

Private Graph Algorithms in the Semi-Honest Model

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Two Party Graph Algorithms

- Parties P_1 and P_2 own graphs G_1 and G_2
- f is a *two-input* graph algorithm
- Compute $f(G_1, G_2)$ without revealing “unnecessary information”

Unnecessary Information

- Intuitively, the protocol should function *as if a trusted third party computed the output*
- We use *simulation* to prove that a protocol is private

The Semi-Honest Model

- A *malicious* adversary can alter his input
- A semi-honest adversary
 - adheres to protocol
 - tries to learn extra information from the message transcript

General Secure Two Party Computation

- *Any* polynomial sized functionality can be made private (in the semi-honest model)
 - Yao's Method
- What are our goals?
 - Yao's Method is inefficient
 - Efficient, private protocols to compute particular graph functionalities
 - Take advantage of information “leaked” by the result

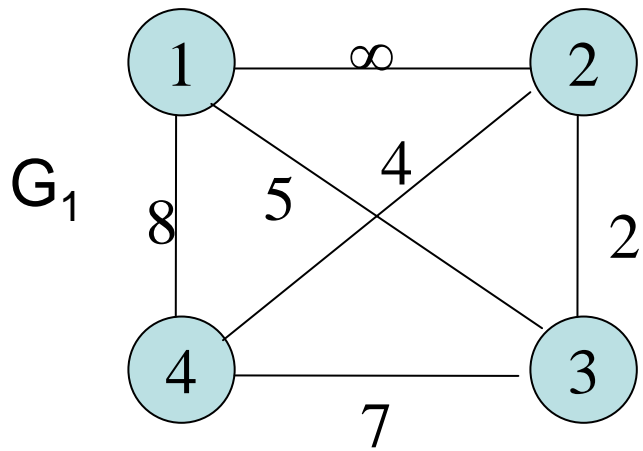
Two-Input Graph Algorithms?

- Graph Isomorphism
- Comparison of graph statistics
 - $f(G_1) > f(G_2)$?
 - max flow, diameter, average degree
- Synthesized Graphs
 - $f(G_1 \bullet G_2)$

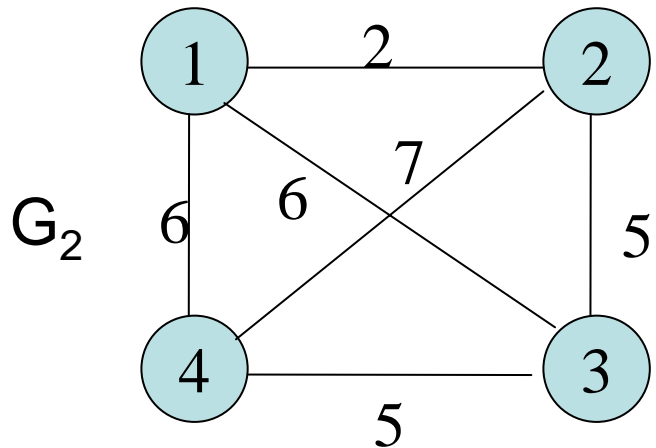
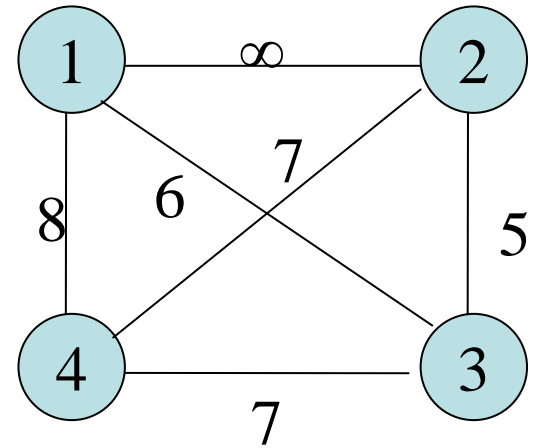
Graph Synthesis

- G_1 and G_2 are weighted complete graphs on the same vertex and edge set
 - $G_1 = (V, E, w_1)$; $G_2 = (V, E, w_2)$
- $\text{gmax}(G_1, G_2) = (V, E, w_{\max})$
 - $w_{\max}(e) = \max(w_1(e), w_2(e))$
- $\text{gmin}(G_1, G_2) = (V, E, w_{\min})$
 - $w_{\min}(e) = \min(w_1(e), w_2(e))$

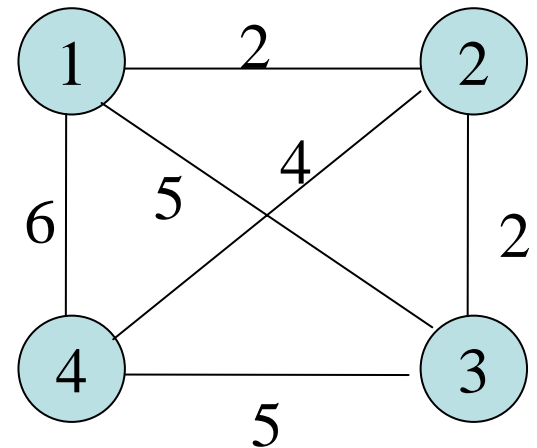
Graph Synthesis



$gmax(G_1, G_2)$



$gmin(G_1, G_2)$



Graph Isomorphism

- Unlikely to find a private protocol
 - No known poly-time algorithm

Comparison of Graph Statistics

1. Compute statistic on own graph
 - Semi-honest participants can't lie
2. Use a private comparison protocol
 - Yao's Millionaire Protocol
 - Yao's method (circuit protocol)

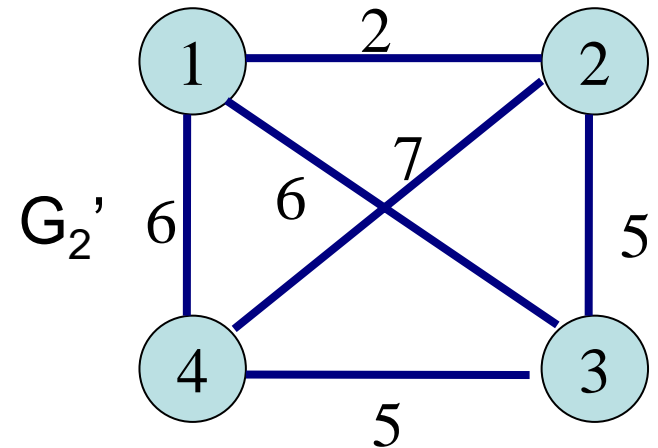
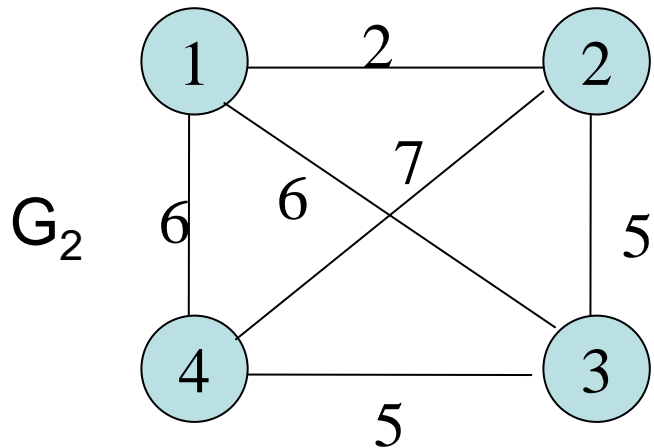
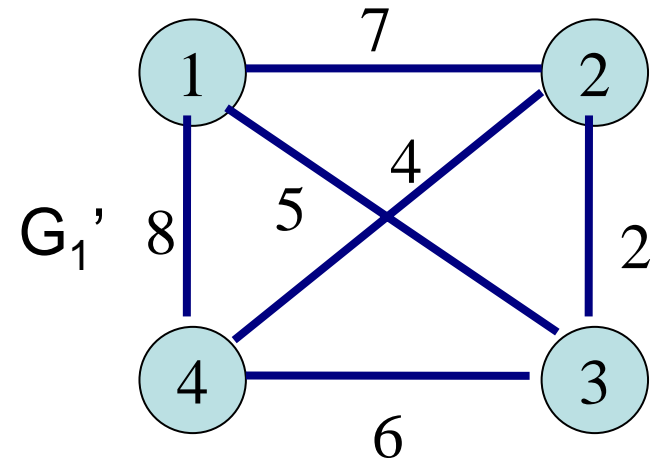
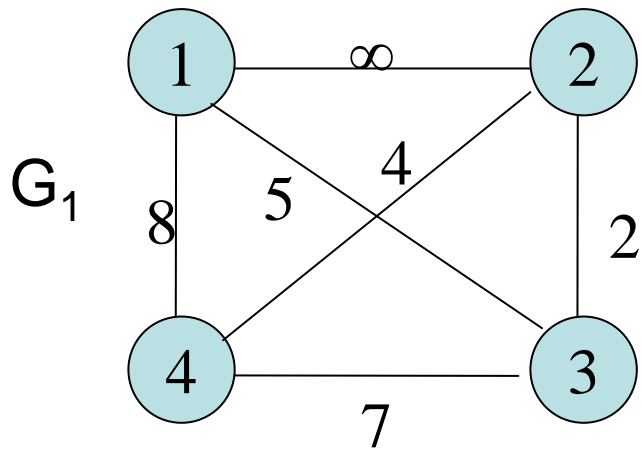
Synthesized Graphs

- This is the interesting case
- All Pairs Shortest Distance and Single Source Shortest Distance both “leak” significant useful information
 - Solved: APSD(gmin), SSSD(gmin)
 - Solved with leaks:
APSD(gmax), SSSD(gmax)

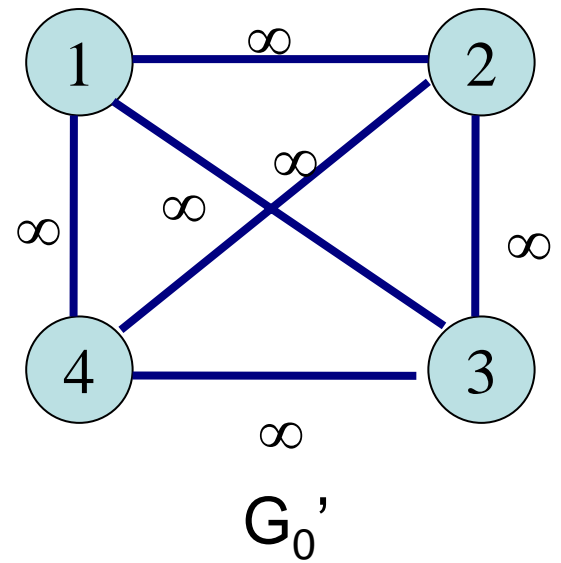
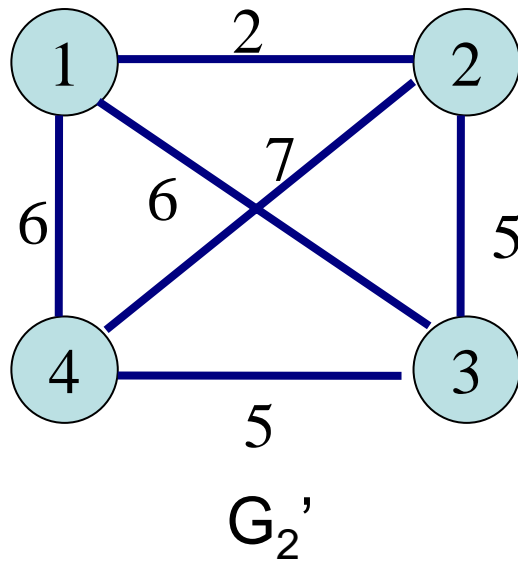
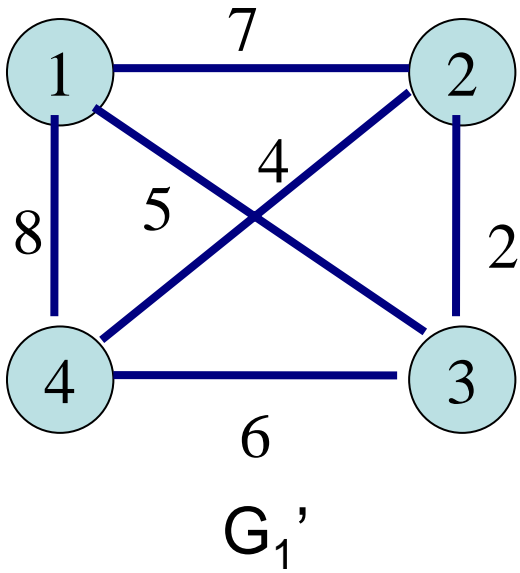
APSD($\text{gmin}(G_1, G_2)$)

- Basic Idea: Add edges to the solution graph in order of smallest to largest
- Private, because we can recover the order from the final solution graph

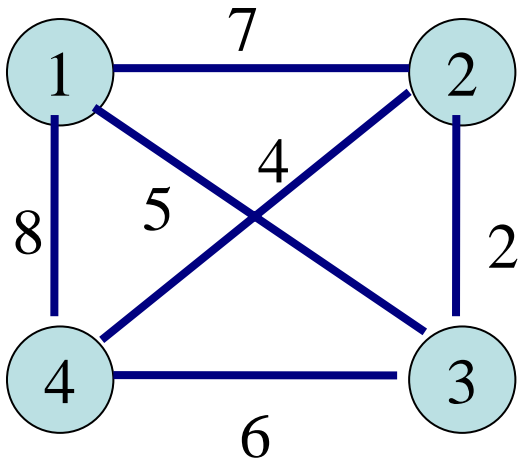
Run APSP on G_1 and G_2



Initialize G_0'

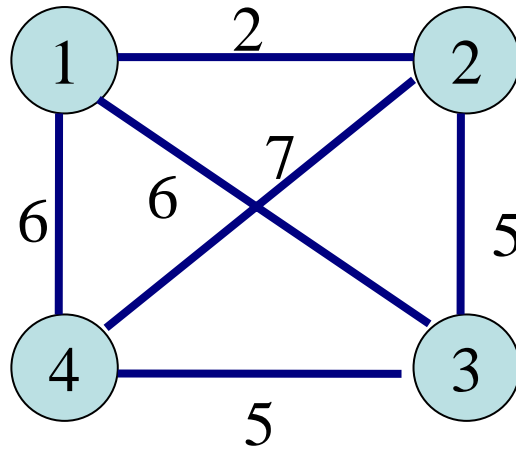


Find shortest blue edge lengths



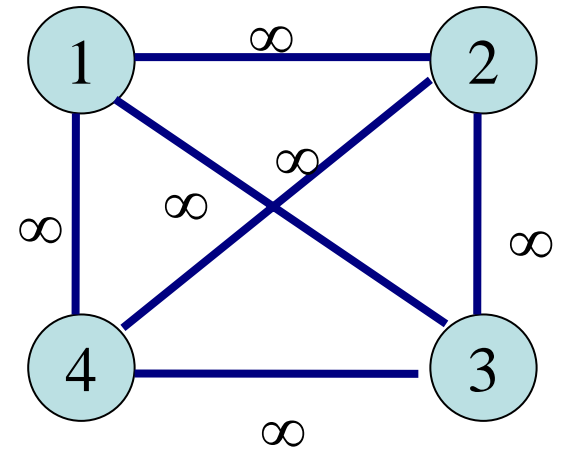
G_1'

$$\min_1 = 2$$



G_2'

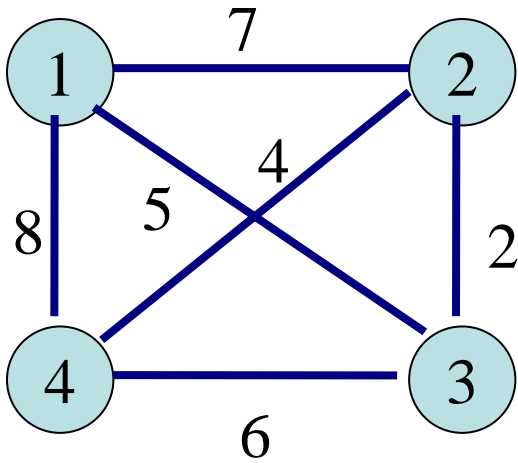
$$\min_2 = 2$$



G_0'

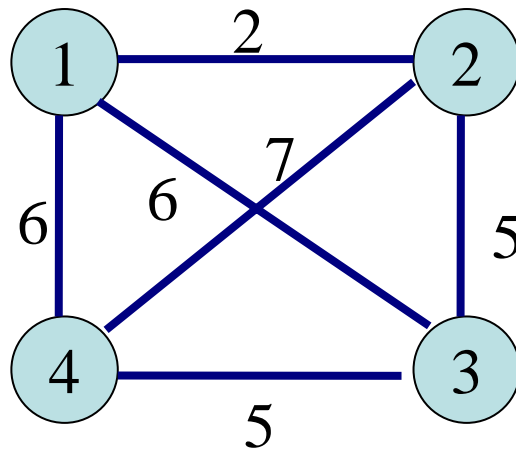
$$\min_0 = \infty$$

Privately find global shortest length



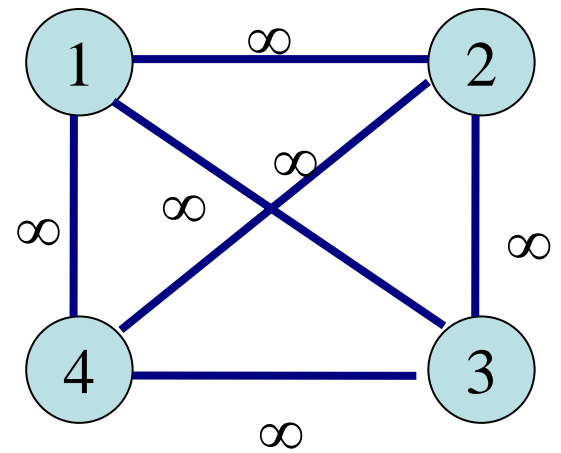
G_1'

$$\text{min1} = 2$$



G_2'

$$\text{min2} = 2$$

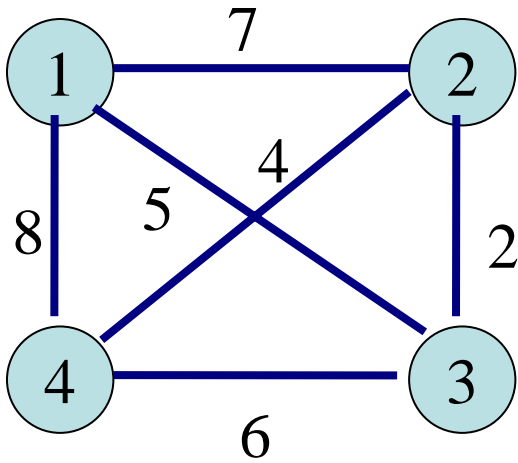


G_0'

$$\text{min0} = \infty$$

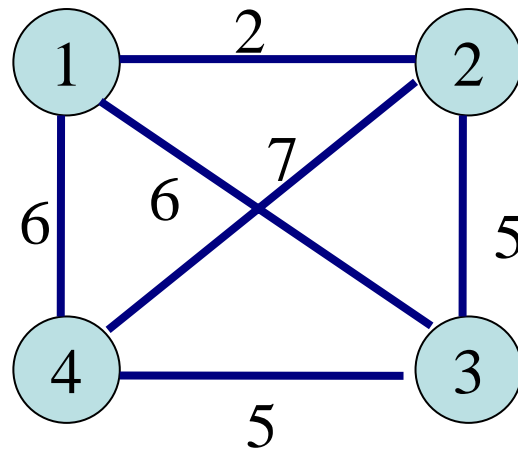
$$\text{bluemin} = \min(\text{min0}, \text{min1}, \text{min2}) = 2$$

Find edges of length blumin



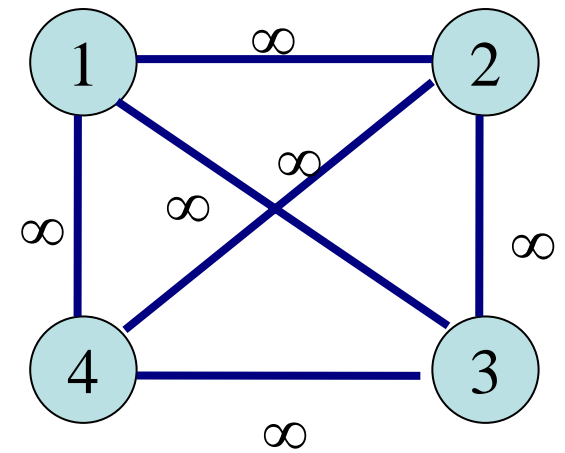
G_1'

$$S_1 = \{e_{23}\}$$



G_2'

$$S_2 = \{e_{12}\}$$

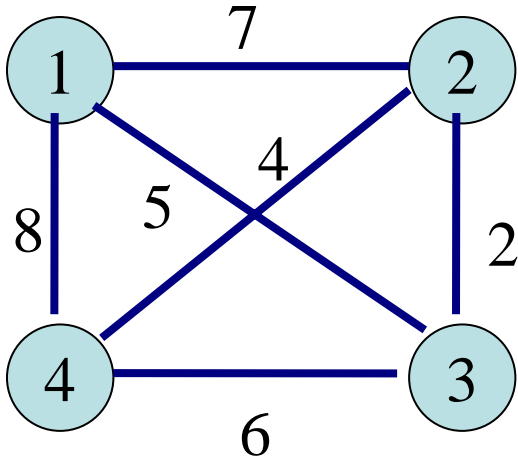


G_0'

$$S_0 = \{\}$$

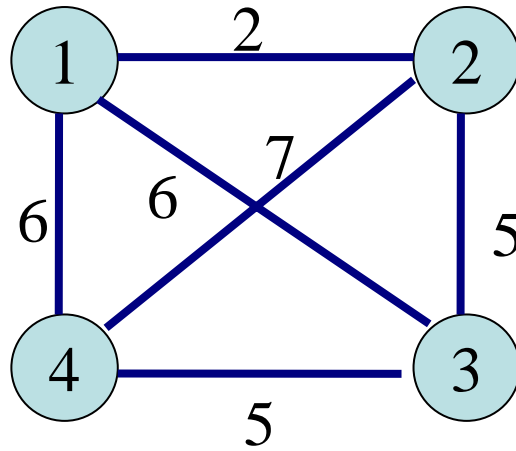
$$\text{blumin} = \min(\min_0, \min_1, \min_2) = 2$$

Privately find all edges of length bluemin



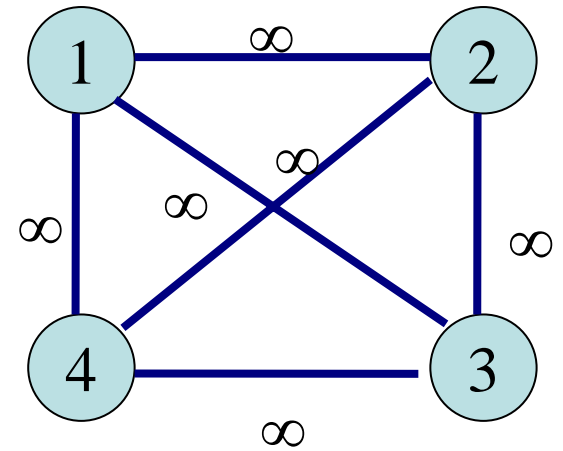
G_1'

$$S_1 = \{e_{23}\}$$



G_2'

$$S_2 = \{e_{12}\}$$

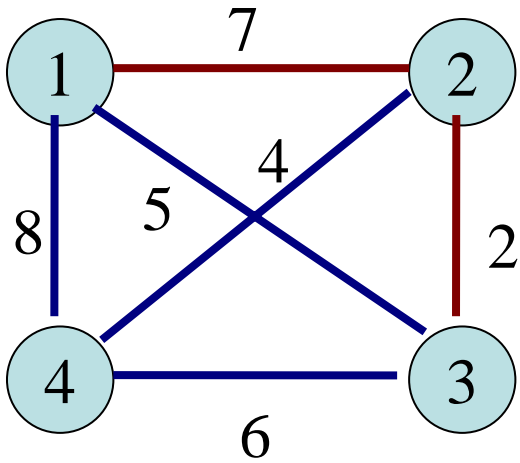


G_0'

$$S_0 = \{\}$$

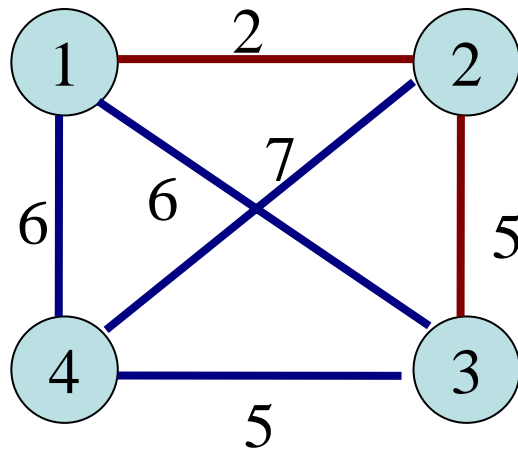
$$S = S_1 \cup S_2 \cup S_3 = \{e_{12}, e_{23}\}$$

Update S edges in G_0'



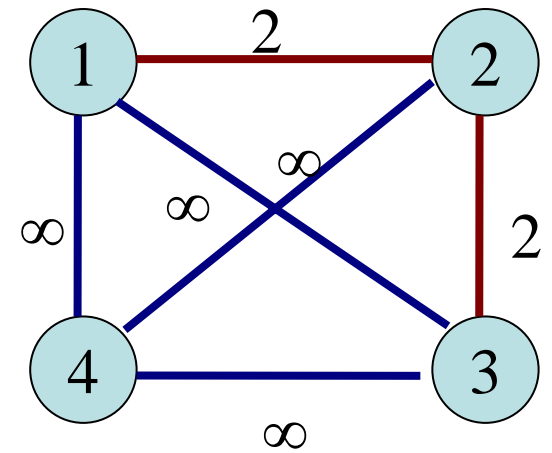
G_1'

$$S_1 = \{e_{23}\}$$



G_2'

$$S_2 = \{e_{12}\}$$

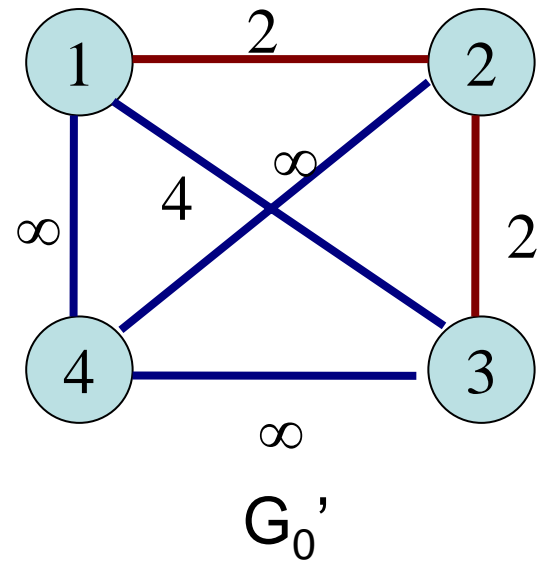
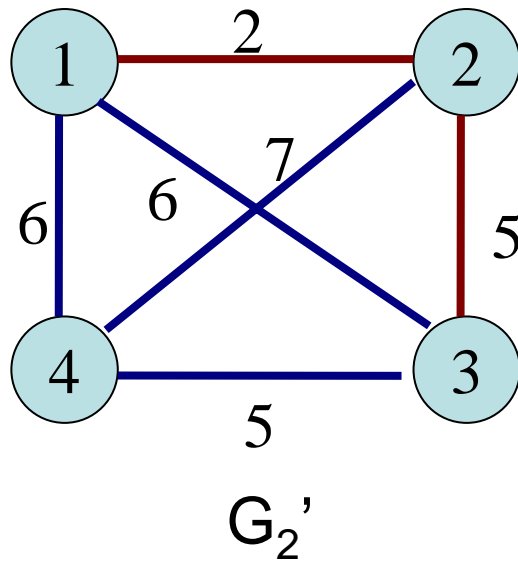
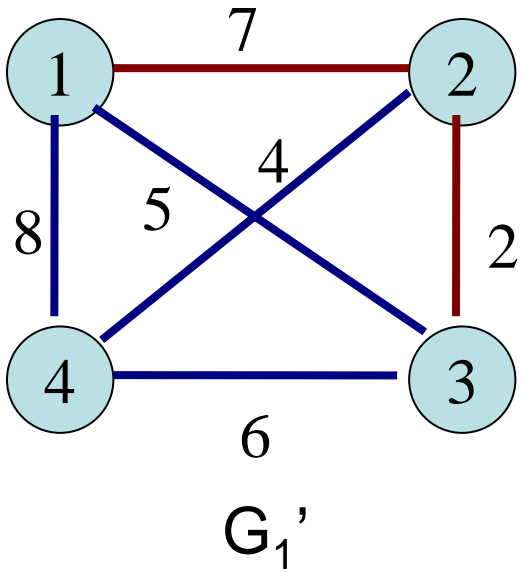


G_0'

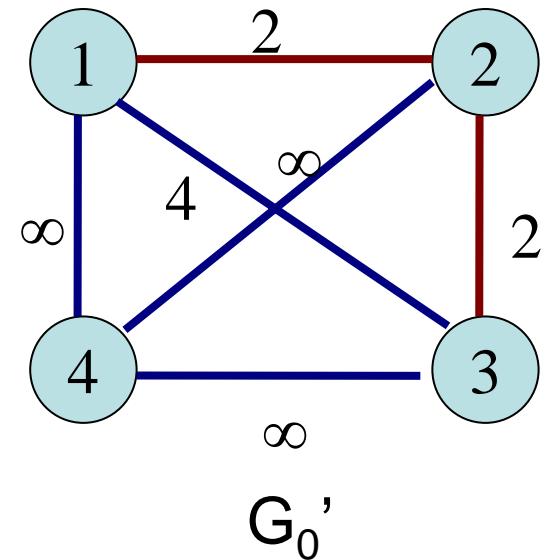
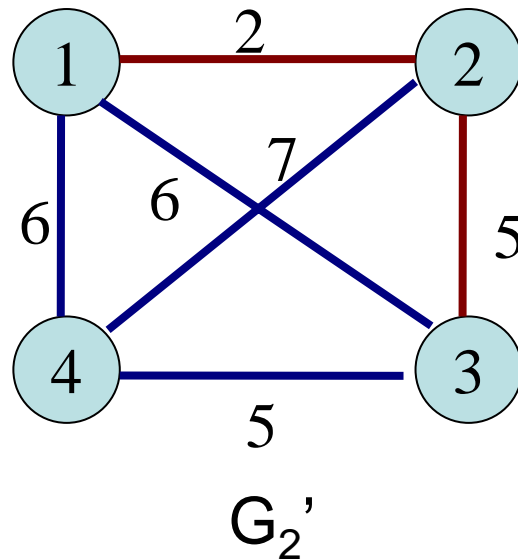
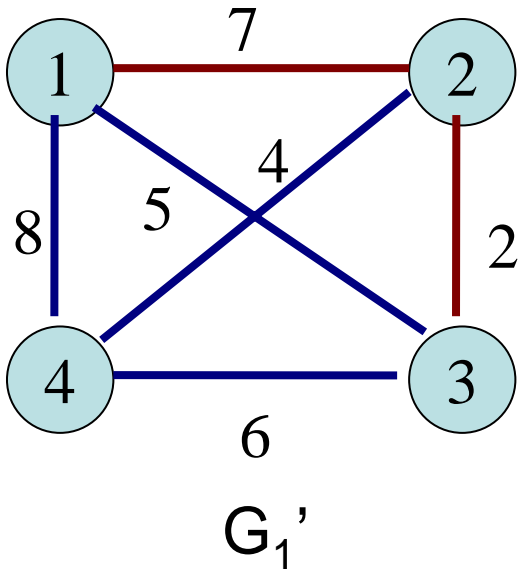
$$S_0 = \{\}$$

$$S = S_1 \cup S_2 \cup S_3 = \{e_{12}, e_{23}\}$$

Run APSP on G_0'

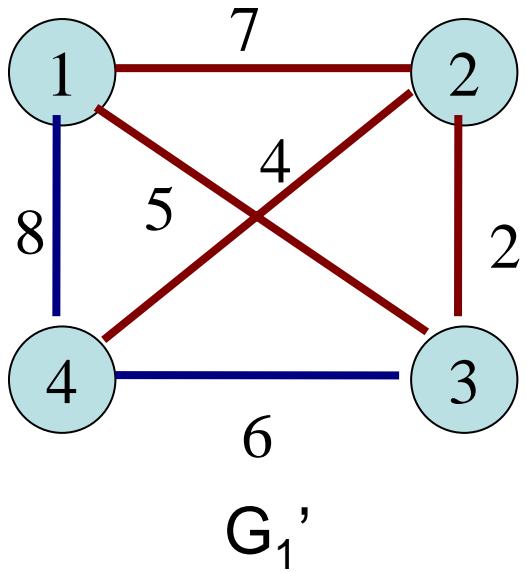


Repeat!

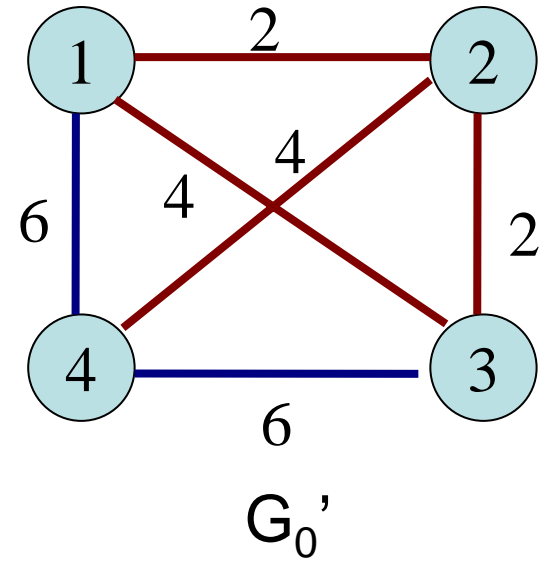
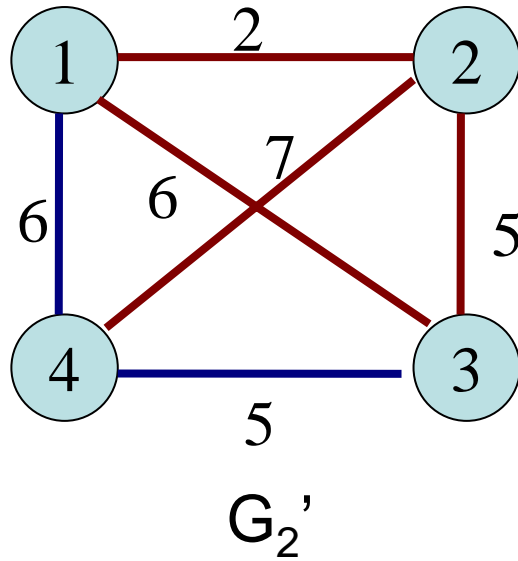


$$\text{bluemin} = \min(\min_0, \min_1, \min_2) = 4$$

$$S = S_1 \cup S_2 \cup S_3 = \{e_{13}, e_{24}\}$$

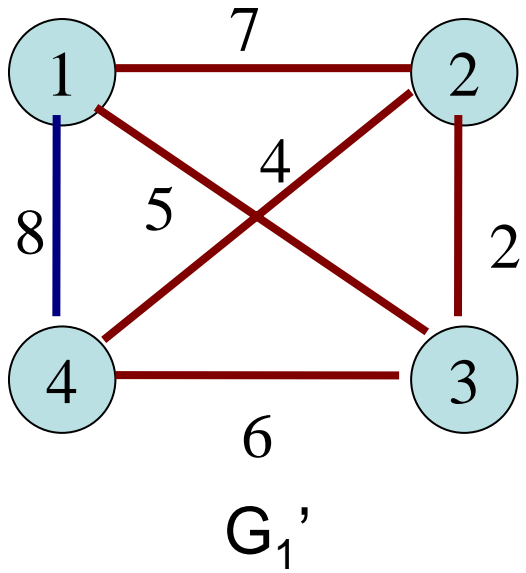


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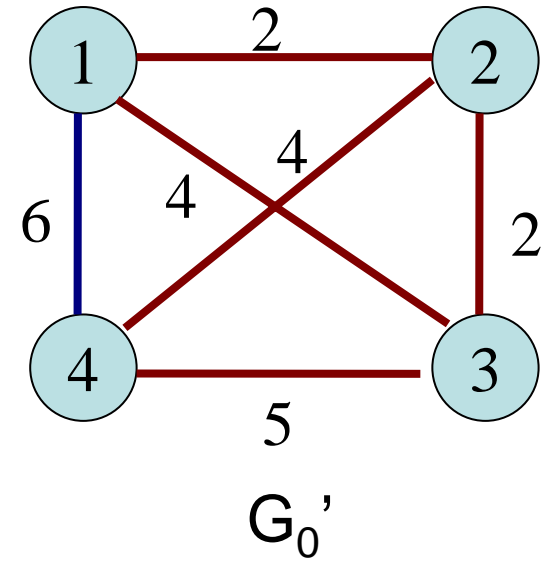
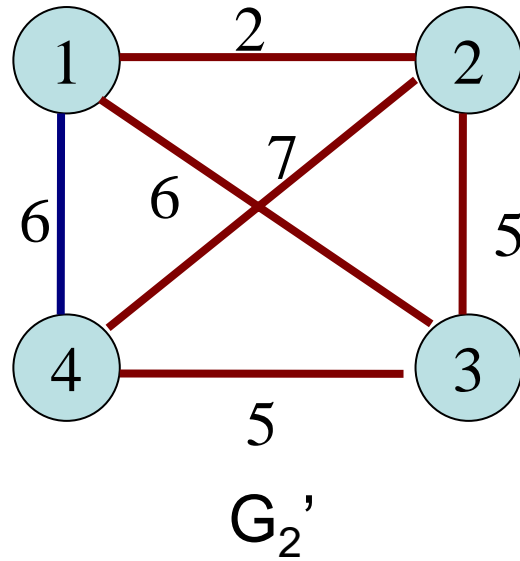


$$\text{bluemin} = \min(\min_0, \min_1, \min_2) = 5$$

$$S = S_1 \cup S_2 \cup S_3 = \{e_{34}\}$$



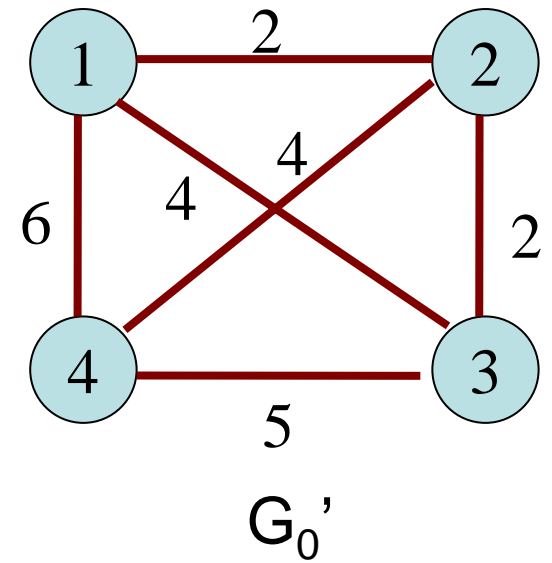
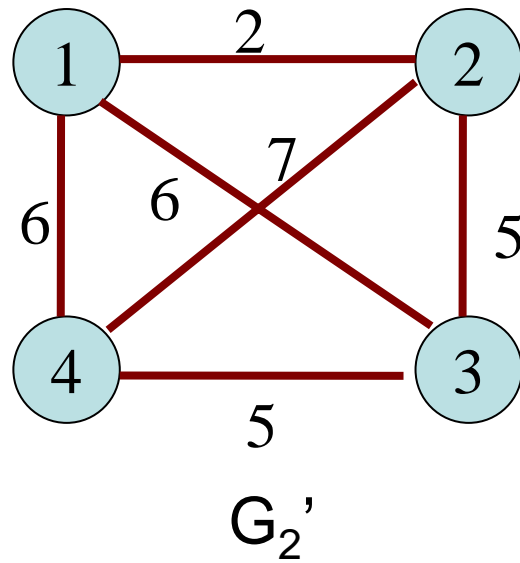
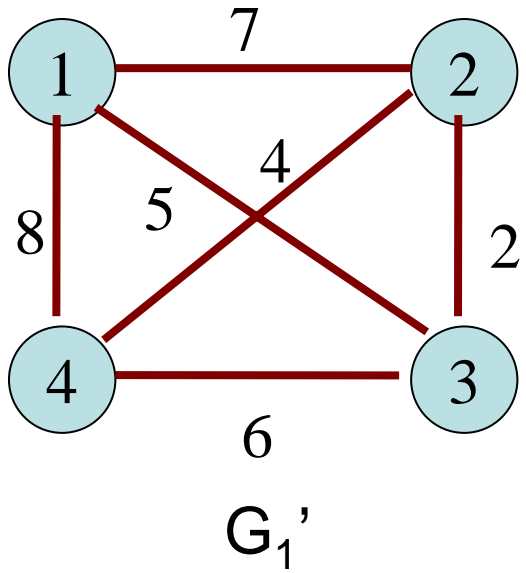
...



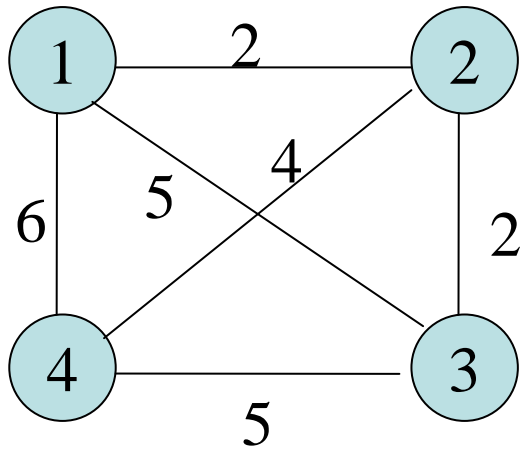
$$\text{bluemin} = \min(\min_0, \min_1, \min_2) = 6$$

$$S = S_1 \cup S_2 \cup S_3 = \{e_{14}\}$$

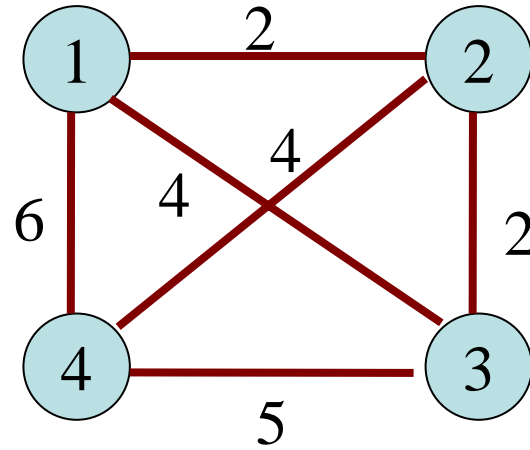
... until all edges are red



The solution is correct!



$\text{gmin}(G_1, G_2)$



G_0'

Other Results

- A similar protocol for $\text{SSSD}(g_{\min})$
 - This isn't free!
- Protocol for special case of $\text{APSD}(g_{\max})$ and $\text{SSSD}(g_{\max})$
 - Input graphs obey triangle inequality
- “Leaky” protocol for $\text{APSD}(g_{\max})$ and $\text{SSSD}(g_{\max})$ in the general case

Final Thought

- Other graph algorithms don't leak enough information
- Questions?