

#### Goals:

Given system of inequalities of the form  $Ax \leq b$ 

- determine if system has an integer solution
- enumerate all integer solutions

### Running example:

$$3x + 4y \ge 16$$

$$4x + 7y \le 56$$

$$4x - 7y \le 20$$

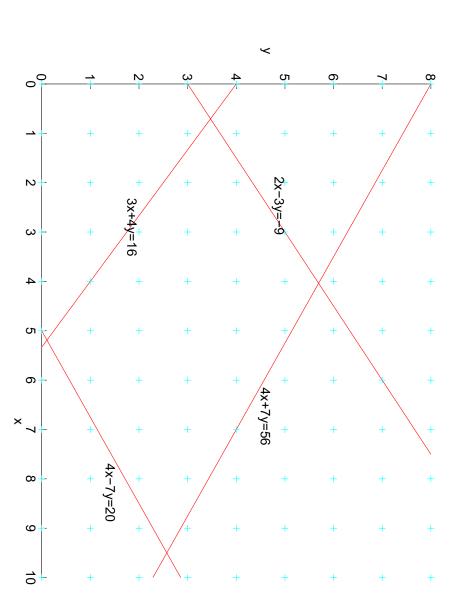
$$2x - 3y \ge -9$$

Lower bounds for x: (1) and (4)

Upper bounds for x: (2) and (3)

Upper bounds for y: (2) and (4) Lower bounds for y: (1) and (3)

### MATLAB graphs:



Code for enumerating integer points in polyhedron: (see graph)

Outer loop: Y, Inner loop: X

D0 Y=
$$\lceil 4/37 \rceil$$
,  $\lfloor 74/13 \rfloor$   
D0 X= $\lceil max(16/3 - 4y/3, -9/2 + 3y/2) \rceil$ ,  $\lfloor min(5 + 7y/4, 14 - 7y/4) \rfloor$   
.....

Outer loop: X, Inner loop: Y

How do we can determine loop bounds?

# inequalities Fourier-Motzkin elimination: variable elimination technique for

$$3x + 4y \ge 16 \tag{5}$$

$$4x + 7y \le 56$$

$$4x - 7y \le 20$$

$$2x - 3y \ge -9$$

(8)

7

(6)

Let us project out x.

First, express all inequalities as upper or lower bounds on x.

$$x \geq 16/3 - 4y/3$$

$$x \leq 14 - 7y/4$$

(9)(10)(11)(12)

$$x \leq 5 + 7y/4$$

$$y \ge -9/2 + 3y/2$$

on x. lower bound on x must be less than or equal to every upper bound For any y, if there is an x that satisfies all inequalities, then every

bounds Generate a new system of inequalities from each pair (upper,lower)

$$5+7y/4 \geq 16/3-4y/3$$
 (Inequalities3, 1)  
 $5+7y/4 \geq -9/2+3y/2$  (Inequalities3, 4)  
 $14-7y/4 \geq 16/3-4y/3$  (Inequalities2, 1)  
 $14-7y/4 \geq -9/2+3y/2$  (Inequalities2, 4)

#### Simplify:

$$y \geq 4/37$$

$$\leq 104/5$$

$$\leq$$
 74/13

$$max(4/37, -38)$$

$$y \le min(104/5, 74/13)$$

$$4/37 \le y \le 74/13$$

This means there are rational solutions to original system of inequalities.

We can now express solutions in closed form as follows:

$$4/37 \leq y \leq 4/37$$

 $max(16/3 - 4y/3, -9/2 + 3y/2) \le x \le min(5 + 7y/4, 14 - 7y/4)$ 

## Iterative step: Fourier-Motzkin elimination: iterative algorithm

- obtain reduced system by projecting out a variable
- if reduced system has a rational solution, so does the original

## Termination: no variables left

Projection along variable x: Divide inequalities into three categories

$$a_1 * y + a_2 * z + \dots \le c_1(no \ x)$$
 $b_1 * x \le c_2 + b_2 * y + b_3 * z + \dots (upper \ bound)$ 
 $d_1 * x \ge c_3 + d_2 * y + d_3 * z + \dots (lower \ bound)$ 

New system of inequalities:

- All inequalities that do not involve x
- Each pair (lower, upper) bounds gives rise to one inequality:

$$b_1[c_3 + d_2 * y + d_3 * z + \ldots] \le d_1[c_2 + b_2 * y + b_3 * z + \ldots]$$

number between  $(x_1, y_1, z_1...)$  satisfies the original system, where  $x_1$  is a rational Theorem: If  $(y_1, z_1, ...)$  satisfies the reduced system, then

and  $min(1/b_1(c_2+b_2y_1+b_3z_1+...),....)$  (over all upper bounds)

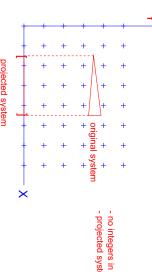
 $max(1/d_1(c_3+d_2y_1+d_3z_1+...),....)$  (over all lower bounds)

Proof: trivial

What can we conclude about integer solutions?

the original system. Corollary: If reduced system has no integer solutions, neither does

does too Not true: Reduced system has integer solutions => original system



no integers in original polyhedron projected system contains integers

not guaranteed to preserve "integrality" (cf. equalities) Key problem: Multiplying one inequality by  $b_1$  and other by  $d_1$  is

reduced system implies integer solution to original system. bound coefficients  $d_i$  happen to be 1, then integer solution to Exact projection: If all upper bound coefficients  $b_i$  or all lower

original system if there exists an integer  $x_1$  between reduced system in FM elimination, then  $(x_1, y_1, z_1...)$  satisfies the Theorem: If  $(y_1, z_1, ...)$  is an integer vector that satisfies the

and  $\lceil max(1/d_1(c_3 + d_2y_1 + d_3z_1 + ...), ....) \rceil$  (over all lower bounds)

 $[min(1/b_1(c_2+b_2y_1+b_3z_1+...),....)]$  (over all upper bounds).

Proof: trivial

that satisfy system as follows: elimination to generate a loop nest to enumerate all integer points Enumeration: Given a system  $Ax \leq b$ , we can use Fourier-Motzkin

- pick an order to eliminate variables (this will be the order of variables from innermost loop to outermost loop)
- eliminate variables in that order to generate upper and lower bounds for loops as shown in theorem in previous slide

of that loop of some loop in the loop nest will be bigger than the upper bound Remark: if polyhedron has no integer points, then the lower bound

elimination to project down to a single variable. Existence: Given a system  $Ax \leq b$ , we can use Fourier-Motzkin

- If the reduced system has no integer solutions, then original system has no integer solutions either.
- If the reduced system has integer solutions and all projections were exact, then original system has integer solutions too.
- If reduced system has integer solutions and some projections has integer solutions. were no exact, be conservative and assume that original system