

CS345H: Programming Languages

Lecture 17: Introduction to Object-Oriented Languages

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Overview

- ▶ Today, we will study Object Oriented Programming
- ▶ Not in the sense of how to use it, but to understand the fundamental aspects of this paradigm
- ▶ We will also explore how to formalize some aspects of it

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Subtyping

- ▶ Before we actually look at object-orientation, let's start with a more basic observation that precedes object oriented thinking
- ▶ If we have expression e_1 with type τ_1 and expression e_2 with type τ_2 , we can sometimes safely use e_2 instead of e_1
- ▶ **Example:** $\tau_1 = \text{Int} \rightarrow \text{Int}$ and $\tau_2 = \alpha \rightarrow \text{Int}$
- ▶ We have seen this before: [Polymorphism](#)

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Another Kind of Polymorphism

- ▶ However, there is also another kind of polymorphism that arises from records called **subtype polymorphism**:
- ▶ Consider the following two data types in Java:

```
class X {
    public int a;
};
class Y {
    public int a;
    public int b;
};
```
- ▶ Here, anyone who expects something of type X can work just as well with something of type Y
- ▶ **Why?** Because the fields of Y are a superset of the fields of X

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Subtyping Polymorphism

- ▶ Intuitively, if B is a subtype of A , it has to be safe to use a B wherever an A is expected
- ▶ But the exact definition of subtype depends on the language!
- ▶ Observe that subtyping is really another kind of polymorphism as it allows us to write code that works with more than one type

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Structural Subtyping

- ▶ One possible definition of subtype: B is a subtype of A if B has all fields A has
- ▶ This is called **structural subtyping**
- ▶ With this, there is no need to declare that one type is a supertype of another, it is just a direct function of the **structure** of a type
- ▶ This is also known as “duck typing”

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Nominative Subtyping

- ▶ Contrasting approach: Programmer must explicitly declare subtype relationships
- ▶ Usually, this is combined with language mechanisms that ensure the resulting type is actually compatible
- ▶ This is known as **nominative subtyping**

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Classes

- ▶ Next Key Idea using in object-oriented languages: Classes
- ▶ A **class** combines data with functions that operate on this data
- ▶ Crucially, the methods operating on data are bundled **together** with the data and can access the data
- ▶ You can view a class as a type that includes both data and functions to manipulate this data
- ▶ We create an **instance** of a class to use it
- ▶ Every instance has its own data

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Classes Example

- ▶ Consider the following program:

```
struct point {
    int x;
    int y;
};
void inc_x(point* p) {
    p->x++;
}

point *p = new point;
inc_x(p);
```

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Classes Example Cont.

- ▶ Here is the same program rewritten with classes:

```
class point {
    int x;
    int y;
    void inc_x() {
        this->x++;
    }
};

point *p = new point;
p->inc_x();
```

- ▶ Here, p is an **object**
- ▶ Object = instantiated class

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Terminology of OO

- ▶ Class: A kind of **type** that combines data and functions to operate on this data
- ▶ Instance of a class: An **object** build from the class type
- ▶ **Important:** There is one class X in a program (even before running it), but potentially many instances of X at run-time
- ▶ Functions in classes are called **methods**
- ▶ Data fields are (sometimes) called **instance variables**
- ▶ **Important:** Every instance of a class has its own set of instance variables!
- ▶ **Constructor:** A special method that is run on instance creation

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Classes Example Extended

- ▶ Here is a slightly extended version of the program

```
class point {
    int x;
    int y;
    point() { this->x = 0; this->y = 0; }
    void inc_x() {
        this->x++;
    }
};

point *p = new point();
point *q = new point();
p->inc_x();
```

- ▶ What is the value of y->x? 0

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Objects and Subtyping

- ▶ Classes combine very powerfully with subtyping
- ▶ But subtyping now also needs to take methods into account!
- ▶ Specifically, any subtype must also have (at least) the same methods as the original type
- ▶ This way, we can use the subtype in any context that expects the original type

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Virtual Methods and Subtyping

- ▶ But this has one problem. In general, we may need to change the implementation of an existing method when defining a subclass
- ▶ Allowing redefinition of methods is known as **virtual methods**
- ▶ This is another key OO feature

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Virtual Method Example

- ▶ Consider the following code:

```
class point {
  int x;
  int y;
  point() { this->x = 0; this->y = 0; }
  virtual void inc_x() {
    this->x++;
  }
};
class bigpoint:public point {
  int x;
  int y;
  virtual void inc_x() {
    this->x+=2;
  }
};
point *p = new bigpoint();
p->inc_x();
```

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Static vs. Dynamic Types

- ▶ Allowing subtyping of classes means that every expression in a program has two kinds of types:
 1. **Static Type**: This is the type computed by the type checker at **compile time**. For example, the static type of `p` in the example is `point`.
 2. **Dynamic Type**: This is the type of object a variable holds at **run-time**. For example, the dynamic type of `p` in the example is `bigpoint`.
- ▶ **The dynamic type decides which virtual method is called!**
- ▶ This is sometimes called the essence of OO

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Relationship between Static and Dynamic Type

- ▶ The dynamic type of an object must always be a **subtype** of its static type in any well-typed program
- ▶ Observe that the **semantics** of a program with virtual methods require knowing the type of the object at run-time
- ▶ This is a new use of typing: So far, we only used types to prevent run-time errors
- ▶ In OOP, we also use types **at run-time** to decide which method to invoke!
- ▶ You can think of this operationally as an implicit run-time check on a type tag that decides which version of a method is called

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Encapsulation

- ▶ The last major OOP feature we have not yet discussed in **encapsulation**
- ▶ This means that class data can be made non-accessible to clients of the class
- ▶ **Example**: Declaring an instance variable `private`
- ▶ Fortunately, this only rejects some programs at compile time but does not change semantics

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Why Encapsulation?

- ▶ But why do we want to restrict access to object data?
- ▶ The idea is to force clients of a class to only rely on its public interface
- ▶ We are therefore free to change the implementation of classes without affecting its clients
- ▶ And we are free to pass any subtype with the same public interface
- ▶ **Actually, old idea:** This is also known as **abstract data types (ADT)** and predates OOP

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Essential OOP Features

- ▶ The following four features are usually considered necessary for a language to be object-oriented:
 1. Subtyping
 2. Classes
 3. Virtual Methods
 4. Encapsulation

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Why OOP?

- ▶ The allure of the OOP model is that it allows you to view your program as a collection of interacting entities (objects) instead of a collection of data and sets of functions
- ▶ This often allows for much cleaner and more extensible code
- ▶ **If the problem you are solving fits into the OO model!**

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Uses of OOP

- ▶ OOP is a great fit for:
 - ▶ GUI toolkits
 - ▶ Data Structures
 - ▶
- ▶ In fact, OOP fits so many problems reasonably well that it has become the **default** paradigm used in most software
- ▶ In fact, many modern languages, such as Java, **force** an OOP style
- ▶ However, not every problem maps well into objects!

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History of OOP

- ▶ The idea of objects and subtyping originated at MIT in the 1950s and 1960s
- ▶ This was in the context of AI research in LISP
- ▶ Over the years, various features that we would call object-oriented today made their way into various LISP dialects

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Object Orientation: History Cont.

- ▶ First object-oriented language: **Simula 67**
 - ▶ Developed by Ole-Johan Dahl and Kristen Nygaard at the Norwegian Computing Center in Oslo
 - ▶ Simula was designed as an **special-purpose** language for **discrete event simulations**
 - ▶ But it was certainly not designed as a general-purpose programming language

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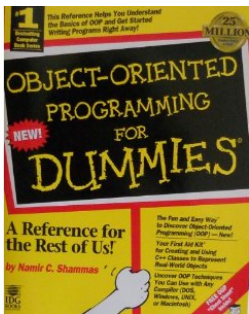
Features Combined in Simula

- ▶ Subtyping
- ▶ Classes
- ▶ virtual methods
- ▶ Garbage Collection

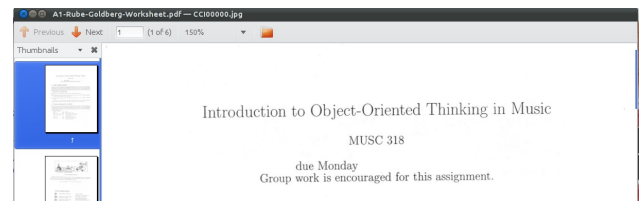
From Simula to Smalltalk

- ▶ Simula was extremely successful in speeding up writing of discrete event simulations
- ▶ In fact, so successful that a group at Xerox PARC decided to create a general-purpose language based on this paradigm called Smalltalk
- ▶ Smalltalk coined the term “object oriented”
- ▶ And lead to a huge wave of OO languages
- ▶ This was an **huge** fad in the late 90's

OO in the Real World



OO in the Real World



Other OO languages

- ▶ C++
- ▶ Java
- ▶ C#
- ▶ Pretty much any newer imperative language (and plenty functional ones as well)

The OO paradigm in the wild

- ▶ Object oriented programming is one of the very few techniques that actually seems to make it easier for humans to build software
- ▶ But it does not work well for every problem
- ▶ However, since it is so pervasive, this tends to be forgotten

Summay

- ▶ We have looked at the four aspects that define object-oriented programming
- ▶ **Next time:** Some issues with semantics and typing in OO languages