

# X15: A Tool For Refactoring Java Software Product Lines

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## ABSTRACT

**X15** is the first tool that can apply common *Object-Oriented* (OO) refactorings to Java *Software Product Lines* (SPLs). **X15** is also the first tool that programmers can write custom scripts (to call refactorings programmatically) to retrofit design patterns into Java SPLs. We motivate and illustrate **X15**'s unique capabilities in this paper.

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## 1 INTRODUCTION

Refactoring is a cornerstone of modern Java software development [11]; it should be a cornerstone of modern Java SPL development too. Surprisingly this is not the case as of 2017.

In 2015, the first refactoring engine for SPLs appeared, called **Morpheus** [17]. It supported SPLs coded in C with embedded *C Preprocessor* (CPP) directives, and offered three refactorings: function inline, lift function, and rename.

Why did it take until 2015? There are many reasons; one is that **Morpheus** requires a *variability-aware compiler* (VAC). A VAC integrates CPP directives into the grammar of a host language, in this case C. Let's call this new language C&CPP. Writing a VAC for C&CPP is not simple. And there is the non-trivial challenge to write a refactoring engine for C&CPP.

Now imagine the difficulty of building a VAC for Java&JPP where JPP stands for the *Java Preprocessor* [3, 10, 13],<sup>1</sup> and then writing a refactoring engine for it (where Java easily supports an order of magnitude more refactorings than the C language). We suspect such a compiler+refactoring tool will never be built because of its difficulty.

**X15** is the first refactoring engine for Java SPLs. It uses a standard Java compiler and Java custom annotations to express CPP-like directives. **X15** leverages a radically new refactoring engine for Java, **R3** [16], which currently supports over 30 OO refactorings. **R3** enables programmers to write refactoring scripts – programmatic invocations of refactorings – to retrofit classical design patterns

<sup>1</sup>Or choose your preferred preprocessor.

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[12]. A robust set of benchmarks showed that **R3** performed at least 10× faster than the *Eclipse Java Development Tools* (JDT) refactoring engine [16]. **X15** inherits these capabilities and is minimally slower than **R3**.

This paper explains the unique capabilities of **X15**. Its technical innovations are explained in our SPLC 2017 paper [15].

## 2 BACK-PROPAGATION CONSIDERED HARMFUL

All SPL tools that we are aware of perform a “projection” operation  $\Pi_c$  on an SPL codebase  $\mathbb{P}$  to produce a separate codebase  $P_c$  for the product with configuration  $c$ . That is,  $\Pi_c$  removes code in  $\mathbb{P}$  that is irrelevant to  $c$  (or composes required code fragments in  $\mathbb{P}$  to produce  $P_c$  [2, 5, 18]):

$$\Pi_c(\mathbb{P}) = P_c \quad (1)$$

Today's SPL tools expect  $P_c$  to be compiled, run, and edited to make repairs. With back-propagation tools, modifications to  $P_c$  can be propagated back automatically to  $\mathbb{P}$ , thereby updating all products that share features/code with  $c$ . Without such tools, changes *have to* must be propagated manually.

*Refactorings break this paradigm.* Figure 1(a) shows an SPL codebase  $\mathbb{P}$  expressed in terms of CPP. When  $\mathbb{P}$  is projected (preprocessed) for configuration  $c$  with feature  $F=\text{true}$ , a separate codebase  $P_c$  for product  $c$  is produced, Figure 1(b). Refactoring  $P_c$ , such as renaming variable  $i$  to  $j$ , modifies  $P_c$  to Figure 1(c). When changes to  $P_c$  are back-propagated to  $\mathbb{P}$ , the binding of reference  $j$  breaks when feature  $F$  is unselected (Figure 1(d)).

```
#ifdef F
int i = 0;
#else
int i = 1;
#endif

int m() { i++; }
```

(a) SPL Codebase  $\mathbb{P}$

```
int i = 0;

int m() { i++; }
```

(b) Product  $P_c$  with Feature  $F \in c$

```
int j = 0;

int m() { j++; }
```

(c) Rename  $i$  to  $j$  in  $P_c$

```
#ifdef F
int j = 0;
#else
int i = 1;
#endif

int m() { j++; }
```

(d) Back-propagation of  $P_c$  to  $\mathbb{P}$

**Figure 1: Refactoring-Unaware Back-Propagation.**

In a nutshell, refactorings of SPL products are *not edits*. As the above example shows, if a programmer wants to  $\mathcal{R}$ -refactor the codebase of an SPL product  $P_c$ , a corresponding  $\mathcal{R}$ -refactoring must be applied to  $\mathbb{P}$  to update all references to refactored entities. This requirement is captured by the following algebraic identities:

$$\mathcal{R}(P_c) = \mathcal{R}(\Pi_c(\mathbb{P})) = \Pi_c(\mathcal{R}(\mathbb{P})) \quad (2)$$

That is,  $\mathcal{R}$ -refactoring codebase  $P_C$  must equal a  $c$ -projection of  $\mathcal{R}$ -refactored codebase  $\mathbb{P}$ . More on this in Section 3.3.

### 3 FEATURES AND CAPABILITIES OF X15

**X15** has 3 key advances that make the refactoring of Java SPLs possible:

- (1) Encoding variability in Java using custom annotations,
- (2) Editing and viewing individual SPL products without creating separate codebases, and
- (3) Implementing the identities in equation (2).

Each is discussed in turn in the following subsections.

#### 3.1 Variability using Java Custom Annotations

**X15** uses Java custom annotations in a simple and intuitive way to encode feature variability.

Figure 2(a) shows a **X15** configuration file, where `Feature` is a reserved word in **X15**. Each feature of the SPL is declared as a unique static boolean variable. A value is assigned to each feature of an SPL to define an SPL configuration. In this example, feature `BLUE` is selected and `RED` is not.

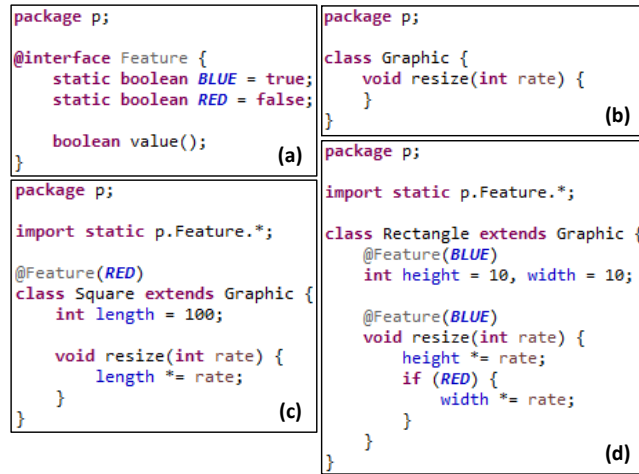


Figure 2: An X15-annotated Java SPL.

Next, every Java package, class, interface, field, and method declaration is annotated with a `Feature` expression to define its *presence condition* (in terms of features) when that declaration is to appear in an SPL product [1]. The class `Square` in Figure 2(c) appears in every SPL product whose configuration includes feature `RED`. If an annotation is absent, the declaration belongs to all SPL products. The `Graphic` class of Figure 2(b) is an example.

Statement blocks are variability-aware if they are enclosed in an `if (feature_expression)` statement. In Figure 2(d), the statement “width \*= rate;” appears in an SPL product whose configuration includes feature `RED`.<sup>2</sup> **X15** infers `Feature` annotations to simplify

<sup>2</sup>Due to the nesting of `Feature` references, the presence condition for statement “width \*= rate;” is `RED ∧ BLUE`. That is, statement “width \*= rate;” appears iff its method `resize` is present (`BLUE` is selected) and if `RED` is also selected (the statement variability qualifier).

variability specification. The precise set of rules **X15** follows is given in Appendix A.

Of course, **X15** needs a feature model for each SPL. **X15** currently uses `guidsl` of the **AHEAD** tool suite [5]; other classical feature modeling tools could have been used [9].

#### 3.2 Projection as Code Folding

All SPL tools have a projection operation ( $\Pi_C$ ) that removes code from  $\mathbb{P}$  that is irrelevant to a product with configuration  $c$ . **X15** uses code folding, a standard IDE functionality [7], to display (project) the code of a product. In Figure 3, code folding hides code fragments of unselected features.

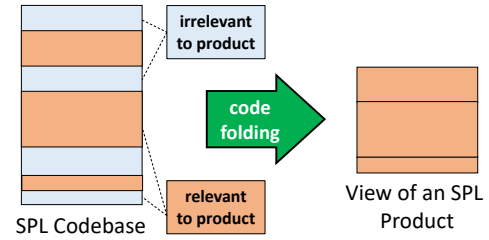


Figure 3: Code Folding reveals  $P_C$ .

How is this done? The standard Java compiler generates ASTs that include subtrees for Java annotations. **X15** determines the current configuration  $c$  by harvesting boolean values from the AST of declaration `@Feature`. Knowing  $c$ , **X15** evaluates the boolean expressions of `@Feature` annotations; if true the annotated AST is pretty-printed, otherwise, the contents of the AST are hidden.

In this way, **X15** folds declarations and statements that are irrelevant to  $P_C$ . An extra benefit, that we will see later, is that *variation points* (VPs) within an SPL codebase are clearly marked by code fold markers ( $\oplus$  meaning folded code and  $\ominus$  indicating expanded code). An **X15** programmer will see where VPs exist in a codebase, and can view – but *not* edit – code-folded regions.

**X15** has a second projection operation that comments-out irrelevant contents of  $\mathbb{P}$  – this version of  $P_C$  is not viewable by a programmer, but is consumed directly by the Java compiler to produce byte codes for  $P_C$ . This allows **X15** users to debug a code-folded view of  $P_C$ , using the commented-out version compiled by `javac`.

Finally, customer-specific adaptations of SPL products are occasionally useful. The question is: should **X15** create a separate codebase for a product (thereby opening a can of worms requiring back-propagation) or should customer-specific adaptations be integrated into the SPL codebase  $\mathbb{P}$  as special features? Currently, we opt for the latter solution.

#### 3.3 Implementing Equation (2)

**X15** follows a common path in SPL tool development: feature algebras axiomatize the semantics of feature composition operations. Tools are then developed to implement these algebras [2, 5, 18].

We started with a known feature algebra [4] and recognized a new axiom that distributes refactorings over sums of features. This axiom was used in the proof of Equation (2) [15].

**X15** implements Equation (2) in the following way: every **X15** user is always modifying the entire SPL codebase  $\mathbb{P}$ . Views merely

restrict edits to a computable subset of  $\mathbb{P}$ . When a programmer invokes a refactoring  $\mathcal{R}$ , the preconditions of  $\mathcal{R}$  are applied to all SPL products, using standard SAT analysis techniques (see Section 4.1). If no violations occur, the code transformation of  $\mathcal{R}$  is applied to  $\mathbb{P}$  to produce  $\mathcal{R}(\mathbb{P})$ . By Equation (2), a c-projection of  $\mathcal{R}(\mathbb{P})$  yields the desired result,  $\mathcal{R}(\mathbb{P}_c)$ . The variability-aware refactoring preconditions that **X15** uses is documented in [15].

Of course, programmers can refactor  $\mathbb{P}$  directly, not requiring a view of a product.

## 4 DEMONSTRATION

**X15** offers four user-interface operations:

- (1) checking dead code,
- (2) checking safe composition,
- (3) code folding to view an SPL product, and
- (4) executing **X15** refactorings on an SPL.

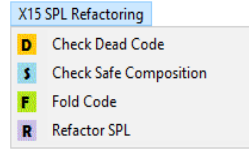


Figure 4: User Interface Operations of **X15**.

The operations are integrated with **X15**'s Eclipse JDT plug-in (Figure 4).

### 4.1 X15 Analyses

**X15** relies on a standard analysis to locate *dead code* – code fragments that belong to no SPL product [1]. **X15** displays a pop-up for any dead code instance found.

**X15** uses a similar standard analysis to guarantee all products of an SPL are type-safe (i.e., compilation is error-free) [8, 19]. **X15** displays a pop-up, much like dead code, for any violation found.

Both are SAT analyses [1], which are leveraged by **X15** to verify refactoring preconditions.

### 4.2 Viewing an SPL Product

As said earlier, **X15** uses code folding to view an SPL product. In Figure 5(a), an if statement is folded in class `Rectangle` because its expression `RED` is false. Figure 5(b) shows an expanded view of `Rectangle`.

For the same reason that `RED` is unselected, class `Square` is folded in Figure 5(c). Figure 5(d) reveals the folded code. **X15** users can fold (or unfold) by clicking either markers  $\oplus$  or  $\ominus$  on the left-hand size of a code fragment. Folded code cannot be edited. Of course, occasionally when a programmer needs to edit folded code, s/he can exit the code-fold mode of **X15** and edit the entire SPL codebase as necessary.

### 4.3 Refactoring of SPL Products

**X15** inherits the **R3**'s ability to allow programmers to write refactoring scripts that automatically retrofit design patterns into a codebase [16]. Figure 6 is an **X15/R3** script that creates a visitor given one of the methods that are to be moved into the created visitor class.<sup>3</sup>

Figure 7 shows the SPL codebase of Figure 2 where a Visitor pattern is retrofitted for the method `resize` [12]. The Visitor design

<sup>3</sup>**R3** scripts can be used without change with **X15**. That is, a script that can be applied to a Java program can also be applied to an **X15** Java SPL.



Figure 5: Code Folding as A View of an SPL Product.

```

1 // member of RMethod class
2 RClass makeVisitor(String N) {
3   RPackage pkg = this.getPackage();
4   RClass v = pkg.newClass(N);
5   RField singleton = v.addSingleton();
6
7   RRelativeList relatives = this.getRelatives();
8   relatives.rename("accept");
9
10  int index = relatives.addparameter(singleton);
11  relatives.moveAndDelegate(index);
12
13  v.getAllMethods().rename("visit");
14
15  return v;
16 }

```

Figure 6: **R3** makeVisitor Method.

pattern moves all methods with the same signature in a class hierarchy to a visitor class, leaving delegates behind. The moved methods are renamed to `visit` and delegates are renamed to `accept`.

Figure 7(d) shows a visitor that has three `visit` methods which are moved from classes `Graphic`, `Rectangle`, and `Square`, respectively. Note that method `resize` of class `Square` is also moved to the Visitor although it does not exist (i.e., its code was folded) in the SPL product (Figure 5(d)). That shows **X15** applies refactorings to the entire SPL codebase, not a single product as explained in Section 3.3.

## 5 RECAP

Modern OO design is at least 20 years old [6]. Refactorings and design patterns have been an integral part of modern OO design for at least that long. SPL development is also about 20 years old, but it lacks tools to refactor OO SPL codebases.

**X15** is a step forward toward improved tools for OO SPL development. It follows and leverages a history of prior work that is

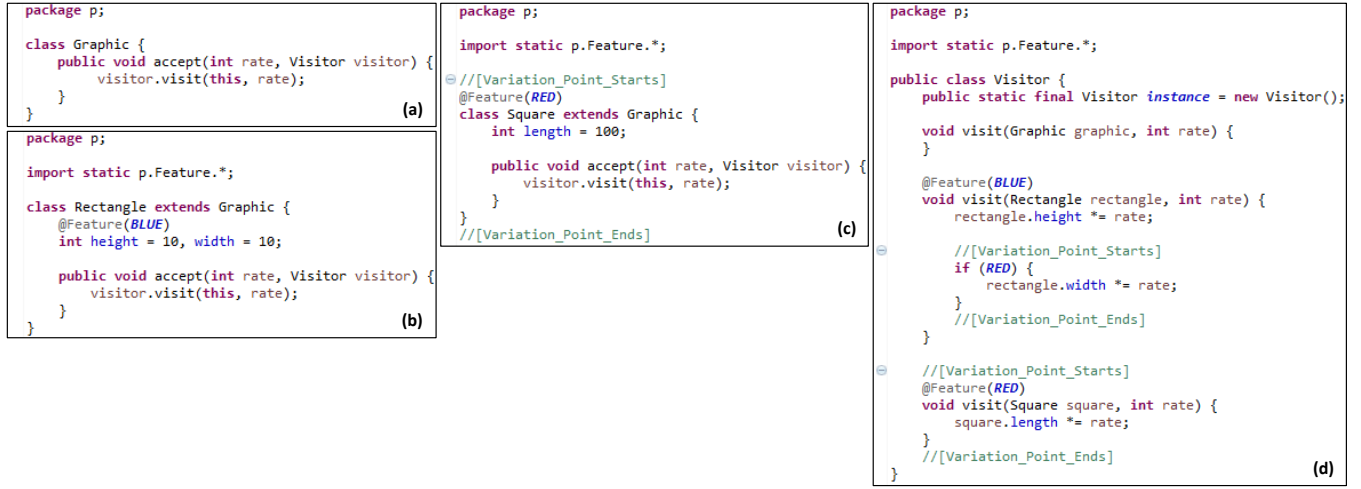


Figure 7: A Visitor Pattern.

detailed in our SPLC 2017 paper [15]. A demo video of **X15** in action is available at our project webpage [20].

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## A ANNOTATION RULES OF X15

Figure 8 lists the declaration types that can be annotated in Java. **X15** allows all annotations except those on formal parameters of methods [14].

```

Package (declared in package-info.java)
Type declaration (i.e., class, interface, enum)
Annotation declaration
Method declaration
Constructor declaration
Field declaration (including enum constant)
Formal and exception parameter
Local variable (including loop variable of for statement
and resource variable of try-with-resources statement)

```

Figure 8: Java Entities that Allow Annotations.

**X15** infers annotation(s) for unannotated entities. (Inferred annotations are not added to the source code.) The purpose of inferring is to avoid annotating every declaration repeatedly when successive annotations are identical. Here are the rules for applying annotations in **X15**. A rule with a lower # has higher priority.

- Rule #1. Use the annotation(s) of the current declaration.
- Rule #2. If Rule #1 is not applicable, use the annotation(s) of the closest preceding declaration in the same scope.
- Rule #3. If Rule #2 is not applicable, use the annotation(s) of the closest enclosing declaration.
- Rule #4. If Rule #3 is not applicable, use a default feature, which is BASE, whose contents is present in every SPL product.

**X15** uses `if (feature_expression) {block}` to conditionally include the codeblock block in a product; `feature_expression` must be true for inclusion to occur; otherwise block is erased along with the `if` statement. This action is independent of the above rules.