Why OO approach for DB?

- Complex Data Handling
 - 1) Design software such as CAD, CASE, production planning
 - 2) Binary files such as image, audio, video

Why OO approach for DB? (Cont.)

Consistent Data Model

Relational: Information regarding entity can be spread across several tables.

OO approach: 1-to-1 relationship between real-life entities and database objects representing them.

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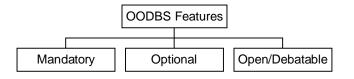
Lack of Specification

Caveat: The paper was written long time ago

- No single specification, everybody had their own spec. and implementation.
- Ill-planned system could emerge and become de-facto standard

The Novel Idea: A Common Spec

• OODBS's should have a common set of characteristics



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Mandatory Features

Mandatory OO Features

+

DB Features

=

Mandatory OODBS Features

Mandatory OO Features

1. Complex Objects

Complex Objects can be built from simpler objects using object constructor.

Example: HTML Code

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Mandatory OO Features (Cont.)

2. Object Identity (OID)
Each Object has a unique identity
Class Person { string name; int age};
Person assistant, manager;

Memory Loc.

assistant = ("John", 24) 000111 maneger = ("John", 24) 000222

Two objects are equal but not identical. Identity maintained by system, not user.

Mandatory OO Features (Cont.)

- 3. Encapsulation
- Knowledge about only specification is sufficient for application developer.
- Implementation change doesn't affect the applications
- Only visible part is interface and allowable functions are only those defined in interface

Mandatory OO Features (Cont.)

- Classes Vs. Types
- 1) Type-checking(static vs. dynamic)

2) new

Class

3) extent (optional)

Mandatory OO Features (Cont.)

4. Classes Vs. Types (Cont.)

Run-time Compile-time $\mathbf{V}\mathbf{s}$. C++

SmallTalk Dynamic type-checking

Class Person {int x;

foo();}

// error

Person john; john.bar(); //run-time

john.bar(); // compile-time // error

Static type-checking

Mandatory OO Features (Cont.)

5. Inheritance

Class Inherit. Vs. Type Inherit.

SmallTalk

Purpose:

Code-reuse

C++

Purpose:

Interface Compliance

Code reuse

Mandatory OO Features (Cont.)

6. Overloading, Late-binding

- · Less work at application programmer level

ostream display(obj Person); ostream display(obj Image); ostream display(obj Matrix);

• Application Programmer:

main () { new person; person.sex='M'; display(person);

new matrix; matrix.size = 4; display(matrix); }

Mandatory OO Features (Cont.)

- 7. Computational Completeness
- DBS should be able to execute any computable expression fed by Data **Manipulation Language**
- SQL is computationally incomplete (e.g., doesn't support recursion)

Mandatory OO Features (Cont.)

8. Extensibility

 Predefined types that would come with a DBS should be extensible

Mandatory OO Features (Cont.)

9. Query Facility

- Application programmers or users should be able to query on data
- Methods: Query language, GUI etc
- 1) High Level, 2) Fast, 3) One type works for all applications

DB Features

- 1) Persistence
 - -- Data survive the execution time of program creating or manipulating it.
- 2) Storage Management
 - -- index mgt., query optimization etc.
 - -- required for managing large DB
- 3) Concurrency
- 4) Recovery

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Optional Features

- 1) Multiple Inheritance
 - -- conflict resolution
- 2) Type Checking and Inferencing
 - -- the more the compile-type checking the better
- 3) Distributed
 - -- distributed among many computers, need distributed data management

Open Features

Open because: OO-ness is debatable, no general consensus

- 1) Programming Paradigm
 - -- Functional/logical/imperative etc.
- 2) Type Formation
 - -- Type formers can be added
- 3) Uniformity
 - -- **Type < > Object < > Method**

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Afterthoughts

- Most DB Developers are new to OODB Concepts, Relational-to-OO conversion cost, departure from set theory
 - -- Can't we extend RDB's to provide the facilities of OODB's?

ODMG-93

- Standard
 - -- Portability Vs. Interoperability

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ODMG Standard (Cont.)

- Portability
 - -- Current OODMS (e.g., O₂, Poet etc.)





ODB Schema

- Relational Schema
 - -- Tables/Data structures, Data Types
- ODB Schema
 - -- Data structures, Data/Object types, methods(**).
- ** Programming Lang. needed to write methods

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Components of ODMG Model

- 1) Data Model
- 2) Language Bindings
- 3) Query Language

Data Model

- ODB Schema
 - -- Three choices: ODL, OMG IDL, OO Language

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Data Model(Cont.)

- ODMG(Object Database Mgt. Group) model is an extension of OMG (Object Mgt. Group.
- OMG group deals with object modeling in general, not particular to DB Systems.

Data Model(Cont.)

- OMG General Object Model
 - -- Class
 - -- Instances
 - -- Methods
 - -- Inheritance
 - -- Encapsulation
 - -- Classic Types (e.g., date, time)

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Data Model(Cont.)

- DB extension to OMG model
 ODMG model
- 1) Relationship
- 2) Collection

Data Model(Cont.)

- Collections
 - -- Container for holding instances of a class. Example: Set< Ref<Person>> where Person is a class
 - -- Collection types are template. Example: Set<T>, where T can be any type

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Data Model(Cont.)

• Relationship (1-to-1)

Class Apartment;

Class Person { String name; Int SSN;

Ref<Apartment> lives_in inverse is_used_by };

Class Apartment

{Ref<Person> is_used_by inverse lives_in };

Data Model(Cont.)

• Relationships (1-to-many, many-to-many)

Class Person {

Set < **Ref**<**Person>** > **parents inverse children**;

List < Ref<Person> > children inverse parents; }

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Object Query Language (OQL)

- Prologue:
 - -- C++ by itself is not enough because
 - Even for a short query, one needs to write, compile, link a C++ file
 - Too much and tedious code

Object Query Language (cont.)

- OQL can help
 - -- Alleviates problems on previous slide
 - -- Extension to host language (e.g., C++). Objects can be manipulated by both.

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Object Query Language (cont.)

- OQL Query Example:
 - -- Select m.year

From Movies m

Where m.title = "Godfather"

· OQL expression results can be assigned directly to host language variables C++ variable

Class Movie { ...}; Set<Movie> oldmovies; oldmovies = SELECT DISTINCT m From Movies m WHERE m.year>1990

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Object Query Language (cont.)

· OQL's output type can be complex

SELECT DISTINCT struct(star1: s1, star2: s2) FROM Stars s1, Stars s2

WHERE s1.address = s2.address AND s1.name<s2.name

· Recursive OQL Query

```
Set < Ref<Person> > Person::ancestors()
Set < Ref<Person> > result;
oql(result, "flatten(select distinct a->ancestors from $1c as a) \
                                       union $1c", parents);
return result;
```

Object Query Language (cont.)

· Compare the ease of data transfer in OQL with the difficulty in SQL (impedance mismatch)

```
void getStudio()
 EXEC SQL BEGIN DECLARE SECTION;
    char studioName[50], stuidoAddr[100];
 EXEC SQL END DECLARE SECTION;
 printf("Enter studio name"); studioName = getline();
 printf("Enter studio addr"); studioAddr = getline();
 EXEC SQL INSERT INTO Studio(name, address)
             VALUES (: studioName, : studioAddr);\\
}
```

Afterthoughts

· Schema changes in OODB result in a system-wide recompilation.

How to solve the problem?

• Is OQL really impedance mismatch free? Syntax, binding stage, programming style are still different.