

ASS Note 48
NIC # 17655

Simon S. Lam
UCLA
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SOME SATELLITE SIMULATION RESULTS

I. INTRODUCTION

In this note we report some simulation results for the slotted ALOHA satellite channel with a large number of small users (Model I in ASS Note 39) and the FCFS reservation scheme (see ASS Note 38 by L.G. Roberts). The simulation models are based upon the following assumptions:

- (1) no random noise errors (the satellite channel has perfect information feedback);
- (2) new messages are generated independently from Poisson distribution(s);
- (3) all ground stations have the same round-trip propagation delay R .

Note that in assumption (2), only the newly generated packets are derived independently from a Poisson distribution; any collisions, random retransmissions and queueing are accounted for in the simulations. Message delays are obtained by time-stamping messages at the time of their creation.

In general, our simulation results indicate the following:

- (1) Under "stable" conditions experimental results agree very well with analytic results (in ASS Notes 38 and 39) thus lending validity to approximations and assumptions made in the mathematical derivations.
- (2) The slotted ALOHA channel with a large number of small users can become unstable even though the input rate to such a channel is significantly below $1/e$. The FCFS reservation scheme depends upon using a fraction of the satellite channel in the ALOHA mode for making reservations and thus is subject to the same effects of instability. Some simple scheme is necessary to monitor the channel traffic and apply controls whenever unstable conditions are detected.

We are currently studying the dynamic behavior of the slotted ALOHA channel and will report our findings in a future note.

II. The Slotted ALOHA Simulations

A start-up period of a few hundred slots is allowed (depending on individual runs) before statistics are gathered to compute the average values for throughput rate S , traffic rate G , traffic distribution and average packet delay D . In addition to the above long-term average values, average values for the corresponding variables are also computed for all consecutive periods of 200 slots. These will be referred to as the short-term averages and they portray the approximate dynamic behavior of the channel. If during a run, all the short-term average values for the traffic rate are below 1, such a run is designated to be a "stable" run.*

"Stable" runs

In Figures 1 and 2 we plot simulation data for "stable" runs together with analytic results for S , G and D . Note the excellent agreement between analytic results and simulation data in almost all cases. In Figures 3-4, we show the comparisons between the probability density function for traffic arrivals obtained by simulation (from "stable" runs) and the Poisson density with the same average rate. Recall in ASS Note 17, it was shown that as $K \rightarrow \infty$ the traffic arrivals are Poisson distributed. Note in the figures that for $K=15$, the agreement is very good; even for $K=5$, the traffic density still approximates a Poisson density very well.

"Unstable" runs

In Figures 5-10, we show situations when during the runtime of a simulation the ALOHA channel becomes unstable (i.e. the average channel traffic rate increases as the throughput rate decreases). In Fig. 5, we see that for $K=5$, the channel traffic rate quickly diverges despite the fact that the input rate = 0.325 and is less than the theoretical maximum of 0.347 corresponding to $G=1$ (see Fig. 1). In Fig. 6 we see that increasing K to 7 enables the channel to accommodate the input rate of 0.325. The channel becomes unstable temporarily but regains stability without any external controls. Similar situations are shown in Figures 7 and 8 for a channel input rate of 0.35 and values of K equal to 15 and 40. In Fig. 9, we see that with $K=2$, the channel becomes unstable for an input rate as low as 0.25. Referring to Fig. 1, for $K=2$ the channel has a theoretical maximum throughput rate $S_{MAX}=0.295$. In Fig. 10, we see that if the input rate is reduced to 0.2, the channel traffic does not diverge. Thus the maximum input rate an ALOHA channel can accommodate is considerably less than $S_{MAX}(K)$ obtained by having the channel traffic rate $G=1$.

* The same run can still become unstable if the runtime is increased.

III. FCFS Reservation Simulations

Program features and assumptions

The scheduling algorithm used in this simulation is the same as outlined by L. G. Roberts in ASS Note 38 with minor differences. The channel is slotted such that for every M slots used for reserved data there is one slot for ALOHA traffic (including reservation requests). Each such ALOHA slot is further divided into V small slots. When there is no reserved data, the channel is switched into the ALOHA mode in which all the slots are used for ALOHA traffic. As soon as a reservation request was successful and the corresponding data packets are ready to transmit, the channel is switched back into the MIXED mode with only 1 ALOHA slot (containing V small slots) for every M data slots. Such a scheme is particularly useful when the input data contain a lot of multipacket messages.

In the simulations, we shall assume the following:

- (1) number of ground stations = 10;
- (2) roundtrip propagation delay $R = 10$;
- (3) input data arrive at σ messages/slot and consist of equal number of single packet and 8-packet messages such that the average message length = 4.5 packets and the average data input rate $S = 4.5\sigma$ packets/slot;
- (4) $V = 6$.

The simulation program can accept message length distributions other than the one in (3) above. The above assumptions are made to coincide with Roberts' so that we can compare the simulation results here with his analytic results. To specify the simulation completely, we need the following further assumptions:

- (5) message arrivals generated in the same slot for a given station are lumped together as a single message;
- (6) the random retransmission interval K in the ALOHA channel is assumed to be equal to $V=6$ so that every blocked ALOHA packet will randomize its retransmission over all the small slots in a single ALOHA slot (this assumption was made for the ease of programming and has been generalized in the simulation program to allow K to be multiples of V);
- (7) newly generated ALOHA packets will also randomize their transmissions over all the small slots within an ALOHA slot rather than transmitting in the first available small slot.

In all of our simulations, the ALOHA traffic in the channel consists of only the reservation requests. In this case, the channel performance is expected to be slightly better than Roberts' analytic results since he included in the ALOHA traffic acknowledgement messages. The simulation program does not provide for the generation of acknowledgement

messages. However, it does provide for the generation of arbitrary ALOHA packets at some specified Poisson rate to use the ALOHA portion of the channel in addition to the reservation request packets.

Comparison of analytic and simulation results

In Fig. 11, we show the tradeoffs between throughput and delay for the FCFS reservation scheme with $M=5$. $E[\text{block delay}]$ is the average message delay for both single and eight-packet messages. It includes both the ALOHA delay for the reservation request packet and the queueing delay in waiting for the reserved data slots. The analytic results for this tradeoff are taken from Fig. 2 in ASS Note 38. We see that the analytic and simulation results agree quite well. The simulation results give a slightly better channel performance than the analytic results because the simulations did not include the use of the ALOHA portion of the channel by acknowledgement packets. Hence the ALOHA packet delays are smaller because of the lower utilization. When the data input rate S is small, the difference in the two curves is negligible since the ALOHA portion of the channel is lightly utilized in both cases. In Fig. 12, we show the average throughput rate, traffic rate and packet delay for the ALOHA portion of the channel as a function of the data rate S .

Delay and stability considerations in determining M

What fraction f of the satellite channel is for the use of ALOHA traffic? The answer is $f = 1-S$ where S is the input data rate. Thus f is not a function of M ! If given S , our main concern is the average message delay, it seems to suggest the use of as large a value of M as possible since the ALOHA utilization does not depend on M whereas the FCFS queueing delay decreases as M increases. However, although the ALOHA utilization is independent of M , increase in the value of M gives rise to the following limitations. If we telescope our attention only on the ALOHA slots, the variance of the input process of the ALOHA channel increases as M increases. As a result the average ALOHA packet delay increases due to the larger variance. In addition, the variance of the traffic in the ALOHA channel is directly proportional to the variance of the input. As the variance of the traffic increases, the ALOHA channel becomes less stable! Since a stable ALOHA channel is necessary for acceptable reservation delays, too large a value of M cannot be used.

In Figs. 13 and 14, we show the channel performance for values of M other than 5. From the simulation data $M=7$ gives the best performance. In Figs. 15 and 16 we show situations when the ALOHA channel under a high data input rate becomes "unstable" for values of M equal to 7 and 10.

IV. Further comments

Operating policies

The operation of a slotted satellite channel in both the ALOHA mode and the FCFS reservation mode requires constant monitoring of the ALOHA channel traffic and the use of external controls (or built-in local SIMP policies) to reduce the channel input data rate and/or increase the random retransmission interval K to offset fluctuations in the channel traffic. From the simulation results, a minimum operating value for K is about 15. (Note in Fig. 2 that the delay curve corresponding to $K=15$ is quite near the optimum). At the same time, the average channel traffic rate G should be kept below 0.7.

An estimate for G

An estimate for G can be obtained by counting the number of empty slots in the ALOHA channel during an appropriate interval. The fraction of empty slots $f_0 \cong e^{-G}$ (see Figs. 3 and 4). Hence the estimate is

$$G \cong -\ln(f_0)$$

Final remarks

Both simulation programs are written in PL/I. Typical runs on the UCLA-CCN IBM 360/91 require approximately 3 CPU seconds execution time each for 10,000 slots of the FCFS reservation simulation and 4000 slots of the slotted ALOHA simulation.

ACKNOWLEDGEMENTS

The author would like to thank E. Walton who coded the FCFS reservation simulation and J. Spencer who coded the slotted ALOHA simulation.

Fig. 1 Throughput as a Function of ALOHA Channel Traffic

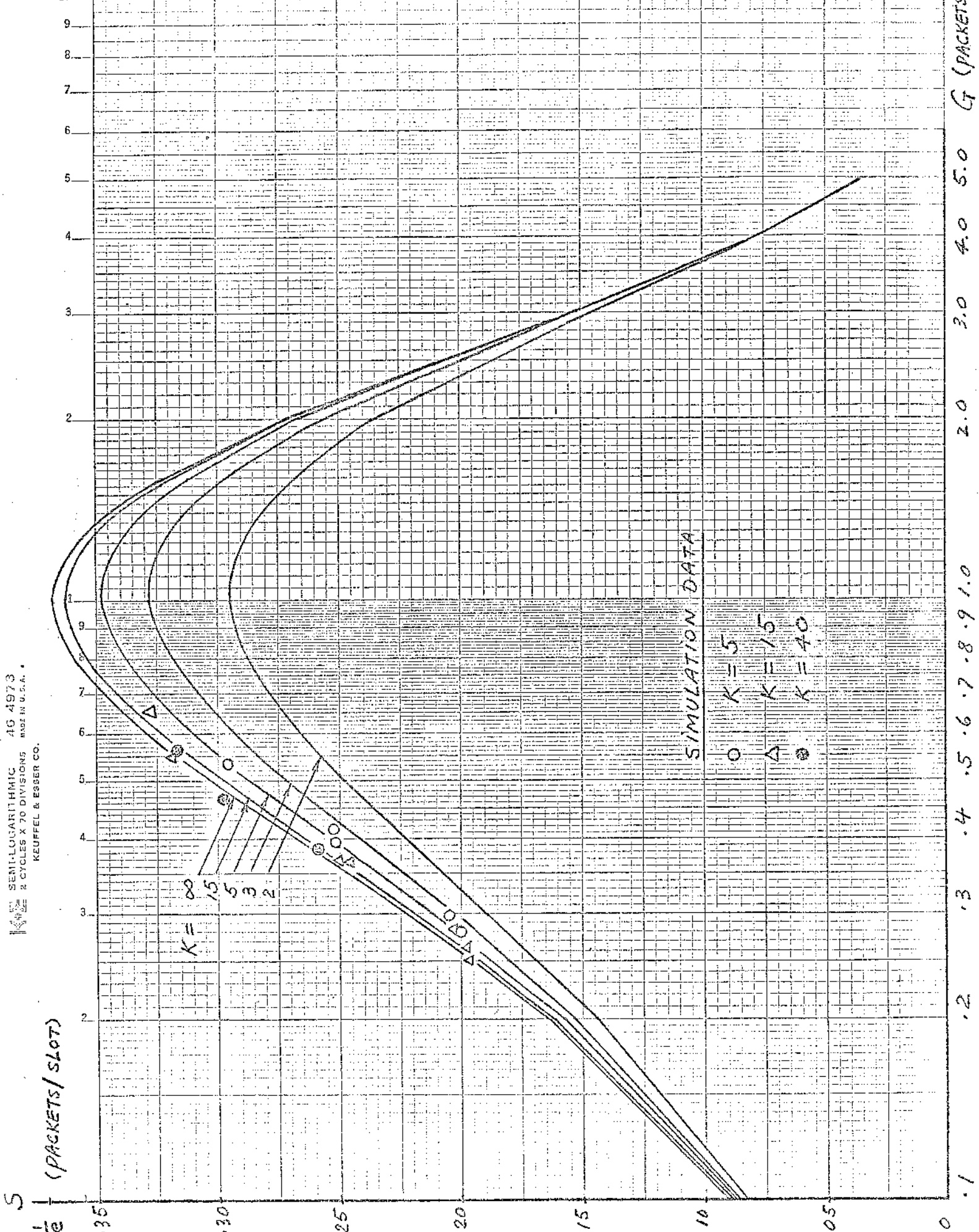
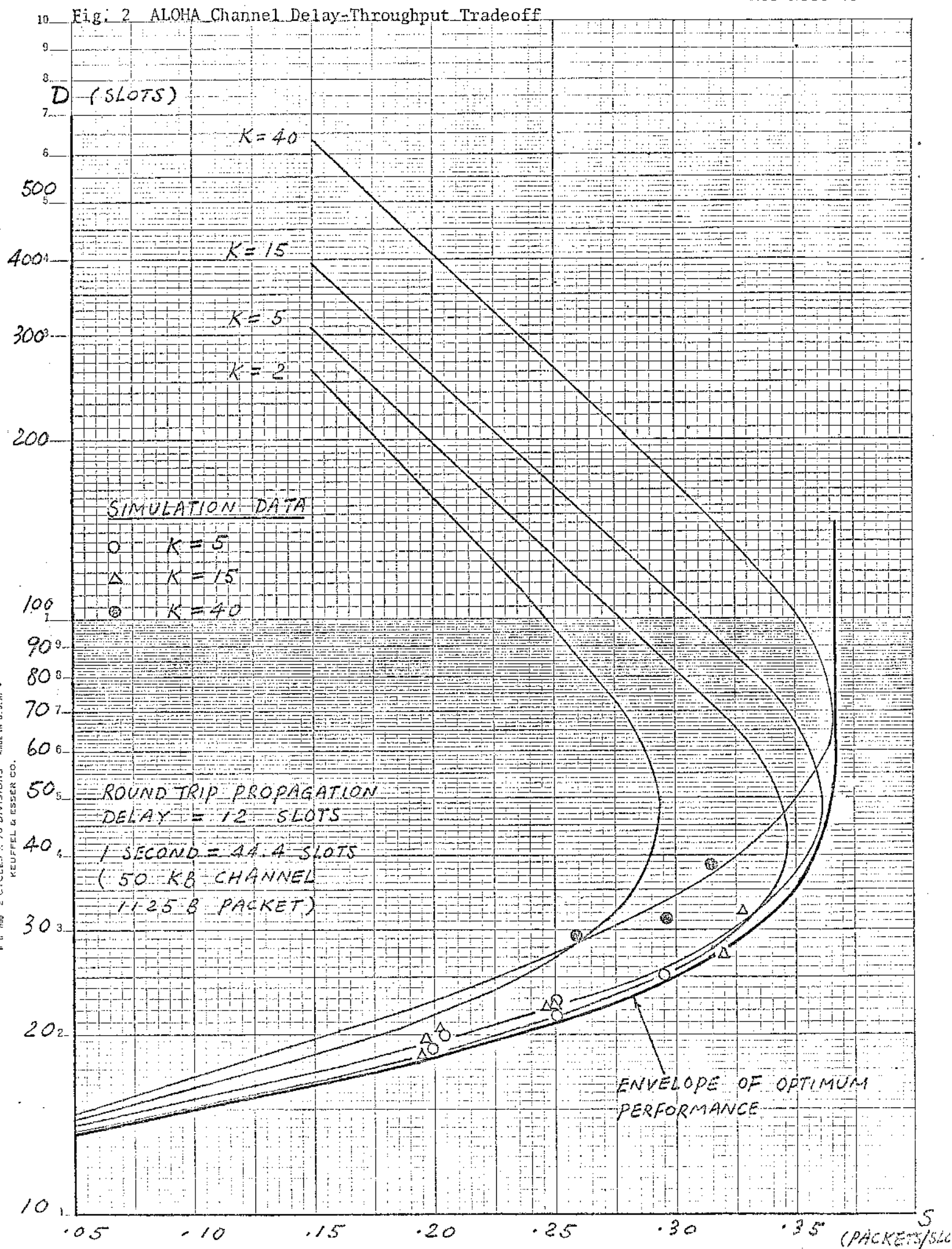


Fig. 2 ALOHA Channel Delay-Throughput Tradeoff



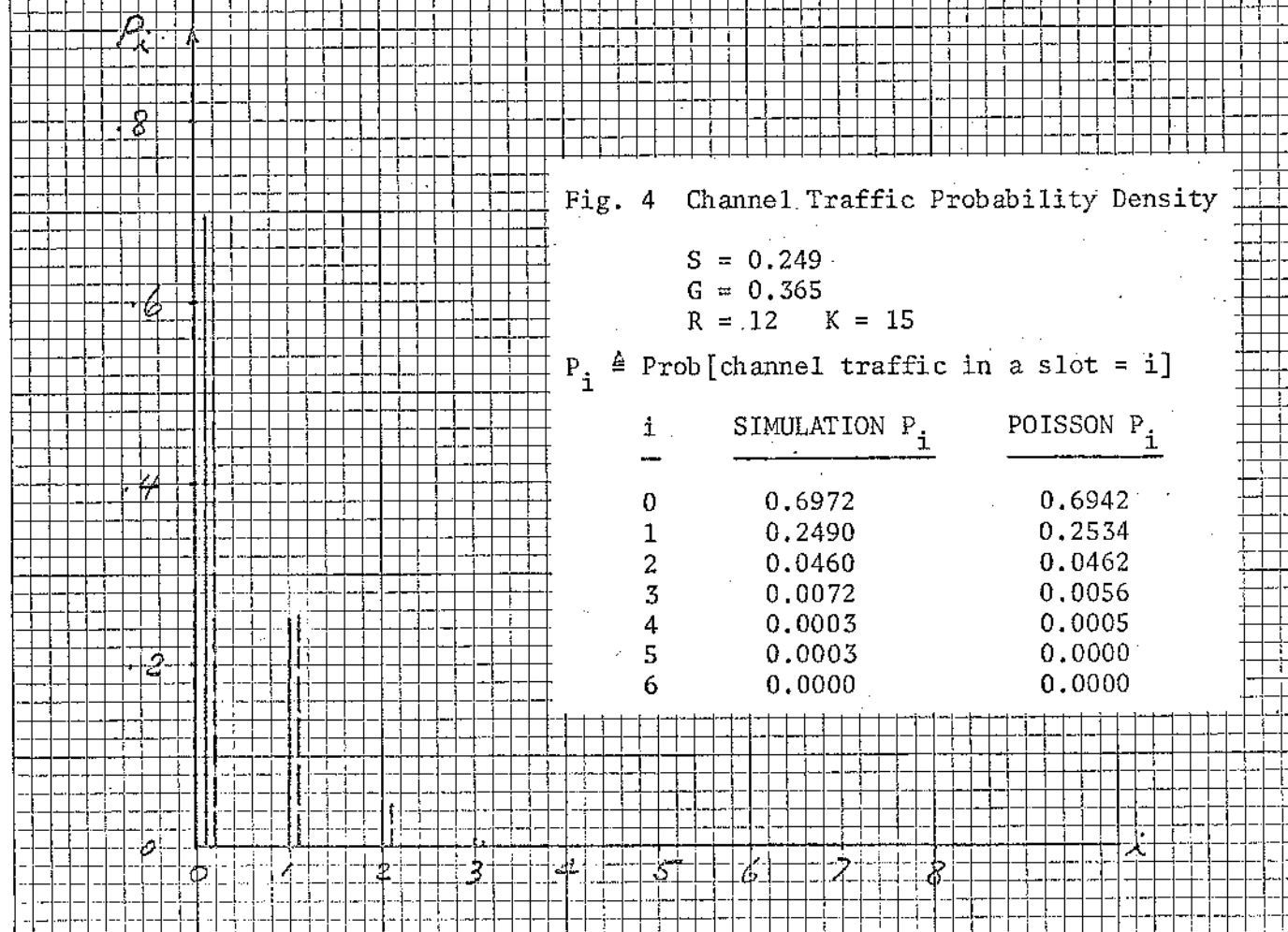
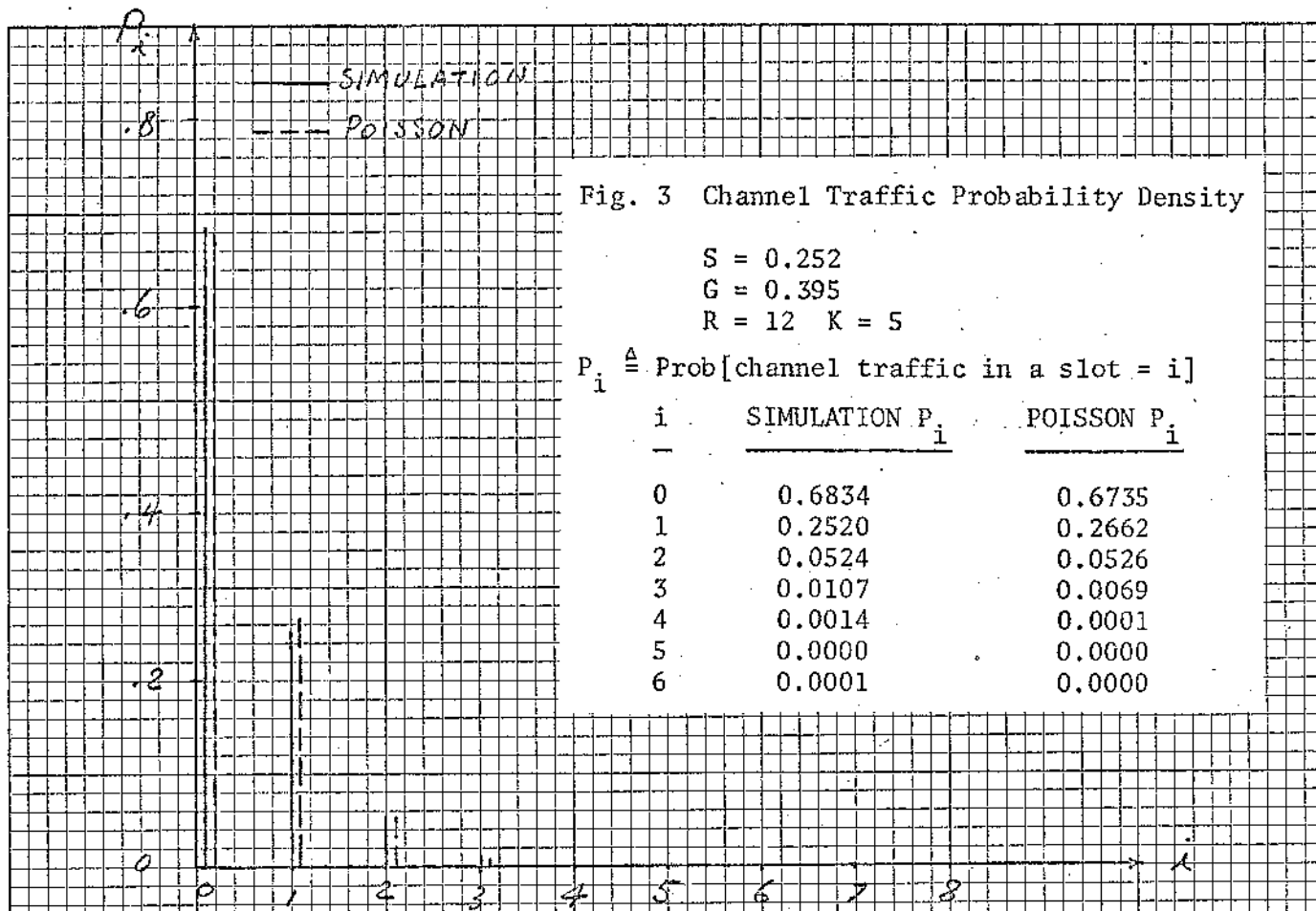


Fig. 5 Dynamic Behavior of a Slotted ALOHA Channel

INPLT PARAMETERS =>=>=>

INPLT RATE = 0.325

PROPAGATION DELAY = 12

K = 5

RUN TIME = 4000

SEED = 64783

AVERAGE VALUES IN 200 TIME SLOT PERIODS

TIME PERIOD	THROUGHPUT RATE = S	TRAFFIC RATE = G	PACKET DELAY = D	FRACTION OF TIME CHANNEL EMPTY
1 - 200	0.315	0.490	20.143	0.600
201 - 400	0.315	0.840	38.635	0.455
401 - 600	0.300	0.735	22.450	0.530
601 - 800	0.250	<u>1.920</u>	67.180	0.195
801 - 1000	0.190	<u>3.175</u>	121.526	0.045
1001 - 1200	0.020	<u>5.535</u>	151.750	0.005
1201 - 1400	0.000	<u>9.805</u>	0.000	0.000
1401 - 1600	0.000	<u>14.590</u>	0.000	0.000

LONG TERM STATISTICS =>=>=>=>=>=>=>

TIME FOR LONG TERM STATS =

Fig. 6 Dynamic Behavior of a Slotted ALOHA Channel

INPUT PARAMETERS =>=>=>

INFLT RATE = 0.325

PROPAGATION DELAY = 12

K = 7

RUN TIME = 4000

SEED = 64783

AVERAGE VALUES IN 200 TIME SLOT PERIODS

TIME PERIOD	THROUGHPUT RATE = S	TRAFFIC RATE = G	PACKET DELAY = D	FRACTION OF TIME CHANNEL EMPTY
1 = 200	0.315	0.455	19.714	0.620
201 = 400	0.310	0.685	27.903	0.525
401 = 600	0.365	0.860	35.562	0.425
601 = 800	0.325	0.930	39.446	0.405
801 = 1000	0.300	0.765	34.283	0.500
1001 = 1200	0.320	<u>1.490</u>	66.484	0.255
1201 = 1400	0.350	<u>1.135</u>	50.700	0.305
1401 = 1600	0.285	<u>1.230</u>	66.579	0.365
1601 = 1800	0.395	<u>1.370</u>	47.886	0.240
1801 = 2000	0.375	<u>1.000</u>	53.587	0.325
2001 = 2200	0.335	0.645	28.701	0.530
2201 = 2400	0.345	0.760	31.565	0.475
2401 = 2600	0.330	0.895	26.955	0.445
2601 = 2800	0.330	<u>1.130</u>	60.712	0.345
2801 = 3000	0.285	0.480	29.860	0.625
3001 = 3200	0.325	0.560	23.062	0.570
3201 = 3400	0.325	0.730	33.708	0.505
3401 = 3600	0.320	0.640	30.750	0.540
3601 = 3800	0.275	0.485	22.691	0.625
3801 = 4000	0.300	0.560	29.517	0.590

LONG TERM STATISTICS =>=>=>=>=>=>=>=>

THROUGHPUT = 0.325

TRAFFIC = 0.840

IDLE SLOTS = 0.464

AVERAGE DELAY = 38.292

TIME FOR LONG TERM STATS = 4000

Fig. 7 Dynamic Behavior of a Slotted ALOHA Channel

INFLT PARAMETERS =>=>=>

INFLT RATE = 0.350

PROPAGATION DELAY = 12

K = 15

RLN TIME = 4000

SEED = 39276

AVERAGE VALUES IN 200 TIME SLOT PERIODS

TIME PERIOD	THROUGHPUT RATE = S	TRAFFIC RATE = G	PACKET DELAY = D	FRACTION OF TIME CHANNEL EMPTY
1 = 200	0.330	0.510	17.924	0.595
201 = 400	0.370	0.605	30.892	0.530
401 = 600	0.360	0.860	32.764	0.425
601 = 800	0.340	0.840	43.147	0.435
801 = 1000	0.315	<u>1.415</u>	58.889	0.250
1001 = 1200	0.380	<u>1.260</u>	72.066	0.260
1201 = 1400	0.325	0.455	37.215	0.610
1401 = 1600	0.355	0.480	20.803	0.590
1601 = 1800	0.275	0.405	20.600	0.665
1801 = 2000	0.360	0.560	27.528	0.550
2001 = 2200	0.330	0.430	18.561	0.620
2201 = 2400	0.310	0.545	22.065	0.580
2401 = 2600	0.335	0.840	44.866	0.455
2601 = 2800	0.320	0.705	34.703	0.500
2801 = 3000	0.325	<u>1.085</u>	43.815	0.350
3001 = 3200	0.310	<u>1.715</u>	67.161	0.180
3201 = 3400	0.105	<u>3.255</u>	147.143	0.050
3401 = 3600	0.015	<u>5.910</u>	220.333	0.000
3601 = 3800	0.000	<u>9.155</u>	0.000	0.000
3801 = 4000	0.000	<u>12.400</u>	0.000	0.000

LONG TERM STATISTICS =>=>=>=>=>=>=>=>

THROUGHPUT = 0.270

TRAFFIC = 2.259

IDLE SLOTS = 0.371

AVERAGE DELAY = 41.287

TIME FOR LONG TERM STATS = 3800

INFLT PARAMETERS =>=>=>

INFLT RATE = 0.250

PROPAGATION DELAY = 12

K = 2

RLN TIME = 4000

SEED = 39276

AVERAGE VALUES IN 200 TIME SLOT PERIODS

TIME PERIOD	THROUGHPUT RATE = S	TRAFFIC RATE = G	PACKET DELAY = D	FRACTION OF TIME CHANNEL EMPTY
1 = 200	0.265	0.370	17.849	0.685
201 = 400	0.245	0.555	24.061	0.630
401 = 600	0.240	0.785	34.188	0.565
601 = 800	0.265	<u>1.035</u>	54.849	0.475
801 = 1000	0.215	<u>1.010</u>	42.581	0.555
1001 = 1200	0.165	<u>1.965</u>	35.788	0.405
1201 = 1400	0.130	<u>3.010</u>	86.615	0.300
1401 = 1600	0.035	<u>5.710</u>	20.857	0.065
1601 = 1800	0.000	<u>9.760</u>	0.000	0.000
1801 = 2000	0.000	<u>13.025</u>	0.000	0.000

LONG TERM STATISTICS =>=>=>=>=>=>=>

TIME FOR LONG TERM STATS =

INPLT PARAMETERS =>=>=>

INPLT RATE = 0.200

PROPAGATION DELAY = 12

K = 2

RUN TIME = 4000

SEED = 44854

AVERAGE VALUES IN 200 TIME SLOT PERIODS

TIME PERIOD	THROUGHPUT RATE = S	TRAFFIC RATE = G	PACKET DELAY = D	FRACTION OF TIME CHANNEL EMPTY
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1 * 200	0.185	0.310	22.081	0.765
201 * 400	0.170	0.280	21.853	0.775
401 * 600	0.185	0.605	31.892	0.650
601 * 800	0.210	0.600	48.333	0.630
801 * 1000	0.195	0.555	28.385	0.660
1001 * 1200	0.225	0.365	29.778	0.715
1201 * 1400	0.185	0.615	24.324	0.635
1401 * 1600	0.200	<u>1.430</u>	70.475	0.375
1601 * 1800	0.155	<u>1.240</u>	80.258	0.505
1801 * 2000	0.205	0.725	102.488	0.600
2001 * 2200	0.145	0.790	43.310	0.610
2201 * 2400	0.265	0.935	59.868	0.485
2401 * 2600	0.220	0.815	39.091	0.565
2601 * 2800	0.270	0.730	50.426	0.580
2801 * 3000	0.165	0.175	19.455	0.830
3001 * 3200	0.185	0.360	16.649	0.730
3201 * 3400	0.230	0.635	32.435	0.605
3401 * 3600	0.215	0.585	40.535	0.645
3601 * 3800	0.185	0.405	34.162	0.725
3801 * 4000	0.225	0.405	23.800	0.650

LONG TERM STATISTICS ==>==>==>==>==>==>

THROUGHPUT = 0.205

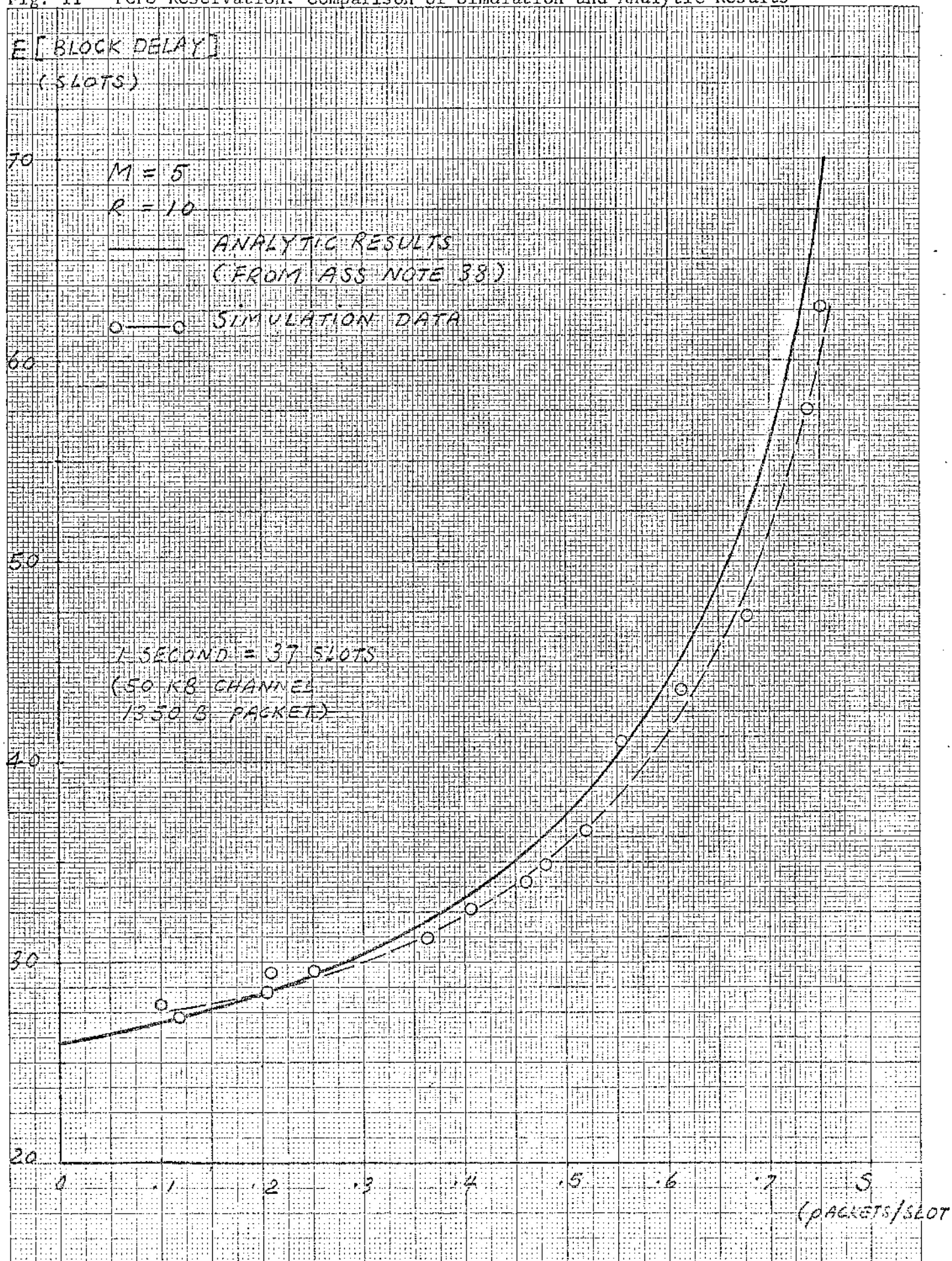
TRAFFIC = 0.669

IDLE SLOTS = 0.623

AVERAGE DELAY = 43.945

TIME FOR LONG TERM STATS = 3400

Fig. 11 FCFS Reservation: Comparison of Simulation and Analytic Results



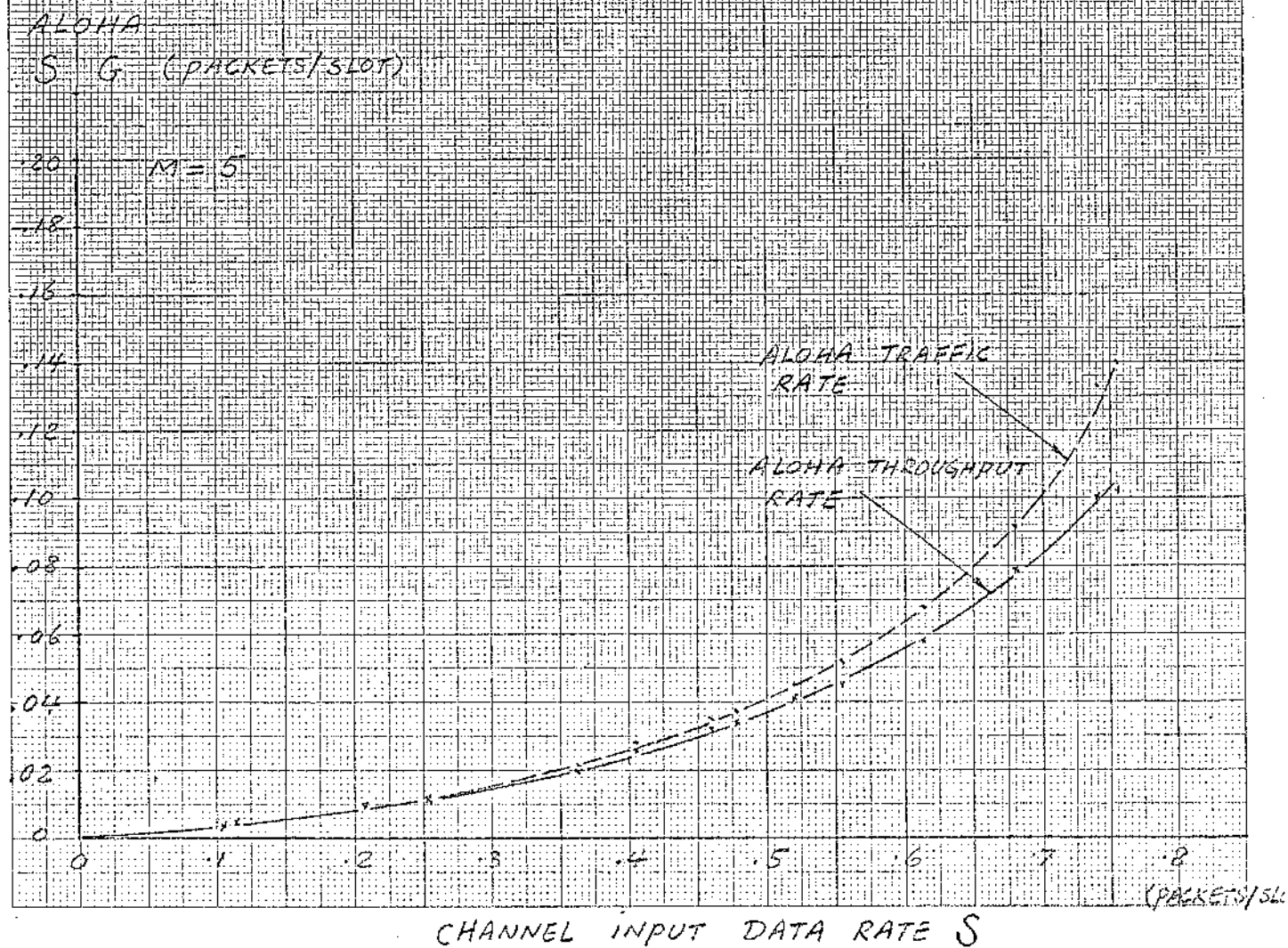
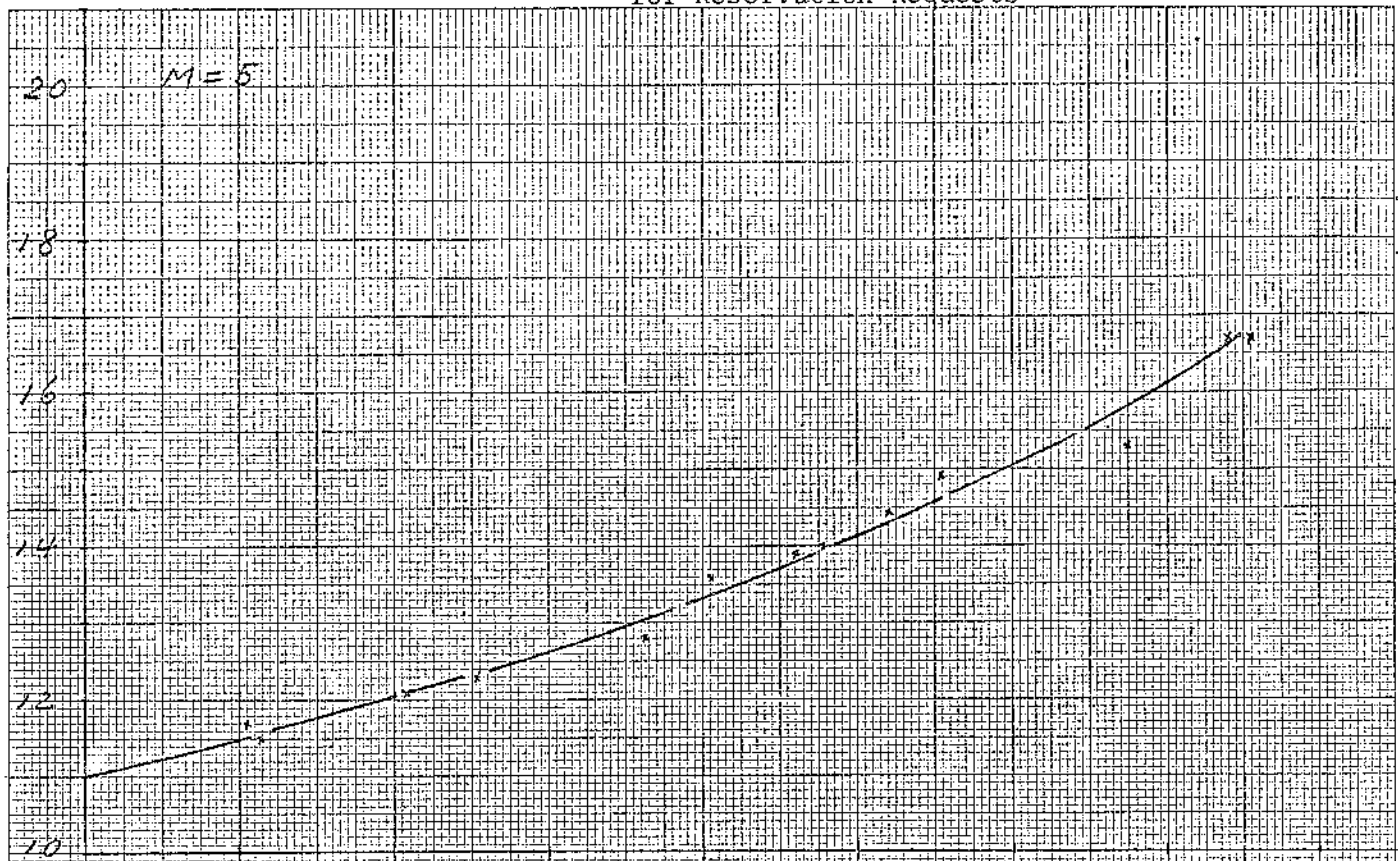
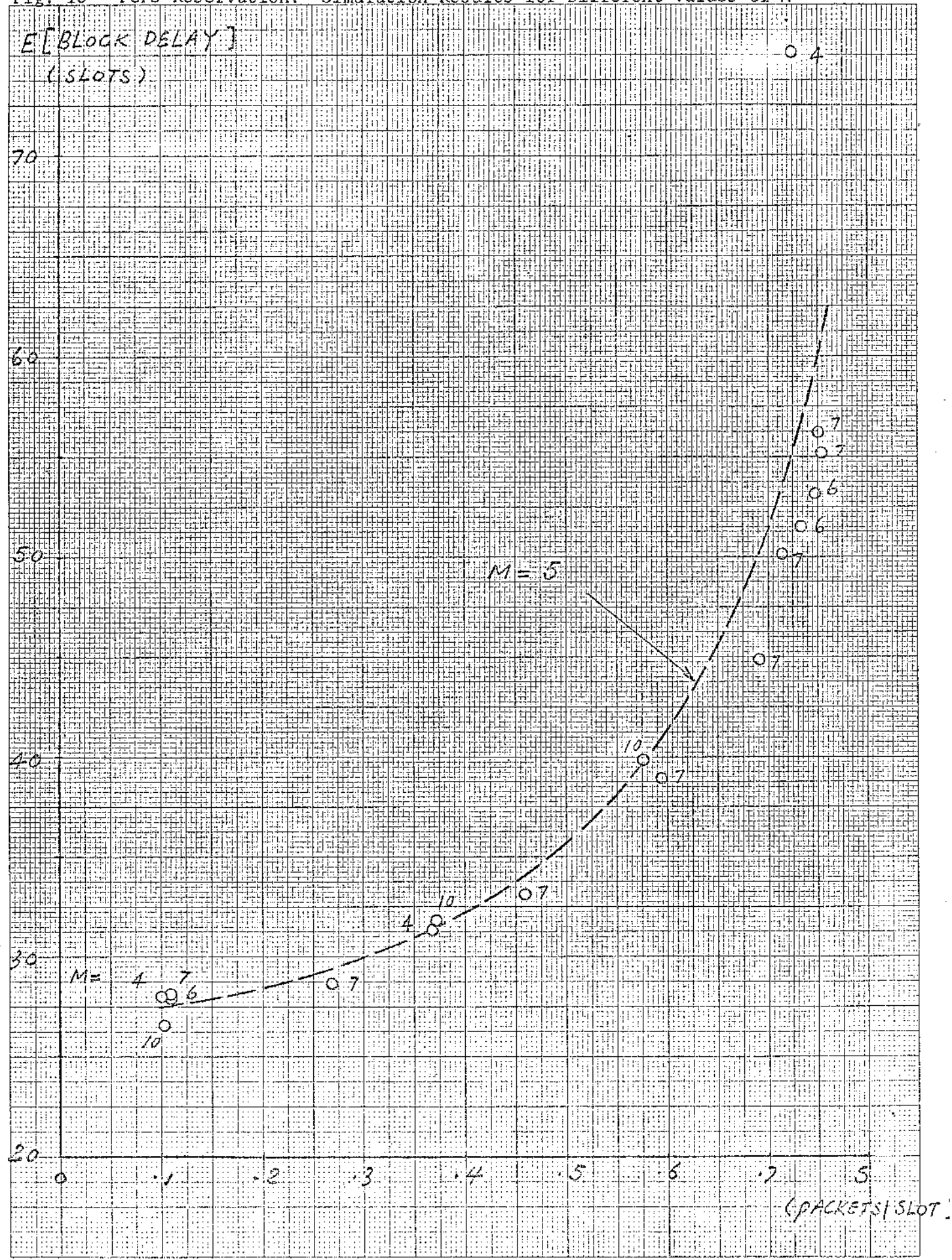


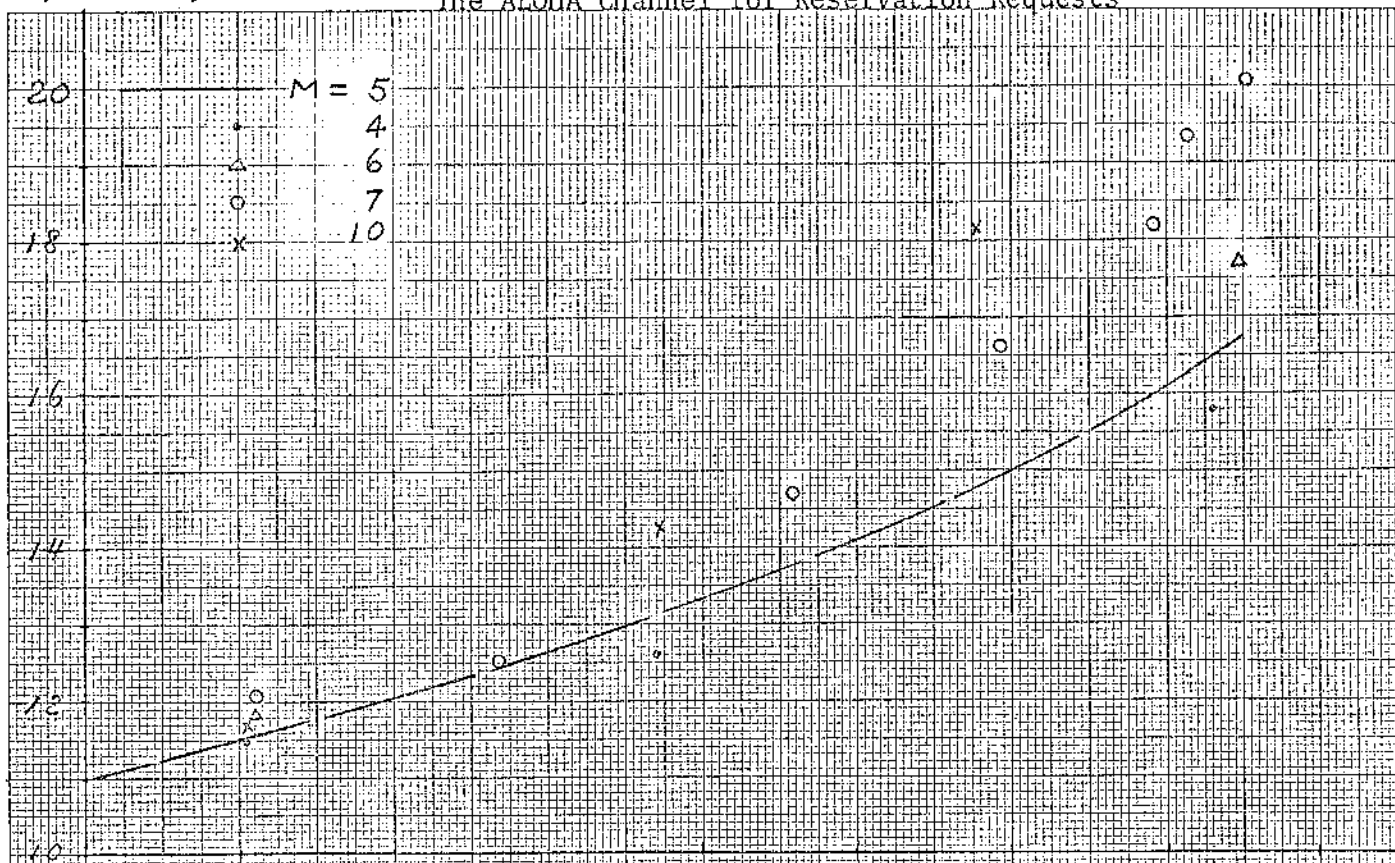
Fig. 13 FCFS Reservation: Simulation Results for Different Values of M



ALOHA PACKET
DELAY (SLOTS)

Fig. 14

Simulation Results for Different Values of M:
The ALOHA Channel for Reservation Requests



ALOHA

S G (PACKETS/SLOT)

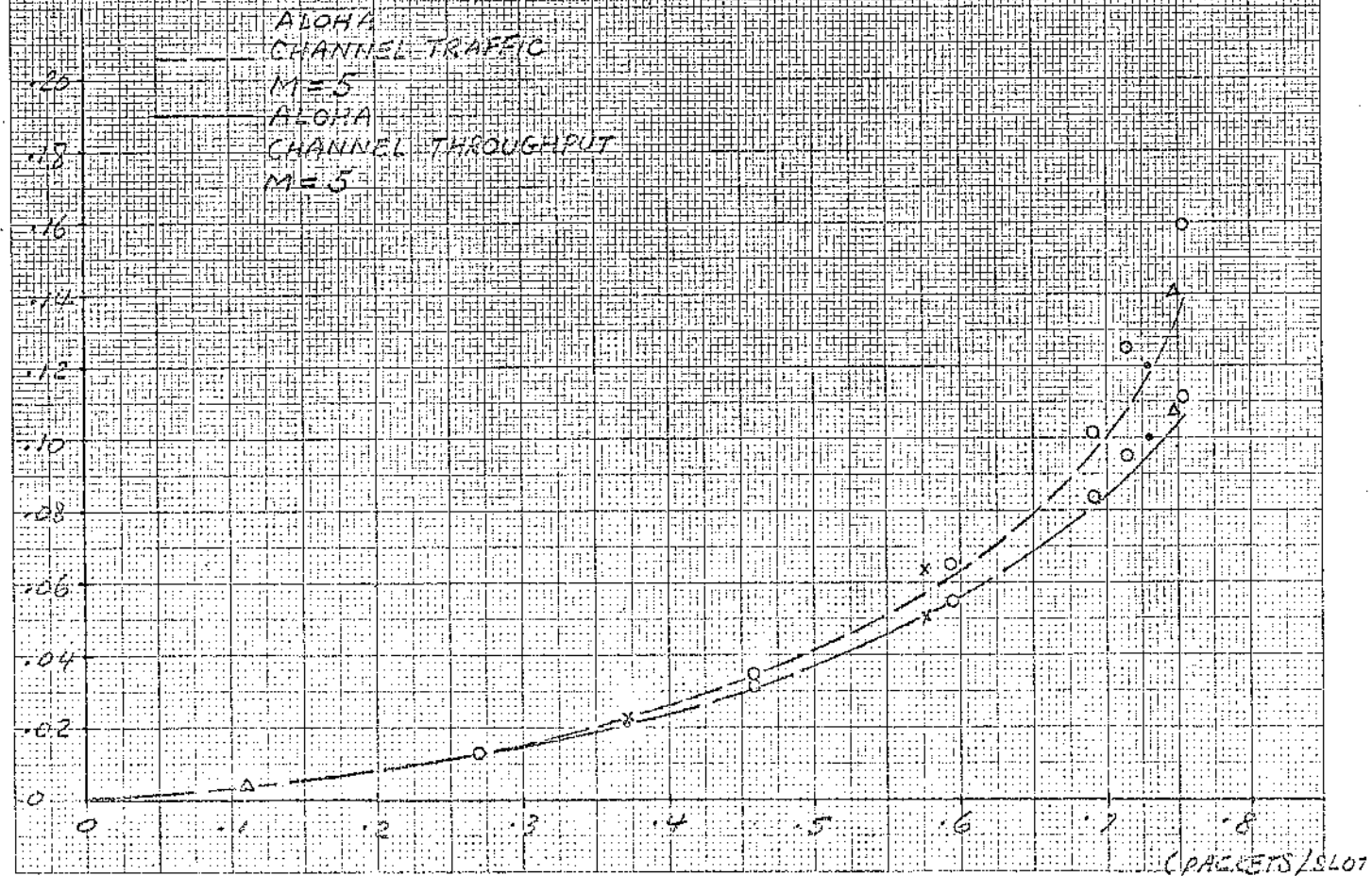
ALOHA
CHANNEL TRAFFIC

M=5

ALOHA

CHANNEL THROUGHPUT

M=5



CHANNEL INPUT DATA RATE S

Fig. 15 FCFS Reservation: Unstable ALOHA Channel

M = 7

Total channel input data rate = 0.75 packets/slot

Round-trip propagation delay R = 10 slots

K = 6

V = 6

No. of ground stations = 10

Period for short-term statistics = 200

SHORT-TERM STATISTICS FOR PERIODS OF 200 SLOTS BEFORE SIMULATION BLOWS UP:

*** AVERAGE DELAYS BETWEEN SLOTS 5427 AND 5627
 AVERAGE RESERVATION DELAY IS 34.522
 TOTAL RESERVATION DELAY IS 794.000
 NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 23.000
 AVERAGE BLOCK DELAY IS 73.650
 TOTAL BLOCK DELAY IS 1573.000
 NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 20.000
 AVERAGE ALOHA UTILIZATION IS 0.044
 ALOHA TRAFFIC AVERAGE IS 0.993
 DATA TO TOTAL SLOTS RATIO IS 0.565

*** AVERAGE DELAYS BETWEEN SLOTS 5628 AND 5828
 AVERAGE RESERVATION DELAY IS 81.444
 TOTAL RESERVATION DELAY IS 1466.000
 NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 18.000
 AVERAGE BLOCK DELAY IS 93.200
 TOTAL BLOCK DELAY IS 1864.000
 NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 20.000
 AVERAGE ALOHA UTILIZATION IS 0.030
 ALOHA TRAFFIC AVERAGE IS 1.298
 DATA TO TOTAL SLOTS RATIO IS 0.510

*** AVERAGE DELAYS BETWEEN SLOTS 5829 AND 6029
 AVERAGE RESERVATION DELAY IS 12.400
 TOTAL RESERVATION DELAY IS 248.000
 NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 20.000
 AVERAGE BLOCK DELAY IS 27.529
 TOTAL BLOCK DELAY IS 468.000
 NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 17.000
 AVERAGE ALOHA UTILIZATION IS 0.023
 ALOHA TRAFFIC AVERAGE IS 0.991
 DATA TO TOTAL SLOTS RATIO IS 0.270

*** AVERAGE DELAYS BETWEEN SLOTS 6030 AND 6230
 AVERAGE RESERVATION DELAY IS 12.278
 TOTAL RESERVATION DELAY IS 221.000
 NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 18.000
 AVERAGE BLOCK DELAY IS 30.500
 TOTAL BLOCK DELAY IS 671.000
 NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 22.000
 AVERAGE ALOHA UTILIZATION IS 0.026
 ALOHA TRAFFIC AVERAGE IS 1.352
 DATA TO TOTAL SLOTS RATIO IS 0.435

STATISTICS FOR THE WHOLE RUN:

*** TOTAL DELAYS BETWEEN SLOTS 0 AND 6200
 AVERAGE RESERVATION DELAY IS 33.650
 TOTAL RESERVATION DELAY IS 33460.000
 NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 1005.000
 AVERAGE BLOCK DELAY IS 77.453
 TOTAL BLOCK DELAY IS 77635.000
 NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 1003.000
 AVERAGE ALOHA UTILIZATION IS 0.100
 ALOHA TRAFFIC AVERAGE IS 0.999
 DATA TO TOTAL SLOTS RATIO IS 0.731

Fig. 16 FCFS Reservation: Unstable ALOHA Channel

M = 10

Total channel input data rate = 0.75 packets/slot

Round trip propagation delay R = 10 slots

K = 6

V = 6

No. of ground stations = 10

Period of short-term statistics = 200

SHORT-TERM STATISTICS FOR PERIODS OF 200 SLOTS BEFORE SIMULATION BLOWS UP:

*** AVERAGE DELAYS BETWEEN SLOTS 7638 AND 7838

AVERAGE RESERVATION DELAY IS 23.129

TOTAL RESERVATION DELAY IS 717.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 31.000

AVERAGE BLOCK DELAY IS 41.222

TOTAL BLOCK DELAY IS 1113.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 27.000

AVERAGE ALOHA UTILIZATION IS 0.058

ALOHA TRAFFIC AVERAGE IS 0.086

DATA TO TOTAL SLOTS RATIO IS 0.560

*** AVERAGE DELAYS BETWEEN SLOTS 7839 AND 8039

AVERAGE RESERVATION DELAY IS 32.515

TOTAL RESERVATION DELAY IS 1073.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 33.000

AVERAGE BLOCK DELAY IS 57.824

TOTAL BLOCK DELAY IS 1966.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 34.000

AVERAGE ALOHA UTILIZATION IS 0.177

ALOHA TRAFFIC AVERAGE IS 0.441

DATA TO TOTAL SLOTS RATIO IS 0.850

*** AVERAGE DELAYS BETWEEN SLOTS 8040 AND 8240

AVERAGE RESERVATION DELAY IS 31.696

TOTAL RESERVATION DELAY IS 729.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 23.000

AVERAGE BLOCK DELAY IS 58.643

TOTAL BLOCK DELAY IS 1642.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 28.000

AVERAGE ALOHA UTILIZATION IS 0.066

ALOHA TRAFFIC AVERAGE IS 0.793

DATA TO TOTAL SLOTS RATIO IS 0.715

*** AVERAGE DELAYS BETWEEN SLOTS 8241 AND 8441

AVERAGE RESERVATION DELAY IS 11.400

TOTAL RESERVATION DELAY IS 228.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 20.000

AVERAGE BLOCK DELAY IS 32.211

TOTAL BLOCK DELAY IS 612.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 19.000

AVERAGE ALOHA UTILIZATION IS 0.037

ALOHA TRAFFIC AVERAGE IS 1.259

DATA TO TOTAL SLOTS RATIO IS 0.555

*** AVERAGE DELAYS BETWEEN SLOTS 8442 AND 8642

AVERAGE RESERVATION DELAY IS 18.360

TOTAL RESERVATION DELAY IS 409.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 25.000

AVERAGE BLOCK DELAY IS 32.042

TOTAL BLOCK DELAY IS 769.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 24.000

AVERAGE ALOHA UTILIZATION IS 0.038

ALOHA TRAFFIC AVERAGE IS 1.578

DATA TO TOTAL SLOTS RATIO IS 0.450

STATISTICS FOR THE WHOLE RUN:

*** TOTAL DELAYS BETWEEN SLOTS 0 AND 8701

AVERAGE RESERVATION DELAY IS 24.980

TOTAL RESERVATION DELAY IS 33998.000

NUMBER OF RESERVATION SAMPLES THIS TIME PERIOD IS 1361.000

AVERAGE BLOCK DELAY IS 49.568

TOTAL BLOCK DELAY IS 67363.000

NUMBER OF BLOCK SAMPLES THIS TIME PERIOD IS 1359.000

AVERAGE ALOHA UTILIZATION IS 0.035

ALOHA TRAFFIC AVERAGE IS 0.287

DATA TO TOTAL SLOTS RATIO IS 0.624