

**A Mechanically-Checked
Correctness Proof of a
Floating-Point Search Program**

Matt Wilding

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Computational Logic Inc.
1717 W. 6th St. Suite 290
Austin, Texas 78703
(512) 322-9951

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1. Introduction

This technical report describes work toward the mechanical certification of floating-point program correctness theorem proofs. A library of useful facts about rational numbers was constructed. A set of axioms about floating-point operations has been proved consistent using a mechanical theorem prover that employs the rational number library. A small floating point program has been developed and mechanically-checked correctness theorems about it constructed.

Floating-point numbers are rational numbers that can be expressed in the form

$$\text{sign} \times \text{value} \times b^{\text{exp}}$$

where *sign* is 1 or -1, *value* is a base *b* number with a number of digits fixed by the floating-point system, and *exp* is in a range fixed by the floating-point system.

Advantages inherent in floating-point numbers were recognized early in the development of modern computers, and many early machines had floating-point capability. [13] Floating-point numbers seemed to make programming easier as operations on very small numbers and very large numbers had the same precision and could usually be handled with little concern about violating range restrictions.

Floating-point arithmetic operations are called *inexact* because they sometimes return values that are "close to" the exact result. An arithmetic operation applied to two rational values that are floating point values may yield a rational value that is not a floating-point value. This inexactness can make analyzing floating-point programs very difficult.

The importance of developing program correctness arguments has become more obvious as programming has matured and programs have become more complex and relied-upon. Techniques for program proof have been developed and applied in many problem domains. There has been much less work has been done to prove the correctness of floating-point programs. Testing practices developed with an understanding of the implementation of floating-point programs are used by some programmers to give them increased confidence in their programs. [15]

Knuth included analysis of floating-point programs in his programming recipes book. [13] One of the

techniques for analyzing floating-point arithmetic to which he refers is Wilkinson's "backward" error analysis. [18] Rather than bound the inexactness of floating-point operations by bounding the result, Wilkinson shows that it is often simpler to view the inexact result of floating-point operations as the exact result of perturbed arguments.

Several researchers have also proposed models of floating-point operations and used them to justify claims about sequences of floating point operations. [6, 8, 14, 17] These efforts work toward putting floating-point arithmetic on more-solid ground. Even so, the complexity of proofs using these systems, and the apparent gap between the floating-point axioms and "realistic" numerical programs has discouraged application of work in this area.

Much of the effort toward precisely specifying floating-point operations has been motivated by the desire to implement floating-point operations correctly. Researchers at Oxford University have formalized a floating-point system expressed using Z and Occam [1]. Their ultimate goal is to construct a floating-point processor that correctly implements the formal description of the floating-point system.

The problem of wedding a model of floating-point operations with program proof rules was examined in Holm's PhD thesis. [10] Holm uses an axiomatization of floating-point operations and Dijkstra's WP calculus. [9] Holm develops enough mathematical machinery to prove some theorems about a searching program. While the searching program example may appear unambitious and though the presentation is quite clean, the proof is long and rather complex.

On another front, proofs about computers and computer programs have been mechanically checked. Boyer's and Moore's NQTHM prover has been used to prove the correctness of a microprocessor, compilers, and operating systems. [3, 4] The checked theorems often have convoluted proofs, but they have the important advantage that NQTHM-checked theorems are presumed to be very reliable.

Theorems about floating-point programs appear to be a good domain for mechanical-checking as theorems about floating-point programs are often not realizable by programmers because of the complexity of their proof. Knuth observes that "many a serious mathematician has attempted to give rigorous analyses of a sequence of floating-point operations, but has found the task to be so formidable he has tried to content himself with plausibility arguments instead" [13]. The many details of a proof about floating-point programs are often

not of great interest, so the disadvantage of many machine-checked proofs - that the proof is obscure - is usually not relevant in this domain.

Nevertheless, to the author's knowledge no mechanically-checked proof about a floating-point program has ever before been constructed.

PC-NQTHM, the Kaufmann Proof-Checker extension of the NQTHM prover, allows the user finer control of the direction of the prover. [12] It is particularly valuable when the underlying theory is immature, as is currently the case with the current development of rational number arithmetic and floating-point operations. Many of the theorems in this project were proved correct using the tools provided by this interactive enhancement.

This technical report describes a proof of a floating-point search program. Section I is this introduction. Section II is a description of the rational number library that is used to accomplish PC-NQTHM proofs about rationals. Section III describes the axiomatization of a floating-point system and the proof of its consistency. Section IV details the development of the search program, and section V describes its correctness proof. Section VI describes work planned for the future.

The appendices list theorems proved in the project. Appendix A presents an example rational number theorem proof. Appendix B contains the rational library. Appendix C contains an axiomatization of an FP system. Appendix D lists the development and proof of the FP searching example.

2. The Rational Library

In this section we describe a rational number library that contains facts that facilitate automatic proofs about rationals using the PC-NQTHM theorem prover. It uses Bevier's hardware libraries [2] (as updated by Bevier and Wilding) that facilitate proofs about integers and natural numbers.

The floating-point numbers are a subset of the rational numbers, and our later development of them will rely on the definition of rational numbers and the rational number operations contained in the subsection DEFINITIONS. The rest of this section describes the theorems proved about rationals that appear to facilitate mechanical proofs about rationals. Though these rational facts are necessary to the construction of mechanical proofs about floating-point numbers, only the definitions are needed to understand what has been proved.

The definition of operations in rational arithmetic is usually well below the level of detail with which mathematicians work. We develop rational arithmetic formally so that we can apply the general purpose PC-NQTHM prover to our problem. The most natural notation for this section is therefore the Lisp-like syntax of the logic of the prover. In later sections a non-Lisp notation will be used when the proof is not so inextricably intertwined with the operation of the prover.

2.1 Definitions

The rational data type is added using the NQTHM add-shell event.

```
(add-shell rational nil rational-formp
  ((numerator (one-of numberp negativep) zero)
   (denominator (one-of numberp) zero)))
```

RATIONALP identifies "proper" rational numbers that are of the correct data type, have integer numerators, and non-zero natural number denominators.

```
(definition rationalp (x)
  (and (rational-formp x)
    (integerp (numerator x))
    (not (zerop (denominator x)))))
```

REDUCE reduces a rational to its least common form.

```
(definition reduce (r)
  (if (rationalp r)
      (if (negativep (numerator r))
          (rational (minus (quotient (negative-guts (numerator r))
                                 (gcd (negative-guts (numerator r))
                                       (denominator r))))
                         (quotient (denominator r)
                                   (gcd (negative-guts (numerator r))
                                         (denominator r))))))
          (rational (quotient (numerator r)
                               (gcd (numerator r) (denominator r))))
                     (quotient (denominator r)
                               (gcd (numerator r) (denominator r))))))
      (rational (quotient (numerator r)
                           (gcd (numerator r) (denominator r))))
                 (quotient (denominator r)
                           (gcd (numerator r) (denominator r))))))
  (rational 0 1)))
```

FIX-RATIONAL maps non-rationalps to rational 0

```
(definition fix-rational (x)
  (if (rationalp x) x (rational 0 1)))
```

RZEROOP identifies non-rationals and rational 0

```
(definition rzeroop (x)
  (or (not (rationalp x))
      (equal (numerator x) 0)))
```

RLESSP is the less-than predicate for rationals.

```
(defn rlessp (x y)
  (let ((a (fix-rational x)) (b (fix-rational y)))
    (ilessp (itimes (numerator a) (denominator b))
             (itimes (numerator b) (denominator a)))))
```

REQUAL is the equality predicate for rationals.

```
(defn requal (x y)
  (let ((a (fix-rational x)) (b (fix-rational y)))
    (equal (itimes (numerator a) (denominator b))
           (itimes (numerator b) (denominator a)))))
```

We next define several mathematical operations. Each is defined to return its result in lowest common terms.

```
(definition simple-rplus (x y)
  (let ((a (fix-rational x)) (b (fix-rational y)))
    (rational (iplus (itimes (numerator a) (denominator b))
                      (itimes (numerator b) (denominator a))))
              (itimes (denominator a) (denominator b)))))

(definition rplus (x y)
  (reduce (simple-rplus x y)))
```

```

(defdefinition simple-rneg (x)
  (let ((a (fix-rational x)))
    (rational (ineg (numerator a))
              (denominator a)))))

(defdefinition rneg (x)
  (reduce (simple-rneg x)))

(defdefinition rdifference (a b)
  (rplus a (rneg b)))

(defdefinition simple-rtimes (x y)
  (let ((a (fix-rational x)) (b (fix-rational y)))
    (rational (itimes (numerator a) (numerator b))
              (times (denominator a) (denominator b)))))

(defdefinition rtimes (x y)
  (reduce (simple-rtimes x y)))

(defn simple-rinverse (r)
  (if (rzerop r)
      (rational 0 1)
      (if (negativep (numerator r))
          (rational (ineg (denominator r))
                    (ineg (numerator r))))
          (rational (denominator r) (numerator r)))))

(defn rinverse (r)
  (reduce (simple-rinverse r)))

(defdefinition rquotient (x y)
  (rtimes x (rinverse y)))

(defdefinition simple-rmagnitude (x)
  (let ((a (fix-rational x)))
    (if (negativep (numerator a))
        (rneg a)
        a)))

(defdefinition rmagnitude (x)
  (reduce (simple-rmagnitude x)))

```

2.2 Rules

A set of useful rules has been developed to prove things about rational numbers. This section lists the rules in the current rational number library called R2. A sample is given of the application of each rule.

The intention of this subsection is to convey some of the strategy of the rationals library. To understand exactly when and how the rules are applied, one needs to examine the actual rules that are given in Appendix B. Each of the rules in this theory has been proved to be truth-preserving by a NQTHM prove-lemma event. Each rule name is followed by an example of the application of the rule to a NQTHM term. "x --> y" signifies that term x is rewritten to term y if the rule is applied.

```

rtimes-rneg-arg2
  (rtimes x (rneg y)) --> (rneg (rtimes x y))

```

```

rtimes-rneg-arg1
  (rtimes (rneg x) y) --> (rneg (rtimes x y))

rtimes-rplus-arg2
  (rtimes x (rplus y z)) --> (rplus (rtimes x y) (rtimes x z))

rtimes-rplus-arg1
  (rtimes (rplus x y) z) --> (rplus (rtimes x z) (rtimes y z))

associativity-of-rtimes
  (rtimes (rtimes x y) z) --> (rtimes x (rtimes y z))

rzerop
  (rzerop x) --> (or (not (rationalp x)) (equal (numerator x) 0))

commutativity-of-rtimes
  (rtimes x y) --> (rtimes y x)

rdifference-rdifference-arg2
  (rdifference (rdifference x y) z) --> (rdifference x (rplus y z))

rdifference-rdifference-arg1
  (rdifference x (rdifference y z)) --> (rdifference (rplus x z) y)

rneg-rplus
  (rneg (rplus x y)) --> (rplus (rneg x) (rneg y))

rplus-reduce-arg2-rewrite
  (rplus x (reduce y)) --> (rplus x y)

rplus-reduce-arg1-rewrite
  (rplus (reduce x) y) --> (rplus x y)

reduce-rneg
  (reduce (rneg x)) --> (rneg x)

reduce-rmagnitude
  (reduce (rmagnitude x)) --> (rmagnitude x)

reduce-rquotient
  (reduce (rquotient x)) --> (rquotient x)

reduce-difference
  (reduce (rdifference x y)) --> (rdifference x y)

reduce-rtimes
  (reduce (rtimes x y)) --> (rtimes x y)

reduce-rplus
  (reduce (rplus x y)) --> (rplus x y)

commutativity2-of-rplus
  (rplus x (rplus y z)) --> (rplus y (rplus x z))

associativity-of-rplus
  (rplus (rplus x y) z) --> (rplus x (rplus y z))

equal-requal-rewrite
  (requal a b) --> (requal a c)      given: (requal b c)

transitivity-of-requal
  (requal a c) --> t                  given: (requal a b) and (requal b c)

fix-rational-of-rationalp
  (fix-rational x) --> x            given: (rationalp x)

nrational-rplus-arg2
  (rplus y x) --> (reduce y)        given: (not (rationalp x))

```

```

nrational-rplus-arg1
  (rplus x y) --> (reduce y)      given: (not (rationalp x))

rationalp-means
  (rational-formp x) --> t      given: (rationalp x)
  (integerp (numerator x)) --> t
  (lessp 0 (denominator x)) --> t

means-rationalp
  (rationalp (rational n d)) --> t given: (integerp n) and (lessp 0 d)

rneg-rneg
  (rneg (rneg x)) --> (reduce x)

rneg-reduce
  (rneg (reduce x)) --> (rneg x)

reduce-0
  (reduce (rational 0 x)) --> (rational 0 1)

numberp-numerator-reduce
  (numberp (numerator (reduce x))) --> (numberp (numerator (fix-rational x)))

reduce-nrationalp
  (reduce x) --> (rational 0 1)      given: (not (rationalp x))

rplus-reduce-arg2
  (rplus x (reduce y)) --> (rplus x y)

rplus-reduce-arg1
  (rplus (reduce x) y) --> (rplus x y)

requal-x-x
  (requal x x) --> t

rplus-requal-arg1
  (requal (rplus x y) (rplus z y)) --> t  given: (requal x z)

commutativity-of-rplus
  (rplus x y) --> (rplus y x))

reduce-reduce
  (reduce (reduce x)) --> (reduce x)

requal-reduce2
  (requal x (reduce y)) --> (requal x y)

requal-reduce1
  (requal (reduce x) y) --> (requal x y)

commutativity-of-requal
  (requal x y) --> (requal y x)

rational-generalization
  adds facts about rationalp's when a rational is generalized

fix-rational-rmagnitude
  (fix-rational (rmagnitude x)) --> (rmagnitude x)

fix-rational-rquotient
  (fix-rational (rquotient x)) --> (rquotient x)

fix-rational-rdifference
  (fix-rational (rdifference x)) --> (rdifference x)

fix-rational-rneg
  (fix-rational (rneg x)) --> (rneg x)

```

```

fix-rational-rtimes
  (fix-rational (rtimes x)) --> (rtimes x)

fix-rational-fix-rational
  (fix-rational (fix-rational x)) --> (fix-rational x)

fix-rational-rplus
  (fix-rational (rplus x y)) --> (rplus x y)

fix-rational-reduce
  (fix-rational (reduce x)) --> (reduce x)

rationalp-rmagnitude
  (fix-rational (rmagnitude x y)) --> (rmagnitude x y)

rationalp-rquotient
  (fix-rational (rquotient x y)) --> (rquotient x y)

rationalp-rdifference
  (fix-rational (rdifference x y)) --> (rdifference x y)

rationalp-rneg
  (fix-rational (rneg x)) --> (rneg x)

rationalp-rtimes
  (rationalp (rtimes x y)) --> t

rationalp-fix-rational
  (rationalp (fix-rational x)) --> t

rationalp-rplus
  (rationalp (rplus x y)) --> t

rationalp-reduce
  (rationalp (reduce x)) --> t

commutativity2-of-rtimes
  (rtimes x (rtimes y z)) --> (rtimes y (rtimes x z))

rdifference
  (rdifference x y) --> (rplus x (rneg y))

correctness-of-cancel-rplus
  (rplus a (rplus b (rneg a))) --> (reduce b)

```

We define the rationals library using an NQTHM deftheory event.

```
(deftheory r2 {list of names above})
```

2.3 An example Use of R2

The automatic proof of the following problem was posed as a challenge problem. (Thanks to Matt Kaufmann who got this problem from Bill Pase at a recent conference.) The NQTHM prover proves the theorem in about 30 seconds using the rationals library. (The output of the theorem-prover when constructing this proof is Appendix A.)

```
square(x) = x * x  
four_squares (a, b, c, d) = square(a) + square(b) + square(c) + square(d)
```

Prove:

```
four_squares (a, b, c, d) * four_squares (r, s, t, u)  
=  
four_squares (a*r + b*s + c*t + d*u, a*s + -b*r + c*u + -d*t,  
a*t + -b*u + -c*r + d*s, a*u + b*t + -c*s + -d*r)
```

3. A Floating-Point System Axiomatization

In this section we present a set of axioms that describes the behavior of floating-point numbers. The axioms are formalized in the Boyer-Moore logic, and mechanically shown to be consistent. The axioms themselves are similar to sets proposed in previous FP work, especially [10].

3.1 The Axioms

Floating-point computation is modelled by axiomatizing some functions. The function FPP is axiomatized as a predicate that identifies floating-point values. The function ROUND is axiomatized to map any value to a floating-point value. FPMINIMUM is axiomatized to be the smallest non-zero positive floating-point value. FPMAXIMUM is axiomatized as the largest floating-point value. FPMINSPACE is a non-zero positive value that is smaller than the distance between any two floating-point numbers. ROUND-MIN and ROUND-MAX bound the inexactness introduced by the ROUND function.

To avoid the NQTHM notation that is difficult for non-Lisp programmers, we'll use a more-standard notation. Numerals are assumed to be of type rational, $<$ is RLESSP, $|x|$ means (RMAGNITUDE x), unary - means RNEG, and function application will be denoted $f(args)$ rather than $(f\ args)$.

The axioms were introduced using the CONSTRAIN mechanism described in [5]. Introducing them in this manner guarantees their consistency. (The model used with CONSTRAIN to demonstrate consistency is presented in subsection 3.2)

Here are the axioms:

0) `fpp (x) --> rationalp (x)`

All floating-point values are rationals.

1) `fpp (0)`

0 is a floating-point number.

2) `fpp (1)`

1 is a floating-point number.

3) `rationalp (x) --> fpp (reduce (x)) = fpp (x)`

If x is rational, then the reduction of x (removing common factors of the numerator and denominator) does not affect whether x is a floating-point number.

4) `fpp (fpmmaximum)`

FPMAXIMUM is a floating-point number

5) `fpp (x) --> fpmmaximum >= |x|`

A floating-point number is not larger in magnitude than FPMAXIMUM.

6) `(fpp (x) and x <> 0) --> |x| >= fpmminimum`

A non-zero floating-point number's magnitude is not less than FPMINIMUM.

7) `fpp (round (x))`

ROUND returns a floating-point value.

8) `rationalp (x) --> fpp (- x)) = fpp (x)`

X is a floating-point value iff -X is.

9) `fpp (y) and x >= y --> round (x) >= y`

Applying ROUND to a value X no less than a floating-point value Y will not return a value less than Y.

10) `fpp (y) and x <= y --> round (x) <= y`

Applying ROUND to a value not greater than floating-point value Y will not return a value greater than Y.

11) `x >= FPMINIMUM and x <= FPMAXIMUM --> round(x) >= (ROUND-MIN * x)`

Values in range to be floating-point numbers when rounded will not be smaller than ROUND-MIN times their original value.

12) `x >= FPMINIMUM and x <= FPMAXIMUM --> round(x) <= (ROUND-MAX * x)`

Values in range to be floating-point numbers when rounded will not be larger than ROUND-MAX times their original value.

13) `0 <= FPMINSPACE`

FPMINSPACE is a positive, non-zero value.

14) `round (- x) = - (round (x))`

Applying ROUND to -X yields the same value as negating the result of applying ROUND to X.

```
15)      fpp (x) and delta > 0 and delta < FPMINSPACE
-->
not fpp(x + delta)
```

Values within FPMINSPACE of a floating-point value not equal to that floating-point value are not floating-point values.

3.2 A Model that Shows the Axioms to be Consistent

The first thing we'd like to show with our set of axioms is that the axioms themselves are consistent. If there is a set of functions that behave in the manner proscribed by the axioms above then the axioms will have been shown to be consistent.¹

Though the axioms are suggestive of a conventional floating-point system, assigning simple functions in the manner described below makes all the axioms true.

```
fpp(x)      (x = 0) or (x = 1) or (x = -1)

fpminspace  1

fpminimum   1

fpmaximum   1

round-max   1

round-min   1

round(x)    if |x| < 1
            0
            if x >= 0
            1
            -1
```

3.3 What is a Floating-Point Program?

The NQTHM logic will be used to express programs. [4] This report contains programs using infix notation to represent terms in the logic in order to assist readers unfamiliar with the logic's Lisp-like notation. Even so, the ultimate authority about what has been proved is in the proof script that has been checked by the theorem prover.

Several axioms were presented in the previous subsection that describe floating-point operations.

¹Though several sets of axioms for floating-point operations have been proposed in the literature, to my knowledge the consistency of a set of axioms has never been addressed. It's not hard, and it is surprising that it has not been handled explicitly before. This is typical of the kind of thing that might be skipped in a written proof but which is required in a mechanically-checked proof.

Programs that manipulate floating point values will use some of the functions described, such as **round** and **fpp**. They will also use non-floating-point functions, such as **if**.

One potential problem with using the NQTHM logic as our programming notation is that some programs that can be expressed in the logic are outside our intended domain. This is not normally a consideration, since programming languages with floating point operations typically exclude unreasonable statements. For example, a program that performs an exact multiplication of two floating-point values is not possible to express in most programming languages, and appears to violate the spirit of what we mean by a "floating-point" program. Precisely what NQTHM functions would commonly be considered floating-point programs is open to debate since what is "generally considered a floating-point program" is not a formal specification for deciding what is a floating-point program and what is not.

For our purposes, the following conditions will be sufficient to claim that a program is a floating-point program.

- Every instance of **rplus**, **rdifference**, and **rtimes** is applied to arguments that are provably **fpp**
- Every instance of **rquotient** is applied to arguments that are provably **fpp** or has a provably **fpp** first argument and a rational constant power of two second argument.
- Every instance of **rplus**, **rdifference**, **rtimes**, and **rquotient** is immediately surrounded by a **round** function.
- Every instance of **rmagnitude** and **rneg** is applied to an argument that is provably **fpp**.
- Every instance of **rlessp** and **requal** is applied to arguments that are provably **fpp**.
- Every instance of a rational operation other than the eight already mentioned - **rplus**, **rdifference**, **rquotient**, **rtimes**, **rneg**, **rmagnitude**, **rlessp**, and **requal** - will be translatable into an operation composed only of those eight that meets the previously-described restrictions and functions commonly built-in to programming languages.
- Every function used in the program is either a floating-point function or a function commonly built-in to programming languages (such as **if**).

Note that the property of being a floating-point program is not purely syntactic since an argument to a rational function may need to be proved to be a floating-point value.

4. An Example Floating-Point Program

In this section we develop a program that finds a zero of an arbitrary floating-point function. We wish a zero-finding program to return a pair of floating-point values that bound a small region containing a zero. This means that the function applied to the two endpoints returns values with opposite signs.

We will solve the zero-searching problem by writing a floating-point program that checks to see if the current region is small enough, and if it isn't finds a midpoint of the region and recursively searches one of the two subregions.

The function **find-func-zero** will look like this:

```
find-func-zero (a,b)
=
  if unreasonable (a,b)
    0
  if close (a,b)
    (a,b)
  let mid := fp-mid (a,b)
    if sign (func (mid)) = sign (func (a))
      find-func-zero (mid,b)
      find-func-zero (a,mid)
```

The following subsections define the functions **func**, **fp-mid**, **close**, and **unreasonable**. **fp-mid** is really a family of search programs that works for different **func**s, and different values of the floating-point constants. Properties we require we add as axioms, using CONSTRAIN events to ensure consistency.

4.1 func

Since **find-func-zero** is supposed to find a zero of an arbitrary floating-point function **func**, we should specify the function **func** as weakly as possible. We can do this by adding the axiom that **func** returns a floating-point value, and not defining precisely what it returns. To insure consistency, we witness the function **func** by the function that returns a floating-point value of 0.

Axiom: **func(x) is a floating-point value.**

4.2 fp-mid

There are several programming choices we can make for finding the midpoint of a region. We will use a very obvious program, namely adding the two values and dividing by 2. Other choices would eliminate the problem of floating-point overflow. For example, if the two values were each divided by 2 and then added

together, overflow could not occur. Our choice will limit the applicability of the program slightly, but our example will demonstrate that the most-obvious programming solution can be specified and proved correct based upon the floating-point axioms.

The axiomatization of floating-point operations suggests another potential problem with this program. The values might be small enough that floating-point-adding them together and exactly-dividing by 2 will yield a value less in absolute value than FPMINIMUM. If this happens, the result is not guaranteed to fall between the two endpoints. By the axioms of floating-point, we can prove that this will not occur if the smaller (in absolute value) endpoint is at least FPMINIMUM/ROUND-MIN (in absolute value.) We wish to work with floating-point values, so since this value is important we axiomatize a constant to be a floating-point value at least as large as this value. The constant will be called MID-BOUND2.

```
proposed axiom: fpp(MID-BOUND2)
    and
    MID-BOUND2 > FPMINIMUM/ROUND-MIN
```

We wish to witness this constant to show that adding an axiom like the one above does not cause an inconsistency. Unfortunately, our lemmas so far do not guarantee that MID-BOUND2 as described above actually exists. It is consistent with our axioms of floating-point that FP MAXIMUM = FPMINIMUM, in which case MID-BOUND2 does not exist. It therefore is necessary to weaken this axiom, and we add the following instead.

```
Axiom: fpp(MID-BOUND2)
    and
    (MID-BOUND2 >= FPMINIMUM/ROUND-MIN)
    or
    (FP MAXIMUM < FPMINIMUM/ROUND-MIN)
```

MID-BOUND2 can now be witnessed with the constant FP MAXIMUM, which insures that we have not added a contradiction by adding the axiom above.

To aid in our eventual proof, we define **fp-mid** to have several cases. Also, we'll write this program with the implicit assumption that $a \leq b$.

```

fp-mid (a b)
=
if (a < MID-BOUND2) and (b > MID-BOUND2)
  MID-BOUND2
  if (a < 0) and (b > 0)
    0
  if (a < - MID-BOUND2) and (b > - MID-BOUND2)
    - MID-BOUND2
  ROUND (ROUND (a + b) / 2)

```

4.3 close

close(a,b) should return true if it is not guaranteed that $a < \text{fp-mid}(a,b) < b$. Otherwise it should return false so that **find-func-zero** returns the smallest range possible.

From the previous discussion, one case where **close(a,b)** should return true is if $0 \leq a \leq b \leq \text{MID-BOUND2}$, or $-\text{MID-BOUND2} \leq a \leq b \leq 0$. (This is not the strictest possible bound, but it is close and the infrequent gain in performance from using the most-strict bound has been judged to be not worth the much greater complexity.)

We can derive a closeness bound for values not close to 0 from the FP axioms too. We will call this bound **MID-BOUND** and axiomatize it in a manner similar to **MID-BOUND2**. The complex terms in this axiom correspond to the inexactness introduced by applying **fp-mid**.²

Axiom: **fpp (MID-BOUND)**
and
[fpp (x)
and
non-negative (x)
and
FPMAXIMUM >= (x * 2)]
-->
[((2 - (ROUND-MAX * ROUND-MAX)) * x) / (ROUND-MAX * ROUND-MAX)
>= ROUND (MID-BOUND * x)]
and
((ROUND-MIN * ROUND-MIN) * x) / (2 - (ROUND-MIN * ROUND-MIN))
>= ROUND (MID-BOUND * x)]

MID-BOUND may be witnessed by 0 to show that the axiom does not cause an inconsistency.

We may now define **close(a,b)**.

²Facts about the bounds introduced in this axiom are proved. They are derived by applying axioms 11 and 12 from section 3 to the function **fp-mid**.

```

close(a,b)
=
  a >= 0
  and
    a >= round (b * MID-BOUND)
    or
    b <= MID-BOUND2
or
  b <= 0
  and
    b >= round (a * MID-BOUND)
    or
    a >= - MID-BOUND2

```

4.4 unreasonable

unreasonable(a,b) is true if and only if an invariant assumed by another part of the program is violated. The simplest such invariants that we've discussed are that a and b must be **fpp** and $a \leq b$.

As described previously, it is consistent with the FP axioms that $\text{FPMAXIMUM} < \text{FPMINIMUM}/\text{ROUND-MIN}$. A realistic floating-point system would not have this property, so we will add its negation to the definition of **unreasonable** and be willing later to accept a correctness theorem that has this unimportant restriction.

Our earlier choice of the midpoint program made the program vulnerable to "floating-point overflow". This means that a floating-point operation takes place whose corresponding exact result is larger in absolute value than FPMAXIMUM . In this example overflow of the midpoint addition operation might cause the midpoint program to return a value not strictly between the endpoints. Our program is therefore not guaranteed to work "correctly" if $|a + b| > \text{FPMAXIMUM}$. (+ here is exact addition.) A somewhat overly-restrictive but sufficient condition that would assure no overflow is $|a| \leq \text{FPMAXIMUM}/2$ and $|b| \leq \text{FPMAXIMUM}/2$. Since we wish to restrict ourselves to FP numbers in our program to test for this condition, we'll axiomatize an fp value to be less than $\text{FPMAXIMUM}/2$.

Axiom: **fpp** (MID-BOUND3)
and
 $\text{MID-BOUND3} \leq \text{FPMAXIMUM}/2$

We're now in a position to define **unreasonable**.

```

unreasonable(a,b)
=
not fpp(a) or not fpp(b) or a > b
or
FPMAXIMUM < FPMINIMUM/ROUND-MIN
or
RMAGNITUDE(a) > MID-BOUND3 or RMAGNITUDE(b) > MID-BOUND3

```

4.5 Is **find-func-zero** a Floating-Point Program?

In the section on axiomatizing floating-point, sufficient conditions for calling a program a floating-point program were stated. Inspection of **find-func-zero** shows that it meets these conditions in all but one respect. FPMINIMUM/ROUND-MIN is used in an **rlessp** term, and there is no guarantee that this value is **fpp**. However, the entire offending term (described above in the description of **unreasonable**) is a boolean constant for any particular floating-point system. That is, a given floating-point system will have numerical values assigned to FPMINIMUM, ROUND-MIN, and FPmaximum, so the truth of the term can be determined before runtime. This violation of one of our sufficient conditions therefore does not cause the program to fail to be a floating-point program.

Thus, **find-func-zero** can sensibly be called an example of a floating-point program.

5. A Floating-Point Program Correctness Theorem

This section makes an informal argument that the zero-finding program works and presents the PC-NQTHM events that represent the zero-finding program and a mechanically-checked statement about its correctness.

5.1 Correctness Argument Outline

The informal correctness theorem is:

```
not unreasonable (a,b)
and
sign (func (a)) <> sign (func (b))
-->
let (lower, upper) := find-func-zero (a,b)
(sign (func (lower)) <> sign (func (upper)))
and
close (lower, upper))
```

The correctness proof of the zero-finding program has many details. The proof script checked by the prover (not including supporting libraries about arithmetic, rationals, and floating-point numbers) is 169,000 bytes long. The following list of lemmas outlines the essential argument. (The name of the corresponding PC-NQTHM events the proof script are given in parenthesis.)

1. If not **unreasonable**(a,b) and not **close**(a,b), then $a < \text{fp-mid}(a,b) < b$ (**fp-mid-fact1**, **fp-mid-fact2**)
2. **fp-mid**(a,b) returns a floating-point value (**fpp-mid**)
3. **find-func-zero-measure**(a,b,l) expects 3 rational numbers and returns the greatest natural number n such that $(b - a) \geq (n * l)$ (**find-func-zero-measure**)
4. If x, y, and m are **fpp**, and $x < m < y$, then **find-func-zero-measure**(a,b,FPMINSPACE) is reduced if a or b is replaced by m (**fpp-means-find-func-ok**)
5. **find-func-zero** terminates (from previous lemmas)
6. **reasonable**(a,b) --> **reasonable**(a , fp-mid(a,b)) and **reasonable(fp-mid(a,b) , b)**
7. **sign(func(a)) <> sign(func(b)) --> sign(func(lower)) <> sign(func(upper))** [where (lower,upper := **find-func-zero**(a,b))]
8. By induction on **find-func-zero** (justified by lemma 5) and lemmas 6 and 7, the correctness theorem holds. (**find-func-zero-returns-a-zero**,**find-func-zero-returns-close-values**)

5.2 An NQTHM Zero-Finding Program and Correctness Theorem

The zero-finding program and its correctness proof have been formalized in the NQTHM logic. [4] The proof checker enhancement of the theorem prover [12] has accepted the program and its proof. Axiomatized functions were added using constrain events that have been proved to keep the prover state consistent. [5]

The following events have been accepted by the theorem prover. Their acceptance constitutes a mechanical proof of the correctness of the program. (Theorem prover hints have been deleted from the events.)

The zero-finding program is expressed in the logic as

```
(defn fp-mid (x y)
  (if (and (rlessp x (mid-bound2))
            (rlessp (mid-bound2) y))
      (mid-bound2)
    (if (and (rlessp x (rational 0 1))
              (rlessp (rational 0 1) y))
        (rational 0 1)
      (if (and (rlessp x (rneg (mid-bound2)))
                (rlessp (rneg (mid-bound2)) y))
          (rneg (mid-bound2))
        (round (rquotient (round (rplus x y))
                           (rational 2 1)))))))

(defn find-func-zero (a b)
  (if (or (not (rlessp a b)) ; not reasonable(a,b)
          (rlessp (mid-bound3) (rmagnitude a))
          (rlessp (mid-bound3) (rmagnitude b))
          (rlessp (fpmaximum)
                  (rquotient (fpminimum) (round-min))))
      (not (fpp a))
      (not (fpp b)))
    (rational 0 1)
  (if (or (and (numberp (numerator a)) ; close (a,b)
                (or (not (rlessp a
                               (round (rtimes b (mid-bound1))))))
                    (not (rlessp (mid-bound2) b))))
      (and (or (rzerop b)
                (negativep (numerator b)))
          (or (not (rlessp (round (rtimes a (mid-bound1))
                                b)))
              (not (rlessp (mid-bound2) (rneg a)))))))
    (cons a b)
  (let ((mid (fp-mid a b)))
    (if (equal (numberp (numerator (func a)))
               (numberp (numerator (func mid)))))
        (find-func-zero mid b)
      (find-func-zero a mid)))))


```

The correctness theorem is expressed as two NQTHM events.

```
(prove-lemma find-func-zero-returns-close-values (rewrite)
  (implies
    (and ; reasonable(a,b)
         (rlessp a b)
         (not (rlessp (mid-bound3) (rmagnitude a)))
         (not (rlessp (mid-bound3) (rmagnitude b)))
         (not (rlessp (fpmaximum)
                       (rquotient (fpminimum) (round-min))))
           (fpp a)
           (fpp b))
      (let ((lower (car (find-func-zero a b)))
            (upper (cdr (find-func-zero a b))))
        (or (and (numberp (numerator lower)) ; close (lower,upper)
                  (or (not (rlessp lower
                                 (round (rtimes upper (mid-bound1))))))
                      (not (rlessp (mid-bound2) upper))))
            (and (or (rzerop upper)
                      (negativep (numerator upper)))
                (or (not (rlessp (round (rtimes lower (mid-bound1))
                                      upper)))
                    (not (rlessp (mid-bound2) (rneg lower))))))))
```

```
(prove-lemma find-func-zero-returns-a-zero (rewrite)
  (implies
    (and (not (equal (numberp (numerator (func a)))
                      (numberp (numerator (func b))))))
         ; reasonable(a,b)
         (rlessp a b)
         (not (rlessp (mid-bound3) (rmagnitude a)))
         (not (rlessp (mid-bound3) (rmagnitude b)))
         (not (rlessp (fpmaximum)
                      (rquotient (fpminimum) (round-min))))
         (fpp a)
         (fpp b))
    (not (equal (numberp (numerator (func (car (find-func-zero a b)))))))
          (numberp (numerator (func (cdr (find-func-zero a b))))))))
```

6. Future Work

This report describes work in an ongoing project whose ultimate goal is to create in a realistic floating-point number system a substantial floating-point program with a mechanically-checked correctness proof. The mid-point example presented here appears to be the first mechanically-checked floating-point program, and is a step toward that goal.

6.1 Planned Work

Much work is planned.

Proofs about floating-point programs require facts about rational arithmetic operations. The rational library described in section 2 will be greatly enhanced and will allow much faster development of proofs. The development of a useful library of rational number facts is an interesting challenge that may have application beyond proofs about floating point arithmetic.

The floating-point system described in section 3 is less-complex than the floating-point system found on most computers. A realistic floating-point system will be formalized to allow proofs about floating point operations.

The system in section 3 was shown to have a trivial model so as to demonstrate the consistency of the axioms of its formalization. The enhanced floating-point system that will be constructed will be proved to have a model that encompasses a significant portion of the IEEE standard for floating-point arithmetic. [11] This will insure that the axioms describe a consistent system that is realistic. Also, any applications proved correct will be able to execute using the model.

An application program that uses the enhanced floating-point formalization will be proved correct. The example that will probably be pursued is a correctness proof of a program that calculates the sine function. [7] The correctness theorem tightly bounds the forward error of the calculation.

6.2 Other Possible Work

The planned work suggests several possible interesting detours.

The floating-point system model could actually be implemented. This entails writing a compiler for a portion of the logic that includes floating point operations. The target language of the compiler would probably

be Piton [16] so that the resulting code could be run on the verified stack. [3] The appeal of verifying a floating-point program down to the level of hardware is very strong.

It appears that backward error analysis of floating-point programs [18] might benefit from algorithms that automate finding the consistency of sets of inequalities. Such a theorem prover might quickly verify the correctness of the statement of backward error theorems. This possibility will be further explored, and if it seems fruitful such a system will be built.

The techniques developed to verify the floating-point program will hopefully be applicable to other examples. Other programs may be verified to evaluate the general effectiveness of the approach.

Appendix A

An example R2 proof

This appendix presents the output of the theorem prover when given the example problem in subsection 2.3.

```
(DEFN SQUARE (X) (RTIMES X X))

Note that (RATIONAL-FORMP (SQUARE X)) is a theorem.

[ 0.0 0.0 0.1 ]
```

SQUARE

```
(DEFN FOUR-SQUARES
  (A B C D)
  (RPLUS (SQUARE A)
    (RPLUS (SQUARE B)
      (RPLUS (SQUARE C) (SQUARE D)))))

From the definition we can conclude that:
(RATIONAL-FORMP (FOUR-SQUARES A B C D))
is a theorem.
```

```
[ 0.1 0.0 0.0 ]
```

FOUR-SQUARES

```
(LEMMA FOUR-SQ NIL
  (EQUAL (RTIMES (FOUR-SQUARES A B C D)
    (FOUR-SQUARES R S T U))
    (FOUR-SQUARES (RPLUS (RTIMES A R)
      (RPLUS (RTIMES B S)
        (RPLUS (RTIMES C T) (RTIMES D U))))))
    (RPLUS (RTIMES A S)
      (RPLUS (RNEG (RTIMES B R)))
      (RPLUS (RTIMES C U)
        (RNEG (RTIMES D T)))))

    (RPLUS (RTIMES A T O)
      (RPLUS (RNEG (RTIMES B U))
        (RPLUS (RNEG (RTIMES C R)))
        (RTIMES D S)))))

    (RPLUS (RTIMES A U)
      (RPLUS (RTIMES B T O)
        (RPLUS (RNEG (RTIMES C S)))
        (RNEG (RTIMES D R))))))

  ((ENABLE-THEORY R2)
    (ENABLE FOUR-SQUARES SQUARE)))
```

This formula can be simplified, using the abbreviation FOUR-SQUARES, to:

```
(EQUAL
  (RTIMES (RPLUS (SQUARE A)
    (RPLUS (SQUARE B)
      (RPLUS (SQUARE C) (SQUARE D)))))

  (RPLUS (SQUARE R)
    (RPLUS (SQUARE S)
      (RPLUS (SQUARE T) (SQUARE U)))))

(RPLUS (SQUARE (RTIMES A R)
  (RPLUS (RTIMES B S)
    (RPLUS (RTIMES C T0) (RTIMES D U)))))

(RPLUS (SQUARE (RPLUS (RTIMES A S)
  (RPLUS (RNG (RTIMES B R))
    (RPLUS (RNG (RTIMES C T1)
      (RPLUS (RTIMES D T0)))))))

(RPLUS (SQUARE (RPLUS (RTIMES A T0)
  (RPLUS (RNG (RTIMES B U))
    (RPLUS (RNG (RTIMES C R)
      (RPLUS (RNG (RTIMES D S)))))))

(SQUARE (RPLUS (RTIMES A U)
  (RPLUS (RTIMES B T0)
    (RPLUS (RNG (RTIMES C S))
      (RPLUS (RNG (RTIMES D R))))))),
```

which simplifies, applying the lemmas ASSOCIATIVITY2-OF-RTIMES, RTIMES-RPLUS-ARG1, COMMUTATIVITY2-OF-RTIMES, RTIMES-RPLUS-ARG2, ASSOCIATIVITY-OF-RPLUS, COMMUTATIVITY2-OF-RPLUS, COMMUTATIVITY-OF-RTIMES, RTIMES-RNEG, REDUCE-RTIMES, RTIMES-RNEG-ARG2, RTIMES-RPLUS-ARG1, and CORRECTNESS-OF-CANCEL-RPLUS, and unfolding the definition of SQUARE, to:

T.

Q.E.D.

[1 . 8 34 . 7 0 . 2]

FOUR-SQ

T

>

This appendix lists the forms that create the rational number library. Some of the events use proof-checker instructions as hints to the prover [12]. These hints have been removed from this listing in the interest of space.

```
(note-lib "/disk1/home/bevier/libs/integers")
(load "/usr/home/bevier/nthm-init.lsp")
(inform (load "/usr/home/wilting/numerical-arithmetic.lisp"))

; Rational Numbers
; -----
; add-shell rational nil rational-form
;   ((numerater (one-of-number negativep) zero)
;    (denominator (one-of-number zero)))
; -----
; deftheory rational-defns
;   (COUNT-RATIONAL-NUMERATOR-DENOMINATOR-ELIM-EQUAL-DENOMINATOR-LESSEQP
;    RATIONAL-NUMERATOR-DENOMINATOR-RATIONAL-RESTRICTION
;    DENOMINATOR-LESSP-NUMERATOR-DENOMINATOR-RATIONAL-NUMERATOR-LESSEQP
;    NUMERATOR-LESSP-NUMERATOR-TYPE-RESTRICTION
;    DENOMINATOR-LESSP-DENOMINATOR-FORMP-DENOMINATOR-RATIONAL-NUMERATOR-LESSEQP
;    DENOMINATOR-NRATIONAL-FORMP-DENOMINATOR-TYPE-RESTRICTION
;    NUMERATOR-NRATIONAL-FORMP-NUMERATOR-RATIONAL
;    DENOMINATOR-NUMERATOR *1*RATIONAL-FORMP
;    RATIONAL-FORMP)

; definition rationalp (x)
;   (and (rationalp x)
;        (integerp (numerater x))
;        (not (zerop (denominator x))))
;      (quotient (minus (quotient (negative-guts (numerater x))
;                                (denominator x)))
;                (gcd (negative-guts (numerater x))
;                     (denominator x))))))

; definition simple-rplus (x y)
;   (let ((a (fix-rational x)) (b (fix-rational y)))
;     (rational (quotient (numerater x) (gcd (numerater x) (denominator x)))
;               (quotient (denominator x) (gcd (numerater x) (denominator x)))
;                         (itimes (numerater b) (denominator a))
;                               (itimes (denominator a) (denominator b)))))

; definition reduce (x)
;   (if (rationalp x)
;       (if (negativep (numerater x))
;           (rational (minus (quotient (negative-guts (numerater x))
;                                 (denominator x)))
;                     (gcd (negative-guts (numerater x))
;                          (denominator x))))))

; definition fix-rational (x)
;   (if (rationalp x) (rational 0 1))

; definition 0_1()
;   (rational 0 1))
```

Appendix B

The Rational Library

```

(defun rplus (x y)
  (reduce (simple-rplus x y)))

(defun simple-rneg (x)
  (let ((a (fix-rational x)))
    (rational (ineg (numerator a))
              (denominator a)))
  (definition rneg (x)
    (reduce (simple-rneg x)))
  (definition rdifference (a b)
    (rplus a (rneg b)))
  (definition rtimes (x y)
    (reduce (simple-rtimes x y)))
  #|
  (definition simple-rquotient (x y)
    (let ((a (fix-rational x)) (b (fix-rational y)))
      (if (negativep (numerator a))
          (if (negativep (numerator b))
              (rational (itimes (numerator a) (numerator b))
                        (times (denominator a) (denominator b))))
              (rational (tines (negative-guts (numerator a)) (denominator a))
                        (tines (negative-guts (numerator b)) (denominator b)))
              (rational (minus (tines (negative-guts (numerator a)) (denominator a)))
                        (tines (numerator b) (denominator a))))
              (if (negativep (numerator b))
                  (rational (minus (times (numerator a) (denominator b))
                                (times (negative-guts (numerator b)) (denominator a))))
                  (rational (times (numerator a) (denominator b))
                            (times (numerator b) (denominator a)))))))
      (if (negativep (numerator a))
          (if (negativep (numerator b))
              (rational (tines (negative-guts (numerator a)) (denominator a))
                        (tines (negative-guts (numerator b)) (denominator b)))
              (rational (minus (times (numerator a) (denominator b))
                                (times (numerator b) (denominator a)))))))
          (rational (times (numerator a) (denominator b))
                    (times (numerator b) (denominator a)))))))
  |#
  (definition rquotient (x y)
    (reduce (simple-rquotient x y)))
  #|
  (definition rzerop (x)
    (or
      (not (rationalp x))
      (equal (numerator x) 0)))
  |#
  (defn simple-rinverse (x)
    (if (zerop x)
        (rational 0 1)
        (if (negativep (numerator x))
            (rational (ineg (denominator x))
                      (ineg (numerator x)))
            (rational (denominator x) (numerator x)))))

  (definition rinverse (x)
    (reduce (simple-rinverse x)))
  (definition quotient (x y)
    (let ((a (fix-rational x))
          (b (fix-rational y)))
      (if (negativep (numerator a))
          (if (negativep (numerator b))
              (rational (itimes (negative-guts (numerator a)) (denominator a))
                        (itimes (negative-guts (numerator b)) (denominator b)))
              (rational (tines (negative-guts (numerator a)) (denominator a))
                        (tines (negative-guts (numerator b)) (denominator b)))
              (rational (minus (tines (negative-guts (numerator a)) (denominator a)))
                        (tines (negative-guts (numerator b)) (denominator b)))))))
          (if (negativep (numerator b))
              (rational (tines (negative-guts (numerator a)) (denominator a))
                        (tines (negative-guts (numerator b)) (denominator b)))
              (rational (minus (times (numerator a) (denominator b))
                                (times (negative-guts (numerator b)) (denominator a)))))))
          (rational (times (numerator a) (denominator b))
                    (times (numerator b) (denominator a)))))))
      (if (negativep (numerator a))
          (if (negativep (numerator b))
              (rational (tines (negative-guts (numerator a)) (denominator a))
                        (tines (negative-guts (numerator b)) (denominator b)))
              (rational (minus (times (numerator a) (denominator b))
                                (times (numerator b) (denominator a)))))))
          (rational (times (numerator a) (denominator b))
                    (times (numerator b) (denominator a)))))))
  |#
  (prove-lemma integerp-minus (rewrite)
    (implies
      (integerp x)
      (equal integerp (minus x)))
    ((enable integerp)))

  (prove-lemma fix-int-on-integers (rewrite)
    (implies
      (fix-int x)
      (equal (fix-int) (minus x)))
    ((enable fix-int)))

  (prove-lemma itimes-minus-arg1 (rewrite)
    (implies
      (equal
        (itimes (ineg x) y)
        (ineg (itimes x y)))
      (equal (fix-int x) x)))
    ((enable itimes lines)))

  (prove-lemma itimes-minus-arg2 (rewrite)
    (implies
      (equal
        (itimes x (ineg y))
        (ineg (itimes x y)))
      (equal (fix-int y) y)))
    ((enable itimes lines)))

  (prove-lemma integerp-if-negativep-non-zero (rewrite)
    (implies
      (and
        (negativep x)
        (not (equal x (minus 0)))))
      (integerp x)))
    ((enable integerp)))
  
```

```

(prove-lemma integerp-if-numberp (rewrite)
  (implies
    (numberp x)
    (integerp x)
    ((enable integerp)))
  ((enable integerp)))

(prove-lemma itimes-negativep-arg1 (rewrite)
  (implies
    (negativep x)
    (equal
      (itimes x y)
      (ineg (itimes (negative-guts x) y))))
    ((enable itimes)))

(prove-lemma itimes-negativep-arg2 (rewrite)
  (implies
    (negativep y)
    (equal
      (itimes x y)
      (ineg (itimes x (negative-guts y))))))
    ((enable itimes negy)))

(prove-lemma equal-ineg-ineg (rewrite)
  (implies
    (and
      (integerp x)
      (integerp y))
    (equal
      (ineg x) (ineg y))
    (equal x y))
    ((enable integerp ineq)))

(prove-lemma itimes-is-times (rewrite)
  (implies
    (and
      (numberp x)
      (numberp y)
      (equal (itimes x y) (times x y)))
    ((enable itimes)))
    (equal (itimes x y) (times x y)))

(prove-lemma iplus-is-plus (rewrite)
  (implies
    (and
      (numberp x)
      (numberp y)
      (equal (iplus x y) (plus x y)))
    ((enable iplus)))
    (equal (iplus x y) (plus x y)))

;;;;;

(prove-lemma rationalp-reduce (rewrite)
  (rationalp (reduce x)))
  ((disable reduce)))

(prove-lemma rationalp-rplus (rewrite)
  (rationalp (rplus x y))
  ((enable reduce rationalp)))

(prove-lemma rationalp-fix-rational (rewrite)
  (rationalp (fix-rational x)))
  ((enable reduce rationalp))

(prove-lemma rationalp-rtimes (rewrite)
  (rationalp (rtimes x y)))
  ((enable reduce rationalp))

(prove-lemma rationalp-rdifference (rewrite)
  (rationalp (rdifference x y)))
  ((enable reduce rationalp))

(prove-lemma rationalp-rquotient (rewrite)
  (rationalp (rquotient x y)))
  ((enable reduce rationalp))

(prove-lemma rationalp-rmagnitude (rewrite)
  (rationalp (rmagnitude x)))
  ((enable reduce rationalp))

(prove-lemma fix-rational-reduce (rewrite)
  (equal (fix-rational (reduce x))
        (reduce x)))
  ((disable reduce)))

(prove-lemma fix-rational-rplus (rewrite)
  (equal (fix-rational (rplus x y))
        (rplus x y)))
  ((disable reduce)))

(prove-lemma fix-rational-rtimes (rewrite)
  (equal (fix-rational (rtimes x y))
        (rtimes x y)))
  ((disable reduce)))

(prove-lemma fix-rational-rdifference (rewrite)
  (equal (fix-rational (rdifference x y))
        (rdifference x y)))
  ((disable reduce)))

(prove-lemma fix-rational-rquotient (rewrite)
  (equal (fix-rational (rquotient x y))
        (rquotient x y)))
  ((disable reduce)))

(prove-lemma fix-rational-rmagnitude (rewrite)
  (equal (fix-rational (rmagnitude x))
        (rmagnitude x)))
  ((disable reduce)))

```

```

(prove-lemma rational-generalization (generalize)
  (and
    (implies
      (rationalp x)
      (integerp (numerator x)))
    (implies
      (rationalp x)
      (numberp (denominator x)))
    (implies
      (rationalp x)
      (not (zerop (denominator x))))))
  ;;; some divides and gcd facts

;(prove-lemma remainder-times-fact1 (rewrite)
;  (implies
;    (equal (times a b) (times c d))
;    (equal (remainder (times c d) a) 0)))
;
;  (prove-lemma gcd-remainder-fact1 (rewrite)
;    (implies
;      (and
;        (lessp 1 b)
;        (equal (gcd a b) 1))
;        (not (equal (remainder a b) 0)))
;      ((enable gcd)))
;
;    (lemma gcd-remainder-fact2 (rewrite)
;      (implies
;        (and
;          (lessp 1 b)
;          (equal (gcd b a) 1))
;          (not (equal (remainder a b) 0)))
;        ((enable commutativity-of-gcd gcd-remainder-fact1)))
;
;      (defn gcd-times1-induct (x y)
;        (if (zerop y)
;          t
;          (gcd-times1-induct x (sub1 y))))
;
;      (prove-lemma gcd-times1 (rewrite)
;        (equal
;          (gcd x (times x y))
;          (fix x)))
;        ((induct (gcd-times1-induct x y)))))

(lemma gcd-times2 (rewrite)
  (equal
    (gcd y (times x y))
    (fix y)))
  ((disable fix)
   (enable commutativity-of-times)
   (use (gcd-times1 (x y) (y x)))))

(prove-lemma gcd-quotient-quotient
  (rewrite)
  (implies (and (lessp 0 a) (lessp 0 b))
    (equal (gcd (quotient a (gcd a b))
      (quotient b (gcd a b))))))

;;; Following until gcd-remainder-times-fact1-proof done with matt k.

;;;;; (disable IDIFFERENCE-IPPLUS-CANONICALIZER1)
(enable lessp)
(enable idifference)
(enable iplus)
(enable itimes)
(enable ineg)
(enable integerp)
(enable fix-int)
(enable zerop)

(defn gcd-factors (x y)
  ;;; returns a and b s.t. a*x+b*y=gcd. Assumes that x and y are non-zero.
  (cond ((zerop x)
    ;;; 0=0+1*y, which is the gcd of 0 and y
    (cons 0 1))
    ((zerop y)
    ;;; 1*x+0*y=x, which is the gcd of x and 0
    (cons 1 0))
    ((lessp x y)
    ;;; if a*x+b*(y-x) = gcd then (a-b)*x+b*y = gcd
    (let ((factors (gcd-factors x (difference x y) y)))
      (cons (difference (car factors) (cdr factors))
        (cdr Factors)))
    (t
      ;;; so (lseq Y X)
      ;;; if a*(x-y)+b*y = gcd then a*x+(b-a)*y = gcd
      (let ((factors (gcd-factors x (difference x y) y)))
        (cons (car Factors)
          (difference (cdr factors) (car factors)))))

(ordinal-lessp (cons (add1 x) (fix y)))))

(prove-lemma gcd-factors-gives-linear-combination ()
  (let ((factors (gcd-factors x y)))
    (let ((a (car factors))
      (b (cdr factors)))
      (implies (and (numberp x) (numberp y))
        (equal (iplus (itimes a x) (itimes b y))
          (gcd x y))))))

(enable remainder)
(enable quotient)

```

```

(prove-lemma divides-product+reduction
  (rewrite)
    (let ((factors (gcd-factors x y)))
      (let ((a (car factors))
            (b (cdr factors)))
        (implies (and (numberp x) (numberp y))
                  (equal (gcd x y) 1)
                  (equal (remainder (times c y) x) 0))
                  (equal (remainder c x) 0)))
    )
  #|
  ax+by = 1
  c(ax+by) = c
  cax+cy = c
  cax + b(cx) = c where q = (quotient cy x)
  so x divides c
  |
  (lemma dpr-hack1 (rewrite)
    (equal (times x (iplus (itimes a x) (itimes b y))))
           (iplus (itimes c (itimes a x))
                  (itimes b (itimes q x))))
    ((enable-theory integers)))
  (prove-lemma dpr-hack2 (rewrite)
    (set (q (quotient (itimes c y) x)))
    (implies (and (numberp x)
                  (numberp y)
                  (numberp c)
                  (equal (remainder (itimes c y) x) 0))
              (equal (itimes q x)
                    (itimes c y)))
    )
  (lemma dpr-hack3 (rewrite)
    (equal (iplus (itimes c (itimes a x)))
           (itimes b (itimes c y)))
    (iplus (itimes c (itimes a x))
           (itimes c (itimes b y))))
    ((enable-theory integers)))
  (lemma dpr-hack4 (rewrite)
    (equal (iplus (itimes c (itimes a x)))
           (itimes c (itimes c (itimes a x)))
           (itimes c (iplus (itimes a x)
                            (itimes b y)))))
    ((enable-theory integers)))
  (lemma dpr-hack5 (rewrite)
    (implies (and (numberp x)
                  (numberp y)
                  (numberp c)
                  (equal (remainder (itimes c y) x) 0))
              (equal (itimes x (iplus (itimes c a)
                           (itimes b (quotient (itimes c y) x)))) 
                     (itimes c (iplus (itimes a x)
                            (itimes b y))))))
    ((enable dpr-hack1 dpr-hack2 dpr-hack3 dpr-hack4)))
  (prove-lemma remainder-0-sufficiency (rewrite)
    (implies (and (numberp x)
                  (numberp c)
                  (equal (itimes x v) c))
              (equal (remainder c x) 0)
              (equal (remainder (remainder c x) 0) t)))
  )
  ;;;;;; we've proved our important gcd fact - let's set on with it:
  (prove-lemma gcd-remainder-times-fact1-proof
    (implies (equal (gcd a b) 1)
              (implies (equal (remainder (times b c) a) 0)
                      (equal (remainder c a) 0)))
    )
  )
  ;;;;;; we've proved our important gcd fact - let's set on with it:
  (prove-lemma times-gcd-fact
    (rewrite)
    (implies (and (equal (gcd a c) 1)
                  (equal (gcd b d) 1))
              (equal (equal (times a b) (times c d))
                    (and (equal (fix a) (fix d))
                          (equal (fix b) (fix c))))))
  )
  (prove-lemma remainder-0-sufficiency (rewrite)
    (implies (and (numberp x)
                  (numberp c)
                  (equal (itimes x v) c))
              (equal (remainder c x) 0)
              (equal (remainder (remainder c x) 0) t)))
  )

```



```

(prove-lemma equal-reduce1 (rewrite)
  (equal
    (equal (reduce x) y)
    (equal x y))
  ((enable integerp fix-int itimes)
   (enable itimes-negativep-arg1 itimes-negativep-arg2)))
(disable commutativity-of-rplus)

(lemma equal-reduce2 (rewrite)
  (equal
    (equal x (reduce y))
    (equal x y))
  ((use (commutativity-of-rplus (x x) (y (reduce y))))
   (commutativity-of-rplus (x x) (y y))
   (rplus-reduce-arg1 (x y) (y x)))))

(lemma equal-simple-rplus-reduce-arg1 (rewrite)
  (equal
    (equal (rplus x (reduce y))
           (simple-rplus (reduce x) y))
    (simple-rplus x y))
  ((use rplus-reduce-arg1)
   (use rplus-reduce-arg2)
   (enable rplus equal-reduce equal-reduce2)))

(lemma equal-simple-rplus-reduce-arg2 (rewrite)
  (equal
    (equal (rplus (reduce x) (reduce y))
           (simple-rplus x (reduce y)))
    (simple-rplus x y))
  ((use rplus-reduce-arg2)
   (enable rplus equal-reduce equal-reduce2)))

(prove-lemma reduce-reduce (rewrite)
  (equal
    (if (rationalp a)
        (reduce (reduce x)) (reduce x))
    (reduce (rational (iplus (itimes (numerator a) (denominator b))
                            (itimes (numerator b)
                                   (denominator a))))))

(prove-lemma rplus-open-up
  (rewrite)
  (equal (rplus a b)
        (if (rationalp a)
            (if (rationalp b)
                (reduce (rational (iplus (itimes (numerator a) (denominator b))
                                         (itimes (numerator b)
                                                (denominator a)))))

                (itimes (denominator a)
                       (denominator b)))
              (reduce (fix-rational a)))
            (reduce (fix-rational b)))))

(disable rplus-open-up)

(lemma commutativity-of-rplus (rewrite)
  (equal (rplus x y)
        (rplus y x))
  ((enable rplus simple-rplus)
   (enable theory arithmetic integers)))
(disable rplus-requal-arg1 (rewrite)
  (implies
    (rplus a b)
    (reqal a b)
    (reqal (rplus a x))
    (reqal (rplus b x)))
    (enable bines regal rplus-open-up equal-reduce equal-reduce2
           itimes-i-times rational-generalization fix-rational rationalp
           integerp integerp-minus fix-int-on-integers
           correctness-of-cancal-equal-times)
    (enable-theory integers arithmetic rationalp-defn)))
(disable rplus-requal-arg1)

(prove-lemma equal-x-x (rewrite)
  (equal x x))

(lemma rplus-reduce-arg1 (rewrite)
  (equal (rplus (reduce x) y)
        (rplus x y))
  ((enable rplus-requal-arg1 equal-reduce equal-reduce2)))

(lemma equal-rplus-reduce (rewrite)
  (equal
    (equal (rplus (reduce x) y)
           (simple-rneg (reduce x)))
    (reduce (simple-rneg x)))
  ((dissable reduce)))

```

```

(prove-lemma rneg-reduce (rewrite)
  (equal
    (rneg (reduce x))
    (rneg x))
  ((enable rneg simple-rneg-reduce reduce-reduce)))

(prove-lemma simple-rneg-simple-rneg (rewrite)
  (equal
    (simple-rneg (simple-rneg x))
    (simple-rneg (simple-rneg x)))
  ((enable rneg simple-rneg-rneg (rewrite)
    (rneg (rneg x))
    (rneg (rneg x))
    x)
   ((enable rneg rneg-reduce rneg-reduce2
     (simple-rneg simple-rneg simple-rneg-reduce2)))
    (lemma rneg-rneg (rewrite)
      (equal
        (rneg (rneg x))
        (reduce x))
      ((use (requal-rneg simple-rneg simple-rneg-reduce)))
       (enable rneg-reduce-reduce-equal rneg reduce-reduce)))
    (prove-lemma rational-p-means (rewrite)
      (implies
        (rationalp x)
        (and
          (rational-formp x)
          (integerp (numerator x))
          (lisp0 (denominator x)))))

    (prove-lemma means-rationalp
      (rewrite)
      (implies (and (integerp n) (lessp 0 d))
        (rationalp (rational n d)))))

    (prove-lemma rational-rplus-arg1 (rewrite)
      (implies
        (not (rationalp x))
        (equal
          (rplus x y)
          (reduce y)))
        ((disable rationalp)))

    (lemma rational-rplus-arg2 (rewrite)
      (implies
        (not (rationalp x))
        (equal
          (rplus y x)
          (reduce y)))
        ((enable commutativity-of-rplus nrational-rplus-arg1))))))

    ;;; equal-times bridge lemmas

    (prove-lemma rational-simple-rplus-arg1 (rewrite)
      (implies
        (not (rationalp x))
        (equal
          (simple-rplus x y)
          (fix-rational y)))
        ((disable rationalp)))

    (prove-lemma equal-times-bridge1
      (rewrite)
      (implies (and (equal (times a b) (times c d))
                    (equal (times a x) (times c y))
                    (not (zerop a)))
                (equal (equal (times b y) (times d x))
                      t)))
        ))
```

```

(lemma equal-times-bridge2 (rewrite)
  (implies (and (equal (times a b) (times c d))
    (equal (times a x) (times c y)))
    (not (zerop a)))
    (equal (equal (times y b) (times d x)) t))
  ((enable equal-times-bridge commutativity-of-times)))

(lemma equal-times-bridge3 (rewrite)
  (implies (and (equal (times a b) (times c d))
    (equal (times a x) (times c y)))
    (not (zerop a)))
    (equal (equal (times b y) (times x d)) t))
  ((enable equal-times-bridge commutativity-of-times)))

(lemma equal-times-bridge4 (rewrite)
  (implies (and (equal (times a b) (times c d))
    (equal (times a x) (times c y)))
    (not (zerop a)))
    (equal (equal (times y b) (times x d)) t))
  ((enable equal-times-bridge commutativity-of-times)))

(prove-lemma transitivity-of-requal-bridge (rewrite)
  (implied
    (and
      (requal a b)
      (requal b c))
      (requal a c)
      (enable itimes integerp neg)
      (disable rationalp)))
  (disable transitivity-of-requal-bridge))

(lemma transitivity-of-requal (rewrite)
  (and
    implies
    (and
      (requal a b)
      (requal b c))
      (requal a c))
      (requal a c))
  (and
    (requal a b)
    (requal b c))
  (requal a c))
  (implies
    (and
      (requal b a)
      (requal a c))
      (requal b a))
  (requal a c))
  (implies
    (and
      (requal a c))
      (requal a c)))
  (implies
    (and
      (requal b a)
      (requal a c))
      (requal b a))
  (requal a c)))
  ((enable transitivity-of-requal-bridge commutativity-of-requal)))

(prove-lemma equal-requal-rewrite
  (rewrite)
  (and (implies (requal b c)
    (equal (requal a b) (requal a c)))
    (implies (requal b c)
      (equal (requal a b) (requal c a)))
      (implies (requal b c)
        (equal (requal b a) (requal c a))))
    )
  )

(disable equal-requal-rewrite)

```

```

; (lemma rdifference-rdifference-arg2 (rewrite)
;   (equal (rdifference x (rdifference y z))
;          (rdifference (rplus x z) y))
;         ((enable rdifference-commutativity-of-rplus commutativity2-of-rplus)))
;   rneg-rplus rneg-rplus rplus rneg rplus rneg)

; (lemma rplus-rneg-arg1 (rewrite)
;   (equal (rplus (rneg x) y)
;          (rplus (rneg y) x))
;         ((enable rdifference-commutativity-of-rplus)))
;   (enable rdifference-commutativity-of-rplus))

; (lemma rplus-rneg-arg2 (rewrite)
;   (equal (rplus x (rneg y))
;          (rplus x (rneg y)))
;         ((enable rdifference-commutativity-of-rplus)))
;   (enable rdifference-commutativity-of-rplus))

; ;;; times, quotient
; (lemma reduce-rquotient (rewrite)
;   (equal (reduce (rquotient x y))
;          (reduce (rquotient y x)))
;         ((enable rquotient reduce-rtimes)))
;   (enable rquotient reduce-rtimes))

; (lemma reduce-rmagnitude (rewrite)
;   (equal (reduce (rmagnitude x))
;          (rmagnitude x))
;         ((enable rmagnitude reduce-reduce)))
;   (enable rmagnitude reduce-reduce))

; (lemma reduce-rneg (rewrite)
;   (equal (reduce (rneg x))
;          (rneg x))
;         ((enable rneg reduce-reduce)))
;   (enable rneg reduce-reduce))

; (lemma rplus-reduce-arg1-rewrite (rewrite)
;   (equal (rplus (reduce x) y)
;          (rplus x (reduce y)))
;         ((use rplus-reduce-arg1)))
;   (enable rplus-reduce-equal reduce-rplus))

; (lemma rplus-reduce-arg2-rewrite (rewrite)
;   (equal (rplus x (reduce y))
;          (rplus x (reduce y)))
;         ((use rplus-reduce-arg2)))
;   (enable rplus-reduce-equal reduce-rplus))

; (prove-lemma rplus-rmag-simpler-times (rewrite)
;   (equal (rplus (simple-rtimes x y))
;          (simple-rtimes (simple-times x y) y))
;         ((enable commutativity-of-rtimes)))
;   (enable commutativity-of-simple-times))

; ;;; appears to work very slowly - speed it up with concept of rzerop
; ; (lemma associativity-of-simple-times (rewrite)
; ;   (equal (simple-rtimes (simple-times x y) z)
; ;          (simple-rtimes x (simple-times y z)))
; ;         ((enable-theory arithmetic rational-defns)))
; ;   (enable simple-rtimes itimes regual fix-rational fix-int integerp))

; (prove-lemma simple-rtimes-rzerop (rewrite)
;   (implies (rzerop x)
;             (and (equal (numerator (simple-rtimes x y)) 0)
;                  (equal (numerator (simple-rtimes y x)) 0)))
;   (enable simple-rtimes itimes regual fix-rational fix-int integerp rzerop fix-rational))

; (lemma associativity-of-simple-times-when-not-rzerop (rewrite)
;   (implies (and (not (rzerop x))
;                 (not (rzerop y))
;                 (not (rzerop z)))
;             (regual (simple-rtimes (simple-times x y) z))
;             (simple-rtimes x (simple-times y z)))
;           ((enable-theory arithmetic rational-defns)))
;   (enable simple-rtimes itimes regual rational fix-int integerp rzerop fix-rational))

; (lemma rneg-rplus (rewrite)
;   (equal (rneg (rplus x y))
;          (rplus (rneg x) (rneg y)))
;         ((enable rplus-reduce-equal reduce-rneg reduce-rplus)))
;   (enable rplus-reduce-equal reduce-rneg reduce-rplus))

; (prove-lemma rplus-rmag (rewrite)
;   (equal (rplus (rplus x y))
;          (rplus (rplus y x)))
;         ((enable rplus-reduce-equal reduce-rplus)))
;   (enable rplus-reduce-equal reduce-rplus))

; (lemma rdifference-rdifference-arg1 (rewrite)
;   (equal (rdifference (rdifference x y))
;          (rdifference (rplus x z) (rneg y)))
;         ((enable rdifference-commutativity-of-rplus commutativity2-of-rplus)))
;   (enable rdifference-commutativity-of-rplus rneg-rplus))

```

```

(prove-lemma numerator-zero-rzerop-bridge (rewrite)
  (implies
    (equal (numerator x) 0)
    (rzerop x)))
  (rzerop x))

(lemma associativity-of-simple-rtimes (rewrite)
  (requl
    (simple-rtimes (simple-rtimes x y) z)
    (simple-rtimes x (simple-rtimes y z)))
  (use (associativity-of-simple-rtimes-when-not-rzerop))
  (enable requl simple-rtimes-rzerop numerator-zero-rzerop-bridge
    fix-rational itimes)
  (enable-theory rational-defns))

(lemma requl-simple-rtimes-requl-arg1 (rewrite)
  (implies
    (requl x y)
    (rzerop x)
    (simple-rtimes x z)
    (simple-rtimes y z)))
  (enable requl simple-rtimes itimes fix-rational fix-int integerp
    rationalp-means means-rationalp
    correctness-of-cancel-equal-times)
  (enable-theory arithmetic rational-defns))

(lemma requl-simple-rtimes-requl-arg2 (rewrite)
  (implies
    (requl x y)
    (rzerop y)
    (simple-rtimes z x))
  (enable requl simple-rtimes-equal-arg1
    commutativity-of-simple-rtimes))

(lemma requl-simple-rtimes-reduce-arg1 (rewrite)
  (requl
    (simple-rtimes (reduce x) y)
    (simple-rtimes x y)))
  (enable requl simple-rtimes-requal-arg1 requl-x-x requl-reduce1))

(lemma requl-simple-rtimes-reduce-arg2 (rewrite)
  (requl
    (simple-rtimes x (reduce y))
    (simple-rtimes x y)))
  (enable requl simple-rtimes-requal-arg2 requl-x-x requl-reduce1))

(lemma requl-simple-rtimes-bridge (rewrite)
  (and
    (equal
      (requl (simple-rtimes (reduce x) y) z)
      (requl (simple-rtimes x y) z))
    (equal
      (requl (simple-rtimes x (reduce y)) z)
      (requl (simple-rtimes x y) (reduce y)))
    (equal
      (requl z (simple-rtimes x (reduce y)))
      (requl z (simple-rtimes x y)))
    (equal
      (use (requl-simple-rtimes-reduce-arg1)
        (requl-simple-rtimes-reduce-arg2)
        (enable equal-requl-rewrite transitivity-of-requl
          commutativity-of-requl)))))

(lemma requl-associativity-of-rtimes (rewrite)
  (equal
    (rtimes (rtimes x y) z)
    (rtimes x (rtimes y z)))
  (enable rtimes requl-simple-rtimes-bridge

(lemma associativity-of-simple-rtimes requl-reduce requl-reduce2))
  (and
    (equal
      (rtimes (rtimes x y) z)
      (rtimes x (rtimes y z)))
    (use (requl-associativity-of-rtimes)
      (enable reduce-rtimes requl-reduce-reduce-equal1)))

(prove-lemma simple-rplus-rzerop (rewrite)
  (implies
    (rzerop x)
    (and
      (equal (simple-rplus x y) y)
      (equal (simple-rplus y x) y)))
  (use (requl-associativity-of-rtimes)
    (enable reduce-rtimes requl-reduce-reduce-equal1)))

(prove-lemma simple-rtimes-type (rewrite)
  (or
    (numberp (numerator x))
    (negativep (numerator x))))
  (negativp (numerator x)))

(prove-lemma rationalp-simple-rtimes (rewrite)
  (implies
    (rationalp (simple-rtimes x y))
    (rationalp (simple-rtimes y x)))
  (disable rationalp))

(prove-lemma fix-rational-simple-rtimes (rewrite)
  (equal
    (fix-rational (simple-rtimes x y))
    (fix-rational (simple-rtimes y x)))
  (fix-rational (simple-rtimes x y)))

(prove-lemma requl-simple-times-simple-rplus-when-not-rzerop
  (rzerop (numerator x)))
  (implies
    (not (rzerop x))
    (requl (simple-times x y))
    (use (requl simple-times simple-rplus when-not-rzerop)
      (enable requl simple-times simple-rplus when-not-rzerop)
      (and (not (rzerop x))
        (not (rzerop y))
        (disable rationalp)))))

(prove-lemma requl-rplus (rewrite)
  (implies
    (requl (simple-rtimes x y))
    (requl (simple-rtimes y x)))
  (requl (simple-rtimes (simple-rplus x y) z)
    (simple-rplus (simple-rtimes x z)
      (simple-rtimes y z)))
  (fix-rational y)))

(prove-lemma requl-rzerop-simple-rplus
  (rzerop x))
  (implies
    (rzerop x)
    (and
      (requl (simple-rplus x y)
        (fix-rational y))
      (requl (simple-rplus y x)
        (fix-rational y)))))

(prove-lemma commutativity-of-simple-rplus (rewrite)
  (equal
    (simple-rplus x y)
    (simple-rplus y x)))
  (simple-rplus y x))

```

```

(prove-lemma simple-rplus-bridge
  (rewrite)
  (implies (or (zerop x) (zerop y) (zerop z))
           (equal (simple-rtimes (simp1-rplus x y) z)
                  (simple-rplus (simple-rtimes x z)
                               (simple-rtimes y z))))
  )

  (lemma rcommut-rtimes-simple-rplus-arg1 (rewrite)
    (equal (rtimes (rdifference x y) z)
           (rdifference (rtimes x z) (rtimes y z)))
    ((enable rdifference rtimes-rplus-arg1 rtimes-rneg-arg1))

    (lemma rtimes-rdifference-arg2 (rewrite)
      (equal (rtimes x (rdifference y z))
             (rdifference (rtimes x y) (rtimes x z)))
      ((enable rdifference rtimes-rplus-arg2 rtimes-rneg-arg2))

      (prove-lemma rneg-rdifference
        (rewrite)
        (equal (rneg (rdifference x y))
               (rdifference y x)))
      )

      (prove-lemma rdifference-rneg-arg2 (rewrite)
        (equal (rdifference (rneg x) (rneg y))
               (rplus x y)))
      ((enable rplus rneg)))

      (disable rdifference-rneg-arg2)
      (disable rneg-rdifference)
      (disable RTIMES-RDIFFERENCE-ARG2)
      (disable RTIMES-RDIFFERENCE-ARG1)

      (disable RTIMES-RNCG-ARG2)
      (disable RTIMES-RNCG-ARG1)

      (disable RTIMES-RNEG-ARG2)
      (disable RTIMES-RNEG-ARG1)

      (disable REVALU-RTIMES-RNEG)
      (disable REVALU-SIMPLE-RTIMES-SIMPLE-RNEG)

      (disable RTIMES-RPLUS-ARG2)
      (disable RTIMES-RPLUS-ARG1)

      (disable RTIMES-RPLUS-3-BRIDGE)
      (disable REVALU-SIMPLE-RTIMES-SIMPLE-RPLUS-ARG1)

      (disable SIMPLE-RPLUS-BRIDGE)
      (disable COMMUTATIVITY-OF-SIMPLE-RPLUS)

      (disable REVALU-RZEROP-SIMPLE-RTIMES)
      (disable REVALU-SIMPLE-RTIMES-SIMPLE-RPLUS-WHEN-NOT-RZEROP)

      (disable FIX-RATIONAL-SIMPLE-RTIMES)
      (disable RATIONALP-SIMPLE-RTIMES)

      (disable NUMERATOR-TYPE)
      (disable SIMPLE-RPLUS-ZERO)

      (disable ASSOCIATIVITY-OF-RTIMES)
      (disable REVALU-ASSOCIATIVITY-OF-RTIMES)

      (disable REVALU-SIMPLE-RTIMES-BRIDGE)
      (disable REVALU-SIMPLE-RTIMES-REDUCE-ARG2)

      (disable REVALU-SIMPLE-RTIMES-REDUCE-ARG1)
      (disable REVALU-SIMPLE-RTIMES-REDUCE-ARG2)

      (disable REVALU-SIMPLE-RTIMES-REDUCE-ARG1)
      (disable REVALU-SIMPLE-RTIMES-REDUCE-ARG2)

      (disable ASSOCIATIVITY-OF-SIMPLE-RTIMES)
      (disable NUMERATOR-ZERO-RZEROP-BRIDGE)

      (disable SIMPLE-RTIMES-ZERO)
      (disable SIMPLE-RTIMES-RZEROP)

      (disable COMMUTATIVITY-OF-RTIMES)
      (disable RPLUS-RNCG-RTIMES)

      (disable COMMUTATIVITY-OF-SIMPLE-RTIMES)
      (disable RPLUS-RNCG-ARG2)

      (disable RDIFFERENCE-RDIFERENCE-ARG2)
      (disable RNEG-RPLUS)

      (disable RDIFFERENCE-RDIFERENCE-ARG1)
      (disable RNEG-RPLUS)

      (disable RDIFFERENCE-RDIFFERENCE-ARG2)
      (disable RPLUS-RNCG-REDUCE)

      (disable RDIFERENCE-REDUCE-ARG2-REWRITE)
      (disable REDUCE-RNCG)

      (disable REDUCE-RMAGNITUDE)
    )
  )
)

```

```

(disable REDUCE-QUOTIENT)
(disable REDUCE-RTIMES)
(disable REDUCE-DIFFERENCE)
(disable REDUCE-RPLUS)
(disable REDUCE-RMINUS)
(disable REDUCE-REDUCE)
(disable RPLUS-REDUCE-ARG2)
(disable RPLUS-REDUCE-ARG1)
(disable RPLUS-REDUCE-FACT1)
(disable COMMUTATIVITY2-OF-RPLUS)
(disable ASSOCIATIVITY-OF-RPLUS)
(disable REDUAL-ASSOCIATIVITY-OF-RPLUS)
(disable EQUAL-REDUAL-SIMPLE-BRIDGE)
(disable EQUAL-REDUAL-REDUCE)
(disable TRANSITIVITY-OF-REDUAL)
(disable EQUAL-TIMES-BRIDGE4)
(disable EQUAL-TIMES-BRIDGE3)
(disable EQUAL-TIMES-BRIDGE2)
(disable EQUAL-TIMES-BRIDGE1)
(disable EQUAL-ASSOCIATIVITY-OF-SIMPLE-RPLUS)
(disable FIX-RATIONAL-SIMPLE-RPLUS)
(disable RATIONAL-SIMPLE-RPLUS)
(disable NEGATIVE-GUTS-RNEG)
(disable FIX-RATIONAL-OF-RATIONALP)
(disable SIMPLE-RPLUS-FIX-RATIONAL-ARG2)
(disable SIMPLE-RPLUS-FIX-RATIONAL-ARG1)
(disable IRATIONAL-SIMPLE-RPLUS-ARG2)
(disable IRATIONAL-SIMPLE-RPLUS-ARG1)
(disable IRATIONAL-RPLUS-ARG2)
(disable IRATIONAL-RPLUS-ARG1)
(disable MEANS-RATIONALP)
(disable RATIONALP-MEANS)
(disable RNEG-RNEG)
(disable REDUAL-RNEG-RNEG)
(disable SIMPLE-RNEG-SIMPLE-RNEG)
(disable RNEG-REDUCE)
(disable SIMPLE-RNEG-REDUCE)
(disable REDUCE-0)
(disable RATIONAL-INEQ-NUMERATOR-REDUCE-BRIDGE)
(disable NUMBERP-INEQ)
(disable NEGATIVEP-INEQ)
(disable NUMBERP-REDUCE)
(disable RPLUS-REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable RPLUS-OPEN-UP)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable COMMUTATIVITY-OF-REDUAL)
(disable REDUAL-REDUCE-REDUCE-EQUAL)
(disable RPLUS-OPEN-UP)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable COMMUTATIVITY-OF-REDUAL)
(disable REDUAL-REDUCE-REDUCE-EQUAL)
(disable RPLUS-OPEN-UP)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable REDUCE-REDUCE)
(disable QUOTIENT-GCD-TIMES-FACT4)
(disable QUOTIENT-GCD-TIMES-FACT3)
(disable QUOTIENT-GCD-TIMES-FACT2)
(disable EQUAL-TIMES-GCD-BRIDGE1)
(disable QUOTIENT-GCD-TIMES-FACT5)
(disable QUOTIENT-GCD-TIMES-FACT4)
(disable QUOTIENT-GCD-TIMES-FACT3)
(disable QUOTIENT-GCD-TIMES-FACT2)
(disable QUOTIENT-GCD-TIMES-FACT1)
(disable EQUAL-TIMES-TIMES-QUOTIENT-ARG2)
(disable TIMES-GCD-FACT)
(disable DIVIDES-PRODUCT-REDUCTION)
(disable REMAINDER-0-SUFFICIENCY)

```

```

(| old definition changed 8/29
  (definition cancel-rplus (x)
    (rdifference-rneg-arg2 rneg-arg1 rtimes-rdifference rtimes-rdifference-arg2
      rtimes-rdifference-arg1 rtimes-rneg-arg2 rtimes-rneg-arg1 rtimes-rplus-arg2
      rtimes-rplus-arg1 associativity-of-rtimes
      rtimes-commutativity-of-rtimes
      rplus-rneg-arg2 rplus-rneg-arg1 rdifference-rdifference-arg2
      rdifference-rcdifference-arg1 rneg-rplus
      rplus-reduce-arg2-rewrite rplus-reduce-arg1-rewrite
      reduce-rtimes reduce-rmagnitude reduce-rquotient reduce-difference
      commutativity2-of-rplus associativity-of-rplus
      rational-rplus-arg2 rational-p-means means-rationalip
      rneg-rmag reduce-rmagnitude reduce-rquotient reduce-rmagnitude
      rplus-reduce-arg2 rplus-reduce-arg1 rneg-xx rplus-requal-arg1
      commutativity-of-rplus reduce-reduce regulareduce2 regulareduce1
      rational-generalization fix-rtional-magnitude
      fix-rtional-quotient fix-rtional-rtimes fix-rtional-rneg
      fix-rtional-rtimes fix-rtional-fix-rtional fix-rtional-rplus
      fix-rtional-reduce rationalip-quotient
      rationalp-rdifference rationalp-rneg rationalp-primes
      rationalp-fix-rtional rationalp-rplus rationalp-reduce )
    x))

  (definition cancel-rplus (x)
    (rplus-reduce-arg2 rplus-reduce-arg1 rneg-xx rplus-requal-arg1
      commutativity-of-rplus reduce-reduce regulareduce2 regulareduce1
      rational-generalization fix-rtional-magnitude
      fix-rtional-quotient fix-rtional-rtimes fix-rtional-rplus
      fix-rtional-reduce rationalip-quotient
      rationalp-rdifference rationalp-rneg rationalp-primes
      rationalp-fix-rtional rationalp-rplus rationalp-reduce )
    x))

  (definition rplus-tree (x)
    (if (nilistp (list 'rational .. .1))
        (if (nilistp (cdr x)) (list 'reduce (car x)))
            (if (nilistp (cddr x)) (list 'rplus (car x) (cadr x))
                (list 'rplus (car x) (rplus-tree (cdr x)))))))
  #|
  (definition rplus-tree (x)
    (if (nilistp x) '(rational 0 1)
        (if (nilistp (cdr x)) (list 'reduce (car x))
            (if (nilistp (cddr x)) (list 'rplus (car x) (cadr x))
                (list 'rplus (car x) (rplus-tree (cdr x)))))))
  #|
```

(definition rplus-tree (x)

(if (and (listp x) (equal (car x) 'rplus))

(append (rplus-fringe (cdr x)) (rplus-fringe (caddr x)))

(cons x nil)))

|#

(definition rplus-fringe (x)

(if (and (listp x) (equal (car x) 'rplus))

(append (rplus-fringe (cdr x)) (rplus-fringe (caddr x)))

(cons x nil)))

|#

(definition split-by-parity (x)

(if (listp x)

(let ((rest (split-by-parity (cdr x))))
 (if (and (equal (caar x) 'rneg)
 (equal (cddar x) 'rneg)
 (equal (cddar x) nil))
 (cons (car rest) (cons (cadr x) (cdr rest)))
 (cons (cons (car x) (car rest)) (cdr rest)))
 (cons nil nil)))

(definition make-negs (x)

(if (listp x)
 (cons (list 'rneg (car x))
 (make-negs (cdr x)))
 nil)))

```

(prove-lemma eval$-reduce (rewrite)
  (implies
    (equal (car x) 'reduce)
    (equal
      (if (equal a b)
          (subl (occurrences a (delete b x))
                (occurrences a x))
        (occurrences a x)))
      ((disable-theory sets-and-bags)))

(prove-lemma occurrences-delete2 (rewrite)
  (implies
    (equal (occurrences a (delete b x))
      (if (equal a b)
          (subl (occurrences a x))
        (occurrences a x)))
      ((disable-theory sets-and-bags)))

(prove-lemma reduce-eval$-rplus-tree (rewrite)
  (implies
    (equal (reduce (eval$ t (rplus-tree y) a))
      (eval$ t (rplus-tree y) a)))
    ((enable-theory r1)))
  )

(prove-lemma eval$-rplus-tree (rewrite)
  (implies
    (equal (eval$ t (rplus-tree (append x y)) a)
      (rplus (eval$ t (rplus-tree x) a) (eval$ t (rplus-tree y) a)))
    ((enable-theory r1)))
  )

(prove-lemma rplus-evals-rplus-tree (rewrite)
  (implies
    (equal (eval$ t (rplus-tree (append x y)) a)
      (rplus (eval$ t (rplus-tree x) a) (eval$ t (rplus-tree y) a)))
    ((enable-theory r1)))
  )

(prove-lemma member-append (rewrite)
  (implies
    (equal
      (member a (append x y))
      (or
        (member a x)
        (member a y)))
      ((enable-theory r1)))
  )

(prove-lemma member-append (rewrite)
  (implies
    (equal
      (member a (append x y))
      (or
        (member a x)
        (member a y)))
      ((enable-theory r1)))
  )

(prove-lemma delete-append (rewrite)
  (implies
    (equal
      (delete x (append a b))
      (if (member x a)
          (append (delete x a) b)
        (append a (delete x b)))))
      ((enable-theory r1)))
  )

(prove-lemma member-subbagp (rewrite)
  (implies
    (and
      (member a x)
      (subbagp x y)
      (member a y)))
      ((enable-theory r1)))
  )

(prove-lemma subbagp-member (rewrite)
  (implies
    (and
      (member a x)
      (subbagp x y)
      (member a y)))
      ((enable-theory r1)))
  )

;; ; next 10 or so events done with Matt K.

(defn badguy (x y)
  (if (listp x)
    (if (member (car x) y)
        (badguy (cdr x) (delete (car x) y))
      (car x))
    0))

(prove-lemma member-occur (rewrite)
  (implies
    (equal (member a x)
      (lessp 0 (occurrences a x))))
      ((enable-theory r1)))
  )

(prove-lemma permutation-a-b (rewrite)
  (implies
    (permutation (append a b) (append b a)))
      ((enable-theory r1)))
  )

; simpler than in library

```

```

(prove-lemma not-subbagp-not-permutation (rewrite)
  (implies
    (not (subbagp x y))
    (and
      (not (permutation x y))
      (not (permutation y x)))))

(prove-lemma permutation-as-subbagp-helper (rewrite)
  (iff
    (permutation x y)
    (and
      (subbagp x y)
      (subbagp y x)))
    ((use (permutation-as-subbagp-helper)))))

(prove-lemma permutation-as-subbagp (rewrite)
  (equal
    (permutation x y)
    (and
      (subbagp x y)
      (subbagp y x)))
    ((use (permutation-as-subbagp-helper)))))

(disable permutation-as-subbagp-helper)
(disable permutation-as-subbagp)

(prove-lemma subbagp-subbagp-necc (rewrite)
  (implies (subbagp x y)
    (not (lessp (occurrences a y) (occurrences a x))))
  ((enable subbagp-wit-lemma member-occur)
  (disable-theory sets-and-bags)))

(prove-lemma subbagp-transitive
  (rewrite)
  (implies (and (subbagp x y) (subbagp y z))
    (subbagp x z))
  ((use (subbagp-necc a (badguy x z)))
    (y z))
  (subbagp-necc a (badguy x z)))
  (dislable subbagp-necc)
  (dislable-theory sets-and-bags)))

(prove-lemma not-member-make-negs-fact (rewrite)
  (implies
    (not (member x y))
    (equal
      (delete (list 'rneg x) (make-negs y))
      (make-negs y))))
  (dislable theory))

(prove-lemma make-negs-delete (rewrite)
  (equal
    (make-negs (badguy x y))
    (badguy (make-negs x) (make-negs y))))
  (enable theory))

(prove-lemma make-negs-badguy (rewrite)
  (equal
    (make-negs (badguy x y))
    (badguy (make-negs x) (make-negs y))))
  (enable theory))

(lemma regual-rplus-x-simple-rneg-x (rewrite)
  (regual (rplus x (simple-rneg x)))
  ((rational 0 1))
  ((enable theory t1)
  (enable regual-simple-rplus-x-simple-rneg-x rplus
    fix-reduced)))
  (regual-rplus simple-rneg ineq iplus itimes regual
    fix-reduced))

(lemma regual-rplus-x-simple-rneg-x (rewrite)
  (regual (rplus x (simple-rneg x)))
  ((rational 0 1))
  ((enable regual-simple-rplus-x-simple-rneg-x rplus
    fix-reduced)))
  (regual-rplus simple-rneg ineq iplus itimes regual
    fix-reduced))

```

```

(prove-lemma equal-difference-rewrite2 (rewrite)
  (equal (rplus x (rneg x)) (rplus x (rneg x)))
  (rational 0 1))
  ((enable rneg rplus-reduce-arg2-rewrite)
    reduce-rplus *1-reduce)
  (use (reqal-rplus-x-simpler-neg-x)))

(prove-lemma rdifference-x-x (rewrite)
  (equal)
  (rdifference x x)
  (rational 0 1))
  ((enable rdifference)
    (enable-theory rl)
    (disable rplus-rdifference-arg1 rplus-rdifference-arg2)
    rplus-rneg-arg1 rplus-rneg-arg2))
  (and
    (not (lessp b a))
    (not (lessp d c)))
  (equal (fix b) (fix d)))))

(prove-lemma equal-times-bridge5 (rewrite)
  (and
    (not (lessp b a))
    (not (lessp d c)))
  (equal (fix b) (fix d)))))

(prove-lemma equal-times-bridge6 (rewrite)
  (and
    (equal (times b a) (times c d))
    (equal (times x a) (times c y)))
    (not (zerop a)))
  (equal (equal (times b y) (times x d)) t)))
  ((enable equal-times-bridge1)))

(prove-lemma equal-plus-difference-rewrite (rewrite)
  (implies
    (and
      (equal (times b a) (times c d))
      (equal (times x a) (times c y)))
      (not (zerop a)))
    (equal (equal (times b y) (times x d)) t)))
  ((enable equal-times-bridge1)))

(prove-lemma rdifference-rplus-hack (rewrite)
  (rdifference)
  (rplus a b) a)
  (equal (rdifference (rplus a b) rplus-rdifference-arg1)
    rplus-rneg-arg1 rplus-rneg-arg2))
  (reduce b))

  )

(prove-lemma rdifference-rplus-hack2 (rewrite)
  (rdifference)
  (rplus a b) a)
  (reduce b))
  ((enable rdifference-rplus-hack commutativity-of-rplus)))

(prove-lemma rplus-rdifference-hack (rewrite)
  (and
    (equal)
    (rplus a (rdifference b a))
    (reduce b))
  (enable rdifference-rplus-hack)
  (equal)
  (rplus (rdifference b a) a)
  (reduce b))
  ((enable-theory rl)))

(prove-lemma equal-difference-hack1 (rewrite)
  (implies
    (equal a (rdifference b c))
    (equal (rplus c a) (reduce b)))
  ((enable-theory rl1)))

(prove-lemma equal-difference-hack2 (rewrite)
  (implies
    (equal (equal (reduce a) (difference b c))
      (equal (rplus c a) (reduce b)))
  )

(prove-lemma equal-difference-rewrite (rewrite)
  (implies
    (equal (fix b) (fix d))
    (equal (difference a b) (difference c d))
    (or
      (and
        (not (lessp b a))
        (not (lessp b c)))
      (equal a c))))
  (lessp (times v w) (times d z1)))))

(prove-lemma lessp-times-bridge6 (rewrite)
  (implies
    (and
      (lessp (times c v) (times x1 z1))
      (equal (times w x1) (times c d))
      (not (zerop c)))
      (not (zerop d)))
    (lessp (times v w) (times d z1)))))

  )

```

```

(prove-lemma equal-difference (rewrite)
  (and
    (equal a (difference a b))
    (and
      (numberp a)
      (or
        (zerop a)
        (zerop b)))
      (equal (difference a b) a)
      (and
        (numberp a)
        (or
          (zerop a)
          (zerop b))))))
  (prove-lemma equal-simplerplus-x-x-rewrite (rewrite)
    (equal
      (requalequal simple-rplus x y) (simple-rplus x z))
    (requalequal y z))
    ((enable simple-rplus fix-rational rplus itimes
           fix-int-on-integers
           integerp-minus equal-times-bridge1
           equal-times-bridge2 equal-times-bridge3
           equal-times-bridge4)
     (enable-theory r1)))
  (lemma equal-rplus-x-x-rewrite (rewrite)
    (equal
      (requalequal (rplus x y) (rplus x z))
      (equal (reduce y) (reduce z))
      (use (requalequal simple-rplus-x-x-rewrite))
      (enable rplus-reduce-reduce-equal rplus)))
  (prove-lemma eval-rplus-rdifference-hack (rewrite)
    (equal
      (eval (rplus x y) (rdifference (rplus x z) v))
      (eval (reduce y) (reduce z))
      (use (requalequal simple-rplus-x-x-rewrite))
      (enable rdifference-associativity-of-rplus reduce-rplus)))
  (prove-lemma eval-rplus-tree-delete (rewrite)
    (equal
      (eval$ t (rplus-tree (delete m x)) a)
      (if (member m x)
        (difference (eval$ t (rplus-tree x)) a)
        (eval$ t (rplus-tree x) a)))
      (enable-theory r1)))
  (prove-lemma permutation-does-not-affect-rplus (rewrite)
    (implies
      (permutation x y)
      (equal
        (eval$ t (rplus-tree x) a)
        (eval$ t (rplus-tree y) a))
      (enable-theory r1)))
  (prove-lemma evals-rplus-tree-zero (rewrite)
    (equal
      (eval$ t (rplus-tree (append (make-negs x) x)) a)
      (rational 0 1)
      ((enable-theory r1)))
  (prove-lemma evals-rplus-tree-zero (rewrite)
    (equal
      (eval$ t (rplus-tree (append (make-negs x) x)) a)
      (rational 0 1)
      ((enable-theory r1)))
  (prove-lemma last-cdr-delete (rewrite)
    (equal
      (last-cdr (delete e x))
      (last-cdr x)))
  (prove-lemma last-cdr-not-listp (rewrite)
    (implies
      (not (listp x))
      (equal (bagdiff x z)
            (cons x (last-cdr z)))))
  (prove-lemma bagdiff-append-arg1 (rewrite)
    (equal
      (bagdiff (append z x))
      (cons x (last-cdr x))))
  (prove-lemma bagdiff-cons-z-z (rewrite)
    (equal
      (bagdiff (cons x z))
      (cons x (last-cdr x))))
  (prove-lemma bagdiff-not-listp (rewrite)
    (implies
      (not (listp x))
      (equal (bagdiff x z)
            (last-cdr x))))
  (prove-lemma bagdiff-car-in (rewrite)
    (implies
      (and
        (listp x)
        (member (car x) z))
      (equal
        (bagdiff x z)
        (bagdiff (cdr x) (delete (car x) z)))))

  (prove-lemma last-cdr-delete (rewrite)
    (equal
      (last-cdr (delete e x))
      (last-cdr x)))

```

```

(prove-lemma member-subbagp2 (rewrite)
  (implies
    (and
      (subbagp x y)
      (member e x))
    (permutation y z)
    (equal (permutation x y) (permutation x z)))
  ((enable subbagp-wit-lemma member-occur)
   (disable-theory sets-and-subbagp)))

(prove-lemma member-subbagp-delete (rewrite)
  (implies
    (not (member e x))
    (equal (subbagp x (delete e y))
          (subbagp x y)))
  ((enable permutation-as-subbagp)))

(prove-lemma subbagp-bagdiff (rewrite)
  (implies
    (subbagp x y)
    (subbagp (bagdiff x z) (bagdiff y z))))
  ((enable permutation-bagdiff-append-helper (rewrite)
    (and
      (equal (permutation (bagdiff (append z x) x) y)
             (permutation (bagdiff (append x z) x) y))
      (equal (permutation (bagdiff (append x z) x) y)
             (permutation y (bagdiff (append z x) x)))
      (permutation y (bagdiff (append x z) x)))
    (enable equal-permutation permutation-append-arg1-arg2-bridge)))))

#|
(prove-lemma bagdiff-append-arg2 (rewrite)
  (equal
    (bagdiff (append x z) z)
    (if (listp x)
        (append (bagdiff x z) (append (bagint x z) (last cdr z)))
        (last cdr z))))
|#

(prove-lemma permutation-bagdiff-append (rewrite)
  (permutation x (bagdiff (append x z) z)))
  ((enable permutation-bagdiff-append-arg2 (rewrite)
    (and
      (permutation (bagdiff x z) (bagdiff y z))
      (permutation (bagdiff x z) (bagdiff y z)))
    (enable permutation-as-subbagp)))))

(lemma permutation-append-arg1-arg2-bridge (rewrite)
  (permutation
    (bagdiff (append x z) z)
    (bagdiff (append z x) z)))
  ((enable permutation-bagdiff-permutation-a-b)))

(prove-lemma subbagp-transitive-bridge-helper (rewrite)
  (implies
    (and
      (subbagp y z)
      (subbagp z y))
    (and
      (permutation
        (bagdiff (append x z) z)
        (bagdiff (append z x) z))
      (enable permutation-bagdiff-permutation-a-b)))
    (permutation
      (bagdiff (append x z) z)
      (bagdiff (append z x) z)))
  ((enable permutation-bagdiff-permutation-a-b)))

(lemma subbagp-transitive-bridge (rewrite)
  (implies
    (and
      (subbagp y z)
      (subbagp z y))
    (and
      (equal (subbagp y x) (subbagp z x))
      (t)
      (equal (subbagp x y) (subbagp x z))
      (t)))
  ((enable subbagp-transitive-bridge-helper)))))

(prove-lemma rnegs-cancel-list (rewrite)
  (implies
    (equal t (rplus-tree (append x (make-negs x))) a)
    ((rational 0 1)))
  ((enable-theory r1)))

(prove-lemma subbagp-append-simplify1 (rewrite)
  (implies
    (subbagp a x)
    (subbagp a (append x y)))))

(prove-lemma subbagp-permutation-equiv (rewrite)
  (implies
    (permutation x y)
    (and
      (subbagp a x)
      (subbagp a y)))
  ((enable permutation-as-subbagp)))

```

```

(prove-lemma subbagp-append-simplify2 (rewrite)
  (implies
    ((subbagp a x)
     (subbagp a (append y x)))
    ((use (subbagp-permutation-equiv) x (append x y)) (y (append y x))))
   (not (equal (car e) 'rneg)))
  (defn all-rnegs (x)
    (if (listp x)
        (and
         (equal (caar x) 'rneg)
         (not (equal (cadar x) 'rneg))
         (equal (cddar x) nil)
         (all-rnegs (cdr x)))
        t))

  (prove-lemma all-rnegs-make-negs-bagint-fact (rewrite)
    (implies (and (all-rnegs (bagint z w))
                  (member v w)))
              (all-rnegs (make-negs (bagint z (delete v w)))))

  (prove-lemma all-rnegs-make-negs-bagint (rewrite)
    (use (subbagp-permutation-equiv) a (append x y) (append b a)))
   (all-rnegs (make-negs (bagint (car (split-by-parity x))
                                 (cdr (split-by-parity y))))))

  (prove-lemma bagint-all-rnegs-car-split-by-parity (rewrite)
    (implies
      ((and
        (subbagp x b)
        (subbagp y a)
        (subbagp (append x y) (append b a)))
       (use (subbagp-permutation-equiv) (a (append x y)) (x (append b a)))))
     (all-rnegs (make-negs (bagint x (car (split-by-parity x))))))

  (prove-lemma bagint-all-rnegs-split-by-parity (rewrite)
    (implies
      ((and
        (subbagp z y)
        (subbagp z nil)
        (equal (bagint (make-negs z) x) nil)))
       (use (subbagp-permutation-equiv) (y (append a b)))))

  (prove-lemma cancelling-from-rplus (rewrite)
    (implies
      ((and
        (subbagp z x)
        (subbagp z y)
        (equal (bagint (make-negs z) x) nil)))
       (all-rnegs x)
       (equal (bagint x (car (split-by-parity y)))))

  (prove-lemma evals-rplus-tree-rplus-fringe (rewrite)
    (implies
      ((eval$ t (rplus-tree (append (bagdiff x z)
                                      (make-negs (bagdiff y z)))) a)
       (evals t (rplus-tree (append x (make-negs y))) a)))
     (all-rnegs-split-by-parity)

  (prove-lemma evals-rplus-tree-rplus-tree (rewrite)
    (implies
      ((eval$ t (rplus-tree (rplus-fringe x)) a)
       (evals t x)
       (induct (rplus-tree-rplus-fringe x))
       (enable-theory r1)))
     (all-rnegs-split-by-parity)

  (prove-lemma subbagp-x-x (rewrite)
    (subbagp x x))

  (prove-lemma rnegs-not-car-split-by-parity (rewrite)
    (implies
      ((and
        (equal (car x) 'rneg)
        (not (equal (cadar x) 'rneg))
        (equal (cddar x) nil)
        (not (member x (car (split-by-parity y))))))
       (permutation-append-split-by-parity-bridge (rewrite)
         (permutation (append (car (split-by-parity x))
                           (make-negs (cdr (split-by-parity x)))) x)
         )))

  (prove-lemma bagint-all-rnegs-make-negs-bagint)
  (disable bagint-all-rnegs-make-negs-split-by-parity)
  (disable ALL-RNESS-MAKE-NEGS-BAGINT)

```

```

(disable ALL-NEGGS-MAKER-NEGS-BAGINT-FACT)
(disable MEMBER-CDR-SPLIT-BY-PARTITY)
(disable PERMUTATION-APPEND-SPLIT-BY-PARTITY-BRIDGE)
(disable RHEGS-NOT-CAR-SPLIT-BY-PARTITY)
(disable SUBAEGP-X-X)
(disable SUBAEGP-APPEND-BRIDGE2)
(disable SUBAEGP-APPEND-BRIDGE)
(disable SUBAEGP-APPEND-SIMPLIFY2)
(disable SUBAEGP-PERMUTATION-EQIV)
(disable SUBAEGP-PERMUTATION-SIMPLIFY1)
(disable PERMUTATION-BAGDIFP-APPEND-HELPER)
(disable EQUAL-PERMUTATION)
(disable SUBAEGP-TRANSITIVE-BRIDGE)
(disable MEMBER-SUBAEGP2)
(disable LAST-CDR-DELETE)
(disable BAGDIFP-CAR-IN)
(disable PERMUTATION-BAGDIFP-APPEND)
(disable SUBAEGP-BAGDIFP)
(disable MEMBER-SUBAEGP-DELETE)
(disable MEMBER-SUBAEGP2)
(disable EQUAL-DIFFERENCE-BRIDGE)
(disable SUBAEGP-NOT-LISTP)
(disable BAGDIFP-NOT-Z)
(disable BAGDIFP-APPEND-ARG1)
(disable BAGDIFP-X-X)
(disable PERMUTATION-TRANSITIVE)
(disable MEMBER-BAGDIFP-APPEND)
(disable MEMBER-CAR-X-X)
(disable EQUAL-RPLUS-DIFFERENCE-HACK)
(disable EQUAL-RPLUS-X-X-REWRITE)
(disable EQUAL-SIMPLE-RPLUS-X-X-REWRITE)
(disable EQUAL-DIFFERENCE)
(disable LESSP-TIMES-BRIDGE1)
(disable EQUAL-PLUS-DIFFERENCE-REWRITE)
(disable EQUAL-TIMES-BRIDGES)
(disable EQUAL-DIFFERENCE-REWRITE2)
(disable EQUAL-DIFFERENCE-REWRITE)
(disable EQUAL-DIFFERENCE-HACK2)
(disable EQUAL-DIFFERENCE-HACK1)
(disable RPLUS-RDIFFERENCE-HACK)
(disable DIFFERENCE-RPLUS-HACK2)
(disable DIFFERENCE-RPLUS-HACK)
(disable RDIFFERENCE-X-X)
(disable RPLUS-X-NEGGS-X)
(disable REQUAL-RPLUS-X-SIMPLE-RNGS-X)
(disable EQUAL-SIMPLE-RPLUS-X-SIMPLE-RNGS-X)
(disable RDIFFERENCE-REDUCE)
(disable SUBAEGP-MAKE-NEGS)
(disable MEMBER-RHEG-MAKE-NESS)
(disable BAGDIFP-BAGDIFP-ARG2)
(disable APPEND-BAGDIFP-ARG2)
(disable APPEND-BAGDIFP-ARG1)
(disable MAKE-NEGS-BAGDIFP)
(disable MAKE-NEGS-DELETE)
(disable MEMBER-MAKE-NEGS-FACT)
(disable SUBAEGP-TRANSITIVE)
(disable SUBAEGP-NECC)
(disable NOT-SUBAEGP-NOT-PERMUTATION)
(disable PERMUTATION-A-B)
(disable SUBAEGP-PERMUTATION)
(disable SUBAEGP-DELETE-CAR2)
(disable SUBAEGP-DELETE-CAR)
(disable SUBAEGP-DELETE-SAME-MEANS)
(disable SUBAEGP-DELETE-SAME)
(disable SUBAEGP-APPEND)
(disable OCCURRENCES-APPEND)
(disable OCCURRENCES-DELETE2)

(deftheory r2
(
  (rtimes-rneg-arg2 rtimes-rneg-arg1 rtimes-rplus-arg2
   rtimes-rplus-arg1 associativity-of-r-times
   rzerop commutativity-of-r-times
   rdifference-rdifference-arg2
   rdifference-rdifference-arg1 rneg-rplus
   rplus-reduce arg2-rewrite rplus-reduce-arg1-rewrite
   reduce-rng reduce-rmagnitude reduce-quotient reduce-difference
   reduce-times reduce-rplus
   commutativity2-of-rplus associativity-of-rplus
   equal-requal-rewrite transitivity-of-requal fix-rational-of-rationalp
   rational-rplus-arg2 rational-rplus-arg1 rationalp-means-means-rationalp
   rneg-rneg rneg-reduce reduce-0 numberp-numerator-reduce reduce-rationalp
   rplus-reduce arg2 rplus-reduce-arg1 regul-x-x
   regul-reduce regul-reduce2 regul-reduce1
   commutativity-of-rplus reduce-rplus
   commutativity-of-requal rational-generalization fix-rational-magnitude
   fix-rational-rquotient fix-rational-rrdifference
   fix-rational-times fix-rational-fix-rational fix-rational-rplus
   fix-rational-reduce rationalp-magnitude rationalp-rquotient
   rationalp-rdifference rationalp-rng rationalp-times
   rationalp-fix-rational rationalp-rplus rationalp-reduce
   commutativity2-of-r-times
   rdifference correctness-of-cancel-rplus)
  correctness-of-cancel-rplus))

```

Appendix C

Floating-point Axiomatization

This appendix lists the forms that introduce the floating-point axioms in a manner guaranteed to be consistent. Some of the events use proof-checker instructions as hints to the prover [12]. These hints have been removed from this listing in the interest of space.

```
(prove-lemma xequal-rationalp-non-zero-numerator (rewrite)
  (and
    (implies
      (and
        (reqnal x y)
        (not (zerop x)))
      (rationalp y))
     (implies
       (and
         (reqnal y x)
         (not (zerop x)))
       (rationalp y)))
    ((enable-theory r1)
     (enable rationalp requal fix-rational itimes fix-int)))
  )

(prove-lemma lessp-times-bridge-bridge
  (rewrite)
  (implies (and (lessp (times a b) (times c d))
                (lessp (times c x) (times a y))
                (lessp (times b x) (times d y)))
            )
  )

(disable lessp-times-bridge-bridge)

(lemma lessp-times-bridge (rewrite)
  (and
    (implies
      (and
        (lessp (times a b) (times c d))
        (lessp (times c x) (times a y)))
      (lessp (times b x) (times d y)))
    (implies
      (and
        (lessp (times b a) (times c d))
        (lessp (times c x) (times a y)))
      (lessp (times b x) (times d y)))
    (implies
      (and
        (lessp (times a b) (times d c))
        (lessp (times c x) (times a y)))
      (lessp (times b x) (times d y)))
    (implies
      (and
        (lessp (times b a) (times d c))
        (lessp (times c x) (times a y)))
      (lessp (times b x) (times d y)))
    (implies
      (enable commutativity-of-times lessp-times-bridge-bridge)))
  )
```

```

(prove-lemma lessp-times-bridge-bridge2
  (rewrite)
    (implies (and (lessp (times a b) (times c d))
                  (not (lessp (times a y) (times c x)))
                  (not (equal y 0))
                  (numberp y))
              (lessp (times b x) (times d y)))
      ())

  (lemma lessp-times-bridge2 (rewrite)
    (and
      (implies
        (and (lessp (times a b) (times c d))
              (not (lessp (times a y) (times c x)))
              (not (equal y 0))
              (numberp y))
        (lessp (times b x) (times d y)))
        (implies
          (and (lessp (times a b) (times c d))
                (not (lessp (times y a) (times c x)))
                (not (equal y 0))
                (numberp y))
            (lessp (times b x) (times d y)))
          (implies
            (and (lessp (times a b) (times c d))
                  (and (lessp (times a b) (times c d))
                        (not (lessp (times a y) (times x c))))
                  (not (equal y 0))
                  (numberp y))
              (lessp (times b x) (times d y)))
            (implies
              (and (lessp (times a b) (times c d))
                    (and (lessp (times a b) (times c d))
                          (not (lessp (times y a) (times x c)))
                          (not (equal y 0))
                          (numberp y))
                      (lessp (times b x) (times d y)))
                (implies
                  (and (lessp (times a b) (times c d))
                        (not (lessp (times a y) (times c x)))
                        (not (equal y 0))
                        (numberp y))
                    (lessp (times b x) (times d y)))
                  (implies
                    (and (lessp (times a b) (times c d))
                          (not (lessp (times a y) (times c x)))
                          (not (equal y 0))
                          (numberp y))
                        (lessp (times b x) (times d y)))
                      (implies
                        (and (lessp (times a b) (times c d))
                              (not (lessp (times a y) (times c x)))
                              (not (equal y 0))
                              (numberp y))
                            (lessp (times b x) (times d y)))
                          (implies
                            (and (lessp (times a b) (times c d))
                                  (not (lessp (times a y) (times c x)))
                                  (not (equal y 0))
                                  (numberp y))
                                (lessp (times b x) (times d y)))
                                  (implies
                                    (and (lessp (times a b) (times c d))
                                          (not (lessp (times a y) (times c x)))
                                          (not (equal y 0))
                                          (numberp y))
                                            (lessp (times b x) (times d y)))
                                              (implies
                                                (and (lessp (times a b) (times c d))
                                                      (not (lessp (times y a) (times x c)))
                                                      (not (equal y 0))
                                                      (numberp y))
                                                        (lessp (times b x) (times d y)))
                                                          (implies
                                                            (and (lessp (times a b) (times c d))
                                                              (not (lessp (times y a) (times x c)))
                                                              (not (equal y 0))
                                                              (numberp y))
                                                                (lessp (times b x) (times d y)))
                                                                (implies
                                                                  (and (lessp (times a b) (times c d))
                                                                    (not (lessp (times y a) (times x c)))
                                                                    (not (equal y 0))
                                                                    (numberp y))
                                                                      (lessp (times b x) (times d y)))
                                                                      (implies
                                                                        (and (lessp (times a b) (times c d))
                                                                          (not (lessp (times y a) (times x c)))
                                                                          (not (equal y 0))
                                                                          (numberp y))
                                                                            (lessp (times b x) (times d y)))
                                                                            (implies
                                                                              (and (lessp (times a b) (times c d))
                                                                                    (not (lessp (times y a) (times x c)))
                                                                                    (not (equal y 0))
                                                                                    (numberp y))
                                                                                  (lessp (times b x) (times d y)))
                                                                                  (implies
                                                                                    (and (lessp (times a b) (times c d))
                                                                                      (not (lessp (times y a) (times x c)))
                                                                                      (not (equal y 0))
                                                                                      (numberp y))
                        (enable commutativity-of-times lessp-times-bridge-bridge3)))))))

  (prove-lemma rationalp-non-integer (rewrite)
    (implies
      (and (not (integerp (numerator x)))
            (not (integerp (denominator x)))
            ((enable rationalP)
             (disable integerP)))
        (enable commutativity-of-times lessp-times-bridge-bridge3)))))

  (prove-lemma rationalp-zeroP (rewrite)
    (implies
      (and (not (rationalp (numerator x)))
            (not (rationalp (denominator x)))
            ((enable rationalP)
             (disable integerP)))
        (enable commutativity-of-times lessp-times-bridge-bridge2)))))

  (prove-lemma rationalp-not-rational-form (rewrite)
    (implies
      (and (not (rationalp x))
            (not (rationalp y)))
        (enable rationalP)
        (disable integerP)))
      (enable commutativity-of-times lessp-times-bridge-bridge2))))
```

```

(prove-lemma rlessp-transitive (rewrite)
  (and
    (implies
      (rlessp a b)
      (rlessp b c))
    (rlessp a c))
  (implies
    (and (rlessp a b)
         (not (rlessp c b)))
    (rlessp a c))
  (implies
    (and (not (rlessp b a))
         (rlessp a c))
    (rlessp a c))
  (implies
    (and (not (rlessp b a))
         (not (rlessp c b)))
    (not (rlessp c a)))
  (enable rlessp illessp itimes fix-int fix-rational
        rationalp-means rational-generalization))

(prove-lemma rlessp-x-x (rewrite)
  (not (rlessp x x))
  (enable rlessp illessp))

(prove-lemma not-rlessp-if-rlessp (rewrite)
  (implies
    (rlessp x y)
    (not (rlessp y x)))
  (enable rlessp illessp))

(prove-lemma not-rlessp=equal (rewrite)
  (implies
    (rationalp x)
    (equal x y))
  (and
    (not (rlessp x y))
    (not (rlessp y x)))
  (enable rlessp illessp)))

(prove-lemma fix-int-numerator (rewrite)
  (implies
    (rationalp x)
    (equal (fix-int (numerator x))
          (numerator x)))
  (enable fix-int rationalp))

(prove-lemma simple-rplus-0 (rewrite)
  (equal
    (numerator (simple-rplus x (simple-rneg x)))
    0)
  (enable simple-rplus simple-rneg ineq fix-rational
        itimes iplus)
  (enable-theory r1))

(prove-lemma regual-simple-rdifference-x-x (rewrite)
  (regual (rdifference x x) (rational 0 1))
  (enable regual rdifference rplus rmagnitude
        regual-simple-rplus-bridge regual-reduce
        fix-int-on-integers fix-rational-simple-rplus
        rationalp-simple-rplus)
  (enable-theory r1))

(prove-lemma regual-simple-rdifference (rewrite)
  (regual (rdifference x y) (rational 0 1))
  (enable regual rdifference rplus rmagnitude
        regual-simple-rplus-bridge regual-reduce
        fix-int-on-integers fix-rational-simple-rplus
        rationalp-simple-rplus)
  (enable-theory r1))

(prove-lemma rmagnitude-reduce (rewrite)
  (rmagnitude (reduce x) (magnitude x))
  (use (regual-simple-rmagnitude-reduce))
  (enable regual-reduce-reduce-equal rmagnitude))

(prove-lemma numberp-numerator-fix-rational (rewrite)
  (equal
    (numberp (numerator (fix-rational x)))
    (not (rationalp x)))
  (or
    (numberp (numerator (fix-rational x)))
    ((enable fix-rational)))
  (enable-theory r1))

# | (prove-lemma rdifference-x-x (rewrite)
#   (equal (rdifference x x) (rational 0 1))
#   ((use (regual-simple-rdifference-x-x)))
#   (enable reduce-reduce rdifference rplus))

# | (prove-lemma regual-fix-rational (rewrite)
#   (and
#     (regual x (fix-rational x))
#     (regual (fix-rational x) x))
#   ((enable theory r1))
#   (enable regual))

# | (prove-lemma negativep-numerator-reduce (rewrite)
#   (equal
#     (negativep (numerator (reduce x)))
#     (negativep (numerator (fix-rational x))))
#   ((enable reduce fix-rational)))

# | (prove-lemma negativep-numerator-fix-rational (rewrite)
#   (and
#     (negativep (numerator (fix-rational x)))
#     (negativep (numerator (fix-rational x))))
#   ((enable fix-rational)))

# | (prove-lemma negativep-numerator-fix-rational (rewrite)
#   (and
#     (negativep (numerator (fix-rational x)))
#     (negativep (numerator (fix-rational x))))
#   ((enable fix-rational)))

# | (prove-lemma rmag-fix-rational (rewrite)
#   (equal
#     (rmag (fix-rational x))
#     (rmag x))
#   ((enable rmag fix-rational)))

# | (prove-lemma rationalp-fix-rational (rewrite)
#   (implies
#     (not (rationalp x))
#     (equal (rmag (fix-rational x)) (rational 0 1)))
#   ((enable rmag fix-rational)))

# | (prove-lemma regual-simple-magnitude-reduce (rewrite)
#   (regual
#     (rmagnitude (reduce x))
#     (simple-rmagnitude (reduce x)))
#   ((enable simple-rmagnitude)
#   (enable-theory r1)))

# | (prove-lemma rmagnitude-reduce (rewrite)
#   (rmagnitude (reduce x) (magnitude x))
#   (use (regual-simple-rmagnitude-reduce))
#   (enable regual-reduce-reduce-equal rmagnitude))

# | (prove-lemma numberp-numerator-fix-rational (rewrite)
#   (equal
#     (numberp (numerator (fix-rational x)))
#     (not (rationalp x)))
#   (or
#     (numberp (numerator (fix-rational x)))
#     ((enable fix-rational)))
#   (enable-theory r1)))

```



```

| (prove-lemma rlessp-0 (rewrite)
|   (numberp (numerator (fpmaximum)))
|     (use fpmaximum-intro))
|     (enable rlessp illessp fix-int integerp itimes))
|
| (lemma negativep-numerator-fpmaximum (rewrite)
|   (not (negativep (numerator (fpmaximum))))
|     ((use (numberp-numerator-fpmaximum))))
|
|   (prove-lemma numberp-numerator-fpminimum (rewrite)
|     (numberp (numerator fpminimum)))
|     ((use fpminimum-intro)))
|
|   (lemma negativep-numerator-fpminimum (rewrite)
|     (not (negativep (numerator fpminimum)))
|       ((use (numberp-numerator-fpminimum))))
|
|   # (prove-lemma rlessp-rmagnitude (rewrite)
|     (and
|       (equal (rlessp x (rmagnitude y))
|         (if (numberp (numerator y))
|           (rlessp x (fix-rational y))
|             (rlessp x (rneg y))))
|       (equal (rlessp (rmagnitude x) y)
|         (if (numberp (numerator (fix-rational x)))
|           (rlessp (fix-rational x) y)
|             (rlessp (rneg x) y)))
|       ((enable rmagnitude simple-rmagnitude)))
|
|   (lemma rlessp-fix-rational (rewrite)
|     (and
|       (equal (rlessp (fix-rational x) y)
|         (numberp (numerator (fix-rational x))))
|       (negativep (numerator (fix-rational y))))))
|
|   (prove-lemma rlessp-neg-pos (rewrite)
|     (and
|       (numberp numerator-fix-rational-0 (rewrite)
|         (equal (numerator (fix-rational rational 0 x))) 0)
|       ((enable fix-rational)))
|
|   (prove-lemma rlessp-neg-pos (rewrite)
|     (and
|       (numberp (numerator (fix-rational x)))
|       (negativep (numerator (fix-rational y))))))
|
|   (prove-lemma rlessp-neg-pos (rewrite)
|     (and
|       (numberp (numerator (fix-rational x)))
|       (negativep (numerator (fix-rational y))))))
|
|   (prove-lemma rlessp-neg-negates-fact (rewrite)
|     (and
|       (numberp (numerator (fix-rational x)))
|       (not (numberp (numerator (rneg x))))))
|
|   (prove-lemma rmag-negates-fact (rewrite)
|     (and
|       (numberp (numerator (fix-rational x)))
|       (not (equal (numerator (rneg x)) 0)))
|       ((enable rmag-negates-fact)))
|
|   (prove-lemma rmag-positive-fact (rewrite)
|     (and
|       (numberp (numerator (fix-rational x)))
|       (not (equal (numerator (rneg x)) 0)))
|       ((use (rmag-negates-fact)))
|       ((enable rmag-negates-fact)))
|
|   (enable rmag simple-rmag rmag)))

```

```

(prove-lemma nrationalp-r-equal (rewrite)
  (implies
    (not (nrationalp x))
    (and
      (implies
        (r-equal x (rneg y))
        (and
          (iff
            (rlessp z (rneg x))
            (rlessp z y))
          (iff
            (rlessp (rneg x) z)
            (rlessp y z)))
        (implies
          (r-equal (rneg y) x)
          (and
            (r-equal (rneg y) 0)
            (iff
              (rlessp z (rneg x))
              (rlessp z y)))
          (implies
            (r-equal y x)
            (and
              (r-equal (numerator (fix-rational y)) 0)))
            (enable-r-equal)))
        (enable-r-equal fix-rational 0zerop)
        (enable-theory r1)))
      (prove-lemma r-equal-0 (rewrite)
        (implies
          (equal (numerator (fix-rational x)) 0)
          (and
            (equal (r-equal x y)
              (equal (numerator (fix-rational y)) 0)))
            (equal (r-equal y x)
              (equal (numerator (fix-rational y)) 0))))
        (enable-r-equal fix-rational 0zerop)
        (enable-theory r1)))
      (enable not-r-lessp-r-equal rlessp-transitive r-equal-rneg)))
    )
    #+
    |#
    (prove-lemma constrain-brigel (rewrite)
      (implies
        (r-equal (rneg (fpminimum)) x)
        (not (rlessp (fpmaximum) x)))
      (use (rneg-positive-fact (x fpminimum)))
      (dissable rneg-positive-fact)
      (enable fpminimum-intro fpmaximum-intro)))
    |#
    (prove-lemma r-equal-rneg-constants (rewrite)
      (and
        (implies
          (r-equal x (rational 1 1))
          (equal (rneg x) (rational -1 1)))
        (implies
          (r-equal (rational 1 1) x)
          (equal (rneg x) (rational 0 1)))
        (implies
          (r-equal (rneg x) (rational 0 1))
          (equal (rneg x) (rational 0 y)))
        (implies
          (r-equal (rational 0 y) x)
          (equal (rneg x) (rational 1 1)))
        (implies
          (r-equal (rational 0 1) x)
          (equal (rneg x) (rational 0 1)))
        (implies
          (r-equal (rational -1 1) x)
          (equal (rneg x) (rational 1 1)))
        (implies
          ((enable-r-equal simple-rneg ineq))
          (reduce-lzercop itimes))
        (enable-theory r1)))
    )
  )
)

```

```

(prove-lemma rlessp-rneg-constants (rewrite)
  (and
    (equal
      (rlessp (rneg x) (rational 1 1))
      (rlessp (rational -1 1) x))
    (equal
      (rlessp (rational 1 1) (rneg x))
      (rlessp x (rational -1 1)))
    (equal
      (rlessp (rneg x) (rational -1 1))
      (rlessp (rational 1 1) x))
    (equal
      (rlessp (rational 1 1) (rneg x))
      (rlessp x (rational -1 1)))
    (equal
      (rlessp (rneg x) (rational 0 1))
      (rlessp (rational 0 1) y))
    (equal
      (rlessp (rational 0 1) x)
      (rlessp (rational -1 1) (rneg x)))
    (equal
      (rlessp (rational -1 1) (rneg x))
      (rlessp x (rational 1 1)))
    (equal
      (rlessp (rneg x) (rational 1 1))
      (rlessp (rational 1 1) x))
    (equal
      (rlessp (rational 1 1) (rneg x))
      (rlessp x (rational 1 1)))
    (equal
      (rlessp (rneg x) (rational 0 1))
      (rlessp (rational 0 1) y))
    (equal
      (rlessp (rational 0 1) x)
      (rlessp (rational -1 1) (rneg x)))
    (equal
      (rlessp (rational -1 1) (rneg x))
      (rlessp x (rational 1 1)))
    (equal
      (rlessp (rneg x) (rational 0 1))
      (rlessp (rational 0 1) y))
    (equal
      (rlessp (rational 0 1) x)
      (rlessp (rational -1 1) (rneg x)))
    (equal
      (rlessp (rational -1 1) (rneg x))
      (rlessp x (rational 1 1)))
    (enable-rtheory r1))

  #+
  | (prove-lemma rationalp-fpmaximum-fpminimum (rewrite)
    (and
      (rationalp (fpmaximum))
      (rationalp (fpminimum)))
    ((use (fpmaximum-intro) (fpminimum-intro)))))

  (prove-lemma constrain-bridge4 (rewrite)
  (and
    (rlessp (fpminimum (rational 1 1))
    (rlessp (fpminimum) fpmaximum)
    (not (rlessp fpmaximum fpminimum)))
    ((use (fpmaximum-intro) (fpminimum-intro)))))

  (prove-lemma negativep-numerator-rneg-fpmaximum (rewrite)
  (and
    (negativep (numerator (rneg (fpmaximum))))
    (use (fpmaximum-intro))
    (enable-rlessp illessp)
    (enable-rtheory r1)))

  (prove-lemma negativep-numerator-rneg-fpmaximum (rewrite)
  (and
    (negativep (numerator (rneg (fpmaximum))))
    (use (fpmaximum-intro))
    (enable-rlessp illessp)
    (enable-rtheory r1)))

  (prove-lemma constrain-bridge2 (rewrite)
  (implies
    (numberp (numerator x))
    (not (rlessp (rneg fpmaximum) x)))
    ((use (fpmaximum-intro))
    (enable-rtheory r1)))

  #+
  | (prove-lemma regular-rneg-pos-pos (rewrite)
    (implies
      (and
        (numberp (numerator x))
        (numberp (numerator y)))
      (and
        (equal
          (rlessp (rneg x) y)
          (rlessp (rational 0 y)))
        (numberp (numerator z))
        (not (equal (numerator (fix-rational z)) 0)))
        ((rlessp (rneg x) z))
        ((enable-rlessp illessp itimes rneg ineg simple-rneg
          fix-int-on-integers integerp-minus)
          (enable-rtheory r1)))))

  (prove-lemma rlessp-zero-means-negativep (rewrite)
  (implies
    (rlessp x (rational 0 y))
    (and
      (rationalp x)
      (negativep (numerator x)))))

  (prove-lemma regular-rneg-zero (rewrite)
  (implies
    (and
      (equal (numerator (fix-rational x)) 0)
      (equal (numerator (fix-rational y)) 0)))
    (enable-rtheory r1)

  (enable-rlessp simple-rneg)))

```

```

(prove-lemma rtimes-1 (rewrite)
  (implies
    (not (zerop x))
    (and
      (equal (rtimes (rational x x) y)
             (reduce y))
      (equal (rtimes y (rational x x))
             (reduce y)))
    ((enable rtimes simple-rtimes itimes iplus int-on-integers iplus itimes
            fix-rationals)
     (enable-theory r1)))
  #+
  (prove-lemma rtimes-minus-1 (rewrite)
    (implies
      (not (zerop x))
      (and
        (equal (rtimes (rational (minus x) x) y)
               (rneg y))
        (equal (rtimes y (rational (minus x) x))
               (rneg y)))
      ((enable rtimes simple-rtimes itimes fix-rationals reduce fix-int)
       (enable-theory r1)))
  #+
  (prove-lemma numberp-numerator-rneg (rewrite)
    (implies
      (not (zerop x))
      (and
        (numberp (numerator (rneg x)))
        (or (negativep (numerator x))
            (zerop x)))
      ((enable rneg simple-rneg neg)
       (enable-theory r1)))
  #+
  (prove-lemma negativep-numerator-rneg (rewrite)
    (implies
      (not (zerop x))
      (and
        (negativep (numerator (rneg x)))
        (rllessp (rational 0 1) x))
      ((enable rneg simple-rneg neg rllessp illessp itimes)
       (enable-theory r1)))
  #+
  (prove-lemma rplus-rzero-bridge (rewrite)
    (implies
      (and
        (not (zerop v))
        (numberp z))
      (equal (reduce (rational (times v z) (times v w)))
             (reduce (rational z w))))
      ((enable reduce rationalP)
       (enable-theory r1)))
  #+
  (prove-lemma rplus-rzero-bridge2 (rewrite)
    (implies
      (not (zerop v))
      (equal (reduce (rational (minus (times d v))
                           (times v w)))
              (reduce (rational (minus d w))))))
      ((enable reduce rationalP)
       (enable-theory r1)))
  
```

```

(prove-lemma rplus-positive (rewrite)
  (and
    (implies
      (and
        (not (zerop z1))
        (not (zerop a))
        (equal (times w x1) (times c d)))
        (equal (equal (times c d z1)
                      (plus a (times w x1 z1)))
               f))
    )
  )

(prove-lemma constant-bridge9-hack2
  (rewrite)
  (implies
    (and
      (not (zerop z1))
      (not (zerop a))
      (not (zerop c))
      (not (zerop d))
      (equal (times w x1) (times c d)))
      (equal (equal (times c d z1)
                    (difference (times w x1 z1) a))
             f))
  )
)

(prove-lemma constrain-bridge9 (rewrite)
  (implies
    (and
      (equal x z)
      (reqn (rplus x y) z)
      (rreop y))
      ((enable rplus reqn iplus itimes rplus izerop
              fix-int-on-integers integerp-minus)
       (enable-theory r1)))
  )
)

(prove-lemma constrain-bridge9-variants (rewrite)
  (and
    (implies
      (and
        (equal x z)
        (equal z (rplus x y))
        (rreop y)))
      ((enable rplus reqn iplus itimes rplus izerop
              fix-int-on-integers integerp-minus)
       (enable-theory r1)))
  )
)

(prove-lemma constrain-super-hack2
  (rewrite)
  (implies
    (and
      (not (equal v 0))
      (not (equal z 0))
      (rllessp (rational z v)
              (rational 1 1)))
      ((not (reqn (rplus x (rational z v)) 0)))
       (use (constrain-bridge9)))
  )
)

```

;

just so it's most recent, redo requl-x-x proof

(lemma requl-x-x-copy (rewrite)

(requl x x)

((enable requl-x-x))

(enable-theory r1))

;

; silliest lemmas ever-

(prove-lemma constrain-super-hack

(rewrite)
 (implies
 (and
 (numberp v)
 (numberp z)
 (not (equal v 0))
 (not (equal z 0))
 (rllessp (rational z v)
 (rational 1 1)))
 ((not (reqn (rplus x (rational z v)) 0)))
 (use (constrain-bridge9)))
)
)

Appendix D fp-mid and a Correctness Proof

```
(lemma rlessp-0-fpminspace (rewrite)
  (rlessp (rational 0 1) (fpminspace))
  ((enable fpp-round-intro)))

(lemma round-rneg (rewrite)
  (implies
    (and
      (fpp x)
      (numberp (numerator (fix-rational delta)))
      (not (equal (numerator (fix-rational delta)) 0))
      (rlessp delta (fpminspace))
      (not (fpp (*plus x delta))))
    ((enable fpp-round-intro)))
  (disable fpp-round-intro))

(lemma fp-minspace (rewrite)
  (implies
    (and
      (fpp x)
      (numberp (numerator (fix-rational delta)))
      (not (equal (numerator (fix-rational delta)) 0))
      (rlessp delta (fpminspace))
      (not (fpp (*plus x delta))))
    ((enable fpp-round-intro)))
  (enable fpp-round-intro)))
```

This appendix lists the forms that introduce **fp-mid**, a program that finds the zero of a function, and a proof of its correctness. Some of the events use proof-checker instructions as hints to the prover [12]. These hints have been removed from this listing in the interest of space.

```

# ; definitions placed in rational.events - for replay remove

(defn simple-rinverse (x)
  (if (zerop x)
    (rational 0 1)
    (if (negativep (numerator x))
      (rational (ineg (denominator x))
                (ineg (numerator x)))
      (rational (denominator x) (numerator x)))))

(defn rinverse (x)
  (reduce (simple-rinverse x)))

(defnition rquotient2 (x y)
  (rtimes x (rinverse y)))
|#
```

; Some arithmetic facts we're missing

```
(disable rinverses)

(lemma rtimes-reduce (rewrite)
  (and
    (equal (rtimes (reduce x) y)
           (rtimes x y))
    (equal (rtimes x (reduce y))
           (rtimes x y)))
  |#)
```

(use regul-simple-rtimes-reduce-arg1)

```
((use regul-simple-rtimes-reduce-arg1)
  (enable rtimes regul-reduce-reduce-equal1))
```

(lemma equal-reduce-reduce (rewrite)

```
(equal
  (equal (reduce x) (reduce y))
  (equal x y))
  |#)
```

(use (regul-reduce-reduce-equal (a x) (b y))))

```
((use (regul-reduce-reduce-equal (a x) (b y))))
```

(prove-lemma times-quotient-gcd-bridge

```
(rewrite)
  (equal (equal (times a (quotient b (gcd d b)))
               (times c (quotient d (gcd d b)))))
        (or (equal (times a b) (times c d))
            (and (zerop b) (zerop d))))
```

```

(prove-lemma rationalp-rinverse (rewrite)
  (equal (times a (quotient b (gcd b d)))
        (times c (quotient d (gcd b d)))))
  (or
    (equal (times a b)
           (times c d))
    (and (zerop b) (zerop d)))
  ((enable commutativity-of-gcd times-quotient-gcd-bridge)))

(prove-lemma regual-simple-rinverse-reduce-bridge-helper (rewrite)
  (regual (simple-rinverse (reduce x)) (simple-rinverse x))
  ((enable theory r2)
   (enable regual itimes fix-int reduce simple-rinverse
         integerp-minus)))
  ((enable commutativity-of-gcd times-quotient-gcd-bridge)))

(prove-lemma regual-simple-rinverse-reduce-bridge (rewrite)
  (and
    (equal
      (regual (simple-rinverse (reduce x)) y)
      (regual (simple-rinverse x y)))
    (equal
      (regual y (simple-rinverse (reduce x)))
      (regual (simple-rinverse x))))
    ((use regual-simple-rinverse-reduce-bridge-helper)
     (disable simple-rinverse regual-simple-rinverse-reduce-bridge-helper)
     (enable equal-regual-reverse commutativity-of-regual)))
    ((enable regual-simple-rinverse-commutativity-of-regual)
     (regual (simple-rinverse (reduce x)) (simple-rneg x))
     (enable regual itimes fix-int reduce simple-rneg ineq
         integerp-minus)))
  ((enable regual-simple-rneg-reduce-bridge-helper (rewrite)
    (regual (simple-rneg (reduce x)) y)
    (enable theory r2)
    (enable regual itimes fix-int reduce simple-rneg ineq
         integerp-minus)))
  ((prove-lemma regual-simple-rneg-reduce-bridge (rewrite)
    (and
      (equal
        (regual (simple-rneg x) y)
        (regual y (simple-rneg (reduce x)))
        (regual y (simple-rneg x)))
      ((use regual-simple-rneg-reduce-bridge-helper)
       (disable simple-rneg regual-simple-rneg-reduce-bridge-helper)
       (enable equal-regual-reverse commutativity-of-regual)))
      (negativep iplus x y)
      (or
        (negativep iplus x y)
        (and
          (negativep x)
          (lessp y (negative-guts x)))
        (and
          (negativep y)
          (lessp x (negative-guts y))))
      ((enable iplus)))
  ((prove-lemma regual-simple-rinverse-simple-rneg (rewrite)
    (regual
      (simple-rinverse (simple-rneg x))
      (simple-rneg (simple-rinverse x)))
    ((enable simple-rinverse simple-rneg regual itimes fix-int-on-integers
      integerp-minus ineq)
     (enable theory r2)))
  ((prove-lemma rinverse-rneg (rewrite)
    (equal
      (rinverse rneg x))
    (rneg (rinverse x))
    ((enable theory r2)
     (enable rinverse rneg)
     (disable simple-rinverse simple-rneg)))
  ((enable simple-rinverse simple-rneg)))

```

```

(prove-lemma itimes-minus-1 (rewrite)
  (implies
    (lessp 0 x)
    (and
      (equal
        (itimes -1 x)
        (minus x))
      (equal
        (itimes x -1)
        (minus x)))
    ((enable itimes fix-int)))

(prove-lemma iplus-x-minus-z-bridge (rewrite)
  (implies
    (and
      (numberp z)
      (numberp z)
      (not (lessp x z)))
    (equal (iplus x (minus z))
          (difference x z)))
  ((enable iplus))

(prove-lemma numberp-numerator-rtimes (rewrite)
  (equal
    (or
      (rzerop x)
      (rzerop y))
    (numberp (numerator (rtimes x y)))))

(prove-lemma numberp-numerator-rinverse (rewrite)
  (equal
    (numberp (numerator (rinverse x)))
    (numberp (numerator (fix-rational x)))))

(prove-lemma rzerop-rlessp-rewrite (rewrite)
  (implies
    (rzerop x)
    (and
      (equal
        (rlessp x y)
        (illessp 0 (numerator (fix-rational y)))))
    (equal
      (rlessp y x)
      (neglativep (numerator (fix-rational y)))))

(prove-lemma rzerop-rtimes (rewrite)
  (implies
    (rzerop x)
    (and
      (equal
        (rtimes x y) (rational 0 1))
      (equal (rtimes y x) (rational 0 1)))
    ((enable rtimes illessp itimes fix-rational fix-int integerp-minus)))

;;;; The zero-finding problem

(constrain func-intro (rewrite)
  (fpp (func x))
  ((func (lambda (x) (rational 0 1)))))

(defn floor (x)
  (if (and
    (rationalp x)
    (numberp (numerator x))
    (quotient (numerator x) (denominator x))
    0))
    0))

;; some new rational functions

(defn rabs (x)
  (if (neglativep (numerator x))
    (rneg x)
    (reduce x)))

(defn rmin (x y)
  (if (rlessp x y)
    x
    y))

(defn rmax (x y)
  (if (rlessp x y)
    y
    x))

```

```

(defun find-func-zero-measure (a b minSpace)
  (floor (quotient (rdifference b a) minSpace)))

(prove-lemma rlessp-rtimes-rational-1-1 (rewrite)
  (equal (rlessp (rtimes x (rinverse y)) (rational 1 1))
        (or
         (rlessp (magnitude x) (magnitude y))
         (not (equal (numberp (numerator x)) (numberp (numerator y))))))
    ((enable-theory r2))
    (rzerop x)
    (rzerop y))
  (enable rlessp rleqsp rquotient rtimes simple-rtimes
        rinverse simple-rinverse itimes fix-int fix-rational rneg ineq
        simple-rneg))

(prove-lemma floor-rplus-minus-1 (rewrite)
  (equal (floor (rplus x (rational -1 1)))
        (sub1 (floor x)))
  ((enable-theory r2))
  (enable rlessp rleqsp fix-int integerp-minus))
  (enable rlessp rleqsp fix-int integerp-minus))

(prove-lemma floor-0 (rewrite)
  (equal (floor x) 0)
  ((enable-theory r2))
  (enable rlessp rleqsp floor fix-int))
  (enable rlessp rleqsp fix-int integerp-minus))

(prove-lemma rlessp-simple-rtimes-reduce-arg1-bridge (rewrite)
  (equal (rlessp (rtimes (reduce x) y) z)
        (rlessp (rtimes x y) z))
  ((enable-theory r2))
  (enable rlessp simple-rtimes-reduce-arg1 (rewrite)
    (lemma rlessp-simple-rtimes-reduce-arg1 (rewrite)
      (equal (rlessp (simple-rtimes (reduce x) y) z)
            (rlessp (simple-rtimes x y) z))
      ((enable-theory r2))
      (use (rlessp-simple-rtimes-reduce-arg1-bridge)
        (enable rtimes rlessp-reduce)))
    (enable rtimes rlessp-reduce)))
  (enable rlessp-simple-rtimes-reduce-arg1-2 (rewrite)
    (equal (rlessp (simple-rtimes y (reduce x)) z)
          (rlessp (simple-rtimes y x) z))
    ((enable commutativity-of-simple-rtimes)))
  (prove-lemma rlessp-simple-rtimes-reduce-arg2-bridge (rewrite)
    (equal (rlessp z (rtimes (reduce x) y))
          (rlessp z (rtimes x y)))
    ((enable-theory r2))
    (enable rlessp simple-rtimes-reduce-arg2 (rewrite)
      (lemma rlessp-simple-rtimes-reduce-arg2 (rewrite)
        (equal (rlessp z (simple-rtimes (reduce x) y))
              (rlessp z (simple-rtimes x y)))
        ((enable-theory r2))
        (use (rlessp-simple-rtimes-reduce-arg2-bridge)
          (enable rtimes rlessp-reduce2)))
    (prove-lemma lessp-sub1 (rewrite)
      (equal (lessp (sub1 x) x)
            (lessp 0 x)))
    (prove-lemma rlessp-simple-rtimes-reduce-arg2-2 (rewrite)
      (equal (rlessp z (simple-rtimes y (reduce x)))
            (rlessp z (simple-rtimes y x)))
      ((enable-theory r2))
      (use (rlessp-simple-rtimes-reduce-arg2-bridge)
        (enable commutativity-of-simple-rtimes)))
    (prove-lemma rlessp-rtimes-reduce-arg2 (rewrite)
      (equal (rlessp (rtimes y (reduce x)))
            (rlessp z (rtimes y x)))
      ((enable-theory r2))
      (use (rlessp-simple-rtimes-reduce-arg2-bridge)
        (enable rtimes rlessp-reduce2)))
    (prove-lemma fix-int-minus0 (rewrite)
      (implies (zerop x)
                (equal (times x y) 0)))
    (prove-lemma difference-0 (rewrite)
      (implies (equal (difference x 0) (fix x)))
                (equal (negativep fix-int minus x)
                      (fix-int (minus x) 0)))
    (prove-lemma negativep-fix-int-minus (rewrite)
      (implies (equal (fix-int (minus x))
                      (fix-int 0))
                (enable fix-int)))
    (prove-lemma times-zero-2 (rewrite)
      (implies (zerop x)
                (equal (times x y) 0)))
    (prove-lemma fix-int-minus0 (rewrite)
      (implies (zerop x)
                (equal (fix-int minus x)
                      (fix-int (minus x) 0)))))))

```

```

(prove-lemma lessp-difference-rewrite (rewrite)
  (equal (lessp (difference x y))
    (lessp x y))
  (and
    (lessp 0 y)
    (lessp 0 x)))
  ((enable-theory t2) arithmetic rational-defns)

(prove-lemma lessp-difference-difference-arg1-rewrite (rewrite)
  (equal)
    (lessp (difference x y))
    (lessp (difference x z)))
  (and
    (lessp z x)
    (lessp z y)))
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma lessp-difference-difference-arg2-rewrite (rewrite)
  (equal (difference x y)
    (difference x z)))
  (and
    (lessp x z)
    (lessp y z)))
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma lessp-difference-plus-rewrite (rewrite)
  (and
    (equal (difference x y)
      (plus x z))
    (or
      (lessp (difference x y)
        (plus x z))
      (lessp 0 z)
      (and
        (not (zerop x))
        (not (zerop y))))))
    (equal
      (lessp (difference x y)
        (plus z x))
      (lessp 0 z)
      (and
        (not (zerop x))
        (not (zerop y))))))
    (not (zerop y)))))

(prove-lemma equal-lessp-bridge
  (rewrite)
  (implies (and (equal (times z zw) (times d v))
    (numberp z)
    (numberp zw)
    (not (equal z 0)))
    (lessp (times v zl) (times d zx1)))
    (lessp (times v zl) (times z zx1)))
  )
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma lessp-0-from-not-zerop (rewrite)
  (numberp x)
  (not (equal x 0)))
  (lessp 0 x))
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma lessp-times-0-rewrite (rewrite)
  (numberp x)
  (not (equal x 0)))
  (lessp 0 x))
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma lessp-times-0-means-zero (rewrite)
  (numberp x)
  (numberp y)
  (numberp z)
  (numberp zw)
  (not (equal z 0)))
  (lessp (times zl zw) (times d zx1)))
  ((enable-theory t2) arithmetic rational-defns))

(prove-lemma fix-times (rewrite)
  (fix (times x y))
  ((times x y)))
  ((enable-theory t2) arithmetic rational-defns))

```

```

(prove-lemma lessp-rplus-difference-bridge (rewrite)
  (implies (and (fix-rational z))
            (lessp 0 (numerator (fix-rational z)))
            (rllessp (rdifference x z) x)
            (enable rllessp simple-rplus rllessp-reduce1 ictimes 1plus
                   *rllessp-reduce2 negative-numerator-reduce fix-rational implies
                   *rllessp and or integer-fix-int rllessp zerop times-zero-2
                   negative-times-minus rllessp neg ling simple -rlng
                   denominator-0 negative-means-not-llessp llessp difference-rewrite
                   llessp-0-fnon-not-zerop rllessp 0-means-zerop
                   negative-guts-numerator-reduce-x-0 llessp-times-0-rewrite fix-times
                   (enable-theory r2 arithmetic rational-defns)
                   (disable-theory t))
            (prove-lemma find-func-lessp-fact1
              (rewrite)
              (implies (and (lessp 0 (numerator (fix-rational z)))
                            (not (rllessp (rdifference x z) a)))
                        (rllessp (find-func-zero-measure a
                           (rdifference x z)
                           (find-func-zero-measure a x z)))
                        )
              )
            (prove-lemma rlessp-rdifference
              (rewrite)
              (implies (and (equal (rllessp (rdifference x y) z)
                                (rllessp x (rplus y z)))
                            (equal (rllessp z (rdifference x y))
                                   (rllessp (rplus y z) x)))
                        )
            (prove-lemma find-func-lessp-fact2 (rewrite)
              (implies (and
                            (lessp 0 (numerator (fix-rational z)))
                            (not (rllessp a (rplus x z)))
                            (rllessp (find-func-zero-measure (rplus x z) a z)
                                   (find-func-zero-measure x a z)))
                            (enable-theory r2 arithmetic)
                            (disable-theory t)
                            (disable difference)
                            (enable quotient find-func-zero-measure rllessp-difference)
                            (use (find-func-lessp-fact1 (x a) (a x) (z z))))
                          )
            (prove-lemma quotient-zero-arg1 (rewrite)
              (implies
                (zerop x)
                (equal (quotient x y) 0)))
            (prove-lemma times-quotient-remainder-fact (rewrite)
              (implies (and
                            (quotient (remainder b a) 0)
                            (equal a (quotient b a))
                            (times a (quotient b a))
                            (fix b)))
                )
            (prove-lemma quotient-reduce (rewrite)
              (implies (and
                            (quotient (numerатор (reduce x))
                                   (denominator (reduce x)))
                            (quotient (numerатор (fix-rational x))
                                   (denominator (fix-rational x))))
                            (enable reduce fix-rational)))
            (prove-lemma plus-zerop-2 (rewrite)
              (implies
                (zerop x)
                (equal
                  (plus y x)
                  (fix y)))
            (prove-lemma quotient-preserves-lessp (rewrite)
              (implies
                (not (llessp b a))
                (equal (llessp (quotient b c)) (quotient a c) t))
                ((induct (double-remainder-induction a b c)))
              )
            (prove-lemma not-llessp-times-means-not-llessp-quotient
              (rewrite)
              (implies (and (not (llessp (times dy nx) (times dx ny)))
                            (not (zerop dy)))
                        (not (zerop dx)))
                        (not (llessp (quotient nx dx)
                           (quotient ny dy))))
              )
            (prove-lemma not-rllessp-neg-pos (rewrite)
              (implies
                (and
                  (not (rllessp x y))
                  (not (numberp (numerator x)))
                  (numberp (numerator y)))
                  (and
                    (not (rationalp x))
                    (implies (rationalp y) (equal (numerator y) 0)))
                    ((enable-theory r2)
                      (enable rllessp))
                  )
                )
            (prove-lemma rlessp-means-not-llessp-times-bridge (rewrite)
              (implies
                (and
                  (rationalp x)
                  (rationalp y)
                  (numberp (numerator x))
                  (numberp (numerator y)))
                  (not (rllessp x y))
                  (equal (llessp (quotient (numerator x) (denominator x))
                               (quotient (numerator y) (denominator y)))
                         f))
                  ((enable rllessp)
                    (enable-theory r2)
                    (disable not-llessp-times-means-not-llessp-quotient)
                    (use (not-llessp-times-means-not-llessp-quotient
                          (nx (numerator x)) (dx (denominator x))
                          (ny (numerator y)) (dy (denominator y))))))
                )
            
```

```

(prove-lemma lessp-floor (rewrite)
  (implies
    (not (lessp x y))
    (equal (lessp (floor x) (floor y)) t))
  (enable lessp))

(disable rlessp-means-not-lessp-times-bridge)

(prove-lemma rationalp-fpminspace (rewrite)
  (rationalp (fpminspace)))
  ((use (rlessp-0-fpminspace))
   (enable-theory r2))
  (disable rlessp-0-fpminspace))

(prove-lemma numerator-fpminspace-0 (rewrite)
  (equal (numerатор (fpminspace)) 0)
  f)
  ((use (rlessp-0-fpminspace))
   (enable-theory r2))
  (enable lessp)
  (disable rlessp-0-fpminspace))
  (disable rlessp-fpminspace))

(prove-lemma numberp-numerator-fpminspace (rewrite)
  (numberp (numerатор (fpminspace)))
  ((use (rlessp-0-fpminspace))
   (enable-theory r2))
  (enable lessp)
  (disable rlessp-0-fpminspace)))
  (disable rlessp-0-fpminspace))

(prove-lemma simple-rplus-x-simple-rneg-y-0 (rewrite)
  (equal (numerатор (simple-rplus x (reduce y))) 0)
  ((enable-theory r2 rational-defs arithmetc))
  (enable reduce simple-rplus iplus itimes fix-int integerp-minus zero))

negativep-numerator-fix-rational implies
numerator-fix-rational implies and not quotient-gcd-fact
lessp-0-from-not-zero lessp and not times-quotient-gcd-bridge
fix-rational-of-rationals integerp integerp-minus
times-quotient-gcd-bridge lessp-times-quotient-gcd-bridge
negativep-guts-minus times-zero-2 fix-rational)))

(prove-lemma numerator-rplus-x-rneg-y-0 (rewrite)
  (equal (numerатор (rplus y (rneg x))) 0)
  ((enable (numerатор (rplus y (rneg x)))) 0)
  ((enable rplus rneg rationalp-simple-rplus)))

(prove-lemma rlessp-1 (rewrite)
  (implies
    (lessp a b)
    (rlessp (rational a b) (rational 1 1)))
  ((enable-theory r2))
  (enable rlessp ilessp itimes iplus fix-int fix-rational)))

(prove-lemma rlessp-numberp-numerator-rplus (rewrite)
  (implies
    (and
      (numberp (numerатор (fix-rational x)))
      (numberp (numerатор (fix-rational y)))
      (numberp (numerатор (rplus x y)))))
    (enable rlessp ilessp times-bridge)

(prove-lemma lessp-rplus-times-bridge
  (rewrite)
  (implies (and (numberp c)
    (numberp x1)
    (numberp (numerатор (fix-rational y)))
    (numberp (numerатор (rplus x y)))))
    (not (equal z 0))
    (numberp z)
    (not (lessp v c)))
    (equal
      (lessp (plus (times x1 z) (times c d z)))
      (plus (times d w)
        (times d z)
        (times d v z)))
    t))

(disable-theory t))

(prove-lemma quotient-gcd-fact (rewrite)
  (implies
    (not (zerop x))
    (and (lessp 0 (quotient x (gcd n y)))
      (lessp 0 (quotient x (gcd y x))))))
  )

(prove-lemma lessp-times-quotient-gcd-bridge
  (rewrite)
  (and (equal (lessp (times a (quotient b (gcd b d)))
    (times c (quotient d (gcd b d))))))
    (lessp (times a b) (times c d))
    (equal (lessp (times a (quotient b (gcd d b)))
      (times c (quotient d (gcd d b))))))
    (lessp (times a b) (times c d)))))

)

```

```

(prove-lemma lessp-floor-t-bridge (rewrite)
  (implies (and (numberp (numerator a))
    (not (rllessp b (rplus (rational 1 1) a)))
    (rationalp b)
    (numberp (numerator b))
    (rationalp a))
    (equal (lessp (quotient (numerator a)
      (denominator a))
      (quotient (numerator b)
        (denominator b)))
      t))
  ((enable-theory r2 arithmetic rational-defns)
  (disable-theory t)
  (enable rllessp simple-rplus rplus itimes fix-int inteserp
    rllessp-rllessp-reduce2 not implies and zeroop
    lessp-plus-times-bridge)))

(prove-lemma rllessp-rplus-rneg (rewrite)
  (and
    (equal
      (rllessp (rplus (rneg x) y) z)
      (rllessp y (rplus x z)))
    (equal
      (rllessp (rplus y (rneg x)) z)
      (rllessp y (rplus x z)))
    (equal
      (rllessp z (rplus (rneg x) y))
      (rllessp (rplus (rplus x z) y)))
    (rllessp z (rplus y (rneg x)))
    (rllessp (rplus (rplus x z) y)))
  ((use rllessp-rplus-rplus (c x) (x (rplus (rneg x) y)) (y z)))
    (rllessp-rplus-rplus (c x) (y (rplus (rneg x) y)) (x z)))
  (disable rllessp-rplus-rplus)
  (disable-theory t)
  (enable rllessp-reduce rllessp-reduce2 implies and)
  (enable-theory r2))

(prove-lemma rllessp-rneg (rewrite)
  (and
    (equal
      (rllessp (rneg x) y)
      (rllessp (rplus 0 1) (rplus x y)))
    (equal
      (rllessp x (rneg y))
      (rllessp (rplus x y) (rational 0 1)))
    ((use rllessp-rplus-rplus (c x) (x (rneg x)) (y y)))
      (rllessp-rplus-rplus (c y) (x x) (y (rneg y))))
  (disable rllessp-rplus-rplus)
  (enable-theory r2)
  (enable rllessp-reduce rllessp-reduce2 implies and)))

; redo this so that it fires first
(prove-lemma rllessp-rplus-x-x (rewrite)
  (equal
    (rllessp (rplus x y) (rplus x z))
    (rllessp y z)))
;
```

```

(disable move-riessp-args-right-correct)

(prove-lemma lessp-floor-t-better
  (rewrite)
  (implies
    (and (not (riessp x y))
         (numberp (numerator x)))
    ((enable-theory r2) (enable rzerop)))

  (prove-lemma lessp-floor-rplus (rewrite)
    (implies
      (and (not (riessp x y))
           (numberp (numerator (rplus x (rneg y))))))
      ((not (riessp x y)))))

  (prove-lemma numberp-numerator-rplus-x-rneg-y-better
    (rewrite)
    (implies
      (and (numberp (numerator (rplus x (rneg y)))))
      ((not (riessp x y)))))

  (prove-lemma riessp-floor-bridge1 (rewrite)
    (implies
      (and (not (riessp x y))
           (equal (remainder (numerator x) (denominator x)) 0)
           (equal (remainder (numerator y) (denominator y)) 0)
           (equal (lessp (floor x) (floor y)) f1)
           ((enable-theory r2) (enable rzerop)))
      ((not (riessp (floor x) (floor y)) f1)
       (negativep-numerator-fix-rational itimes floor implies lessp
         quotient-zero-argl zerop rationalp-zero rationalp-non-integer
         not and integerp correctness-of-cancel-lessp-times fix)
       (enable-theory r2 arithmetic rational-defns)
       (enable-theory t2)
       (disable-theory t))))))

  (lemma riimes-monotonic (rewrite)
    (implies
      (and (numberp (numerator a))
           (not (riessp x y)))
      ((and
        (not (riessp (rtimes a x) (rtimes a y)))
        (not (riessp (rtimes a x) (rtimes y a)))
        (not (riessp (rtimes x a) (rtimes a y)))
        (not (riessp (rtimes x a) (rtimes y a))))))
      ((enable-theory r2)
       (enable riessp-rtimes-x-x fix-rational)))
      (enable-theory r2)
      (enable riessp-rtimes-x-x fix-rational))))))

  (lemma riessp-rneg-rneg (rewrite)
    (implies
      (and (not (riessp x y))
           (numberp (numerator b)))
      ((equal (riessp (rneg x) (rneg y))
             (riessp y x))
       ((enable-theory r2 arithmetic rational-defns)
        (enable riessp rneg simple-rneg itimes fix-rational itimes
          riessp-reduce1 riessp-reduce2 integerp fix-int)))
      (enable find-func-zero-measure-monotonic (rewrite)
        (implies
          (and
            (not (riessp x y))
            (numberp (numerator b)))
            ((equal (riessp (find-func-zero-measure a x b)
                           (find-func-zero-measure a y b))
                   (equal (riessp (find-func-zero-measure a b)
                                 (find-func-zero-measure x a b)) f)))
            ((enable-theory r2)
             (enable lessp-rplus find-func-zero-measure rtimes-monotonic
               riessp-reduce1 riessp-reduce2 integerp fix-int)
             (enable numberp-numerator-fix-rational)
             (numberp-numerator-fix-rational))))))

  (lemma riessp-means-not-equal (rewrite)
    (implies
      (and (riessp a b)
           (not (equal a b))
           ((enable-theory r2 arithmetic rational-defns)
            (enable riessp rneg simple-rneg itimes fix-rational itimes
              riessp-reduce1 riessp-reduce2 integerp fix-int))))))

```

```

(lemma fp-rationalp (rewrite) ; was axiom - changed 3-26-90
  (implies
    (fpp x)
    (rationalp x))
  (enable fpp-round-intro))

(prove-lemma rationalp-round (rewrite)
  (rationalp (round x)))

(prove-lemma round-0 (rewrite)
  (and
    (implies
      (fpp x)
      (equal (numerator x) 0))
    (equal (numerator (round x)) 0))
  (use (not-round-down-past y (rational 0 1)) (x x))
  (not-round-up-past y (rational 0 1)) (x x)))
  (enable theory r2)
  (enable rationalp illessP))

(prove-lemma fix-rational-round (rewrite)
  (equal
    (fix-rational (round x))
    (round x))
  (enable fix-rational)))

(lemma rtimes-rinverse (rewrite)
  (equal (rtimes (rinverse x) (rtimes y))
    (rinverse (rtimes x y)))
  (enable-theory r2 rational-defns arithmetic)
  (enable rtimes simple-rinverse simple-rtimes reduce itimes
    fix-int-ing equal-reduce-reduce real
    equal-simple-rtimes-bridge
    fix-rational numerator-fix-rational-0 rationalp-0)))
  (enable theory r2))

(lemma reduce-rinverse (rewrite)
  (equal
    (reduce (rinverse x))
    (rinverse x))
  (enable rinverse)
  (enable-theory r2)))

(prove-lemma rtimes-rinverse-hack
  (rewrite)
  (and (equal (rtimes a (rinverse (rtimes a b))))
    (if (zerop a)
      (rationalp 0 1)
      (rinverse b))
    (equal (rtimes (rinverse a) c))
      (if (zerop b)
        (rtimes (rinverse (rtimes b a)) c))
        (rationalp 0 1)
        (rtimes (rinverse a) c)))))

(prove-lemma rtimes-rinverse-hack3
  (rewrite)
  (and (equal (rtimes (rinverse a) b))
    (if (zerop b)
      (rationalp 0 1)
      (rtimes (rinverse a) b))
    (equal (rtimes b (rinverse a) c))
      (if (zerop b)
        (rtimes (rinverse (rtimes b a)) c))
        (rationalp 0 1)
        (rtimes (rinverse a) c)))))

(prove-lemma reduce-rzerop (rewrite)
  (implies
    (zerop x)
    (equal (reduce x) (rational 0 1)))
  (enable reduce zzerop))

(prove-lemma equal-numerator-rinverse-0 (rewrite)
  (equal
    (equal (numerator (rinverse x)) 0)
    (enable rzerop rinverse simple-rinverse neg)
    (enable-theory r2)))

(prove-lemma rinverse-0 (rewrite)
  (implies
    (zerop x)
    (equal (rinverse x) (rational 0 1)))
  (enable rinverse simple-rinverse reduce-0)))

(prove-lemma rzerop-rneg (rewrite)
  (equal
    (rzerop (rneg x))
    (equal (rzerop x) (rneg x)))
  (enable rzerop rinverse simple-rinverse neg)

(prove-lemma rzerop-rinverse (rewrite)
  (equal
    (rzerop (rinverse x))
    (rzerop x)))
  (enable-theory r2)
  ((enable-theory r2) (enable rzerop rneg simple-rneg neg)))
  (enable rzerop rinverse simple-rinverse neg)))

```

```

(lemma rzerop-rtimes-rewrite (rewrite)
  (equal
    (rzerop (rtimes a b))
    (or
      (rzerop)
      (rzerop b)))
  ((enable-theory r2 rational-defs arithmetic-integers)
   (enable rzerop rtimes simple-rtimes itimes fix-int-on-integers
         integerp-ninjs rationalp-fix-rational reduce-0
         *fix-int-negativep-numerator-fix-rational integerp
         rationalp-not-rational-formp numerator-reduce-0))

  (prove-lemma rlessp-rinverse-hack
    (rewrite)
    (implies (and (equal (numberp (numerator x))
                         (numberp (numerator y)))
                  (not (rzerop x))
                  (not (rzerop y)))
                  (equal (rlessp (trivese x) y)
                         (rlessp (trivese y) x)))
              )
    )

  (lemma round-min-bounds (rewrite)
    (and
      (not (rlessp (rational 1 1) (round-min)))
      (rlessp (rational 99 100) (round-min)))
    ((enable fpp-round-intro)))
  (lemma round-max-bounds (rewrite)
    (and
      (not (rlessp (round-max) (rational 1 1)))
      (rlessp (round-max) (rational 101 100)))
    ((enable fpp-round-intro)))
  (prove-lemma rzerop-round-max (rewrite)
    (not (rzerop (round-max)))
    ((use (round-max-bounds))
     (enable rzerop)
     (disable round-max-bounds)))
  (lemma lessp-times-3 (rewrite)
    (implies
      (and
        (not (lessp a c))
        (not (lessp b d)))
        (equal (lessp (times a b) (times c d)) f)))
  (prove-lemma rlessp-rtimes-simple-hack
    (rewrite)
    (implies (and (numberp (numerator (fix-rational a)))
                  (numberp (numerator (fix-rational b)))
                  (numberp (numerator (fix-rational c)))
                  (numberp (numerator (fix-rational d)))
                  (not (rlessp c a))
                  (not (rlessp d b)))
                  (not (rlessp (rtimes c d) (times a b)))))

  (prove-lemma rlessp-rtimes-round-max (rewrite)
    (numberp (numerator (round-max)))
    ((use (round-max-bounds))
     (enable rlessp itimes)
     (enable-theory r2)))
  (prove-lemma rationalp-round-min (rewrite)
    (rationalp (round-min))
    ((use (round-min-bounds))
     (enable round-min-bounds)))

  (lemma numberp-numerator-round-min (rewrite)
    (numberp (numerator (round-min)))
    ((use (round-min-bounds))
     (enable round-min-bounds times-add1)
     (enable rlessp itimes numberp-numerator-fix-rational
           negativep-numerator-fix-rational *fix-rational
           fix-int-numerator fix-int-on-integers integerp-minus
           rationalp-round-min fix-rational-of-rationalp
           negative-guts-numerator-0
           lessp-times-0-rewrite)
     (enable-theory r2 arithmetic rational-defns)))
  (prove-lemma round-max-hack
    (rewrite)
    (implies (numberp (numerator a))
              (not (rlessp (rtimes (trivese (times (round-max) (round-max))) a))
                    )
    )
  (PROVE-LEMMA ROUND-MIN-HACK
    (REWRITE)
    (implies (numberp (numerator (RPLUS (RNEG (RTIMES (ROUND-MIN) (ROUND-MIN)) (RATIONAL 2 1)))) a))
              (not (rlessp (rtimes (trivese (times (rational 2 1) s)) a)))
    )
  )

```

```

(constrain mid-bound1-intro (rewrite)
  (and
    (fpp (mid-bound1))
    (implies
      (and
        (fpp x)
        (numberp (numerатор x))
        (not (rllessp (fpmaximум) (rtimes (ratиональ 2 1) x))))
      (and
        (not (rllessp (rquotient (rtimes (rdifference (ratиональ 2 1)
          (rtimes (round-max) (round-max)))) x)
          (rtimes (round-max) (round-max)))))))
        (not (rllessp (rquotient (rtimes (round-min) (round-min))
          (round (rtimes (mid-bound1) x)))))))
      (not (rllessp (rquotient (rtimes (rtimes (round-min) (round-min)
        (rtimes (round-max) (round-max)))) x)
        (rdifference (ratиональ 2 1)
          (rtimes (rtimes (mid-bound1) x)))))))
        (round (rtimes (mid-bound1) x)))))))
      ((mid-bound1 (lambda nil (ratиональ 0 1)))))))
    (enable-theory r2 ground-zero rational-defns)
    (disable-theory t)
    (enable rquotient rzero-rtimes fpp-0 round-0 rllessp-0
      fix-rational-round negativep-numerator
      numberp-numerator-rtplus-x-rneg-y-better reduce-rational-p
      rtimes-r-inverse-hack2 rllessp-x-x rtimes-r-inverse-hack
      rtimes-r-inverse-hack3 reduce-rzero rllessp-reduce
      rllessp-reduce2 rationalp-rllessp *1*invverse
      numberp-numerator-rtimes numberp-numerator-rtinverse
      numberp-numerator-round-min round-min-hack
      rationalp-not-rational-form round-max-hack
      (disable rtimes-r-inverse-invverse)))
    (disable rtimes-r-inverse-invverse)))

(prove-lemma numberp-numerator-round (rewrite)
  (implies
    (and
      (rationalp x)
      (numberp (numerатор x))
      (numberp (numerатор (round x))))
      ((use (not-round-down-past (y (ratиональ 0 1)) (x x)))
        (enable not-round-down-past)
        (enable-theory r2)))
    (enable rlessp-rlessp-requal (rewrite)
      (implies
        (and
          (not (rllessp a b))
          (equal a b))
        (enable rllessp rllessp requal itimes)
        (enable-theory arithmeticc)))
    (prove-lemma round-of-small-fact
      (rewrite)
      (implies
        (and (rationalp x)
          (not (rllessp fpminimум x)))
        (or (rzerop (round x))
          (equal (round x) (fpminimум)))))))
  )
)
```

```

(prove-lemma equal-numerator-round-min-0 (rewrite)
  (and
    (equal
      (rtimes (rational 2 1) x)
      (rplus x x))
    (equal
      (rtimes x (rational 2 1))
      (rplus x x)))
  (enable-theory r2 arithmetic integers rational-defns)
  (enable rplus simple-rplus equal-reduce-reduce regual itimes
    fix-int-on-integers integerp-minus rtimes simple-times
    iplus fix-rational *1*fix-int *1*rationalP
    integerp-if-numberP))

(lemma rlessp-rtimes-cancel (rewrite)
  (implies
    (and
      (numberP (numerator (fix-rational a)))
      (numberP (numerator (fix-rational b))))
    (and
      (equal
        (rlessp (rtimes a b))
        (rlessp a (rational 1 1))
        (not (zerop b)))
      (equal
        (rlessp (rtimes b a))
        (rlessp a (rational 1 1))
        (not (zerop b)))
      (equal
        (rlessp (rtimes a b))
        (rlessp (rtimes a b))
        (and (zerop b))
        (not (zerop b)))
      (equal
        (rlessp b (rtimes a b))
        (rlessp (rtimes a b))
        (not (zerop b)))
      (equal
        (rlessp b (rtimes b a))
        (rlessp b (rtimes b a))
        (and (zerop b))
        (not (zerop b)))))

    (enable-theory r2 arithmetic integers rational-defns)
    (enable rplus simple-rplus equal-reduce-reduce regual itimes
      fix-int-on-integers integerp-minus rtimes simple-times
      iplus fix-rational *1*fix-int *1*rationalP rlessp
      integerp-if-numberP correctness-of-cancel-lessp-times))

  (prove-lemma bound-fact
    (rewrite)
    (implies (and (not (rlessp b a))
      (not (rlessp (fpmmaximum) (rplus a b)))
      (not (rlessp a
        (rquotient (fpmminimum) (round-min)))))

    (not (rlessp (rquotient (round (rplus a b))
      (rational 2 1))
      (fpmminimum)))))

  (lemma rlessp-round-fpmmaximum (rewrite)
    (implies
      (and
        (numberP (numerator (fix-rational x)))
        (rationalP x))
      (not (rlessp (fpmmaximum) (round x)))
      (enable fpp-round rmagnitude-positive numberP-numerator-round
        fix-rational-of-rationalP fpp-rationalP rlessp-reduce2)
      (use (fpp-bounded-fpmmaximum (x (round x)))))

    )
  )

```

```

(prove-lemma fp-mid-fact1
  (rewrite)
  (implies (and (numberp (numerator a))
                (numberp (numerator b)))
    (not (rllessp (fmaximun)
                  (rimes a (rational 2 1)))))
    (not (rllessp (fmaximun)
                  (rimes b (rational 2 1)))))
    (not (rllessp (fmaximun)
                  (round (rimes b (mid-bound1))))))
    (not (rllessp (fmaximun)
                  (round (rimes b (mid-bound2))))))
    (not (rllessp (fmaximun)
                  (round (rtimes b (mid-bound1))))))
    (not (rllessp (fmaximun)
                  (round (rtimes b (mid-bound2))))))
    (not (rllessp (fmaximun)
                  (quotient (fminimum) (round-min))))))
    (not (rllessp (fp-mid a b))
        (and (rllessp (fp-mid a b))
             (rllessp a (fp-mid a b)))))

  )
  (prove-lemma fix-rational-rzerop (rewrite)
    (implies
      (rllessp a)
      (rzerop x)
      (equal (numerator (fix-rational x)) 0)
      ((enable rzerop reduce-0 fix-rational1)))
    )
  )

  (constrain mid-bound2-intro (rewrite)
    (and (fpp (mid-bound2))
        (or (rllessp (fmaximun) (quotient (fminimum) (round-min)))
            (not (rllessp (mid-bound2)
                          (quotient (fminimum) (round-min))))))

    (mid-bound2 (lambda () (fmaximun))))
  )

  (prove-lemma numberp-numerator-mid-bound2
    (rewrite)
    (implies (not (rllessp (fmaximun)
                           (quotient (fminimum) (round-min)))))

      (numberp (numerator (mid-bound2)))
    )
  )

  (constrain mid-bound3-intro (rewrite)
    (and (fpp (mid-bound3))
        (not (rllessp (quotient (fmaximun) (rational 2 1))
                      (mid-bound3)))
        (mid-bound3 (lambda () (rational 0 1)))
        ((enable quotient)))
    )
  )

  (defn fp-mid (&rest y)
    (if (and (rllessp x (mid-bound2))
              (rllessp (mid-bound2) y))
        (mid-bound2)
        (if (and (rllessp x (rational 0 1))
                  (rllessp (rational 0 1) y))
            (rational 0 1)
            (if (and (rllessp x (rneg (mid-bound2)))
                      (rllessp (rneg (mid-bound2)) y))
                (rneg (mid-bound2))
                (round (quotient (round (rplus x y)
                                         (rational 2 1)))))))
        (lema fpp-mid (rewrite)
          (fpp (fp-mid x y))
          ((enable fp-mid fpp-round-intro mid-bound2-intro)))
        )
      (prove-lemma rlessp-mid-bound3-hack
        (rewrite)
        (implies (not (rllessp (mid-bound3) x))
          (not (rllessp fmaximun)
              (rimes x (rational 2 1))))))
    )
  )

```

```

(prove-lemma fp-mid-bounded-fact
  (rewrite)
  (implies (and (fpp x)
                (fpp m)
                (fpp y)
                (flessp x m)
                (rlessp m y))
    (and (equal (rlessp (find-func-zero-measure x m
                                                (fpminspace))
                        (find-func-zero-measure x y
                            (fpminspace)))
                  t)
          (equal (rlessp (find-func-zero-measure m y
                                                (fpminspace))
                        (find-func-zero-measure x y
                            (fpminspace)))
                  t))
        ((use (fpp-means-find-func-ok)))))

(prove-lemma fp-mid-fact3
  (rewrite)
  (implies (and (not (numberp (numerator a)))
                (not (or (zerop b)
                          (negativep (numerator b)))))

                (fpp a)
                (flessp a b)
                (and (rlessp (fp-mid a b) b)
                     (rlessp a (fp-mid a b)))))

    (def find-func-zero
      (a b)
      (if (or (not (rlessp a b))
              (rlessp (mid-bound3) (rmagnitude a))
              (rlessp (mid-bound3) (rmagnitude b))
              (rlessp (fpmaximum)
                      (rquotient (fpminimum) (round-min))))
          (not (fpp a)))
          (not (fpp b))
          (ratiional 0 1)
          (if (or (and (numberp (numerator a))
                        (or (not (rlessp a
                                          (round (rtimes b (mid-bound1))))))

                        (not (rlessp (mid-bound2) b)))
                        (and (or (zerop b)
                                  (negativep (numerator b)))
                            (or (not (rlessp (round (rtimes a (mid-bound1)))))

                                (cons a b)
                                (let ((mid (fp-mid a b)))
                                  (if (equal (numberp (numerator (func a)))
                                            (numberp (numerator (func mid))))
                                      (find-func-zero mid b)
                                      (find-func-zero a mid)))
                                    ((rlessp (find-func-zero-measure a b
                                                    (fpminspace)))))

                                (prove-lemma rlessp-rmagnitude-x-x
                                  (rewrite)
                                  (not (rlessp (rmagnitude x) x))
                                )
                              )
                            )
                          )
                        )
                      )
                    )
                  )
                )
              )
            )
          )
        )
      )
    )
  )
)

```

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