

1. (Recursive Programming; 28 points)

```

(a)   tie ([], [])          = []
       tie (x:xs) (y: ys) = (x,y): (tie (xs,ys))

       untie []             = ([], [])
       untie (x,y): zs = ((x:xs), (y: ys))
                           where (xs,ys) = untie zs

```

(b) Define function `gpos` of two arguments k and n , both integers and $n \geq 0$, where `(gpos k n)` returns a list of n consecutive integers starting at k . Then,

```
pos n = gpos 1 n
```

The definition of `gpos` is given by

```

gpos k 0 = []
gpos k (n+1) = k: (gpos (k+1) n)

```

```

(c)   check f g zs = glt (f zs) (g zs)
       glt xs [] = True
       glt [] ys = False
       glt (x: xs) (y: ys) = (x >= y) && (glt xs ys)

```

2. (Finite State Machines)

(a) The following machine accepts strings with even number of zeroes.

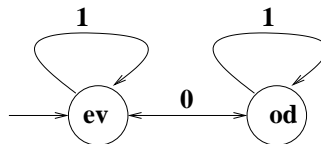


Figure 1: Machine that accepts strings with even number of zeroes

The following machine accepts strings ending in “11”.

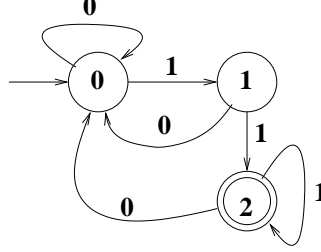


Figure 2: Machine that accepts strings ending in “11”

Combine the two machines.

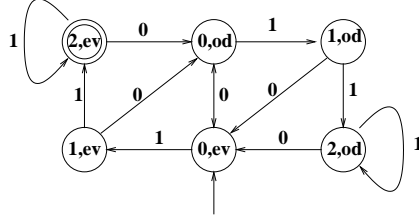


Figure 3: Combination of the two machines given above

(b) Define the following predicates.

- $emp(x)$: string x is empty
- $p0(x)$: last symbol of string x is 0
- $p1(x)$: last symbol of string x is 1
- $eq(x)$: number of zero blocks in x equals the number of one blocks
- $lt(x)$: number of zero blocks in x is 1 less than the number of one blocks
- $gt(x)$: number of zero blocks in x is 1 more than the number of one blocks

Annotate as follows:

- A: $emp \wedge eq$
- B: $p0 \wedge gt$
- C: $p1 \wedge lt$
- D: $p1 \wedge eq$
- E: $p0 \wedge eq$

Proof for transition:

- C to E: $p1(x) \wedge lt(x) \Rightarrow p0(x0) \wedge eq(x0)$
- E to E: $p0(x) \wedge eq(x) \Rightarrow p0(x0) \wedge eq(x0)$
- E to C: $p0(x) \wedge eq(x) \Rightarrow p1(x1) \wedge lt(x1)$

(c) $0(0|1)^*1$