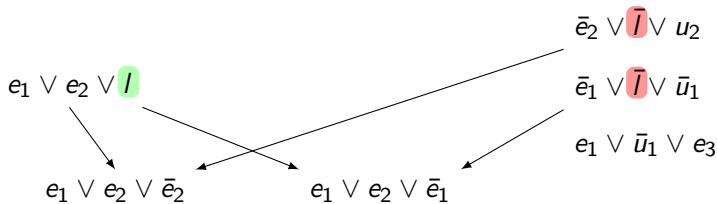
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- By **successively removing complementary literals** from resolution steps, we obtain a valid QRAT proof.

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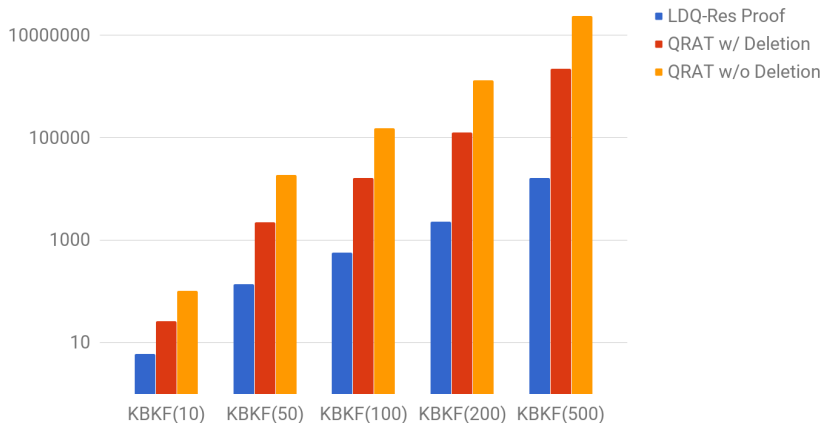
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# Kleine Bünig Formulas (KBKF): LDQ-Res to QRAT

File size of generated proofs: LDQ-Res (Egly et al. 2013) to QRAT with and without deletion.



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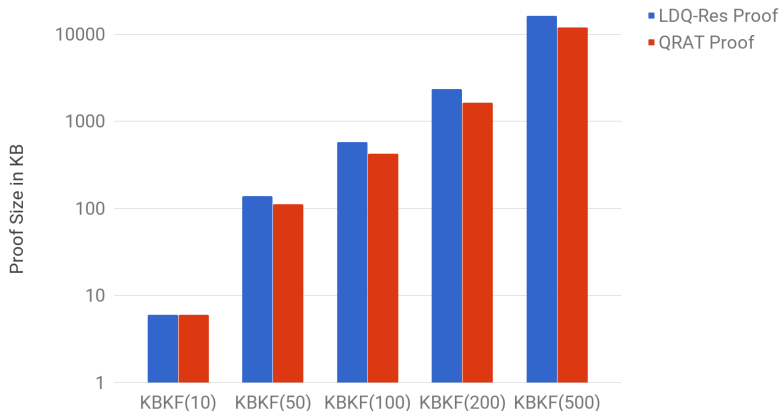
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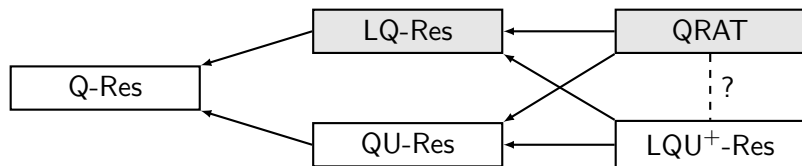
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# Kleine Bünig Formulas (KBKF): QRAT vs. LDQ-Res

File size of hand-crafted proofs: LDQ-Res (Egly et al. 2013) vs. QRAT.



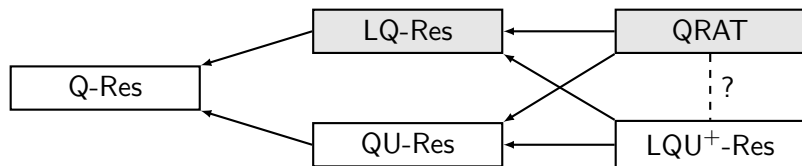
# Complexity Landscape: QRAT and Resolution Systems



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# Complexity Landscape: QRAT and Resolution Systems



- Open question: Can QRAT simulate  $LQU^+$ -Res?
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