

0x1A Great Papers in Computer Security

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<http://www.cs.utexas.edu/~shmat/courses/cs380s/>

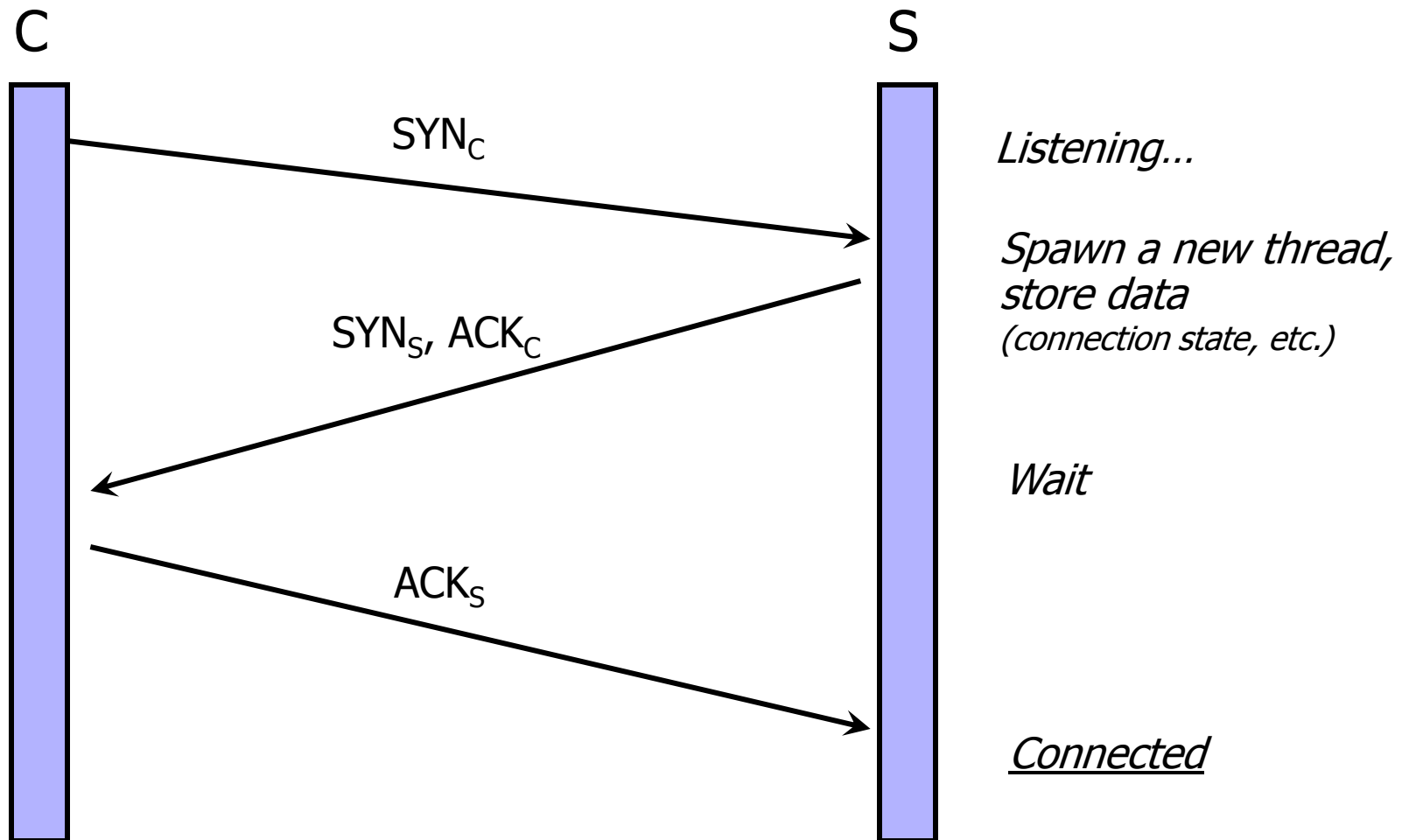
D. Bernstein

SYN cookies

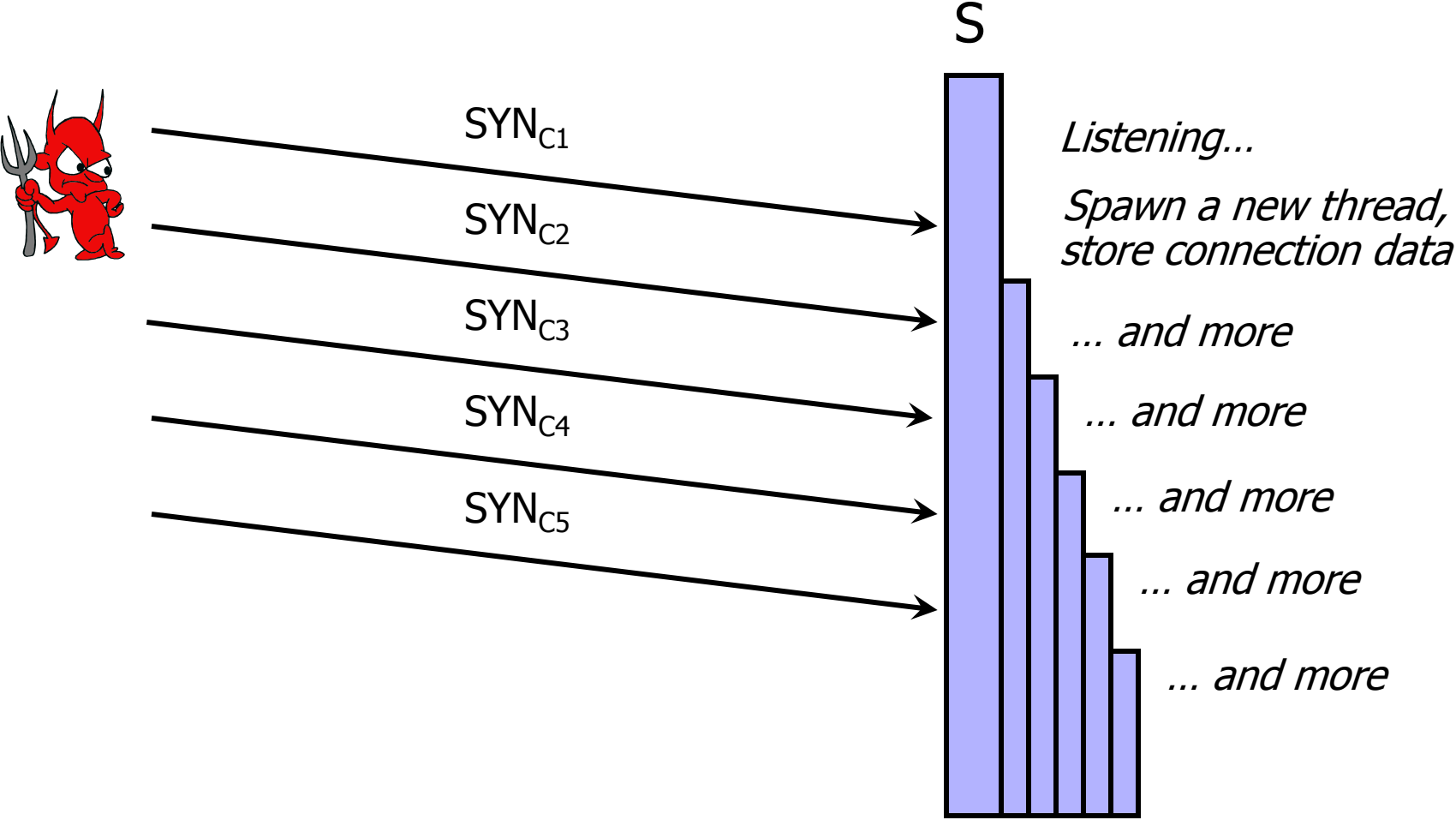
(1996)



TCP Handshake



SYN Flooding Attack



SYN Flooding Explained

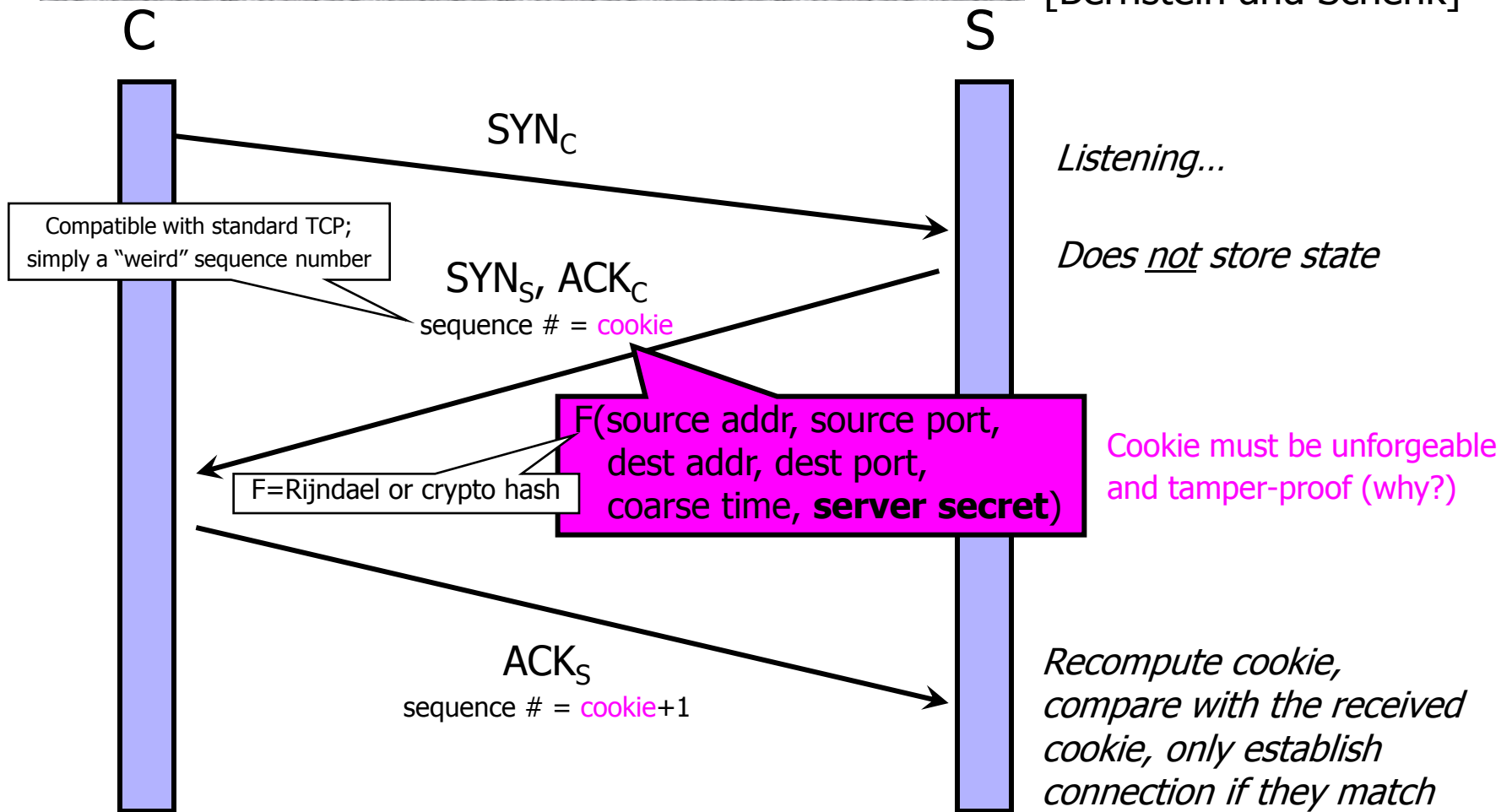
- ◆ Attacker sends many connection requests with spoofed source addresses
- ◆ Victim allocates resources for each request
 - New thread, connection state maintained until timeout
 - Fixed bound on half-open connections
- ◆ Once resources exhausted, requests from legitimate clients are denied
- ◆ This is a classic denial of service attack
 - Common pattern: it costs nothing to TCP initiator to send a connection request, but TCP responder must spawn a thread for each request - **asymmetry!**

Preventing Denial of Service

- ◆ DoS is caused by asymmetric state allocation
 - If responder opens new state for each connection attempt, attacker can initiate thousands of connections from bogus or forged IP addresses
- ◆ **Cookies** ensure that the responder is stateless until initiator produced at least two messages
 - Responder's state (IP addresses and ports of the connection) is stored in a cookie and sent to initiator
 - After initiator responds, cookie is regenerated and compared with the cookie returned by the initiator

SYN Cookies

[Bernstein and Schenk]



More info: <http://cr.yp.to/syncookies.html>

Anti-Spoofing Cookies: Basic Pattern

- ◆ Client sends request (message #1) to server
- ◆ Typical protocol:
 - Server sets up connection, responds with message #2
 - Client may complete session or not - potential DoS!
- ◆ Cookie version:
 - Server responds with hashed connection data instead of message #2
 - Client confirms by returning hashed data
 - If source IP address is spoofed, attacker can't confirm
 - Need an extra step to send postponed message #2, except in TCP (SYN-ACK already there)