CS313H Logic, Sets, and Functions: Honors Fall 2012

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Good Morning, Colleagues



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Are there any questions?

Logistics

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- Module C1 is longer than most

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- Manage your sleep and time

• "I don't like the way we start so abruptly with "Good morning... Any questions?" Often I don't have a list of specific questions in mind but if we spent a few minutes briefly overviewing the new material we were supposed to have picked up in the module, the context would bring up questions."

Some questions

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 - Removing vertex of **min** degree (not max as on webpage)

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- Modern day Königsberg any better? (Nick Walther)

Eulerian and Hamiltonian Graphs

1. Eulerian Path?

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- 1. Eulerian Path?
- 2. Eulerian Circuit?
- 3. Hamiltonian Path?
- 4. Hamiltonian Circuit?

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- Induction on the number of vertices
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- Consider the longest simple path $W=v_0,...,v_r$ (there are a finite number of these paths, so the longest one certainly exists; if it's not unique pick any old one). Prove that W is an Eulerian circuit.

Assignments for Thursday

• Module C1