

A short note on universal quantification

With the dummy  $x$  ranging over some anonymous range and using

$$(0) \quad [(\underline{A}x:: p.x) \wedge (\underline{A}x:: q.x) \equiv (\underline{A}x:: p.x \wedge q.x)]$$

we shall show

$$(1) \quad [(\underline{A}x:: p.x \equiv q.x) \Rightarrow ((\underline{A}x:: p.x) \equiv (\underline{A}x:: q.x))]$$

$$(2) \quad [(\underline{A}x:: p.x \Rightarrow q.x) \Rightarrow ((\underline{A}x:: p.x) \Rightarrow (\underline{A}x:: q.x))]$$

From propositional calculus we shall use

$$(3) \quad [X \Rightarrow (P \equiv Q) \equiv X \wedge P \equiv X \wedge Q]$$

$$(4) \quad [(P \equiv Q) \wedge P \equiv P \wedge Q]$$

Proof of (3)  $X \Rightarrow (P \equiv Q)$

$$= \{ \text{pred. calc} \}$$

$$\neg X \vee (\neg P \equiv \neg Q)$$

$$= \{ \text{distribution} \}$$

$$\neg X \vee \neg P \equiv \neg X \vee \neg Q$$

$$= \{ \text{de Morgan} \}$$

$$X \wedge P \equiv X \wedge Q \quad \cdot \text{ (End of Proof)}$$

Proof of (4)  $(P \equiv Q) \wedge P$

$$= \{ \text{p.c.} \}$$

$$(\neg P \neq Q) \wedge P$$

$$= \{ \text{distribution} \}$$

$$\neg P \wedge P \neq Q \wedge P$$

$$= \{ \text{p.c.} \}$$

$$P \wedge Q \quad \cdot \text{ (End of Proof)}$$

Proof of (1) Denoting for brevity's sake universal quantification over  $x$  by " $\llbracket \quad \rrbracket$ ", we observe

$$\begin{aligned} & \llbracket p \equiv q \rrbracket \Rightarrow (\llbracket p \rrbracket \equiv \llbracket q \rrbracket) \\ & = \{ (3) \} \\ & \llbracket p \equiv q \rrbracket \wedge \llbracket p \rrbracket \equiv \llbracket p \equiv q \rrbracket \wedge \llbracket q \rrbracket \\ & = \{ (0) \text{ twice} \} \\ & \llbracket (p \equiv q) \wedge p \rrbracket \equiv \llbracket (p \equiv q) \wedge q \rrbracket \\ & = \{ (4) \text{ twice} \} \\ & \llbracket p \wedge q \rrbracket \equiv \llbracket q \wedge p \rrbracket \\ & = \{ p.c \} \\ & \text{true} \end{aligned}$$

(End of Proof)

Proof of (2) With the same notational convention

$$\begin{aligned} & \llbracket p \Rightarrow q \rrbracket \\ & = \{ p.c \} \\ & \llbracket p \wedge q \equiv p \rrbracket \\ & \Rightarrow \{ (1) \} \\ & \llbracket p \wedge \overline{q} \rrbracket \equiv \llbracket p \rrbracket \\ & = \{ (0) \} \\ & \llbracket p \rrbracket \wedge \llbracket q \rrbracket \equiv \llbracket p \rrbracket \\ & = \{ p.c \} \\ & \llbracket p \rrbracket \Rightarrow \llbracket q \rrbracket \end{aligned}$$

(End of Proof.)

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