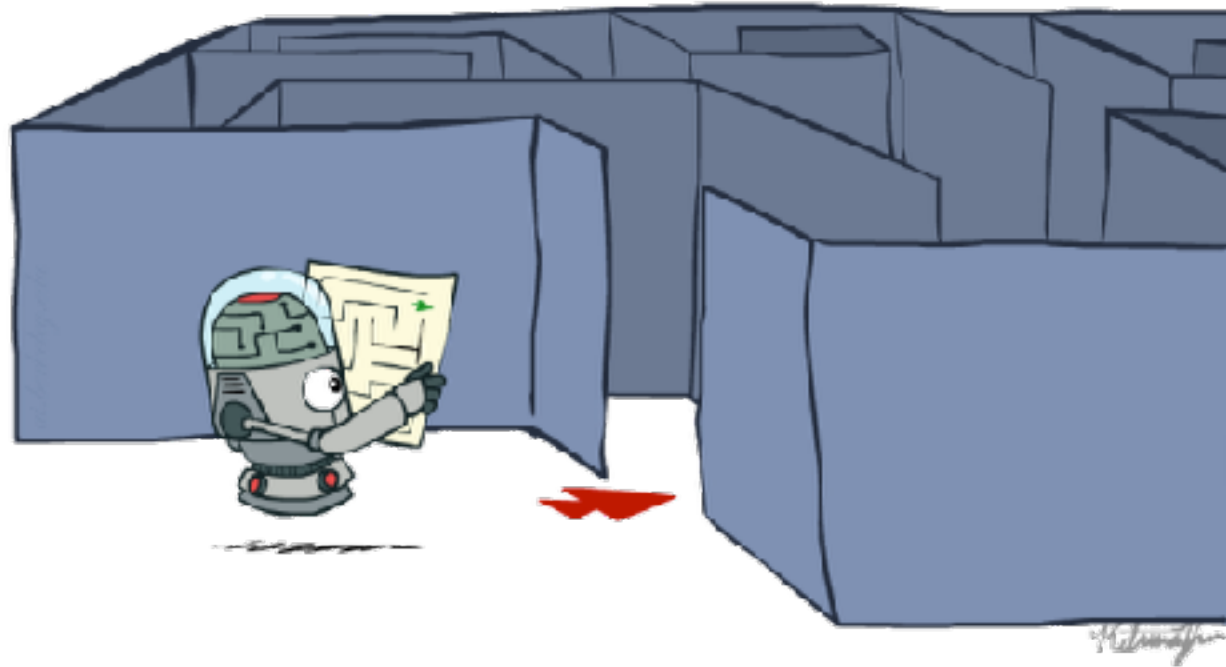


# CS 343H: Honors Artificial Intelligence

## Search

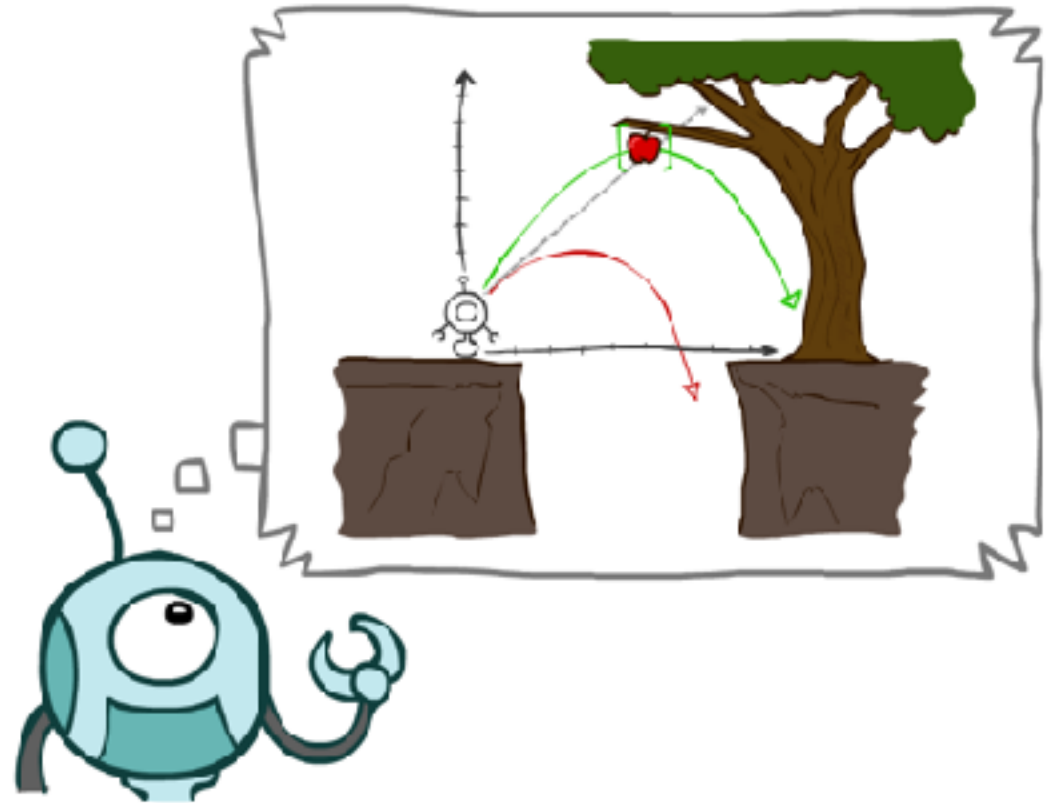


Prof. Peter Stone

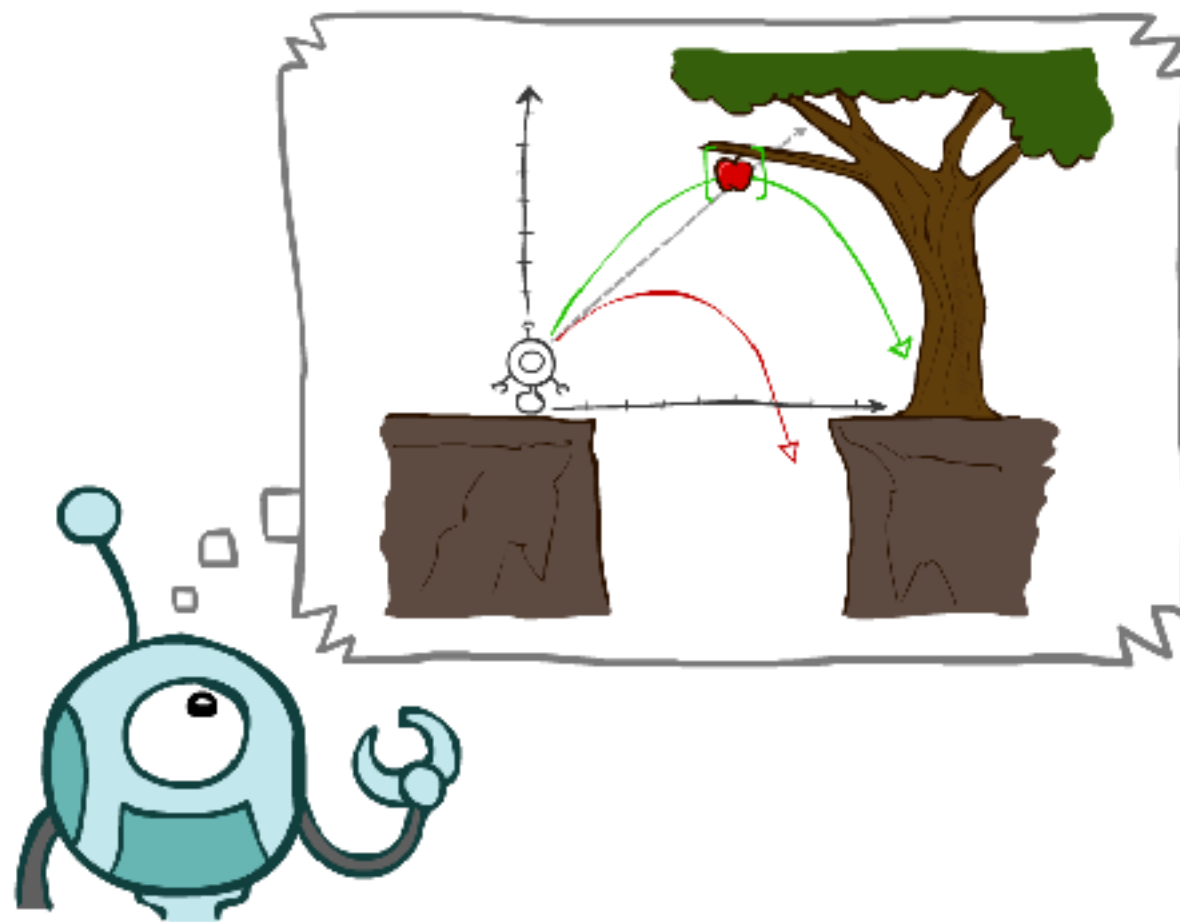
University of Texas at Austin

# Today

- Agents that Plan Ahead
- Search Problems
- Uninformed Search Methods
  - Depth-First Search
  - Breadth-First Search
  - Uniform-Cost Search

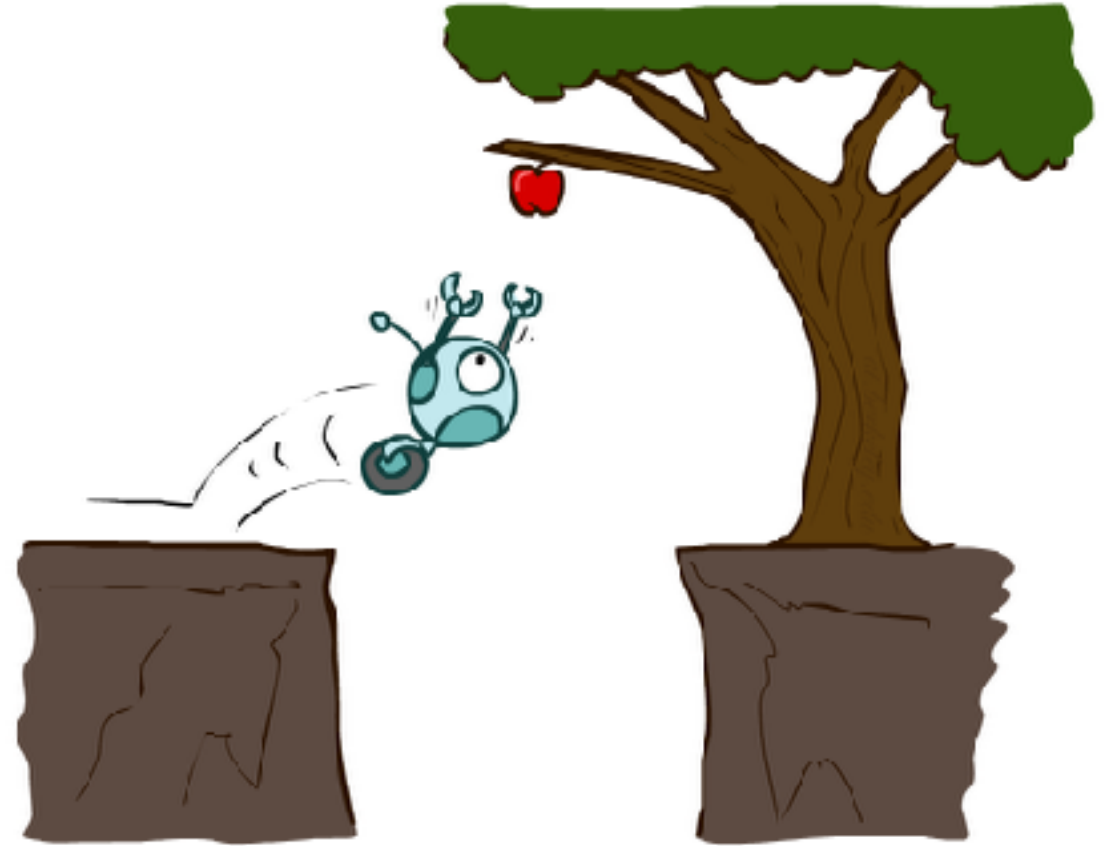


# Agents that Plan

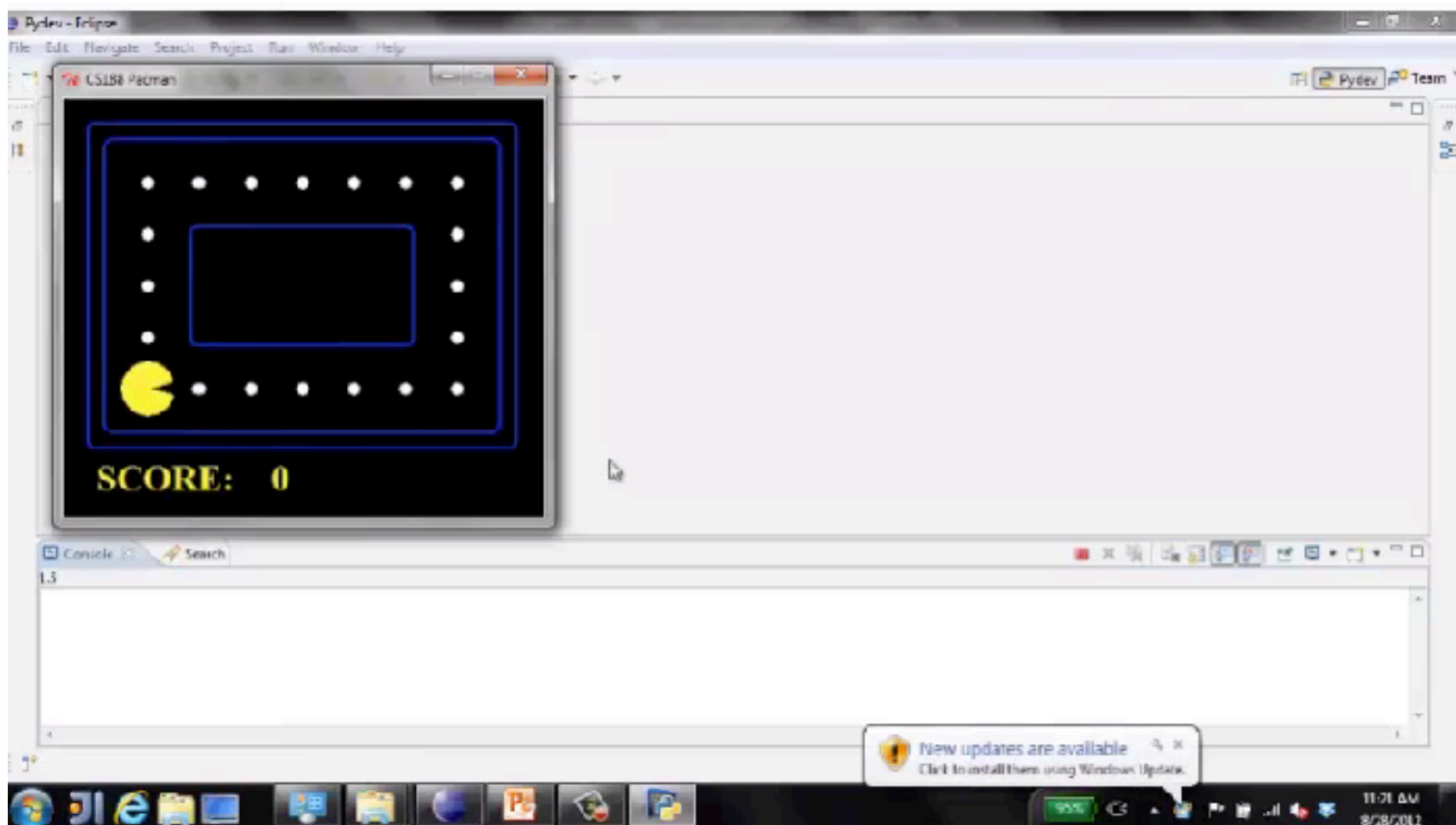


# Reflex Agents

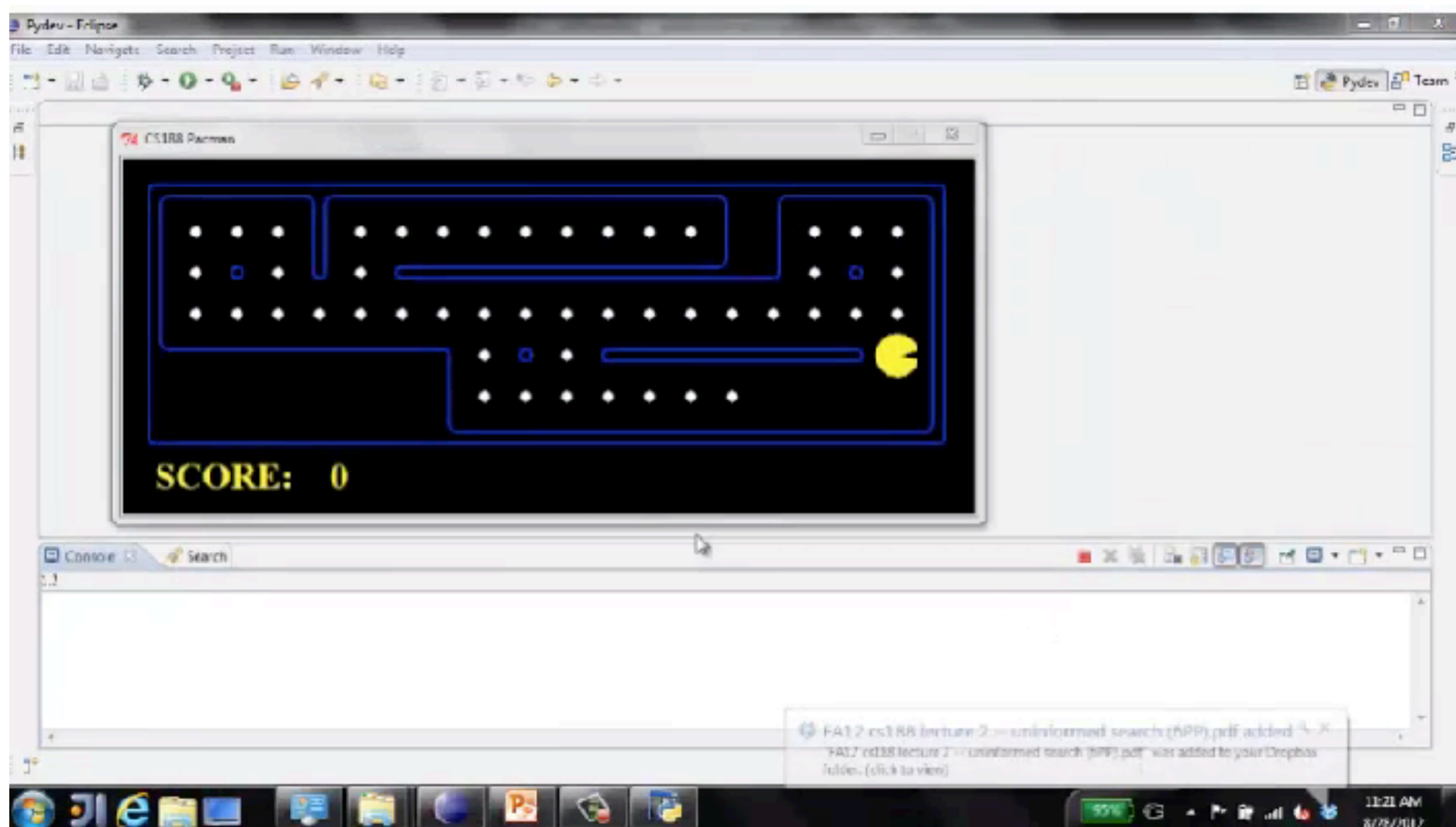
- Reflex agents:
  - Choose action based on current percept (and maybe memory)
  - May have memory or a model of the world's current state
  - Do not consider the future consequences of their actions
  - Consider how the world IS
- Can a reflex agent be rational?



# Video of Demo Reflex — Success

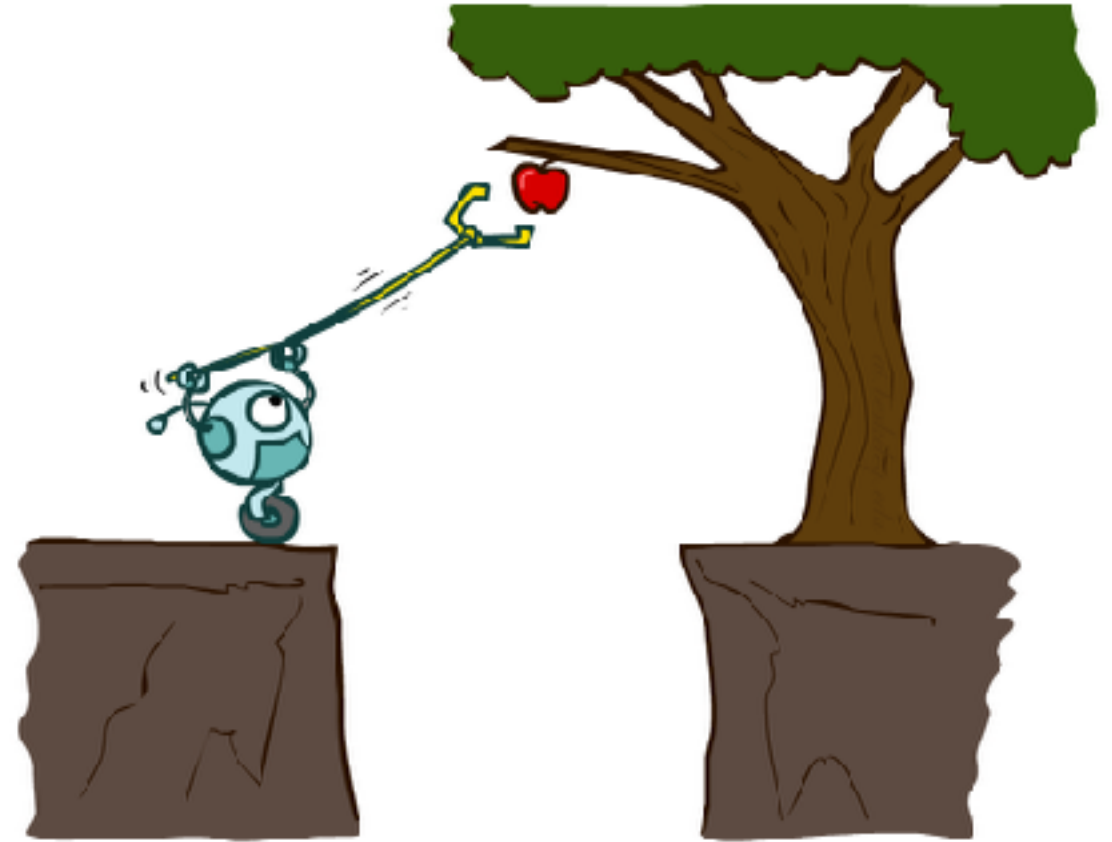


# Video of Demo Reflex — Stuck

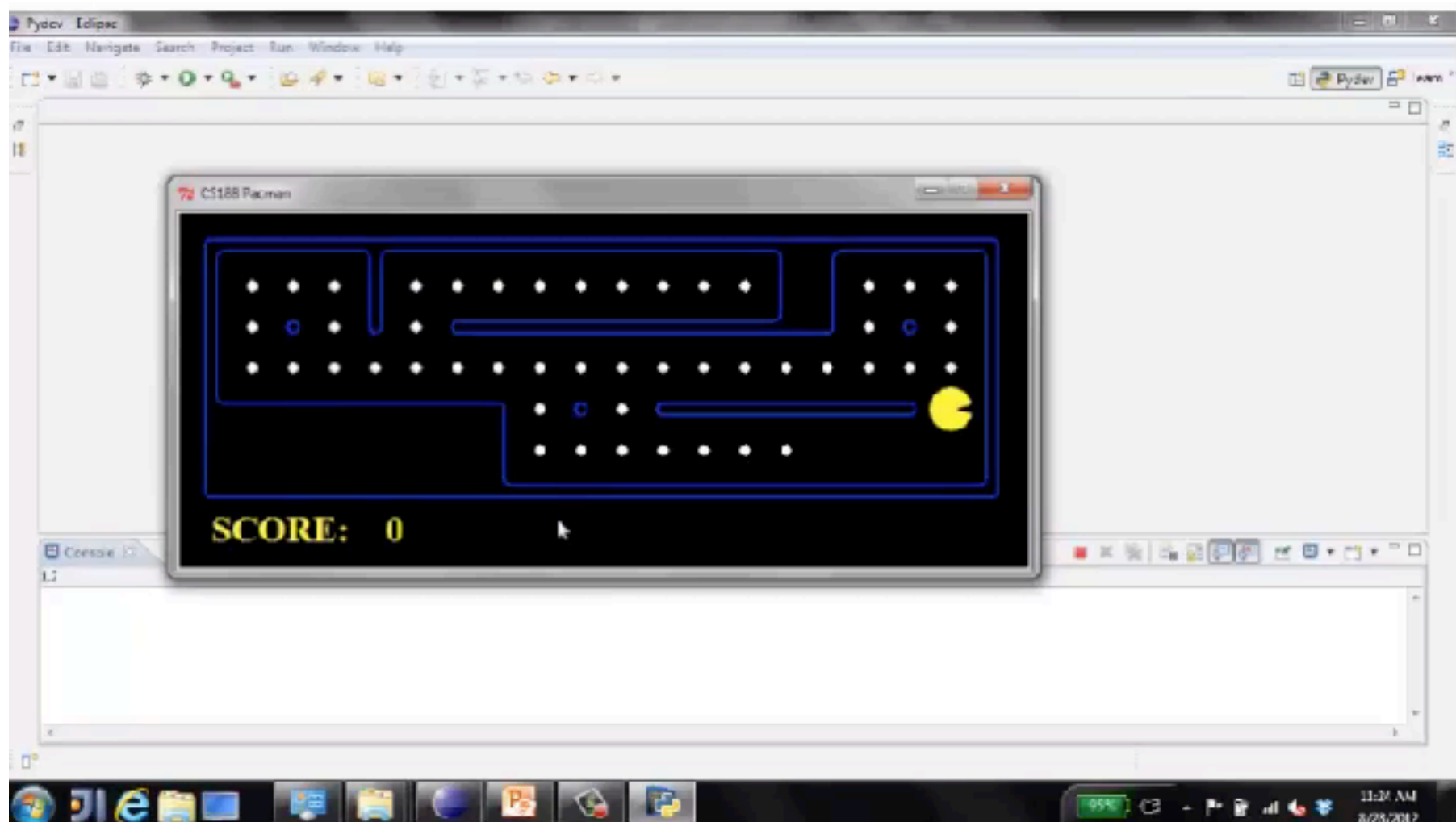


# Planning Agents

- Planning agents:
  - Ask “what if”
  - Decisions based on (hypothesized) consequences of actions
  - Must have a model of how the world evolves in response to actions
  - Must formulate a goal (test)
  - Consider how the world **WOULD BE**
- Optimal vs. complete planning
- Planning vs. replanning

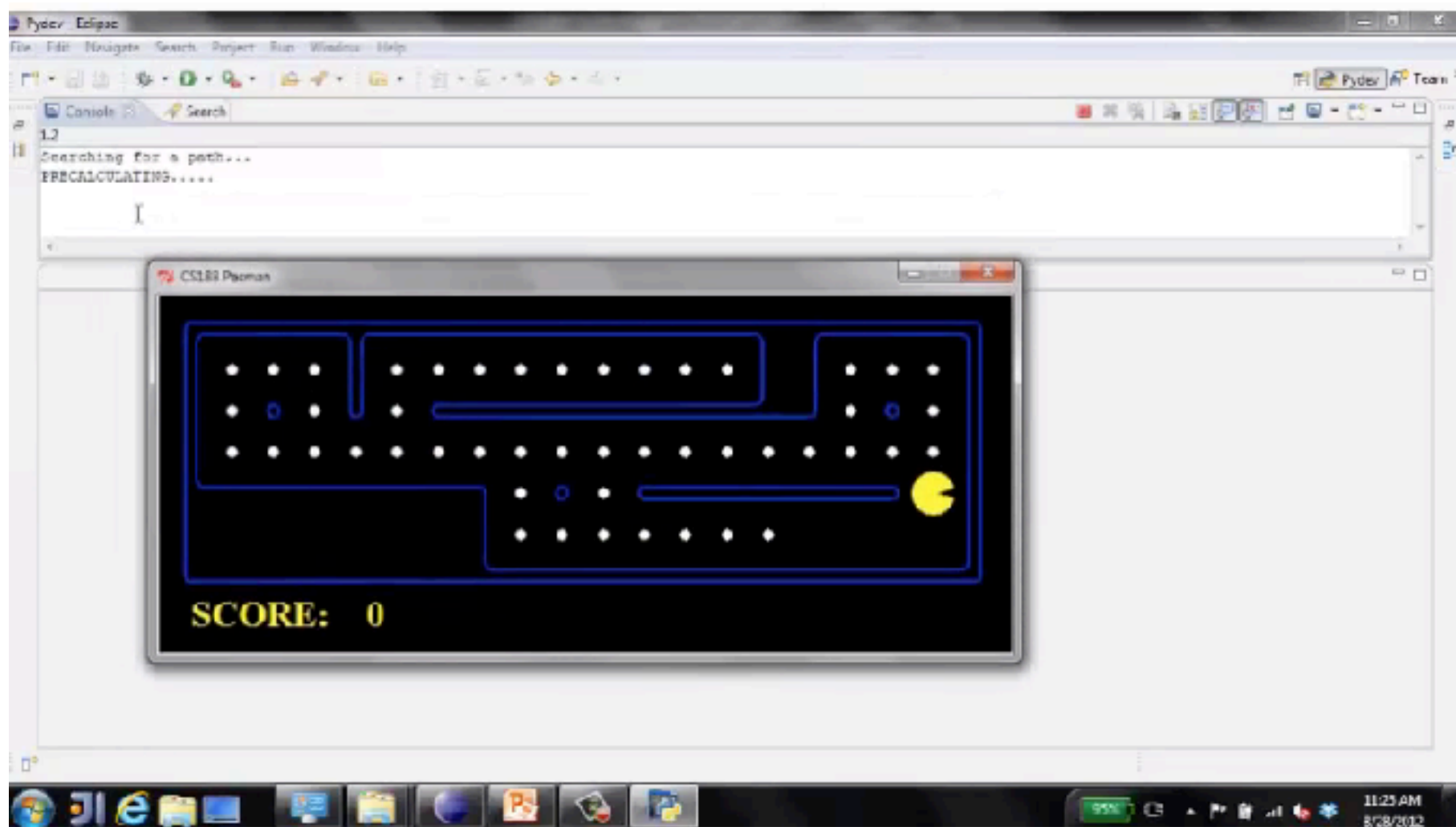


# Video of Demo — Suboptimal local replanning





# Video of Demo — Globally optimal planning



# Search Problems

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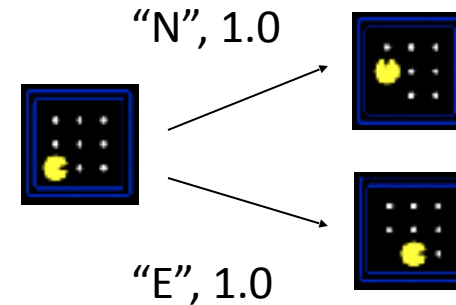
# Search Problems

- A **search problem** consists of:

- A state space



- A successor function  
(with actions, costs)



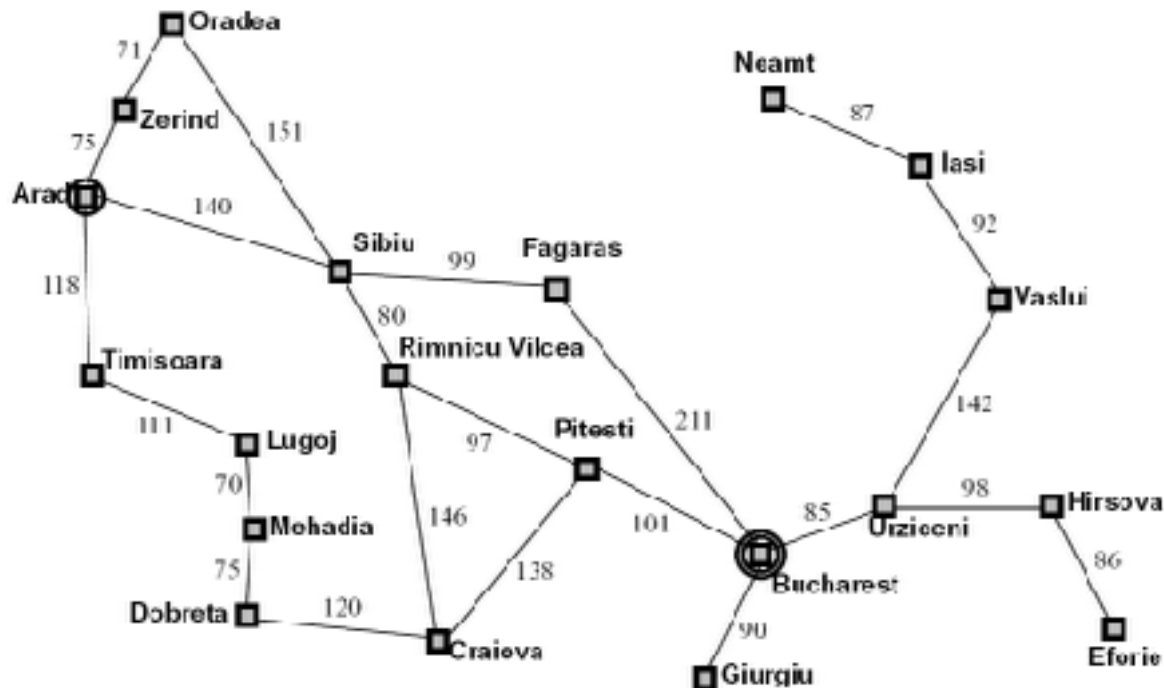
- A start state and a goal test
- A **solution** is a sequence of actions (a plan) which transforms the start state to a goal state

# Search Problems Are Models

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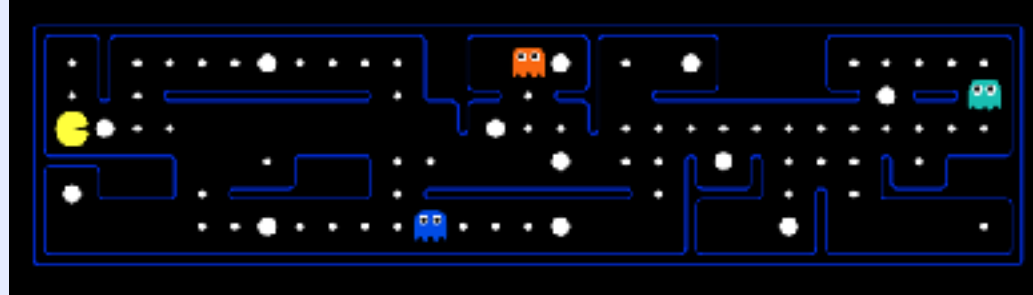
# Example: Traveling in Romania



- State space:
  - Cities
- Successor function:
  - Roads: Go to adjacent city with cost = distance
- Start state:
  - Arad
- Goal test:
  - Is state == Bucharest?
- Solution?

# What's in a State Space?

The **world state** includes every last detail of the environment



A **search state** keeps only the details needed for planning (abstraction)

## ■ Problem: Pathing

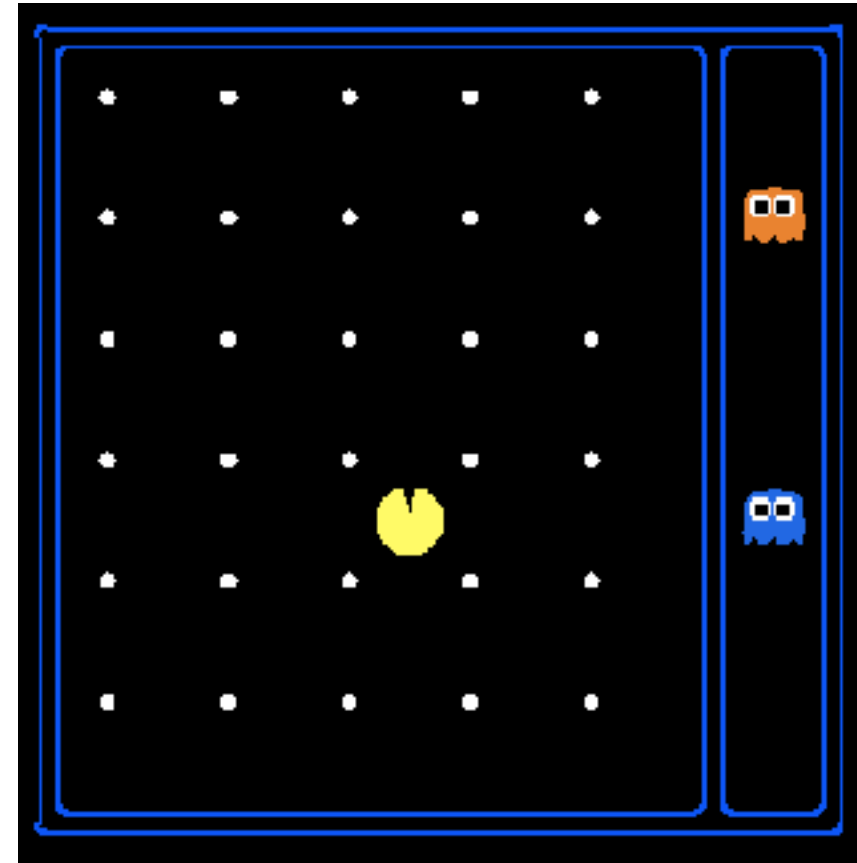
- States:  $(x,y)$  location
- Actions: NSEW
- Successor: update location only
- Goal test: is  $(x,y)=\text{END}$

## ■ Problem: Eat-All-Dots

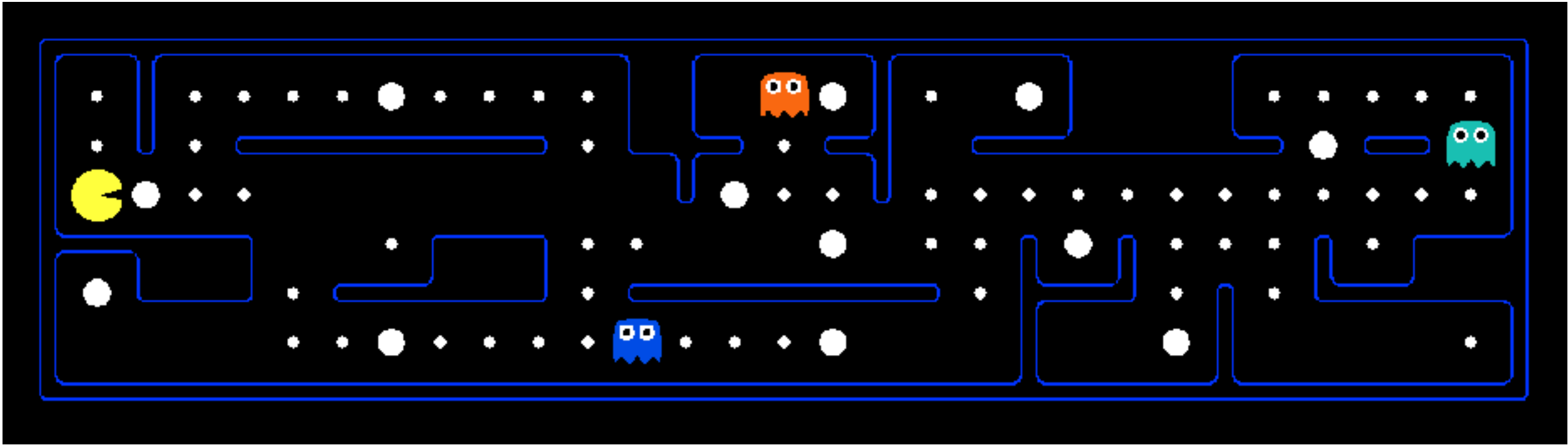
- States:  $\{(x,y), \text{dot booleans}\}$
- Actions: NSEW
- Successor: update location and possibly a dot boolean
- Goal test: dots all false

# State Space Sizes?

- World state:
  - Agent positions: 120
  - Food count: 30
  - Ghost positions: 12
  - Agent facing: NSEW
- How many
  - World states?  
 $120 \times (2^{30}) \times (12^2) \times 4$  (> 74 trillion!)
  - States for pathing?  
120
  - States for eat-all-dots?  
 $120 \times (2^{30})$  (> 128 billion)



# Quiz: Safe Passage

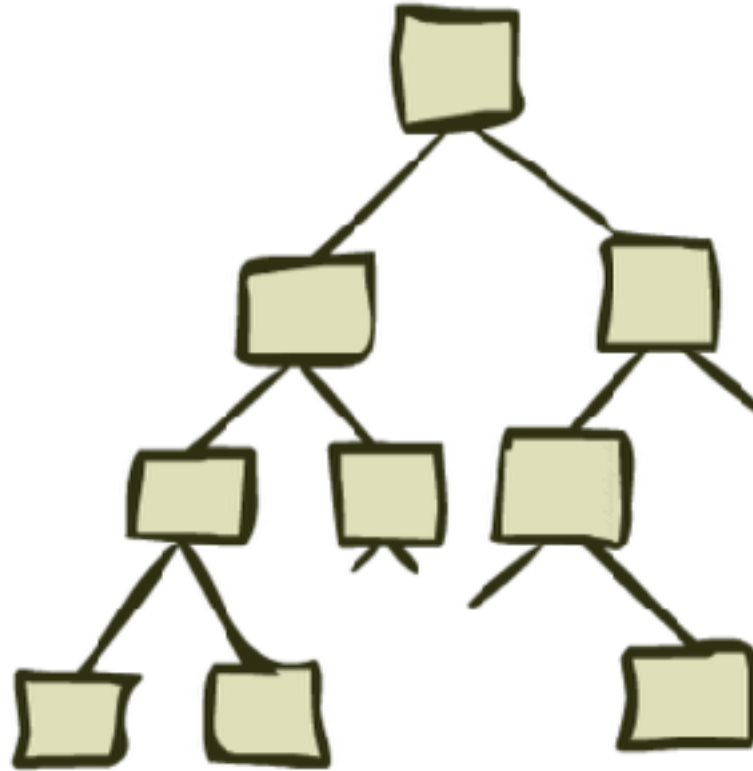


- Problem: eat all dots while keeping the ghosts perma-scared
- What does the state space have to specify?
  - (agent position, dot booleans, power pellet booleans, remaining scared time)



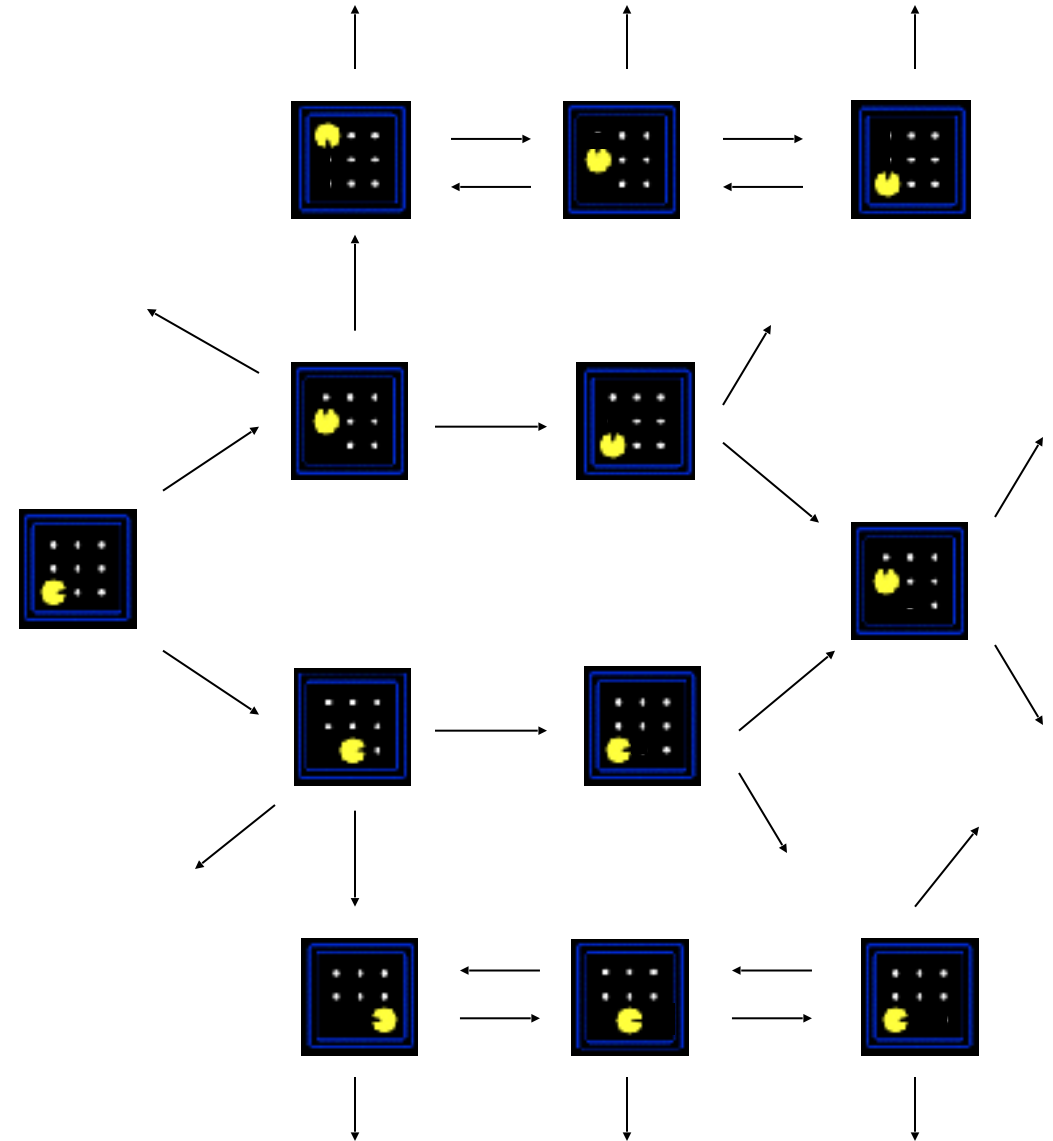
# State Space Graphs and Search Trees

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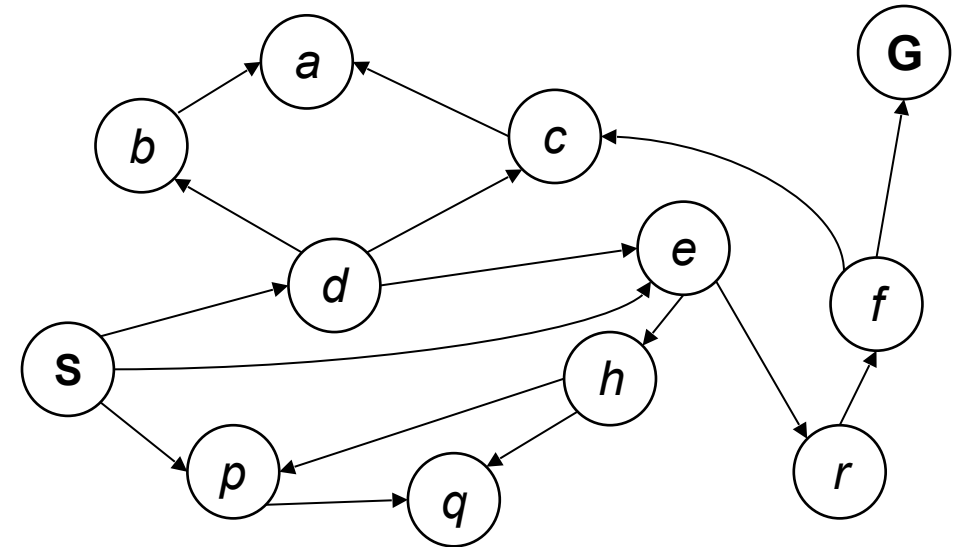
# State Space Graphs

- State space graph: A mathematical representation of a search problem
  - Nodes are (abstracted) world configurations
  - Arcs represent successors (action results)
  - The goal test is a set of goal nodes (maybe only one)
- In a state space graph, each state occurs only once!
- We can rarely build this full graph in memory (it's too big), but it's a useful idea



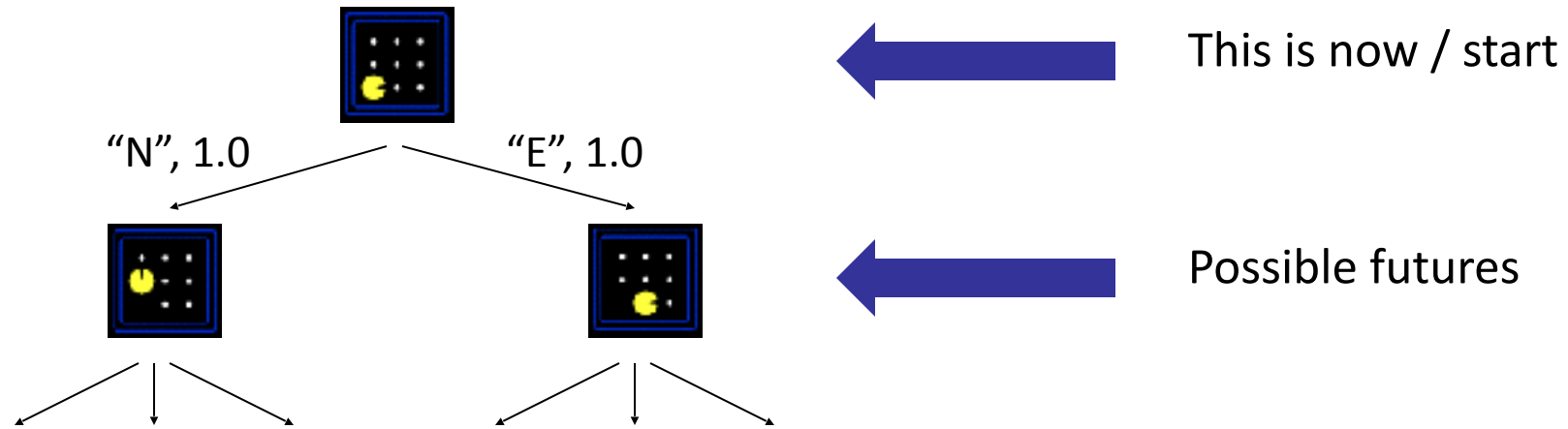
# State Space Graphs

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*Tiny search graph for a tiny search problem*

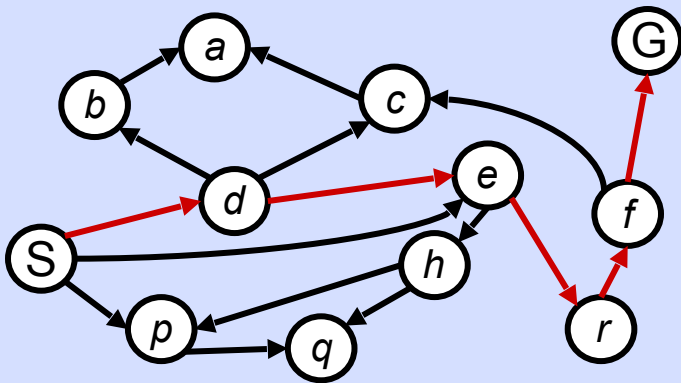
# Search Trees



- A search tree:
  - A “what if” tree of plans and their outcomes
  - The start state is the root node
  - Children correspond to successors
  - Nodes show states, but correspond to PLANS that achieve those states
  - For most problems, we can never actually build the whole tree

# State Space Graphs vs. Search Trees

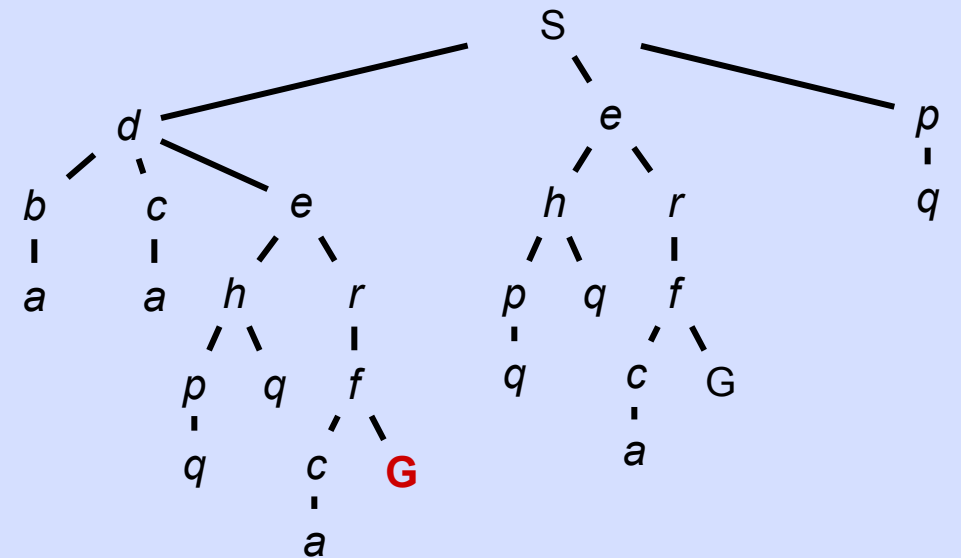
## State Space Graph



*Each NODE in the search tree is an entire PATH in the state space graph.*

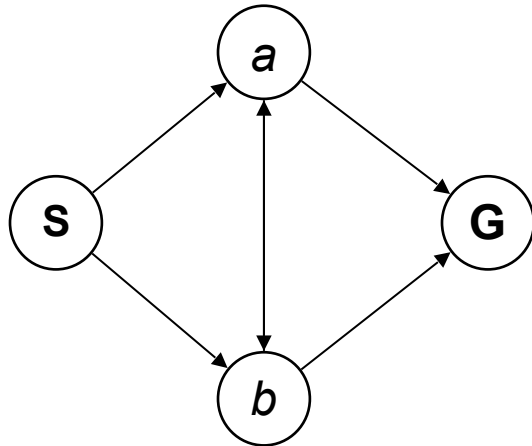
*We construct both on demand – and we construct as little as possible.*

## Search Tree



# Quiz: State Space Graphs vs. Search Trees

Consider this 4-state graph:



How big is its search tree (from S)?



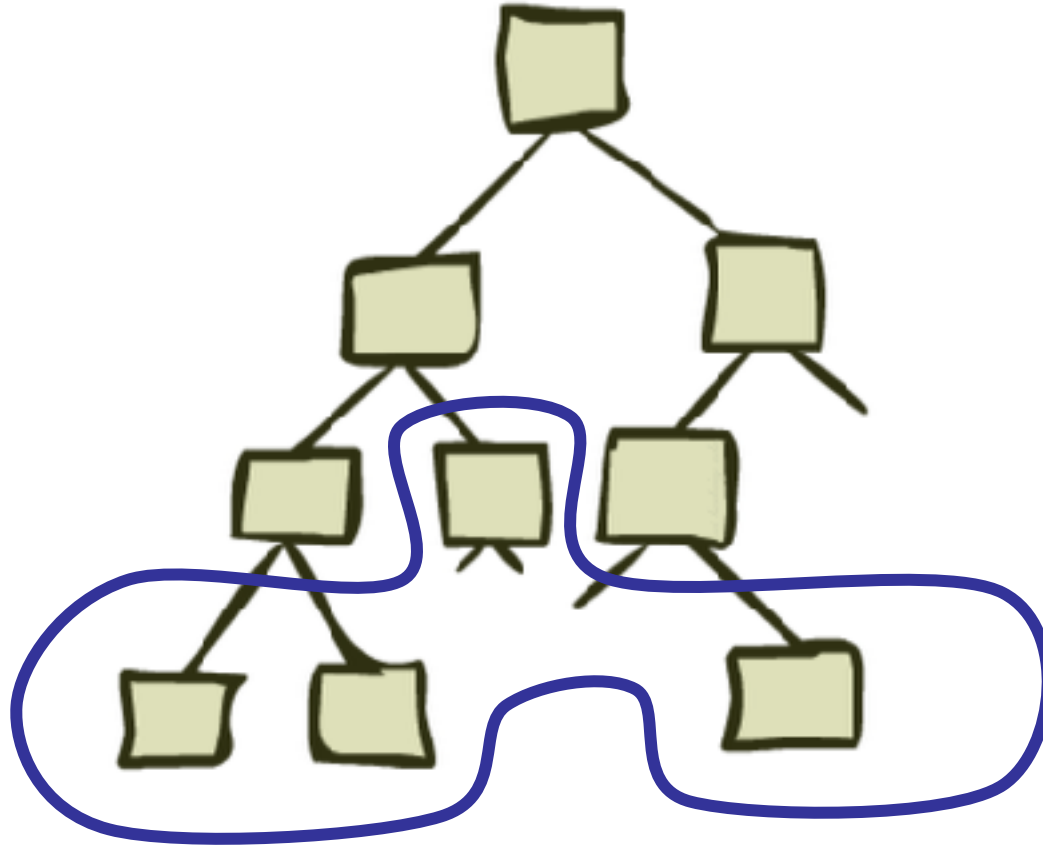
Important: Lots of repeated structure in the search tree!

So why would we ever use a search tree?

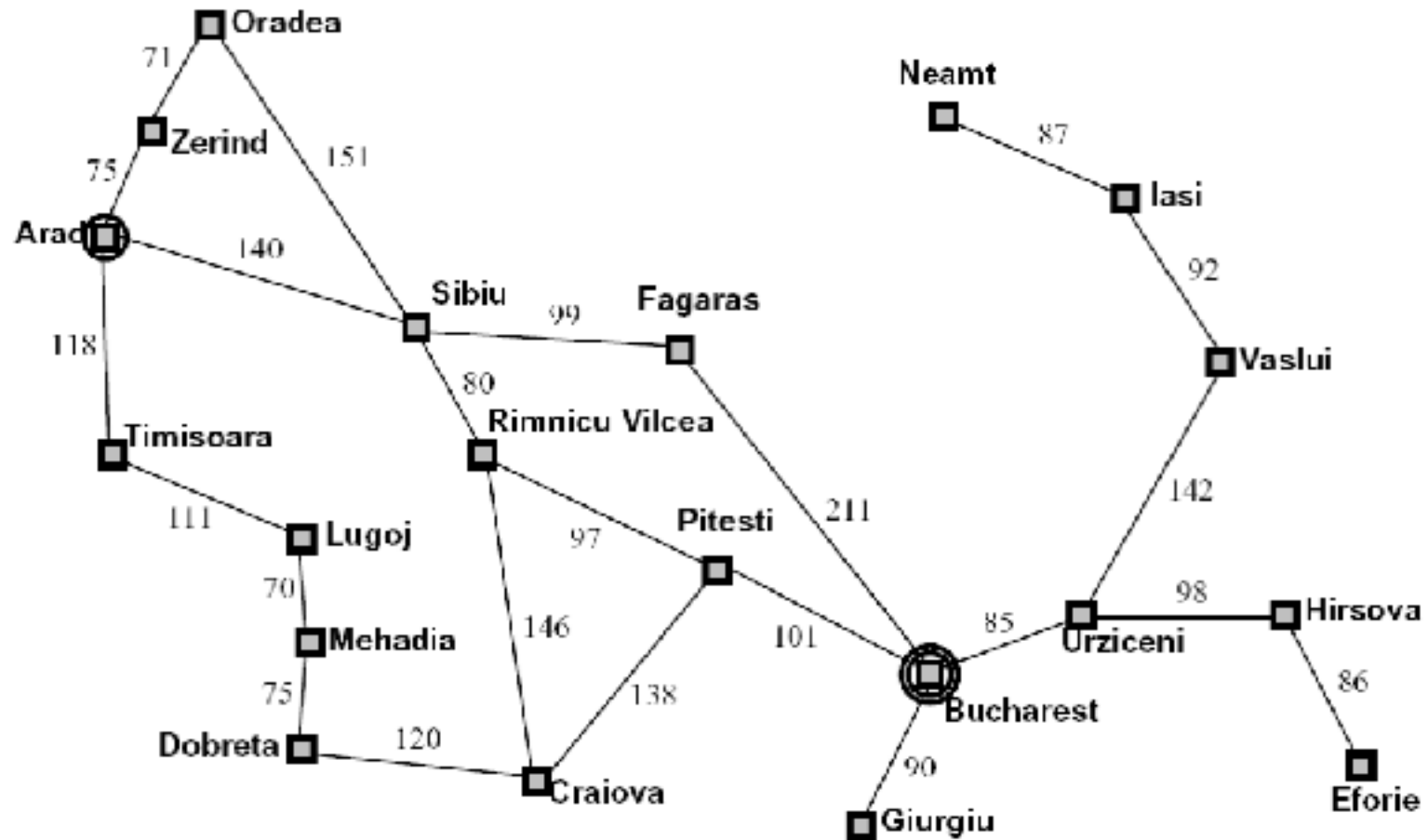
- 1) Cannot store “closed list” (previously visited nodes)
- 2) Graph happens to be a tree, so no reason to store closed list

# Tree Search

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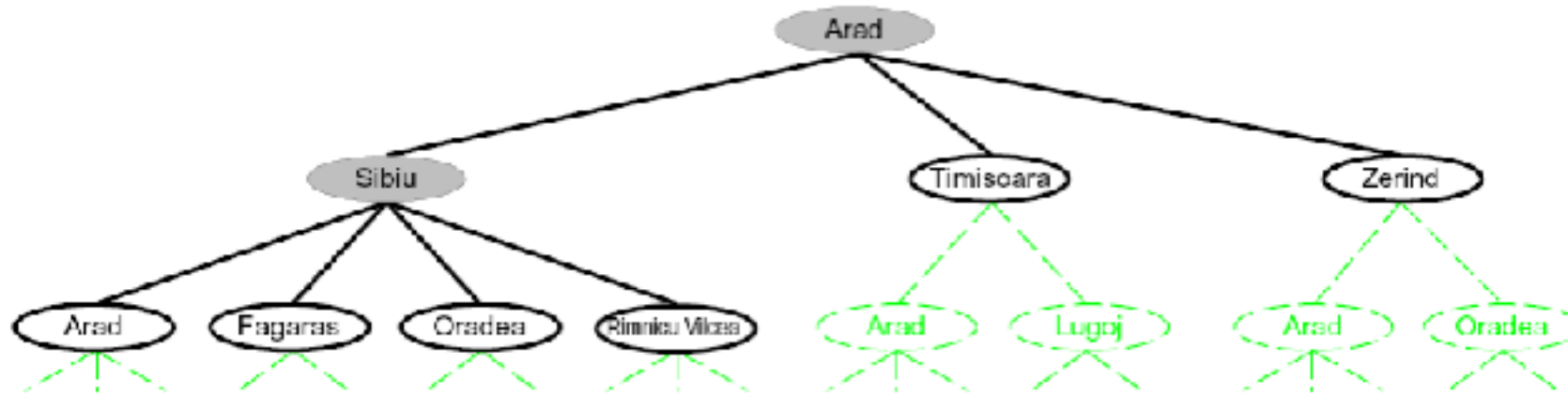


# Search Example: Romania





# Searching with a Search Tree



## ■ Search:

- Expand out potential plans (tree nodes)
- Maintain a **fringe** of partial plans under consideration
- Try to expand as few tree nodes as possible

# General Tree Search

```
function TREE-SEARCH( problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state of problem
  loop do
    if there are no candidates for expansion then return failure
    choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add the resulting nodes to the search tree
  end
```

- Important ideas:
  - Fringe
  - Expansion
  - Exploration strategy
- Main question: which fringe nodes to explore?

# Depth-First Search

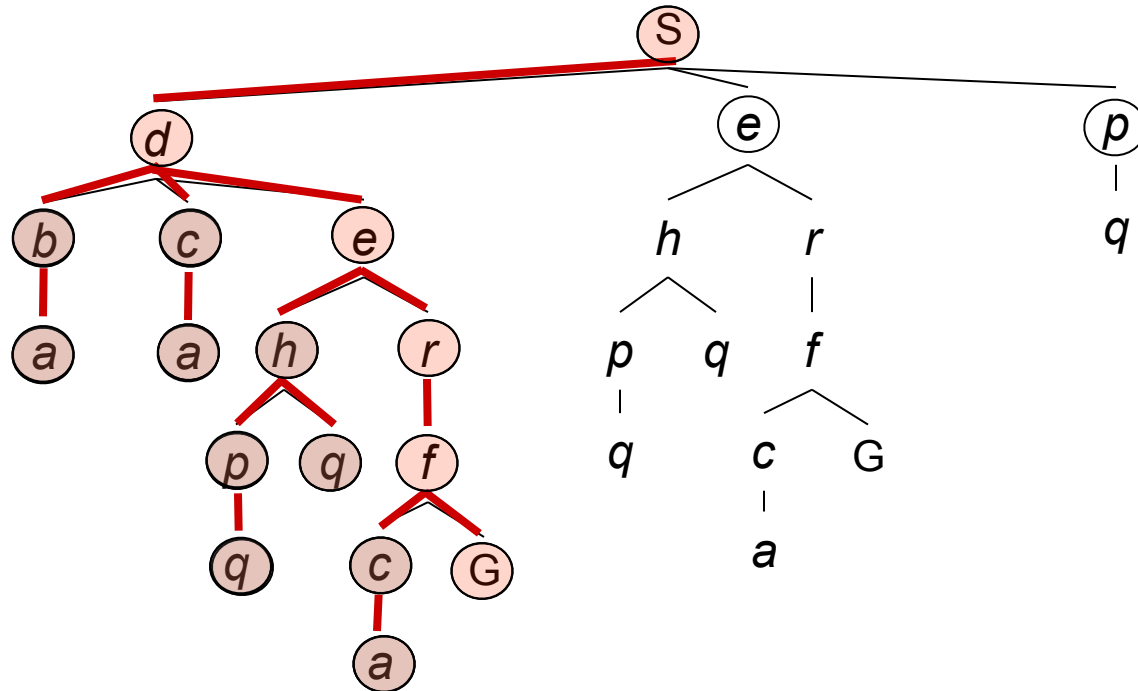
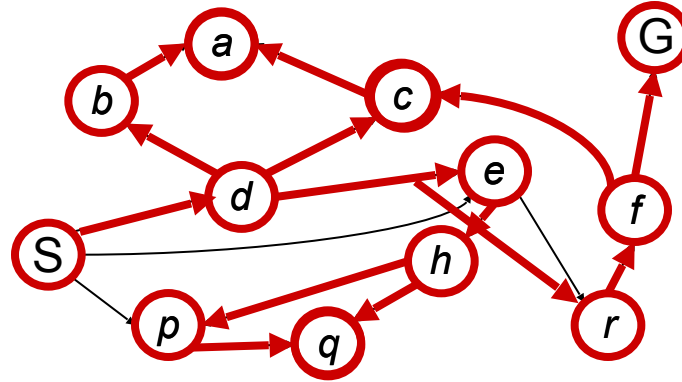
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# Depth-First Search

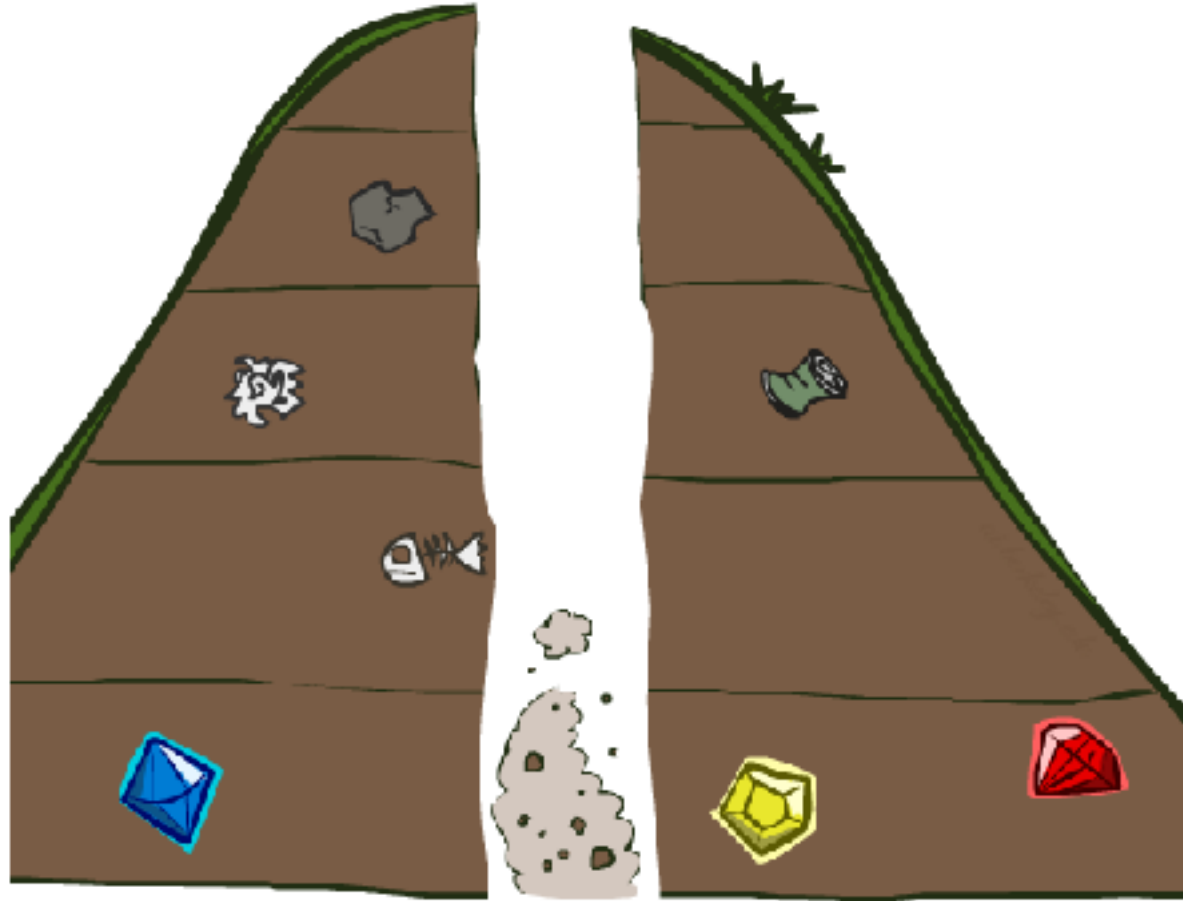
*Strategy: expand a  
deepest node first*

*Implementation: Fringe  
is a LIFO stack*



# Search Algorithm Properties

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# Search Algorithm Properties

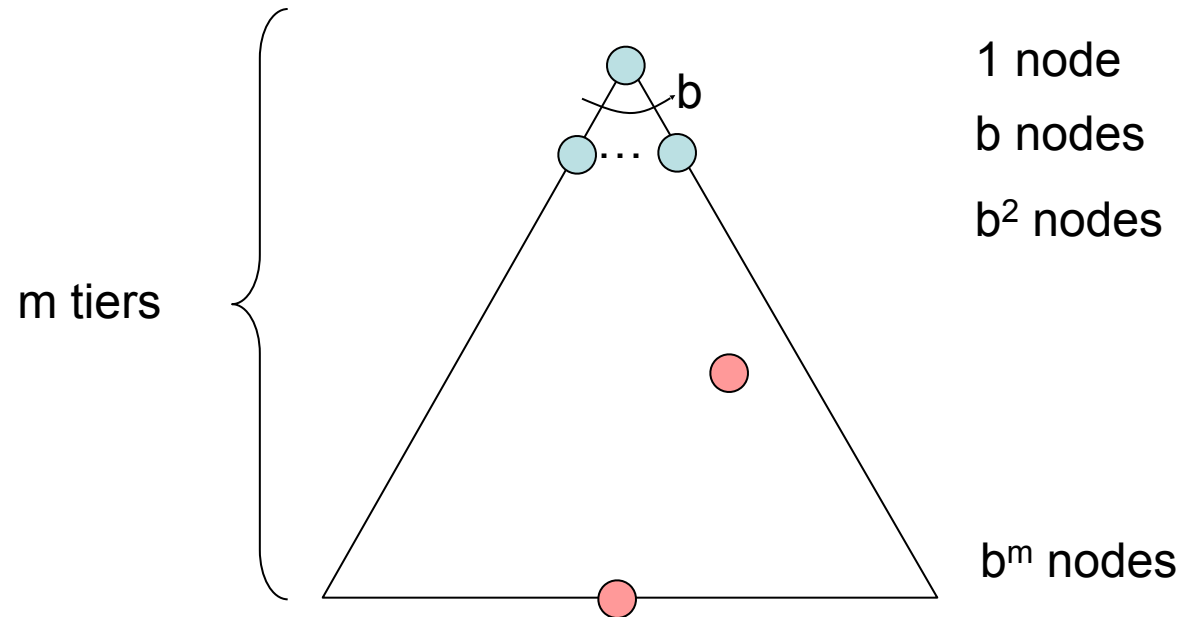
- Complete: Guaranteed to find a solution if one exists?
- Optimal: Guaranteed to find the least cost path?
- Time complexity?
- Space complexity?

- Cartoon of search tree:

- $b$  is the branching factor
- $m$  is the maximum depth
- solutions at various depths

- Number of nodes in entire tree?

- $1 + b + b^2 + \dots + b^m = O(b^{m+1})$



# Depth-First Search (DFS) Properties

- What nodes DFS expand?

- Some left prefix of the tree.
- Could process the whole tree!
- If  $m$  is finite, takes time  $O(b^m)$

- How much space does the fringe take?

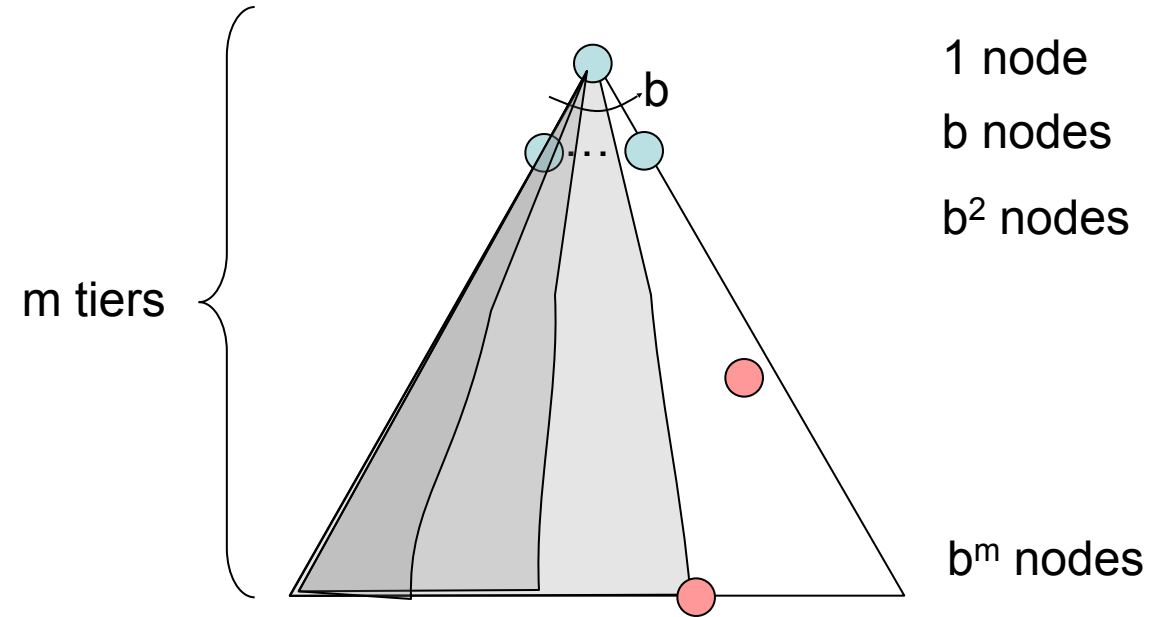
- Only has siblings on path to root, so  $O(bm)$

- Is it complete?

- $m$  could be infinite, so only if we prevent cycles (more later)

- Is it optimal?

- No, it finds the “leftmost” solution, regardless of depth or cost



# Breadth-First Search

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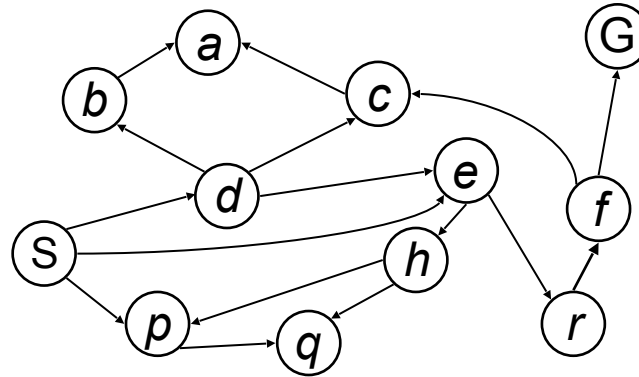




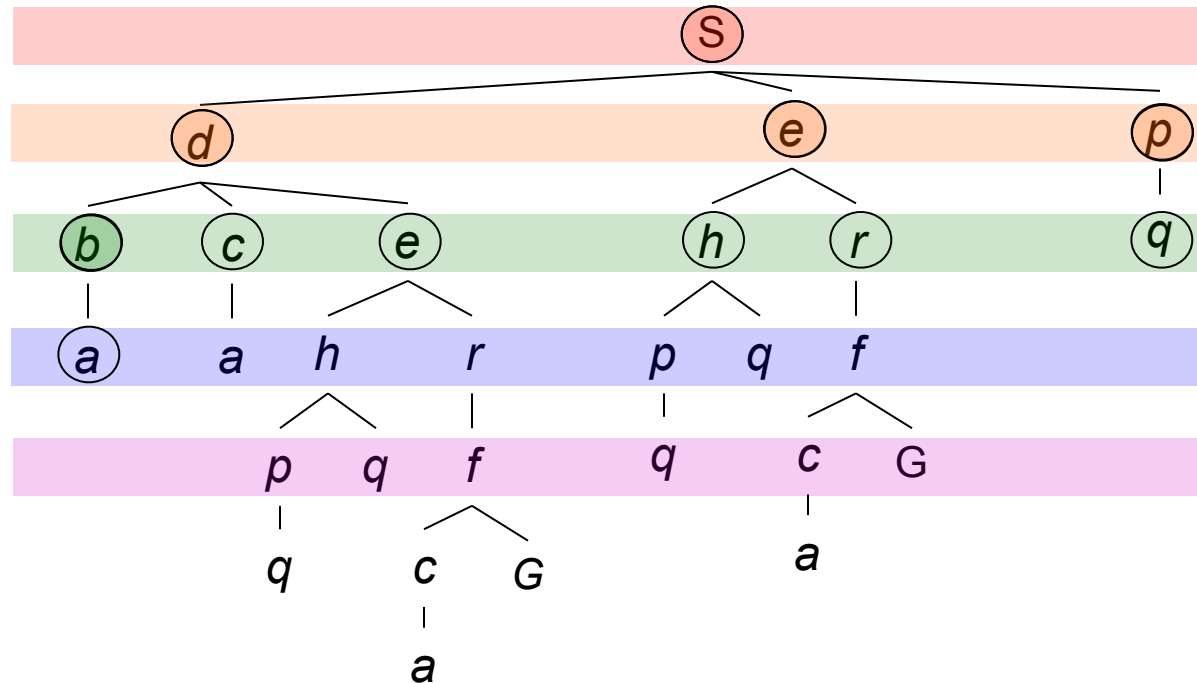
# Breadth-First Search

*Strategy: expand a shallowest node first*

*Implementation: Fringe is a FIFO queue*

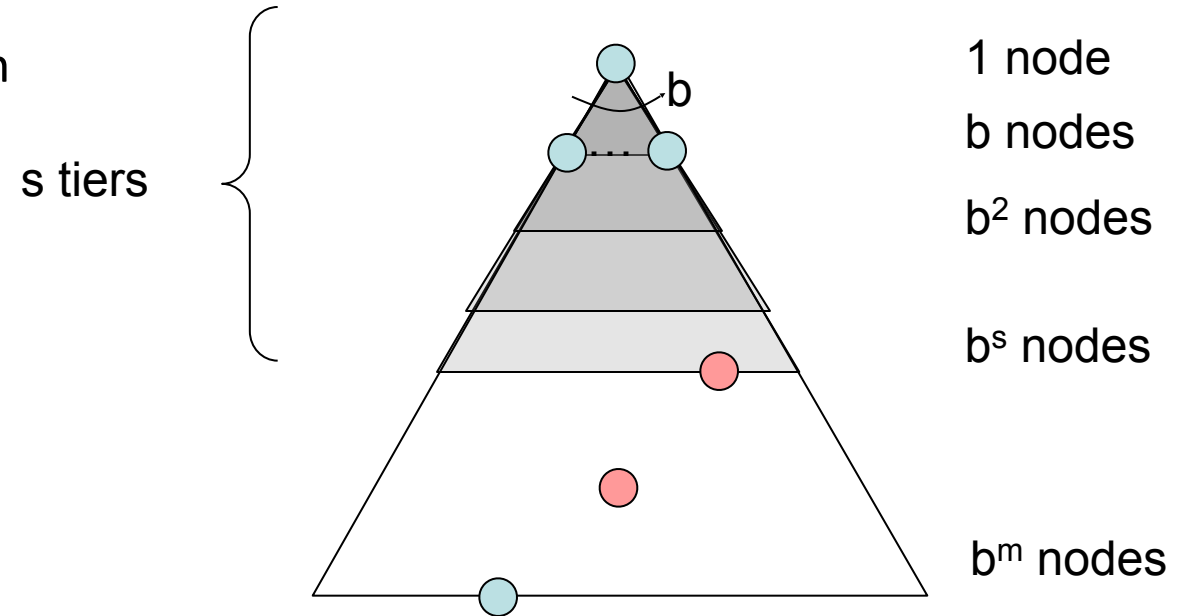


Search  
Tiers



# Breadth-First Search (BFS) Properties

- What nodes does BFS expand?
  - Processes all nodes above shallowest solution
  - Let depth of shallowest solution be  $s$
  - Search takes time  $O(b^s)$
- How much space does the fringe take?
  - Has roughly the last tier, so  $O(b^s)$
- Is it complete?
  - $s$  must be finite if a solution exists, so yes!
- Is it optimal?
  - Only if costs are all 1 (more on costs later)



# Quiz: DFS vs BFS

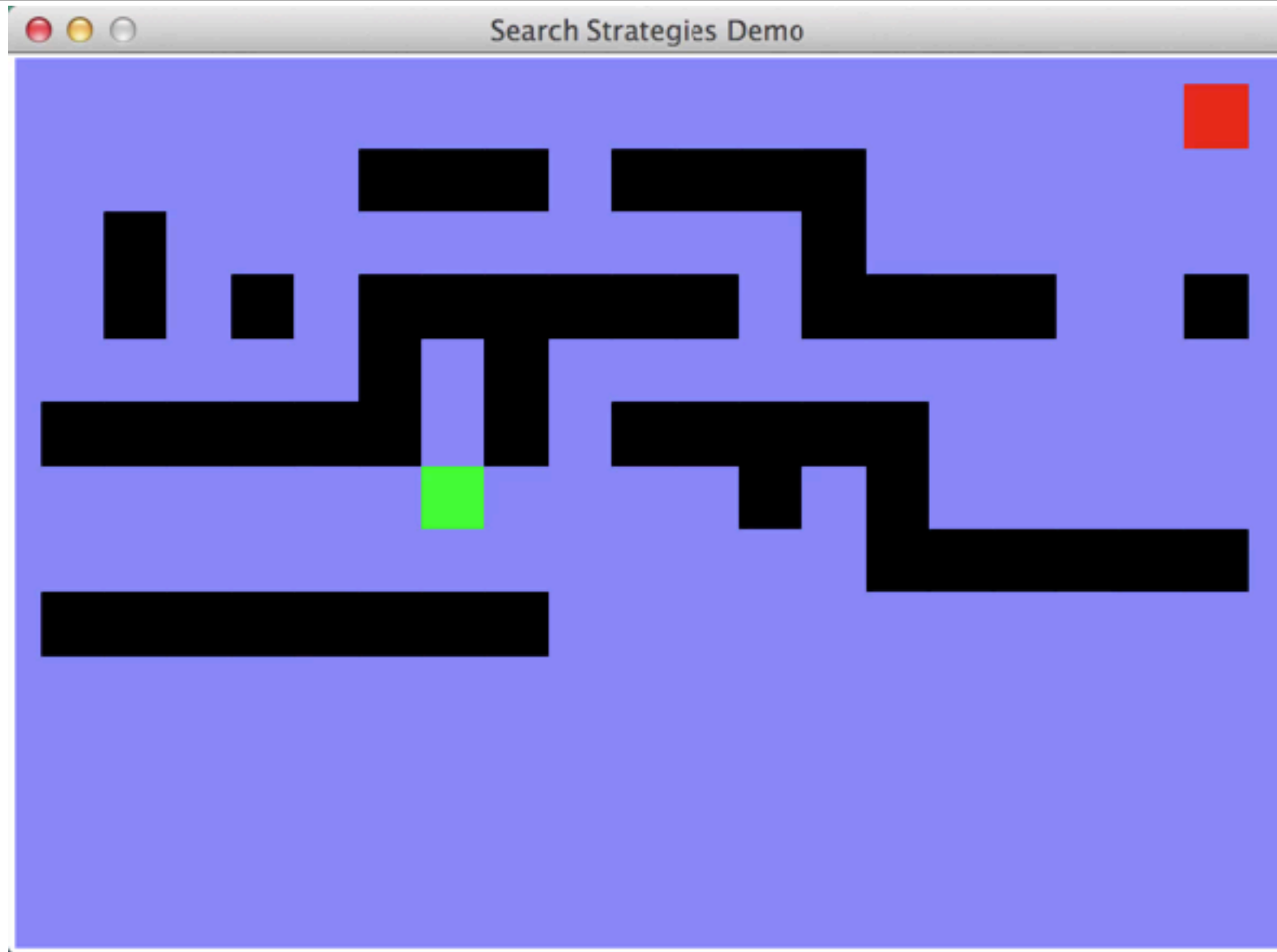


# Quiz: DFS vs BFS

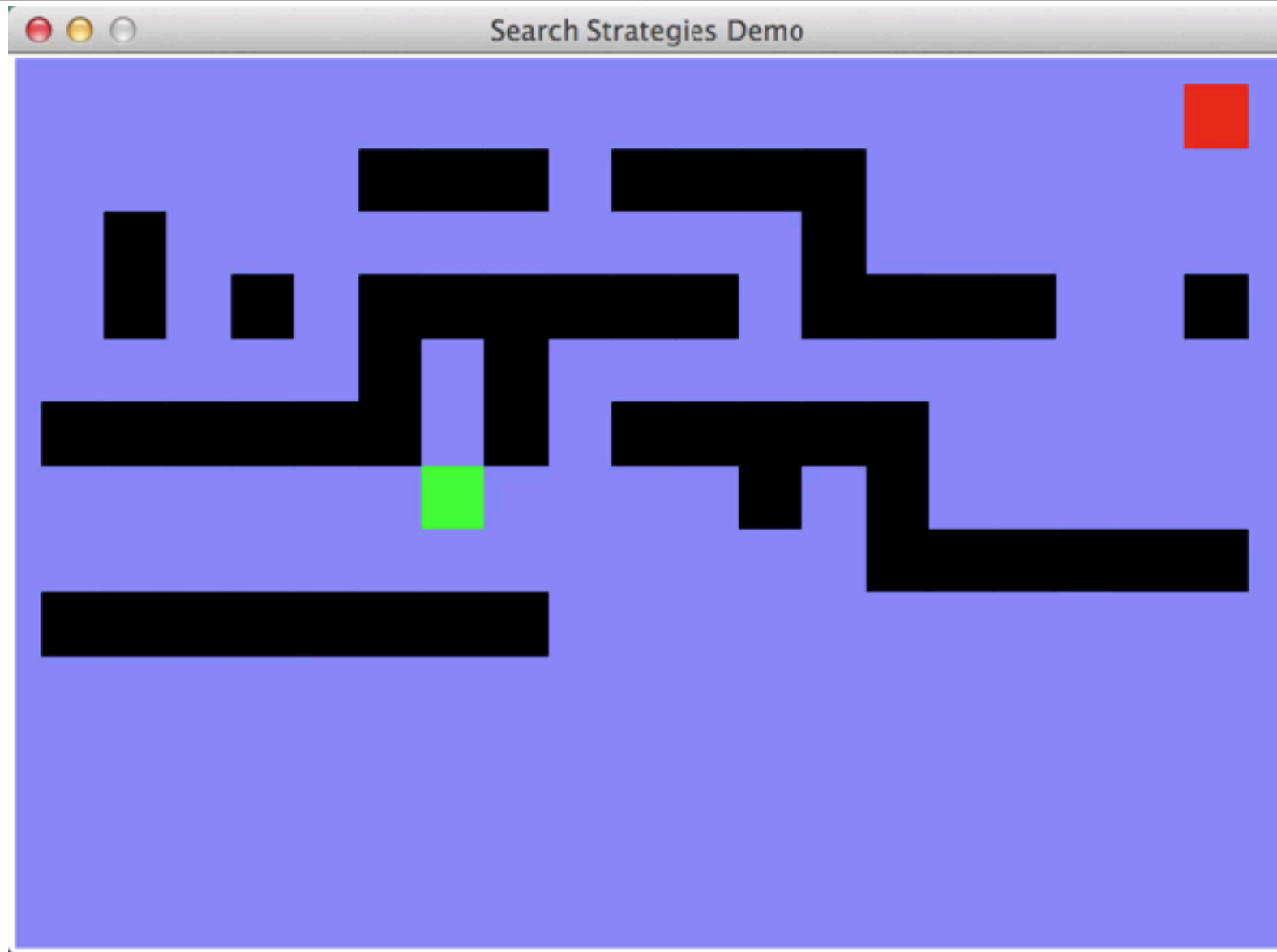
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- When will BFS outperform DFS?
- When will DFS outperform BFS?
- What is the worst case for each?

# Video of Demo Maze Water DFS/BFS (part 1)

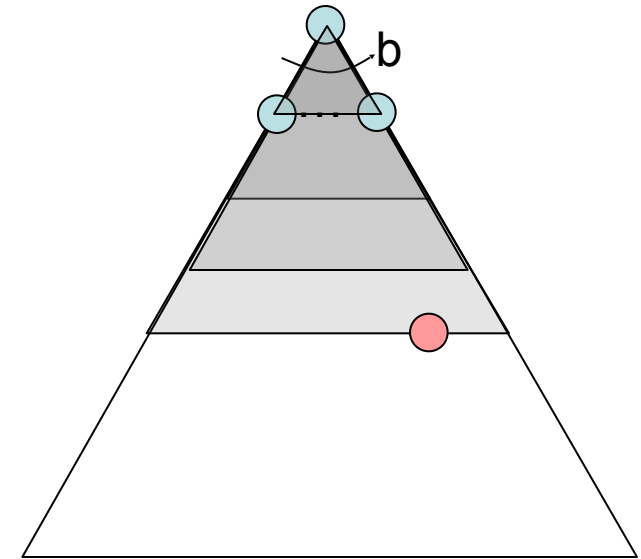


# Video of Demo Maze Water DFS/BFS (part 2)

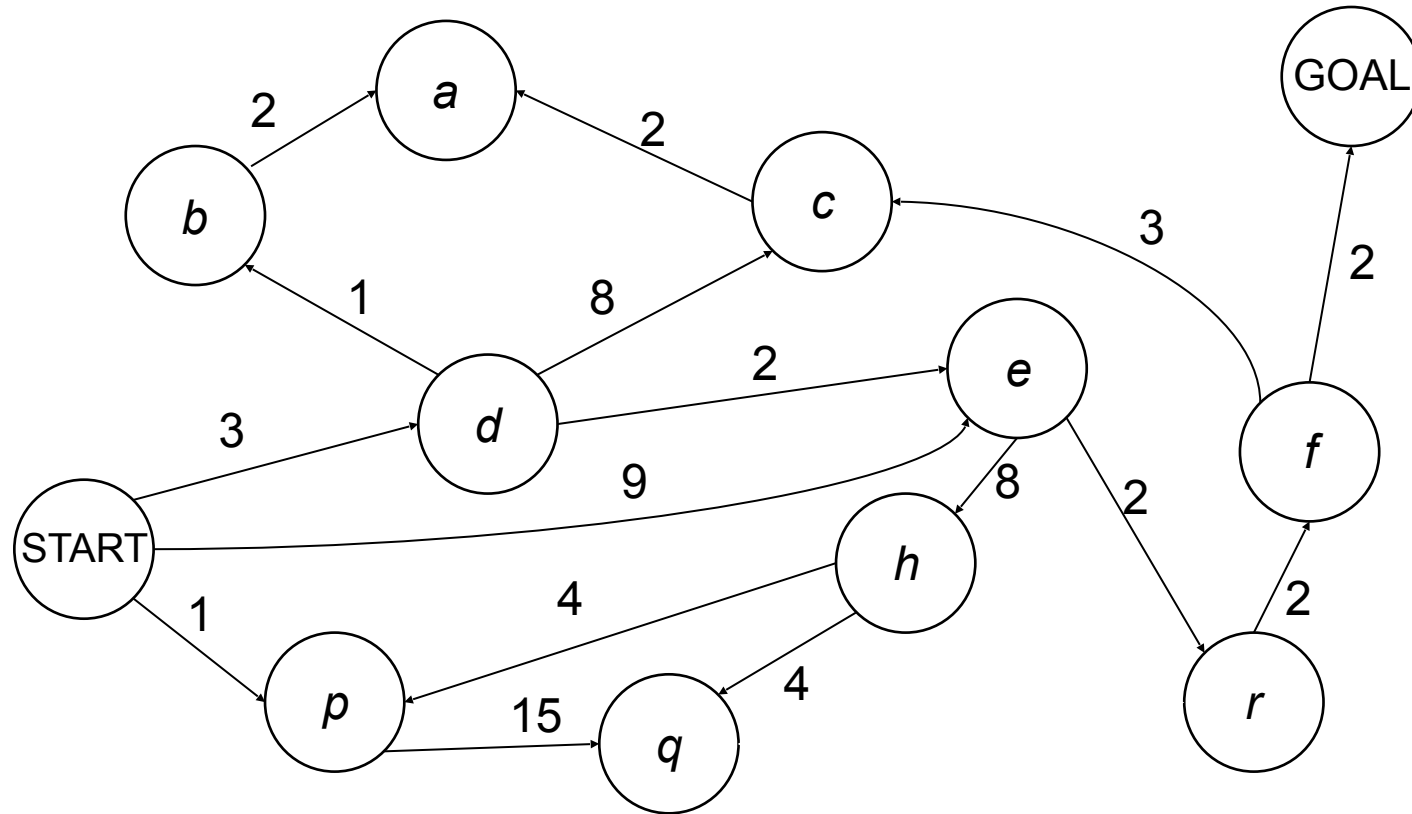


# Iterative Deepening

- Idea: get DFS's space advantage with BFS's time / shallow-solution advantages
  - Run a DFS with depth limit 1. If no solution...
  - Run a DFS with depth limit 2. If no solution...
  - Run a DFS with depth limit 3. ....
- Isn't that wastefully redundant?
  - Generally most work happens in the lowest level searched, so not so bad!



# Cost-Sensitive Search



BFS finds the shortest path in terms of number of actions. It does not find the least-cost path. We will now cover a similar algorithm which does find the least-cost path.



# Uniform Cost Search

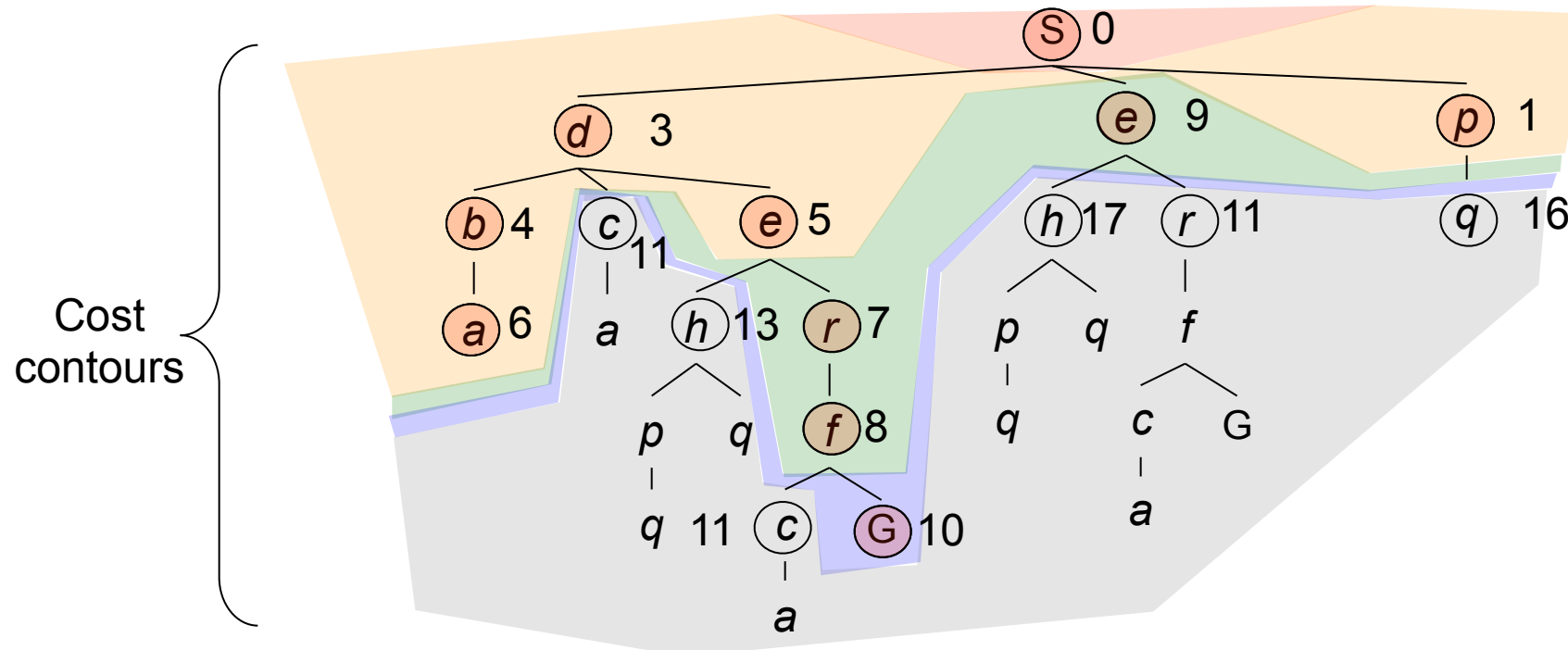
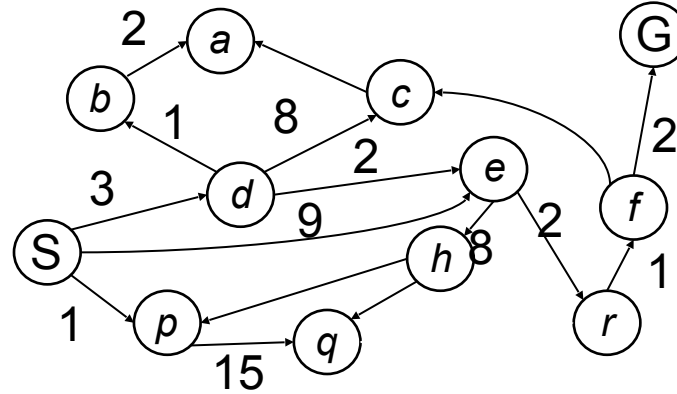
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# Uniform Cost Search

*Strategy: expand a  
cheapest node first:*

*Fringe is a priority queue  
(priority: cumulative cost)*



# Uniform Cost Search (UCS) Properties

- What nodes does UCS expand?

- Processes all nodes with cost less than cheapest solution!
- If that solution costs  $C^*$  and arcs cost at least  $\epsilon$ , then the “effective depth” is roughly  $C^*/\epsilon$
- Takes time  $O(b^{C^*/\epsilon})$  (exponential in effective depth)

- How much space does the fringe take?

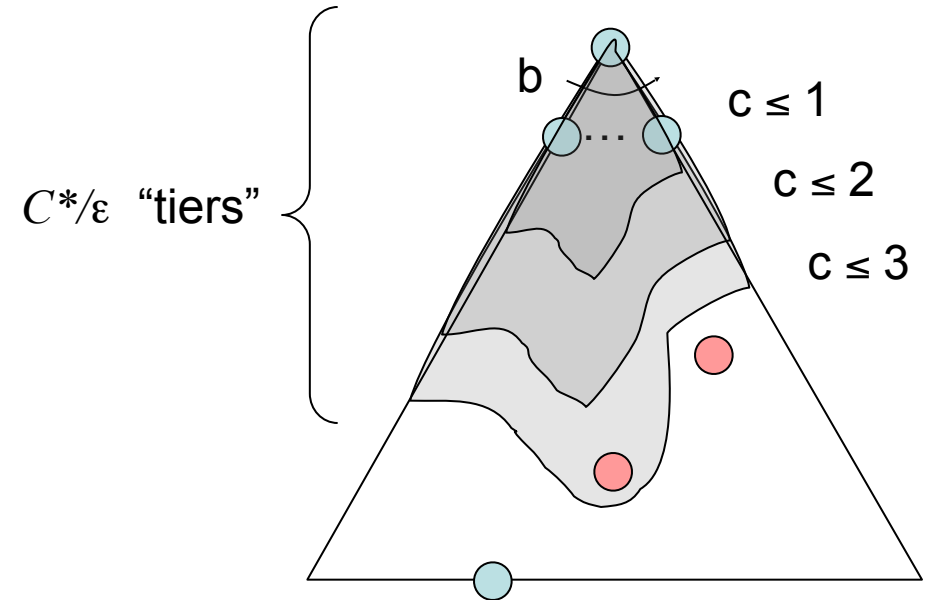
- Has roughly the last tier, so  $O(b^{C^*/\epsilon})$

- Is it complete?

- Assuming best solution has a finite cost and minimum arc cost is positive, yes!

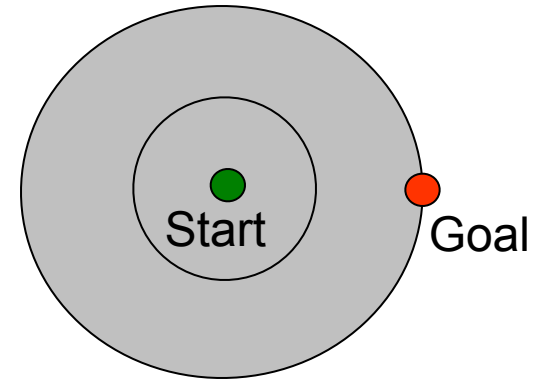
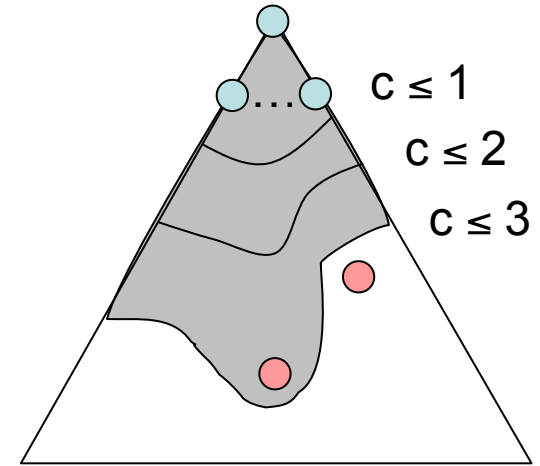
- Is it optimal?

- Yes! (Proof next lecture via  $A^*$ )

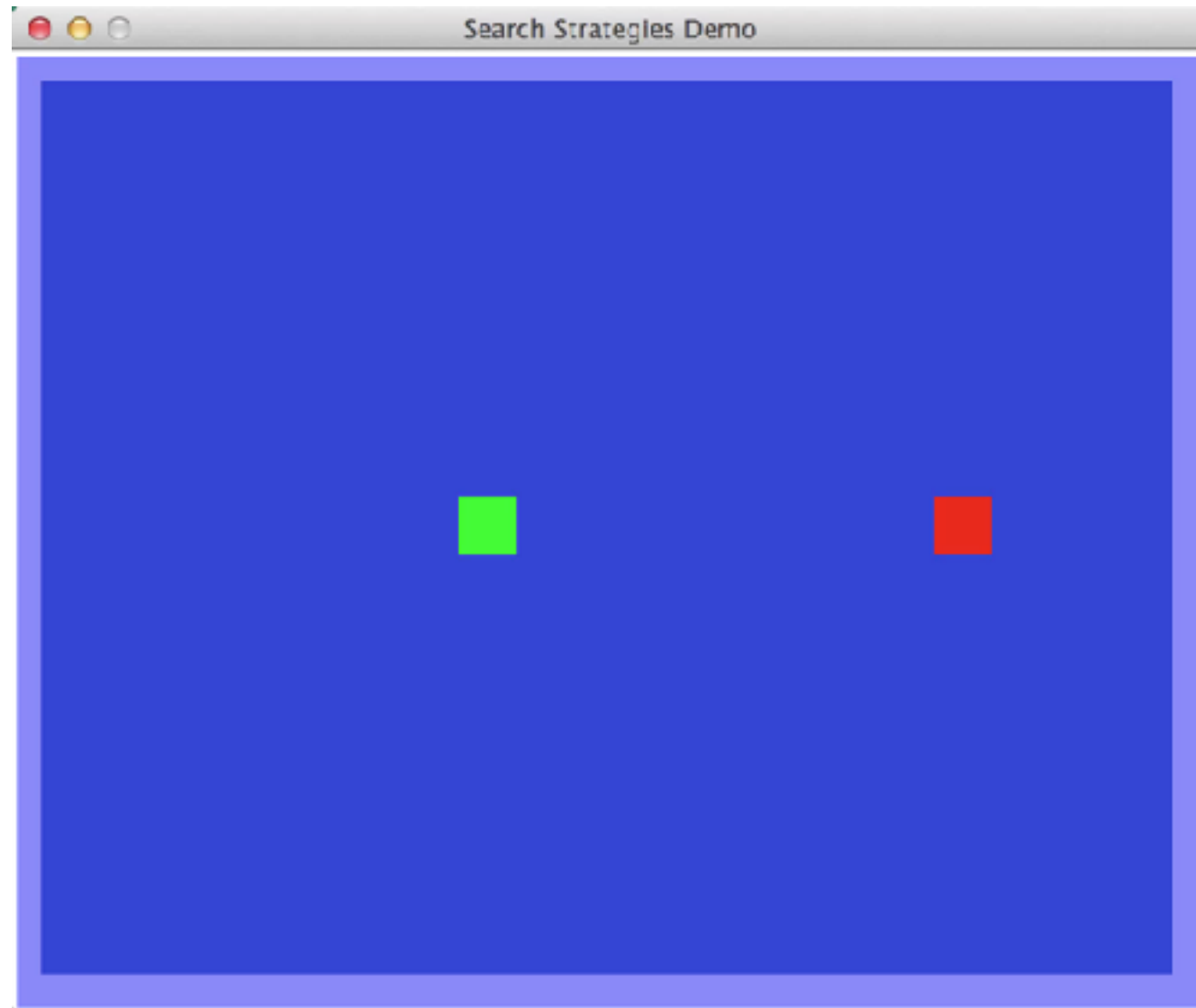


# Uniform Cost Issues

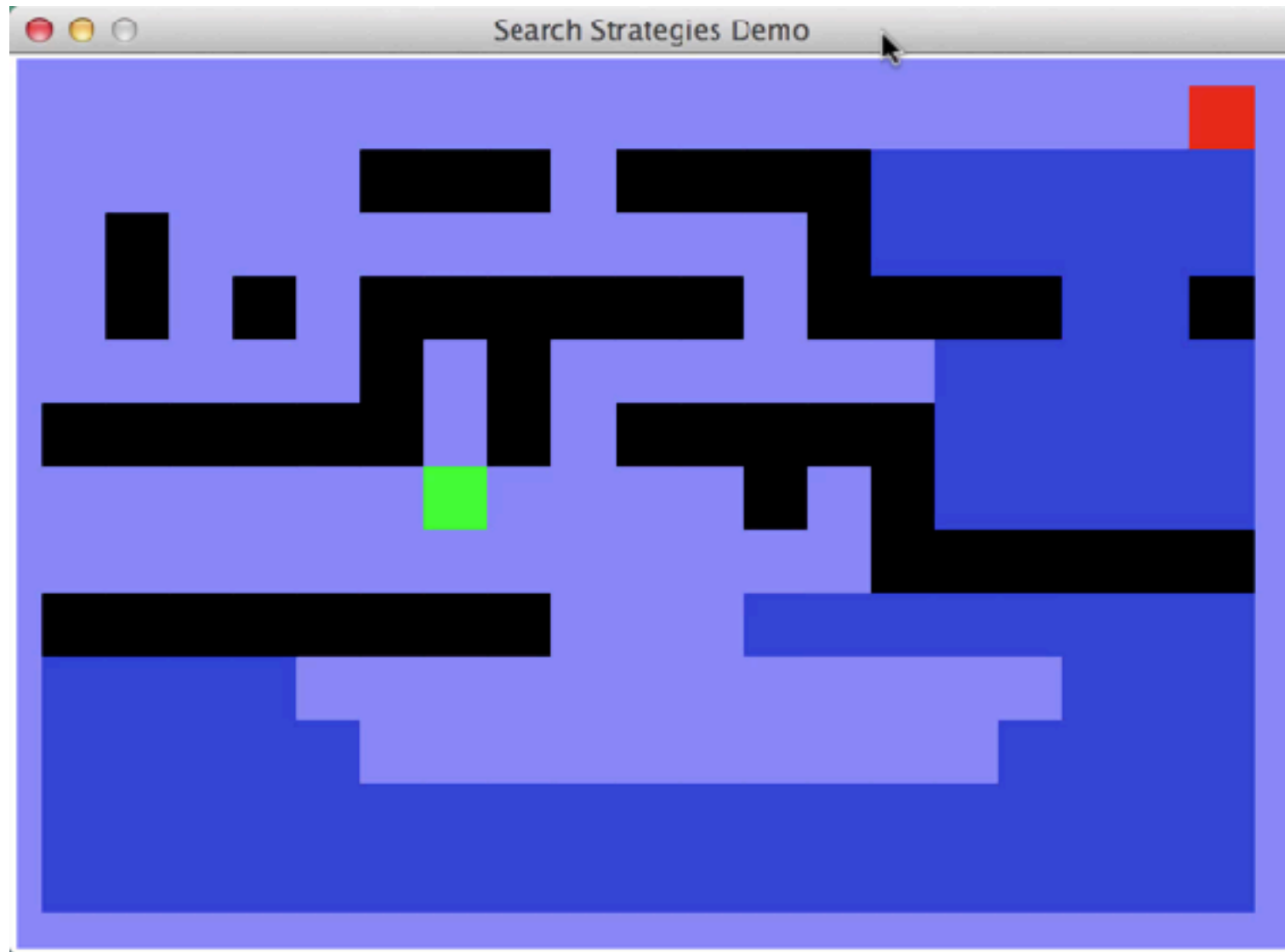
- Remember: UCS explores increasing cost contours
- The good: UCS is complete and optimal!
- The bad:
  - Explores options in every “direction”
  - No information about goal location
- We’ll fix that soon!



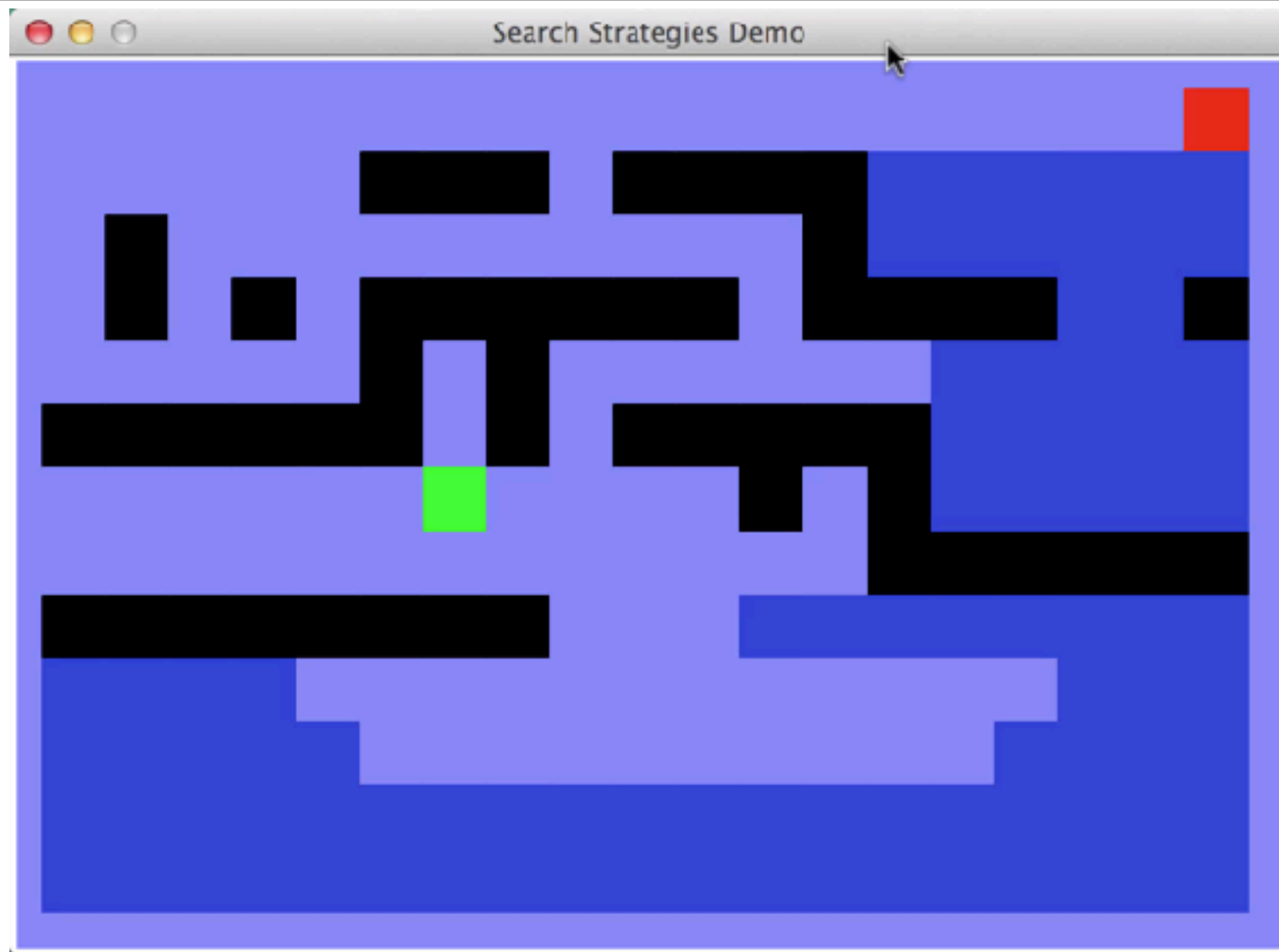
# Video of Demo Empty UCS



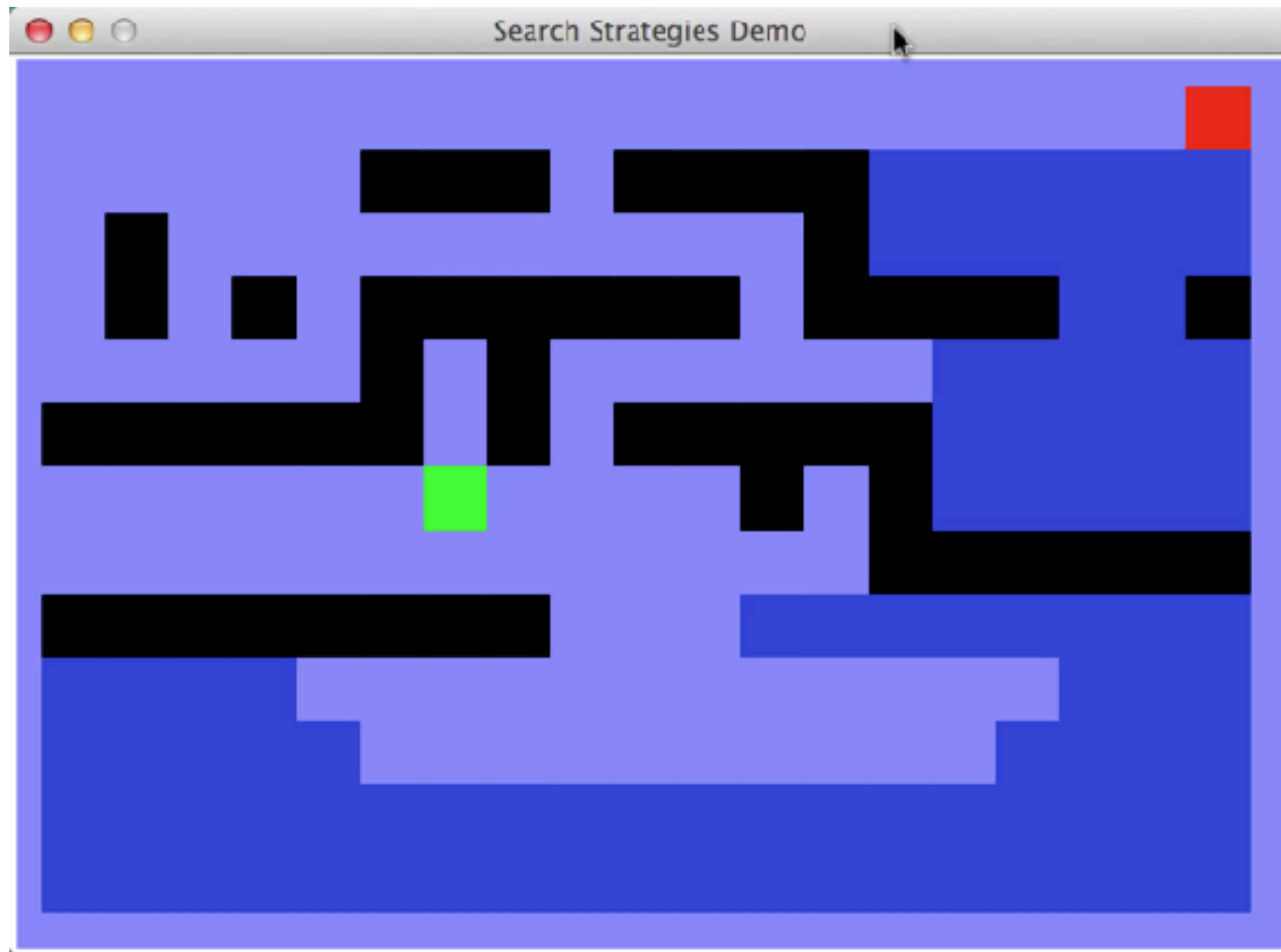
# Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 1)



## Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 2)



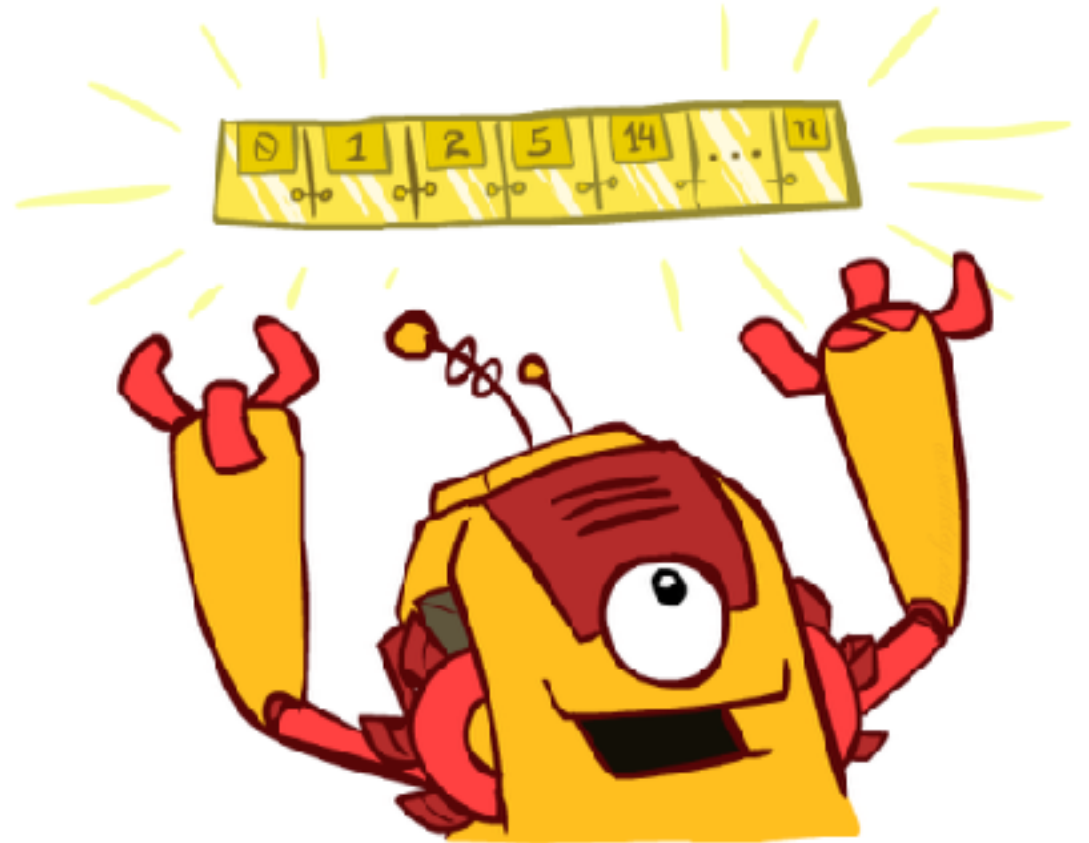
## Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 3)





# The One Queue

- All these search algorithms are the same except for fringe strategies
  - Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
  - Practically, for DFS and BFS, you can avoid the  $\log(n)$  overhead from an actual priority queue, by using stacks and queues
  - Can even code one implementation that takes a variable queuing object



# Search and Models

- Search operates over models of the world
  - The agent doesn't actually try all the plans out in the real world!
  - Planning is all “in simulation”
  - Your search is only as good as your models...

