# Teaching Teammates in Ad Hoc Teams

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Joint work with

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- Autonomous agents
- Robotics
- Machine learning (RL)
- Multiagent systems

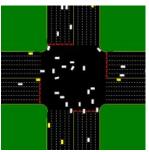
- Autonomous agents
- Robotics
- Machine learning (RL)
- Multiagent systems
  - e-commerce
  - mechanism design

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# **Teamwork**



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- Typical scenario: pre-coordination
  - People practice together
  - Robots given coordination languages, protocols
  - "Locker room agreement" (Stone & Veloso, '99)

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  - Unknown teammates (programmed by others)

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Goal: Create a good team player

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Goal: Create a good team player

- Minimal representative scenarios
  - One teammate, no communication
  - Fixed and known behavior

### **Scenarios**

• Cooperative normal form game (w/ Kaminka & Rosenschein)

M1	$b_0$	$b_1$	$b_2$
$a_0$	25	1	0
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$a_2$	0	33	40

Cooperative k-armed bandit









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### **Formalism**

- Agent A in our control: actions  $a_0, a_1, \dots a_{x-1}$
- Agent B reacts in a fixed way:  $b_0, b_1, \ldots, b_{y-1}$

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- Payoff from joint action  $(a_i, b_j)$ :  $m_{i,j}$
- ullet Highest payoff  $m^*$  always at  $(a_{x-1},b_{y-1})$
- Agent B's default action:  $b_0$

- Agent A's goal: action sequence with highest reward
  - Undiscounted, medium-term (finite)
  - Depends on Agent B's strategy

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Agent B best response  $\Longrightarrow$  can do better

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• Reward sequence: 25

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- Reward sequence: 25, 0

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- Reward sequence: 25, 25, 25, ...
- Reward sequence: 25, 0, 40, 40, 40, ...
- How?

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- Reward sequence: 25, 25, 25, ...
- Reward sequence: 25, 0, 40, 40, 40, ...
- Reward sequence: 25, 10

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- Reward sequence: 25, 25, 25, ...
- Reward sequence: 25, 0, 40, 40,... (65 from 1st 3)
- Reward sequence: 25, 10, 33, 40, ... (68 from 1st 3)

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Agent B best response  $\Longrightarrow$  or even better

M1	$b_0$	$b_1$	$b_2$
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- Reward sequence: 25, 25, 25, ...
- Reward sequence: 25, 0, 40, 40,...
- Reward sequence: 25, 10, 33, 40, ...

**Cost**: 15+40=**55** 

**Cost**: 15+30+7=**52** 

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- Reward sequence: 25, 25, 25, ...
- Reward sequence: 25, 0, 40, 40,... Cost: 55, **Length: 2**
- Reward sequence: 25, 10, 33, 40, ... Cost: 52, Length: 3

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- Example: mem = 4,  $\epsilon = 0.1$ 
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  - $-BR(a_1, a_0, a_1, a_1) = b_1$

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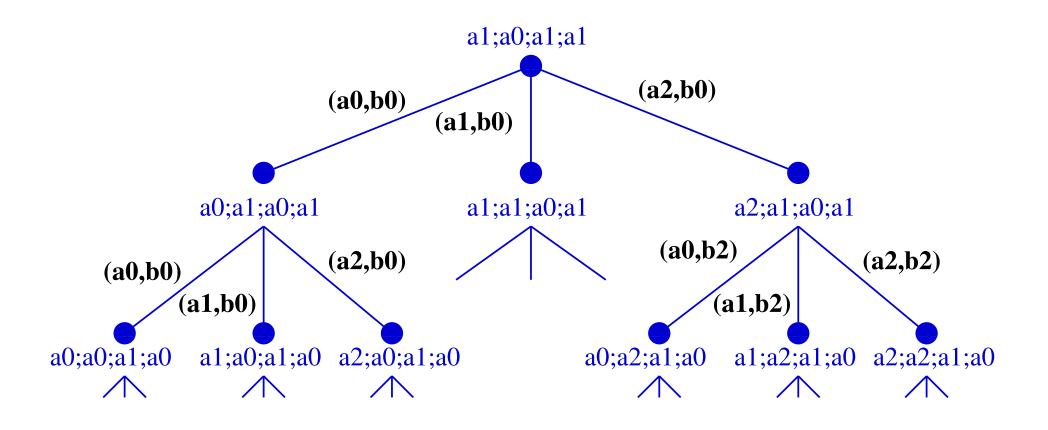
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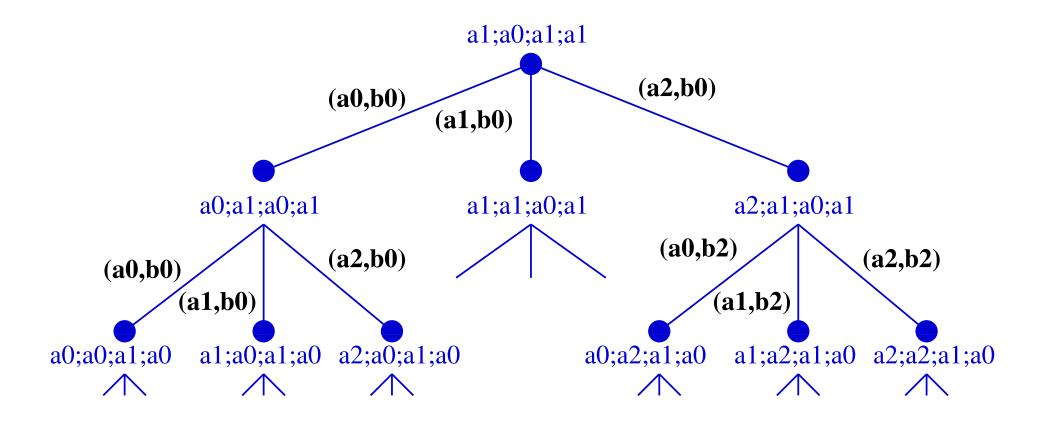
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  - Agent A: action determines payoff and next history

#### **Extensive Form Version**



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Stick with iterated normal form for presentation, algorithms

#### **Questions**

- Can we find the optimal action sequence efficiently?
- How long can the optimal action sequences be?

#### Cases

- Deterministic teammate, 1-step memory (mem= 1,  $\epsilon = 0$ )
- Longer teammate memory  $(mem>1, \epsilon=0)$
- Teammate non-determinism (mem> 1,  $\epsilon$  > 0)

- Define  $S_n^*(a_i, b_j)$  = optimal sequence of length n
- ullet Define  $S_0^*(a_i,b_j)$  to be cost 0 if  $m_{i,j}=m_*$ , else  $\infty$

$$S_0^*(a_2,b_2)$$
 Cost 0

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$S_0^*(a_2,b_2)$	Cost 0
$S_2^*(a_0, b_0)$	Cost 15+40
$S_2^*(a_1, b_0)$	Cost 30+7
$S_2^*(a_2, b_0)$	Cost 40

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- ullet Find  $S_n^*(a_i,b_j)$  using  $S_{n-1}^*$ 's  $(O(d),d=\dim(M))$ 
  - Either  $S_{n-1}^*(a_i,b_j)$  or
  - Best sequence that prepends  $(a_i, b_j)$  to  $S_{n-1}^*(a_{act}, b_{BR(a_i)})$

$$S_3^*(a_0,b_0)$$
 ?  $S_2^*(a_0,b_0)$  Cost 55  $S_2^*(a_1,b_0)$  Cost 37  $S_2^*(a_2,b_0)$  Cost 40

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$S_3^*(a_0,b_0)$	Cost 52
$S_2^*(a_0, b_0)$	Cost 55
$S_2^*(a_1, b_0)$	Cost 37+15
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  - Best sequence that prepends  $(a_i, b_j)$  to  $S_{n-1}^*(a_{act}, b_{BR(a_i)})$
- Sufficient to calculate  $S_n^*(a_i, b_0), \forall i < x$

 $loop O(d^2)$ 

(O(d), d = dim(M))

- How high do we need to let n get?

$$S_3^*(a_0, b_0)$$
 Cost 52  
 $S_2^*(a_0, b_0)$  Cost 55  
 $S_2^*(a_1, b_0)$  Cost 37+15  
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M2	$b_0$	$b_1$	$b_2$		$b_{y-3}$	$b_{y-2}$	$b_{y-1}$
$a_0$	$100 - \delta$	0	0	• • •	0	0	0
$a_1$	$100-2\delta$	$100 - \delta$	0		:	0	0
$a_2$	0	$100-2\delta$	$100 - \delta$			1	0
:	;		100	200			
$a_{x-3}$	0	:		200	$100 - \delta$	0	0
$a_{x-2}$	0	0	:		$100-2\delta$	$100 - \delta$	0
$a_{x-1}$	a	0	0		0	$100-2\delta$	100

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$a_0$	$100 - \delta$	0	0	• • •	0	0	0
$a_1$	$100-2\delta$	$100 - \delta$	0		:	0	0
$a_2$	0	$100-2\delta$	$100 - \delta$			1	0
:	:		200	200			:
$a_{x-3}$	0	1		$(\mathcal{O}_{\mathcal{A}}}}}}}}}}$	$100 - \delta$	0	0
$a_{x-2}$	0	0	:		$100-2\delta$	$100 - \delta$	0
$a_{x-1}$	a	0	0		0	$100-2\delta$	100

#### **Questions**

- ullet Find the optimal action sequence efficiently?  $O(d^3)$
- Maximum length of optimal sequences?  $\min(x, y)$

#### Cases

- Deterministic teammate, 1-step memory (mem= 1,  $\epsilon = 0$ )
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History  $[a_2; a_1; a_0]$ Response  $b_2$ 

M3	$b_0$	$b_1$	$b_2$
$a_0$	0	30	50
$a_1$	41	20	0
$a_2$	99	20	100

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NP-hard: reduction from Hamiltonian Path (Littman)

- Theorem:  $\exists M$  with optimal seq.  $(\min(x,y)-1)*mem+1$
- Conjecture: No seq. longer than  $(\min(x,y)-1)*mem+1$

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- Conjecture: No seq. longer than  $(\min(x,y)-1)*mem+1$ 
  - Can only prove no seq. longer than  $\min(x,y) * x^{mem-1}$

M2	$b_0$	$b_1$	$b_2$		$b_{y-3}$	$b_{y-2}$	$b_{y-1}$
$a_0$	$100 - \delta$	0	0	• • •	0	0	0
$a_1$	$100-2\delta$	$100 - \delta$	0		:	0	0
$a_2$	0	$100-2\delta$	$100 - \delta$			1	0
:	:		200	100			
$a_{x-3}$	0	:		$(\mathcal{O}_{\mathcal{A}_{\mathcal{A}}})$	$100 - \delta$	0	0
$a_{x-2}$	0	0	:		$100-2\delta$	$100 - \delta$	0
$a_{x-1}$	Ø	0	0		0	$100-2\delta$	100

### **Questions**

Find the optimal action sequence efficiently?

no

Maximum length of optimal sequences?

#### Cases

- Deterministic teammate, 1-step memory (mem= 1,  $\epsilon = 0$ )
- Longer teammate memory

(mem> 1,  $\epsilon = 0$ )

Teammate non-determinism

(mem> 1,  $\epsilon > 0$ )

#### Questions

- Find the optimal action sequence efficiently?
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- Longer teammate memory  $(mem>1, \epsilon=0)$
- Teammate non-determinism (mem> 1,  $\epsilon > 0$ )

- $EV(a_i, b_j) = (1 \epsilon)m_{i,j} + \frac{\epsilon}{y}(\sum_{k=0}^{y-1} m_{i,k})$ 
  - Cost now sum of  $m^* EV(a_i, b_j)$  over sequence

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$$\epsilon=0$$
:  $m_*$  at  $(a_3,b_3)$  L( $S^*$ )=10  $\epsilon=0.1$ :  $m_*$  at  $(a_3,b_3)$  L( $S^*$ )=8  $\epsilon=0.3$ :  $m_*$  at  $(a_3,b_3)$  L( $S^*$ )=3  $\epsilon=0.4$ :  $m_*$  at  $(a_2,b_2)$  L( $S^*$ )=3

M4	$b_0$	$b_1$	$b_2$	$b_3$
$a_0$	25	0	0	0
$a_1$	88	90	99	80
$a_2$	70	98	99	80
$a_3$	70	70	98	100

#### **Teammate Non-Determinism**

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- Algorithm and theorems hold unchanged
  - Except when  $\epsilon = 1$

#### **Questions**

Find the optimal action sequence efficiently?

no

Maximum length of optimal sequences?

#### Cases

- Deterministic teammate, 1-step memory (mem= 1,  $\epsilon = 0$ )
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mem=1	1	2	3	4	5	6	7	8	9	10
$3 \times 3$	104	852	44							

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mem=1	]	2	3	4	5	6	7	8	9	10
$3 \times 3$	104	852	44							
$4 \times 4$	12	825	158	5						
$5 \times 5$	3	662	316	19	0					
$6 \times 6$	0	465	489	45	1	0				
$7 \times 7$	0	349	565	81	5	0	0			
$8 \times 8$	0	236	596	159	8	1	0	0		
$9 \times 9$	0	145	640	193	20	2	0	0	0	
$10 \times 10$	0	72	636	263	29	0	0	0	0	0

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$9 \times 9$	0	145	640	193	20	2	0	0	0	
10 × 10	0	72	636	263	29	0	0	0	0	0

mem=3	1	2	3	4	5	6	7	8	9	10	11
$3 \times 3$	98	178	344	340	28	8	4	0	0	0	0
$4 \times 4$	15	76	266	428	134	60	21	0	0	0	0
$5 \times 5$	1	19	115	408	234	145	71	7	0	0	0
$6 \times 6$	0	0	22	282	272	222	164	27	11	0	0
$7 \times 7$	0	0	5	116	293	282	220	55	17	10	1

# **Robot Experiments**

• In progress...



#### **Related Work**

#### **Game Theory**

Multiagent learning (Claus & Boutilier, '98), (Littman, '01),

(Conitzer & Sandholm, '03), (Powers & Shoham, '05), (Chakraborty & Stone, '08)

- Economic repeated games (Hart & Mas-Colell, '00), (Neyman & Okada, '00)
- Fictitious play (Brown, '51)
- Adaptive play (Young, '93)

#### **Opponent Modeling**

- Intended plan recognition (Sidner, '85), (Lochbaum, '91), (Carberry, '01)
- SharedPlans (Grosz & Kraus, '96)
- Recursive Modeling (Vidal & Durfee, '95)

#### **Ad Hoc Teams**

- Ad hoc team player is an individual
  - Unknown teammates (programmed by others)
- May or may not be able to communicate
- Teammates likely sub-optimal: no control





Goal: Create a good team player

- Minimal representative scenarios
  - One teammate, no communication
  - Fixed and known behavior: best response

#### **Scenarios**

• Cooperative normal form game (w/ Kaminka & Rosenschein)

M1	$b_0$	$b_1$	$b_2$
$a_0$	25	1	0
$a_1$	10	30	10
$a_2$	0	33	40

• Cooperative *k*-armed bandit













- Random value from a distribution
- ullet Expected value  $\mu$

 $Arm_*$ 



 $Arm_1$ 



 $Arm_2$ 





- Agent A: teacher
  - Knows payoff distributions
  - Objective: maximize expected sum of payoffs



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- Agent B: learner
  - Can only pull Arm<sub>1</sub> or Arm<sub>2</sub>



- Agent A: teacher
  - Knows payoff distributions
  - Objective: maximize expected sum of payoffs
  - If alone, always Arm<sub>\*</sub>
- Agent B: learner
  - Can only pull Arm<sub>1</sub> or Arm<sub>2</sub>
  - Selects arm with highest observed sample average

 $Arm_*$ 



 $Arm_1$ 



 $Arm_2$ 





- Alternate actions (teacher first)
- Results of all actions fully observable (to both)



- Alternate actions (teacher first)
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- Number of rounds remaining finite, known to teacher



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Objective: maximize expected sum of payoffs

- $\mu_i$ : expected payoff of Arm<sub>i</sub> ( $i \in \{1, 2, *\}$ )
  - Assume  $\mu_* > \mu_1 > \mu_2$ : only interesting case

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- r: number of rounds left

Which arm should the teacher pull, as a function of r and all the  $\mu_i$ ,  $n_i$ , and  $\bar{x}_i$ ?

$$\mu_* = 10.0$$



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 3$$

 $\mu_* = 10.0$ 



 $\mu_1 = 9.0$ 



 $\mu_2 = 5.0$ 



r = 3

i	$m_i$	$n_i$	$\bar{x_i}$
1	0.0	0	
2	0.0	0	

9.8

 $\mu_* = 10.0$ 



 $\mu_1 = 9.0$ 



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r = 3

i	$m_i$	$n_i$	$\bar{x_i}$
1	0.0	0	
2	0.0	0	

9.8 7.0

$$\mu_* = 10.0$$



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 2$$

i	$m_i$	$n_i$	$\bar{x_i}$
1	0.0	0	
2	7.0	1	7.0

 $\mu_* = 10.0$ 



 $\mu_1 = 9.0$ 



 $\mu_2 = 5.0$ 



r = 2

i	$m_i$	$n_i$	$\bar{x_i}$
1	0.0	0	
2	7.0	1	7.0

10.3

 $\mu_* = 10.0$ 



 $\mu_1 = 9.0$ 



 $\mu_2 = 5.0$ 



r = 2

i	$m_i$	$n_i$	$\bar{x_i}$
1	0.0	0	
2	7.0	1	7.0

10.3

6.0

$$\mu_* = 10.0$$



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 1$$

i	$m_i$	$n_i$	$\bar{x_i}$
1	6.0	1	6.0
2	7.0	1	7.0

 $\mu_* = 10.0$ 



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 1$$

i	$m_i$	$n_i$	$\bar{x_i}$
1	6.0	1	6.0
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• Teacher Arm<sub>1</sub> expected value:

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- Teacher Arm<sub>1</sub> expected value:
  - Define  $\eta$ : probability Arm<sub>1</sub> returns > 8
  - Assume:  $\eta > \frac{1}{2}$

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  - EV:  $\mu_1 + \eta \mu_1 + (1 \eta)\mu_2$

$$\mu_* = 10.0$$



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  - EV:  $\mu_1 + \eta \mu_1 + (1 \eta)\mu_2 > 9 + \frac{9}{2} + \frac{5}{2} = 16$

$$\mu_* = 10.0$$



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 1$$

i	$m_i$	$n_i$	$\bar{x_i}$
1	6.0	1	6.0
2	7.0	1	7.0

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- Teacher Arm<sub>\*</sub> expected value:
  - EV:  $\mu_* + \mu_2$

$$\mu_* = 10.0$$



$$\mu_1 = 9.0$$



$$\mu_2 = 5.0$$



$$r = 1$$

i	$m_i$	$n_i$	$\bar{x_i}$
1	6.0	1	6.0
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  - Assume:  $\eta > \frac{1}{2}$
  - EV:  $\mu_1 + \eta \mu_1 + (1 \eta)\mu_2 > 9 + \frac{9}{2} + \frac{5}{2} = 16$
- Teacher Arm<sub>\*</sub> expected value:
  - EV:  $\mu_* + \mu_2 = 15$

•  $\bar{x_1} < \bar{x_2} \Longrightarrow \mathbf{no}$ 

- $\bar{x_1} < \bar{x_2} \Longrightarrow \mathbf{no}$ 
  - Sequence of values from Arm<sub>2</sub>:  $u_0, u_1, u_2, \dots$

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  - Also possible:  $\mu_*$

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  - Also possible:  $\mu_*, u_0, \mu_*$

- $\bar{x_1} < \bar{x_2} \Longrightarrow \mathbf{no}$ 
  - Sequence of values from Arm<sub>2</sub>:  $u_0, u_1, u_2, \dots$
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- ullet  $ar{x_1} < ar{x_2} \Longrightarrow \mathbf{no}$ 
  - Sequence of values from Arm<sub>2</sub>:  $u_0, u_1, u_2, \dots$
  - Optimal from Arm<sub>2</sub>:  $u_0, a, b, c, d, e, \dots w, x, y, z$
  - Also possible:  $\mu_*, u_0, \mu_*, a, b, c, d, e, \dots, w, x$
- $\bar{x_1} > \bar{x_2} \Longrightarrow$  ?

- ullet  $ar{x_1} < ar{x_2} \Longrightarrow \mathbf{no}$ 
  - Sequence of values from Arm<sub>2</sub>:  $u_0, u_1, u_2, \dots$
  - Optimal from Arm<sub>2</sub>:  $u_0, a, b, c, d, e, \dots w, x, y, z$
  - Also possible:  $\mu_*, u_0, \mu_*, a, b, c, d, e, \dots, w, x$
- $\bar{x_1} > \bar{x_2} \Longrightarrow$  ?
  - Subtle, but still no

- ullet  $ar{x_1} < ar{x_2} \Longrightarrow \mathbf{no}$ 
  - Sequence of values from Arm<sub>2</sub>:  $u_0, u_1, u_2, \dots$
  - Optimal from Arm<sub>2</sub>:  $u_0, a, b, c, d, e, \dots w, x, y, z$
  - Also possible:  $\mu_*, u_0, \mu_*, a, b, c, d, e, \dots, w, x$
- $\bar{x_1} > \bar{x_2} \Longrightarrow$  ?
  - Subtle, but still **no**
  - Challenge: prove it!

• Same proof

- Same proof
  - Sequence of values from Arm<sub>1</sub>:  $v_0, v_1, v_2, \dots$
  - Optimal from Arm<sub>1</sub>:  $v_0, a, b, c, d, e, \dots w, x, y, z$
  - Also possible:  $\mu_*, v_0, \mu_*, a, b, c, d, e, \dots, w, x$

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  - Sequence of values from Arm<sub>1</sub>:  $v_0, v_1, v_2, \dots$
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- ullet Only need to consider Arm $_1$  when  $ar{x_1} < ar{x_2}$ 
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- ullet Only need to consider Arm $_1$  when  $ar{x_1} < ar{x_2}$ 
  - Depends on distributions
  - Consider binary and normal

 $\longrightarrow \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1-p \end{cases}$ 



```
 \longrightarrow \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1-p \\ \mu_i = p_i & m_i = \text{number of 1's so far} \end{cases}
```







 $p_1$ 



 $p_2$ 







 $p_*$ 

>

 $p_1$ 

>

 $p_2$ 

- Consider teaching if
  - 1.  $\bar{x_1} < \bar{x_2}$







 $p_*$ 

 $p_1$ 

 $p_2$ 

1. 
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$$p_*$$

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$$p_1$$

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92

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Consider teaching if

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Teach iff conditions 1, 2, and  $p_* - p_1 < p_1(p_1 - p_2)$ 

### **Algorithm for Optimal Teacher Action**

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- ullet  $O(r^5)$  in both memory and runtime

#### **Arms with Normal Distributions**

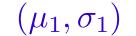


$$\longrightarrow$$
  $N(\mu, \sigma)$ 











 $(\mu_2,\sigma_2)$ 





 $(\mu_*, \sigma_*)$   $(\mu_1, \sigma_1)$ 



 $(\mu_2,\sigma_2)$ 

• Cost of teaching:  $\mu_* - \mu_1$ 



$$(\mu_*, \sigma_*)$$



 $(\mu_*, \sigma_*)$   $(\mu_1, \sigma_1)$ 



 $(\mu_2,\sigma_2)$ 

- Cost of teaching:  $\mu_* \mu_1$
- Benefit of teaching if successful:  $\mu_1 \mu_2$





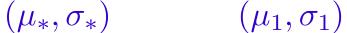
 $(\mu_*, \sigma_*)$   $(\mu_1, \sigma_1)$ 



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- Cost of teaching:  $\mu_* \mu_1$
- Benefit of teaching if successful:  $\mu_1 \mu_2$  $(\bar{x_1} < \bar{x_2})$







$$(\mu_1,\sigma_1)$$



 $(\mu_2,\sigma_2)$ 

- Cost of teaching:  $\mu_* \mu_1$
- Benefit of teaching if successful:  $\mu_1 \mu_2$  $(\bar{x_1} < \bar{x_2})$
- Probability it's successful:  $1 \Phi_{\mu_1,\sigma_1}(\bar{x_2}(n_1+1) \bar{x_1}n_1)$ 
  - Cumulative probability that pulling Arm<sub>1</sub> causes  $\bar{x_1} > \bar{x_2}$



$$(\mu_*,\sigma_*)$$



$$(\mu_*, \sigma_*) \qquad (\mu_1, \sigma_1)$$



 $(\mu_2,\sigma_2)$ 

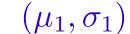
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Teach iff 
$$1 - \Phi_{\mu_1,\sigma_1}(\bar{x_2}(n_1+1) - \bar{x_1}n_1) > \frac{\mu_* - \mu_1}{\mu_1 - \mu_2}$$











 $(\mu_2,\sigma_2)$ 

- Can solve computationally nested integral
- Not exactly, nor efficiently







 $(\mu_1,\sigma_1)$ 



 $(\mu_2,\sigma_2)$ 

- Can solve computationally nested integral
- Not exactly, nor efficiently
- Can you find an efficient algorithm?

Evaluating teacher heuristics

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  - False! (binary and normal)

















 $\mathsf{Arm}_2$ 





 $Arm_3 \cdots Arm_z$ 





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Arm<sub>\*</sub>



 $Arm_1$ 



 $Arm_2$ 



 $Arm_3$ 



 $\cdots$  Arm<sub>z</sub>



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 $Arm_1$ 



 $Arm_2$ 



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 $\mathsf{Arm}_1$ 



 $Arm_2$ 



 $Arm_3$ 



 $Arm_z$ 



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Arm<sub>\*</sub>



 $\mathsf{Arm}_1$ 



 $Arm_2$ 



 $Arm_3$ 



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  - **Surprising:** May be best to teach with  ${\sf Arm}_j$  for j>i (teach with  ${\sf Arm}_2$ , even though  $\bar{x}_1>\bar{x}_2>\bar{x}_3$ )

What if the teacher doesn't know the distributions?

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  - Exploration vs. exploitation

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- What if there are multiple learners?

- Ad hoc team player is an individual
  - Unknown teammates (programmed by others)
- May or may not be able to communicate
- Teammates likely sub-optimal: no control





Goal: Create a good team player

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Goal: Create a good team player

- So far: Minimal representative scenarios
- Future: Unknown teammate behavior, communication, incomplete teacher knowledge,... much more!

# Acknowledgements

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