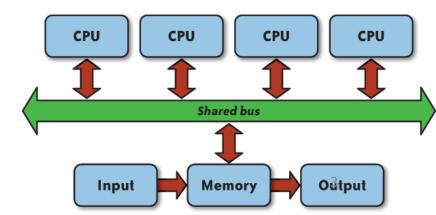
Parallel Systems Events and Futures

Chris Rossbach + Calvin Lin

CS380p

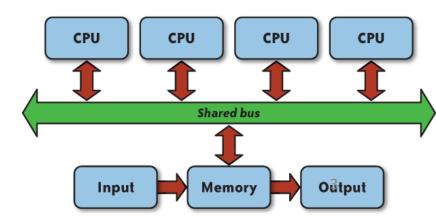
Outline for Today

- Asynchronous Programming Models
 - Events
 - Futures

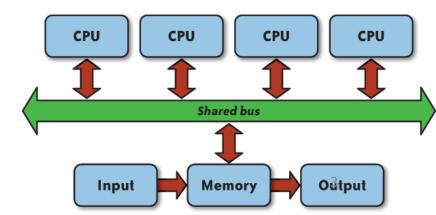


• Concrete model:

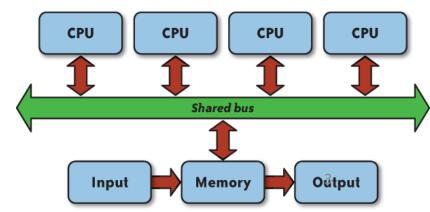
• CPU(s) execute instructions sequentially



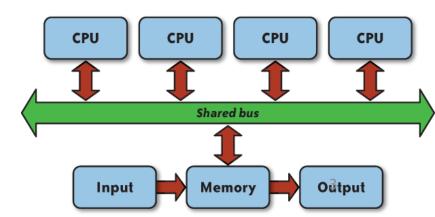
- Concrete model:
 - CPU(s) execute instructions sequentially
- Dimensions:
 - How to specify computation
 - How to specify communication
 - How to specify coordination/control transfer



- Concrete model:
 - CPU(s) execute instructions sequentially
- Dimensions:
 - How to specify computation
 - How to specify communication
 - How to specify coordination/control transfer
- Techniques/primitives
 - Threads/Processes
 - Message passing vs shared memory
 - Preemption vs Non-preemption



- Concrete model:
 - CPU(s) execute instructions sequentially
- Dimensions:
 - How to specify computation
 - How to specify communication
 - How to specify coordination/control transfer
- Techniques/primitives
 - Threads/Processes
 - Message passing vs shared memory
 - Preemption vs Non-preemption
- Dimensions/techniques not always orthogonal



"Task" == "Flow of Control" "Stack" == Task State

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- Preemptive
 - Interleave on uniprocessor
 - Overlap on multiprocessor

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 - One at a time, no conflict

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- Cooperative
 - Yields at well-defined points
 - E.g. wait for long-running I/O

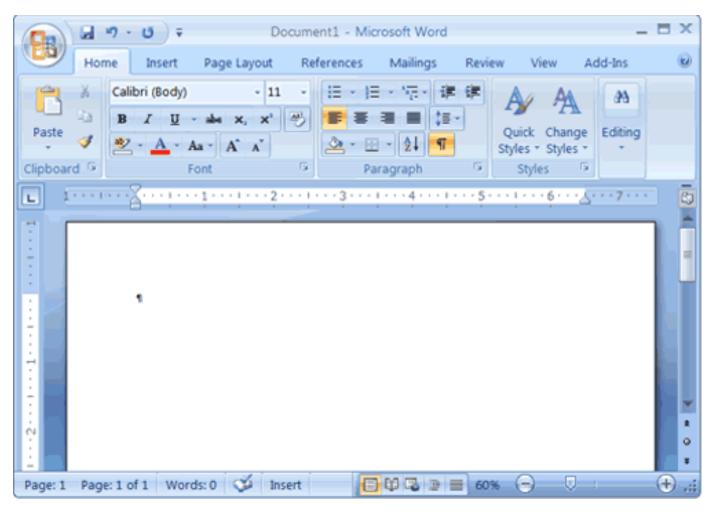
"Task" == "Flow of Control" "Stack" == Task State

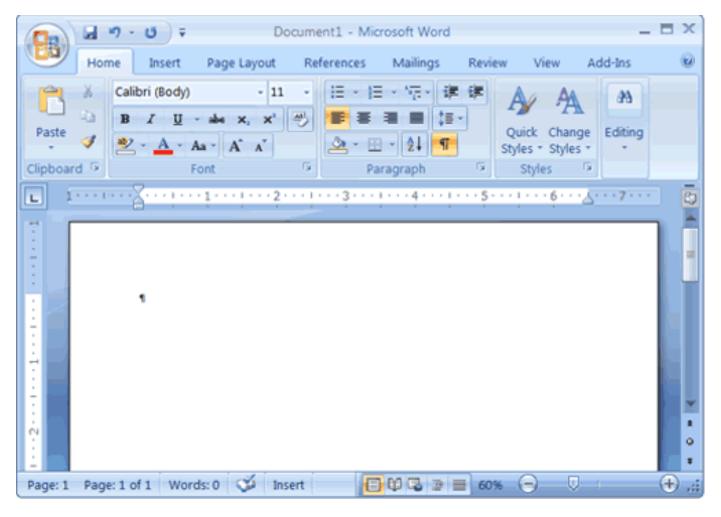
Task Management

- Preemptive
 - Interleave on uniprocessor
 - Overlap on multiprocessor
- Serial
 - One at a time, no conflict
- Cooperative
 - Yields at well-defined points
 - E.g. wait for long-running I/O

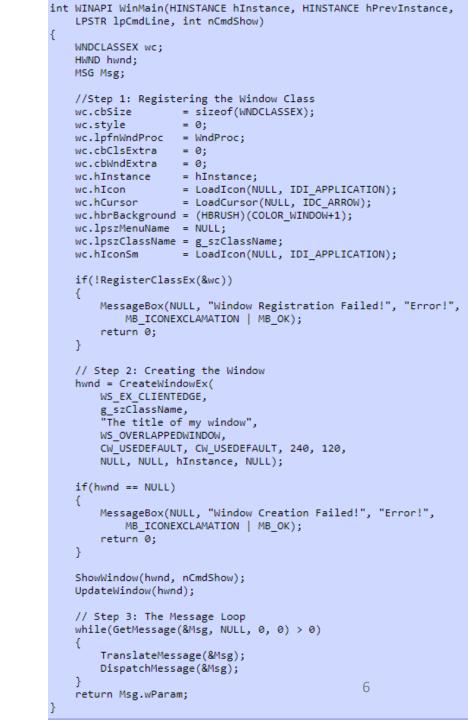
Stack Management

- Manual
 - Inherent in Cooperative
 - Changing at quiescent points
- Automatic
 - Inherent in pre-emptive
 - Downside: Hidden concurrency assumptions





do { WaitForSomething(); RespondToThing(); until(forever);





// Step 2: Creating the Window
hwnd = CreateWindowEx(
 WS_EX_CLIENTEDGE,
 g_szClassName,
 "The title of my window",
 WS_OVERLAPPEDWINDOW,
 CW_USEDEFAULT, CW_USEDEFAULT, 240, 120,
 NULL, NULL, hInstance, NULL);

int WINAPI WinMain(HINSTANCE hInstance, HINSTANCE hPrevInstance, LPSTR lpCmdLine, int nCmdShow) WNDCLASSEX wc; HWND hwnd; MSG Msg; //Step 1: Registering the Window Class wc.cbSize = sizeof(WNDCLASSEX); wc.style = 0; wc.lpfnWndProc = WndProc; wc.cbClsExtra = 0: wc.cbWndExtra = 0; wc.hInstance = hInstance; = LoadIcon(NULL, IDI APPLICATION); wc.hIcon = LoadCursor(NULL, IDC ARROW); wc.hCursor wc.hbrBackground = (HBRUSH)(COLOR WINDOW+1); wc.lpszMenuName = NULL; wc.lpszClassName = g szClassName; wc.hIconSm = LoadIcon(NULL, IDI_APPLICATION); if(!RegisterClassEx(&wc)) MessageBox(NULL, "Window Registration Failed!", "Error!", MB_ICONEXCLAMATION | MB_OK); return 0; Step 2: Creating the Window hwnd = CreateWindowEx(WS_EX_CLIENTEDGE, g szClassName, "The title of my window", WS OVERLAPPEDWINDOW, CW USEDEFAULT, CW USEDEFAULT, 240, 120, NULL, NULL, hInstance, NULL); if(hwnd == NULL) MessageBox(NULL, "Window Creation Failed!", "Error!", MB_ICONEXCLAMATION | MB_OK); return 0; ShowWindow(hwnd, nCmdShow); UpdateWindow(hwnd); // Step 3: The Message Loop while(GetMessage(&Msg, NULL, 0, 0) > 0) TranslateMessage(&Msg); DispatchMessage(&Msg);

return Msg.wParam;

6

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Events+Futures

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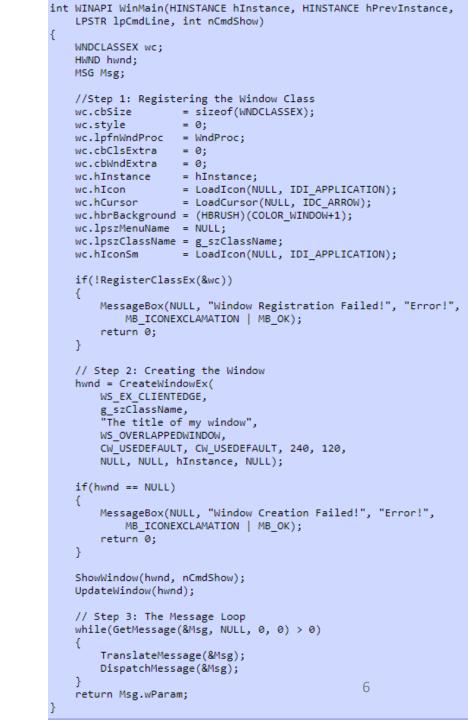
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```
switch (message)
        //case WM COMMAND:
          // handle menu selections etc.
       //break;
       //case WM PAINT:
          // draw our window - note: you must paint something here or not trap it!
       //break;
        case WM DESTROY:
             PostQuitMessage(0);
        break;
        default:
            // We do not want to handle this message so pass back to Windows
            // to handle it in a default way
            return DefWindowProc(hWnd, message, wParam, lParam);
   CS380P
                                                           Events+Futures
```

```
int WINAPI WinMain(HINSTANCE hInstance, HINSTANCE hPrevInstance,
    LPSTR lpCmdLine, int nCmdShow)
    WNDCLASSEX wc;
    HWND hwnd;
    MSG Msg;
    //Step 1: Registering the Window Class
    wc.cbSize
                     = sizeof(WNDCLASSEX);
    wc.style
                     = 0;
    wc.lpfnWndProc
                   = WndProc;
    wc.cbClsExtra
                     = 0;
    wc.cbWndExtra
                     = 0;
    wc.hInstance
                     = hInstance;
    wc.hIcon
                     = LoadIcon(NULL, IDI_APPLICATION);
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    wc.lpszMenuName = NULL;
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    if(hwnd == NULL)
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    ShowWindow(hwnd, nCmdShow);
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    while(GetMessage(&Msg, NULL, 0, 0) > 0)
       TranslateMessage(&Msg);
        DispatchMessage(&Msg);
                                                  6
    return Msg.wParam;
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int WINAPI WinMain(HINSTANCE hInstance, HINSTANCE hPrevInstance, LPSTR lpCmdLine, int nCmdShow)

Symbolic

WNDCLASSEX wc; HWND hwnd; MSG Msg;

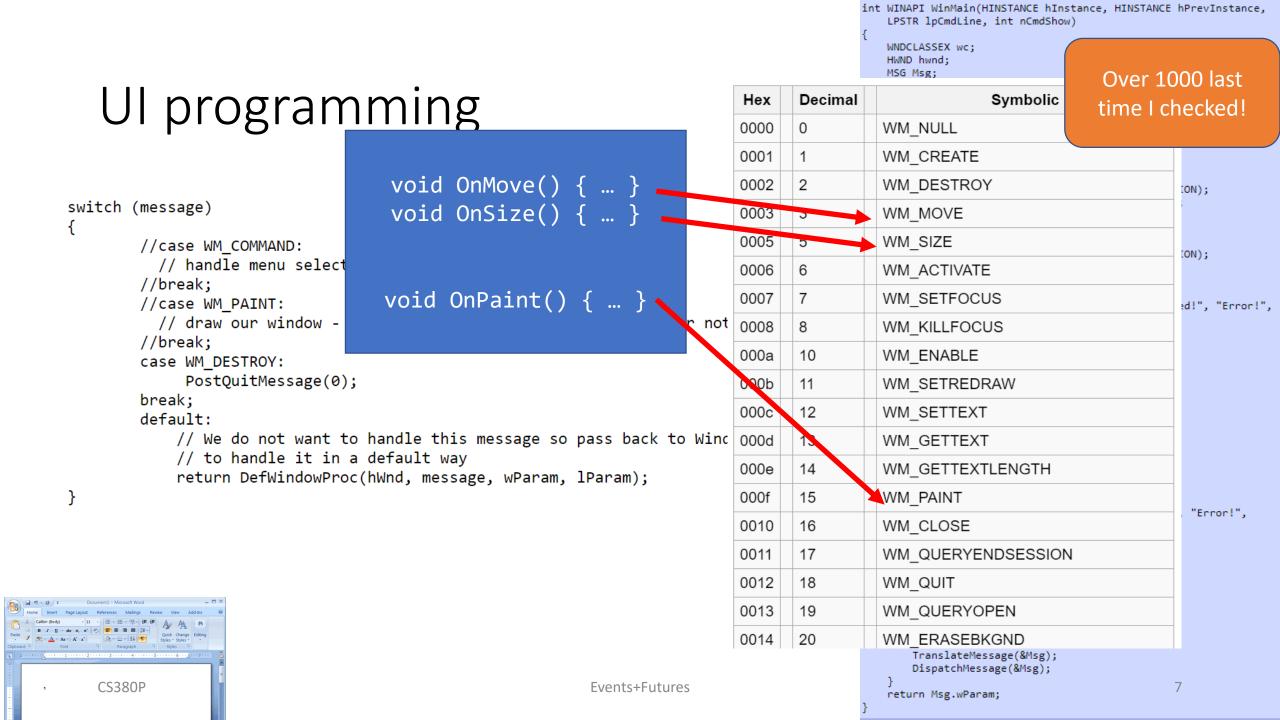
Decimal

Hex

UI programming

61 p1 60 ann 10	0000	0	WM_NULL		
	0001	1	WM_CREATE		
	0002	2	WM_DESTROY	:ON);	
switch (message) r	0003	3	WM_MOVE		
ו //case WM_COMMAND:	0005	5	WM_SIZE	(ON);	
<pre>// handle menu selections etc. // handle menu selections etc.</pre>	0006	6	WM_ACTIVATE		
//break; //case WM_PAINT:	0007	7	WM_SETFOCUS	d!", "Error!",	
<pre>// draw our window - note: you must paint something here or not // here here</pre>	0008	8	WM_KILLFOCUS		
//break; case WM_DESTROY:	000a	10	WM_ENABLE		
PostQuitMessage(0);	000b	11	WM_SETREDRAW		
break; default:	000c	12	WM_SETTEXT		
<pre>// We do not want to handle this message so pass back to Winc // to handle it in a default way return DefWindowProc(hWnd, message, wParam, lParam); }</pre>	000d	13	WM_GETTEXT		
	000e	14	WM_GETTEXTLENGTH		
	000f	15	WM_PAINT		
	0010	16	WM_CLOSE	"Error!",	
	0011	17	WM_QUERYENDSESSION		
	0012	18	WM_QUIT		
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里 - Au ×, x* (P) ■ ■ (E) - Apr / Au Apr / Au	0014	20	WM_ERASEBKGND		
CS380P Events+Futures			TranslateMessage(&Msg); DispatchMessage(&Msg); } return Msg.wParam; }	7	

		nt WINAPI WinMain(HINSTANCE hIns LPSTR lpCmdLine, int nCmdShow	INAPI WinMain(HINSTANCE hInstance, HINSTANCE hPrevInstance, PSTR lpCmdLine, int nCmdShow)					
			·	WNDCLASSEX wc; HWND hwnd; MSG Msg;	Over 1(000 last		
UI programming		Hex	Decimal	Symbolic	time I checked!			
<pre>switch (message) { //case WM_COMMAND: // handle menu selections etc. //break; //case WM_PAINT: // draw our window - note: you must paint something here or not //break; case WM_DESTROY:</pre>	Jeranning	0000	0	WM_NULL		checked.		
		0001	1	WM_CREATE		:ON);		
		0002	2	WM_DESTROY				
		0003	3	WM_MOVE				
	0005	5	WM_SIZE		[ON];			
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	0012	18	WM_QUIT					
	0013	19	WM_QUERYOPEN					
		0014	20	WM_ERASEBKGND TranslateMessage(&Msg);				
• CS380P	Events+Futures		}	DispatchMessage(&Msg); } return Msg.wParam;		7		



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          switch(message) {
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Pros

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Pros

• Simple imperative programming

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Pros

CS380P

- Simple imperative programming
- Good fit for uni-processor

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Pros

- Simple imperative programming
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Cons •Awkward/verbose

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```

Pros

- Simple imperative programming
- Good fit for uni-processor

Cons

- •Awkward/verbose
- Obscures available parallelism

```
⊟winmain(...) {
      while(true) {
3
          message = GetMessage();
          switch(message) {
5
          case WM LONGRUNNING CPU HOG: HogCPU(); break;
          case WM HIGH LATENCY IO: BlockForALongTime(); break;
6
          case WM DO QUICK IMPORTANT THING: HopeForTheBest(); break;
8
9
10
```

Pros

- Simple imperative programming
- Good fit for uni-processor

Cons

- •Awkward/verbose
- Obscures available parallelism

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 9
10
```

```
How can we
  ⊟winmain(...) {
     while(true) {
2
                                             parallelize
3
         message = GetMessage();
4
         switch(message) {
                                                this?
5
         case WM THIS: DoThis(); break;
         case WM THAT: DoThat(); break;
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         case WM OTHERTHING: DoOtherThing(); break;
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```









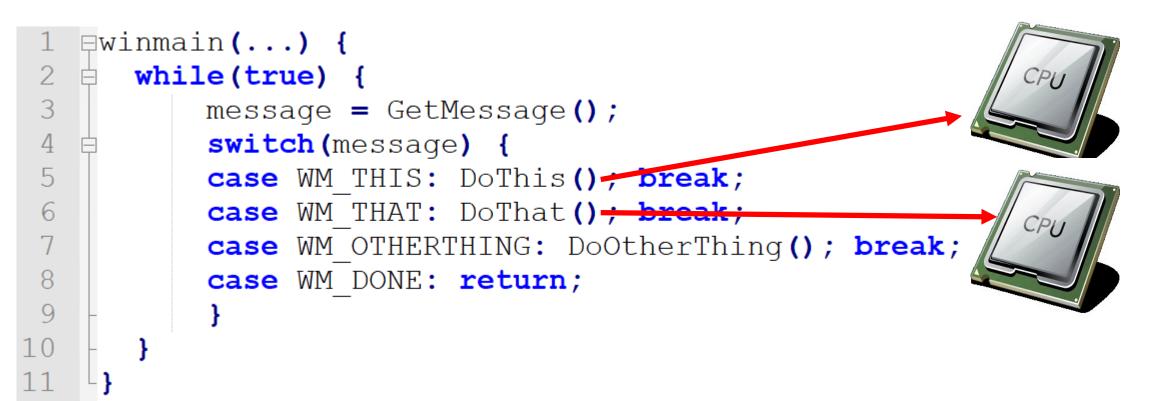
Parallel UI Implementation 1

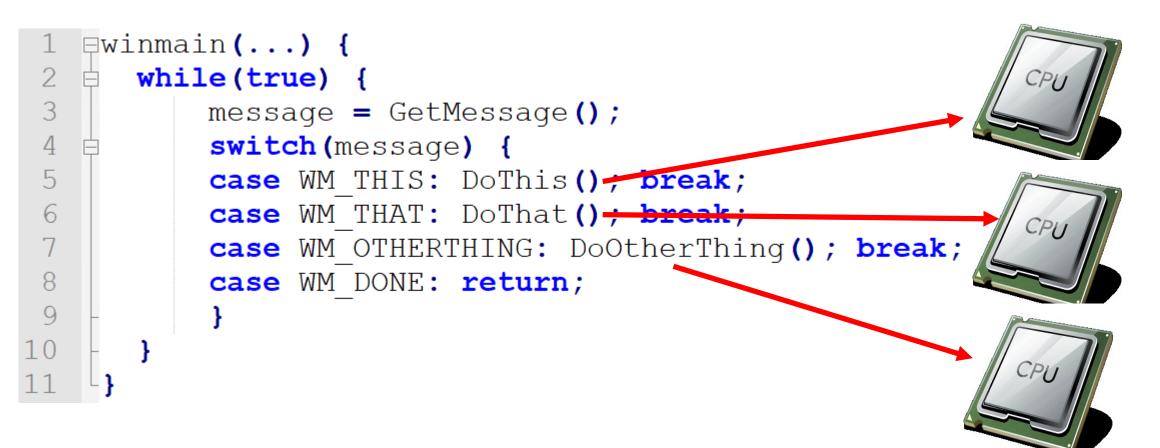
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 7
          case WM OTHERTHING: DoOtherThing(); break;
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          case WM DONE: return;
 9
           }
10
```

Parallel UI Implementation 1

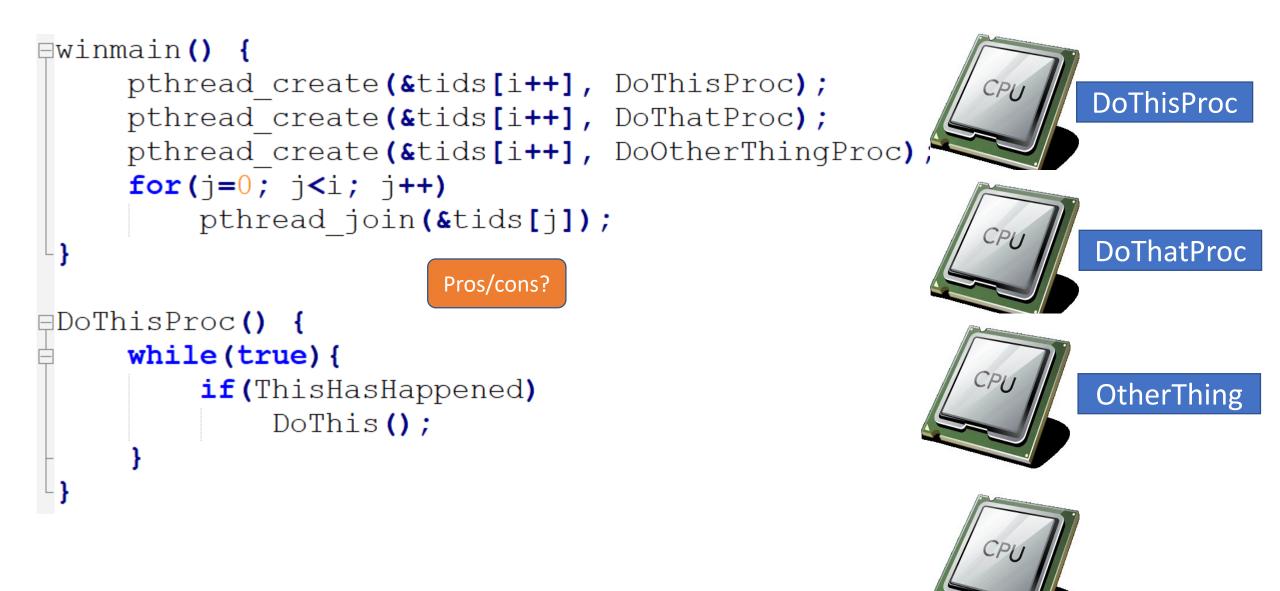
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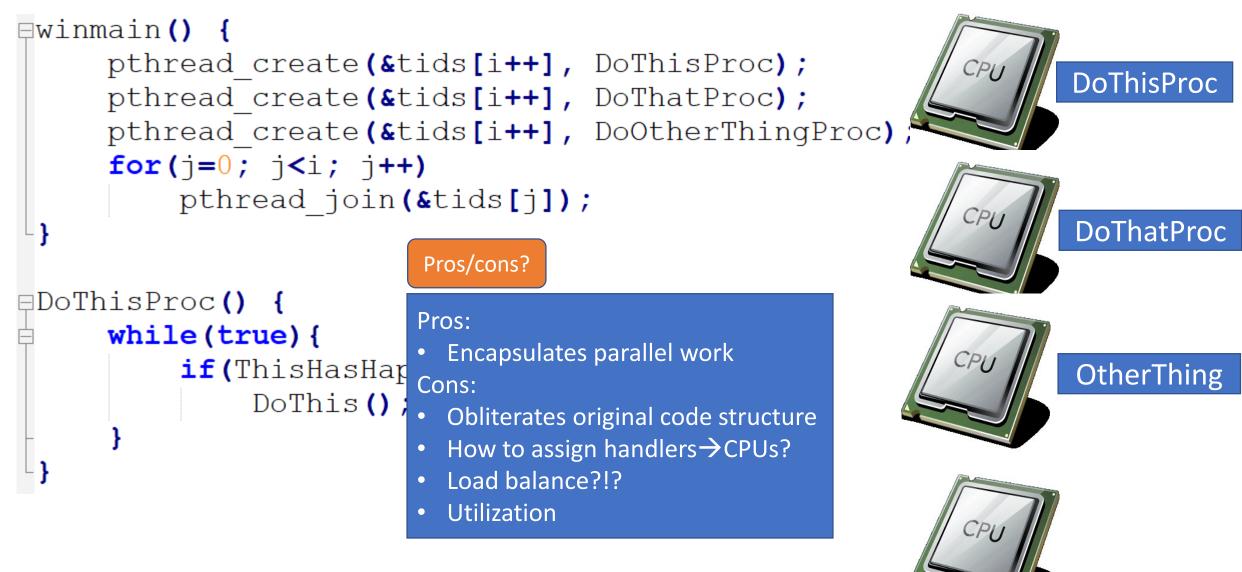
Parallel UI Implementation 1



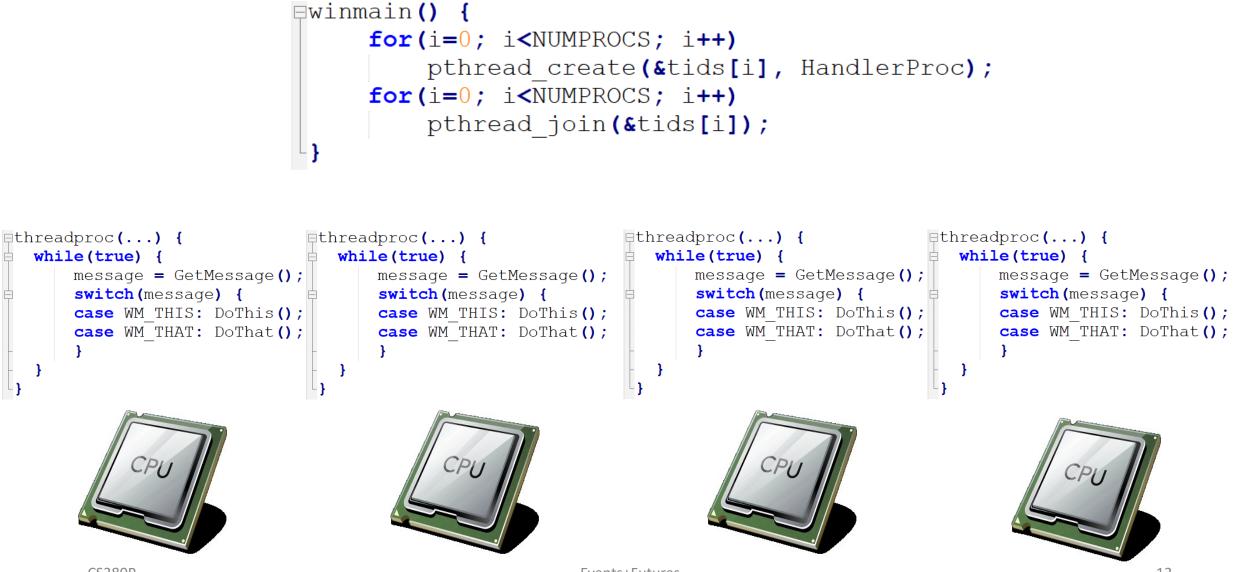


```
⊟winmain() {
     pthread create(&tids[i++], DoThisProc);
                                                               DoThisProc
     pthread create(&tids[i++], DoThatProc);
     pthread create(&tids[i++], DoOtherThingProc)
     for(j=0; j<i; j++)</pre>
         pthread join(&tids[j]);
                                                               DoThatProc
⊟DoThisProc() {
     while(true) {
         if (ThisHasHappened)
                                                               OtherThing
              DoThis();
```



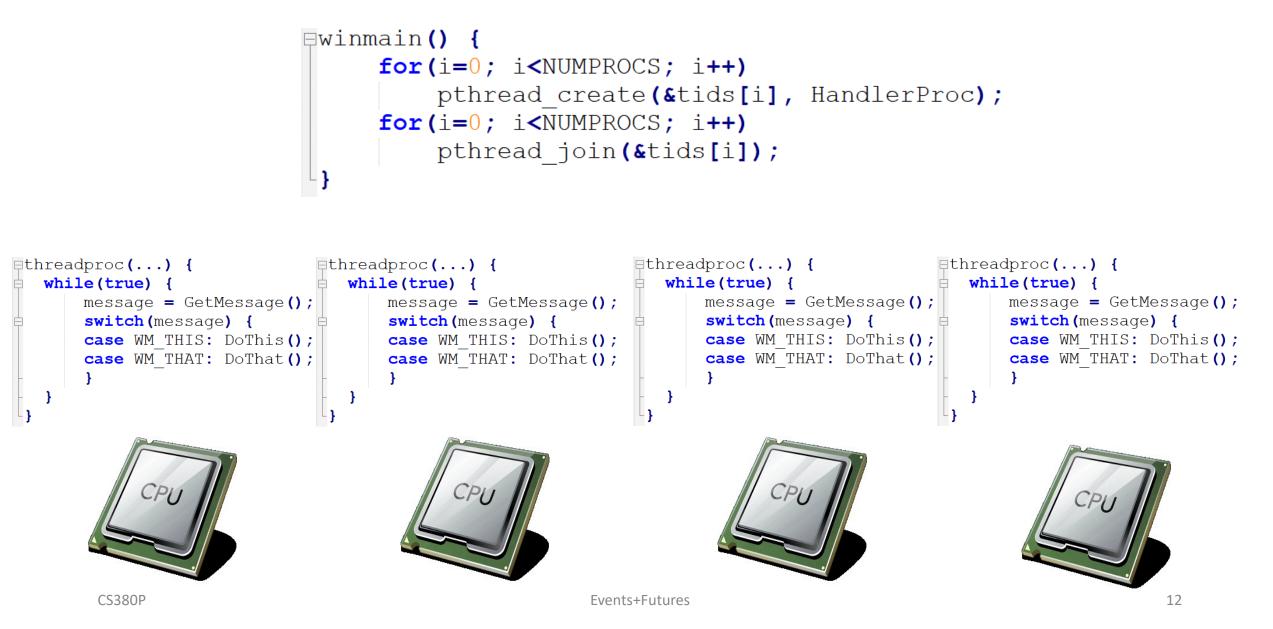


11



Events+Futures



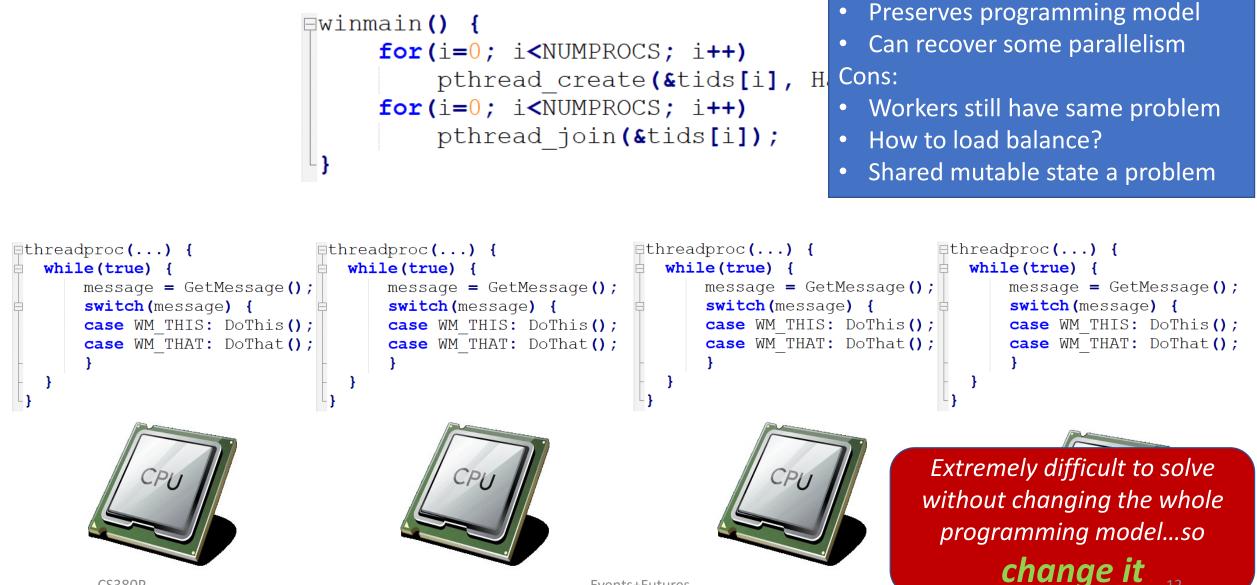


```
Preserves programming model
                          ⊟winmain() {
                                                                               Can recover some parallelism
                                  for(i=0; i<NUMPROCS; i++)</pre>
                                       pthread create (&tids[i], H. Cons:
                                  for(i=0; i<NUMPROCS; i++)</pre>
                                                                             • Workers still have same problem
                                       pthread join(&tids[i]);
                                                                               How to load balance?
                                                                             • Shared mutable state a problem
                                                          □threadproc(...) {
                                                                                       □threadproc(...) {
⊟threadproc(...) {
                            □threadproc(...) {
  while(true) {
                               while(true) {
                                                             while(true) {
                                                                                          while(true) {
      message = GetMessage();
                                  message = GetMessage();
                                                                message = GetMessage();
                                                                                             message = GetMessage();
                                                                                             switch(message) {
      switch(message) {
                                   switch(message) {
                                                                switch(message) {
                                                                case WM THIS: DoThis();
                                                                                             case WM THIS: DoThis();
      case WM THIS: DoThis();
                                   case WM THIS: DoThis();
      case WM THAT: DoThat();
                                   case WM THAT: DoThat();
                                                                case WM THAT: DoThat();
                                                                                             case WM THAT: DoThat();
       CS380P
```

Pros/cons?

Pros:

Events+Futures



Pros/cons?

Pros:

Events+Futures

Event-based Programming: Motivation

Event-based Programming: Motivation

- Threads have a *lot* of down-sides:
 - Tuning parallelism for different environments
 - Load balancing/assignment brittle
 - Shared state requires locks \rightarrow
 - Priority inversion
 - Deadlock
 - Incorrect synchronization
 - ...

Event-based Programming: Motivation

- Threads have a *lot* of down-sides:
 - Tuning parallelism for different environments
 - Load balancing/assignment brittle
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 - Incorrect synchronization
 - ...

• Events: restructure programming model to have no threads!

• Programmer *only writes events*

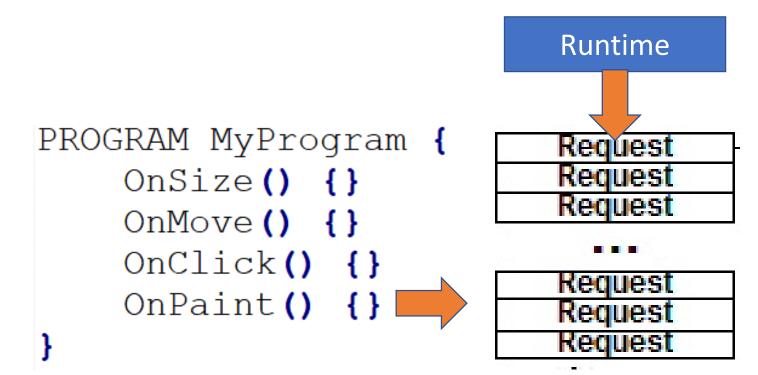
- Programmer *only writes events*
- Event: an object queued for a module (think future/promise)

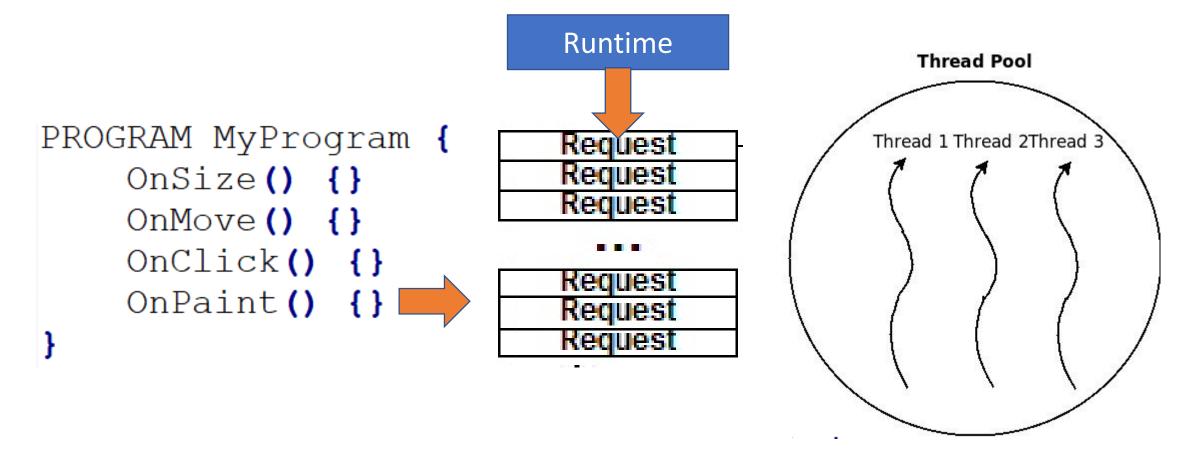
- Programmer *only writes events*
- Event: an object queued for a module (think future/promise)
- Basic primitives
 - create_event_queue(handler) → event_q
 - enqueue_event(event_q, event-object)
 - Invokes handler (eventually)

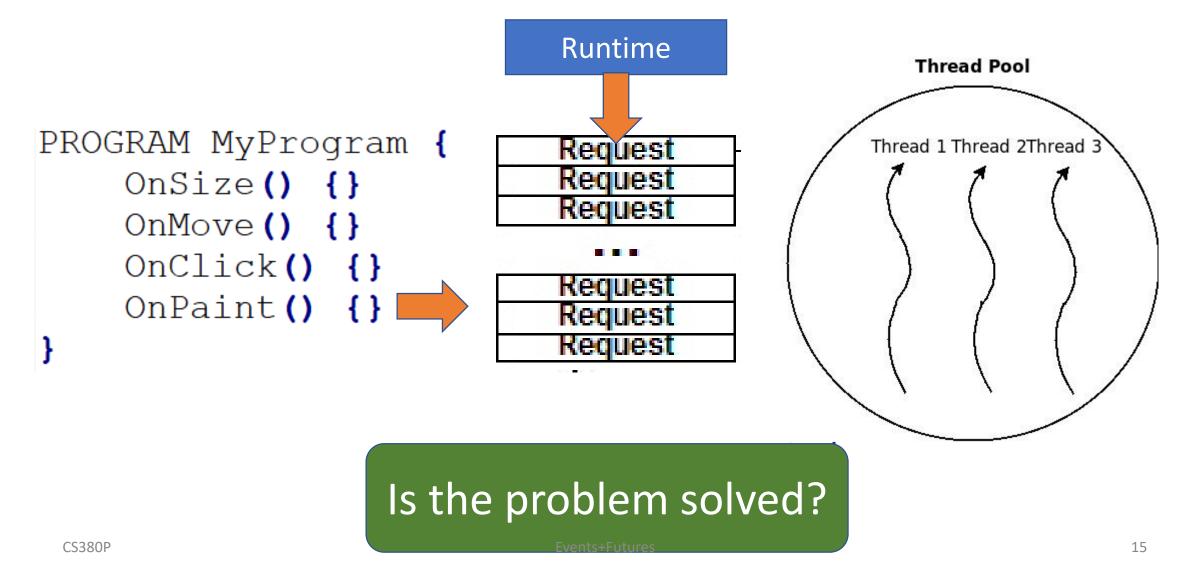
- Programmer *only writes events*
- Event: an object queued for a module (think future/promise)
- Basic primitives
 - create_event_queue(handler) → event_q
 - enqueue_event(event_q, event-object)
 - Invokes handler (eventually)
- Scheduler decides which event to execute next
 - E.g. based on priority, CPU usage, etc.

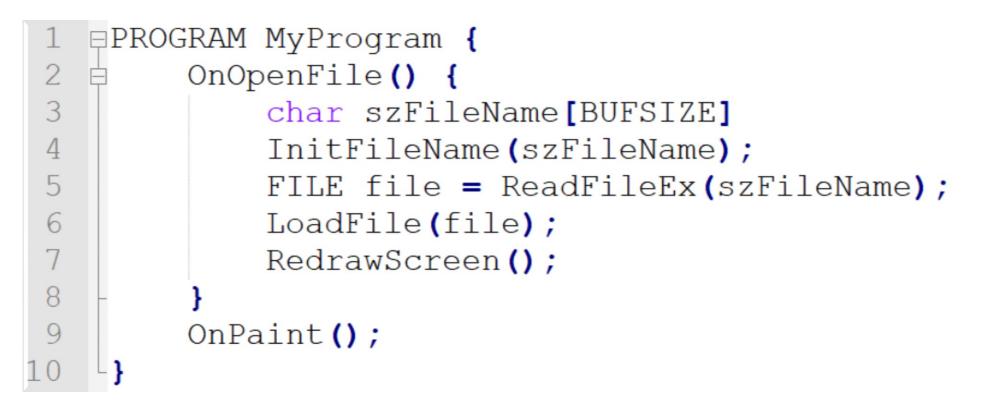
```
switch (message)
{
        //case WM COMMAND:
          // handle menu selections etc.
       //break;
        //case WM_PAINT:
         // draw our window - note: you must paint something here or not trap it!
        //break;
        case WM DESTROY:
             PostQuitMessage(0);
        break;
        default:
            // We do not want to handle this message so pass back to Windows
           // to handle it in a default way
            return DefWindowProc(hWnd, message, wParam, lParam);
}
```

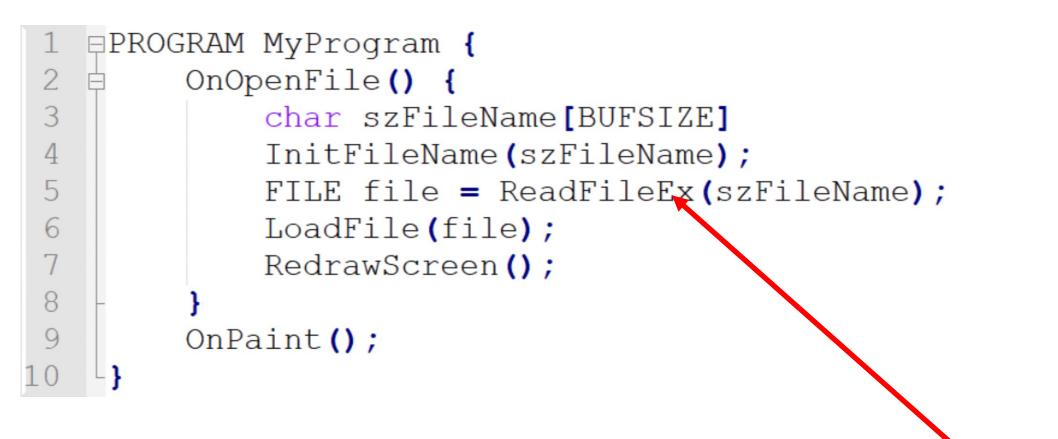
```
PROGRAM MyProgram {
    OnSize() {}
    OnMove() {}
    OnClick() {}
    OnPaint() {}
```



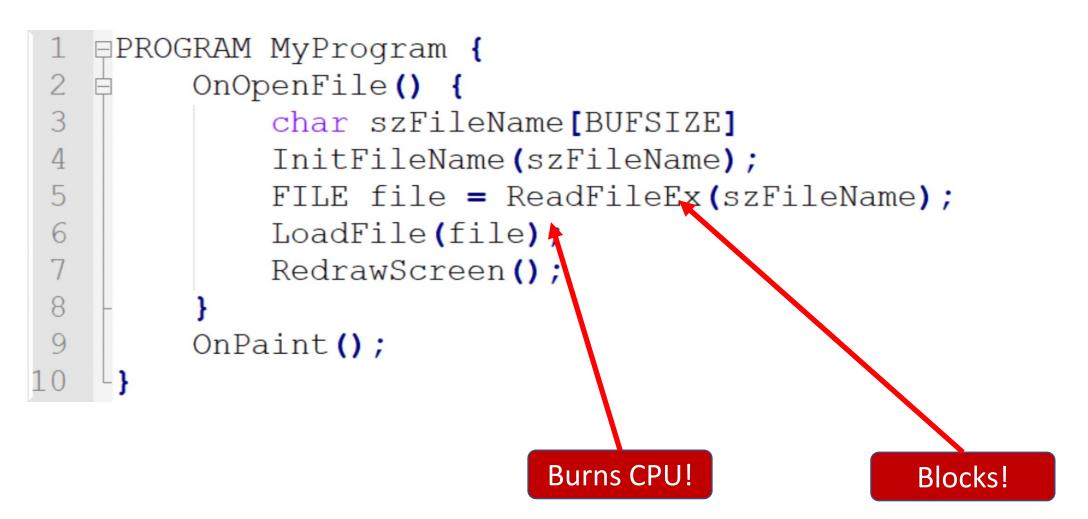


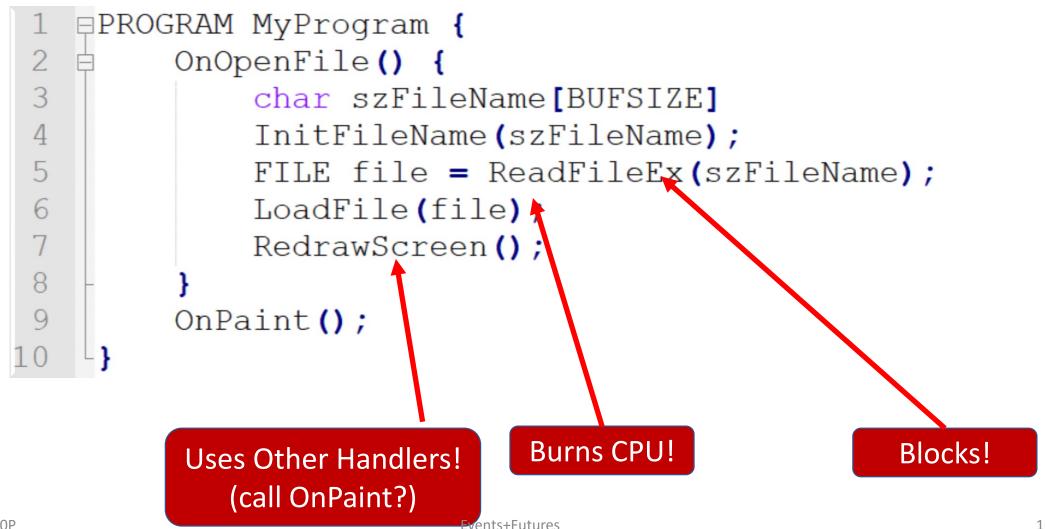






Blocks!

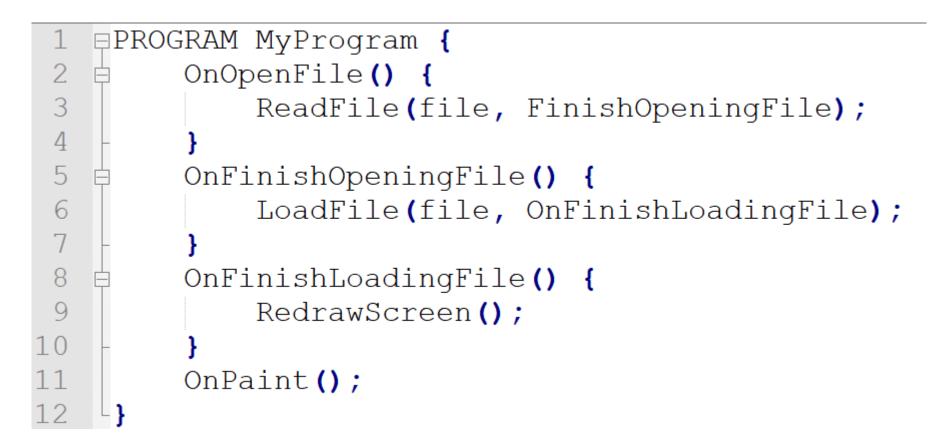


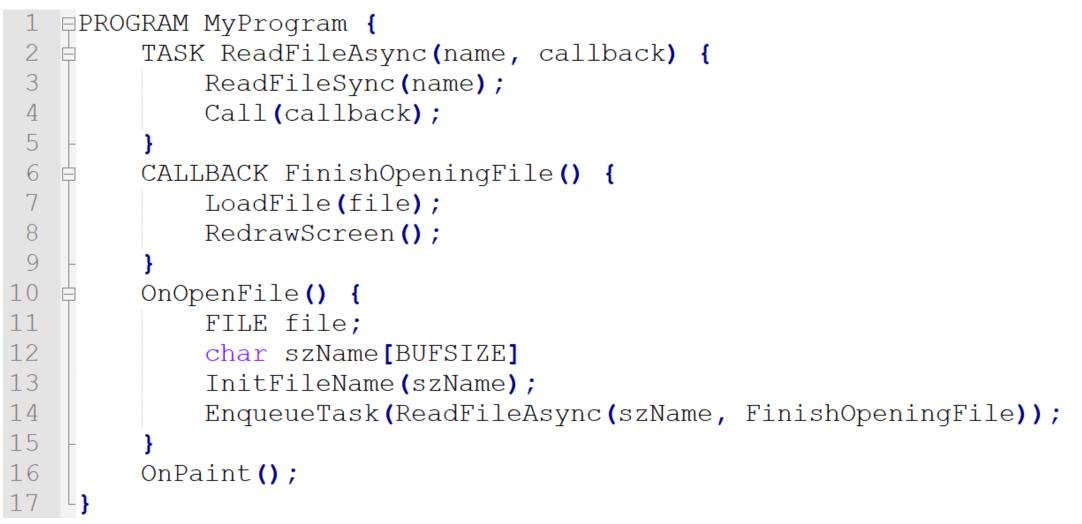


No problem! Just use more events/handlers, right?

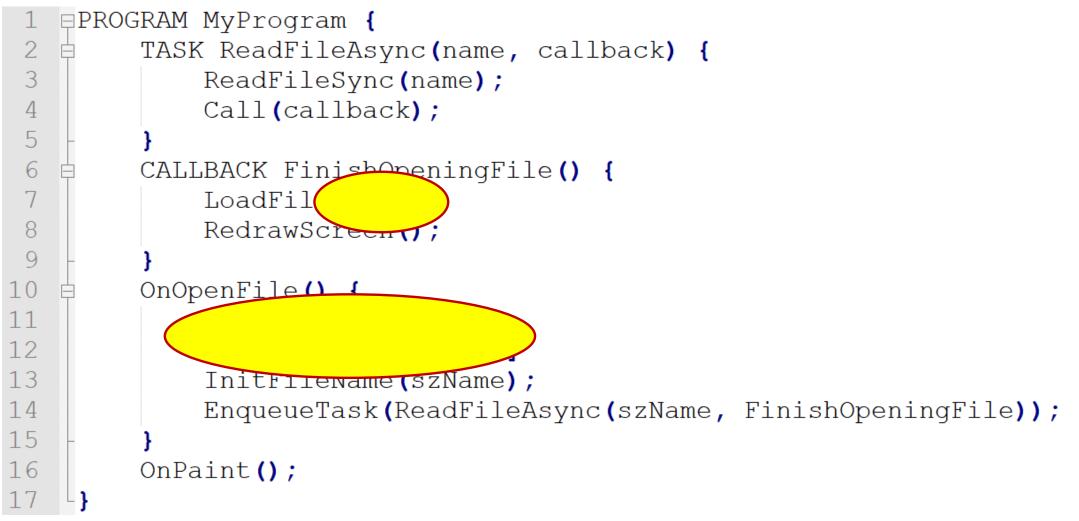
```
⊟PROGRAM MyProgram {
 2
         TASK ReadFileAsync(name, callback) {
 3
             ReadFileSync(name);
             Call(callback);
 4
 5
         CALLBACK FinishOpeningFile() {
 6
             LoadFile(file);
             RedrawScreen();
 8
 9
10
         OnOpenFile() {
11
             FILE file;
12
             char szName[BUFSIZE]
13
             InitFileName(szName);
14
             EnqueueTask(ReadFileAsync(szName, FinishOpeningFile));
15
16
         OnPaint();
                                 Events+Futures
                                                                        17
```

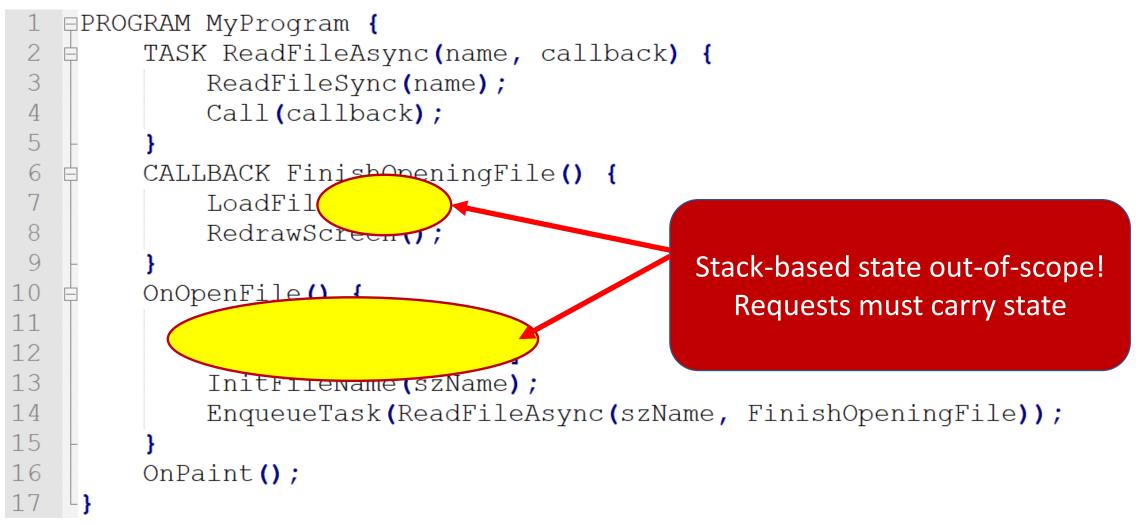
Continuations, BTW











Threads vs Events

- Thread Pros
 - Overlap I/O and computation
 - While looking sequential
 - Intermediate state on stack
 - Control flow naturally expressed
- Thread Cons
 - Synchronization required
 - Overflowable stack
 - Stack memory pressure

- Event Pros
 - Easier to create well-conditioned system
 - Easier to express dynamic change in level of parallelism

- Event Cons
 - Difficult to program
 - Control flow between callbacks obscure
 - When to deallocate memory
 - Incomplete language/tool/debugger support
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 - Computation not complete yet
- Construct (future X)
 - immediately returns value
 - concurrently executes X

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1 static void runAsyncExample() {
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 - Anonymous function
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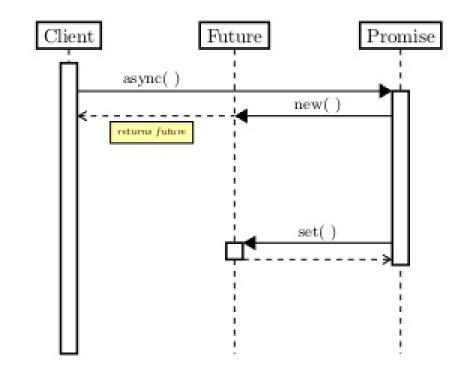
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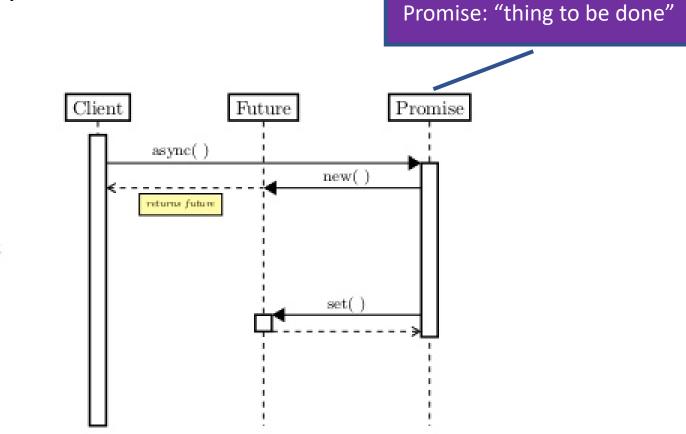
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- Where (on what thread) does the lambda expression run?

Futures and Promises: Why two kinds of objects?

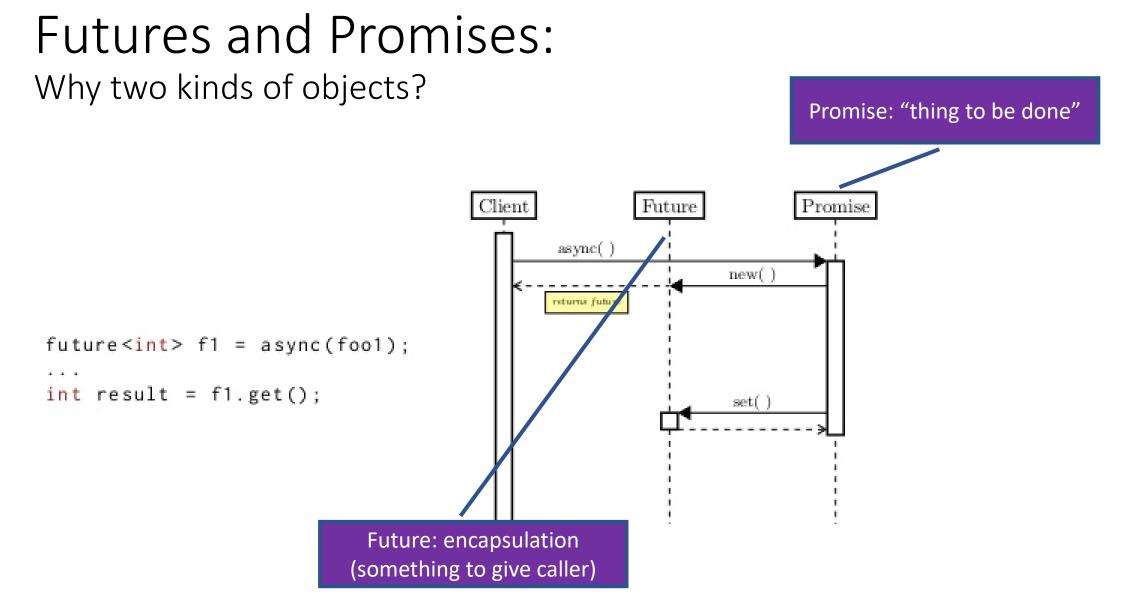


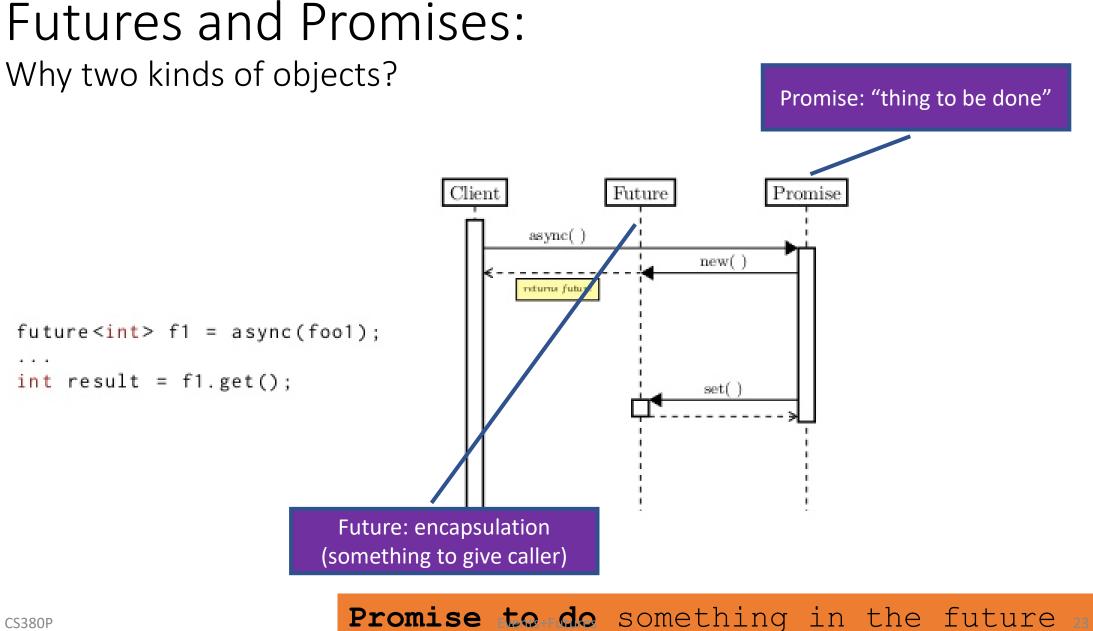
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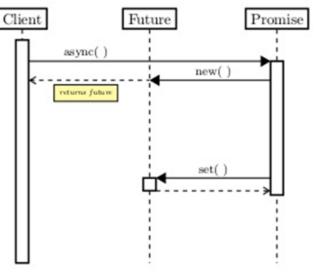


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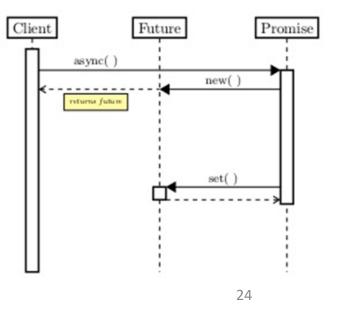


- Future: read-only reference to uncompleted value
- Promise: single-assignment variable that the future refers to
- Promises *complete* the future with:
 - Result with success/failure
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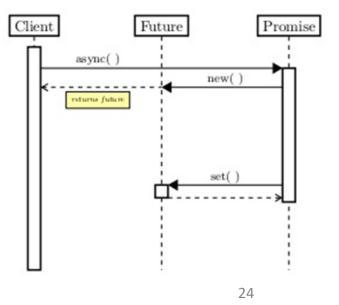
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Algol	Thunk	Address of async result
Java	CompletableFuture <t></t>	Future <t></t>
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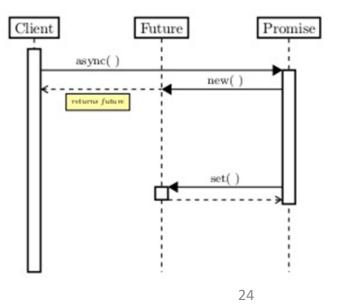


Mnemonic: Promise to *do* something Make a promise *for* the future

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 - Compiler: abstractions are *language-level objects*
 - Runtime: scheduler, task queues, thread-pools are *transparent*

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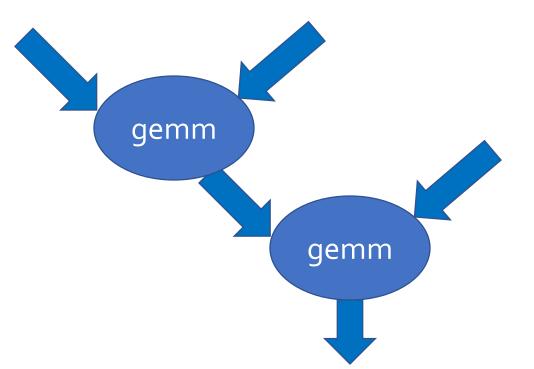
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Dataflow: a better abstraction?



- nodes \rightarrow computation
- edges \rightarrow communication
- Expresses parallelism explicitly
- Minimal specification of data movement: runtime does it.
- asynchrony is a runtime concern (not programmer concern)
- No specification of compute \rightarrow device mapping: like threads! CS380P Events+Futures