

VERIFICATION OF THE STEINING PROTOCOL

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This report contains the transcript of a mechanical verification of the Stenning protocol [Stenning 76]. A description of this protocol, as well as complete documentation on the methods used, can be found in a separate report [DiVito 82]. Reference to this report is necessary, since the following material is not self-contained. The transcripts themselves were produced by the Boyer-Moore theorem prover [Boyer 79].

VC Proof Log

27-Jun-82 09:54:20

Proof of VC TRANSPORT #1

(IMPLIES (AND (SENDER-EXT SOURCE ACK-IN PKT-OUT)
 (RECEIVER-EXT PKT-IN SINK ACK-OUT)
 (FOLLOWS PKT-IN PKT-OUT)
 (FOLLOWS ACK-IN ACK-OUT))
 (INITIAL SINK SOURCE)))

This formula can be simplified, using the abbreviations SENDER, EXP, AND, and IMPLIES, to:

```

(implies
  (and (consistent (quote segno))
        (quote equal)
        (pktno.out))
    (equal source
          (apply (quote msg)
                 (range (latest (quote segno) pkt.out) 0)))
    (receiver-ext pkt.in sink ack.out)
    (follows pkt.in pkt.out)
    (follows ack.in ack.out))
  (initial sink source)),

```

which we simplify, applying CONSISTENT.FOLLOWS, INITIAL.FAPPY, INITIAL.CONSEC.FOLLOWS, NUMBERE.SEGNO, FOLLOWS.GATESI.CONSISTENT, PHAPP.GATESI, INITIAL.RANGE, and INITIAL.TRANS, and opening up RECEIVER.SXT and INITIAL, to:

1

Q.E.D.

13069 conses
9.963 seconds
0.0 seconds, garbage collection time

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VC Proof Log

27-Jun-82 09:55:11

Proof of VC "SENDER#1"

(SENDER-INT (NULL)
 (NULL)
 (NULL)
 (QUOTE IDLE)
 0 0
 (NULL))
 0)

This formula can be simplified, using the abbreviation SENDER.INT, to the following six new conjectures:

Case 6. (PMAPP QUOTE (QUOTE NULL)),

which simplifies, opening up P4APP, etc.

1

Case 5 (NUMBER 9).

this signifies, clearly, to

2

Case 4. (FOLLOWS RANGE QUOTE (QUOTE NULL))
CHART (QUOTE NULL)).

which simplifies, opening up the definitions of RANGE and FOLLOWS,
etc.

5

Case 3. (CONSISTENT (QUOTE SECOND)
 (QUOTE EQUAL)
 (QUOTE (QUOTE NULU))))

which we simplify, opening up SEQP and CONSISTENT, to:

T.

Case 2. (EQUAL (QUOTE (1QUOTE NULL))
 (FAPPLY (QUOTE MSG))
 (RANGE (LATEST (QUOTE SEQNO)
 (QUOTE (1QUOTE NULL))))).

This simplifies, opening up the definitions of SEQP, LATEST,
 RANGE, FAPPLY, and EQUAL, to:

T.

Case 1. (IF
 (EQUAL 0 0)
 (IF (EQUAL (QUOTE (1QUOTE NULL))
 (QUOTE (1QUOTE NULL)))
 (EQUAL (QUOTE (1QUOTE NULL))
 (QUOTE (1QUOTE NULL)))
 F)
 (IF (SEQP (QUOTE (1QUOTE NULL)))
 (IF (EQUAL (DDM (LST (LATEST (QUOTE SEQNO)
 (QUOTE (1QUOTE NULL))))
 (SUB1 0))
 (IF (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO)
 (QUOTE (1QUOTE NULL))))
 (SUB1 0))
 (IF (SEQP (QUOTE (1QUOTE NULL)))
 (EQUAL (DDM (LST (QUOTE (1QUOTE NULL))))
 (SUB1 0))
 T)
 F)
 F))
 F)).

This simplifies, opening up the function EQUAL, to:

T.

Q.E.D.

1052 conses
 2.57 seconds

0.0 seconds, garbage collection time

=====

++++++

Proof of VC "SENDER#2"

(IMPLIES (SENDER.INT SOURCE ACK.IN PKT.OUT STATE UNACK NEXT
 QUEUE TOT)
 (SENDER.EXIT SOURCE ACK.IN PKT.OUT))

This conjecture can be simplified, using the abbreviations
 SENDER.INT and IMPLIES, to:

```
s   (IMPLIES
      (AND
        (PAPP QUEUE)
        (NUMBERP NEXT)
        (FOLLOWS (RANGE QUEUED) PKT.OUT)
        (CONSTSTENT (QUOTE SEQNO)
                     (QUOTE EQUAL)
                     PKT.OUT))
      (EQUAL SOURCE
            (FAPPY (QUOTE ASSG)
                  (RANGE (LATEST (QUOTE SEQNO) PKT.OUT))))
      (IF
        (EQUAL SEXT 0)
        (IF (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
            (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
            F)
        (IF (SEQP PKT.OUT)
            (IF (EQUAL (DOM (LSI (LATEST (QUOTE SEQNO) PKT.OUT)))
                      (SUB1 NEXT))
                (IF (EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
                          (SUB1 NEXT))
                    (IF (SEQP QUEUE)
                        (EQUAL (DOM (LSI QUEUE)) (SUB1 NEXT))
                        T)
                    )
                )
            )
        )
    )
)
```

```

      F)
      F)
      F))
(SENDER.EXT SOURCE ACK.IN PKT.OUT),

```

which simplifies, opening up SENDER.EXT, LATEST, RANGE, SEQ#,
FAPPY, EQUAL, and CONSISTENT, to:

C.

Q.E.D.

15996 conses
 13.576 seconds
 3.261 seconds, garbage collection time

+++++

Proof of VC "SENDER#3"

```

(IMPLIES (AND (SENDER.INT SOURCE ACK.IN PKT.OUT STATE UNACK
                           NEXT QUEUE TOT)
                  (NUMBERP ACK)
                  (LESSP UNACK ACK))
          (SENDER.INT SOURCE
                  (APR ACK.IN ACK)
                  PKT.OUT STATE UNACK NEXT
                  (UPPER QUEUE ACK)
                  TOT))

```

This conjecture can be simplified, using the abbreviations
 SENDER.INT, AND, and IMPLIES, to:

```

(IMPLIES
 (AND
  (PMAPP QUEUED)
  (NUMBERP NEXT))

```

```

(FOLLOW.S (RANGE.QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.OUT)
(EQUAL SOURCE
  (FAPPLY (QUOTE MSG)
    (RANGE (LATEST (QUOTE SEQNO) PKT.OUT))))
(IF
  (EQUAL NEXT 0)
  (IF (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
    (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
    F)
  (IF (SEQP PKT.OUT)
    (IF (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
      (SUB1 NEXT))
      (IF (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
        (SUB1 NEXT))
        (IF (SFQP QUEUE)
          (EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
          T)
        F)
      F)
    F))
  (NUMBERP ACK)
  (LESSP UNACK ACK))
(SENDER.INT SOURCE
  (APP ACK-IN ACK)
  PKT.OUT STATE UNACK NEXT
  (UPPER QUEUE ACK)
  TOT)),

```

which simplifies, rewriting with the lemmas LST.UPPER,
 FOLLOW.RANGE.UPPER, and PHAPP.UPPER, and expanding the functions
 UPPER, FOLLOW, SFQP, RANGE, PHAPP, SENDER.INT, LATEST, FAPPLY,
 EQUAL, and CONSISTENT, to:

F.

Q.E.D.

24249 conses
 20.704 seconds
 0.0 seconds, garbage collection time

+++++

Proof of VC "SENDER#4"

```
(IMPLIES (AND (SENDER.INT SOURCE ACK.IN PKT.OUT STATE UNACK
NEXT QUEUE TOT)
(NUMBERP ACK)
(EQUAL ACK NEXT))
(SENDER.INT SOURCE
(APR ACK.IN ACK)
PKT.OUT
(CQUOTE IDLE)
ACK NEXT
(NULL)
TOT))
```

This conjecture can be simplified, using the abbreviations
SENDER.INT, AND, and IMPLIES, to:

```
(IMPLIES
(AND
(PHAPP QUEUE)
(NUMBERP NEXT)
(FOLLOWAS RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
(CQUOTE EQUAL)
PKT.OUT))
(EQUAL SOURCE
(CAPPLY (QUOTE MSG)
(RANGE (LATEST (QUOTE SEQNO) PKT.OUT))))
(IF
(EQUAL NEXT 0)
(IF (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
(EQUAL QUEUE (QUOTE (1QUOTE NULL)))
F)
(IF (SEQP PKT.OUT)
(IF (EQUAL (CDOM (LST (LATEST (QUOTE SEQNO) PKT.OUT))))
```

```

          (SUB1 NEXT))
(CIF EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
          (SUB1 NEXT))
(CIF (SEQP QUED)
     (EQUAL (CDR (CST QUED)) (SUB1 NEXT))
     T)
   F)
  F)
(F))
(CUMAFRP ACK)
(EQUAL ACK NEXT)
(SENDER.INT SOURCE
  (APR ACK.IN ACK)
  PKT.OUT
  (QUOTE IDLE)
  ACK NEXT
  (QUOTE (QUOTE NULG))
  TUT)),

```

which simplifies, opening up FUBBOWS, SEQP, RANGE, PMAPP,
SENDER.INT, LATEST, FAPPY, EQUAL, and CONSISTENT, to:

F.

Q.E.D.

21684 conses
 18.532 seconds
 3.342 seconds, garbage collection time

Proof of VC "SENDER#5"

```

(IMPRES (SENDER.INT SOURCE ACK.IN PKT.OUT STATE UNACK NEXT
                      QUEUE TOP)
        (SENDER.INT SOURCE ACK.IN

```

```
(JOIN PKT.OUT (RANGE QUEUE))
STATE UNACK NEXT QUEUE TOT))
```

This conjecture can be simplified, using the abbreviations SENDER.INT and IMPLIES, to:

```
(IMPLIES
  (AND
    (PMAFP.QUEUE)
    (NUMBERP.NEXT)
    (FOLLOWS (RANGE QUEUE) PKT.OUT)
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.OUT)
    (EQUAL SOURCE
      (FAPPLY (QUOTE MSG)
        (RANGE (LATEST (QUOTE SEQNO) PKT.OUT))))
    (IF
      (EQUAL NEXT 0)
      (IF (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
          (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
          F)
      (IF (SEQP PKT.OUT)
          (IF (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            (SUB1 NEXT))
              (IF (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
                (SUB1 NEXT))
                  (IF (SFQP.QUEUE)
                      (EQUAL (DOM (LS1.QUEUE)) (SUB1 NEXT))
                      T)
                  F)
              F)
          F)))
      (SENDER.INT SOURCE ACK.IN
        (JOIN PKT.OUT (RANGE QUEUE))
        STATE UNACK NEXT QUEUE TOT)),
```

which simplifies, rewriting with the lemmas HIGHEST.JOIN.FOLLOWS, PISFQP.FAPPLY, FOLLOWS.FAPPLY, FAPPLY.JOIN, LATEST.JOIN.FOLLOWS, NUMBERP.SEQNO, CONSISTENT.JOIN.FOLLOWS, FOLLOWS.JOIN.2, and PMAFP.NLST, and expanding the functions RANGE, SEQ, JOIN, FOLLOWS, PMAFP, SENDER.INT, LATEST, FAPPLY, EQUAL, and CONSISTENT, to six new conjectures:

Case 6. COMPUTES

```
(AND (PMAPP QUEUE)
     (NUMBERP NEXT)
     (FOLLOWS (RANGE QUEUE) PKT.OUT)
     (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.OUT)
     (NOT (EQUAL NEXT 0))
     (SEQP PKT.OUT)
     (EQUAL (DM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            (SUB1 NEXT))
     (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
            (SUB1 NEXT))
     (NOT (SEQP QUEUE)))
     (EQUAL QUEUE (QUOTE (1QUOTE NULL)))),
```

which again simplifies, opening up the definition of PMAPP, to:

T.

Case 5. COMPUTES

```
(AND (PMAPP QUEUE)
     (NUMBERP NEXT)
     (FOLLOWS (RANGE QUEUE) PKT.OUT)
     (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.OUT)
     (NOT (EQUAL NEXT 0))
     (SEQP PKT.OUT)
     (EQUAL (DM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            (SUB1 NEXT))
     (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
            (SUB1 NEXT))
     (EQUAL (DM (LST QUEUE)) (SUB1 NEXT))
     (NOT (SEQP QUEUE)))
     (EQUAL QUEUE (QUOTE (1QUOTE NULL)))),
```

which we again simplify, opening up PMAPP, to:

T.

Case 4. COMPUTES

```
(AND (PMAPP QUEUE)
     (NUMBERP NEXT))
```

```

(FOLLOWING (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL))
(PKT.OUT)
(NOT (EQUAL NEXT 0))
(SEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
  (SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
  (SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(SEQP QUEUE)
(SEQP (NLST QUEUE)))
(CMPAIRP (LST (NLST QUEUE))).

```

However this simplifies again, expanding PMAPP, to:

T.

Case 3. (IMPLIES

```

(AND (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOWING (RANGE QUEUE) PKT.OUT)
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL))
  (PKT.OUT)
  (NOT (EQUAL NEXT 0))
  (SEQP PKT.OUT)
  (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
    (SUB1 NEXT))
  (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
    (SUB1 NEXT))
  (EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
  (SEQP QUEUE)
  (SEQP (NLST QUEUE)))
  (LESSP (DOM (LST (NLST QUEUE)))
    (DOM (LST QUEUE))))).

```

But this simplifies again, unfolding the definition of PMAPP, to:

T.

Case 2. (IMPLIES

```
(AND (PMAPP QUEUE))
```

```

(NUMBERP NEXT)
(FOLLOWS (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.OUT)
(NOT (EQUAL NEXT 0))
(SEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
  (SUB1 NEXT))
(EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
  (SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(SEQP QUEUE)
(NOT (SEQP (LST QUEUE))))
(EQUAL (NLST QUEUE)
  (QUOTE (1QUOTE NIL)))).

```

This simplifies again, unfolding the function PMAPP, to:

T.

Case 1. (IMPLIES

```

(AND (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOWS (RANGE QUEUE) PKT.OUT)
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.OUT)
  (NOT (EQUAL NEXT 0))
  (SEQP PKT.OUT)
  (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
    (SUB1 NEXT))
  (EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
    (SUB1 NEXT))
  (EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
  (SEQP QUEUE))
  (PAIRP (LST QUEUE))).

```

However this simplifies again, unfolding the function PMAPP, to:

T.

Q.E.D.

26435 conses
 24.43 seconds
 3.375 seconds, garbage collection time

 ++++++ ++++++ ++++++ ++++++ ++++++ ++++++ ++++++

Proof of VC "SENDER#6"

```
(IMPLIES (SENDER.INT SOURCE ACK.IN PKT.OUT STATE UNACK NEXT
  QUEUE TOP)
  (SENDER.INT (APP SOURCE MESS)
    ACK.IN
    (APP PKT.OUT (PACKET MESS NEXT))
    (QUOTE BUSY)
    UNACK
    (ADD1 NEXT)
    (WITH THE QUEUE NEXT (PACKET MESS NEXT))
    (PLUS TOP DELTA)))
```

This conjecture can be simplified, using the abbreviations SENDER.INT and IMPLIES, to:

```
IMPLIES
(AND
  (PAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOWS (RANGE QUEUE) PKT.OUT)
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.OUT)
  (EQUAL SOURCE
    (FAPPLY (QUOTE MSG)
      (RANGE (GATEST (QUOTE SEQNO) PKT.OUT))))
  (IF
    (EQUAL NEXT 0)
    (IF (EQUAL PKT.OUT (QUOTE (QUOTE NULL)))
      (EQUAL QUEUE (QUOTE (QUOTE NULL))))
```

```

    F)
  (IF (SEQP PKT.OUT)
      (IF (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
                  (SUB1 NEXT))
           (IF (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
                      (SUB1 NEXT))
               (IF (SEQP QUEUE)
                   (EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
                   T)
               F)
            F)
        F))
  (SENDER.INT (APR SOURCE MESS)
    ACK.IN
    (APR PKT.OUT (PACKET MESS NEXT))
    (QUOTE BUSY)
    UNACK
    (ADD1 NEXT)
    (ATTHE QUEUE NEXT (PACKET MESS NEXT))
    (PLUS TOT DELTA))).
```

This simplifies, applying the lemmas LST.NSEQP, ATTHE.LESSP.DOM.LST, DOM.WPATH, SUB1.ADD1, LST.ATTHE, PMAPP.LATEST, APPLY2.EQUAL, SEQNO.PACKET, APPLY1.SEQNO, FOLLOW.S.PAIR, FOLLOW.S.APR, FOLLOW.S.SAME, FOLLOW.S.TRANS, RNG.WPAIR, LST.APR, NLST.APR, PMAPP.ATTHE, ATTHE.EQUAL.DOM.LST, NLST.NSEQP, and CONSISTENT.APR.NOT.IN, and opening up LESSP, EQUAL, DOM, HIGHEST, MAX, FAPPLY, LATEST, CONSISTENT, CONSISTENT2, IN, RANGE, PMAPP, SENDER.INT, SEQP, and FOLLOW, to the following 23 new goals:

Case 23. (IMPLIES

```

  (AND (PMAPP QUEUE)
       (NUMBERP NEXT)
       (FOLLOW (RANGE QUEUE) PKT.OUT)
       (CONSISTENT (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    PKT.OUT)
       (NOT (EQUAL NEXT 0))
       (SEQP PKT.OUT)
       (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
              (SUB1 NEXT)))
       (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
              (SUB1 NEXT)))
       (NOT (SEQP QUEUE))))
```

```
(EQUAL OUTQUE (QUOTE (QUOTE NULL)))),
```

which we again simplify, expanding the definition of PMAPP, to:

T.

Case 22.(IMPLTES

```
(AND (PMAPP OUTQUE)
      (NUMBERP NEXT)
      (FOLLOWS (RANGE QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (NOT (EQUAL NEXT 0))
      (SEQP PKT.OUT)
      (EQUAL (DDM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
             (SUB1 NEXT))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
             (SUB1 NEXT))
      (NOT (SEQP QUEUE)))
      (CONSISTENT2 (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT
                   (PACKET MESS NEXT))),
```

which again simplifies, using linear arithmetic, applying
SEQNO.PACKET, APPLY1.SEQNO, NUMBERP.SEQNO, and
CONSISTENT2.LESSP.HIGHEST, and opening up the functions PMAPP,
RANGE, SEQP, and FOLLOWS, to:

T.S

Case 21.(IMPLTES

```
(AND (PMAPP OUTQUE)
      (NUMBERP NEXT)
      (FOLLOWS (RANGE QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (NOT (EQUAL NEXT 0))
      (SEQP PKT.OUT)
      (EQUAL (DDM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
             (SUB1 NEXT))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
```

```

(SUB1 NEXT))
(NOT (SEQP QUEUE)))
(EQUAL (APR (FAPPY (QUOTE MSG)
(RANGE (LATEST (QUOTE SEQNO) PKT.OUT))
MESS)
(FAPPY (QUOTE MSG)
(RANGE (WITHE (LATEST (QUOTE SEQNO) PKT.OUT)
NEXT
(PACKET MESS NEXT))))))).

```

This again simplifies, using linear arithmetic, applying the lemmas PMAPP.LATEST, WITHE.LESSP.DOM.LST, RNG.MPAIR, LST.APR, MST.APR, MSG.PACKET, and APPYL1.MSG, and unfolding PMAPP, RANGE, SEQP, FOLLOWS, and FAPPY, to:

I.

Case 20. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWS (RANGE QUEUE) PKT.OUT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.OUT))
 (NOT (EQUAL NEXT 0))
 (SEQP PKT.OUT)
 (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
 (SUB1 NEXT))
 (EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
 (SUB1 NEXT))
 (NOT (SEQP QUEUE))
 (NOT (LESSP NEXT
 (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT))))))
 (EQUAL (DOM (MPAIR NEXT (PACKET MESS NEXT)))
 NEXT)),

which we again simplify, rewriting with DOM.MPAIR, and expanding the definitions of PMAPP, RANGE, SEQP, and FOLLOWS, to:

I.

Case 19. (IMPLIES
 (AND (PMAPP QUEUE)

```

(NUMBERP NEXT)
(FOLLOWS (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
             (QUOTE EQUAL)
             PKT.OUT)
(NOT (EQUAL NEXT 0))
(SEQP PKT.OUT)
(EQUAL (CDR (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
        (SUB1 NEXT))
(EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
        (SUB1 NEXT))
(NOT (SEQP QUEUE))
(GESSP NEXT
       (CDR (LST (LATEST (QUOTE SEQNO) PKT.OUT))))))
(EQUAL (CDR (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
       NEXT)).

```

But this again simplifies, using linear arithmetic, to:

T.

Case 18. (IMPLIES

```

(AND (PMAPP QUEUE)
     (NUMBERP NEXT)
     (FOLLOWS (RANGE QUEUE) PKT.OUT)
     (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.OUT)
     (NOT (EQUAL NEXT 0))
     (SEQP PKT.OUT)
     (EQUAL (CDR (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            (SUB1 NEXT))
     (EQUAL (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
            (SUB1 NEXT))
     (NOT (SEQP QUEUE))
     (GESSP (HIGHEST (FAPPY (QUOTE SEQNO) PKT.OUT))
            NEXT))
     (EQUAL NEXT (SUB1 NEXT))),

```

which again simplifies, using linear arithmetic, to:

T.

Case 17. (IMPLIES

```

(AND (PMAPP QUEUE)
      (NUMBERP NEXT)
      (FOLLOWS (RANGE QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (NOT (EQUAL NEXT 0))
      (SEQP PKT.OUT)
      (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
              (SUB1 NEXT)))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
              (SUB1 NEXT)))
      (NOT (SEQP QUEUE))
      (NOT (LESSP (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
                  NEXT)))
      (EQUAL NEXT (SUB1 NEXT))).

```

This again simplifies, using linear arithmetic, to:

T.

Case 1b. (IMPLIES

```

(AND (PMAPP QUEUE)
      (NUMBERP NEXT)
      (FOLLOWS (RANGE QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (NOT (EQUAL NEXT 0))
      (SEQP PKT.OUT)
      (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
              (SUB1 NEXT)))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
              (SUB1 NEXT)))
      (NOT (SEQP QUEUE))
      (SEQP (FAPPLY (QUOTE SEQNO) PKT.OUT))
      (NOT (LESSP (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
                  NEXT)))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
              NEXT)),

```

which we again simplify, using linear arithmetic, to:

T.

Case 15. (IMPLIES

```

    (AND (PMAAPP.QUEUE)
        (NUMBERP.NEXT)
        (FOLLOWS.RANGE.QUEUE) PKT.OUT)
    (CONSISTENT.(QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.OUT)
    (NOT.(EQUAL NEXT 0))
    (SEQP.PKT.OUT)
    (EQUAL.(DUM(LST(LATEST.(QUOTE SEQNO) PKT.OUT)))
        (SUB1.NEXT))
    (EQUAL.(HIGHEST.(FAPPLY.(QUOTE SEQNO) PKT.OUT))
        (SUB1.NEXT))
    (EQUAL.(DUM(LST.QUEUE)) (SUB1.NEXT)))
    (FOLLOWS.RANGE.(WITH.QUEUE.NEXT.(PACKET.MESS.NEXT)))
    (APR.PKT.OUT.(PACKET.MESS.NEXT))),
```

which again simplifies, using linear arithmetic, applying WITH.LESSP.DUM.LST, RNG.MPAIR, LST.APR, NLST.APR, FOLLOWS.TRANS, FOLLOWS.SAME, FOLLOWS.APR, and FOLLOWS.APR.IN, and expanding RANGE and IN, to:

T.

Case 14. (IMPLIES

```

    (AND (PMAAPP.QUEUE)
        (NUMBERP.NEXT)
        (FOLLOWS.RANGE.QUEUE) PKT.OUT)
    (CONSISTENT.(QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.OUT)
    (NOT.(EQUAL NEXT 0))
    (SEQP.PKT.OUT)
    (EQUAL.(DUM(LST(LATEST.(QUOTE SEQNO) PKT.OUT)))
        (SUB1.NEXT))
    (EQUAL.(HIGHEST.(FAPPLY.(QUOTE SEQNO) PKT.OUT))
        (SUB1.NEXT))
    (EQUAL.(DUM(LST.QUEUE)) (SUB1.NEXT)))
    (CONSISTENT2.(QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.OUT)
    (PACKET.MESS.NEXT))).
```

This again simplifies, using linear arithmetic and applying

SECOND.PACKET, APPLY1.SEQNO, NUMBERP.SEQNO, and
CONSISTENT2.LESSP.HIGHEST, to:

T.S

Case 13.(IMPLTES

```
(AND (PMAAPP.QUEUE)
      (NUMBERP.NEXT)
      (FOLLOWS (RANGE.QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (NOT (EQUAL NEXT 0))
      (SEGP PKT.OUT)
      (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
             (SUB1 NEXT))
      (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
             (SUB1 NEXT))
      (EQUAL (DOM (LST.QUEUE)) (SUB1 NEXT)))
      (EQUAL (APR (FAPPLY (QUOTE MSG)
                           (RANGE (LATEST (QUOTE SEQNO) PKT.OUT)))
                  MSG)
             (FAPPLY (QUOTE MSG)
                     (RANGE (WITH (LATEST (QUOTE SEQNO) PKT.OUT)
                                NEXT
                                (PACKET MSG NEXT))))))),
```

However this simplifies again, using linear arithmetic, applying P1APP.LATEST, WITH.LESSP.DOM.LST, RNG.MPAIR, LST.APR, NLST.APR,
MSG.PACKET, and APPLY1.MSG, and expanding the functions RANGE
and FAPPLY, to:

I.

Case 12.(IMPLTES

```
(AND
  (PMAAPP.QUEUE)
  (NUMBERP.NEXT)
  (FOLLOWS (RANGE.QUEUE) PKT.OUT)
  (CONSISTENT (QUOTE SEQNO)
               (QUOTE EQUAL)
               PKT.OUT)
  (NOT (EQUAL NEXT 0))
  (SEGP PKT.OUT))
```

```

EQUAL (QUM (LST (QATEST (QUOTE SEQNO) PKT,OUT)))
      (SUB1 NEXT))
EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT,OUT))
      (SUB1 NEXT))
EQUAL (QUM (LST QUEUE)) (SUB1 NEXT))
(NOT (LESSP NEXT
              (QUM (LST (QATEST (QUOTE SEQNO) PKT,OUT))))))
EQUAL (QUM (MPATR NEXT (PACKET MESS NEXT)))
      (NEXT)).

```

This simplifies again, rewriting with `DOMPAIR`, to:

三

Case 11. COMPUTES

```

(CARD (PMAPIP-QUEUE)
  (NUMBERP NEXT)
  (FOLLOWS (RANGE-QUEUE) PKT-OUT)
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT-OUT)
  (NOT (EQUAL NEXT 0))
  (SEQP PKT-OUT)
  (EQUAL (CDOM (LST (LATEST (QUOTE SEQNO) PKT-OUT)))
    (SUB1 NEXT)))
  (EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT-OUT))
    (SUB1 NEXT)))
  (EQUAL (CDOM (LST QUEUE)) (SUB1 NEXT))
  (NOT (LESSP NEXT (CDOM (LST QUEUE)))))
  (EQUAL (CDOM (MPAIR NEXT (PACKET-MESS NEXT)))
    NEXT))).

```

This simplifies again, applying DOMPAIR, to:

三

Case 10. IMPLIES

```

(CSEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
(SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
(SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(LESSP NEXT
(DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT))))))
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
NEXT)),

```

which again simplifies, using linear arithmetic, to:

T.

Case 9. (IMPLIES

```

(AND (PMAAPP QUEUE)
(CNUMBERP NEXT)
(FOLLOWSP (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.OUT)
(NOT (EQUAL NEXT 0))
(CSEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
(SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
(SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(LESSP (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
NEXT))
(EQUAL NEXT NEXT)),

```

which again simplifies, using linear arithmetic, to:

T.

Case 8. (IMPLIES

```

(AND (PMAAPP QUEUE)
(CNUMBERP NEXT)
(FOLLOWSP (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.OUT))

```

```

(NOT (EQUAL NEXT 0))
(SEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
(SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
(SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(NOT (SEQP (FAPPLY (QUOTE SEQNO) PKT.OUT))))
(EQUAL NEXT NEXT)).

```

This again simplifies, using linear arithmetic, to:

T.

Case 7. IMPLIES

```

(AND (PMAPP QUEUE)
(NUMBERP NEXT)
(FOLLOWSP RANGE QUEUED PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.OUT)
(NOT (EQUAL NEXT 0))
(SEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
(SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
(SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(GEOR (FAPPLY (QUOTE SEQNO) PKT.OUT))
(NOT (LESSP (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
NEXT)))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
NEXT)),

```

which again simplifies, using linear arithmetic, to:

T.

Case 6. IMPLIES

```

(AND (PMAPP QUEUE)
(NUMBERP NEXT)
(FOLLOWSP RANGE QUEUED PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL))

```

```

          PKT.OUT)
(SEQP PKT.OUT)
(EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
        (SUB1 NEXT))
(EQUAL (HIGHEST (FAPPLY (QUOTE SEQNO) PKT.OUT))
        (SUB1 NEXT))
(EQUAL (DOM (LST QUEUE)) (SUB1 NEXT))
(LESSP NEXT (DOM (LST QUEUE))))
(EQUAL (DOM (LST QUEUE)) (NEXT)),

```

which again simplifies, using linear arithmetic, to:

T.

Case 5. (IMPLIES

```

(AND (PMAPP QUEUE)
      (NUMBERP NEXT)
      (FOLLOWS (RANGE QUEUE) PKT.OUT)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT.OUT)
      (EQUAL NEXT 0)
      (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
      (EQUAL QUEUE (QUOTE (1QUOTE NULL))))
      (EQUAL (APR (QUOTE (1QUOTE NULL)) MESS)
            (FAPPLY (QUOTE MESS)
                    (RANGE (WITHE (LATEST (QUOTE SEQNO) PKT.OUT)
                                  NEXT
                                  (PACKET MESS NEXT))))).

```

But this again simplifies, rewriting with WITHE.ZERO.2, RNG.MPAIR, LST.APK, NLST.APR, MSG.PACKET, APPLY1.MSG, and APR.EQUAL, and opening up the definitions of PMAPP, NUMBERP, RANGE, FOLLOWS, SEQP, CONSISTENT, LATEST, DOMAIN, IN, JOIN, and FAPPLY, to:

```

(EQUAL (QUOTE (1QUOTE NULL))
      (FAPPLY (QUOTE MSG)
              (QUOTE (1QUOTE NULL))).

```

This again simplifies, opening up the definitions of SEQP, FAPPLY, and EQUAL, to:

T.

Case 4. (IMPLIES

```
(AND (PMAPP QUEUE)
     (NUMBERP NEXT)
     (FOLLOWS (RANGE QUEUE) PKT.OUT)
     (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.OUT)
     (EQUAL NEXT 0)
     (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
     (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
     (NOT (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
                  0)))
     (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            NEXT)).
```

But this again simplifies, expanding the definitions of PMAPP, NUMBERP, RANGE, FOLLOWS, SEQP, CONSISTENT, LATEST, LST, DOM, and EQUAL, to:

T.

Case 3. (IMPLIES

```
(AND (PMAPP QUEUE)
     (NUMBERP NEXT)
     (FOLLOWS (RANGE QUEUE) PKT.OUT)
     (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.OUT)
     (EQUAL NEXT 0)
     (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
     (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
     (EQUAL (DOM (LST (LATEST (QUOTE SEQNO) PKT.OUT)))
            0))
     (EQUAL (DOM (MPAIR NEXT (PACKET MESS NEXT)))
            NEXT)).
```

But this simplifies again, applying DOM.MPAIR, and expanding the functions PMAPP, NUMBERP, RANGE, FOLLOWS, SEQP, CONSISTENT, LATEST, LST, DOM, and EQUAL, to:

T.

Case 2. (IMPLIES (AND (PMAPP QUEUE))

```

(NUMBERP NEXT)
(FOLLOWs (RANGE QUEUE) PKT.OUT)
(CONSISTENT (QUOTE SEQNO)
             (QUOTE EQUAL)
             PKT.OUT)
(EQUAL NEXT 0)
(EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
(EQUAL QUEUE (QUOTE (1QUOTE NULL)))
(NOT (SEQP (FAPPLY (QUOTE SEQNO) PKT.OUT))))
(EQUAL NEXT NEXT)),

```

which again simplifies, using linear arithmetic, to:

F.

Case 1. (IMPLIES (AND (PHAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWs (RANGE QUEUE) PKT.OUT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.OUT)
 (EQUAL NEXT 0)
 (EQUAL PKT.OUT (QUOTE (1QUOTE NULL)))
 (EQUAL QUEUE (QUOTE (1QUOTE NULL)))
 (SEQP (FAPPLY (QUOTE SEQNO) PKT.OUT)))
 (EQUAL (HIGHFST (FAPPLY (QUOTE SEQNO) PKT.OUT))
 NEXT)),

which again simplifies, expanding PHAPP, NUMBERP, RANGE, FOLLOWs, SEQP, CONSISTENT, and FAPPLY, to:

F.

Q.E.D.

101452 conses
 94.17 seconds
 10.246 seconds, garbage collection time

VC Proof Log

27-Jun-82 10:06:29

```
+++++
```

Proof of VC "RECEIVER\\$1"

```
(RECEIVER,TNT (NULL)
 (NULL)
 (NULL)
 ())
 (NULL))
```

This conjecture can be simplified, using the abbreviation RECEIVER,TNT, to three new conjectures:

Case 3. (P#APP (QUOTE (1QUOTE NULL))),

which simplifies, expanding the definition of P#APP, to:

T.

Case 2. (NUMBERP 0),

which we simplify, clearly, to:

T.

Case 1. (IMPLIES

```
(CONSISTENT (QUOTE SECOND)
 (QUOTE EQUAL)
 (QUOTE (1QUOTE NULL)))
 (IF
 (FOLLOWSD (QUOTE (1QUOTE NULL))
 (UPPER (LATEST (QUOTE SECOND)
 (QUOTE (1QUOTE NULL)))))
 1))
 (IF
 (LESSP (REACH (LATEST (QUOTE SECOND)
 (QUOTE (1QUOTE NULL)))))
 0)
 F
 (IF
 (IV 0
```

```

(DOMAIN (LATEST (QUOTE SEQNO)
                  (QUOTE (1QUOTE NULL)))))

(IF
  (DESSP 0 0)
  (EQUAL
    (QUOTE (1QUOTE NULL))
    (FAPPLY (QUOTE MSG)
            (RANGE (LOWER (LATEST (QUOTE SEQNO)
                                    (QUOTE (1QUOTE NULL)))
                        (SUB1 0))))))

  F)
(IF (EQUAL 0 0)
  (EQUAL (QUOTE (1QUOTE NULL))
        (QUOTE (1QUOTE NULL)))
  F))
F),

```

which simplifies, expanding the functions SEQN, CONSISTENT,
LATEST, UPPER, FOLLOW, RFACH, DESSP, DOMAIN, IN, and EQUAL, to:

T.

Q.E.D.

1444 conses
4.3 seconds
0.0 seconds, garbage collection time

=====

Proof of VC "RECEIVER#2"

```

(IMPLIES (RECEIVER.INT PKT.IN SINK ACK.OUT NEXT QUEUE)
          (RECEIVER.EXT PKT.IN SINK ACK.OUT))

```

This formula can be simplified, using the abbreviations
RECEIVER,EXT, RECEIVER,INT, and IMPLIES, to:

```

(CMAPLIES
  (AND
    (PAPP-QUEUE)
    (NUMBERP-NEXT)
    (IF
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT-IN)
      (IF
        (FOLLOWING-QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
                 (ADD1 NEXT)))
        (IF
          (LESSP (REACH (LATEST (QUOTE SEQNO) PKT-IN))
                 NEXT)
          F
          (IF
            (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN)))
            (IF
              (LESSP 0 NEXT)
              (EQUAL SINK
                (FAPPY (QUOTE MSG)
                      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT-IN)
                                    (SUB1 NEXT))))))
              F)
            (IF (EQUAL NEXT 0)
                (EQUAL SINK (QUOTE (1QUOTE NULL)))
                F)))
          F)
        T)
      (CONSISTENT (QUOTE SEQNO)
                   (QUOTE EQUAL)
                   PKT-IN))
    (IF
      (IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN)))
      (INITIAL-SINK
        (FAPPY (QUOTE MSG)
              (RANGE (COASEC (LATEST (QUOTE SEQNO) PKT-IN))))),
      (EQUAL SINK (QUOTE (1QUOTE NULL)))),)

```

which simplifies, expanding the functions EQUAL and LESSP, to:

```

(THROUES
  (AND
    (PMAPP.QUEUE)
    (NUMBERP.NEXT)
    (FOLLOWS.QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  NEXT)))
    (IF 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEXT 0))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
              (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                            (SUB1 NEXT))))))
    (CONSISTENT (QUOTE SEQNO)
                 (QUOTE EQUAL)
                 PKT.IN)))
  (INITIAL.SINK
    (FAPPLY (QUOTE MSG)
            (RANGE (CONSEC (LATEST (QUOTE SEQNO) PKT.IN)))))),

```

which again simplifies, using linear arithmetic and applying INITIAL.RANGE, PMAPP.LATEST, INITIAL.LOWER.CONSEC, and INITIAL.FAPPLY, to:

T.

Q.E.D.

42883 conses
 35.007 seconds
 6.695 seconds, garbage collection time

+++++

Proof of VC "RECEIVER#3"

```
(IMPLIES (AND (RECEIVER.INT PKT.IN SINK ACK,OUT NEXT QUEUE)
    (LESSP NEXT (SEQNO PKT))
    (PKTP PKT))
    (RECEIVER.INT (APR PKT.IN PKT)
        SINK ACK,OUT NEXT QUEUE))
```

This conjecture can be simplified, using the abbreviations RECEIVER.INT, AND, and IMPLIES, to the conjecture: \$

```
(IMPLIES
  (AND
    (PAPP QUEUE)
    (NUMBERP NEXT))
  (IF
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.IN)
    (IF
      (FOLLOWS QUOTE
        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
          (ADD1 NEXT)))
      (IF
        (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
          NEXT)
        F
        (IF
          (EQ 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
          (IF
            (LESSP 0 NEXT)
            (EQUAL SINK
              (FAPPY (QUOTE MSG)
                (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                  (SUB1 NEXT))))))
            F)
          (IF (EQUAL NEXT 0)
            (EQUAL SINK (QUOTE (QUOTE NULL)))
            F)))
        F)
      T)
    (LESSP NEXT (SEQNO PKT)))
```

```
(PKTP PKT))
( RECEIVER.INT (APR PKT.IN PKT)
    SINK ACK.OUT NEXT QUEUE)),
```

which we simplify, applying LST.APR, NLST.APR, APPLY2.EQUAL, APPLY1.SEQNO, PMAPP.NLST, IN.DOMAIN.WITHIN.IF, SUB1.ADD1, UPPEK.WITHIN.IF, PMAPP.LATEST, and LOWER.WITHIN.IF, and expanding the definitions of EQUAL, LESSP, CONSISTENT, CONSISTENT2, PMAPP, RECEIVER.INT, and LATEST, to 21 new conjectures:

```
Case 21.(IMPLIES (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (NOT (CONSISTENT (QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.IN)))
    (LESSP NEXT (SECOND PKT))
    (PKTP PKT)
    (NOT (SEQP QUEUE)))
    (EQUAL QUEUE (QUOTE (1QUOTE NULL))))),
```

which we again simplify, opening up PMAPP, to:

T.

```
Case 20.(IMPLIES (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (NOT (CONSISTENT (QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.IN)))
    (LESSP NEXT (SECOND PKT))
    (PKTP PKT)
    (SEQP QUEUE)
    (SEQP (NLST QUEUE)))
    (EQUALP (LST (NLST QUEUE))))).
```

However this simplifies again, expanding PMAPP, to:

T.

```
Case 19.(IMPLIES (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (NOT (CONSISTENT (QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.IN)))
```

```

(GLESSP NEXT (SECOND PKT))
(PKTP PKT)
(SEQP QUEUE)
(NOT (SEQP (NLST QUEUE)))
(GLESSP (DOM (LST (NLST QUEUE)))
(DOM (LST QUEUE))).

```

But this simplifies again, unfolding the definition of PMAPP, to:

T.

```

Case 18.(IMPLIES (AND (PMAPP QUEUE)
                        (NUMBERP NEXT)
                        (NOT (CONSISTENT (QUOTE SEQNO)
                                          (QUOTE EQUAL)
                                          (PKT.IND)))
                        (GLESSP NEXT (SECOND PKT))
                        (PKTP PKT)
                        (SEQP QUEUE)
                        (NOT (SEQP (NLST QUEUE))))
                        (EQUAL (NLST QUEUE)
                               (QUOTE (1QUOTE NULL))))).

```

This simplifies again, unfolding the function PMAPP, to:

T.

```

Case 17.(IMPLIES (AND (PMAPP QUEUE)
                        (NUMBERP NEXT)
                        (NOT (CONSISTENT (QUOTE SEQNO)
                                          (QUOTE EQUAL)
                                          (PKT.IND)))
                        (GLESSP NEXT (SECOND PKT))
                        (PKTP PKT)
                        (SEQP QUEUE))
                        (XPAIRP (LST QUEUE))).

```

However this simplifies again, unfolding the function PMAPP, to:

T.

```

Case 16.(IMPLIES
            (AND (PMAPP QUEUE)
                  (NUMBERP NEXT))

```

```

(FOOLLOWING-QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
         (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT-IN))
                NEXT)))
  (NOT (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN))))
  (EQUAL NEXT 0)
  (EQUAL SINK (QUOTE (1QUOTE NULL)))
  (NOT (EQUAL (SEQNO PKT) 0))
  (PKTP PKT)
  (NOT (SEQP QUEUE)))
  (EQUAL QUEUE (QUOTE (1QUOTE NULL)))),
```

which we again simplify, expanding the definition of PMAPP, to:

T.

Case 15. (TMAPPIES

```

  (AND (PMAPP QUEUE)
        (NUMBERP NEXT)
        (FOOLLOWING-QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
                 (ADD1 NEXT)))
          (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT-IN))
                       NEXT)))
          (NOT (IN 0
                    (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN))))
          (EQUAL NEXT 0)
          (EQUAL SINK (QUOTE (1QUOTE NULL)))
          (NOT (EQUAL (SEQNO PKT) 0))
          (PKTP PKT)
          (SEQP QUEUE)
          (SEQP (NLST QUEUE)))
          (MPAIRP (LST (NLST QUEUE))))),
```

which again simplifies, unfolding the definition of PMAPP, to:

T.

Case 14. (TMAPPIES

```

  (AND (PMAPP QUEUE)
        (NUMBERP NEXT)
        (FOOLLOWING-QUEUE
```

```

        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
          (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    NEXT))
  (NOT (IN 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
  (EQUAL NEXT 0)
  (EQUAL SINK (QUOTE (1QUOTE NULL)))
  (NOT (EQUAL (SEQNO PKT) 0))
  (PKTP PKT)
  (SEQP QUEUE)
  (SEQP (NLST QUEUE)))
  (LESSP (DOM (LST (NLST QUEUE)))
    (DO1 (LST QUEUE))).

```

This again simplifies, unfolding the function PMAPP, to:

T.

Case 13. (IMPLIES

```

  (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT)))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
        NEXT))
      (NOT (IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
      (EQUAL NEXT 0)
      (EQUAL SINK (QUOTE (1QUOTE NULL)))
      (NOT (EQUAL (SEQNO PKT) 0))
      (PKTP PKT)
      (SEQP QUEUE)
      (NOT (SEQP (NLST QUEUE))))
      (EQUAL (NLST QUEUE)
        (QUOTE (1QUOTE NULL))))),

```

which again simplifies, opening up the definition of PMAPP, to:

T.

Case 12. (IMPLIES

```

  (AND (PMAPP QUEUE))

```

```

(CNUMBERP NEXT)
(FOOLLOWING-QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
         (ADD1 NEXT)))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT-IN))
               NEXT)))
(NOT (IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN))))
(EQUAL NEXT 0)
(EQUAL SINK (QUOTE (1QUOTE NULL)))
(NOT (EQUAL (SEQNO PKT) 0))
(PKTP PKT)
(SEQP QUEUE))
(CPAIRP (LST QUEUE))).

```

This again simplifies, expanding the definition of PMAPP, to:

T. S. S.

Case 11. (IMPLIES)

```

(AND (PMAPP-QUEUE)
      (NUMBERP NEXT)
      (FOLLOWS-QUEUE
        (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
               (ADD1 NEXT)))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT-IN))
                    NEXT)))
      (NOT (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT-IN))))
      (EQUAL NEXT 0)
      (EQUAL SINK (QUOTE (1QUOTE NIL)))
      (NOT (EQUAL (SEQNO PKT) 0))
      (PKTP PKT)
      (CONSISTENT2 (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT-IN PKT))
      (CONSISTENT (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT-IN))
      (NOT (LESSP (SUB1 (SEQNO PKT)) NEXT)))
      (FOLLOWS-QUEUE
        (WITH (UPPER (LATEST (QUOTE SEQNO) PKT-IN)
                  (ADD1 NEXT))
              (SEQNO PKT)))
      )
    )
  )
)

```

PKT))),

which again simplifies, applying the lemmas PMAPP.UPPER, PMAPP.LATEST, FOLLOW.SWIPE.IF, IN.DOMAIN.UPPER, IN.UPPER, and DOM.MPAIR, and expanding the definitions of NUMBERP, EQUAL, and LESSP, to:

```
(IMPLIES
  (AND (PMAPP QUEUE)
    (FOLLOW QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        1))
    (NOT (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
    (NOT (EQUAL (SEQNO PKT) 0))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
      (QUOTE EQUAL)
      (PKT.IN PKT))
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      (PKT.IN))
    (NOT (LESSP (SEQNO PKT) 1))
    (IN (SEQNO PKT)
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (IN (SEQNO PKT) (DOMAIN QUEUE)))
    (IN (MPAIR (SEQNU PKT) PKT)
      (LATEST (QUOTE SEQNO) PKT.IN))).

```

This simplifies again, applying NUMBERP.SEQNO, APPLY1.SEQNO, and IN.MPAIR.LATEST.IN, to:

T.

Case 10.(IMPLIES

```
(AND
  (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOW QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    NEXT)))
  (IN 0
    (IN (MPAIR (SEQNU PKT) PKT)
      (LATEST (QUOTE SEQNO) PKT.IN))).

```

```

(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (SUB1 NEXT)))))

(LESSP NEXT (SEQNO PKT))
(PKTP PKT)
(NOT (SEQP QUEUE)))
(EQUAL QUEUE (QUOTE (1QUOTE NULL))),
```

which again simplifies, expanding the function PMAPP, to:

T.

Case 9. IMPLIES

```

(AND
  (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOWS QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    NEXT)))
  (IN 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL NEXT 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        (SUB1 NEXT)))))

  (LESSP NEXT (SEQNO PKT))
  (PKTP PKT)
  (SEQP QUEUE)
  (SEQP (QUOTE NULLQUEUE)))
  (PAIRP (LST (QUOTE NULLQUEUE))))).
```

However this simplifies again, expanding the definition of PMAPP, to:

T.

Case 8. IMPLIES

```
(AND
```

```

(PMAPP QUEUE)
(NUMBERP NEXT)
(FULLQS QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT,IN)
    (ADD1 NEXT)))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
  NEXT)))
(= 0
  (DOMAIN (LATEST (QUOTE SEQNO) PKT,IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
      (SUB1 NEXT)))))

(LESSP NEXT (SEQNO PKT))
(PKTP PKT)
(SEQP QUEUE)
(SEQP (NLST QUEUE)))
(LESSP (DM (LST (NLST QUEUE)))
  (DM (LST QUEUE))).

```

This again simplifies, opening up the definition of PMAPP, to:

T.S

Case 7. (IMPLIES

```

(AND
  (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FULLQS QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT,IN)
      (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
    NEXT)))
  (= 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT,IN)))
  (NOT (EQUAL NEXT 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
        (SUB1 NEXT)))))

  (LESSP NEXT (SEQNO PKT))
  (PKTP PKT)
  (SEQP QUEUE)

```

```
(NOT (SEQP (NLST QUEUE)))
(EQUAL (NLST QUEUE)
      (QUOTE (QUOTE NULL)))).
```

But this again simplifies, expanding the definition of PMAPP, to:

T.

Case 5. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 NEXT)))))
 (LESSP NEXT (SEQNO PKT))
 (PKTP PKT)
 (SEQP QUEUE)
 (MPATRP (LST QUEUE))).

But this simplifies again, expanding the definition of PMAPP, to:

T.

Case 5. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))

```

(NOT (EQUAL NEXT 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (SUB1 NEXT)))))

(LESSP NEXT (SEQNO PKT))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN)
(EQUAL (SEQNO PKT) (SUB1 NEXT))
(NOT (IN (PAIR (SEQNO PKT) PKT)
  (LATEST (QUOTE SEQNO) PKT.IN)))
(EQUAL (SUB1 NEXT) 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (APR (QUOTE (QUOTE VULN))
      (PAIR 0 PKT))))),

```

which we again simplify, using linear arithmetic, to:

T.

Case 4. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      NEXT))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEXT 0))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
          (SUB1 NEXT)))))

    (LESSP NEXT (SEQNO PKT))
    (PKTP PKT)
  )
)
```

```

(CONSISTENT2 (QUOTE SEQNO)
              (QUOTE EQUAL)
              PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
              (QUOTE EQUAL)
              PKT.IN)
(EQUAL (SEQNO PKT) (SUB1 NEXT))
(NOT (IN (MPAIR (SEQNO PKT) PKT)
           (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SUB1 NEXT) 0)))
(EQUAL SINK
      (FAPPLY (QUOTE MSG)
              (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                           (SUB1 (SUB1 NEXT)))
                            (MPAIR (SEQNO PKT) PKT))))).

```

but this again simplifies, using linear arithmetic, to:

T.

Case 3. (INPUTS

```

  (AND
    (P4APP QUEUE)
    (NUMBERP TEXT)
    (FUDOWNS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                 NEXT)))
    (IN 0
      (DOMATN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEXT 0)))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
              (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                           (SUB1 NEXT))))))
    (LESSP NEXT (SEQNO PKT))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN PKT)
    (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN)
  )

```

```

(CNOT (EQUAL (SEQNO PKT) (SUB1 NEXT)))
(CNOT (LESSP (SUB1 NEXT) (SEQNO PKT))))))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (WITHE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
      (SUB1 NEXT))
      (SEQNO PKT)
      PKT))))).

```

However this simplifies again, using linear arithmetic, to:

. T.

Case 2. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT,IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
 NEXT)))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT,IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG)
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
 (SUB1 NEXT)))))))
 (LESSP NEXT (SEQNO PKT))
 (PKTP PKT)
 (CONSTANT2 (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT,IN PKT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT,IN)))
 (NOT (LESSP (REACH (WITHE (LATEST (QUOTE SEQNO) PKT,IN)
 (SEQNO PKT)
 PKT)))
 NEXT))),

which we again simplify, using linear arithmetic and applying
 LESSP,REACH,WITHE,Z and PMAPP,LATEST, to:

T.S

Case 1. (IMPLIES
 (AND
 (PMAPP.QUEUE)
 (NUMBERP.NEXT)
 (FOLLOWS.QUEUE
 (UPPER.(LATEST.(QUOTE SEQNO).PKT.IN)
 (ADD1.NEXT)))
 (NOT.LESSP.(REACH.(LATEST.(QUOTE SEQNO).PKT.IN))
 (NEXT)))
 (IN.0
 (DOMAIN.(LATEST.(QUOTE SEQNO).PKT.IN)))
 (NOT.(EQUAL.NEXT.0))
 (EQUAL.SINK
 (FAPPLY.(QUOTE MSG)
 (RANGE.(LOWER.(LATEST.(QUOTE SEQNO).PKT.IN)
 (SUB1.NEXT))))))
 (LESSP.NEXT.(SEQNO.PKT))
 (PKTP.PKT)
 (CONSTSTENT2.(QUOTE SEQNO)
 (QUOTE.EQUAL)
 (PKT.IN.PKT))
 (CONSTSTENT.(QUOTE SEQNO)
 (QUOTE.EQUAL)
 (PKT.IN))
 (NOT.(EQUAL.(SEQNO.PKT).0))
 (NOT.LESSP.(SUB1.(SEQNO.PKT).NEXT)))
 (FOLLOWS.QUEUE
 (WITHE.CUPPER.(LATEST.(QUOTE SEQNO).PKT.IN)
 (ADD1.NEXT))
 (SEQNO.PKT)
 (PKT))).

But this simplifies again, applying PMAPP.UPPER, PMAPP.LATEST,
 FOLLOWS.WITHE.IF, IN.DOMAIN.UPPER, SUB1.ADD1, IN.UPPER, DOM.MPAIR,
 NUMBERP.SEQNO, APPLY1.SEQNU, and IN.MPAIR.LATEST.II, and
 expanding the definition of LESSP, to:

T.

Q.E.D.

179171 conses
164.418 seconds
20.28 seconds, garbage collection time

Proof of VC "RECEIVER #4"

(IMPLIES (AND (RECEIVER.INT PKT.IN SINK ACK.OUT NEXT QUEUE)
 (LESSP (SECOND PKT) NEXT)
 (PKTP PKT))
 (RECEIVER.INT (APR PKT.IN PKT)
 SINK
 (APR ACK.OUT NEXT)
 (NEXT QUEUE)))

This formula can be simplified, using the abbreviations RECEIVES, INT., AND, and IMPLIES, to: $\neg S \rightarrow (R \wedge I)$

```

(IMPUTES
  (AND
    (PMAAPP QUEUE)
    (NUMBERP NEXT))
  (IF
    (CONSISTENT QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT,IN)
  (IF
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT,IN)
        (ADD1 NEXT)))
    (IF
      (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
        NEXT)
      F
      (IF
        (EQ 0

```

```

        (PDRATH (LATEST (QUOTE SEQNO) PKT.IN)))
(IF
  (LESSP V NEXT)
  (EQUAL SINK
    (FAPPY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                    (SUB1 NEXT))))))
  F)
(IF (EQUAL NEXT 0)
  (EQUAL SINK (QUOTE (QUOTE NULL)))
  F))
F)
T)
(LESSP (SEQNO PKT) NEXT)
(PKTP PKT))
(RECEIVER.INT (APR PKT.IN PKT)
  SINK
  (APR ACK.OUT NEXT)
  NEXT QUEUE)).

```

This simplifies, applying LST.APR, NLST.APR, APPLY2.EQUAL, APPLY1.SEQNO, PMAPP.NLST, LOWER.WITH.E.IF, IN.DOMAIN.WITH.E.IF, SUB1.ADD1, UPPER.WITH.E.IF, and PMAPP.LATEST, and expanding the functions EQUAL, LESSP, CONSISTENT, CONSISTENT2, PMAPP, RECEIVER.INT, and LATEST, to the following 15 new goals:

Case 15. (IMPLIES (AND (PMAPP QUEUE)
 (NUMBERP NEXT)
 (NOT (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)))
 (LESSP (SEQNO PKT) NEXT)
 (PKTP PKT)
 (NOT (SEQP QUEUE)))
 (EQUAL QUEUE (QUOTE (QUOTE NULL))))).

This simplifies again, opening up the definition of PMAPP, to:

T.

Case 14. (IMPLIES (AND (PMAPP QUEUE)
 (NUMBERP NEXT)
 (NOT (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)))
 (LESSP (SEQNO PKT) NEXT)
 (PKTP PKT)
 (NOT (SEQP QUEUE)))
 (EQUAL QUEUE (QUOTE (QUOTE NULL))))).

```

                           PKT.IN))
(LESSP (SEQNO PKT) NEXT)
(PKTP PKT)
(SEQP QUEUE)
(SEQP (NLST QUEUE)))
(NPATKP (LST (NLST QUEUE))),
```

which again simplifies, opening up the definition of PMAPP, to:

T.

```

Case 13.(IMPLIES (AND (PMAPP QUEUE)
                        (NUMBERP NEXT)
                        (NOT (CONSISTENT (QUOTE SEQNO)
                                          (QUOTE EQUAL)
                                          PKT.IN)))
                        (LESSP (SEQNO PKT) NEXT)
                        (PKTP PKT)
                        (SEQP QUEUE)
                        (SEQP (NLST QUEUE)))
                        (LESSP (DOM (LST (NLST QUEUE)))
                               (DOM (LST QUEUE)))),
```

which we again simplify, opening up PMAPP, to:

T.

```

Case 12.(IMPLIES (AND (PMAPP QUEUE)
                        (NUMBERP NEXT)
                        (NOT (CONSISTENT (QUOTE SEQNO)
                                          (QUOTE EQUAL)
                                          PKT.IN)))
                        (LESSP (SEQNO PKT) NEXT)
                        (PKTP PKT)
                        (SEQP QUEUE)
                        (NOT (SEQP (NLST QUEUE))))
                        (EQUAL (NLST QUEUE)
                               (QUOTE (1QUOTE NULL)))).
```

However this simplifies again, expanding PMAPP, to:

T.

```
Case 11.(IMPLIES (AND (PMAPP QUEUE)
```

```

  (NUMBERP NEXT)
  (NOT (CONSISTENT (QUOTE SEQNO)
                     (QUOTE EQUAL)
                     PKT.IN))
  (LESSP (SEQNO PKT) NEXT)
  (PKTP PKT)
  (SEQP QUEUE))
  (MAPAPP (USE QUEUE))).

```

But this simplifies again, unfolding the definition of PMAPP, to:

T.

Case 10. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN)
                         NEXT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEAT 0))
    (EQUAL SINK
      (MAPAPP (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                      (SUB1 NEXT)))))
    (LESSP (SEQNO PKT) NEXT)
    (PKTP PKT)
    (NOT (SEQP QUEUE)))
    (EQUAL QUEUE (QUOTE (QUOTE VALUE))))).

```

This simplifies again, unfolding the function PMAPP, to:

T.

Case 9. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN)
                         NEXT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEAT 0))
    (EQUAL SINK
      (MAPAPP (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                      (SUB1 NEXT)))))
    (EQUAL QUEUE (QUOTE (QUOTE VALUE))))).

```

```

          (ADD1 NEXT)))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
NEXT))
(CIN 0
  (DOMAIN (LATEST (QUOTE SEQNO) PKT,IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
  (FAPPY (QUOTE MSG))
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
(SUB1 NEXT)))))

(LESSP (SEQNO PKT) NEXT)
(PKTP PKT)
(SEQP QUEUED)
(SEQP (NLST QUEUED))
(CPATRP (LST (NLST QUEUED))).

```

However this simplifies again, unfolding the function PMAPP, to:

T.

Case 8. (IMPLIES
 (AND
 (PMAPP QUEUED)
 (NUMBERP NEXT)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT,IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT,IN))
NEXT))
 (CIN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT,IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT,IN)
(SUB1 NEXT)))))

 (LESSP (SEQNO PKT) NEXT)
 (PKTP PKT)
 (SEQP QUEUED)
 (SEQP (NLST QUEUED))
 (LESSP (QUM (LST (NLST QUEUED)))
 (QUM (LST QUEUED))))).

which we again simplify, expanding the definition of PMAPP, to:

T.

Case 7. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWSP QUEUE
 (CUPPER (LATEST (QUOTE SEQNO) PKT.IN))
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN))
 (SUB1 NEXT))))
 (LESSP (SEQNO PKT) NEXT)
 (PKTP PKT)
 (SEQP QUEUE)
 (NOT (SEQP (NLST QUEUE))))
 (EQUAL (NLST QUEUE)
 (QUOTE (1QUOTE NULL)))),

which again simplifies, unfolding the definition of PMAPP, to:

T.

Case 6. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (NUMBERP NEXT)
 (FOLLOWSP QUEUE
 (CUPPER (LATEST (QUOTE SEQNO) PKT.IN))
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))

```

(RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
(SUB1 NEXT)))))

(LESSP (SEQNO PKT) NEXT)
(PKTP PKT)
(SQDP QUEUE))
(CMPAIRP (LST QUEUE))).

```

This again simplifies, unfolding the function PMAPP, to:

I.s

Case 5. CTMPLTES

```

(AND
  (PMAPP QUEUE)
  (NUMBERP NEXT)
  (FULLQNS QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    NEXT))
  (IN 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL NEXT 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG))
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (SUB1 NEXT)))))

  (LESSP (SEQNO PKT) NEXT)
  (PKTP PKT)
  (CONSISTENT2 (QUOTE SEQNO)
    (QUOTE EQUAL)
    (PKT.IN PKT))
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    (PKT.IN))
  (EQUAL (SEQNO PKT) (SUB1 NEXT))
  (NOT (IN (CMPAIR (SEQNO PKT) PKT)
    (LATEST (QUOTE SEQNO) PKT.IN)))
  (EQUAL (SUB1 NEXT) 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG))
    (RANGE (APR (QUOTE (QUOTE NULL))
      (PAIR 0 PKT))))),

```

which again simplifies, using linear arithmetic, applying the lemmas NUMBERP.SEQNO, APPY1.SEQNO, and IN.MPAIR.LATEST.IN, and expanding the definitions of EQUAL and LESSP, to:

T.

Case 4. (IMPLIES

```
(AND
  (PMAPP OUTUE)
  (NUMBERP NEXT)
  (FOLLOWS OUTUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
           (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
               NEXT)))
  (EQV 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL NEXT 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG)
            (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                          (SUB1 NEXT)))))
  (LESSP (SEQNO PKT) NEXT)
  (PKTP PKT)
  (CONSISTENT2 (QUOTE SEQNO)
                (QUOTE EQUAL)
                PKT.IN PKT)
  (CONSISTENT (QUOTE SEQNO)
                (QUOTE EQUAL)
                PKT.IN)
  (EQUAL (SEQNO PKT) (SUB1 NEXT))
  (NOT (IN (MPAIR (SEQNO PKT) PKT)
            (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL (SUB1 NEXT) 0)))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG)
            (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                           (SUB1 (SUB1 NEXT))))
                  (MPAIR (SEQNO PKT) PKT))))).
```

This simplifies again, using linear arithmetic and applying NUMBERP.SEQNO, IN.DOMAIN.REACH, PMAPP.LATEST, APPY1.SEQNO, and IN.MPAIR.LATEST.IN, to:

T.S

Case 3. (IMPLIES
 (AND
 (PMAPP.QUEUE)
 (NUMBERP.NEXT)
 (FOLLOWS.QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT)))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 NEXT))))
 (LESSP (SEQNO PKT) NEXT)
 (PKTP PKT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN PKT))
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN))
 (NOT (EQUAL (SEQNO PKT) (SUB1 NEXT)))
 (NOT (LESSP (SUB1 NEXT) (SEQNO PKT))))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (WITH (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 NEXT))
 (SEQNO PKT)
 (PKT))))).

But this again simplifies, using linear arithmetic and applying
 IN.MPAIR.LATEST.IN, APPLY1.SEQNO, IN.DOMAIN.REACH, NUMBERP.SEQNO,
 DOM.MPAIR, IN.LOWER, PMAPP.LOWER, PMAPP.LATEST, and
 WITH.IN.MPAIR, to:

T.

Case 2. (IMPLIES
 (AND

```

(PMAPP.QUEUE)
(NUMBERP.NEXT)
(FULLDAS.QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
           (ADD1 NEXT)))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
              NEXT)))
(IN 0
    (DOMATV (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
    (FAPPY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                      (SUB1 NEXT))))))
(LESSP (SERNO PKT) NEXT)
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN))
(NOT (LESSP (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
                           (SEQNO PKT)
                           PKT)))
              NEXT))).
```

However this again simplifies, using linear arithmetic and applying LESSP.REACH.WITHE.2 and PMAPP.LATEST!, to:

I.

Case 1. (IMPLIES
 (AND
 (PMAPP.QUEUE)
 (NUMBERP.NEXT)
 (FULLDAS.QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT)))
 (IN 0
 (DOMATV (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0))))

```
(EQUAL SINK  
  (APPEND (QUOTE MSG)  
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT-IN)  
      (SUB1 NEXT))))  
(LESSP (SEQNO PKT) NEXT)  
(PKTP PKT)  
(CONSISTENT? (QUOTE SEQNO)  
  (QUOTE EQUAL)  
  PKT-IN PKT)  
(CONSISTENT (QUOTE SEQNO)  
  (QUOTE EQUAL)  
  PKT-IN)  
(NOT (EQUAL (SEQNO PKT) 0))  
(NOT (LESSP (SUB1 (SEQNO PKT)) NEXT)))  
(FOLLOWS QUEUE  
  (WITH-CURRENT (LATEST (QUOTE SEQNO) PKT-IN)  
    (ADD1 NEXT))  
  (SEQNO PKT)  
  PKT))),
```

which again simplifies, using linear arithmetic, to:

1

Geodes.

```
126660 conses  
114.513 seconds  
13.528 seconds, garbage collection time
```

Proof of VC "RECEIVER #5"

(IMPLIES (AND (RECEIVER INT PKT IN SINK ACK OUT NEXT QUEUE)
 (LESSP NEXT (SEND PKT))
 (NOT (IN (SEND PKT) (DOMAIN QUEUED))))

```

(PKTP PKT))
( RECEIVER.INT (APR PKT.IN PKT)
  STK ACK.OUT NEXT
  (WITH THE QUEUE (SEQNO PKT) PKT)))

```

This formula can be simplified, using the abbreviations NOT,
RECEIVER.INT, AND, and IMPLIES, to: \$\$\$

```

(IMPLIES
  (AND
    (P=APP QUEUE)
    (NUMBERP NEXT)
    (IF
      (CONSISTENT (QUOTE SEQNO)
        (QUOTE EQUAL)
        PKT.IN)
      (IF
        (FOLLOWS QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
            (ADD1 NEXT)))
        (IF
          (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
            NEXT)
          F
          (IF
            (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
            (IF
              (LESSP 0 NEXT)
              (EQUAL STK
                (FAPPLY (QUOTE MSG))
                (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                  (SUB1 NEXT))))))
            F)
            (IF (EQUAL NEXT 0)
              (EQUAL SINK (QUOTE (QUOTE NULL)))
              F1))
            F)
          T)
        (LESSP NEXT (SEQNO PKT))
        (NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
        (PKTP PKT))
        ( RECEIVER.INT (APR PKT.IN PKT)
          SINK ACK.OUT NEXT

```

(WITHQUEUE (SEQNO PKT) PKT))),

which we simplify, applying LST.APP, NLST.APP, APPLY2.EQUAL, APPBY1.SEQNO, PMAPP.WITHQUEUE, IN.DOMAIN.WITHQUEUE.IF, SUB1.ADD1, UPPER.WITHQUEUE.IF, PMAPP.LATEST, and LOWER.WITHQUEUE.IF, and opening up EQUAL, LESSP, CONSISTENT, CONSISTENT2, RECEIVER.IN, and LATEST, to the following nine new goals:

Case 9. (IMPLIES

```

  (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWSP QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATESTI (QUOTE SEQNO) PKT.IN))
      NEXT)))
    (NOT (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
    (EQUAL NEXT 0)
    (EQUAL SINK (QUOTE (QUOTE NULL)))
    (NOT (EQUAL (SEQNO PKT) 0))
    (NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
      (QUOTE EQUAL)
      (PKT.IN PKT))
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      (PKT.IN))
    (LESSP (SUB1 (SEQNO PKT)) NEXT))
    (FOLLOWSP (WITHQUEUE (SEQNO PKT) PKT)
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT)))),
```

which we again simplify, using linear arithmetic, to:

T.S\$

Case 8. (IMPLIES

```

  (AND (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWSP QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT))))
```

```

(SUB1 NEXT)))))

(LESSP NEXT (SEQNO PKT))
(NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN))
(EQUAL (SEQNO PKT) (SUB1 NEXT))
(NOT (IN (MPAIR (SEQNO PKT) PKT)
  (LATEST (QUOTE SEQNO) PKT.IN)))
(EQUAL (SUB1 NEXT) 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSGG)
    (RANGE (APR (QUOTE (QUOTE NULL))
      (MPAIR 0 PKT))))),

```

which again simplifies, using linear arithmetic, to:

T.

Case 6. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWSP QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      NEXT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEXT 0))
    (EQUAL SINK
      (FAPPLY (QUOTE MSGG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
          (SUB1 NEXT))))))
    (LESSP NEXT (SEQNO PKT))
    (NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
      (QUOTE EQUAL))

```

```

        PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.IN)
(EQUAL (SEQNO PKT) (SUB1 NEXT))
(NOT (IN (MPAIR (SEQNO PKT) PKT)
    (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SUB1 NEXT) 0)))
(EQUAL SINK
    (FAPPLY (QUOTE MSG))
    (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        (SUB1 (SUB1 NEXT))))
        (MPAIR (SEQNO PKT) PKT)))).
```

However this simplifies again, using linear arithmetic, to:

T.

Case 5. (IMPLIES

- (AND
- (PMAAPP QUEUE)
- (NUMBERP NEXT)
- (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 NEXT)))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 NEXT))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL NEXT 0)))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 NEXT)))))

 (LESSP NEXT (SEQNO PKT))
 (NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
 (PKTP PKT)
 (CONSISTENT2 (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN PKT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)
 (NOT (EQUAL (SEQNO PKT) (SUB1 NEXT))))

```

(NOT (LESSP (SUB1 NEXT) (SEQNO PKT))))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (WITHE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (SUB1 NEXT))
      (SEQNO PKT)
      PKT)))),,

```

which again simplifies, using linear arithmetic, to:

I.

Case 4. IMPLIES

```

(AND
  (PAAPP QUEUE)
  (NUMBERP NEXT)
  (FOLLOWs QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 NEXT)))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    NEXT))
  (IN 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL NEXT 0))
  (EQUAL SINK
    (FAPPLY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        (SUB1 NEXT)))))

  (LESSP NEXT (SEQNO PKT))
  (NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
  (PKTP PKT)
  (CONSTANT2 (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.IN PKT)
  (CONSISTENT (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.IN))
  (NOT (LESSP (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
    (SEQNO PKT)
    PKT)))
    NEXT))).

```

But this simplifies again, using linear arithmetic and rewriting with LESSP.REACH.WITHE.2 and PAAPP.LATEST, to:

I.

Case 3. (IMPLIES

(AND

(PNAAPP QUEUE)

(NUMBERP NEXT)

(FOLLOWS QUEUE

 (UPPER (LATEST (QUOTE SEQNO) PKT.IN))

 (ADD1 NEXT)))

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))

 NEXT)))

(IN 0

 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))

(NOT (EQUAL NEXT 0))

(EQUAL SINK

 (FAPPLY (QUOTE MSG))

 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN))

 (SUB1 NEXT))))

(LESSP NEXT (SEQNO PKT))

(NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))

(PKTP PKT)

(CONSISTENT2 (QUOTE SEQNO)

 (QUOTE EQUAL)

 PKT.IN PKT)

(CONSISTENT (QUOTE SEQNO)

 (QUOTE EQUAL)

 PKT.IN)

(EQUAL (SEQNO PKT) 0))

(FOLLOWS (WITHE QUEUE (SEQNO PKT) PKT)

 (UPPER (LATEST (QUOTE SEQNO) PKT.IN))

 (ADD1 NEXT))).

This again simplifies, using linear arithmetic, to:

I.

Case 2. (IMPLIES

(AND

(PNAAPP QUEUE)

(NUMBERP NEXT)

(FOLLOWS QUEUE

 (UPPER (LATEST (QUOTE SEQNO) PKT.IN))

 (ADD1 NEXT)))

```

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
               NEXT))
(IN 0
   (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                  (SUB1 NEXT)))))

(LESSP NEXT (SEQNO PKT))
(NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
(PKTP PKT)
(CONSTSENT2 (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN PKT)
(CONSTSENT (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN)
(LESSP (SUB1 (SEQNO PKT)) NEXT))
(FOLLOWS (WITHQUEUE (SEQNO PKT) PKT)
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                  (ADD1 NEXT))),,

```

which again simplifies, using linear arithmetic, to:

T.SS

Case 1. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
              (ADD1 NEXT)))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                 NEXT))
    (IN 0
       (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL NEXT 0))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                      (SUB1 NEXT)))))

    (LESSP NEXT (SEQNO PKT)))

```

```

(NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN))
(NOT (EQUAL (SEQNO PKT) 0))
(NOT (LESSP (SUB1 (SEQNO PKT)) NEXT)))
(FOLLOWs (WITHE.QUEUE (SEQNO PKT) PKT)
  (WITHE.UPPER (LATEST (QUOTE SEQNO) PKT.IN)
    (ADD1 NEXT))
  (SEQNO PKT)
  PKT)),

```

which again simplifies, rewriting with RNG.MPAIR, DOM.MPAIR,
 IN.WITHE.MPAIRP, SUB1.ADD1, IN.DOMAIN.UPPER, FOLLOWs.WITHE.IF,
 PMAPP.WITHE, PMAPP.LATEST, PMAPP.UPPER, and
 FOLLOWs.WITHE.IN.MPAIR, and expanding LESSP, to:

T.

Q.E.D.

176501 conses
 161.821 seconds
 23.681 seconds, garbage collection time

+++++

Proof of VC "RECEIVER#6"

```

IMPLIES (AND (RECEIVER.INT PKT.IN SINK ACK.OUT NEXT QUEUE)
  (NOT (IN (ADD1 NEXT) (DOMAIN QUEUE)))
  (EQUAL (SEQNO PKT) NEXT)
  (PKTP PKT))

```

```
(RECEIVER.INT (APR PKT.IN PKT)
  (APR SINK (MSSG PKT))
  (APR ACK.OUT (ADD1 NEXT))
  (ADD1 NEXT)
  QUEUED))
```

This formula can be simplified, using the abbreviations NOT, RECEIVER.INT, AND, and IMPLIES, to:

```
(IMPLIES
  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT))
  (IF
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.IN)
    (IF
      (FOLLOWS QUEUE
        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
          (ADD1 NEXT)))
      (IF
        (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
          NEXT)
        F
        (IF
          (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
          (IF
            (LESSP 0 NEXT)
            (EQUAL SINK
              (FAPPLY (QUOTE MSSG)
                (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                  (SUB1 NEXT)))))
            F)
            (IF (EQUAL NEXT 0)
              (EQUAL SINK (QUOTE (QUOTE NULL)))
              F)))
          F)
        T)
      (NOT (IN (ADD1 NEXT) (DOMAIN QUEUE)))
      (EQUAL (SEQNO PKT) NEXT)
      (PKTP PKT))
    (RECEIVER.INT (APR PKT.IN PKT)))
```

```

(APR.SINK (MSG.PKT))
(APR.ACK.OUT (ADD1.NEXT))
(ADD1.NEXT)
QUEUE)),

```

which simplifies, applying LST.APR, NLST.APR, APPLY2.EQUAL, APPLY1.SEQNO, PMAPP.NLST, LOWER.WITH.E.IF, IN.DOMAIN.WITH.E.IF, EQUAL.REACH.ZERO, SUB1.ADD1, FOLLOWS.UPPER.ADD1, UPPER.WITH.E.IF, and PMAPP.LATEST, and expanding the definitions of EQUAL, LESSP, CONSISTENT, CONSISTENT2, PMAPP, RECEIVER.INT, and LATEST, to 23 new conjectures:

Case 23. (IMPLIES (AND (PMAPP QUEUE)

```

(NOT (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  (PKT.IN)))
 (NOT (IN (ADD1 (SEQNO PKT))
           (DOMAIN QUEUE)))
 (PKTP PKT)
 (NOT (SEQP QUEUE)))
 (EQUAL QUEUE (QUOTE (1QUOTE NULL))))),

```

which again simplifies, expanding the function PMAPP, to:

T.

Case 22. (IMPLIES (AND (PMAPP QUEUE)

```

(NOT (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  (PKT.IN)))
 (NOT (IN (ADD1 (SEQNO PKT))
           (DOMAIN QUEUE)))
 (PKTP PKT)
 (SEQP QUEUE)
 (SEQP (NLST QUEUE)))
 (PAIRP (LST (NLST QUEUE)))),)

```

which we again simplify, expanding the definition of PMAPP, to:

T.

Case 21. (IMPLIES (AND (PMAPP QUEUE)

```

(NOT (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL))

```

```

          PKT.IN))
(NOT (IN (ADD1 (SEQNO PKT))
          (DOMAIN QUEUE)))
(PKTP PKT)
(SEQP QUEUE)
(SEQP (NLST QUEUE)))
(CLESSP (DOM (LST (NLST QUEUE)))
          (DOM (LST QUEUE))),
```

which again simplifies, expanding the function PMAPP, to:

T.

Case 20.(IMPLIES (AND (PMAPP QUEUE)
 (NOT (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)))
 (NOT (IN (ADD1 (SEQNO PKT))
 (DOMAIN QUEUE)))
 (PKTP PKT)
 (SEQP QUEUE)
 (NOT (SEQP (NLST QUEUE))))
 (EQUAL (NLST QUEUE)
 (QUOTE (QUOTE NULL)))),

which again simplifies, expanding PMAPP, to:

T.

Case 19.(INPUTS (AND (PMAPP QUEUE)
 (NOT (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)))
 (NOT (IN (ADD1 (SEQNO PKT))
 (DOMAIN QUEUE)))
 (PKTP PKT)
 (SEQP QUEUE))
 (MPATRP (LST QUEUE))),

which we again simplify, unfolding PMAPP, to:

T.

Case 18.(INPUTS

```

(AND (PMAAPP QUEUE)
     (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 (SEQNO PKT))))
     (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))
     (NOT (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
     (EQUAL (SEQNO PKT) 0)
     (EQUAL SINK (QUOTE (1QUOTE NULL)))
     (NOT (IN (ADD1 (SEQNO PKT))
              (DOMAIN QUEUE)))
     (PKTP PKT)
     (NOT (SEQP QUEUE)))
     (EQUAL QUEUE (QUOTE (1QUOTE NULL)))),

```

which we again simplify, unfolding the definition of PMAAPP, to:

T.

Case 17.(IMPLIES

```

(AND (PMAAPP QUEUE)
     (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 (SEQNO PKT))))
     (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))
     (NOT (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
     (EQUAL (SEQNO PKT) 0)
     (EQUAL SINK (QUOTE (1QUOTE NULL)))
     (NOT (IN (ADD1 (SEQNO PKT))
              (DOMAIN QUEUE)))
     (PKTP PKT)
     (SEQP QUEUE)
     (SEQP (NLST QUEUE)))
     (MPAIRP (LST (NLST QUEUE))))),

```

This simplifies again, unfolding the function PMAAPP, to:

T.

Case 16.(IMPLIES

```
(AND (PMAAPP QUEUE))
```

```

(FOLLOWERS QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
    (ADD1 (SEQNO PKT))))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    (SEQNO PKT)))
  (NOT (IN 0
    (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (EQUAL (SEQNO PKT) 0)
  (EQUAL SINK (QUOTE (1QUOTE NULL)))
  (NOT (IN (ADD1 (SEQNO PKT))
    (DOMAIN QUEUE)))
  (PKTP PKT)
  (SEQP QUEUE)
  (SEQP (NLST QUEUE)))
  (LESSP (DOM (LST (NLST QUEUE)))
    (DOM (LST QUEUE))).

```

However this simplifies again, unfolding the function PMAPP, to:

t.

Case 15. (IMPLIES

```

  (AND (PMAPP QUEUE)
    (FOLLOWERS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
        (SEQNO PKT)))
      (NOT (IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
      (EQUAL (SEQNO PKT) 0)
      (EQUAL SINK (QUOTE (1QUOTE NULL)))
      (NOT (IN (ADD1 (SEQNO PKT))
        (DOMAIN QUEUE)))
      (PKTP PKT)
      (SEQP QUEUE)
      (NOT (SEQP (NLST QUEUE))))
      (EQUAL (NLST QUEUE)
        (QUOTE (1QUOTE NULL)))),

```

which we again simplify, expanding the definition of PMAPP, to:

f.

Case 14.(IMPLIES

```

(AND (PMAAPP QUEUE)
      (FOLLOWS QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                  (ADD1 (SEQNO PKT))))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))))

      (NOT (IN 0
                  (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))))

      (EQUAL (SEQNO PKT) 0)
      (EQUAL SINK (QUOTE (1QUOTE NULL)))
      (NOT (IN (ADD1 (SEQNO PKT))
                  (DOMAIN QUEUE)))))

      (PKTP PKT)
      (SEQP QUEUE))
      (MPAIRP (LST QUEUE))),
```

which again simplifies, unfolding the definition of PMAAPP, to:

T.S

Case 13.(IMPLIES

```

(AND (PMAAPP QUEUE)
      (FOLLOWS QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                  (ADD1 (SEQNO PKT))))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))))

      (NOT (IN 0
                  (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))))

      (EQUAL (SEQNO PKT) 0)
      (EQUAL SINK (QUOTE (1QUOTE NULL)))
      (NOT (IN (ADD1 (SEQNO PKT))
                  (DOMAIN QUEUE)))))

      (PKTP PKT)
      (CONSISTENT2 (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN PKT))
      (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN))

      (IN (MPATR (SEQNO PKT) PKT)
          (LATEST (QUOTE SEQNO) PKT.TN)))
      (EQUALS (APR SINK (MSG PKT)))
```

```
(FAPPLY (QUOTE MSSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                      (SEQNO PKT))))).
```

This again simplifies, rewriting with PMAPP.LATEST and IN.DOMAIN.MPAIR, and opening up the definitions of EQUAL and LESSP, to:

1.6

Case 12. (IMPLIES

```
(AND (PMAPP QUEUE)
      (FOLLOWS QUEUE
        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
               (ADD1 (SEQNO PKT))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                   (SEQNO PKT)))
      (NOT (IN 0
              (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
      (EQUAL (SEQNO PKT) 0)
      (EQUAL SINK (QUOTE (1QUOTE NULL)))
      (NOT (IN (ADD1 (SEQNO PKT))
                (DOMAIN QUEUE)))
      (PKTP PKT)
      (CONSISTENT2 (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT.IN PKT))
      (CONSISTENT (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT.IN))
      (NOT (IN (MPAIR (SEQNO PKT) PKT)
                (LATEST (QUOTE SEQNO) PKT.IN)))
      (EQUAL (APR.SINK (MSSG PKT))
      (FAPPLY (QUOTE MSSG)
              (RANGE (APR (QUOTE (1QUOTE NULL))
                           (MPAIR 0 PKT)))))).
```

But this again simplifies, applying PMAPP.LATEST, IN.MPAIR.DOMAIN, RUG.MPAIR, LST.APR, NLST.APR, APPLY1.MSSG, and APR.EQUAL, and expanding EQUAL, LESSP, RANGE, and FAPPLY, to:

```
(IMPLIES
  (AND (PMAPP QUEUE)
    (FOLLOWS QUEUE
```

```

(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 1))
(CNOT (IN 0
  (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
(EQUAL (SEQNO PKT) 0)
(CNOT (IN 1 (DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  (PKT.IN))
(EQUAL (QUOTE (1QUOTE NULL))
  (FAPPLY (QUOTE MSG)
    (QUOTE (1QUOTE NULL)))),

```

which we again simplify, expanding the functions SEQN, FAPPLY, and EQUAL, to:

T.s

Case 11. (IMPLTES

```

  (AND (PNAAPP QUEUE)
    (FOLLOWNS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
        (SEQNO PKT)))
      (NOT (IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
        (EQUAL (SEQNO PKT) 0)
        (EQUAL SINK (QUOTE (1QUOTE NULL)))
        (NOT (IN (ADD1 (SEQNO PKT))
          (DOMAIN QUEUE)))
          (PKTP PKT)
          (CONSISTENT2 (QUOTE SEQNO)
            (QUOTE EQUAL)
            (PKT.IN PKT))
          (CONSISTENT (QUOTE SEQNO)
            (QUOTE EQUAL)
            (PKT.IN)))
        (NOT
          (LESSP (SUB1 (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)

```

```
(SEQNO PKT)
PKT))))  

(SEQNO PKT)))).
```

This again simplifies, rewriting with PMAPP.LATEST and
WITH.EQD.2, and expanding the definitions of EQUAL and LESSP,
to:

T.

Case 10. (IMPLIES

```
(AND
  (PMAPP.QUEUE)
  (FOLLOWS.QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
           (ADD1 (SEQNO PKT))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                 (SEQNO PKT)))
        (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
        (NOT (EQUAL (SEQNO PKT) 0)))
        (EQUAL SINK
          (FAPPLY (QUOTE ASIG)
                  (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                                (SUB1 (SEQNO PKT))))))
        (NOT (IN (ADD1 (SEQNO PKT))
                  (DOMAIN.QUEUE)))
        (PKTP PKT)
        (NOT (SEQP.QUEUE)))
      (EQUAL.QUEUE (QUOTE (QUOTE NULL))))).
```

which again simplifies, expanding the definition of PMAPP, to:

T.

Case 9. (IMPLIES

```
(AND
  (PMAPP.QUEUE)
  (FOLLOWS.QUEUE
    (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
           (ADD1 (SEQNO PKT))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                 (SEQNO PKT)))
        (IN 0
```

```

(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0))
(EQUAL SINK
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (SUB1 (SEQNO PKT))))))
(NOT (IN (ADD1 (SEQNO PKT))
  (DOMAIN QUEUE)))
(PKTP PKT)
(SEQP QUEUE)
(SEQP (NLST QUEUE))
(CPATRP (LST (NLST QUEUE))).

```

But this again simplifies, unfolding the function PMAPP, to:

T.S

Case 8. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN)
 (SEQNO PKT))))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG)
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SEQNO PKT))))))
 (NOT (IN (ADD1 (SEQNO PKT))
 (DOMAIN QUEUE)))
 (PKTP PKT)
 (SEQP QUEUE)
 (SEQP (NLST QUEUE))
 (LFSSP (DUM (LST (NLST QUEUE)))
 (DOM (LST QUEUE))))).

This again simplifies, expanding the definition of PMAPP, to:

T.

Case 7. (IMPLIES

```

    (AND
      (PMAPP QUEUE)
      (FOLLOWS QUEUE
        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
               (ADD1 (SEQNO PKT))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))
      (TN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
      (NOT (EQUAL (SEQNO PKT) 0))
      (EQUAL SINK
        (FAPPLY (QUOTE MSG)
          (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                        (SUB1 (SEQNO PKT))))))
      (NOT (IN (ADD1 (SEQNO PKT))
                (DOMAIN QUEUE)))
      (PKTP PKT)
      (SEQP QUEUE)
      (NOT (SEQP (NLST QUEUE))))
      (EQUAL (NLST QUEUE)
        (QUOTE (1QUOTE NULL)))),
    )
```

which we again simplify, unfolding the function PMAPP, to:

T.

Case 6. (IMPLIES

```

    (AND
      (PMAPP QUEUE)
      (FOLLOWS QUEUE
        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
               (ADD1 (SEQNO PKT))))
      (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                  (SEQNO PKT)))
      (TN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
      (NOT (EQUAL (SEQNO PKT) 0))
      (EQUAL SINK
        (FAPPLY (QUOTE MSG)
          (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                        (SUB1 (SEQNO PKT))))))
      (NOT (IN (ADD1 (SEQNO PKT))
                (DOMAIN QUEUE)))
      (PKTP PKT)
      (SEQP QUEUE)
      (NOT (SEQP (NLST QUEUE))))
      (EQUAL (NLST QUEUE)
        (QUOTE (1QUOTE NULL)))),
    )
```

```
(PKTP PKT)
(SEQP QUEUE)
(CPATRP (LST QUEUE)).
```

This simplifies again, unfolding the function PMAPP, to:

T.S

```
Case 5. (IMPLIES
  (AND
    (PMAPP QUEUE)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      (SEQNO PKT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL (SEQNO PKT) 0))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
          (SUB1 (SEQNO PKT)))))))
    (NOT (IN (ADD1 (SEQNO PKT))
      (DOMAIN QUEUE)))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.IN PKT)
    (CONSISTENT (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.IN)
    (IN (MPAIR (SEQNO PKT) PKT)
      (LATEST (QUOTE SEQNO) PKT.IN)))
    (EQUAL (APR SINK (MSG PKT))
      (FAPPLY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
          (SEQNO PKT)))))),
```

which again simplifies, clearly, to the new formula:

```
(IMPLIES
  (AND (PMAPP QUEUE)
    (FOLLOWS QUEUE
```

```

(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
       (ADD1 (SEQNO PKT)))))

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
             (SEQNO PKT)))))

(IN 0
   (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
  (NOT (EQUAL (SEQNO PKT) 0)))
  (NOT (IN (ADD1 (SEQNO PKT))
            (DOMAIN QUEUE))))
  (PKTP PKT)
  (CONSISTENT2 (QUOTE SEQNO)
                (QUOTE EQUAL)
                (PKT.IN PKT))
  (CONSISTENT (QUOTE SEQNO)
                (QUOTE EQUAL)
                (PKT.IN))
  (IN (MPAIR (SEQNO PKT) PKT)
      (LATEST (QUOTE SEQNO) PKT.IN)))

(EQUAL
  (APP (FAPPLY (QUOTE MSSG)
                (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                               (SUB1 (SEQNO PKT))))))

  (MSSG PKT))
  (FAPPLY (QUOTE MSSG)
          (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                         (SEQNO PKT)))))).
```

Applying the lemmas MSSG/SEQNO.ELIM and SUB1.ELIM, replace PKT by (PACKET Z X) to eliminate (SEQNO PKT) and (MSSG PKT) and X by (ADD1 V) to eliminate (SUB1 X). We rely upon the type restriction lemma noted when SEQNO was introduced and the type restriction lemma noted when SUB1 was introduced to restrict the new variables. The result is:

```

(IMPLIES
  (AND (NUMBERP V)
       (PKAPP QUEUE)
       (FOLLOWS QUEUE
                  (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                         (ADD1 (ADD1 V)))))
       (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                   (ADD1 V)))))

  (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
```

```

(NOT (EQUAL (ADD1 V) 0))
(NOT (IN (ADD1 (ADD1 V)) (DOMAIN QUEUE)))
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN
  (PACKET Z (ADD1 V)))
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN)
(IN (MPAIR (ADD1 V) (PACKET Z (ADD1 V)))
  (LATEST (QUOTE SEQNO) PKT.IN)))
(EQUAL
  (APR (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      V))))
  Z)
  (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 V))))),

```

which further simplifies, rewriting with SUB1.ADD1,
 EQUAL.REACH.ZERO, PMAPP.LATEST, LOWER.ADD1, RNG.MPAIR, LST.APR,
 NLST.APR, MSG.PACKET, and APPLY1.MSG, and opening up the
 functions LESSP, RANGE, and FAPPLY, to:

T.S

Case 4. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      (SEQNO PKT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL (SEQNO PKT) 0)))
    (EQUAL SINK
      (FAPPLY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
          (SUB1 (SEQNO PKT))))))
    (NOT (IN (ADD1 (SEQNO PKT))
      (DOMAIN QUEUE)))

```

```

(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN)
(NOT (IN (MPAIR (SEQNO PKT) PKT)
  (LATEST (QUOTE SEQNO) PKT.IN)))
EQUAL
(APR SINK (MSSG PKT))
(FAPPLY (QUOTE MSSG)
  (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
    (SUB1 (SEQNO PKT))))
    (MPAIR (SEQNO PKT) PKT)))),

```

which again simplifies, applying RNG.MPAIR, LST.APR, NLST.APR, and APPLY1.MSSG, and opening up the definitions of RANGE and FAPPLY, to:

T.S

Case 3. (IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT)))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      (SEQNO PKT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL (SEQNO PKT) 0)))
  EQUAL SINK
    (FAPPLY (QUOTE MSSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        (SUB1 (SEQNO PKT)))))
    (NOT (IN (ADD1 (SEQNO PKT))
      (DOMAIN QUEUE)))
  (PKTP PKT)
  (CONSISTENT2 (QUOTE SEQNO)
    (QUOTE EQUAL)
    PKT.IN PKT)
  (CONSISTENT (QUOTE SEQNO))

```

```

(QUOTE EQUAL)
PKT.IN))
(NOT
(LESSP (SUB1 (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
(SEQNO PKT)
PKT)))
(SEQNO PKT))).

```

This again simplifies, using linear arithmetic and applying LESSP.REACH.WITHE.2 and PMAPP.LATEST, to:

```

(IMPLIES
(AND (EQUAL (REACH (LATEST (QUOTE SEQNO) PKT.IN))
(SEQNO PKT))
(PMAPP QUEUE)
(FOLLOWS QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 (SEQNO PKT))))
(NOT (LESSP (SEQNO PKT) (SEQNO PKT)))
(CTM 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0))
(NOT (IN (ADD1 (SEQNO PKT))
(DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN)))
(NOT
(LESSP (SUB1 (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
(SEQNO PKT)
PKT)))
(SEQNO PKT))).

```

This simplifies again, using linear arithmetic and applying LESSP.SUB1.REACH.WITHE and PMAPP.LATEST, to:

```

(IMPLIES
(AND (EQUAL (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
(SEQNO PKT)
PKT)))

```

```

    0)
(EQUAL (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      (SEQNO PKT))
(PMAPP.QUEUE)
(FOLLOWNS.QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))
(NOT (LESSP (SECOND PKT) (SECOND PKT)))
(IN 0
  (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0))
(NOT (IN (ADD1 (SEQNO PKT))
          (DOMAIN.QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN)))
(NOT
  (LESSP (SUB1 (REACH (#ITHE (LATEST (QUOTE SEQNO) PKT.IN)
                           (SEQNO PKT)
                           PKT)))
            (SECOND PKT)))),

```

which we again simplify, using linear arithmetic and applying LESSP.REACH.#ITHE.2 and PMAPP.LATEST, to:

T.S

Case 2. (IMPLIES
 (AND
 (PMAPP.QUEUE)
 (FOLLOWNS.QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 (SEQNO PKT))))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0)))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG)))
)
)

```

(RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
(SUB1 (SEQNO PKT)))))

(NOT (IN (ADD1 (SEQNO PKT))
(DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN)
(LESSP (SUB1 (SEQNO PKT))
(ADD1 (SEQNO PKT))))
(FOLLOWS QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 (ADD1 (SEQNO PKT))))),

```

which we again simplify, applying SUB1.ADD1, PMAPP.LATEST, and FOLLOWS.UPPER.ADD1, and expanding the function LESSP, to:

T.

Case 1. (IMPLIES

```

(CAND
(PMAPP QUEUE)
(FOLLOWS QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 (SEQNO PKT))))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
(SEQNO PKT)))
(IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0))
(EQUAL SINK
(FAPPLY (QUOTE MSG)
(RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
(SUB1 (SEQNO PKT)))))))
(NOT (IN (ADD1 (SEQNO PKT))
(DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO))

```

```

        (QUOTE EQUAL)
        PKT.IN)
(CNOT (LESSP (SUB1 (SEQNO PKT)) ,
              (ADD1 (SEQNO PKT)))))

;FOLDUPS.QUEUE
(*ITHE (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                (ADD1 (ADD1 (SEQNO PKT)))))

                (SEQNO PKT)
                PKT))),
```

which we again simplify, using linear arithmetic, to:

1

Q.E.D.

390914 conses
360.125 seconds
44.228 seconds, garbage collection time

Proof of VC "RECEIVER#7"

```

(IMPLIES
  (AND (RECEIVER.TNT PKT.IN SINK ACK.OUT NEXT QUEUE)
       (IN (ADD1 NEXT) (DOMAIN QUEUE))
       (EQUAL (SECOND PKT) NEXT)
       (PKTP PKT))
   (RECEIVER.INI (APR PKT.IN PKT)
                 (JOIN SINK
                       (APL (MSSG PKT)
                            (FAPPLY (QUOTE MSSG)
                                   (RANGE (CONSEC QUEUE))))))
   (APR ACK.OUT (REACH QUEUE))
   (REACH QUEUE)
   (UPPER QUEUE (REACH QUEUE))))
```

This conjecture can be simplified, using the abbreviations RECEIVER.INT, AND, IMPLIES, and JOIN.APR.APR, to:\$\$\$\$

```

(IMPLIES
  (AND
    (PMAPP QUEUE)
    (NUMBERP NEXT))
  (IF
    (CONSISTENT (QUOTE SEQNO)
                 (QUOTE EQUAL)
                 PKT.IN)
    (IF
      (FOLLOWS QUEUE
        (CUPPER (LATEST (QUOTE SEQNO) PKT.IN)
                (ADD1 NEXT)))
      (IF
        (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
               NEXT)
        F
        (IF
          (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
          (IF
            (LESSP 0 NEXT)
            (EQUAL SINK
                  (FAPPLY (QUOTE MSSG)
                          (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                                      (SUB1 NEXT))))))
            F)
          (IF (EQUAL NEXT 0)
              (EQUAL SINK (QUOTE (1QUOTE NULL)))
              F)))
        F)
      T)
    (IN (ADD1 NEXT) (DOMAIN QUEUE))
    (EQUAL (SEQNO PKT) NEXT)
    (PKTP PKT))
  RECEIVER.INT (APR PKT.IN PKT)
    (JOIN (APR SINK (MSSG PKT))
          (FAPPLY (QUOTE MSSG)
                  (RANGE (CONSEC QUEUE))))
    (APR ACK.OUT (REACH QUEUE))
    (REACH QUEUE))

```

(UPPER QUEUE (REACH QUEUE))),

which simplifies, using linear arithmetic, rewriting with the lemmas LST.APR, NLST.APR, APPLY2.EQUALS, APPY1.SEQNO, PHAPP.UPPER, SUB1.ADD1, FIRST.DOMAIN.FOLLOWS.UPPER, LOWER.SUB1.REACH.3, IN.DOMAIN.FOLLOWS.UPPER, LOWER.WITH.E.IF, LESSP.ZERO.REACH, IN.DOMAIN.WITH.E.IF, FOLLOWS.UPPER.ADD1, IN.DOMAIN.UPPER, IN.REACH.DOMAIN, FOLLOWS.TRANS, FOLLOWS.SAME, FOLLOWS.UPPER, FOLLOWS.UPPER.UPPER, UPPER.WITH.E.IF, PHAPP.LATEST, WITH.E.JOIN.LOWER, FIRST.DOMAIN.CONSEC, and PHAPP.CONSEC, and expanding the functions EQUAL, LESSP, CONSISTENT, CONSISTENT2, RECEIVER.IN, and LATEST, to 15 new conjectures:

Case 15.(IMPLIES

```

    (AND (PMAPP QUEUE)
        (FOLLOWS QUEUE
            (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                (ADD1 (SEQNO PKT))))
            (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                (SEQNO PKT)))
            (NOT (IN 0
                (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
            (EQUAL (SEQNO PKT) 0)
            (EQUAL SINK (QUOTE (1QUOTE NULL)))
            (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
            (PKTP PKT)
            (CONSISTENT2 (QUOTE SEQNO)
                (QUOTE EQUAL)
                PKT.IN PKT)
            (CONSISTENT (QUOTE SEQNO)
                (QUOTE EQUAL)
                PKT.IN)
            (EQUAL (SEQNO PKT)
                (SUB1 (REACH QUEUE)))
            (NOT (IN (MPAIR (SEQNO PKT) PKT)
                (LATEST (QUOTE SEQNO) PKT.IN)))
            (EQUAL (SUB1 (REACH QUEUE)) 0))
            (EQUAL (JOIN (APR SINK (MSSG PKT))
                (FAPPLY (QUOTE MSSG)
                    (RANGE (CONSEC QUEUE))))
            (FAPPLY (QUOTE MSSG)
                (RANGE (APR (QUOTE (1QUOTE NULL)
                    (MPAIR 0 PKT))))),

```

which again simplifies, applying the lemma IN.ADD1.DOMAIN, and expanding the functions EQUAL and LESSE, to:

1

Case 14. IMPLIES

```

(CAND (MPAIR SEQNO)
      (FOLLOWS QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                  (ADD1 (SEQNO PKT)))))

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
             (SEQNO PKT)))

(NOT (IN 0
          (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))))

(EQUAL (SEQNO PKT) 0)
(EQUAL SINK (QUOTE (QUOTE NULL)))
(IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
(PKTP PKT)

(CONSISTENT2 (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN PKT))

(CONSISTENT (QUOTE SEQNO)
              (QUOTE EQUAL)
              (PKT.IN))

(EQUAL (SEQNO PKT)
       (SUB1 (REACH QUEUE)))
(NOT (IN (MPAIR (SEQNO PKT) PKT)
          (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SUB1 (REACH QUEUE)) 0)))

(EQUAL
  (JOIN (APR SINK (MSSG PKT))
        (FAPPLY (QUOTE MSSG)
                (RANGE (CONSEC QUEUE))))
  (FAPPLY (QUOTE MSSG)
          (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                           (SUB1 (SUB1 (REACH QUEUE)))))))
          (MPAIR (SEQNO PKT))))),

```

which again simplifies, using linear arithmetic, to:

15

Case 13. (IMPLIES)

(AND (PMAAPP QUEUE))

```

(FOLLOWERS.QUEUE
  (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
         (ADD1 (SEQNO PKT))))
  (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
              (SEQNO PKT)))
  (NOT (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
  (EQUAL (SEQNO PKT) 0)
  (EQUAL SINK (QUOTE (QUOTE NULL)))
  (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
  (PKTP PKT)
  (CONSISTENT2 (QUOTE SEQNO)
                (QUOTE EQUAL)
                (PKT.IN PKT))
  (CONSISTENT (QUOTE SEQNO)
                (QUOTE EQUAL)
                (PKT.IN))
  (EQUAL (SEQNO PKT)
         (SUB1 (REACH QUEUE)))
  (IN (MPAIR (SEQNO PKT) PKT)
      (LATEST (QUOTE SEQNO) PKT.IN)))
  (EQUAL
    (JOIN (APR STNK (MSSG PKT))
          (FAPPLY (QUOTE MSSG)
                  (RANGE (CONSEC QUEUE))))
    (FAPPLY (QUOTE MSSG)
            (RANGE (JOIN (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                                (SEQNO PKT))
                          (CONSEC QUEUE))))).

```

However this simplifies again, using linear arithmetic, applying LESSP.REACH.REACH.1, IN.DOMAIN.MPAIR, PMAPP.LATEST, IN.DOMAIN.REACH, and IN.ADD1.DOMAIN, and unfolding EQUAL and LESSP, to:

T.S

Case 12. (IMPLIES

```

  (AND (PMAPP.QUEUE)
        (FOLLOWERS.QUEUE
          (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                 (ADD1 (SEQNO PKT))))
        (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                    (SEQNO PKT)))

```

```

(NOT (IN 0
          (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
(EQUAL (SEQNO PKT) 0)
(EQUAL SINK (QUOTE (QUOTE NULL)))
(IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
              (QUOTE EQUAL)
              PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
              (QUOTE EQUAL)
              PKT.IN)
(NOT (EQUAL (SEQNO PKT)
            (SUB1 (REACH QUEUE)))))

(EQUAL
  (JOIN (APR SINK (MSSG PKT))
        (FAPPLY (QUOTE MSSG)
                (RANGE (CONSEC QUEUE)))))

(FAPPLY
  (QUOTE MSSG)
  (RANGE (WITHE (JOIN (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                           (SEQNO PKT))
                        (CONSEC QUEUE))
                     (SEQNO PKT)
                     PKT)))),,

```

which we again simplify, rewriting with PMAPP.LATEST,
 LOWER.ZERO.2, PMAPP.CONSEC, JOIN.NULL.PMAPP, WITHE.ZERO.2,
 IN.DOMAIN.FOLLOW.SUPER, IN.DOMAIN.CONSEC.NOT, RANGE.JOIN,
 RNG.XPAIR, LST.APR, NLST.APR, FAPPLY.JOIN, and APPYL1.MSSG, and
 expanding the definitions of EQUAL, LESSP, RANGE, and FAPPLY, to:

```

\$      (IMPLIES
          (AND (PMAPP QUEUE)
               (FOLLOWS QUEUE
                  (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                         1)))
               (NOT (IN 0
                      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
               (EQUAL (SEQNO PKT) 0)
               (IN 1 (DOMAIN QUEUE))
               (PKTP PKT)
               (CONSISTENT2 (QUOTE SEQNO)
                             (QUOTE EQUAL)))

```

```

    PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN)
(NOT (EQUAL 0 (SUB1 (REACH QUEUE))))))
(EQUAL (JOIN (APR (QUOTE (1QUOTE NULL)) (MSG PKT))
(FAPPLY (QUOTE MSG)
(RANGE (CONSEC QUEUE))))))
(JOIN (APR (FAPPLY (QUOTE MSG)
(QUOTE (1QUOTE NULL)))
(MSG PKT))
(FAPPLY (QUOTE MSG)
(RANGE (CONSEC QUEUE))))),

```

which again simplifies, opening up SEQP and FAPPLY, to:

T.

Case 11.(IMPLIES

```

(AND (PMAPP QUEUE)
(FOLLOWS QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 (SEQNO PKT))))))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
(SEQNO PKT))))
(NOT (IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
(EQUAL (SEQNO PKT) 0)
(EQUAL STINK (QUOTE (1QUOTE NULL)))
(IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN))
(SEQP QUEUE)).

```

This again simplifies, applying the lemmas UPPER.ZERO, PMAPP.LATEST, and FOLLOWS.UPPER.ADD1, and opening up the definitions of PMAPP, FOLLOWS, IN, DOMAIN, EQUAL, and LESSP, to:

T.6

Case 10. (IMPLIES

```

    (AND (PMAPP QUEUE)
        (FOLLOWS QUEUE
            (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                (ADD1 (SEQNO PKT))))
        (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
            (SEQNO PKT)))
        (NOT (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
        (EQUAL (SEQNO PKT) 0)
        (EQUAL SINK (QUOTE (1QUOTE NULL)))
        (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
        (PKTP PKT)
        (CONSISTENT2 (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN PKT)
        (CONSISTENT (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN))
        (NOT (LESSP (REACH (WITHE (LATEST (QUOTE SEQNO) PKT.IN)
            (SEQNO PKT)
            PKT))
            (REACH QUEUE)))).

```

This again simplifies, rewriting with PMAPP.LATEST, WITHE.ZERO.2, and REACH.JOIN.APR.NULL.2, and expanding EQUAL and LESSP, to:

```

(implies
    (AND (PMAPP QUEUE)
        (FOLLOWS QUEUE
            (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                1))
        (NOT (IN 0
            (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN))))
        (EQUAL (SEQNO PKT) 0)
        (IN 1 (DOMAIN QUEUE))
        (PKTP PKT)
        (CONSISTENT2 (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN PKT)
        (CONSISTENT (QUOTE SEQNO)
            (QUOTE EQUAL)
            PKT.IN))

```

```
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
    (REACH QUEUE)))).
```

However this simplifies again, using linear arithmetic and applying LESSP, REACH, REACH.i and PMAPP,LATEST, to:

T.

Case 9. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 (SEQNO PKT)))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SEQNO PKT))))))
 (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
 (PKTP PKT)
 (CONSISTENT2 (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN PKT))
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN))
 (EQUAL (SEQNO PKT)
 (SUB1 (REACH QUEUE)))
 (NOT (IN (PAIR (SEQNO PKT) PKT)
 (LATEST (QUOTE SEQNO) PKT.IN)))
 (EQUAL (SUB1 (REACH QUEUE)) 0))
 (EQUAL (JOIN (APR SINK (MSG PKT))
 (FAPPLY (QUOTE MSG))
 (RANGE (CONSEC QUEUE))))
 (FAPPLY (QUOTE MSG))
 (RANGE (APR (QUOTE (QUOTE NULL))
 (PAIR 0 PKT))))).

But this again simplifies, using linear arithmetic, to:

T.S

Case 8. (INPUTS
 (AND
 (PKTIN QUEUE)
 (FOLLOWS SEQNO
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 (SEQNO PKT))))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0)))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG))
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SEQNO PKT))))))
 (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
 (PKTP PKT)
 (CONSISTENT2 (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN PKT))
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 (PKT.IN))
 (EQUAL (SEQNO PKT)
 (SUB1 (REACH QUEUE)))
 (NOT (IN (PAIR (SEQNO PKT) PKT)
 (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SUB1 (REACH QUEUE)) 0)))
 (EQUAL
 (JOIN (APR SINK (MSG PKT))
 (FAPPLY (QUOTE MSG))
 (RANGE (CONSEC QUEUE))))
 (FAPPLY (QUOTE MSG))
 (RANGE (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SUB1 (REACH QUEUE))))
 (PAIR (SEQNO PKT) PKT))))).

This again simplifies, applying IN.ADD1.DOMAIN, to:

T.

Case 7. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (FOLLOWS QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 (SEQNO PKT)))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG)
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SEQNO PKT))))))
 (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
 (PKTP PKT)
 (CONSISTENT2 (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN PKT)
 (CONSISTENT (QUOTE SEQNO)
 (QUOTE EQUAL)
 PKT.IN)
 (EQUAL (SEQNO PKT)
 (SUB1 (REACH QUEUE)))
 (IN (MPAIR (SEQNO PKT) PKT)
 (LATEST (QUOTE SEQNO) PKT.IN)))
 (EQUAL
 (JOIN (APR SINK (MSG PKT))
 (FAPPLY (QUOTE MSG)
 (RANGE (CONSEC QUEUE))))
 (FAPPLY (QUOTE MSG)
 (RANGE (JOIN (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SEQNO PKT))
 (CONSEC QUEUE))))),

which we again simplify, applying IN.ADD1.DOMAIN, to:

T.S

Case 6. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (FOLLOWS QUEUE

```

        (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (SEQNO PKT))))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
(SEQNO PKT)))
(IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0))
(EQUAL SINK
(FAPPLY (QUOTE MSG)
(RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
(SUB1 (SEQNO PKT)))))))
(IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
(PKT.IN PKT))
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
(PKT.IN))
(NOT (EQUAL (SEQNO PKT)
(SUB1 (REACH QUEUE))))
(LESSP (SUB1 (REACH QUEUE))
(SEQNO PKT)))
(EQUAL
(JOIN (APR STNK (MSG PKT))
(FAPPLY (QUOTE MSG)
(RANGE (CONSEC QUEUE))))
(FAPPLY (QUOTE MSG)
(RANGE (JOIN (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
(SEQNO PKT))
(CONSEC QUEUE))))),

```

which again simplifies, applying RANGE.JOIN and FAPPLY.JOIN, to:

```

(implies
(AND (PAPP QUEUE)
(FOLLOWS QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 (SEQNO PKT))))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
(SEQNO PKT)))
(IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0)))

```

```

(IN (ADD1 (SECOND PKT)) (DOMAIN QUEUE))
(PKTP PKT)
(UNSTSTNT? (QUOTE SECOND)
    (QUOTE EQUAL)
    (PKT.IN PKT))
(CONSISTENT (QUOTE SECOND)
    (QUOTE EQUAL)
    (PKT.IN))
(NOT (EQUAL (SECOND PKT)
    (SUB1 (REACH QUEUE))))
(LESSP (SUB1 (REACH QUEUE))
    (SECOND PKT)))

(EQUAL
  (JOIN
    (APR (FAPPY (QUOTE MSG)
        (RANGE (LOWER (LATEST (QUOTE SECOND) PKT.IN)
            (SUB1 (SECOND PKT)))))

        (MSG PKT))
      (FAPPY (QUOTE MSG)
        (RANGE (CONSEC QUEUE))))))

  (JOIN (FAPPY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SECOND) PKT.IN)
          (SECOND PKT)))))

    (FAPPY (QUOTE MSG)
      (RANGE (CONSEC QUEUE))))).

```

Appealing to the lemmas MSG/SECOND.ELIM and SUB1.ELIM, replace PKT by (PACKET Z X) to eliminate (SECOND PKT) and (MSG PKT) and X by (ADD1 V) to eliminate (SUB1 X). we rely upon the type restriction lemma noted when SECOND was introduced and the type restriction lemma noted when SUB1 was introduced to constrain the new variables. This generates:

```

(IMPLIES
  (AND (NUMBERP V)
    (PMAFP QUEUE)
    (FOLLOWS QUEUE
      (UPPER (LATEST (QUOTE SECOND) PKT.IN)
          (ADD1 (ADD1 V)))))
    (NOT (LESSP (REACH (LATEST (QUOTE SECOND) PKT.IN))
        (ADD1 V))))
  (IN 0
    (DOMAIN (LATEST (QUOTE SECOND) PKT.IN)))
  (NOT (EQUAL (ADD1 V) 0)))

```

```

(IN (ADD1 (ADD1 V)) (DOMAIN QUEUE))
(CONSISTENT2 (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN
  (PACKET Z (ADD1 V)))
(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN)
(NOT (EQUAL (ADD1 V) (SUB1 (REACH QUEUE))))
(LESSP (SUB1 (REACH QUEUE)) (ADD1 V)))
(EQUAL
  (JOIN
    (FAPPY (FAPPY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        V)))
      Z)
    (FAPPY (QUOTE MSG)
      (RANGE (CONSEC QUEUE))))
    (JOIN (FAPPY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 V))))
      (FAPPY (QUOTE MSG)
        (RANGE (CONSEC QUEUE))))))).

```

This simplifies further, using linear arithmetic and rewriting with LESSP.SUB1.REACH.FOLLOWs and FMAPP.LATEST, to:

```

(IMPLIES
  (AND (EQUAL (REACH QUEUE) 0)
    (NUMBERP V)
    (FMAPP QUEUE)
    (FOLLOWs QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
        (ADD1 (ADD1 V)))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
      (ADD1 V))))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL (ADD1 V) 0)))
    (IN (ADD1 (ADD1 V)) (DOMAIN QUEUE))
    (CONSISTENT2 (QUOTE SEQNO)
      (QUOTE EQUAL)
      PKT.IN
      (PACKET Z (ADD1 V))))
    )
  )

```

```

(CONSISTENT (QUOTE SEQNO)
  (QUOTE EQUAL)
  PKT.IN)
(NOT (EQUAL (ADD1 V) (SUB1 (REACH QUEUE))))
(LESSP (SUB1 (REACH QUEUE)) (ADD1 V)))
(EQUAL
  (JOIN
    (APR (FAPPLY (QUOTE MSG)
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
        V))))
    ?)
    (FAPPLY (QUOTE MSG)
      (RANGE (CONSEC QUEUE))))
  (JOIN (FAPPLY (QUOTE MSG)
    (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
      (ADD1 V))))
    (FAPPLY (QUOTE MSG)
      (RANGE (CONSEC QUEUE))))),

```

which again simplifies, applying EQUAL.REACH.ZERO,
 FOLLOWSP.QUEUE, PMAPP.LATEST, and SUB1.ADD1, and expanding
 the definitions of PMAPP, FOLLOWSP, IN, SEQN, DOMAIN, and LESSP,
 to:

T.S

Case 5. (IMPLIES
 (AND
 (PMAPP QUEUE)
 (FOLLOWSP QUEUE
 (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
 (ADD1 (SEQNO PKT))))
 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
 (SEQNO PKT))))
 (IN 0
 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
 (NOT (EQUAL (SEQNO PKT) 0)))
 (EQUAL SINK
 (FAPPLY (QUOTE MSG)
 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
 (SUB1 (SEQNO PKT)))))))
 (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
 (PKTP PKT)
 (CONSISTENT2 (QUOTE SEQNO)))
)

```

        (QUOTE EQUAL)
        (PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
        (QUOTE EQUAL)
        (PKT.IN)
(NOT (EQUAL (SEQNO PKT)
        (SUB1 (REACH QUEUE))))))
(NOT (LESSP (SUB1 (REACH QUEUE))
        (SEQNO PKT)))))

(EQUAL
(JOIN (APR SINK (MSG PKT))
        (FAPPLY (QUOTE MSG)
                (RANGE (CONSEC QUEUE)))))

(FAPPLY
        (QUOTE MSG)
        (RANGE (JOIN (APR (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                (SUB1 (SEQNO PKT))))))
                (MPAIR (SEQNO PKT) PKT)))
        (CONSEC QUEUE)))).
```

But this simplifies again, rewriting with RANGE.JOIN, RNG.MPAIR, LST.APR, NLST.APR, FAPPLY.JOIN, and APPY1.MSG, and opening up RANGE and FAPPLY, to:

I.s

Case 4. (IMPLIES

```

        (AND
        (PMAAPP QUEUE)
        (FOLLOWQUEUE
            (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                    (ADD1 (SEQNO PKT)))))

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
        (SEQNO PKT)))))

(IN 0
        (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL (SEQNO PKT) 0)))

(EQUAL SINK
        (FAPPLY (QUOTE MSG)
                (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                    (SUB1 (SEQNO PKT))))))

(IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)))
```

```

          (QUOTE EQUAL)
          PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
          (QUOTE EQUAL)
          PKT.IN))
(SEQP QUEUE)),

```

which we again simplify, rewriting with PMAPP.LATEST and FOLLOW.SUPER.ADD1, and opening up the definitions of PMAPP, FOLLOW, IN, SEQNO, and DOMAIN, to:

T.

Case 3. IMPLIES

```

  (AND
    (PMAPP QUEUE)
    (FOLLOW QUEUE
      (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
             (ADD1 (SEQNO PKT))))
    (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
                 (SEQNO PKT)))
    (IN 0
      (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
    (NOT (EQUAL (SEQNO PKT) 0)))
    (EQUAL SEQNO
      (FAPPLY (QUOTE MSG))
      (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                    (SUB1 (SEQNO PKT))))))
    (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
    (PKTP PKT)
    (CONSISTENT2 (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN PKT)
    (CONSISTENT (QUOTE SEQNO)
                  (QUOTE EQUAL)
                  PKT.IN))
    (NOT (LESSP (REACH (WITH (LATEST (QUOTE SEQNO) PKT.IN)
                           (SEQNO PKT)
                           PKT))
                  (REACH QUEUE))).

```

This simplifies again, using linear arithmetic and applying LESSP.REACH.WITH.EQUAL, PMAPP.LATEST, FOLLOW.TRANS, FOLLOW.LATEST.CONSISTENT, FOLLOW.SAME, NUMBERP, SEQNO,

FOLLOW.SUPPER, FOLLOW.WITH.E.IF, and IN.DOMAIN.FOLLOW.SUPPER, to:

I.S

Case 2. (IMPLIES

(AND

(PMAPP.QUEUE)

(FOLLOW.S QUEUE

 (UPPER (LATEST (QUOTE SEQNO) PKT.IN))

 (ADD1 (SEQNO PKT))))

 (NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))

 (SEQNO PKT))))

 (IN 0

 (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))

 (NOT (EQUAL (SEQNO PKT) 0))

 (EQUAL SINK

 (FAPPLY (QUOTE MSG))

 (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN))

 (SUB1 (SEQNO PKT))))))

 (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))

 (PKTP PKT)

 (CONSISTENT2 (QUOTE SEQNO))

 (QUOTE EQUAL)

 (PKT.IN PKT))

 (CONSISTENT (QUOTE SEQNO))

 (QUOTE EQUAL)

 (PKT.IN))

 (LESSP (SUB1 (SEQNO PKT))

 (REACH QUEUE)))

 (FOLLOW.S (UPPER.QUEUE (REACH.QUEUE))

 (UPPER (LATEST (QUOTE SEQNO) PKT.IN))

 (ADD1 (REACH.QUEUE))))).

But this again simplifies, applying PMAPP.UPPER,
 FOLLOW.UPPER.UPPER, FOLLOW.UPPER, NUMBERP.SEQNO, FOLLOW.SAME,
 FOLLOW.LATEST.CONSISTENT, FOLLOW.TRANS, PMAPP.LATEST,
 IN.REACH.DOMAIN, IN.DOMAIN.UPPER, and FOLLOW.UPPER.ADD1, to:

I.SS

Case 1. (IMPLIES

(AND

(PMAPP.QUEUE)

(FOLLOW.S QUEUE

```

(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
       (ADD1 (SEQNO PKT))))
(CNOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
              (SEQNO PKT)))
      (IN 0
          (DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
      (CNOT (EQUAL (SEQNO PKT) 0)))
      (EQUAL SINK
             (FAPPLY (QUOTE MSG)
                     (RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
                                  (SUB1 (SEQNO PKT)))))))
      (IN (ADD1 (SEQNO PKT)) (DOMAIN QUEUE))
      (PKTP PKT)
      (CONSISTENT2 (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT.IN PKT))
      (CONSISTENT (QUOTE SEQNO)
                    (QUOTE EQUAL)
                    (PKT.IN))
      (NOT (LESSP (SUB1 (SEQNO PKT))
                  (REACH QUEUE))))
      (FOLLOWS (UPPER QUEUE (REACH QUEUE))
                (WITHE (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
                               (ADD1 (REACH QUEUE)))
                        (SEQNO PKT)
                        PKT))).
```

This simplifies again, using linear arithmetic, applying FOLLOW.S.UPPER.ADD1, IN.DOMAIN.UPPER, PMAPP.LATEST, IN.DOMAIN.FOLLOW.S.UPPER, FOLLOW.S.TRANS, FOLLOW.LATEST.CONSISTENT, FOLLOW.SAME, NUMBERP.SEQNO, FOLLOW.UPPER, FOLLOW.UPPER.UPPER, PMAPP.UPPER, FOLLOW.WITHE.IF, and SUB1.ADD1, and expanding LESSP, to:

T.

Q.E.D.

434606 conses
 393.51 seconds
 54.115 seconds, garbage collection time

```

(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
NEXT))
(NOT (IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(EQUAL NEXT 0)
(EQUAL SINK (QUOTE (1QUOTE NULL)))
(NOT (EQUAL (SEQNO PKT) 0)))
(NOT (IN (SEQNO PKT) (DOMAIN QUEUE)))
(PKTP PKT)
(CONSISTENT2 (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN PKT)
(CONSISTENT (QUOTE SEQNO)
(QUOTE EQUAL)
PKT.IN)
(NOT (LESSP (SUB1 (SEQNO PKT)) NEXT)))
(FOLLOW (WITHQUEUE (SEQNO PKT) PKT)
(*ITHE (UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 NEXT))
(SEQNO PKT)
PKT))),
```

which again simplifies, applying the lemmas RNG.MPAIR, DOM.MPAIR, IN.WITHQUEUE, IN.DOMAIN.UPPER, FOLLOW.WITHQUEUE, PMAPP.WITHQUEUE, PMAPP.LATEST, PMAPP.UPPER, and FOLLOW.WITHQUEUE, and expanding the functions NUMBERP, EQUALP, and LESSP, to:

T.

Case 7. (IMPLIES

```

(AND
(PMAPP QUEUE)
(NUMBERP NEXT)
(FOLLOW QUEUE
(UPPER (LATEST (QUOTE SEQNO) PKT.IN)
(ADD1 NEXT)))
(NOT (LESSP (REACH (LATEST (QUOTE SEQNO) PKT.IN))
NEXT))
(IN 0
(DOMAIN (LATEST (QUOTE SEQNO) PKT.IN)))
(NOT (EQUAL NEXT 0))
(EQUAL SINK
(CAPPLY (QUOTE MSG)
(RANGE (LOWER (LATEST (QUOTE SEQNO) PKT.IN)
```

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