



## Facilitating analysis by tailoring the programming model to the problem at hand

### Concurrent Programming

- Concurrency is pervasive and useful
- But it adds complexity:  
deadlocks, race conditions, starvation
- A tradeoff...

### Retaining “Just Enough” Concurrency

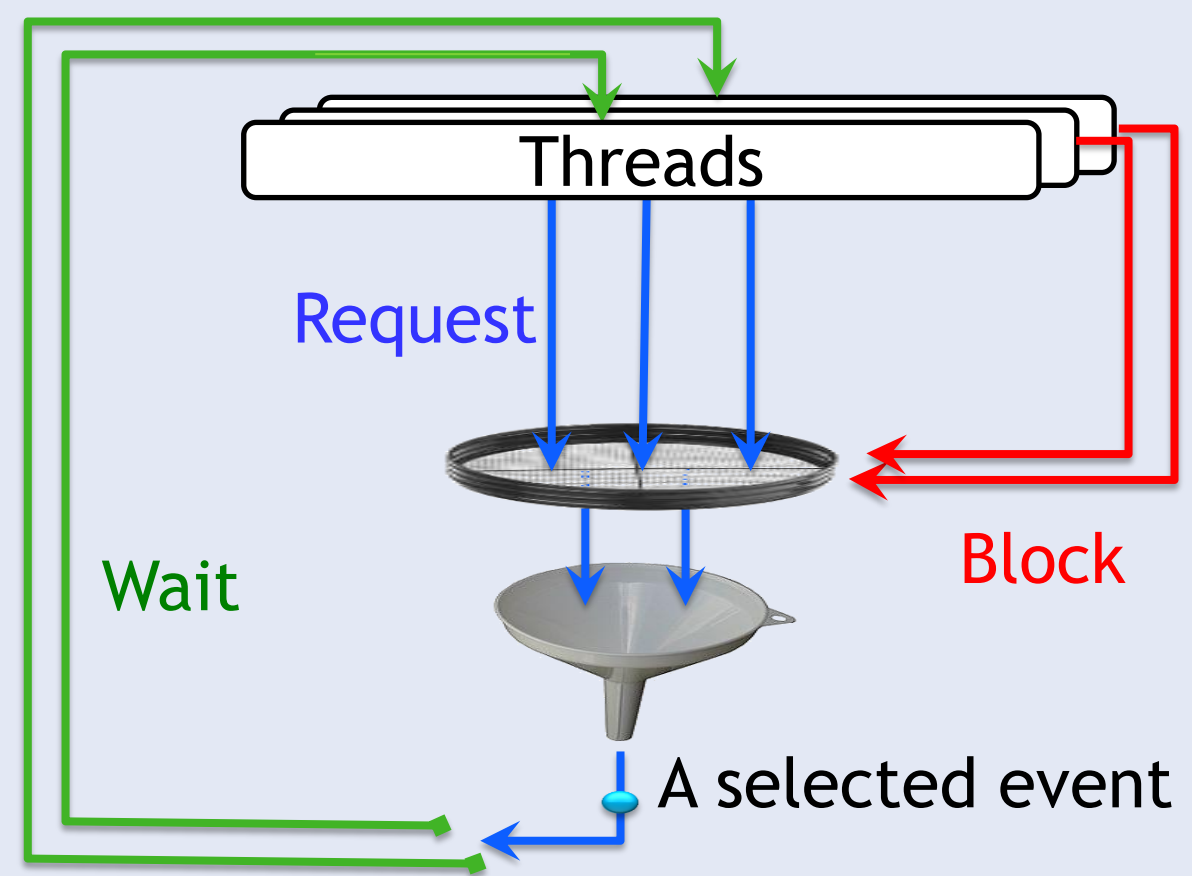
- Tailor the model to the task at hand
- Only pick the required concurrency idioms
- Solve the problem efficiently, while keeping the program simple

### The Request / Wait / Block (RWB) Model

Interweaving parallel behavior threads

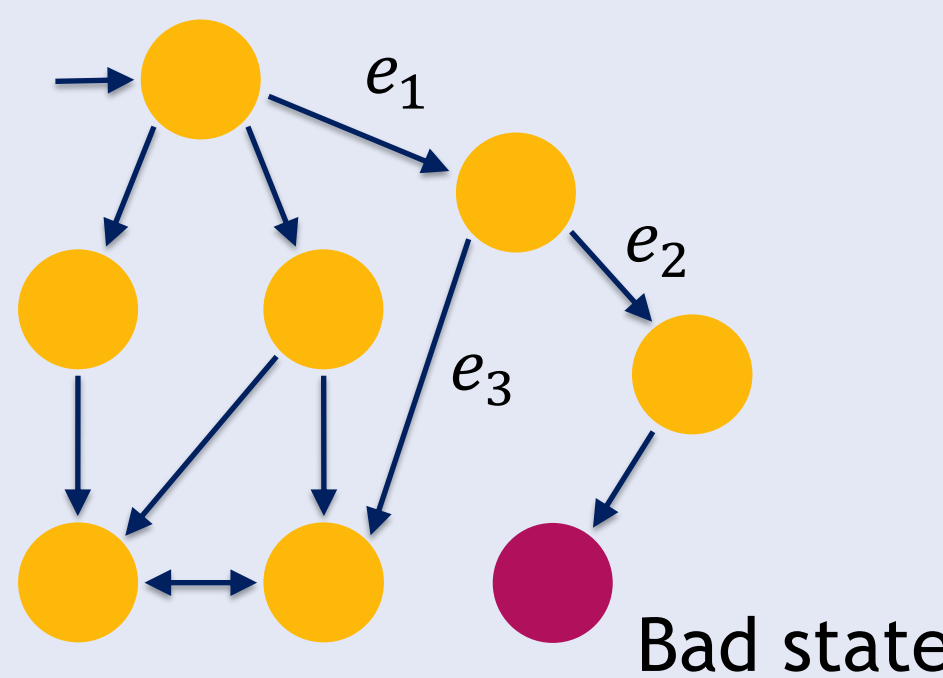
<b>Declaration</b>	Threads <b>request</b> , <b>wait-for</b> and <b>block</b> events
<b>Event Selection</b>	An event that is <b>requested</b> and not <b>blocked</b> is triggered
<b>Notification</b>	Resume threads that <b>requested</b> or <b>waited-for</b> the event

### The RWB Execution Cycle



### The Blocking Idiom Facilitates Program Repair

- Safety violation: a **bad** state is reachable
- A patch **blocks** bad transitions, without introducing deadlocks
- Incremental, non-intrusive repair



- Fix violation by adding the patch thread:
1. **wait-for**  $e_1$
  2. **wait-for**  $e_3$  while **blocking**  $e_2$

### Each Idiom Affords Unique Descriptive Succinctness and Makes Programs Smaller

- Smaller programs are easier to maintain and verify
- Each of **requesting**, **waiting-for** and **blocking** render some programs exponentially smaller
- Example:  $L_n = (0^{n-1} \cdot (0 + 1))^\omega$
- RWB implementation size:  $O(\log^2 n \cdot \log \log n)$
- Size without **blocking**:  $\Omega(n)$

RWB implementation for  $(0^5 \cdot (0 + 1))^\omega$

